

Slides from FYS3150/4150 Lectures

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Aug 31, 2014

Overview of week 37

Linear Algebra and differential equations.

- Monday: Repetition from last week
- Discussion of some Armadillo examples and matrix operations.
- Discussion of classes in C++
- Iterative methods, Gauss-Seidel. Cubic spline and interpolation (chapter 6).
- Tuesday: Eigenvalue problems
- We start discussing Jacobi's algorithm, chapter 7.
- Presentation of project 2.
- Computer-Lab: Project 1 and project 2.

Iterative methods, Chapter 6

- Direct solvers such as Gauss elimination and LU decomposition discussed last week.
- Iterative solvers such as Basic iterative solvers, Jacobi, Gauss-Seidel, Successive over-relaxation. These methods are easy to parallelize, as we will see later. Much used in solutions of partial differential equations.
- Other iterative methods such as Krylov subspace methods with Generalized minimum residual (GMRES) and Conjugate gradient etc will not be discussed.

Iterative methods, Jacobi's method

It is a simple method for solving

$$\hat{A}\mathbf{x} = \mathbf{b},$$

where \hat{A} is a matrix and \mathbf{x} and \mathbf{b} are vectors. The vector \mathbf{x} is the unknown.

It is an iterative scheme where we start with a guess for the unknown, and after $k + 1$ iterations we have

$$\mathbf{x}^{(k+1)} = \hat{D}^{-1}(\mathbf{b} - (\hat{L} + \hat{U})\mathbf{x}^{(k)}),$$

with $\hat{A} = \hat{D} + \hat{U} + \hat{L}$ and \hat{D} being a diagonal matrix, \hat{U} an upper triangular matrix and \hat{L} a lower triangular matrix.

If the matrix \hat{A} is positive definite or diagonally dominant, one can show that this method will always converge to the exact solution.

Iterative methods, Jacobi's method

We can demonstrate Jacobi's method by this 4×4 matrix problem. We assume a guess for the vector elements $x_i^{(0)}$, a guess which represents our first iteration. The new values are obtained by substitution

$$\begin{aligned} x_1^{(1)} &= (b_1 - a_{12}x_2^{(0)} - a_{13}x_3^{(0)} - a_{14}x_4^{(0)})/a_{11} \\ x_2^{(1)} &= (b_2 - a_{21}x_1^{(0)} - a_{23}x_3^{(0)} - a_{24}x_4^{(0)})/a_{22} \\ x_3^{(1)} &= (b_3 - a_{31}x_1^{(0)} - a_{32}x_2^{(0)} - a_{34}x_4^{(0)})/a_{33} \\ x_4^{(1)} &= (b_4 - a_{41}x_1^{(0)} - a_{42}x_2^{(0)} - a_{43}x_3^{(0)})/a_{44}, \end{aligned}$$

which after $k + 1$ iterations reads

$$\begin{aligned} x_1^{(k+1)} &= (b_1 - a_{12}x_2^{(k)} - a_{13}x_3^{(k)} - a_{14}x_4^{(k)})/a_{11} \\ x_2^{(k+1)} &= (b_2 - a_{21}x_1^{(k)} - a_{23}x_3^{(k)} - a_{24}x_4^{(k)})/a_{22} \\ x_3^{(k+1)} &= (b_3 - a_{31}x_1^{(k)} - a_{32}x_2^{(k)} - a_{34}x_4^{(k)})/a_{33} \end{aligned}$$

Iterative methods, Jacobi's method

We can generalize the above equations to

$$x_i^{(k+1)} = (b_i - \sum_{j=1, j \neq i}^n a_{ij}x_j^{(k)})/a_{ii}$$

or in an even more compact form as

$$\mathbf{x}^{(k+1)} = \hat{D}^{-1}(\mathbf{b} - (\hat{L} + \hat{U})\mathbf{x}^{(k)}),$$

with $\hat{A} = \hat{D} + \hat{U} + \hat{L}$ and \hat{D} being a diagonal matrix, \hat{U} an upper triangular matrix and \hat{L} a lower triangular matrix.

Iterative methods, Gauss-Seidel's method

Our 4×4 matrix problem

$$\begin{aligned}x_1^{(k+1)} &= (b_1 - a_{12}x_2^{(k)} - a_{13}x_3^{(k)} - a_{14}x_4^{(k)})/a_{11} \\x_2^{(k+1)} &= (b_2 - a_{21}x_1^{(k)} - a_{23}x_3^{(k)} - a_{24}x_4^{(k)})/a_{22} \\x_3^{(k+1)} &= (b_3 - a_{31}x_1^{(k)} - a_{32}x_2^{(k)} - a_{34}x_4^{(k)})/a_{33} \\x_4^{(k+1)} &= (b_4 - a_{41}x_1^{(k)} - a_{42}x_2^{(k)} - a_{43}x_3^{(k)})/a_{44},\end{aligned}$$

can be rewritten as

$$\begin{aligned}x_1^{(k+1)} &= (b_1 - a_{12}x_2^{(k)} - a_{13}x_3^{(k)} - a_{14}x_4^{(k)})/a_{11} \\x_2^{(k+1)} &= (b_2 - a_{21}x_1^{(k+1)} - a_{23}x_3^{(k)} - a_{24}x_4^{(k)})/a_{22} \\x_3^{(k+1)} &= (b_3 - a_{31}x_1^{(k+1)} - a_{32}x_2^{(k+1)} - a_{34}x_4^{(k)})/a_{33} \\x_4^{(k+1)} &= (b_4 - a_{41}x_1^{(k+1)} - a_{42}x_2^{(k+1)} - a_{43}x_3^{(k+1)})/a_{44},\end{aligned}$$

which allows us to utilize the preceding solution (forward

Iterative methods, Gauss-Seidel's method

We can generalize

$$\begin{aligned}x_1^{(k+1)} &= (b_1 - a_{12}x_2^{(k)} - a_{13}x_3^{(k)} - a_{14}x_4^{(k)})/a_{11} \\x_2^{(k+1)} &= (b_2 - a_{21}x_1^{(k+1)} - a_{23}x_3^{(k)} - a_{24}x_4^{(k)})/a_{22} \\x_3^{(k+1)} &= (b_3 - a_{31}x_1^{(k+1)} - a_{32}x_2^{(k+1)} - a_{34}x_4^{(k)})/a_{33} \\x_4^{(k+1)} &= (b_4 - a_{41}x_1^{(k+1)} - a_{42}x_2^{(k+1)} - a_{43}x_3^{(k+1)})/a_{44},\end{aligned}$$

to the following form

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j>i} a_{ij}x_j^{(k)} - \sum_{j<i} a_{ij}x_j^{(k+1)} \right), \quad i = 1, 2, \dots, n.$$

The procedure is generally continued until the changes made by an iteration are below some tolerance.

The convergence properties of the Jacobi method and the

Iterative methods, Successive over-relaxation

Given a square system of n linear equations with unknown \mathbf{x} :

$$\hat{\mathbf{A}}\mathbf{x} = \mathbf{b}$$

where:

$$\hat{\mathbf{A}} = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}, \quad \mathbf{x} = \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_n \end{bmatrix}.$$

Iterative methods, Successive over-relaxation

Then \mathbf{A} can be decomposed into a diagonal component \mathbf{D} , and strictly lower and upper triangular components \mathbf{L} and \mathbf{U} :

$$\hat{\mathbf{A}} = \hat{\mathbf{D}} + \hat{\mathbf{L}} + \hat{\mathbf{U}},$$

where

$$\mathbf{D} = \begin{bmatrix} a_{11} & 0 & \cdots & 0 \\ 0 & a_{22} & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & a_{nn} \end{bmatrix}, \quad \mathbf{L} = \begin{bmatrix} 0 & 0 & \cdots & 0 \\ a_{21} & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & 0 \end{bmatrix}, \quad \mathbf{U} = \begin{bmatrix} 0 & a_{12} & \cdots & a_{1n} \\ 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \cdots & 0 \end{bmatrix}.$$

The system of linear equations may be rewritten as:

$$(\mathbf{D} + \omega\mathbf{L})\mathbf{x} = \omega\mathbf{b} - [\omega\mathbf{U} + (\omega - 1)\mathbf{D}]\mathbf{x}$$

for a constant $\omega > 1$.

Iterative methods, Successive over-relaxation

The method of successive over-relaxation is an iterative technique that solves the left hand side of this expression for x_i , using previous value for x on the right hand side. Analytically, this may be written as:

$$\mathbf{x}^{(k+1)} = (\mathbf{D} + \omega\mathbf{L})^{-1}(\omega\mathbf{b} - [\omega\mathbf{U} + (\omega - 1)\mathbf{D}]\mathbf{x}^{(k)}).$$

However, by taking advantage of the triangular form of $(\mathbf{D} + \omega\mathbf{L})$, the elements of $\mathbf{x}^{(k+1)}$ can be computed sequentially using forward substitution:

$$x_i^{(k+1)} = (1-\omega)x_i^{(k)} + \frac{\omega}{a_{ii}} \left(b_i - \sum_{j>i} a_{ij}x_j^{(k)} - \sum_{j<i} a_{ij}x_j^{(k+1)} \right), \quad i = 1, 2, \dots, n.$$

The choice of relaxation factor is not necessarily easy, and depends upon the properties of the coefficient matrix. For symmetric, positive-definite matrices it can be proven that $0 < \omega < 2$ will lead to convergence, but we are generally interested in faster

Cubic Splines, Chapter 6

Cubic spline interpolation is among one of the mostly used methods for interpolating between data points where the arguments are organized as ascending series. In the library program we supply such a function, based on the so-called cubic spline method to be described below.

A spline function consists of polynomial pieces defined on subintervals. The different subintervals are connected via various continuity relations.

Assume we have at our disposal $n + 1$ points x_0, x_1, \dots, x_n arranged so that $x_0 < x_1 < x_2 < \dots < x_{n-1} < x_n$ (such points are called knots). A spline function s of degree k with $n + 1$ knots is defined as follows

- On every subinterval $[x_{i-1}, x_i]$ s is a polynomial of degree $\leq k$.
- s has $k - 1$ continuous derivatives in the whole interval $[x_0, x_n]$.

Splines

As an example, consider a spline function of degree $k = 1$ defined as follows

$$s(x) = \begin{cases} s_0(x) = a_0x + b_0 & x \in [x_0, x_1) \\ s_1(x) = a_1x + b_1 & x \in [x_1, x_2) \\ \dots & \dots \\ s_{n-1}(x) = a_{n-1}x + b_{n-1} & x \in [x_{n-1}, x_n] \end{cases}$$

In this case the polynomial consists of series of straight lines connected to each other at every endpoint. The number of continuous derivatives is then $k - 1 = 0$, as expected when we deal with straight lines. Such a polynomial is quite easy to construct given $n + 1$ points x_0, x_1, \dots, x_n and their corresponding function values.

Splines

The most commonly used spline function is the one with $k = 3$, the so-called cubic spline function. Assume that we have in addition to the $n + 1$ knots a series of functions values $y_0 = f(x_0), y_1 = f(x_1), \dots, y_n = f(x_n)$. By definition, the polynomials s_{i-1} and s_i are then supposed to interpolate the same point i , i.e.,

$$s_{i-1}(x_i) = y_i = s_i(x_i),$$

with $1 \leq i \leq n - 1$. In total we have n polynomials of the type

$$s_i(x) = a_{i0} + a_{i1}x + a_{i2}x^2 + a_{i3}x^3,$$

yielding $4n$ coefficients to determine.

Splines

Every subinterval provides in addition the $2n$ conditions

$$y_i = s(x_i),$$

and

$$s(x_{i+1}) = y_{i+1},$$

to be fulfilled. If we also assume that s' and s'' are continuous, then

$$s'_{i-1}(x_i) = s'_i(x_i),$$

yields $n - 1$ conditions. Similarly,

$$s''_{i-1}(x_i) = s''_i(x_i),$$

results in additional $n - 1$ conditions. In total we have $4n$ coefficients and $4n - 2$ equations to determine them, leaving us with 2 degrees of freedom to be determined.

Splines

Using the last equation we define two values for the second derivative, namely

$$s''_i(x_i) = f_i,$$

and

$$s''_i(x_{i+1}) = f_{i+1},$$

and setting up a straight line between f_i and f_{i+1} we have

$$s''_i(x) = \frac{f_i}{x_{i+1} - x_i}(x_{i+1} - x) + \frac{f_{i+1}}{x_{i+1} - x_i}(x - x_i),$$

and integrating twice one obtains

$$s_i(x) = \frac{f_i}{6(x_{i+1} - x_i)}(x_{i+1} - x)^3 + \frac{f_{i+1}}{6(x_{i+1} - x_i)}(x - x_i)^3 + c(x - x_i) + d(x_{i+1} - x)$$

Splines

Using the conditions $s_i(x_i) = y_i$ and $s_i(x_{i+1}) = y_{i+1}$ we can in turn determine the constants c and d resulting in

$$s_i(x) = \frac{f_i}{6(x_{i+1} - x_i)}(x_{i+1} - x)^3 + \frac{f_{i+1}}{6(x_{i+1} - x_i)}(x - x_i)^3 + \left(\frac{y_{i+1}}{x_{i+1} - x_i} - \frac{f_{i+1}(x_{i+1} - x_i)}{6}\right)(x - x_i) + \left(\frac{y_i}{x_{i+1} - x_i} - \frac{f_i(x_{i+1} - x_i)}{6}\right)(x_{i+1} - x) \quad (1)$$

Splines

How to determine the values of the second derivatives f_i and f_{i+1} ? We use the continuity assumption of the first derivatives

$$s'_{i-1}(x_i) = s'_i(x_i),$$

and set $x = x_i$. Defining $h_i = x_{i+1} - x_i$ we obtain finally the following expression

$$h_{i-1}f_{i-1} + 2(h_i + h_{i-1})f_i + h_if_{i+1} = \frac{6}{h_i}(y_{i+1} - y_i) - \frac{6}{h_{i-1}}(y_i - y_{i-1}),$$

and introducing the shorthands $u_i = 2(h_i + h_{i-1})$, $v_i = \frac{6}{h_i}(y_{i+1} - y_i) - \frac{6}{h_{i-1}}(y_i - y_{i-1})$, we can reformulate the problem as a set of linear equations to be solved through e.g., Gaussian elimination

Splines

Gaussian elimination

$$\begin{bmatrix} u_1 & h_1 & 0 & \dots & & & \\ h_1 & u_2 & h_2 & 0 & \dots & & \\ 0 & h_2 & u_3 & h_3 & 0 & \dots & \\ \dots & \dots & \dots & \dots & \dots & \dots & \\ & & & 0 & h_{n-3} & u_{n-2} & h_{n-2} \\ & & & & 0 & h_{n-2} & u_{n-1} \end{bmatrix} \begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ \dots \\ f_{n-2} \\ f_{n-1} \end{bmatrix} = \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ \dots \\ v_{n-2} \\ v_{n-1} \end{bmatrix}$$

Note that this is a set of tridiagonal equations and can be solved through only $O(n)$ operations.

Splines

The functions supplied in the program library are *spline* and *splint*. In order to use cubic spline interpolation you need first to call

```
spline(double x[], double y[], int n, double yp1,
       double yp2, double y2[])
```

This function takes as input $x[0, \dots, n-1]$ and $y[0, \dots, n-1]$ containing a tabulation $y_i = f(x_i)$ with $x_0 < x_1 < \dots < x_{n-1}$ together with the first derivatives of $f(x)$ at x_0 and x_{n-1} , respectively. Then the function returns $y2[0, \dots, n-1]$ which contains the second derivatives of $f(x_i)$ at each point x_i . n is the number of points. This function provides the cubic spline interpolation for all subintervals and is called only once.

Splines

Thereafter, if you wish to make various interpolations, you need to call the function

```
splint(double x[], double y[], double y2a[], int n,
      double x, double *y)
```

which takes as input the tabulated values $x[0, \dots, n-1]$ and $y[0, \dots, n-1]$ and the output $y2a[0, \dots, n-1]$ from *spline*. It returns the value y corresponding to the point x .

Overview of week 38

Overview of week 38

Eigenvalue problems and project 2.

- Monday: Repetition from last week and discussion of project 2
- Jacobi's algorithm
- Tuesday:
- Householder's algorithm and Francis' algorithm
- Iterative methods for symmetric matrices: Lanczos' algorithm

Eigenvalue Solvers

Let us consider the matrix \mathbf{A} of dimension n . The eigenvalues of \mathbf{A} is defined through the matrix equation

$$\mathbf{A}\mathbf{x}^{(\nu)} = \lambda^{(\nu)}\mathbf{x}^{(\nu)},$$

where $\lambda^{(\nu)}$ are the eigenvalues and $\mathbf{x}^{(\nu)}$ the corresponding eigenvectors. Unless otherwise stated, when we use the wording eigenvector we mean the right eigenvector. The left eigenvector is defined as

$$\mathbf{x}^{(\nu)\top}\mathbf{A} = \lambda^{(\nu)}\mathbf{x}^{(\nu)\top}$$

The above right eigenvector problem is equivalent to a set of n equations with n unknowns x_i .

Eigenvalue Solvers

The eigenvalue problem can be rewritten as

$$(\mathbf{A} - \lambda^{(\nu)}\mathbf{I})\mathbf{x}^{(\nu)} = 0,$$

with \mathbf{I} being the unity matrix. This equation provides a solution to the problem if and only if the determinant is zero, namely

$$|\mathbf{A} - \lambda^{(\nu)}\mathbf{I}| = 0,$$

which in turn means that the determinant is a polynomial of degree n in λ and in general we will have n distinct zeros.

Eigenvalue Solvers

The eigenvalues of a matrix $\mathbf{A} \in \mathbb{C}^{n \times n}$ are thus the n roots of its characteristic polynomial

$$P(\lambda) = \det(\lambda \mathbf{I} - \mathbf{A}),$$

or

$$P(\lambda) = \prod_{i=1}^n (\lambda_i - \lambda).$$

The set of these roots is called the spectrum and is denoted as $\lambda(\mathbf{A})$. If $\lambda(\mathbf{A}) = \{\lambda_1, \lambda_2, \dots, \lambda_n\}$ then we have

$$\det(\mathbf{A}) = \lambda_1 \lambda_2 \dots \lambda_n,$$

and if we define the trace of \mathbf{A} as

$$\text{Tr}(\mathbf{A}) = \sum_{i=1}^n a_{ii}$$

Abel-Ruffini Impossibility Theorem

The Abel-Ruffini theorem (also known as Abel's impossibility theorem) states that there is no general solution in radicals to polynomial equations of degree five or higher.

The content of this theorem is frequently misunderstood. It does not assert that higher-degree polynomial equations are unsolvable. In fact, if the polynomial has real or complex coefficients, and we allow complex solutions, then every polynomial equation has solutions; this is the fundamental theorem of algebra. Although these solutions cannot always be computed exactly with radicals, they can be computed to any desired degree of accuracy using numerical methods such as the Newton-Raphson method or Laguerre method, and in this way they are no different from solutions to polynomial equations of the second, third, or fourth degrees.

The theorem only concerns the form that such a solution must take. The content of the theorem is that the solution of a higher-degree equation cannot in all cases be expressed in terms of the polynomial coefficients with a finite number of operations of

Abel-Ruffini Impossibility Theorem

The Abel-Ruffini theorem says that there are some fifth-degree equations whose solution cannot be so expressed. The equation $x^5 - x + 1 = 0$ is an example. Some other fifth degree equations can be solved by radicals, for example $x^5 - x^4 - x + 1 = 0$. The precise criterion that distinguishes between those equations that can be solved by radicals and those that cannot was given by Galois and is now part of Galois theory: a polynomial equation can be solved by radicals if and only if its Galois group is a solvable group.

