

Let  $p^*$  be an approximation to a number  $p$  (not necessarily its floating point approximation). Then:

- **Absolute error:**  $|p - p^*|$
- **Relative error:**  $\frac{|p - p^*|}{|p|}$  (provided  $p \neq 0$ )

**Significant Digits.** The number  $p^*$  is said to approximate  $p$  to  $t$  significant digits (or figures) if  $t$  is the largest nonnegative integer such that

$$\frac{|p - p^*|}{|p|} \leq 5 \times 10^{-t}.$$

$$\begin{aligned}\cos(h) &= 1 - \frac{h^2}{2!} + \frac{h^4}{4!} - \frac{h^6}{6!} + \dots \\ \sin(h) &= h - \frac{h^3}{3!} + \frac{h^5}{5!} - \frac{h^7}{7!} + \dots \\ \tan(h) &= h + \frac{h^3}{3} + \frac{2h^5}{15} + \frac{17h^7}{315} + \dots \\ \sinh(x) &= \frac{e^x - e^{-x}}{2} \\ \cosh(x) &= \frac{e^x + e^{-x}}{2} \\ e^h &= 1 + h + \frac{h^2}{2!} + \frac{h^3}{3!} + \frac{h^4}{4!} + \dots \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots\end{aligned}$$

Suppose  $\{\beta_n\}_{n \geq 1}$  is a sequence known to converge to 0 and  $\{\alpha_n\}_{n \geq 1}$  a sequence that converges to a value  $\alpha \in \mathbb{R}$ .

If there exists a constant  $K > 0$  such that

$$|\alpha_n - \alpha| \leq K|\beta_n| \quad \text{for large } n,$$

then we say that  $\{\alpha_n\}$  converges to  $\alpha$  with **rate of convergence**  $\mathcal{O}(\beta_n)$ . This is denoted as:

$$\alpha_n = \alpha + \mathcal{O}(\beta_n).$$

**Order of Convergence.** Let  $\{p_n\}_{n=0}^\infty$  be a sequence with  $\lim_{n \rightarrow \infty} p_n = p$  and  $p_n \neq p$  for all  $n$ . If there exist positive constants  $\lambda$  and  $\alpha$  such that

$$\lim_{n \rightarrow \infty} \frac{|p_{n+1} - p|}{|p_n - p|^\alpha} = \lambda,$$

then we say that  $\{p_n\}_{n=0}^\infty$  converges to  $p$  with order  $\alpha$  and constant  $\lambda$ .

- In general, the larger  $\alpha$ , the more rapid the convergence.
- The constant  $\lambda$  affects this speed, but generally it is not as important as the order  $\alpha$ .

The following statements are equivalent for a matrix  $A \in \mathbb{R}^{n \times n}$ :

- $\det(A) \neq 0$
- The system  $Ax = 0$  has a unique solution  $x = 0$  (i.e., the column vector with all zero entries)
- The system  $Ax = b$  has a unique solution for any  $n$ -dimensional column vector  $b$
- The matrix  $A$  is *nonsingular*, that is,  $A^{-1}$  exists

If GE can be performed on the system  $Ax = b$  without row interchanges, then the matrix  $A$  can be factored as:

$$A = LU, \quad \Rightarrow \quad L(Ux) = b$$

Let  $A \in \mathbb{R}^{n \times n}$ .  $A$  is said to be **strictly diagonally dominant** if

$$|a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}}^n |a_{ij}| \quad \text{for all } i = 1, \dots, n,$$

A strictly diagonally dominant matrix  $A$  is nonsingular. Moreover, in this case, GE can be performed on any linear system  $Ax = b$  without row or column interchanges, and the computations will be robust to round-off errors.

If  $A \in \mathbb{R}^{n \times n}$  is **positive definite**, then:

- $A$  is *symmetric*, i.e.  $A = A^\top$ .
- $\mathbf{x}^\top A \mathbf{x} > 0$  for all nonzero  $\mathbf{x} \in \mathbb{R}^n$ .
- $A$  is *nonsingular*.
- $a_{ii} > 0$  for each  $i = 1, \dots, n$ .
- $\max_{1 \leq i, j \leq n} |a_{ij}| \leq \max_{1 \leq i \leq n} |a_{ii}|$ .
- $(a_{ij})^2 < a_{ii}a_{jj}$  for each  $i \neq j$ .

