

Intéraction Logiciel Architecture

Les performances et leur mesure

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Motivation

Argumentation



Illustration

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Partie1 Intro

Argumentation



Illustration

Metrics How to Measure Performance ?

What scale ?

- Application level (TPS, Frame/s, .../)
- Run to completion
- Method level
- Instruction level

Tools

- Wall clock
- gettimeofday
- Performance counters

Full application or function call ?

- Cold start / warmup
- Statistic / multiple calls
- How many calls

Vocabulaire : Cycle par seconds / Flops

Units

Mips Million operation per second

Flops Floating point operation per second

<http://www.top500.org>

Flops/Watt Floating point operation per watt per second

<http://www.green500.org>

IPC : Instructions per Cycle

How to mesure

- Analytique (wall clock)
- Instrumentation
 - “Portable” (`gettimeofday()`)
 - Hardware (hardware performance counter)

Vocabulaire Peak / Sustained

Notions

Peak performance : maximal theoretical performance, assuming no bubble

Sustain performance : real achieved performance, on a real benchmark

Cost

- What is the percentage you're ready to lose ? 90% 95% ?
- How many are you ready to pay (time, money) to minimise this loss ?

<http://www.top500.org>

Vocabulaire Lois

- Moore http://en.wikipedia.org/wiki/Moore's_law
- Amdahl http://en.wikipedia.org/wiki/Amdahl's_law
- Memory bound http://en.wikipedia.org/wiki/IO_bound
- CPU bound http://en.wikipedia.org/wiki/CPU_bound

Vocabulaire Speedup

Notion

Speedup $S = 100 * \frac{T_{seq}}{T_{opt}}$

Can be between

- 1 and N processor
- 1 and vectorized
- non optimized versus optimized version

Mesure Quality what are the execution conditions : data set, computer workload, reproducibly, ...

Computer science has to use “human science” tools & methodology

Amdahl2

Argumentation

Assume that a task has two independent parts, A and B. B takes roughly 25% of the time of the whole computation. By working very hard, one may be able to make this part 5 times faster, but this only reduces the time for the whole computation by a little. In contrast, one may need to perform less work to make part A be twice as fast. This will make the computation much faster than by optimizing part B, even though B's speed-up is greater by ratio, (5x versus 2x)

Illustration

Two independent parts

A **B**

Original process



Make **B** 5x faster



Vocabulaire Amdahl

Argumentation

The speedup of a program using multiple processors in parallel computing is limited by the sequential fraction of the program. For example, if 95% of the program can be parallelized, the theoretical maximum speedup using parallel computing would be 20x as shown in the diagram, no matter how many processors are used.

Illustration