Today, in the modern algebraic context, we say that second, third and fourth degree polynomial equations can always be solved by radicals because the symmetric groups S_2, S_3 and S_4 are solvable groups, whereas S_n is not solvable for $n \geq 5$.

Eigenvalue Solvers

In the present discussion we assume that our matrix is real and symmetric, that is $\mathbf{A} \in \mathbb{R}^{n \times n}$. The matrix \mathbf{A} has n eigenvalues $\lambda_1 \dots \lambda_n$ (distinct or not). Let \mathbf{D} be the diagonal matrix with the eigenvalues on the diagonal

$$\mathbf{D} = \begin{pmatrix} \lambda_1 & 0 & 0 & 0 & \dots & 0 & 0 \\ 0 & \lambda_2 & 0 & 0 & \dots & 0 & 0 \\ 0 & 0 & \lambda_3 & 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & \dots & \dots & \dots & \lambda_{n-1} & 0 \\ 0 & \dots & \dots & \dots & \dots & 0 & \lambda_n \end{pmatrix}.$$

If \mathbf{A} is real and symmetric then there exists a real orthogonal matrix \mathbf{S} such that

$$\mathbf{S}^T \mathbf{A} \mathbf{S} = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_n),$$

and for $j = 1 : n$ we have $\mathbf{A} \mathbf{S}(:, j) = \lambda_j \mathbf{S}(:, j)$.

Eigenvalue Solvers

To obtain the eigenvalues of $\mathbf{A} \in \mathbb{R}^{n \times n}$, the strategy is to perform a series of similarity transformations on the original matrix \mathbf{A} , in order to reduce it either into a diagonal form as above or into a tridiagonal form.

We say that a matrix \mathbf{B} is a similarity transform of \mathbf{A} if

$$\mathbf{B} = \mathbf{S}^T \mathbf{A} \mathbf{S}, \quad \text{where} \quad \mathbf{S}^T \mathbf{S} = \mathbf{S}^{-1} \mathbf{S} = \mathbf{I}.$$

The importance of a similarity transformation lies in the fact that the resulting matrix has the same eigenvalues, but the eigenvectors are in general different.

Eigenvalue Solvers

To prove this we start with the eigenvalue problem and a similarity transformed matrix \mathbf{B} .

$$\mathbf{A} \mathbf{x} = \lambda \mathbf{x} \text{ and } \mathbf{B} = \mathbf{S}^T \mathbf{A} \mathbf{S}.$$

We multiply the first equation on the left by \mathbf{S}^T and insert $\mathbf{S}^T \mathbf{S} = \mathbf{I}$ between \mathbf{A} and \mathbf{x} . Then we get

$$(\mathbf{S}^T \mathbf{A} \mathbf{S})(\mathbf{S}^T \mathbf{x}) = \lambda \mathbf{S}^T \mathbf{x}, \quad (2)$$

which is the same as

$$\mathbf{B}(\mathbf{S}^T \mathbf{x}) = \lambda (\mathbf{S}^T \mathbf{x}).$$

The variable λ is an eigenvalue of \mathbf{B} as well, but with eigenvector $\mathbf{S}^T \mathbf{x}$.

Eigenvalue Solvers

The basic philosophy is to

either apply subsequent similarity transformations (direct method) so that

$$\mathbf{S}_N^T \dots \mathbf{S}_1^T \mathbf{A} \mathbf{S}_1 \dots \mathbf{S}_N = \mathbf{D}, \quad (3)$$

or apply subsequent similarity transformations so that \mathbf{A} becomes tridiagonal (Householder) or upper/lower triangular (**QR** method). Thereafter, techniques for obtaining eigenvalues from tridiagonal matrices can be used.

or use so-called power methods

or use iterative methods (Krylov, Lanczos, Arnoldi). These methods are popular for huge matrix problems.

Gaussian Elimination and Tridiagonal matrices, project 1

In project 1 we rewrote our original differential equation in terms of a discretized equation with approximations to the derivatives as

$$-\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} = f(x_i, u(x_i)),$$

with $i = 1, 2, \dots, n$. We need to add to this system the two boundary conditions $u(a) = u_0$ and $u(b) = u_{n+1}$. If we define a matrix

$$\mathbf{A} = \frac{1}{h^2} \begin{pmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & \dots & \dots & \dots \\ & & & -1 & 2 & -1 \\ & & & & -1 & 2 \end{pmatrix}$$

and the corresponding vectors $\mathbf{u} = (u_1, u_2, \dots, u_n)^T$ and $\mathbf{f}(\mathbf{u}) = f(x_1, x_2, \dots, x_n, u_1, u_2, \dots, u_n)^T$ we can rewrite the differential equation including the boundary conditions as a system

Student project 2, part a-b)

We are first interested in the solution of the radial part of Schrödinger's equation for one electron. This equation reads

$$-\frac{\hbar^2}{2m} \left(\frac{1}{r^2} \frac{d}{dr} r^2 \frac{d}{dr} - \frac{l(l+1)}{r^2} \right) R(r) + V(r)R(r) = ER(r).$$

In our case $V(r)$ is the harmonic oscillator potential $(1/2)kr^2$ with $k = m\omega^2$ and E is the energy of the harmonic oscillator in three dimensions. The oscillator frequency is ω and the energies are

$$E_{nl} = \hbar\omega \left(2n + l + \frac{3}{2} \right),$$

with $n = 0, 1, 2, \dots$ and $l = 0, 1, 2, \dots$.

Since we have made a transformation to spherical coordinates it means that $r \in [0, \infty)$. The quantum number l is the orbital momentum of the electron.

Then we substitute $R(r) = (1/r)u(r)$ and obtain

Student project 2, part a-b)

We introduce a dimensionless variable $\rho = (1/\alpha)r$ where α is a constant with dimension length and get

$$-\frac{\hbar^2}{2m\alpha^2} \frac{d^2}{d\rho^2} u(\rho) + \left(V(\rho) + \frac{l(l+1)}{\rho^2} \frac{\hbar^2}{2m\alpha^2} \right) u(\rho) = Eu(\rho).$$

In project 2 we set $l = 0$. Inserting $V(\rho) = (1/2)k\alpha^2\rho^2$ we end up with

$$-\frac{\hbar^2}{2m\alpha^2} \frac{d^2}{d\rho^2} u(\rho) + \frac{k}{2} \alpha^2 \rho^2 u(\rho) = Eu(\rho).$$

We multiply thereafter with $2m\alpha^2/\hbar^2$ on both sides and obtain

$$-\frac{d^2}{d\rho^2} u(\rho) + \frac{mk}{\hbar^2} \alpha^4 \rho^2 u(\rho) = \frac{2m\alpha^2}{\hbar^2} Eu(\rho).$$

Student project 2, part a-b)

We have from the previous slide

$$-\frac{d^2}{d\rho^2} u(\rho) + \frac{mk}{\hbar^2} \alpha^4 \rho^2 u(\rho) = \frac{2m\alpha^2}{\hbar^2} Eu(\rho).$$

The constant α can now be fixed so that

$$\frac{mk}{\hbar^2} \alpha^4 = 1,$$

or

$$\alpha = \left(\frac{\hbar^2}{mk} \right)^{1/4}.$$

Defining

$$\lambda = \frac{2m\alpha^2}{\hbar^2} E,$$

we can rewrite Schrödinger's equation as

Student project 2, part a-b)

We use the by now standard expression for the second derivative of a function u

$$u'' = \frac{u(\rho+h) - 2u(\rho) + u(\rho-h)}{h^2} + O(h^2), \quad (4)$$

where h is our step. Next we define minimum and maximum values for the variable ρ , $\rho_{\min} = 0$ and ρ_{\max} , respectively. You need to check your results for the energies against different values ρ_{\max} , since we cannot set $\rho_{\max} = \infty$.

Student project 2, part a-b)

With a given number of steps, n_{step} , we then define the step h as

$$h = \frac{\rho_{\text{max}} - \rho_{\text{min}}}{n_{\text{step}}}.$$

Define an arbitrary value of ρ as

$$\rho_i = \rho_{\text{min}} + ih \quad i = 0, 1, 2, \dots, n_{\text{step}}$$

we can rewrite the Schrödinger equation for ρ_i as

$$-\frac{u(\rho_i + h) - 2u(\rho_i) + u(\rho_i - h)}{h^2} + \rho_i^2 u(\rho_i) = \lambda u(\rho_i),$$

or in a more compact way

$$-\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} + \rho_i^2 u_i = -\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} + V_i u_i = \lambda u_i,$$

where $V_i = \rho_i^2$ is the harmonic oscillator potential.

Student project 2, part a-b)

Define first the diagonal matrix element

$$d_i = \frac{2}{h^2} + V_i,$$

and the non-diagonal matrix element

$$e_i = -\frac{1}{h^2}.$$

In this case the non-diagonal matrix elements are given by a mere constant. *All non-diagonal matrix elements are equal.* With these definitions the Schrödinger equation takes the following form

$$d_i u_i + e_{i-1} u_{i-1} + e_{i+1} u_{i+1} = \lambda u_i,$$

where u_i is unknown. We can write the latter equation as a matrix eigenvalue problem

$$\begin{pmatrix} d_1 & e_1 & 0 & 0 & \dots & 0 & 0 \end{pmatrix} \begin{pmatrix} u_1 \end{pmatrix} = \lambda \begin{pmatrix} u_1 \end{pmatrix}$$

Student project 2, part a-b)

or if we wish to be more detailed, we can write the tridiagonal matrix as

$$\begin{pmatrix} \frac{2}{h^2} + V_1 & -\frac{1}{h^2} & 0 & 0 & \dots & 0 & 0 \\ -\frac{1}{h^2} & \frac{2}{h^2} + V_2 & -\frac{1}{h^2} & 0 & \dots & 0 & 0 \\ 0 & -\frac{1}{h^2} & \frac{2}{h^2} + V_3 & -\frac{1}{h^2} & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & \dots & \dots & \dots & \frac{2}{h^2} + V_{n_{\text{step}}-2} & -\frac{1}{h^2} \\ 0 & \dots & \dots & \dots & \dots & -\frac{1}{h^2} & \frac{2}{h^2} + V_{n_{\text{step}}} \end{pmatrix} \begin{pmatrix} u_1 \\ u_2 \\ u_3 \\ \dots \\ u_{n_{\text{step}}-2} \\ u_{n_{\text{step}}-1} \\ u_{n_{\text{step}}} \end{pmatrix} = \lambda \begin{pmatrix} u_1 \\ u_2 \\ u_3 \\ \dots \\ u_{n_{\text{step}}-2} \\ u_{n_{\text{step}}-1} \\ u_{n_{\text{step}}} \end{pmatrix} \quad (6)$$

Recall that the solutions are known via the boundary conditions at $i = n_{\text{step}}$ and at the other end point, that is for ρ_0 . The solution is zero in both cases.

Student project 2, part c)

We are going to study two electrons in a harmonic oscillator well which also interact via a repulsive Coulomb interaction. Let us start with the single-electron equation written as

$$-\frac{\hbar^2}{2m} \frac{d^2}{dr^2} u(r) + \frac{1}{2} k r^2 u(r) = E^{(1)} u(r),$$

where $E^{(1)}$ stands for the energy with one electron only. For two electrons with no repulsive Coulomb interaction, we have the following Schrödinger equation

$$\left(-\frac{\hbar^2}{2m} \frac{d^2}{dr_1^2} - \frac{\hbar^2}{2m} \frac{d^2}{dr_2^2} + \frac{1}{2} k r_1^2 + \frac{1}{2} k r_2^2 \right) u(r_1, r_2) = E^{(2)} u(r_1, r_2).$$

Student project 2, part c)

Note that we deal with a two-electron wave function $u(r_1, r_2)$ and two-electron energy $E^{(2)}$.

With no interaction this can be written out as the product of two single-electron wave functions, that is we have a solution on closed form.

We introduce the relative coordinate $\mathbf{r} = \mathbf{r}_1 - \mathbf{r}_2$ and the center-of-mass coordinate $\mathbf{R} = 1/2(\mathbf{r}_1 + \mathbf{r}_2)$. With these new coordinates, the radial Schrödinger equation reads

$$\left(-\frac{\hbar^2}{m} \frac{d^2}{dr^2} - \frac{\hbar^2}{4m} \frac{d^2}{dR^2} + \frac{1}{4} k r^2 + k R^2 \right) u(r, R) = E^{(2)} u(r, R).$$

Student project 2, part c)

The equations for r and R can be separated via the ansatz for the wave function $u(r, R) = \psi(r)\phi(R)$ and the energy is given by the sum of the relative energy E_r and the center-of-mass energy E_R , that is

$$E^{(2)} = E_r + E_R.$$

We add then the repulsive Coulomb interaction between two electrons, namely a term

$$V(r_1, r_2) = \frac{\beta e^2}{|\mathbf{r}_1 - \mathbf{r}_2|} = \frac{\beta e^2}{r},$$

with $\beta e^2 = 1.44 \text{ eVnm}$.

Student project 2, part c)

Adding this term, the r -dependent Schrödinger equation becomes

$$\left(-\frac{\hbar^2}{m} \frac{d^2}{dr^2} + \frac{1}{4}kr^2 + \frac{\beta e^2}{r}\right)\psi(r) = E_r\psi(r).$$

This equation is similar to the one we had previously in parts (a) and (b) and we introduce again a dimensionless variable $\rho = r/\alpha$. Repeating the same steps, we arrive at

$$-\frac{d^2}{d\rho^2}\psi(\rho) + \frac{mk}{4\hbar^2}\alpha^4\rho^2\psi(\rho) + \frac{m\alpha\beta e^2}{\rho\hbar^2}\psi(\rho) = \frac{m\alpha^2}{\hbar^2}E_r\psi(\rho).$$

Student project 2, part c)

We want to manipulate this equation further to make it as similar to that in (a) as possible. We define a 'frequency'

$$\omega_r^2 = \frac{1}{4} \frac{mk}{\hbar^2} \alpha^4,$$

and fix the constant α by requiring

$$\frac{m\alpha\beta e^2}{\hbar^2} = 1$$

or

$$\alpha = \frac{\hbar^2}{m\beta e^2}.$$

Student project 2, part c)

Defining

$$\lambda = \frac{m\alpha^2}{\hbar^2} E,$$

we can rewrite Schrödinger's equation as

$$-\frac{d^2}{d\rho^2}\psi(\rho) + \omega_r^2\rho^2\psi(\rho) + \frac{1}{\rho}\psi(\rho) = \lambda\psi(\rho).$$

Student project 2, part c)

We treat ω_r as a parameter which reflects the strength of the oscillator potential.

Here we will study the cases $\omega_r = 0.01$, $\omega_r = 0.5$, $\omega_r = 1$, and $\omega_r = 5$ for the ground state only, that is the lowest-lying state.

With no repulsive Coulomb interaction you should get a result which corresponds to the relative energy of a non-interacting system. Make sure your results are stable as functions of ρ_{\max} and the number of steps.

We are only interested in the ground state with $l = 0$. We omit the center-of-mass energy.

For specific oscillator frequencies, the above equation has closed answers, see the article by M. Taut, Phys. Rev. A 48, 3561 - 3566 (1993). The article can be retrieved from the following web address http://prola.aps.org/abstract/PRA/v48/i5/p3561_1.

Eigenvalue Solvers

One speaks normally of two main approaches to solving the eigenvalue problem.

The first is the formal method, involving determinants and the *characteristic polynomial*. This proves how many eigenvalues there are, and is the way most of you learned about how to solve the eigenvalue problem, but for matrices of dimensions greater than 2 or 3, it is rather impractical.

The other general approach is to use similarity or unitary transformations to reduce a matrix to diagonal form. Almost always this is done in two steps: first reduce to for example a *tridiagonal* form, and then to diagonal form. The main algorithms we will discuss in detail, Jacobi's and Householder's (so-called direct method) and Lanczos algorithms (an iterative method), follow this methodology.

Diagonalization methods, direct methods

Direct or non-iterative methods.

...require for matrices of dimensionality $n \times n$ typically $O(n^3)$ operations. These methods are normally called standard methods and are used for dimensionalities $n \sim 10^5$ or smaller. A brief

historical overview

Year	n	
1950	$n = 20$	(Wilkinson)
1965	$n = 200$	(Forsythe et al.)
1980	$n = 2000$	Lapack
1995	$n = 20000$	Lapack
2012	$n \sim 10^5$	Lapack

shows that in the course of 60 years the dimension that direct diagonalization methods can handle has increased by almost a factor of 10^4 . However, it pales beside the progress achieved by computer hardware, from flops to petaflops, a factor of almost 10^{15} . We see clearly played out in history the $O(n^3)$ bottleneck of direct matrix algorithms. Sloppily speaking, when $n \sim 10^4$ is cubed we have $O(10^{12})$ operations, which is smaller than the 10^{15} increase in flops.

Diagonalization methods

Why iterative methods?

If the matrix to diagonalize is large and sparse, direct methods simply become impractical, also because many of the direct methods tend to destroy sparsity. As a result large dense matrices may arise during the diagonalization procedure. The idea behind iterative methods is to project the n -dimensional problem in smaller spaces, so-called Krylov subspaces. Given a matrix \hat{A} and a vector \hat{v} , the associated Krylov sequences of vectors (and thereby subspaces) $\hat{v}, \hat{A}\hat{v}, \hat{A}^2\hat{v}, \hat{A}^3\hat{v}, \dots$, represent successively larger Krylov subspaces.

Matrix	$\hat{A}\hat{x} = \hat{b}$	$\hat{A}\hat{x} = \lambda\hat{x}$
$\hat{A} = \hat{A}^*$	Conjugate gradient	Lanczos
$\hat{A} \neq \hat{A}^*$	GMRES etc	Arnoldi

Important Matrix and vector handling packages

The Numerical Recipes codes have been rewritten in Fortran 90/95 and C/C++ by us. The original source codes are taken from the widely used software package LAPACK, which follows two other popular packages developed in the 1970s, namely EISPACK and LINPACK.

- LINPACK: package for linear equations and least square problems.
- LAPACK: package for solving symmetric, unsymmetric and generalized eigenvalue problems. From LAPACK's website <http://www.netlib.org> it is possible to download for free all source codes from this library. Both C/C++ and Fortran versions are available.
- BLAS (I, II and III): (Basic Linear Algebra Subprograms) are routines that provide standard building blocks for performing basic vector and matrix operations. Blas I is vector operations, II vector-matrix operations and III matrix-matrix operations. Highly parallelized and efficient codes, all available for download from <http://www.netlib.org>.

Eigenvalue Solvers, Jacobi

Consider an $(n \times n)$ orthogonal transformation matrix

$$S = \begin{pmatrix} 1 & 0 & \dots & 0 & 0 & \dots & 0 & 0 \\ 0 & 1 & \dots & 0 & 0 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & \cos \theta & 0 & \dots & 0 & \sin \theta \\ 0 & 0 & \dots & 0 & 1 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & \dots & -\sin \theta & 0 & \dots & 1 & \cos \theta \\ 0 & 0 & \dots & 0 & \dots & \dots & 0 & 1 \end{pmatrix}$$

with property $S^T = S^{-1}$. It performs a plane rotation around an angle θ in the Euclidean n -dimensional space.