**Cauchy-Schwarz Inequality.** For each  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^n$ ,

$$|\mathbf{x}^\top \mathbf{y}| = \left| \sum_{i=1}^n x_i y_i \right| \leq \left( \sum_{i=1}^n x_i^2 \right)^{1/2} \left( \sum_{i=1}^n y_i^2 \right)^{1/2} = \|\mathbf{x}\|_2 \|\mathbf{y}\|_2.$$

**Matrix Norms and Distances.** A *matrix norm* on  $\mathbb{R}^{n \times n}$  is a function  $\|\cdot\|: \mathbb{R}^{n \times n} \rightarrow \mathbb{R}$  such that for any  $\mathbf{A}, \mathbf{B} \in \mathbb{R}^{n \times n}$  and any  $\alpha \in \mathbb{R}$ :

- $\|\mathbf{A}\| \geq 0$ ,
- $\|\mathbf{A}\| = 0$  if and only if  $\mathbf{A} = \mathbf{0}$ ,
- $\|\alpha \mathbf{A}\| = |\alpha| \|\mathbf{A}\|$ ,
- $\|\mathbf{A} + \mathbf{B}\| \leq \|\mathbf{A}\| + \|\mathbf{B}\|$  (triangle inequality),
- $\|\mathbf{AB}\| \leq \|\mathbf{A}\| \cdot \|\mathbf{B}\|$  (submultiplicative property).

**Spectral radius**  $\rho(A)$  of a matrix  $A$  is defined by

$$\rho(A) = \max |\lambda|,$$

where  $\lambda$  is an eigenvalue of  $\mathbf{A}$ .

Let  $A \in \mathbb{R}^{n \times n}$ . Then,

$$\|A\|_2 = \sqrt{\rho(A^\top A)}.$$

Furthermore, for any natural norm  $\|\cdot\|$ , we have

$$\rho(A) \leq \|A\|.$$

**Condition Number ( $\ell_2$  norm).** Let  $A \in \mathbb{R}^{n \times n}$  be nonsingular. The condition number of  $A$  (w.r.t.  $\ell_2$ ) is

$$\kappa(A) = \|A\|_2 \|A^{-1}\|_2 = \sqrt{\frac{\lambda_{\max}(A^\top A)}{\lambda_{\min}(A^\top A)}},$$

where  $\lambda_{\max}(A^\top A)$  and  $\lambda_{\min}(A^\top A)$  are the largest and smallest eigenvalues (in absolute value) of  $A^\top A$ .

- Well-conditioned if  $\kappa(A) \approx 1$ .
- Ill-conditioned if  $\kappa(A) \gg 1$ .

$$A = D - L - U \quad \text{for iterative}$$

1. **Jacobi:**

$$x^{(k+1)} = D^{-1}(L + U)x^{(k)} + D^{-1}b$$

2. **Gauss-Seidel:**

$$x^{(k+1)} = (D - L)^{-1}Ux^{(k)} + (D - L)^{-1}b$$

3. **Successive Over-Relaxation (SOR):**

$$x^{(k+1)} = (D - \omega L)^{-1}[(1 - \omega)D + \omega U]x^{(k)} + \omega(D - \omega L)^{-1}b$$

For any  $\mathbf{x}^{(0)} \in \mathbb{R}^n$ , the sequence  $\{\mathbf{x}^{(k)}\}_{k=1}^\infty$  defined by the iteration

$$\mathbf{x}^{(k+1)} = T\mathbf{x}^{(k)} + \mathbf{c}$$

converges to the unique solution of

$$\mathbf{x} = T\mathbf{x} + \mathbf{c}$$

if and only if  $\rho(T) < 1$ .

**Newton's Method.**

Given a function  $f(x) \in C^2[a, b]$ , Newton's method generates a sequence:

$$p_n = p_{n-1} - \frac{f(p_{n-1})}{f'(p_{n-1})}, \quad n = 1, 2, \dots$$

**Convergence:** Quadratic if  $f(p) = 0$ ,  $f'(p) \neq 0$ , and  $p_0$  is close to the root.

**Secant Method (derivative-free):**

$$p_n = p_{n-1} - f(p_{n-1}) \cdot \frac{p_{n-1} - p_{n-2}}{f(p_{n-1}) - f(p_{n-2})}$$

**Convergence:** Superlinear, with order approximately  $\alpha = \frac{1+\sqrt{5}}{2} \approx 1.618$

**Modified Newton Method (for root of multiplicity  $m$ ).**

If  $f(x)$  has a root of multiplicity  $m$  at  $p$ , standard Newton converges only linearly.