Eigenvalue Solvers, Jacobi

It means that its matrix elements that differ from zero are given by

$$s_{kk} = s_{ll} = \cos \theta, s_{kl} = -s_{lk} = -\sin \theta, s_{ij} = -s_{ji} = 1 \quad i \neq k \quad i \neq l,$$

A similarity transformation

$$B = S^T A S,$$

results in

$$\begin{aligned} b_{ik} &= a_{ik} \cos \theta - a_{il} \sin \theta, i \neq k, i \neq l \\ b_{il} &= a_{il} \cos \theta + a_{ik} \sin \theta, i \neq k, i \neq l \\ b_{kk} &= a_{kk} \cos^2 \theta - 2a_{kl} \cos \theta \sin \theta + a_{ll} \sin^2 \theta \\ b_{ll} &= a_{ll} \cos^2 \theta + 2a_{kl} \cos \theta \sin \theta + a_{kk} \sin^2 \theta \\ b_{kl} &= (a_{kk} - a_{ll}) \cos \theta \sin \theta + a_{kl} (\cos^2 \theta - \sin^2 \theta) \end{aligned}$$

Eigenvalue Solvers, Jacobi

The main idea is thus to reduce systematically the norm of the off-diagonal matrix elements of a matrix A

$$\text{off}(A) = \sqrt{\sum_{i=1}^n \sum_{j=1, j \neq i}^n a_{ij}^2}.$$

To demonstrate the algorithm, we consider the simple 2×2 similarity transformation of the full matrix. The matrix is symmetric, we single out $1 \leq k < l \leq n$ and use the abbreviations $c = \cos \theta$ and $s = \sin \theta$ to obtain

$$\begin{pmatrix} b_{kk} & 0 \\ 0 & b_{ll} \end{pmatrix} = \begin{pmatrix} c & -s \\ s & c \end{pmatrix} \begin{pmatrix} a_{kk} & a_{kl} \\ a_{lk} & a_{ll} \end{pmatrix} \begin{pmatrix} c & s \\ -s & c \end{pmatrix}.$$

Eigenvalue Solvers, Jacobi

We require that the non-diagonal matrix elements $b_{kl} = b_{lk} = 0$, implying that

$$a_{kl}(c^2 - s^2) + (a_{kk} - a_{ll})cs = b_{kl} = 0.$$

If $a_{kl} = 0$ one sees immediately that $\cos \theta = 1$ and $\sin \theta = 0$.

The Frobenius norm of an orthogonal transformation is always preserved. The Frobenius norm is defined as

$$\|A\|_F = \sqrt{\sum_{i=1}^n \sum_{j=1}^n |a_{ij}|^2}.$$

This means that for our 2×2 case we have

$$2a_{kl}^2 + a_{kk}^2 + a_{ll}^2 = b_{kk}^2 + b_{ll}^2,$$

which leads to

Eigenvalue Solvers, Jacobi

Defining the quantities $\tan \theta = t = s/c$ and

$$\cot 2\theta = \tau = \frac{a_{ll} - a_{kk}}{2a_{kl}},$$

we obtain the quadratic equation (using $\cot 2\theta = 1/2(\cot \theta - \tan \theta)$)

$$t^2 + 2\tau t - 1 = 0,$$

resulting in

$$t = -\tau \pm \sqrt{1 + \tau^2},$$

and c and s are easily obtained via

$$c = \frac{1}{\sqrt{1 + t^2}},$$

and $s = tc$. Choosing t to be the smaller of the roots ensures that $|\theta| \leq \pi/4$ and has the effect of minimizing the difference between the matrices \mathbf{B} and \mathbf{A} since

Eigenvalue Solvers, Jacobi algo

Step 1. Choose a tolerance ϵ , making it a small number, typically 10^{-8} or smaller.

Step 2. Setup a while-test where one compares the norm of the newly computed off-diagonal matrix elements

$$\text{off}(\mathbf{A}) = \sqrt{\sum_{i=1}^n \sum_{j=1, j \neq i}^n a_{ij}^2} > \epsilon.$$

Step 3. Now choose the matrix elements a_{kl} so that we have those with largest value, that is $|a_{kl}| = \max_{i \neq j} |a_{ij}|$.

Step 4. Compute thereafter $\tau = (a_{ll} - a_{kk})/2a_{kl}$, $\tan \theta$, $\cos \theta$ and $\sin \theta$.

Step 5. Compute thereafter the similarity transformation for this set of values (k, l) , obtaining the new matrix

$$\mathbf{B} = \mathbf{S}(k, l, \theta)^T \mathbf{A} \mathbf{S}(k, l, \theta).$$

Step 6. Compute the new norm of the off-diagonal matrix

Jacobi's method, an example to convince you about the algorithm

We specialize to a symmetric 3×3 matrix \mathbf{A} . We start the process as follows (assuming that $a_{23} = a_{32}$ is the largest non-diagonal) with $c = \cos \theta$ and $s = \sin \theta$

$$\mathbf{B} = \begin{pmatrix} 1 & 0 & 0 \\ 0 & c & -s \\ 0 & s & c \end{pmatrix} \begin{pmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{pmatrix} \begin{pmatrix} 1 & 0 & 0 \\ 0 & c & s \\ 0 & -s & c \end{pmatrix}.$$

We will choose the angle θ in order to have $a_{23} = a_{32} = 0$. We get (symmetric matrix)

$$\mathbf{B} = \begin{pmatrix} a_{11} & a_{12}c - a_{13}s & a_{12}s + a_{13}c \\ a_{12}c - a_{13}s & a_{22}c^2 + a_{33}s^2 - 2a_{23}sc & (a_{22} - a_{33})sc + a_{23}(c^2 - s^2) \\ a_{12}s + a_{13}c & (a_{22} - a_{33})sc + a_{23}(c^2 - s^2) & a_{22}s^2 + a_{33}c^2 + 2a_{23}sc \end{pmatrix}.$$

Note that a_{11} is unchanged! As it should.

Jacobi's method, an example to convince you about the algorithm

We have

$$\mathbf{B} = \begin{pmatrix} a_{11} & a_{12}c - a_{13}s & a_{12}s + a_{13}c \\ a_{12}c - a_{13}s & a_{22}c^2 + a_{33}s^2 - 2a_{23}sc & (a_{22} - a_{33})sc + a_{23}(c^2 - s^2) \\ a_{12}s + a_{13}c & (a_{22} - a_{33})sc + a_{23}(c^2 - s^2) & a_{22}s^2 + a_{33}c^2 + 2a_{23}sc \end{pmatrix}$$

or

$$\begin{aligned} b_{11} &= a_{11} \\ b_{12} &= a_{12} \cos \theta - a_{13} \sin \theta, 1 \neq 2, 1 \neq 3 \\ b_{13} &= a_{13} \cos \theta + a_{12} \sin \theta, 1 \neq 2, 1 \neq 3 \\ b_{22} &= a_{22} \cos^2 \theta - 2a_{23} \cos \theta \sin \theta + a_{33} \sin^2 \theta \\ b_{33} &= a_{33} \cos^2 \theta + 2a_{23} \cos \theta \sin \theta + a_{22} \sin^2 \theta \\ b_{23} &= (a_{22} - a_{33}) \cos \theta \sin \theta + a_{23}(\cos^2 \theta - \sin^2 \theta) \end{aligned}$$

We will fix the angle θ so that $b_{23} = 0$.

Jacobi's method, an example to convince you about the algorithm

We get then a new matrix

$$\mathbf{B} = \begin{pmatrix} b_{11} & b_{12} & b_{13} \\ b_{12} & b_{22} & 0 \\ b_{13} & 0 & b_{33} \end{pmatrix}.$$

We repeat then assuming that b_{12} is the largest non-diagonal matrix element and get a new matrix

$$\mathbf{C} = \begin{pmatrix} c & -s & 0 \\ s & c & 0 \\ 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} b_{11} & b_{12} & b_{13} \\ b_{12} & b_{22} & 0 \\ b_{13} & 0 & b_{33} \end{pmatrix} \begin{pmatrix} c & s & 0 \\ -s & c & 0 \\ 0 & 0 & 1 \end{pmatrix}.$$

We continue this process till all non-diagonal matrix elements are zero (ideally). You will notice that performing the above operations that the matrix element b_{23} which was previous zero becomes different from zero. This is one of the problems which slows down the jacobi procedure.

Jacobi's method, an example to convince you about the algorithm

The more general expression for the new matrix elements are

$$\begin{aligned} b_{ii} &= a_{ii}, i \neq k, i \neq l \\ b_{ik} &= a_{ik} \cos \theta - a_{il} \sin \theta, i \neq k, i \neq l \\ b_{il} &= a_{il} \cos \theta + a_{ik} \sin \theta, i \neq k, i \neq l \\ b_{kk} &= a_{kk} \cos^2 \theta - 2a_{kl} \cos \theta \sin \theta + a_{ll} \sin^2 \theta \\ b_{ll} &= a_{ll} \cos^2 \theta + 2a_{kl} \cos \theta \sin \theta + a_{kk} \sin^2 \theta \\ b_{kl} &= (a_{kk} - a_{ll}) \cos \theta \sin \theta + a_{kl}(\cos^2 \theta - \sin^2 \theta) \end{aligned}$$

This is what we will need to code.

Jacobi code example

Main part

```
// we have defined a matrix A and a matrix R for the eigenvector,
// The final matrix R has the eigenvectors in its row elements, it
// for the diagonal elements in the beginning, zero else.
....
double tolerance = 1.0E-10;
int iterations = 0;
while ( maxnondiag > tolerance && iterations <= maxiter)
{
    int p, q;
    maxnondiag = offdiag(A, p, q, n);
    Jacobi_rotate(A, R, p, q, n);
    iterations++;
}
...
```

Jacobi code example

Finding the max nondiagonal element

```
// the offdiag function
double offdiag(double **A, int p, int q, int n);
{
    double max;
    for (int i = 0; i < n; ++i)
    {
        for (int j = i+1; j < n; ++j)
        {
            double aij = fabs(A[i][j]);
            if (aij > max)
            {
                max = aij; p = i; q = j;
            }
        }
    }
    return max;
}
...
```

Jacobi code example

Finding the new matrix elements

```
void Jacobi_rotate ( double ** A, double ** R, int k, int l, int n)
{
    double s, c;
    if ( A[k][l] != 0.0 ) {
        double t, tau;
        tau = (A[l][l] - A[k][k]) / (2*A[k][l]);
        if ( tau >= 0 ) {
            t = 1.0 / (tau + sqrt(1.0 + tau*tau));
        } else {
            t = -1.0 / (-tau + sqrt(1.0 + tau*tau));
        }

        c = 1 / sqrt(1+t*t);
        s = c*t;
    } else {
        c = 1.0;
        s = 0.0;
    }
}
```

Jacobi code example

```
double a_kk, a_ll, a_ik, a_il, r_ik, r_il;
a_kk = A[k][k];
a_ll = A[l][l];
A[k][k] = c*c*a_kk - 2.0*c*s*A[k][l] + s*s*a_ll;
A[l][l] = s*s*a_kk + 2.0*c*s*A[k][l] + c*c*a_ll;
A[k][l] = 0.0; // hard-coding non-diagonal elements by hand
A[l][k] = 0.0; // same here
for (int i = 0; i < n; i++) {
    if ( i != k && i != l ) {
        a_ik = A[i][k];
        a_il = A[i][l];
        A[i][k] = c*a_ik - s*a_il;
        A[k][i] = A[i][k];
        A[i][l] = c*a_il + s*a_ik;
        A[l][i] = A[i][l];
    }
}
```

Jacobi code example

and finally the new eigenvectors

```
r_ik = R[i][k];
r_il = R[i][l];

R[i][k] = c*r_ik - s*r_il;
R[i][l] = c*r_il + s*r_ik;

return;
} // end of function jacobi_rotate
```

Week 39

Overview of week 39

Eigenvalue problems and differential equations (change of plans).

- Monday: Brief repetition from last week, with discussion of project 2.
- Discussion of Householder's and Francis' algorithm (not finished last week)
- Discussion of Lanczos' method (also optional part of project 2)
- Tuesday:
 - Introduction to differential equations
 - Differential equations, general properties.
 - Runge-Kutta methods The material on differential equations is covered by chapters 8, 9 and 10.

Eigenvalue Solvers, Householder

The first step consists in finding an orthogonal matrix \mathbf{S} which is the product of $(n-2)$ orthogonal matrices

$$\mathbf{S} = \mathbf{S}_1 \mathbf{S}_2 \dots \mathbf{S}_{n-2},$$

each of which successively transforms one row and one column of \mathbf{A} into the required tridiagonal form. Only $n-2$ transformations are required, since the last two elements are already in tridiagonal form. In order to determine each \mathbf{S}_i let us see what happens after the first multiplication, namely,

$$\mathbf{S}_1^T \mathbf{A} \mathbf{S}_1 = \begin{pmatrix} a_{11} & e_1 & 0 & 0 & \dots & 0 & 0 \\ e_1 & a'_{22} & a'_{23} & \dots & \dots & \dots & a'_{2n} \\ 0 & a'_{32} & a'_{33} & \dots & \dots & \dots & a'_{3n} \\ 0 & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & a'_{n2} & a'_{n3} & \dots & \dots & \dots & a'_{nn} \end{pmatrix}$$

where the primed quantities represent a matrix \mathbf{A}' of dimension

Eigenvalue Solvers, Householder

The factor e_1 is a possibly non-vanishing element. The next transformation produced by \mathbf{S}_2 has the same effect as \mathbf{S}_1 but now on the submatrix \mathbf{A}' only

$$(\mathbf{S}_1 \mathbf{S}_2)^T \mathbf{A} \mathbf{S}_1 \mathbf{S}_2 = \begin{pmatrix} a_{11} & e_1 & 0 & 0 & \dots & 0 & 0 \\ e_1 & a'_{22} & e_2 & 0 & \dots & \dots & 0 \\ 0 & e_2 & a''_{33} & \dots & \dots & \dots & a''_{3n} \\ 0 & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & a''_{n3} & \dots & \dots & \dots & a''_{nn} \end{pmatrix}$$

Note that the effective size of the matrix on which we apply the transformation reduces for every new step. In the previous Jacobi method each similarity transformation is in principle performed on the full size of the original matrix.

Eigenvalue Solvers, Householder

After a series of such transformations, we end with a set of diagonal matrix elements

$$a_{11}, a'_{22}, a''_{33} \dots a^{(n-1)}_{nn},$$

and off-diagonal matrix elements

$$e_1, e_2, e_3, \dots, e_{n-1}.$$

The resulting matrix reads

$$\mathbf{S}^T \mathbf{A} \mathbf{S} = \begin{pmatrix} a_{11} & e_1 & 0 & 0 & \dots & 0 & 0 \\ e_1 & a'_{22} & e_2 & 0 & \dots & 0 & 0 \\ 0 & e_2 & a''_{33} & e_3 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & \dots & \dots & \dots & a^{(n-1)}_{n-2} & e_{n-1} \\ 0 & \dots & \dots & \dots & \dots & e_{n-1} & a^{(n-1)}_{n-1} \end{pmatrix}.$$

Eigenvalue Solvers, Householder

It remains to find a recipe for determining the transformation \mathbf{S}_n . We illustrate the method for \mathbf{S}_1 which we assume takes the form

$$\mathbf{S}_1 = \begin{pmatrix} 1 & \mathbf{0}^T \\ \mathbf{0} & \mathbf{P} \end{pmatrix},$$

with $\mathbf{0}^T$ being a zero row vector, $\mathbf{0}^T = \{0, 0, \dots\}$ of dimension $(n-1)$. The matrix \mathbf{P} is symmetric with dimension $((n-1) \times (n-1))$ satisfying $\mathbf{P}^2 = \mathbf{I}$ and $\mathbf{P}^T = \mathbf{P}$. A possible choice which fulfills the latter two requirements is

$$\mathbf{P} = \mathbf{I} - 2\mathbf{u}\mathbf{u}^T,$$

where \mathbf{I} is the $(n-1)$ unity matrix and \mathbf{u} is an $n-1$ column vector with norm $\mathbf{u}^T \mathbf{u}$ (inner product).

Eigenvalue Solvers, Householder

Note that $\mathbf{u}\mathbf{u}^T$ is an outer product giving a matrix of dimension $((n-1) \times (n-1))$. Each matrix element of \mathbf{P} then reads

$$P_{ij} = \delta_{ij} - 2u_i u_j,$$

where i and j range from 1 to $n-1$. Applying the transformation \mathbf{S}_1 results in

$$\mathbf{S}_1^T \mathbf{A} \mathbf{S}_1 = \begin{pmatrix} a_{11} & (\mathbf{P}\mathbf{v})^T \\ \mathbf{P}\mathbf{v} & \mathbf{A}' \end{pmatrix},$$

where $\mathbf{v}^T = \{a_{21}, a_{31}, \dots, a_{n1}\}$ and \mathbf{P} must satisfy $(\mathbf{P}\mathbf{v})^T = \{k, 0, 0, \dots\}$. Then

$$\mathbf{P}\mathbf{v} = \mathbf{v} - 2\mathbf{u}(\mathbf{u}^T \mathbf{v}) = k\mathbf{e}, \quad (7)$$

with $\mathbf{e}^T = \{1, 0, 0, \dots, 0\}$.

Eigenvalue Solvers, Householder

Solving the latter equation gives us \mathbf{u} and thus the needed transformation \mathbf{P} . We do first however need to compute the scalar k by taking the scalar product of the last equation with its transpose and using the fact that $\mathbf{P}^2 = \mathbf{I}$. We get then

$$(\mathbf{P}\mathbf{v})^T \mathbf{P}\mathbf{v} = k^2 = \mathbf{v}^T \mathbf{v} = |\mathbf{v}|^2 = \sum_{i=2}^n a_{i1}^2,$$

which determines the constant $k = \pm v$.

Eigenvalue Solvers, Householder

Now we can rewrite Eq. (7) as

$$\mathbf{v} - k\mathbf{e} = 2\mathbf{u}(\mathbf{u}^T \mathbf{v}),$$

and taking the scalar product of this equation with itself and obtain

$$2(\mathbf{u}^T \mathbf{v})^2 = (\mathbf{v}^2 \pm a_{21}v), \quad (8)$$

which finally determines

$$\mathbf{u} = \frac{\mathbf{v} - k\mathbf{e}}{2(\mathbf{u}^T \mathbf{v})}.$$

In solving Eq. (8) great care has to be exercised so as to choose those values which make the right-hand largest in order to avoid loss of numerical precision. The above steps are then repeated for every transformations till we have a tridiagonal matrix suitable for obtaining the eigenvalues.

Eigenvalue Solvers, Householder, brute force

Our Householder transformation has given us a tridiagonal matrix. We discuss here how one can use Jacobi's iterative procedure to obtain the eigenvalues. Let us specialize to a 4×4 matrix. The tridiagonal matrix takes the form

$$\mathbf{A} = \begin{pmatrix} d_1 & e_1 & 0 & 0 \\ e_1 & d_2 & e_2 & 0 \\ 0 & e_2 & d_3 & e_3 \\ 0 & 0 & e_3 & d_4 \end{pmatrix}.$$

As a first observation, if any of the elements e_i are zero the matrix can be separated into smaller pieces before diagonalization. Specifically, if $e_1 = 0$ then d_1 is an eigenvalue.

Eigenvalue Solvers, Householder

Thus, let us introduce a transformation \mathbf{S}_1 which operates like

$$\mathbf{S}_1 = \begin{pmatrix} \cos \theta & 0 & 0 & \sin \theta \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ \cos \theta & 0 & 0 & \cos \theta \end{pmatrix}$$

Eigenvalue Solvers, Householder

Then the similarity transformation

$$\mathbf{S}_1^T \mathbf{A} \mathbf{S}_1 = \mathbf{A}' = \begin{pmatrix} d'_1 & e'_1 & 0 & 0 \\ e'_1 & d_2 & e_2 & 0 \\ 0 & e_2 & d_3 & e'_3 \\ 0 & 0 & e'_3 & d'_4 \end{pmatrix}$$

produces a matrix where the primed elements in \mathbf{A}' have been changed by the transformation whereas the unprimed elements are unchanged. If we now choose θ to give the element $a'_{21} = e'_1 = 0$ then we have the first eigenvalue $= a_{11} = d_1$. (This is actually what you are doing in project 2!!)

Eigenvalue Solvers, Householder's and Francis' algorithm

This procedure can be continued on the remaining three-dimensional submatrix for the next eigenvalue. Thus after few transformations we have the wanted diagonal form.

What we see here is just a special case of the more general procedure developed by Francis in two articles in 1961 and 1962. Using Jacobi's method is not very efficient either.

The algorithm is based on the so-called **QR** method (or just **QR**-algorithm). It follows from a theorem by Schur which states that any square matrix can be written out in terms of an orthogonal matrix \hat{Q} and an upper triangular matrix \hat{U} . Historically R was used instead of U since the wording right triangular matrix was first used.

Eigenvalue Solvers, Householder's and Francis' algorithm

The method is based on an iterative procedure similar to Jacobi's method, by a succession of planar rotations. For a tridiagonal matrix it is simple to carry out in principle, but complicated in detail!

Schur's theorem

$$\hat{A} = \hat{Q} \hat{U},$$

is used to rewrite any square matrix into a unitary matrix times an upper triangular matrix. We say that a square matrix is similar to a triangular matrix.

Householder's algorithm which we have derived is just a special case of the general Householder algorithm. For a symmetric square matrix we obtain a tridiagonal matrix.

There is a corollary to Schur's theorem which states that every Hermitian matrix is unitarily similar to a diagonal matrix.

Eigenvalue Solvers, Householder's and Francis' algorithm

It follows that we can define a new matrix

$$\hat{\lambda}\hat{Q} = \hat{Q}\hat{U}\hat{Q},$$

and multiply from the left with \hat{Q}^{-1} we get

$$\hat{Q}^{-1}\hat{\lambda}\hat{Q} = \hat{B} = \hat{U}\hat{Q},$$

where the matrix \hat{B} is a similarity transformation of \hat{A} and has the same eigenvalues as \hat{A} .

Eigenvalue Solvers, Householder's and Francis' algorithm

Suppose \hat{A} is the triangular matrix we obtained after the Householder transformation,

$$\hat{A} = \hat{Q}\hat{U},$$

and multiply from the left with \hat{Q}^{-1} resulting in

$$\hat{Q}^{-1}\hat{A} = \hat{U}.$$

Suppose that \hat{Q} consists of a series of planar Jacobi like rotations acting on sub blocks of \hat{A} so that all elements below the diagonal are zeroed out

$$\hat{Q} = \hat{R}_{12}\hat{R}_{23}\dots\hat{R}_{n-1,n}.$$

Eigenvalue Solvers, Householder's and Francis' algorithm

A transformation of the type \hat{R}_{12} looks like

$$\hat{R}_{12} = \begin{pmatrix} c & s & 0 & 0 & 0 & \dots & 0 & 0 & 0 \\ -s & c & 0 & 0 & 0 & \dots & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & \dots & 0 & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & 0 & 0 & \dots & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & 0 & 1 \end{pmatrix}$$

Eigenvalue Solvers, Householder's and Francis' algorithm

The matrix \hat{U} takes then the form

$$\hat{U} = \begin{pmatrix} x & x & x & 0 & 0 & \dots & 0 & 0 & 0 \\ 0 & x & x & x & 0 & \dots & 0 & 0 & 0 \\ 0 & 0 & x & x & x & \dots & 0 & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & 0 & 0 & \dots & x & x & x \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & x & x \\ 0 & 0 & 0 & 0 & 0 & \dots & 0 & 0 & x \end{pmatrix}$$

which has a second superdiagonal.

Eigenvalue Solvers, Householder's and Francis' algorithm

We have now found \hat{Q} and \hat{U} and this allows us to find the matrix \hat{B} which is, due to Schur's theorem, unitarily similar to a triangular matrix (upper in our case) since we have that

$$\hat{Q}^{-1}\hat{\lambda}\hat{Q} = \hat{B},$$

from Schur's theorem the matrix \hat{B} is triangular and the eigenvalues the same as those of \hat{A} and are given by the diagonal matrix elements of \hat{B} . Why?

Our matrix $\hat{B} = \hat{U}\hat{Q}$.

Another iterative procedure

The matrix \hat{A} is transformed into a tridiagonal form and the last step is to transform it into a diagonal matrix giving the eigenvalues on the diagonal.

The eigenvalues of a matrix can be obtained using the characteristic polynomial

$$P(\lambda) = \det(\lambda\mathbf{I} - \mathbf{A}) = \prod_{i=1}^n (\lambda_i - \lambda),$$

which rewritten in matrix form reads

$$P(\lambda) = \begin{pmatrix} d_1 - \lambda & e_1 & 0 & 0 & \dots & 0 & 0 \\ e_1 & d_2 - \lambda & e_2 & 0 & \dots & 0 & 0 \\ 0 & e_2 & d_3 - \lambda & e_3 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & \dots & \dots & \dots & \dots & d_{N_{\text{step}}-2} - \lambda & e_{N_{\text{step}}-1} \\ 0 & \dots & \dots & \dots & \dots & e_{N_{\text{step}}-1} & d_{N_{\text{step}}-1} \end{pmatrix}$$

Eigenvalue Solvers, Householder

We can solve this equation in an iterative manner. We let $P_k(\lambda)$ be the value of k subdeterminant of the above matrix of dimension $n \times n$. The polynomial $P_k(\lambda)$ is clearly a polynomial of degree k . Starting with $P_1(\lambda)$ we have $P_1(\lambda) = d_1 - \lambda$. The next polynomial reads $P_2(\lambda) = (d_2 - \lambda)P_1(\lambda) - e_1^2$. By expanding the determinant for $P_k(\lambda)$ in terms of the minors of the n th column we arrive at the recursion relation

$$P_k(\lambda) = (d_k - \lambda)P_{k-1}(\lambda) - e_{k-1}^2 P_{k-2}(\lambda).$$

Together with the starting values $P_1(\lambda)$ and $P_2(\lambda)$ and good root searching methods we arrive at an efficient computational scheme for finding the roots of $P_n(\lambda)$. However, for large matrices this algorithm is rather inefficient and time-consuming.

Eigenvalue Solvers, Householder functions

The programs which performs these transformations are
matrix $\mathbf{A} \rightarrow$ tridiagonal matrix \rightarrow diagonal matrix

C++:

```
void trd2(double S**a, int n, double d[], double e[])
void tqli(double d[], double[], int n, double S**Sz)
```

Fortran:

```
CALL tred2(a, n, d, e)
CALL tqli(d, e, n, z)
```

Using Lapack to solve linear algebra and eigenvalue problems

Suppose you wanted to solve a general system of linear equations $\mathbf{Ax} = \mathbf{b}$, where \mathbf{A} is an $n \times n$ square matrix and x and b are n -element column vectors. You opt to use the routine **dgesv**. The man page abstract obtained with

Lapack example

```
#include <iostream>
#define MAX 10
using namespace std;
int main(){
    // Values needed for dgesv
    int n;
    int nrhs = 1;
    double a[MAX][MAX];
    double b[1][MAX];
    int lda = MAX;
    int ldb = MAX;
    int ipiv[MAX];
    int info;
    // Other values
    int i,j;
```

Lapack example

```
// Read the values of the matrix
cout << "Enter n \n";
cin >> n;
cout << "On each line type a row of the matrix A followed by one
for(i = 0; i < n; i++){
    cout << "row " << i << " ";
    for(j = 0; j < n; j++)std::cin >> a[j][i];
    cin >> b[0][i];
```

Lapack example

```
// Solve the linear system
dgesv(n, nrhs, &a[0][0], lda, ipiv, &b[0][0], ldb, &info);
// Check for success
if(info == 0)
{
    // Write the answer
    cout << "The answer is\n";
    for(i = 0; i < n; i++)
        cout << "b[" << i << "]\t" << b[0][i] << "\n";
}
else
{
    // Write an error message
    cerr << "dgesv returned error " << info << "\n";
}
return info;
```

Lanczos' iteration

Basic features with a real symmetric matrix (and normally huge $n > 10^6$ and sparse) \hat{A} of dimension $n \times n$:

Lanczos' algorithm generates a sequence of real tridiagonal matrices T_k of dimension $k \times k$ with $k \leq n$, with the property that the extremal eigenvalues of T_k are progressively better estimates of \hat{A} ' extremal eigenvalues.

The method converges to the extremal eigenvalues.

The similarity transformation is

$$\hat{T} = \hat{Q}^T \hat{A} \hat{Q},$$

with the first vector $\hat{Q} \hat{e}_1 = \hat{q}_1$.

Lanczos' iteration

We are going to solve iteratively

$$\hat{T} = \hat{Q}^T \hat{A} \hat{Q},$$

with the first vector $\hat{Q} \hat{e}_1 = \hat{q}_1$. We can write out the matrix \hat{Q} in terms of its column vectors

$$\hat{Q} = [\hat{q}_1 \hat{q}_2 \dots \hat{q}_n].$$

Lanczos' iteration

The matrix

$$\hat{T} = \hat{Q}^T \hat{A} \hat{Q},$$

can be written as

$$\hat{T} = \begin{pmatrix} \alpha_1 & \beta_1 & 0 & \dots & \dots & 0 \\ \beta_1 & \alpha_2 & \beta_2 & 0 & \dots & 0 \\ 0 & \beta_2 & \alpha_3 & \beta_3 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & 0 \\ \dots & \dots & \dots & \beta_{n-2} & \alpha_{n-1} & \beta_{n-1} \\ 0 & \dots & \dots & 0 & \beta_{n-1} & \alpha_n \end{pmatrix}$$

Lanczos' iteration

Using the fact that $\hat{Q} \hat{Q}^T = \hat{I}$, we can rewrite

$$\hat{T} = \hat{Q}^T \hat{A} \hat{Q},$$

as

$$\hat{Q} \hat{T} = \hat{A} \hat{Q},$$

and if we equate columns (recall from the previous slide)

$$\hat{T} = \begin{pmatrix} \alpha_1 & \beta_1 & 0 & \dots & \dots & 0 \\ \beta_1 & \alpha_2 & \beta_2 & 0 & \dots & 0 \\ 0 & \beta_2 & \alpha_3 & \beta_3 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots & 0 \\ \dots & \dots & \dots & \beta_{n-2} & \alpha_{n-1} & \beta_{n-1} \\ 0 & \dots & \dots & 0 & \beta_{n-1} & \alpha_n \end{pmatrix}$$

we obtain

Lanczos' iteration

We have thus

$$\hat{A} \hat{q}_k = \beta_{k-1} \hat{q}_{k-1} + \alpha_k \hat{q}_k + \beta_k \hat{q}_{k+1},$$

with $\beta_0 \hat{q}_0 = 0$ for $k = 1 : n - 1$. Remember that the vectors \hat{q}_k are orthonormal and this implies

$$\alpha_k = \hat{q}_k^T \hat{A} \hat{q}_k,$$

and these vectors are called Lanczos vectors.

Lanczos' iteration, the algorithm

We have thus

$$\hat{A} \hat{q}_k = \beta_{k-1} \hat{q}_{k-1} + \alpha_k \hat{q}_k + \beta_k \hat{q}_{k+1},$$

with $\beta_0 \hat{q}_0 = 0$ for $k = 1 : n - 1$ and

$$\alpha_k = \hat{q}_k^T \hat{A} \hat{q}_k.$$

If

$$\hat{r}_k = (\hat{A} - \alpha_k \hat{I}) \hat{q}_k - \beta_{k-1} \hat{q}_{k-1},$$

is non-zero, then

$$\hat{q}_{k+1} = \hat{r}_k / \beta_k,$$

with $\beta_k = \pm \|\hat{r}_k\|_2$.

A simple implementation of the Lanczos algorithm

```

r_0 = q_1; beta_0=1; q_0=0; int k = 0;
while (beta_k != 0)
    q_{k+1} = r_k/beta_k
    k = k+1
    alpha_k = q_k^T A q_k
    r_k = (A-alpha_k I) q_k - beta_{k-1}q_{k-1}
    beta_k = || r_k ||_2
end while

```

Differential equations program

- Ordinary differential equations, Runge-Kutta method, chapter 8
- Ordinary differential equations with boundary conditions: one-variable equations to be solved by shooting and Green's function methods, chapter 9
- We can solve such equations by a finite difference scheme as well, turning the equation into an eigenvalue problem. Still one variable. Done in projects 1 and 2.
- If we have more than one variable, we need to solve partial differential equations, see Chapter 10
- Fourier transforms and Fast Fourier transforms if we get time. Project 3 deals with ordinary differential equations (most likely the solar system).

Differential Equations, chapter 8

The order of the ODE refers to the order of the derivative on the left-hand side in the equation

$$\frac{dy}{dt} = f(t, y). \quad (9)$$

This equation is of first order and f is an arbitrary function. A second-order equation goes typically like

$$\frac{d^2y}{dt^2} = f(t, \frac{dy}{dt}, y). \quad (10)$$

A well-known second-order equation is Newton's second law

$$m \frac{d^2x}{dt^2} = -kx, \quad (11)$$

where k is the force constant. ODE depend only on one variable

Differential Equations

partial differential equations like the time-dependent Schrödinger equation

$$i\hbar \frac{\partial \psi(\mathbf{x}, t)}{\partial t} = -\frac{\hbar^2}{2m} \left(\frac{\partial^2 \psi(\mathbf{r}, t)}{\partial x^2} + \frac{\partial^2 \psi(\mathbf{r}, t)}{\partial y^2} + \frac{\partial^2 \psi(\mathbf{r}, t)}{\partial z^2} \right) + V(\mathbf{x})\psi(\mathbf{x}, t), \quad (12)$$

may depend on several variables. In certain cases, like the above equation, the wave function can be factorized in functions of the separate variables, so that the Schrödinger equation can be rewritten in terms of sets of ordinary differential equations. These equations are discussed in chapter 10. Involve boundary conditions in addition to initial conditions.

Differential Equations

We distinguish also between linear and non-linear differential equation where e.g.,

$$\frac{dy}{dt} = g^3(t)y(t), \quad (13)$$

is an example of a linear equation, while

$$\frac{dy}{dt} = g^3(t)y(t) - g(t)y^2(t), \quad (14)$$

is a non-linear ODE.

Differential Equations

Another concept which dictates the numerical method chosen for solving an ODE, is that of initial and boundary conditions. To give an example, if we study white dwarf stars or neutron stars we will need to solve two coupled first-order differential equations, one for the total mass m and one for the pressure P as functions of ρ

$$\frac{dm}{dr} = 4\pi r^2 \rho(r)/c^2,$$

and

$$\frac{dP}{dr} = -\frac{Gm(r)}{r^2} \rho(r)/c^2.$$

where ρ is the mass-energy density. The initial conditions are dictated by the mass being zero at the center of the star, i.e., when $r = 0$, yielding $m(r = 0) = 0$. The other condition is that the pressure vanishes at the surface of the star.

In the solution of the Schrödinger equation for a particle in a potential, we may need to apply boundary conditions as well, such

Differential Equations

In many cases it is possible to rewrite a second-order differential equation in terms of two first-order differential equations. Consider again the case of Newton's second law in Eq. (11). If we define the position $x(t) = y^{(1)}(t)$ and the velocity $v(t) = y^{(2)}(t)$ as its derivative

$$\frac{dy^{(1)}(t)}{dt} = \frac{dx(t)}{dt} = y^{(2)}(t), \quad (15)$$

we can rewrite Newton's second law as two coupled first-order differential equations

$$m \frac{dy^{(2)}(t)}{dt} = -kx(t) = -ky^{(1)}(t), \quad (16)$$

and

$$\frac{dy^{(1)}(t)}{dt} = y^{(2)}(t). \quad (17)$$

Differential Equations, Finite Difference

These methods fall under the general class of one-step methods. The algorithm is rather simple. Suppose we have an initial value for the function $y(t)$ given by

$$y_0 = y(t = t_0). \quad (18)$$

We are interested in solving a differential equation in a region in space $[a, b]$. We define a step h by splitting the interval in N sub intervals, so that we have

$$h = \frac{b - a}{N}. \quad (19)$$

With this step and the derivative of y we can construct the next value of the function y at

$$y_1 = y(t_1 = t_0 + h), \quad (20)$$

and so forth.

Differential Equations

If the function is rather well-behaved in the domain $[a, b]$, we can use a fixed step size. If not, adaptive steps may be needed. Here we concentrate on fixed-step methods only. Let us try to generalize the above procedure by writing the step y_{i+1} in terms of the previous step y_i

$$y_{i+1} = y(t = t_i + h) = y(t_i) + h\Delta(t_i, y_i(t_i)) + O(h^{p+1}), \quad (21)$$

where $O(h^{p+1})$ represents the truncation error. To determine Δ , we Taylor expand our function y

$$y_{i+1} = y(t = t_i + h) = y(t_i) + h(y'(t_i) + \dots + y^{(p)}(t_i) \frac{h^{p-1}}{p!}) + O(h^{p+1}), \quad (22)$$

where we will associate the derivatives in the parenthesis with

$$\Delta(t_i, y_i(t_i)) = (y'(t_i) + \dots + y^{(p)}(t_i) \frac{h^{p-1}}{p!}). \quad (23)$$

Differential Equations

We define

$$y'(t_i) = f(t_i, y_i) \quad (24)$$

and if we truncate Δ at the first derivative, we have

$$y_{i+1} = y(t_i) + hf(t_i, y_i) + O(h^2), \quad (25)$$

which when complemented with $t_{i+1} = t_i + h$ forms the algorithm for the well-known Euler method. Note that at every step we make an approximation error of the order of $O(h^2)$, however the total error is the sum over all steps $N = (b - a)/h$, yielding thus a global error which goes like $NO(h^2) \approx O(h)$.

Differential Equations

To make Euler's method more precise we can obviously decrease h (increase N). However, if we are computing the derivative f numerically by e.g., the two-steps formula

$$f'_{2c}(x) = \frac{f(x+h) - f(x)}{h} + O(h),$$

we can enter into roundoff error problems when we subtract two almost equal numbers $f(x+h) - f(x) \approx 0$. Euler's method is not recommended for precision calculation, although it is handy to use in order to get a first view on how a solution may look like. As an example, consider Newton's equation rewritten in Eqs. (16) and (17). We define $y_0 = y^{(1)}(t = 0)$ and $v_0 = y^{(2)}(t = 0)$. The first steps in Newton's equations are then

$$y_1^{(1)} = y_0 + hv_0 + O(h^2) \quad (26)$$

and

Differential Equations

The Euler method is asymmetric in time, since it uses information about the derivative at the beginning of the time interval. This means that we evaluate the position at $y_1^{(1)}$ using the velocity at $y_0^{(2)} = v_0$. A simple variation is to determine $y_{n+1}^{(1)}$ using the velocity at $y_{n+1}^{(2)}$, that is (in a slightly more generalized form)

$$y_{n+1}^{(1)} = y_n^{(1)} + hy_{n+1}^{(2)} + O(h^2) \quad (28)$$

and

$$y_{n+1}^{(2)} = y_n^{(2)} + ha_n + O(h^2). \quad (29)$$

The acceleration a_n is a function of $a_n(y_n^{(1)}, y_n^{(2)}, t)$ and needs to be evaluated as well. This is the Euler-Cromer method.