To restore quadratic convergence:

$$p_n = p_{n-1} - \frac{mf(p_{n-1})}{f'(p_{n-1})}, \quad (\text{if multiplicity } m \text{ is known})$$

**Alternative:** Apply Newton to  $\mu(x) = \frac{f(x)}{f'(x)}$ . This leads to:

$$p_n = p_{n-1} - \frac{f(p_{n-1})f'(p_{n-1})}{(f'(p_{n-1}))^2 - f(p_{n-1})f''(p_{n-1})}$$

**Note:** Requires second derivative  $f''(x)$ ; restores quadratic convergence.

### Lagrange Interpolation Formula.

Given  $n+1$  distinct data points  $(x_0, f(x_0)), (x_1, f(x_1)), \dots, (x_n, f(x_n))$ , the unique polynomial  $P(x) \in \mathbb{P}_n$  that interpolates the data is:

$$P(x) = \sum_{j=0}^n f(x_j) L_j(x)$$

where the Lagrange basis polynomials  $L_j(x)$  are defined by:

$$L_j(x) = \prod_{\substack{0 \leq k \leq n \\ k \neq j}} \frac{x - x_k}{x_j - x_k}$$

**Numerical Differentiation (Equally Spaced Nodes,  $n=2$ ).** If  $x_1 = x_0 + h$  and  $x_2 = x_0 + 2h$ , then:

$$f'(x_0) = \frac{1}{h} \left( -\frac{3}{2}f(x_0) + 2f(x_0 + h) - \frac{1}{2}f(x_0 + 2h) \right) + \frac{h^2}{3} f^{(3)}(\xi(x_0)),$$

$$f'(x_1) = \frac{1}{h} \left( -\frac{1}{2}f(x_1 - h) + \frac{1}{2}f(x_1 + h) \right) - \frac{h^2}{6} f^{(3)}(\xi(x_1)),$$

$$f'(x_2) = \frac{1}{h} \left( \frac{1}{2}f(x_2 - 2h) - 2f(x_2 - h) + \frac{3}{2}f(x_2) \right) + \frac{h^2}{3} f^{(3)}(\xi(x_2)).$$

**Divided Differences & Newton Form (compact).** Let distinct nodes  $x_0, \dots, x_n$ . Define divided differences recursively

$$f[x_i] = f(x_i), \quad f[x_i, \dots, x_{i+k}] = \frac{f[x_{i+1}, \dots, x_{i+k}] - f[x_i, \dots, x_{i+k-1}]}{x_{i+k} - x_i}$$

The Newton interpolant is

$$P_n(x) = \sum_{k=0}^n f[x_0, \dots, x_k] \prod_{j=0}^{k-1} (x - x_j),$$

**Cubic Spline Interpolation Types.** Given data points  $(x_0, y_0), \dots, (x_n, y_n)$ , a cubic spline consists of piecewise cubics  $S_i(x)$  with continuous first and second derivatives.

- **Natural (Free) Spline:**

$$S''(x_0) = 0 \text{ and } S''(x_n) = 0$$

- **Clamped (Complete) Spline:**

$$S'(x_0) = f'(x_0) \text{ and } S'(x_n) = f'(x_n) \text{ (derivatives specified)}$$

- **Not-a-Knot Spline:**

$$S''' \text{ is continuous at } x_1 \text{ and } x_{n-1} \text{ (i.e., } S_0'''(x_1) = S_1'''(x_1), \\ S_{n-2}'''(x_{n-1}) = S_{n-1}'''(x_{n-1}))$$

### Numerical Differentiation via Lagrange Interpolation.

Let  $x_0, x_1, \dots, x_n \in [a, b]$  and suppose  $f \in C^{n+1}[a, b]$ . Then the Lagrange interpolation formula gives:

$$f(x) = \sum_{k=0}^n f(x_k) L_k(x) + \frac{f^{(n+1)}(\xi(x))}{(n+1)!} (x - x_0)(x - x_1) \cdots (x - x_n)$$

for some  $\xi(x) \in [a, b]$ , where  $L_k(x)$  are the Lagrange basis polynomials. Differentiating and evaluating at  $x = x_j$ , we obtain:

$$f'(x_j) = \sum_{k=0}^n f(x_k) L'_k(x_j) + \frac{f^{(n+1)}(\xi)}{(n+1)!} \prod_{\substack{0 \leq k \leq n \\ k \neq j}} (x_j - x_k)$$

This is an  $(n+1)$ -point formula for estimating  $f'(x_j)$  from values  $f(x_0), \dots, f(x_n)$ .

### Gaussian Quadrature (5-point Rule).

Gaussian quadrature approximates

$$\int_{-1}^1 f(x) dx \approx \sum_{i=1}^n w_i f(x_i)$$

where  $x_i$  are the roots of the  $n$ th Legendre polynomial  $P_n(x)$ , and  $w_i$  are weights.

**Change of Interval:** To apply over interval  $[a, b]$ , transform:

$$\int_a^b f(x) dx = \frac{b-a}{2} \int_{-1}^1 f\left(\frac{b-a}{2}t + \frac{a+b}{2}\right) dt$$

where the *nodes*  $x_i$  are the  $n$  distinct zeros of the Legendre polynomial  $P_n(x)$  (orthogonal on  $[-1, 1]$  with weight 1), and the *weights* are

$$w_i = \frac{2}{(1 - x_i^2) [P'_n(x_i)]^2}.$$

### Legendre polynomials (Bonnet's recursion).

$$P_0(x) = 1, \quad P_1(x) = x,$$

$$(n+1)P_{n+1}(x) = (2n+1)xP_n(x) - nP_{n-1}(x) \quad (n \geq 1).$$

Useful derivative identity (for weights):

$$P'_n(x) = \frac{n}{1-x^2} (P_{n-1}(x) - xP_n(x)).$$

**Degree of Precision (Exactness).** A quadrature rule

$$\int_a^b f(x) dx \approx \sum_{i=1}^n w_i f(x_i)$$

is said to have *degree of precision*  $m$  if it integrates exactly every polynomial  $p(x)$  with  $\deg p \leq m$ , but fails for some polynomial of degree  $m+1$ .

For Gaussian quadrature with  $n$  nodes, the degree of precision is  $2n-1$ .

**Lipschitz Condition.** A function  $f(t, y)$  satisfies a Lipschitz condition in  $y$  if there exists  $L > 0$  such that

$$|f(t, y_1) - f(t, y_2)| \leq L|y_1 - y_2| \quad \text{for all } y_1, y_2 \in \mathbb{R}, a \leq t \leq b.$$

Then we say  $f$  is *Lipschitz in  $y$* . Or, sufficient condition  $\left| \frac{\partial f}{\partial y}(t, y) \right| \leq L$

**Existence and Uniqueness Theorem.** Let  $D = \{(t, y) \in \mathbb{R}^2 : a \leq t \leq b, -\infty < y < \infty\}$ . Suppose  $f(t, y)$  is continuous on  $D$ , and Lipschitz in  $y$ . Then the initial value problem:

$$y'(t) = f(t, y), \quad y(a) = \alpha, \quad a \leq t \leq b,$$

has a unique solution  $y(t)$  on  $[a, b]$ .

**Well-posedness:** The same conditions (continuity and Lipschitz) imply the IVP is *well posed*: small changes in inputs (e.g. initial value or  $f(t, y)$ ) lead to small changes in solution  $y(t)$ .

### Euler's Method (for IVPs).

Consider the initial value problem (IVP):

$$y'(t) = f(t, y), \quad y(a) = \alpha, \quad a \leq t \leq b.$$

Let  $h > 0$  be the *step size*, assumed constant for simplicity:

$$h = \frac{b-a}{N}, \quad t_j = a + jh, \quad j = 0, 1, \dots, N.$$

We seek approximations  $y_j \approx y(t_j)$  via:

$$y_{j+1} = y_j + h f(t_j, y_j), \quad j = 0, 1, \dots, N-1.$$

### Remarks:

- Local truncation error:  $\mathcal{O}(h^2)$
- Global truncation error:  $\mathcal{O}(h)$  (first-order method)
- Variable step sizes  $h_j$  may also be used when needed.

### Runge-Kutta Methods (for IVPs).

Consider the IVP

$$y'(t) = f(t, y), \quad y(a) = \alpha, \quad a \leq t \leq b$$

with step size  $h = (b-a)/N$ ,  $t_i = a + ih$ .

*General explicit two-stage RK method:*

$$k_1 = hf(t_i, w_i), \\ k_2 = hf(t_i + \alpha h, w_i + \beta k_1), \\ w_{i+1} = w_i + a k_1 + b k_2.$$

Choose  $a, b, \alpha, \beta$  to maximize local order.

Two common  $\mathcal{O}(h^2)$  choices:

- **Midpoint method:**  $a = 0, b = 1, \alpha = \beta = 1/2$

$$k_1 = hf(t_i, w_i), \quad k_2 = hf\left(t_i + \frac{h}{2}, w_i + \frac{k_1}{2}\right), \quad w_{i+1} = w_i + k_2.$$

- **Modified Euler:**  $a = b = \frac{1}{2}, \alpha = \beta = 1$

$$k_1 = hf(t_i, w_i), \quad k_2 = hf(t_i + h, w_i + k_1), \quad w_{i+1} = w_i + \frac{k_1 + k_2}{2}.$$