Differential Equations

Let us then include the second derivative in our Taylor expansion. We have then

$$\Delta(t_i, y_i(t_i)) = f(t_i) + \frac{h}{2} \frac{df(t_i, y_i)}{dt} + O(h^3). \quad (30)$$

The second derivative can be rewritten as

$$y'' = f' = \frac{df}{dt} = \frac{\partial f}{\partial t} + \frac{\partial f}{\partial y} \frac{\partial y}{\partial t} = \frac{\partial f}{\partial t} + \frac{\partial f}{\partial y} f \quad (31)$$

and we can rewrite Eq. (22) as

$$y_{i+1} = y(t = t_i + h) = y(t_i) + hf(t_i) + \frac{h^2}{2} \left(\frac{\partial f}{\partial t} + \frac{\partial f}{\partial y} f \right) + O(h^3), \quad (32)$$

which has a local approximation error $O(h^3)$ and a global error $O(h^2)$.

Differential Equations

These approximations can be generalized by using the derivative f to arbitrary order so that we have

$$y_{i+1} = y(t = t_i + h) = y(t_i) + hf(t_i) + \dots + f^{(p-1)}(t_i, y_i) \frac{h^{p-1}}{p!} + O(h^{p+1}) \quad (33)$$

These methods, based on higher-order derivatives, are in general not used in numerical computation, since they rely on evaluating derivatives several times. Unless one has analytical expressions for these, the risk of roundoff errors is large.

Differential Equations

The most obvious improvements to Euler's and Euler-Cromer's algorithms, avoiding in addition the need for computing a second derivative, is the so-called midpoint method. We have then

$$y_{n+1}^{(1)} = y_n^{(1)} + \frac{h}{2} (y_{n+1}^{(2)} + y_n^{(2)}) + O(h^2) \quad (34)$$

and

$$y_{n+1}^{(2)} = y_n^{(2)} + ha_n + O(h^2), \quad (35)$$

yielding

$$y_{n+1}^{(1)} = y_n^{(1)} + hy_n^{(2)} + \frac{h^2}{2} a_n + O(h^3) \quad (36)$$

implying that the local truncation error in the position is now $O(h^3)$, whereas Euler's or Euler-Cromer's methods have a local error of $O(h^2)$.

Differential Equations

Thus, the midpoint method yields a global error with second-order accuracy for the position and first-order accuracy for the velocity. However, although these methods yield exact results for constant accelerations, the error increases in general with each time step.

One method that avoids this is the so-called half-step method. Here we define

$$y_{n+1/2}^{(2)} = y_{n-1/2}^{(2)} + ha_n + O(h^2), \quad (37)$$

and

$$y_{n+1}^{(1)} = y_n^{(1)} + hy_{n+1/2}^{(2)} + O(h^2). \quad (38)$$

Note that this method needs the calculation of $y_{1/2}^{(2)}$. This is done using e.g., Euler's method

$$y_{1/2}^{(2)} = y_0^{(2)} + ha_0 + O(h^2). \quad (39)$$

Differential Equations

Another method which one may encounter is the Euler-Richardson method with

$$y_{n+1}^{(2)} = y_n^{(2)} + ha_{n+1/2} + O(h^2), \quad (40)$$

and

$$y_{n+1}^{(1)} = y_n^{(1)} + hy_{n+1/2}^{(2)} + O(h^2). \quad (41)$$

The program program2.cpp includes all of the above methods.

Overview of week 40

Ordinary differential equations (ODEs) and Partial differential equations (PDEs).

- Monday: Repetition from last week
 - Runge-Kutta methods, with adaptive methods as well
 - Examples with codes
 - Tuesday:
 - Discussion of project 3
 - Diffusion equation, implicit, explicit (if we get there!!)
- Chapter 9, ODEs with boundary conditions will not be discussed.

Differential Equations, Runge-Kutta methods

Runge-Kutta (RK) methods are based on Taylor expansion formulae, but yield in general better algorithms for solutions of an ODE. The basic philosophy is that it provides an intermediate step in the computation of y_{i+1} .

To see this, consider first the following definitions

$$\frac{dy}{dt} = f(t, y), \quad (42)$$

and

$$y(t) = \int f(t, y) dt, \quad (43)$$

and

$$y_{i+1} = y_i + \int_{t_i}^{t_{i+1}} f(t, y) dt. \quad (44)$$

Differential Equations, Runge-Kutta methods

To demonstrate the philosophy behind RK methods, let us consider the second-order RK method, RK2. The first approximation consists in Taylor expanding $f(t, y)$ around the center of the integration interval t_i to t_{i+1} , i.e., at $t_i + h/2$, h being the step. Using the midpoint formula for an integral, defining $y(t_i + h/2) = y_{i+1/2}$ and $t_i + h/2 = t_{i+1/2}$, we obtain

$$\int_{t_i}^{t_{i+1}} f(t, y) dt \approx hf(t_{i+1/2}, y_{i+1/2}) + O(h^3). \quad (45)$$

This means in turn that we have

$$y_{i+1} = y_i + hf(t_{i+1/2}, y_{i+1/2}) + O(h^3). \quad (46)$$

Differential Equations, Runge-Kutta methods

However, we do not know the value of $y_{i+1/2}$. Here comes thus the next approximation, namely, we use Euler's method to approximate $y_{i+1/2}$. We have then

$$y_{(i+1/2)} = y_i + \frac{h}{2} \frac{dy}{dt} = y(t_i) + \frac{h}{2} f(t_i, y_i). \quad (47)$$

This means that we can define the following algorithm for the second-order Runge-Kutta method, RK2.

$$k_1 = hf(t_i, y_i), \quad (48)$$

$$k_2 = hf(t_{i+1/2}, y_i + k_1/2), \quad (49)$$

with the final value

$$y_{i+1} \approx y_i + k_2 + O(h^3). \quad (50)$$

Differential Equations, Runge-Kutta methods

The difference between the previous one-step methods is that we now need an intermediate step in our evaluation, namely $t_i + h/2 = t_{(i+1/2)}$ where we evaluate the derivative f . This involves more operations, but the gain is a better stability in the solution.

Differential Equations, Runge-Kutta methods

The fourth-order Runge-Kutta, RK4, which we will employ in the solution of various differential equations below, has the following algorithm

$$k_1 = hf(t_i, y_i), \quad (51)$$

$$k_2 = hf(t_i + h/2, y_i + k_1/2), \quad (52)$$

$$k_3 = hf(t_i + h/2, y_i + k_2/2) \quad (53)$$

$$k_4 = hf(t_i + h, y_i + k_3) \quad (54)$$

with the final value

$$y_{i+1} = y_i + \frac{1}{6} (k_1 + 2k_2 + 2k_3 + k_4). \quad (55)$$

Simple Example, Block tied to a Wall

Our first example is the classical case of simple harmonic oscillations, namely a block sliding on a horizontal frictionless surface. The block is tied to a wall with a spring. If the spring is not compressed or stretched too far, the force on the block at a given position x is

$$F = -kx.$$

Simple Example, Block tied to a Wall

The negative sign means that the force acts to restore the object to an equilibrium position. Newton's equation of motion for this idealized system is then

$$m \frac{d^2 x}{dt^2} = -kx,$$

or we could rephrase it as

$$\frac{d^2 x}{dt^2} = -\frac{k}{m}x = -\omega_0^2 x,$$

with the angular frequency $\omega_0^2 = k/m$.

The above differential equation has the advantage that it can be solved analytically with solutions on the form

$$x(t) = A \cos(\omega_0 t + \nu),$$

where A is the amplitude and ν the phase constant. This provides in turn an important test for the numerical solution and the

Simple Example, Block tied to a Wall

With the position $x(t)$ and the velocity $v(t) = dx/dt$ we can reformulate Newton's equation in the following way

$$\frac{dx(t)}{dt} = v(t),$$

and

$$\frac{dv(t)}{dt} = -\omega_0^2 x(t).$$

We are now going to solve these equations using the Runge-Kutta method to fourth order discussed previously.

Simple Example, Block tied to a Wall

Before proceeding however, it is important to note that in addition to the exact solution, we have at least two further tests which can be used to check our solution.

Since functions like \cos are periodic with a period 2π , then the solution $x(t)$ has also to be periodic. This means that

$$x(t + T) = x(t),$$

with T the period defined as

$$T = \frac{2\pi}{\omega_0} = \frac{2\pi}{\sqrt{k/m}}.$$

Observe that T depends only on k/m and not on the amplitude of the solution.

Simple Example, Block tied to a Wall

In addition to the periodicity test, the total energy has also to be conserved.

Suppose we choose the initial conditions

$$x(t=0) = 1 \quad m \quad v(t=0) = 0 \quad m/s,$$

meaning that block is at rest at $t=0$ but with a potential energy

$$E_0 = \frac{1}{2} k x(t=0)^2 = \frac{1}{2} k.$$

The total energy at any time t has however to be conserved, meaning that our solution has to fulfil the condition

$$E_0 = \frac{1}{2} k x(t)^2 + \frac{1}{2} m v(t)^2.$$

Simple Example, Block tied to a Wall

An algorithm which implements these equations is included below.

- Choose the initial position and speed, with the most common choice $v(t=0) = 0$ and some fixed value for the position.
- Choose the method you wish to employ in solving the problem.
- Subdivide the time interval $[t_i, t_f]$ into a grid with step size $h = \frac{t_f - t_i}{N}$, where N is the number of mesh points.
- Calculate now the total energy given by $E_0 = \frac{1}{2} k x(t=0)^2 = \frac{1}{2} k$.
- The Runge-Kutta method is used to obtain x_{i+1} and v_{i+1} starting from the previous values x_i and v_i .
- When we have computed $x(v)_{i+1}$ we upgrade $t_{i+1} = t_i + h$.
- This iterative process continues till we reach the maximum time t_f .
- The results are checked against the exact solution. Furthermore, one has to check the stability of the numerical solution against the chosen number of mesh points N .

Simple Example, Block tied to a Wall

```

y[0] = initial_x;           // initial position
y[1] = initial_v;           // initial velocity
t=0;                         // initial time
E0 = 0.5*y[0]*y[0]+0.5*y[1]*y[1]; // the initial total energy
// now we start solving the differential
// equations using the RK4 method
while (t <= tmax){
    derivatives(t, y, dydt); // initial derivatives
    runge_kutta_4(y, dydt, n, t, h, yout, derivatives);
    for (i = 0; i < n; i++) {
        y[i] = yout[i];
    }
    t += h;
    output(t, y, E0); // write to file
}

```

Simple Example, Block tied to a Wall

```
// this function sets up the derivatives for this special case
void derivatives(double t, double *y, double *dydt)
{
    dydt[0]=y[1]; // derivative of x
    dydt[1]=-y[0]; // derivative of v
} // end of function derivatives
```

Runge-Kutta methods, code

```
void runge_kutta_4(double *y, double *dydx, int n,
                  double x, double h,
                  double *yout, void (*derive)(double, double *, double *))
{
    int i;
    double xh,hh,h6;
    double *dym, *dyt, *yt;
    // allocate space for local vectors
    dym = new double [n];
    dyt = new double [n];
    yt = new double [n];
    hh = h*0.5;
    h6 = h/6.;
    xh = x+hh;
```

Runge-Kutta methods, code

```
for (i = 0; i < n; i++) {
    yt[i] = y[i]+hh*dydx[i];

    (*derive)(xh,yt,dyt); // computation of k2
    for (i = 0; i < n; i++) {
        yt[i] = y[i]+hh*dyt[i];

        (*derive)(xh,yt,dym); // computation of k3
        for (i=0; i < n; i++) {
            yt[i] = y[i]+h*dym[i];
            dym[i] += dyt[i];

            (*derive)(x+h,yt,dyt); // computation of k4
            // now we upgrade y in the array yout
            for (i = 0; i < n; i++){
                yout[i] = y[i]+h6*(dydx[i]+dyt[i]+2.0*dym[i]);

                delete [] dym;
                delete [] dyt;
                delete [] yt;
            } // end of function Runge-kutta 4
```

The classical pendulum

The angular equation of motion of the pendulum is given by Newton's equation and with no external force it reads

$$ml \frac{d^2\theta}{dt^2} + mgsin(\theta) = 0, \quad (56)$$

with an angular velocity and acceleration given by

$$v = l \frac{d\theta}{dt}, \quad (57)$$

and

$$a = l \frac{d^2\theta}{dt^2}. \quad (58)$$

More on the Pendulum

We do however expect that the motion will gradually come to an end due a viscous drag torque acting on the pendulum. In the presence of the drag, the above equation becomes

$$ml \frac{d^2\theta}{dt^2} + \nu \frac{d\theta}{dt} + mgsin(\theta) = 0, \quad (59)$$

where ν is now a positive constant parameterizing the viscosity of the medium in question. In order to maintain the motion against viscosity, it is necessary to add some external driving force. We choose here a periodic driving force. The last equation becomes then

$$ml \frac{d^2\theta}{dt^2} + \nu \frac{d\theta}{dt} + mgsin(\theta) = Asin(\omega t), \quad (60)$$

with A and ω two constants representing the amplitude and the angular frequency respectively. The latter is called the driving frequency.

More on the Pendulum

We define

$$\omega_0 = \sqrt{g/l},$$

the so-called natural frequency and the new dimensionless quantities

$$\hat{t} = \omega_0 t,$$

with the dimensionless driving frequency

$$\hat{\omega} = \frac{\omega}{\omega_0},$$

and introducing the quantity Q , called the *quality factor*,

$$Q = \frac{mg}{\omega_0 \nu},$$

and the dimensionless amplitude

More on the Pendulum

we have

$$\frac{d^2\theta}{dt^2} + \frac{1}{Q} \frac{d\theta}{dt} + \sin(\theta) = \hat{A} \cos(\hat{\omega} \hat{t}).$$

This equation can in turn be recast in terms of two coupled first-order differential equations as follows

$$\frac{d\theta}{dt} = \hat{v},$$

and

$$\frac{d\hat{v}}{dt} = -\frac{\hat{v}}{Q} - \sin(\theta) + \hat{A} \cos(\hat{\omega} \hat{t}).$$

These are the equations to be solved. The factor Q represents the number of oscillations of the undriven system that must occur before its energy is significantly reduced due to the viscous drag. The amplitude \hat{A} is measured in units of the maximum possible

Classes for ODE methods

It can be very useful to make a Class which contains all possible methods discussed. In Fortran we can use the MODULE keyword in order to can methods and keep the variables private and hidden from other parts of our code. This allows for a generalization which can be used to tackle other ODEs as well.

Classes for ODE methods

In program2.cpp of chapter 8 we have canned the following methods

```
void euler();
void euler_cromer();
void midpoint();
void euler_richardson();
void half_step();
void rk2(); //runge-kutta-second-order
void rk4_step(double,double*,double*,double); // we need it in func
void rk4(); //runge-kutta-fourth-order
void asc(); //runge-kutta-fourth-order with adaptive stepsize contr
```

Classes for ODE methods

```
class pendulum
{
private:
    double Q, A_roof, omega_0, omega_roof,g; //
    double y[2]; //for the initial-values of phi and v
    int n; // how many steps
    double delta_t,delta_t_roof;

public:
    void derivatives(double,double*,double*);
    void initialise();
    void euler();
    void euler_cromer();
    void midpoint();
    void euler_richardson();
    void half_step();
    void rk2(); //runge-kutta-second-order
    void rk4_step(double,double*,double*,double); // we need it in f
    void rk4(); //runge-kutta-fourth-order
    void asc(); //runge-kutta-fourth-order with adaptive stepsize co
};
```

Classes for ODE methods

```
void pendulum::derivatives(double t, double* in, double* out)
{ /* Here we are calculating the derivatives at (dimensionless) time t
   'in' are the values of phi and v, which are used for the calculation
   The results are given to 'out' */

    out[0]=in[1]; //out[0] = (phi)' = v
    if(Q)
        out[1]=-in[1]/(((double)Q)-sin(in[0])+A_roof*cos(omega_roof*t));
    else
        out[1]=-sin(in[0])+A_roof*cos(omega_roof*t); //out[1] = (phi)''
```

Classes for ODE methods

```
int main()
{
    pendulum testcase;
    testcase.initialise();
    testcase.euler();
    testcase.euler_cromer();
    testcase.midpoint();
    testcase.euler_richardson();
    testcase.half_step();
    testcase.rk2();
    testcase.rk4();
    return 0;
} // end of main function
```

Classes for ODE methods

In Fortran we would use

```
MODULE pendulum
USE CONSTANTS
IMPLICIT NONE
REAL(DP), PRIVATE :: Q, A_roof, omega_0, omega_roof, g
REAL(DP), PRIVATE :: y(2) ! for the initial-values of p
INTEGER, PRIVATE :: n ! how many steps
REAL(DP), PRIVATE :: delta_t, delta_t_roof

CONTAINS
SUBROUTINE derivatives(...)
SUBROUTINE initialise(...)
SUBROUTINE euler(...)
SUBROUTINE euler_cromer(...)
SUBROUTINE midpoint(...)
etc
END MODULE pendulum
```

The report: how to write a good scientific/technical report

What should it contain? A typical structure.

- An introduction where you explain the aims and rationale for the physics case and what you have done. At the end of the introduction you should give a brief summary of the structure of the report
- Theoretical models and technicalities. This is the methods section.
- Results and discussion
- Conclusions and perspectives
- Appendix with extra material
- Bibliography Keep always a good log of what you do.

The report

What should I focus on? Introduction.

You don't need to answer all questions in a chronological order. When you write the introduction you could focus on the following aspects

- Motivate the reader, the first part of the introduction gives always a motivation and tries to give the overarching ideas
- What I have done
- The structure of the report, how it is organized etc

The report

What should I focus on? Methods sections.

- Describe the methods and algorithms
- You need to explain how you implemented the methods and also say something about the structure of your algorithm and present some parts of your code
- You should plug in some calculations to demonstrate your code, such as selected runs used to validate and verify your results. The latter is extremely important!! A reader needs to understand that your code reproduces selected benchmarks and reproduces previous results, either numerical and/or well-known closed form expressions.

The report

What should I focus on? Results.

- Present your results
- Give a critical discussion of your work and place it in the correct context.
- Relate your work to other calculations/studies
- An eventual reader should be able to reproduce your calculations if she/he wants to do so. All input variables should be properly explained.
- Make sure that figures and tables should contain enough information in their captions, axis labels etc so that an eventual reader can gain a first impression of your work by studying figures and tables only.

The report

What should I focus on? Conclusions.

- State your main findings and interpretations
- Try as far as possible to present perspectives for future work
- Try to discuss the pros and cons of the methods and possible improvements

The report

What should I focus on? additional material.

- Additional calculations used to validate the codes
- Selected calculations, these can be listed with few comments
- Listing of the code if you feel this is necessary You can consider moving parts of the material from the methods section to the appendix. You can also place additional material on your webpage.

The report

What should I focus on? References.

- Give always references to material you base your work on, either scientific articles/reports or books.
- *Wikipedia is not accepted as a scientific reference.* Under no circumstances.
- Refer to articles as: name(s) of author(s), journal, volume (boldfaced), page and year in parenthesis.
- Refer to books as: name(s) of author(s), title of book, publisher, place and year, eventual page numbers

Overview of week 41

Ordinary differential equations (ODEs) and Partial differential equations (PDEs).

- Monday: Repetition from last week
- Adaptive Runge-Kutta methods and stiff equations
- Examples with codes
- Discussion of project 3
- Begin partial differential equations, discussion of the diffusion equation
- Tuesday:
- Diffusion equation in one spatial dimension, implicit and explicit scheme and the Crank-Nicolson scheme Chapter 9, ODEs with boundary conditions will not be discussed.

Student project 3: An object oriented example program for project 3

See <http://folk.uio.no/mhjensen/compphys/programs/chapter08/cpp/solarsystem>.

- <http://folk.uio.no/mhjensen/compphys/programs/chapter08/cpp/solarsystem/solarsystem.cpp>
- <http://folk.uio.no/mhjensen/compphys/programs/chapter08/cpp/solarsystem/planet.cpp>
- <http://folk.uio.no/mhjensen/compphys/programs/chapter08/cpp/solarsystem/planet.h>
- <http://folk.uio.no/mhjensen/compphys/programs/chapter08/cpp/solarsystem/constants.cpp>
- <http://folk.uio.no/mhjensen/compphys/programs/chapter08/cpp/solarsystem/constants.h>

Adaptive methods

In case the function to integrate varies slowly or fast in different integration domains, adaptive methods are normally used. One strategy is always to decrease the step size. As we have seen earlier, this leads to more CPU cycles and may lead to loss or numerical precision. An alternative is to use higher-order RK methods for example. However, this leads again to more cycles, furthermore, there is no guarantee that higher-order leads to an improved error.

Adaptive methods

Assume the exact result is \tilde{x} and that we are using an RKM method. Suppose we run two calculations, one with h (called x_1) and one with $h/2$ (called x_2). Then

$$\tilde{x} = x_1 + Ch^{M+1} + O(h^{M+2}),$$

and

$$\tilde{x} = x_2 + 2C(h/2)^{M+1} + O(h^{M+2}),$$

with C a constant. Note that we calculate two halves in the last equation. We get then

$$|x_1 - x_2| = Ch^{M+1}(1 - \frac{1}{2^M}).$$

yielding

$$C = \frac{|x_1 - x_2|}{(1 - 2^{-M})h^{M+1}}.$$

Adaptive methods

With RK4 the expressions become

$$\tilde{x} = x_2 + \epsilon + O((h)^6),$$

with

$$\epsilon = \frac{|x_1 - x_2|}{15}.$$

The estimate is one order higher than the original RK4. But this method is normally rather inefficient since it requires a lot of computations. We solve typically the equation three times at each time step. However, we can compare the estimate ϵ with some by us given accuracy ξ . We can then ask the question: what is, with a given x_j and t_j , the largest possible step size \tilde{h} that leads to a truncation error below ξ ? We want

$$C\tilde{h} \leq \xi,$$

which leads to

Adaptive methods

With

$$\tilde{h} = h \left(\frac{\xi}{\epsilon} \right)^{1+1/M}.$$

we can design the following algorithm:

- If the two answers are close, keep the approximation to h .
- If $\epsilon > \xi$ we need to decrease the step size in the next time step.
- If $\epsilon < \xi$ we need to increase the step size in the next time step. A much used algorithm is the so-called RKF45 which uses a combination of a fourth and fifth order RK methods.

Adaptive methods, RKF45

At each step, two different approximations for the solution are made and compared. If the two answers are in close agreement, the approximation is accepted. If the two answers do not agree to a specified accuracy, the step size is reduced. If the answers agree to more significant digits than required, the step size is increased. Each step requires the use of the following six values:

$$k_1 = hf(t_k, y_k),$$

$$k_2 = hf\left(t_k + \frac{1}{4}h, y_k + \frac{1}{4}k_1\right),$$

$$k_3 = hf\left(t_k + \frac{3}{8}h, y_k + \frac{3}{32}k_1 + \frac{9}{32}k_2\right),$$

$$k_4 = hf\left(t_k + \frac{12}{13}h, y_k + \frac{1932}{2197}k_1 + \frac{7200}{2197}k_2 + \frac{7296}{2197}k_3\right),$$

Adaptive methods, RKF45

Then an approximation to the solution of the ODE is made using a Runge-Kutta method of order 4:

$$y_{k+1} = y_k + \frac{25}{216}k_1 + \frac{1408}{2565}k_3 + \frac{2197}{4101}k_4 - \frac{1}{5}k_5,$$

where the four function values k_1 , k_3 , k_4 , and k_5 are used.

Notice that k_2 is not used here. A better value for the solution is determined using a Runge-Kutta method of order 5:

$$z_{k+1} = y_k + \frac{16}{135}k_1 + \frac{6656}{12825}k_3 + \frac{28561}{56430}k_4 - \frac{9}{50}k_5 + \frac{2}{55}k_6.$$

The optimal time step αh is then determined by

$$\alpha = \left(\frac{\xi h}{2|z_{k+1} - y_{k+1}|} \right)^{1/4},$$

with ξ our defined tolerance.

Partial Differential Equations, chapter 10

General 2+1-dim PDE

$$A(x, y) \frac{\partial^2 U}{\partial x^2} + B(x, y) \frac{\partial^2 U}{\partial x \partial y} + C(x, y) \frac{\partial^2 U}{\partial y^2} = F(x, y, U, \frac{\partial U}{\partial x}, \frac{\partial U}{\partial y})$$

Examples

$$B = C = 0,$$

give e.g., 1+1-dim diffusion equation

$$A \frac{\partial^2 U}{\partial x^2} = \frac{\partial U}{\partial t}$$

and is an example of a parabolic PDE

Partial Differential Equations

More examples 2+1-dim wave equation

$$A \frac{\partial^2 U}{\partial x^2} + C \frac{\partial^2 U}{\partial y^2} = \frac{\partial^2 U}{\partial t^2}$$

Poisson's (Laplace's $\rho = 0$) equation

$$\nabla^2 U(\mathbf{x}) = -4\pi\rho(\mathbf{x}).$$

Heat/Diffusion Equation

Diffusion equation

$$\frac{\kappa}{C\rho} \nabla^2 T(\mathbf{x}, t) = \frac{\partial T(\mathbf{x}, t)}{\partial t}$$

$$\frac{\kappa}{C\rho(\mathbf{x}, t)} \nabla^2 T(\mathbf{x}, t) = \frac{\partial T(\mathbf{x}, t)}{\partial t}$$

Explicit Scheme for the Diffusion Equation

In one dimension we have thus the following equation

$$\nabla^2 u(x, t) = \frac{\partial u(x, t)}{\partial t},$$

or

$$u_{xx} = u_t,$$

with initial conditions, i.e., the conditions at $t = 0$,

$$u(x, 0) = g(x) \quad 0 \leq x \leq L$$

with $L = 1$ the length of the x -region of interest. The boundary conditions are

$$u(0, t) = a(t) \quad t \geq 0,$$

and

$$u(L, t) = b(t) \quad t \geq 0,$$

where $a(t)$ and $b(t)$ are two functions which depend on time only, while $g(x)$ depends only on the position x .

Explicit Scheme, Forward Euler

$$u_t \approx \frac{u_{i,j+1} - u_{i,j}}{\Delta t},$$

and

$$u_{xx} \approx \frac{u_{i+1,j} - 2u_{i,j} + u_{i-1,j}}{\Delta x^2}.$$

The one-dimensional diffusion equation can then be rewritten in its discretized version as

$$\frac{u_{i,j+1} - u_{i,j}}{\Delta t} = \frac{u_{i+1,j} - 2u_{i,j} + u_{i-1,j}}{\Delta x^2}.$$

Defining $\alpha = \Delta t / \Delta x^2$ results in the explicit scheme

$$u_{i,j+1} = \alpha u_{i-1,j} + (1 - 2\alpha) u_{i,j} + \alpha u_{i+1,j}.$$

Explicit Scheme

$$V_{j+1} = AV_j$$

with

$$A = \begin{pmatrix} 1 - 2\alpha & \alpha & 0 & 0 \dots \\ \alpha & 1 - 2\alpha & \alpha & 0 \dots \\ \dots & \dots & \dots & \dots \\ 0 \dots & 0 \dots & \alpha & 1 - 2\alpha \end{pmatrix}$$

yielding

$$V_{j+1} = AV_j = \dots = A^j V_0$$

The explicit scheme, although being rather simple to implement has a very weak stability condition given by

$$\Delta t / \Delta x^2 \leq 1/2$$

Implicit Scheme

Choose now

$$u_t \approx \frac{u(x_i, t_j) - u(x_i, t_j - k)}{k}$$

and

$$u_{xx} \approx \frac{u(x_i + h, t_j) - 2u(x_i, t_j) + u(x_i - h, t_j)}{h^2}$$

Define $\alpha = k / h^2$. Gives

$$u_{i,j-1} = -\alpha u_{i-1,j} + (1 - 2\alpha) u_{i,j} - \alpha u_{i+1,j}$$

Here $u_{i,j-1}$ is the only unknown quantity.

Have

$$AV_j = V_{j-1}$$

with

Brute Force Implicit Scheme, inefficient algo

```
! now invert the matrix
CALL matinv( a, ndim, det)
DO i = 1, m
  DO l=1, ndim
    u(1) = DOT_PRODUCT(a(1,:),v(:))
  ENDDO
  v = u
  t = i*k
  DO j=1, ndim
    WRITE(6,*) t, j*h, v(j)
  ENDDO
ENDDO
```

Brief Summary of the Explicit and the Implicit Methods

Explicit is straightforward to code, but avoid doing the matrix vector multiplication since the matrix is tridiagonal.

$$u_t \approx \frac{u(x, t + \Delta t) - u(x, t)}{\Delta t} = \frac{u(x_i, t_j + \Delta t) - u(x_i, t_j)}{\Delta t}$$

The implicit method can be applied in a brute force way as well as long as the element of the matrix are constants.

$$u_t \approx \frac{u(x, t) - u(x, t - \Delta t)}{\Delta t} = \frac{u(x_i, t_j) - u(x_i, t_j - \Delta t)}{\Delta t}$$

However, it is more efficient to use a linear algebra solver for tridiagonal matrices.

Crank-Nicolson

$$\frac{\theta}{\Delta x^2} (u_{i-1,j} - 2u_{i,j} + u_{i+1,j}) + \frac{1-\theta}{\Delta x^2} (u_{i+1,j-1} - 2u_{i,j-1} + u_{i-1,j-1}) = \frac{1}{\Delta t}$$

which for $\theta = 0$ yields the forward formula for the first derivative and the explicit scheme, while $\theta = 1$ yields the backward formula and the implicit scheme. These two schemes are called the backward and forward Euler schemes, respectively. For $\theta = 1/2$ we obtain a new scheme after its inventors, Crank and Nicolson.

Crank Nicolson

Using our previous definition of $\alpha = \Delta t / \Delta x^2$ we can rewrite the latter equation as

$$-\alpha u_{i-1,j} + (2 + 2\alpha) u_{i,j} - \alpha u_{i+1,j} = \alpha u_{i-1,j-1} + (2 - 2\alpha) u_{i,j-1} + \alpha u_{i+1,j-1},$$

or in matrix-vector form as

$$(2\hat{I} + \alpha\hat{B}) V_j = (2\hat{I} - \alpha\hat{B}) V_{j-1},$$

where the vector V_j is the same as defined in the implicit case while the matrix \hat{B} is

$$\hat{B} = \begin{pmatrix} 2 & -1 & 0 & 0 \dots \\ -1 & 2 & -1 & 0 \dots \\ \dots & \dots & \dots & \dots \\ 0 \dots & 0 \dots & & 2 \end{pmatrix}$$

Analysis of diffusion equation

We start with the forward Euler scheme and Taylor expand $u(x, t + \Delta t)$, $u(x + \Delta x, t)$ and $u(x - \Delta x, t)$

$$\begin{aligned} u(x + \Delta x, t) &= u(x, t) + \frac{\partial u(x, t)}{\partial x} \Delta x + \frac{\partial^2 u(x, t)}{2 \partial x^2} \Delta x^2 + \mathcal{O}(\Delta x^3), \\ u(x - \Delta x, t) &= u(x, t) - \frac{\partial u(x, t)}{\partial x} \Delta x + \frac{\partial^2 u(x, t)}{2 \partial x^2} \Delta x^2 + \mathcal{O}(\Delta x^3), \\ u(x, t + \Delta t) &= u(x, t) + \frac{\partial u(x, t)}{\partial t} \Delta t + \mathcal{O}(\Delta t^2). \end{aligned} \quad (61)$$

Analysis of diffusion equation

With these Taylor expansions the approximations for the derivatives takes the form

$$\begin{aligned} \left[\frac{\partial u(x, t)}{\partial t} \right]_{\text{approx}} &= \frac{\partial u(x, t)}{\partial t} + \mathcal{O}(\Delta t), \\ \left[\frac{\partial^2 u(x, t)}{\partial x^2} \right]_{\text{approx}} &= \frac{\partial^2 u(x, t)}{\partial x^2} + \mathcal{O}(\Delta x^2). \end{aligned} \quad (62)$$

It is easy to convince oneself that the backward Euler method must have the same truncation errors as the forward Euler scheme.

Analysis of diffusion equation

For the Crank-Nicolson scheme we also need to Taylor expand $u(x + \Delta x, t + \Delta t)$ and $u(x - \Delta x, t + \Delta t)$ around $t' = t + \Delta t/2$.

$$\begin{aligned} u(x + \Delta x, t + \Delta t) &= u(x, t') + \frac{\partial u(x, t')}{\partial x} \Delta x + \frac{\partial u(x, t')}{\partial t} \frac{\Delta t}{2} + \frac{\partial^2 u(x, t')}{2 \partial x^2} \frac{\partial^2 u(x, t')}{\partial x \partial t} \frac{\Delta t}{2} \Delta x + \mathcal{O}(\Delta t^3) \\ u(x - \Delta x, t + \Delta t) &= u(x, t') - \frac{\partial u(x, t')}{\partial x} \Delta x + \frac{\partial u(x, t')}{\partial t} \frac{\Delta t}{2} + \frac{\partial^2 u(x, t')}{2 \partial x^2} \frac{\partial^2 u(x, t')}{\partial x \partial t} \frac{\Delta t}{2} \Delta x + \mathcal{O}(\Delta t^3) \\ u(x + \Delta x, t) &= u(x, t') + \frac{\partial u(x, t')}{\partial x} \Delta x - \frac{\partial u(x, t')}{\partial t} \frac{\Delta t}{2} + \frac{\partial^2 u(x, t')}{2 \partial x^2} \frac{\partial^2 u(x, t')}{\partial x \partial t} \frac{\Delta t}{2} \Delta x + \mathcal{O}(\Delta t^3) \\ u(x - \Delta x, t) &= u(x, t') - \frac{\partial u(x, t')}{\partial x} \Delta x - \frac{\partial u(x, t')}{\partial t} \frac{\Delta t}{2} + \frac{\partial^2 u(x, t')}{2 \partial x^2} \end{aligned}$$

Analysis of diffusion equation

We now insert these expansions in the approximations for the derivatives to find

$$\begin{aligned} \left[\frac{\partial u(x, t')}{\partial t} \right]_{\text{approx}} &= \frac{\partial u(x, t')}{\partial t} + \mathcal{O}(\Delta t^2), \\ \left[\frac{\partial^2 u(x, t')}{\partial x^2} \right]_{\text{approx}} &= \frac{\partial^2 u(x, t')}{\partial x^2} + \mathcal{O}(\Delta x^2). \end{aligned} \quad (64)$$

Analysis of diffusion equation

The following table summarizes the three methods.

Scheme:	Truncation Error:	Stability requirements:
Crank-Nicolson	$\mathcal{O}(\Delta x^2, \Delta t^2)$	Stable for all Δt and Δx
Backward Euler	$\mathcal{O}(\Delta x^2, \Delta t)$	Stable for all Δt and Δx
Forward Euler	$\mathcal{O}(\Delta x^2, \Delta t)$	$\Delta t \leq \frac{1}{2} \Delta x^2$

Analysis of diffusion equation

It cannot be repeated enough, it is always useful to find cases where one can compare the numerics and the developed algorithms and codes with analytic solution. The above case is also particularly simple. We have the following partial differential equation

$$\nabla^2 u(x, t) = \frac{\partial u(x, t)}{\partial t},$$

with initial conditions

$$u(x, 0) = g(x) \quad 0 < x < L.$$

Analysis of diffusion equation

The boundary conditions are

$$u(0, t) = 0 \quad t \geq 0, \quad u(L, t) = 0 \quad t \geq 0,$$

We assume that we have solutions of the form (separation of variable)

$$u(x, t) = F(x)G(t), \quad (65)$$

which inserted in the partial differential equation results in

$$\frac{F''}{F} = \frac{G'}{G}, \quad (66)$$

where the derivative is with respect to x on the left hand side and with respect to t on right hand side. This equation should hold for all x and t . We must require the rhs and lhs to be equal to a constant.

Analysis of diffusion equation

We call this constant $-\lambda^2$. This gives us the two differential equations,

$$F'' + \lambda^2 F = 0; \quad G' = -\lambda^2 G, \quad (67)$$

with general solutions

$$F(x) = A \sin(\lambda x) + B \cos(\lambda x); \quad G(t) = C e^{-\lambda^2 t}. \quad (68)$$

Analysis of diffusion equation

To satisfy the boundary conditions we require $B = 0$ and $\lambda = n\pi/L$. One solution is therefore found to be

$$u(x, t) = A_n \sin(n\pi x/L) e^{-n^2 \pi^2 t/L^2}. \quad (69)$$

But there are infinitely many possible n values (infinite number of solutions). Moreover, the diffusion equation is linear and because of this we know that a superposition of solutions will also be a solution of the equation. We may therefore write

$$u(x, t) = \sum_{n=1}^{\infty} A_n \sin(n\pi x/L) e^{-n^2 \pi^2 t/L^2}. \quad (70)$$

Analysis of diffusion equation

The coefficient A_n is in turn determined from the initial condition. We require

$$u(x, 0) = g(x) = \sum_{n=1}^{\infty} A_n \sin(n\pi x/L). \quad (71)$$

The coefficient A_n is the Fourier coefficients for the function $g(x)$. Because of this, A_n is given by (from the theory on Fourier series)

$$A_n = \frac{2}{L} \int_0^L g(x) \sin(n\pi x/L) dx. \quad (72)$$

Different $g(x)$ functions will obviously result in different results for A_n .

Overview of week 42

Partial differential equations (chapter 10) and begin numerical integration (chapter 5).

- Monday: Repetition from last week
- Discussion of project 3, with an emphasis on object orientation
- Diffusion equation in two spatial dimensions
- Poisson's and Laplace's equations
- Tuesday:
- Wave equation in one and two dimensions.
- If we get time, we will start with numerical integration

Laplace's and Poisson's equations

Laplace's equation reads

$$\nabla^2 u(\mathbf{x}) = u_{xx} + u_{yy} = 0.$$

with possible boundary conditions $u(x, y) = g(x, y)$ on the border. There is no time-dependence. Choosing equally many steps in both directions we have a quadratic or rectangular grid, depending on whether we choose equal steps lengths or not in the x and the y directions. Here we set $\Delta x = \Delta y = h$ and obtain a discretized version

$$u_{xx} \approx \frac{u(x+h, y) - 2u(x, y) + u(x-h, y)}{h^2},$$

and

$$u_{yy} \approx \frac{u(x, y+h) - 2u(x, y) + u(x, y-h)}{h^2},$$

Laplace's and Poisson's equations

$$u_{xx} \approx \frac{u_{i+1,j} - 2u_{i,j} + u_{i-1,j}}{h^2},$$

and

$$u_{yy} \approx \frac{u_{i,j+1} - 2u_{i,j} + u_{i,j-1}}{h^2},$$

which gives when inserted in Laplace's equation

$$u_{i,j} = \frac{1}{4} [u_{i,j+1} + u_{i,j-1} + u_{i+1,j} + u_{i-1,j}]. \quad (73)$$

This is our final numerical scheme for solving Laplace's equation. Poisson's equation adds only a minor complication to the above equation since in this case we have

$$u_{xx} + u_{yy} = -\rho(\mathbf{x}),$$

and we need only to add a discretized version of $\rho(\mathbf{x})$ resulting in

$$u_{i,j} = \frac{1}{4} [u_{i,j+1} + u_{i,j-1} + u_{i+1,j} + u_{i-1,j}] + \rho_{i,j}. \quad (74)$$

Solution Approach

The way we solve these equations is based on an iterative scheme we discussed in connection with linear algebra, namely the so-called Jacobi, Gauss-Seidel and relaxation methods. The steps are rather simple. We start with an initial guess for $u_{i,j}^{(0)}$ where all values are known. To obtain a new solution we solve Eq. (73) or Eq. (74) in order to obtain a new solution $u_{i,j}^{(1)}$. Most likely this solution will not be a solution to Eq. (73). This solution is in turn used to obtain a new and improved $u_{i,j}^{(2)}$. We continue this process till we obtain a result which satisfies some specific convergence criterion.

Code example for the two-dimensional diff equation/Laplace

```
int DiffusionJacobi(int N, double dx, double dt,
    double **A, double **q, double abstol){
    int i,j,k;
    int maxit = 100000;
    double sum;
    double **Aold = CreateMatrix(N,N);

    double D = dt/(dx*dx);
```

Code example for the two-dimensional diff equation/Laplace

```
for(i=1; i<N-1; i++)
    for(j=1; j<N-1; j++)
        Aold[i][j] = 1.0;
/* Boundary Conditions -- all zeros */
for(i=0; i<N; i++){
    A[0][i] = 0.0;
    A[N-1][i] = 0.0;
    A[i][0] = 0.0;
    A[i][N-1] = 0.0;
```

Code example for the two-dimensional diff equation/Laplace

```
for(k=0; k<maxit; k++){
    for(i = 1; i<N-1; i++){
        for(j=1; j<N-1; j++){
            A[i][j] = dt*q[i][j] + Aold[i][j] +
            D*(Aold[i+1][j] + Aold[i][j+1] - 4.0*Aold[i][j] +
            Aold[i-1][j] + Aold[i][j-1]);

            sum = 0.0;
            for(i=0; i<N; i++){
                for(j=0; j<N; j++){
                    sum += (Aold[i][j]-A[i][j])*(Aold[i][j]-A[i][j]);
                }
            }
            Aold[i][j] = A[i][j];

            if(sqrt(sum)<abstol){DestroyMatrix(Aold,N,N);
                return k;
```

Two-dimensional wave equation

Consider first the two-dimensional wave equation for a vibrating square membrane given by the following initial and boundary conditions

$$\begin{cases} \lambda \left(\frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} \right) = \frac{\partial^2 u}{\partial t^2} & x, y \in [0, 1], t \geq 0 \\ u(x, y, 0) = \sin(\pi x) \sin(2\pi y) & x, y \in (0, 1) \\ u = 0 \text{ boundary} & t \geq 0 \\ \partial u / \partial t|_{t=0} = 0 & x, y \in (0, 1) \end{cases}.$$

The boundary is defined by $x = 0$, $x = 1$, $y = 0$ and $y = 1$. Here we set $\lambda = 1$.

Two-dimensional wave equation

Our equations depend on three variables whose discretized versions are now

$$\begin{cases} t_l = l \Delta t & l \geq 0 \\ x_i = i \Delta x & 0 \leq i \leq n_x \\ y_j = j \Delta y & 0 \leq j \leq n_y \end{cases},$$

and we will let $\Delta x = \Delta y = h$ and $n_x = n_y$ for the sake of simplicity. We have now the following discretized partial derivatives

$$u_{xx} \approx \frac{u'_{i+1,j} - 2u'_{i,j} + u'_{i-1,j}}{h^2},$$

and

$$u_{yy} \approx \frac{u'_{i,j+1} - 2u'_{i,j} + u'_{i,j-1}}{h^2},$$

and

$$u_{tt} \approx \frac{u'^{l+1}_{i,j} - 2u'^l_{i,j} + u'^{l-1}_{i,j}}{\Delta t^2}.$$

Two-dimensional wave equation

We merge this into the discretized 2 + 1-dimensional wave equation as

$$u'^{l+1}_{i,j} = 2u'^l_{i,j} - u'^{l-1}_{i,j} + \frac{\Delta t^2}{h^2} (u'^l_{i+1,j} - 4u'^l_{i,j} + u'^l_{i-1,j} + u'^l_{i,j+1} + u'^l_{i,j-1}),$$

where again we have an explicit scheme with $u'^{l+1}_{i,j}$ as the only unknown quantity. It is easy to account for different step lengths for x and y . The partial derivative is treated in much the same way as for the one-dimensional case, except that we now have an additional index due to the extra spatial dimension, viz., we need to compute $u'^{-1}_{i,j}$ through

$$u'^{-1}_{i,j} = u^0_{i,j} + \frac{\Delta t}{2h^2} (u^0_{i+1,j} - 4u^0_{i,j} + u^0_{i-1,j} + u^0_{i,j+1} + u^0_{i,j-1}),$$

in our setup of the initial conditions.

Code example for the two-dimensional wave equation

We show here how to implement the two-dimensional wave equation

```
// After initializations and declaration of variables
for ( int i = 0; i < n; i++ ) {
    u[i] = new double [n];
    uLast[i] = new double [n];
    uNext[i] = new double [n];
    x[i] = i*h;
    y[i] = x[i];

    // initializing
    for ( int i = 0; i < n; i++ ) { // setting initial step
        for ( int j = 0; j < n; j++ ) {
            uLast[i][j] = sin(PI*x[i])*sin(2*PI*y[j]);
```

Code example for the two-dimensional wave equation

```
for ( int i = 1; i < (n-1); i++ ) { // setting first step using
    for ( int j = 1; j < (n-1); j++ ) {
        u[i][j] = uLast[i][j] - ((tStep*tStep)/(2.0*h*h))*
        (4*uLast[i][j] - uLast[i+1][j] - uLast[i-1][j] - uLast[i][j+1] - uLast[i][j-1]);

        u[i][0] = 0; // setting boundaries once and for all
        u[i][n-1] = 0;
        u[0][i] = 0;
        u[n-1][i] = 0;

        uNext[i][0] = 0;
        uNext[i][n-1] = 0;
        uNext[0][i] = 0;
        uNext[n-1][i] = 0;
    }
}
```

Code example for the two-dimensional wave equation

```
// iterating in time
double t = 0.0 + tStep;
int iter = 0;

while ( t < tFinal ) {
    iter ++;
    t = t + tStep;

    for ( int i = 1; i < (n-1); i++ ) { // computing next step
        for ( int j = 1; j < (n-1); j++ ) {
            uNext[i][j] = 2*u[i][j] - uLast[i][j] - ((tStep*tStep)/(h*h))*
            (4*u[i][j] - u[i+1][j] - u[i-1][j] - u[i][j+1] - u[i][j-1]);
        }
    }
}
```

Code example for the two-dimensional wave equation

```
for ( int i = 1; i < (n-1); i++ ) { // shifting results down
    for ( int j = 1; j < (n-1); j++ ) {
        uLast[i][j] = u[i][j];
        u[i][j] = uNext[i][j];
    }
}
```

Closed form solution of the wave equation

We develop here the analytic solution for the 2 + 1 dimensional wave equation with the following boundary and initial conditions

$$\begin{cases} c^2(u_{xx} + u_{yy}) = u_{tt} & x, y \in (0, L), t > 0 \\ u(x, y, 0) = f(x, y) & x, y \in (0, L) \\ u(0, 0, t) = u(L, L, t) = 0 & t > 0 \\ \partial u / \partial t|_{t=0} = g(x, y) & x, y \in (0, L) \end{cases}.$$

Closed form solution of the wave equation

Our first step is to make the ansatz

$$u(x, y, t) = F(x, y)G(t),$$

resulting in the equation

$$FG_{tt} = c^2(F_{xx}G + F_{yy}G),$$

or

$$\frac{G_{tt}}{c^2 G} = \frac{1}{F}(F_{xx} + F_{yy}) = -\nu^2.$$

The lhs and rhs are independent of each other and we obtain two differential equations

$$F_{xx} + F_{yy} + F\nu^2 = 0,$$

and

Closed form solution of the wave equation

We can in turn make the following ansatz for the x and y dependent part

$$F(x, y) = H(x)Q(y),$$

which results in

$$\frac{1}{H}H_{xx} = -\frac{1}{Q}(Q_{yy} + Q\nu^2) = -\kappa^2.$$

Since the lhs and rhs are again independent of each other, we can separate the latter equation into two independent equations, one for x and one for y, namely

$$H_{xx} + \kappa^2 H = 0,$$

and

$$Q_{yy} + \rho^2 Q = 0,$$

Closed form solution of the wave equation

The second step is to solve these differential equations, which all have trigonometric functions as solutions, viz.

$$H(x) = A \cos(\kappa x) + B \sin(\kappa x),$$

and

$$Q(y) = C \cos(\rho y) + D \sin(\rho y).$$

The boundary conditions require that $F(x, y) = H(x)Q(y)$ are zero at the boundaries, meaning that $H(0) = H(L) = Q(0) = Q(L) = 0$. This yields the solutions

$$H_m(x) = \sin\left(\frac{m\pi x}{L}\right) \quad Q_n(y) = \sin\left(\frac{n\pi y}{L}\right),$$

or

$$F_{mn}(x, y) = \sin\left(\frac{m\pi x}{L}\right) \sin\left(\frac{n\pi y}{L}\right).$$

Closed form solution of the wave equation

With $\rho^2 = \nu^2 - \kappa^2$ and $\lambda = c\nu$ we have an eigenspectrum $\lambda = c\sqrt{\kappa^2 + \rho^2}$ or $\lambda_{mn} = c\pi/L\sqrt{m^2 + n^2}$. The solution for G is

$$G_{mn}(t) = B_{mn} \cos(\lambda_{mn} t) + D_{mn} \sin(\lambda_{mn} t),$$

with the general solution of the form

$$u(x, y, t) = \sum_{mn=1}^{\infty} u_{mn}(x, y, t) = \sum_{mn=1}^{\infty} F_{mn}(x, y) G_{mn}(t).$$

Closed form solution of the wave equation

The final step is to determine the coefficients B_{mn} and D_{mn} from the Fourier coefficients. The equations for these are determined by the initial conditions $u(x, y, 0) = f(x, y)$ and $\partial u / \partial t|_{t=0} = g(x, y)$. The final expressions are

$$B_{mn} = \frac{2}{L} \int_0^L \int_0^L dx dy f(x, y) \sin\left(\frac{m\pi x}{L}\right) \sin\left(\frac{n\pi y}{L}\right),$$

and

$$D_{mn} = \frac{2}{L} \int_0^L \int_0^L dx dy g(x, y) \sin\left(\frac{m\pi x}{L}\right) \sin\left(\frac{n\pi y}{L}\right).$$

Inserting the particular functional forms of $f(x, y)$ and $g(x, y)$ one obtains the final analytic expressions.

Two-dimensional wave equation

We can check our results as function of the number of mesh points and in particular against the stability condition

$$\Delta t \leq \frac{1}{\sqrt{\lambda}} \left(\frac{1}{\Delta x^2} + \frac{1}{\Delta y^2} \right)^{-1/2}$$

where Δt , Δx and Δy are the chosen step lengths. In our case $\Delta x = \Delta y = h$. How do we find this condition? In one dimension we can proceed as we did for the diffusion equation.

Two-dimensional wave equation

The analytic solution of the wave equation in 2 + 1 dimensions has a characteristic wave component which reads

$$u(x, y, t) = A \exp(i(k_x x + k_y y - \omega t))$$

Then from

$$u_{xx} \approx \frac{u_{i+1,j}^l - 2u_{i,j}^l + u_{i-1,j}^l}{\Delta x^2},$$

we get, with $u_i = \exp i k x_i$

$$u_{xx} \approx \frac{u_i}{\Delta x^2} (\exp i k \Delta x - 2 + \exp(-i k \Delta x)),$$

or

$$u_{xx} \approx 2 \frac{u_i}{\Delta x^2} (\cos(k \Delta x) - 1) = -4 \frac{u_i}{\Delta x^2} \sin^2(k \Delta x / 2)$$

We get similar results for t and y .

Two-dimensional wave equation

We have

$$\lambda \left(\frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} \right) = \frac{\partial^2 u}{\partial t^2},$$

resulting in

$$\lambda \left(-4 \frac{u_{ij}^l}{\Delta x^2} \sin^2(k_x \Delta x / 2) - 4 \frac{u_{ij}^l}{\Delta y^2} \sin^2(k_y \Delta y / 2) \right) = -4 \frac{u_{ij}^l}{\Delta t^2} \sin^2(\omega \Delta t / 2)$$

resulting in

$$\sin(\omega \Delta t / 2) = \pm \sqrt{\lambda} \Delta t \left(\frac{1}{\Delta x^2} \sin^2(k_x \Delta x / 2) + \frac{1}{\Delta y^2} \sin^2(k_y \Delta y / 2) \right)^{1/2}.$$

The squared sine functions can at most be unity. The frequency ω is real and our wave is neither damped nor amplified.

Two-dimensional wave equation

We have

$$\sin(\omega\Delta t/2) = \pm\sqrt{\lambda}\Delta t \left(\frac{1}{\Delta x^2} \sin^2(k_x\Delta x/2) + \frac{1}{\Delta y^2} \sin^2(k_y\Delta y/2) \right)^{1/2}.$$

The squared sine functions can at most be unity. ω is real and our wave is neither damped nor amplified. The numerical ω must also be real which is the case when $\sin(\omega\Delta t/2)$ is less than or equal to unity, meaning that

$$\Delta t \leq \frac{1}{\sqrt{\lambda}} \left(\frac{1}{\Delta x^2} + \frac{1}{\Delta y^2} \right)^{-1/2}.$$

Two-dimensional wave equation

We modify now the wave equation in order to consider a 2 + 1 dimensional wave equation with a position dependent velocity, given by

$$\frac{\partial^2 u}{\partial t^2} = \nabla \cdot (\lambda(x, y) \nabla u).$$

If λ is constant, we obtain the standard wave equation discussed in the two previous points. The solution $u(x, y, t)$ could represent a model for water waves. It represents then the surface elevation from still water. We can model λ as

$$\lambda = gH(x, y),$$

with g being the acceleration of gravity and $H(x, y)$ is the still water depth.

The function $H(x, y)$ simulates the water depth using for example measurements of still water depths in say a fjord or the north sea. The boundary conditions are then determined by the coast lines as

Two-dimensional wave equation

You can discretize

$$\nabla \cdot (\lambda(x, y) \nabla u) = \frac{\partial}{\partial x} \left(\lambda(x, y) \frac{\partial u}{\partial x} \right) + \frac{\partial}{\partial y} \left(\lambda(x, y) \frac{\partial u}{\partial y} \right),$$

as follows using again a quadratic domain for x and y :

$$\frac{\partial}{\partial x} \left(\lambda(x, y) \frac{\partial u}{\partial x} \right) \approx \frac{1}{\Delta x} \left(\lambda_{i+1/2,j} \left[\frac{u_{i+1,j}^l - u_{i,j}^l}{\Delta x} \right] - \lambda_{i-1/2,j} \left[\frac{u_{i,j}^l - u_{i-1,j}^l}{\Delta x} \right] \right)$$

and

$$\frac{\partial}{\partial y} \left(\lambda(x, y) \frac{\partial u}{\partial y} \right) \approx \frac{1}{\Delta y} \left(\lambda_{i,j+1/2} \left[\frac{u_{i,j+1}^l - u_{i,j}^l}{\Delta y} \right] - \lambda_{i,j-1/2} \left[\frac{u_{i,j}^l - u_{i,j-1}^l}{\Delta y} \right] \right)$$

Two-dimensional wave equation

How did we do that? Look at the derivative wrt x only: First we compute the derivative

$$\frac{d}{dx} \left(\lambda(x) \frac{du}{dx} \right) \Big|_{x=x_i} \approx \frac{1}{\Delta x} \left(\lambda \frac{du}{dx} \Big|_{x=x_{i+1/2}} - \lambda \frac{du}{dx} \Big|_{x=x_{i-1/2}} \right),$$

where we approximated it at the midpoint by going half a step to the right and half a step to the left. Then we approximate

$$\lambda \frac{du}{dx} \Big|_{x=x_{i+1/2}} \approx \lambda_{x_{i+1/2}} \frac{u_{i+1} - u_i}{\Delta x},$$

and similarly for $x = x_i - 1/2$.

Overview of week 43

Numerical integration (chapter 5).

- Monday: Repetition from last week
- Numerical integration, Trapezoidal and Simpson's rules (equal step methods)
- Gaussian quadrature (better methods)
- Tuesday:
- Gaussian quadrature, continued
- Our first encounter with parallelization, this week MPI, next week also OpenMP with examples Next week we start also with Monte Carlo methods, which will keep us busy till the end of November.

Numerical integration and Equal Step Methods

Generalities.

Choose a step size

$$h = \frac{b-a}{N}$$

where N is the number of steps and a and b the lower and upper limits of integration.

Choose then to stop the Taylor expansion of the function $f(x)$ at a certain derivative.

With these approximations to $f(x)$ perform the integration.

$$\int_a^b f(x) dx = \int_a^{a+2h} f(x) dx + \int_{a+2h}^{a+4h} f(x) dx + \dots + \int_{b-2h}^b f(x) dx.$$

The strategy then is to find a reliable Taylor expansion for $f(x)$ in the smaller sub intervals. Consider e.g., evaluating $\int_{-h}^{+h} f(x) dx$

Equal Step Methods

Trapezoidal Rule.

Taylor expansion

$$f(x) = f_0 + \frac{f_h - f_0}{h}x + O(x^2),$$

for $x = x_0$ to $x = x_0 + h$ and

$$f(x) = f_0 + \frac{f_0 - f_{-h}}{h}x + O(x^2),$$

for $x = x_0 - h$ to $x = x_0$. The error goes like $O(x^2)$. If we then evaluate the integral we obtain

$$\int_{-h}^{+h} f(x)dx = \frac{h}{2}(f_h + 2f_0 + f_{-h}) + O(h^3),$$

which is the well-known trapezoidal rule. Local error $O(h^3) = O((b-a)^3/N^3)$, and the *global error* goes like $\approx O(h^2)$.

Equal Step Methods

Trapezoidal Rule.

Easy to implement numerically through the following simple algorithm

- Choose the number of mesh points and fix the step.
- calculate $f(a)$ and $f(b)$ and multiply with $h/2$
- Perform a loop over $n = 1$ to $n - 1$ ($f(a)$ and $f(b)$ are known) and sum up the terms $f(a+h) + f(a+2h) + f(a+3h) + \dots + f(b-h)$. Each step in the loop corresponds to a given value $a + nh$.
- Multiply the final result by h and add $hf(a)/2$ and $hf(b)/2$.

Trapezoidal Rule

```
double trapezoidal_rule(double a, double b, int n,
                       double (*func)(double))
{
    double trapez_sum;
    double fa, fb, x, step;
    int j;
    step=(b-a)/((double) n);
    fa=(*func)(a)/2.;
    fb=(*func)(b)/2.;
    trapez_sum=0.;
    for (j=1; j <= n-1; j++){
        x=j*step+a;
        trapez_sum+=(*func)(x);
    }
    trapez_sum=(trapez_sum+fa+fb)*step;
    return trapez_sum;
} // end function for trapezoidal rule
```

Trapezoidal Rule

Pay attention to the way we transfer the name of a function. This gives us the possibility to define a general trapezoidal method, where we give as input the name of the function.

```
double trapezoidal_rule(double a, double b, int n,
                       double (*func)(double))
```

We call this function simply as something like this

```
integral = trapezoidal_rule(a, b, n, mysuperduperfunction);
```

Equal Step Methods

Simpson.

The first and second derivatives are given by

$$\frac{f_h - f_{-h}}{2h} = f'_0 + \sum_{j=1}^{\infty} \frac{f_0^{(2j+1)}}{(2j+1)!} h^{2j},$$

and

$$\frac{f_h - 2f_0 + f_{-h}}{h^2} = f''_0 + 2 \sum_{j=1}^{\infty} \frac{f_0^{(2j+2)}}{(2j+2)!} h^{2j},$$

results in $f(x) = f_0 + \frac{f_h - f_{-h}}{2h}x + \frac{f_h - 2f_0 + f_{-h}}{h^2}x^2 + O(x^3)$. Inserting this formula in the integral

$$\int_{-h}^{+h} f(x)dx = \frac{h}{3}(f_h + 4f_0 + f_{-h}) + O(h^5),$$

which is Simpson's rule.

Equal Step Methods

Simpson's rule.

Note that the improved accuracy in the evaluation of the derivatives gives a better error approximation, $O(h^5)$ vs. $O(h^3)$. But this is just the *local error approximation*. Using Simpson's rule we arrive at the composite rule

$$I = \int_a^b f(x)dx = \frac{h}{3}(f(a) + 4f(a+h) + 2f(a+2h) + \dots + 4f(b-h) + f(b))$$

with a global error which goes like $O(h^4)$. Algo

- Choose the number of mesh points and fix the step.
- calculate $f(a)$ and $f(b)$
- Perform a loop over $n = 1$ to $n - 1$ ($f(a)$ and $f(b)$ are known) and sum up the terms $4f(a+h) + 2f(a+2h) + 4f(a+3h) + \dots + 4f(b-h)$. Odd values of n give 4 as factor while even values yield 2 as factor.
- Multiply the final result by $\frac{h}{3}$.

Equal Step Methods

The basic idea behind all integration methods is to approximate the integral

$$I = \int_a^b f(x) dx \approx \sum_{i=1}^N \omega_i f(x_i),$$

where ω and x are the weights and the chosen mesh points, respectively. Simpson's rule gives

$$\omega : \{h/3, 4h/3, 2h/3, 4h/3, \dots, 4h/3, h/3\},$$

for the weights, while the trapezoidal rule resulted in

$$\omega : \{h/2, h, h, \dots, h, h/2\}.$$

In general, an integration formula which is based on a Taylor series using N points, will integrate exactly a polynomial P of degree $N - 1$. That is, the N weights ω_n can be chosen to satisfy N linear equations

Equal Step Methods, Polynomials and Newton-Cotes

Given $n + 1$ distinct points $x_0, \dots, x_n \in [a, b]$ and $n + 1$ values y_0, \dots, y_n there exists a unique polynomial p_n with the property

$$p_n(x_j) = y_j \quad j = 0, \dots, n$$

In the Lagrange representation this interpolation polynomial is given by

$$p_n = \sum_{k=0}^n l_k y_k,$$

with the Lagrange factors

$$l_k(x) = \prod_{\substack{i=0 \\ i \neq k}}^n \frac{x - x_i}{x_k - x_i} \quad k = 0, \dots, n$$

Example: $n = 1$

Equal Step Methods, Polynomials and Newton-Cotes

The polynomial interpolatory quadrature of order n with equidistant quadrature points $x_k = a + kh$ and step $h = (b - a)/n$ is called the Newton-Cotes quadrature formula of order n . The integral is

$$\int_a^b f(x) dx \approx \int_a^b p_n(x) dx = \sum_{k=0}^n w_k f(x_k)$$

with

$$w_k = h \frac{(-1)^{n-k}}{k!(n-k)!} \int_0^n \prod_{\substack{j=0 \\ j \neq k}}^n (z - j) dz,$$

for $k = 0, \dots, n$.

Equal Step Methods, Polynomials and Newton-Cotes

The local error for the trapezoidal rule is

$$\int_a^b f(x) dx - \frac{b-a}{2} [f(a) + f(b)] = -\frac{h^3}{12} f^{(2)}(\xi),$$

and the global error (composite formula)

$$\int_a^b f(x) dx - T_h(f) = -\frac{b-a}{12} h^2 f^{(2)}(\xi).$$

For Simpson's rule we have

$$\int_a^b f(x) dx - \frac{b-a}{6} [f(a) + 4f((a+b)/2) + f(b)] = -\frac{h^5}{90} f^{(4)}(\xi),$$

and the global error

$$\int_a^b f(x) dx - S_h(f) = -\frac{b-a}{180} h^4 f^{(4)}(\xi).$$

Gaussian Quadrature

Methods based on Taylor series using $n + 1$ points will integrate exactly a polynomial P of degree n . If a function $f(x)$ can be approximated with a polynomial of degree n

$$f(x) \approx P_n(x),$$

with $n + 1$ mesh points we should be able to integrate exactly the polynomial P_n .

Gaussian quadrature methods promise more than this. We can get a better polynomial approximation with order greater than $n + 1$ to $f(x)$ and still get away with only $n + 1$ mesh points. More precisely, we approximate

$$f(x) \approx P_{2n+1}(x),$$

and with only $n + 1$ mesh points these methods promise that

$$\int f(x) dx \approx \int P_{2n+1}(x) dx = \sum_{k=1}^n P_{2n+1}(x_k) w_k$$

Gaussian Quadrature

What we have done till now is called Newton-Cotes quadrature. The numerical approximation goes like $O(h^n)$, where n is method-dependent.

A greater precision for a given amount of numerical work can be achieved if we are willing to give up the requirement of equally spaced integration points. In Gaussian quadrature (hereafter GQ), both the mesh points and the weights are to be determined. The points will not be equally spaced. The theory behind GQ is to obtain an arbitrary weight ω through the use of so-called orthogonal polynomials. These polynomials are orthogonal in some interval say e.g., $[-1, 1]$. Our points x_i are chosen in some optimal sense subject only to the constraint that they should lie in this interval. Together with the weights we have then $2(n + 1)$ ($n + 1$ the number of points) parameters at our disposal.

Gaussian Quadrature

Even though the integrand is not smooth, we could render it smooth by extracting from it the weight function of an orthogonal polynomial, i.e., we are rewriting

$$I = \int_a^b f(x) dx = \int_a^b W(x) g(x) dx \approx \sum_{i=0}^n \omega_i g(x_i),$$

where g is smooth and W is the weight function, which is to be associated with a given orthogonal polynomial.

Gaussian Quadrature

Weight Functions.

The weight function W is non-negative in the integration interval $x \in [a, b]$ such that for any $n \geq 0$ $\int_a^b |x|^n W(x) dx$ is integrable.

The naming weight function arises from the fact that it may be used to give more emphasis to one part of the interval than another.

Weight function	Interval	Polynomial
$W(x) = 1$	$x \in [a, b]$	Legendre
$W(x) = e^{-x^2}$	$-\infty \leq x \leq \infty$	Hermite
$W(x) = e^{-x}$	$0 \leq x \leq \infty$	Laguerre
$W(x) = 1/(\sqrt{1-x^2})$	$-1 \leq x \leq 1$	Chebyshev

Legendre

$$I = \int_{-1}^1 f(x) dx$$

$$C(1-x^2)P - m_l^2 P + (1-x^2) \frac{d}{dx} \left((1-x^2) \frac{dP}{dx} \right) = 0.$$

C is a constant. For $m_l = 0$ we obtain the Legendre polynomials as solutions, whereas $m_l \neq 0$ yields the so-called associated Legendre polynomials. The corresponding polynomials P are

$$L_k(x) = \frac{1}{2^k k!} \frac{d^k}{dx^k} (x^2 - 1)^k \quad k = 0, 1, 2, \dots,$$

which, up to a factor, are the Legendre polynomials L_k . The latter fulfil the orthogonality relation

$$\int_{-1}^1 L_i(x) L_j(x) dx = \frac{2}{2i+1} \delta_{ij},$$

and the recursion relation

Laguerre

$$I = \int_0^\infty f(x) dx = \int_0^\infty x^\alpha e^{-x} g(x) dx.$$

These polynomials arise from the solution of the differential equation

$$\left(\frac{d^2}{dx^2} - \frac{d}{dx} + \frac{\lambda}{x} - \frac{l(l+1)}{x^2} \right) \mathcal{L}(x) = 0,$$

where l is an integer $l \geq 0$ and λ a constant. They fulfil the orthogonality relation

$$\int_{-\infty}^\infty e^{-x} \mathcal{L}_n(x)^2 dx = 1,$$

and the recursion relation

$$(n+1)\mathcal{L}_{n+1}(x) = (2n+1-x)\mathcal{L}_n(x) - n\mathcal{L}_{n-1}(x).$$

Hermite

In a similar way, for an integral which goes like

$$I = \int_{-\infty}^\infty f(x) dx = \int_{-\infty}^\infty e^{-x^2} g(x) dx.$$

we could use the Hermite polynomials in order to extract weights and mesh points. The Hermite polynomials are the solutions of the following differential equation

$$\frac{d^2 H(x)}{dx^2} - 2x \frac{dH(x)}{dx} + (\lambda - 1)H(x) = 0.$$

They fulfil the orthogonality relation

$$\int_{-\infty}^\infty e^{-x^2} H_n(x)^2 dx = 2^n n! \sqrt{\pi},$$

and the recursion relation

$$H_{n+1}(x) = 2xH_n(x) - 2nH_{n-1}(x).$$

Gaussian Quadrature, general Properties

A quadrature formula

$$\int_a^b W(x) f(x) dx \approx \sum_{i=0}^n \omega_i f(x_i),$$

with $n+1$ distinct quadrature points (mesh points) is called a Gaussian quadrature formula if it integrates all polynomials $p \in P_{2n+1}$ exactly, that is

$$\int_a^b W(x) p(x) dx = \sum_{i=0}^n \omega_i p(x_i),$$

It is assumed that $W(x)$ is continuous and positive and that the integral

$$\int_a^b W(x) dx,$$

exists. Note that the replacement of $f \rightarrow Wg$ is normally a better

Numerical integration: A simple example

We want to compute

$$I = \int_0^{\infty} x \exp(-x) \sin x = \frac{1}{2},$$

using brute force Trapezoidal rule, Simpson's rule, Gauss-Legendre, Gauss-Laguerre and Gauss-Legendre again but with a smarter mapping.

- Before we start it can be useful to study the integrand.
- How should we pick the integration limits?

Simple integral

Approximate

$$\int_0^{\infty} f(x) dx \approx \int_0^{\Lambda} f(x) dx$$

```
int n;  
double a, b, alf, xx;  
cout << "Read in the number of integration points" << endl;  
cin >> n;  
cout << "Read in integration limits" << endl;  
cin >> a >> b;
```

Simple integral

```
// reserve space in memory for vectors containing the mesh points  
// weights and function values for the use of the gauss-legendre  
// method  
double *x = new double [n];  
double *w = new double [n];  
// Gauss-Laguerre is old-fashioned translation of FT7 --> C++  
// arrays start at 1 and end at n  
double *xgl = new double [n+1];  
double *wgl = new double [n+1];
```

Simple integral

Note the parameter alf in $x^{\text{alpha}} \exp -x$

```
// set up the mesh points and weights  
gauleg(a, b, x, w, n);  
// set up the mesh points and weights  
alf = 1.0; // <--- Note alf  
gauss_laguerre(xgl, wgl, n, alf);  
// evaluate the integral with the Gauss-Legendre method  
// Note that we initialize the sum. Here brute force gauleg  
double int_gauss = 0;  
for (int i = 0; i < n; i++){  
    int_gauss += w[i] * int_function(x[i]);  
}
```

Simple integral, importance integration/sampling

```
// evaluate the integral with the Gauss-Laguerre method  
// Note that we initialize the sum  
double int_gausslag = 0;  
for (int i = 1; i <= n; i++){  
    int_gausslag += wgl[i] * sin(xgl[i]);  
}
```

Simple integral

```
// evaluate the integral with the Gauss-Laguerre method  
// Here we change the mesh points with a mapping  
// Need to call gauleg from -1 to +1  
gauleg(-1.0, 1.0, x, w, n);  
double pi_4 = acos(-1.0) * 0.25;  
for (int i = 0; i < n; i++){  
    xx = pi_4 * (x[i] + 1.0);  
    r[i] = tan(xx);  
    s[i] = pi_4 / (cos(xx) * cos(xx)) * w[i];  
  
    double int_gausslegimproved = 0;  
    for (int i = 0; i < n; i++){  
        int_gausslegimproved += s[i] * int_function(r[i]);  
    }
```

A six-dimensional integral, project 3 2010

The ansatz for the wave function for two electrons is given by the product of two 1s wave functions as

$$\Psi(\mathbf{r}_1, \mathbf{r}_2) = e^{-\alpha(r_1+r_2)}.$$

Note that it is not possible to find a closed form solution to Schrödinger's equation for two interacting electrons in the helium atom.

The integral we need to solve is the quantum mechanical expectation value of the correlation energy between two electrons, namely

$$\left\langle \frac{1}{|\mathbf{r}_1 - \mathbf{r}_2|} \right\rangle = \int_{-\infty}^{\infty} d\mathbf{r}_1 d\mathbf{r}_2 e^{-2\alpha(r_1+r_2)} \frac{1}{|\mathbf{r}_1 - \mathbf{r}_2|} = \frac{5\pi^2}{16^2} = 0.192765711. \quad (75)$$

Note that our wave function is not normalized. There is a normalization factor missing.

Brute force, six-dimensional integral, Gauss-Legendre

```
double *x = new double [N];
double *w = new double [N];
// set up the mesh points and weights
gauleg(a,b,x,w, N);

// evaluate the integral with the Gauss-Legendre method
// Note that we initialize the sum
double int_gauss = 0.;
for (int i=0; i<N; i++){
  for (int j = 0; j<N; j++){
    for (int k = 0; k<N; k++){
      for (int l = 0; l<N; l++){
        for (int m = 0; m<N; m++){
          for (int n = 0; n<N; n++){
            int_gauss+=w[i]*w[j]*w[k]*w[l]*w[m]*w[n]
              *int_function(x[i],x[j],x[k],x[l],x[m],x[n]);
          }}}}}
}
```

The six-dimensional integral with Gauss-Legendre

```
// this function defines the function to integrate
double int_function(double x1, double y1, double z1, double x2, double y2, double z2)
{
  double alpha = 2.;
  // evaluate the different terms of the exponential
  double exp1=-2*alpha*sqrt(x1*x1+y1*y1+z1*z1);
  double exp2=-2*alpha*sqrt(x2*x2+y2*y2+z2*z2);
  double deno=sqrt(pow((x1-x2),2)+pow((y1-y2),2)+pow((z1-z2),2));
  if(deno < pow(10.,-6.)) { return 0;}
  else return exp(exp1+exp2)/deno;
} // end of function to evaluate
```

Switch to spherical coordinates

Useful to change to spherical coordinates

$$d\mathbf{r}_1 d\mathbf{r}_2 = r_1^2 dr_1 r_2^2 dr_2 d\cos(\theta_1) d\cos(\theta_2) d\phi_1 d\phi_2,$$

and

$$\frac{1}{r_{12}} = \frac{1}{\sqrt{r_1^2 + r_2^2 - 2r_1 r_2 \cos(\beta)}}$$

with

$$\cos(\beta) = \cos(\theta_1) \cos(\theta_2) + \sin(\theta_1) \sin(\theta_2) \cos(\phi_1 - \phi_2))$$

Switch to spherical coordinates

This means that our integral becomes

$$\int_0^{\infty} r_1^2 dr_1 \int_0^{\infty} r_2^2 dr_2 \int_0^{\pi} d\cos(\theta_1) \int_0^{\pi} d\cos(\theta_2) \int_0^{2\pi} d\phi_1 \int_0^{2\pi} d\phi_2 \times$$

$$\frac{e^{-2\alpha(r_1+r_2)}}{\sqrt{r_1^2 + r_2^2 - 2r_1 r_2 \cos(\theta_1) \cos(\theta_2) + \sin(\theta_1) \sin(\theta_2) \cos(\phi_1 - \phi_2)}}$$

since

$$\frac{1}{r_{12}} = \frac{1}{\sqrt{r_1^2 + r_2^2 - 2r_1 r_2 \cos(\beta)}}$$

with

$$\cos(\beta) = \cos(\theta_1) \cos(\theta_2) + \sin(\theta_1) \sin(\theta_2) \cos(\phi_1 - \phi_2))$$

Adaptive methods

Before we abandon totally methods like the trapezoidal rule, we mention briefly how an adaptive integration method can be implemented.

The above methods are all based on a defined step length, normally provided by the user, dividing the integration domain with a fixed number of subintervals. This is rather simple to implement may be inefficient, in particular if the integrand varies considerably in certain areas of the integration domain. In these areas the number of fixed integration points may not be adequate. In other regions, the integrand may vary slowly and fewer integration points may be needed.

Adaptive methods

In order to account for such features, it may be convenient to first study the properties of integrand, via for example a plot of the function to integrate. If this function oscillates largely in some specific domain we may then opt for adding more integration points to that particular domain. However, this procedure needs to be repeated for every new integrand and lacks obviously the advantages of a more generic code.

Adaptive methods

Assume that we want to compute an integral using say the trapezoidal rule. We limit ourselves to a one-dimensional integral. Our integration domain is defined by $x \in [a, b]$. The algorithm goes as follows

Step 1. We compute our first approximation by computing the integral for the full domain. We label this as $I^{(0)}$. It is obtained by calling our previously discussed function `trapezoidal_rule` as

```
I0 = trapezoidal_rule(a, b, n, function);
```

Step 2. We split the integration in two, with $c = (a + b)/2$. We compute then the two integrals $I^{(1L)}$ and $I^{(1R)}$

```
I1L = trapezoidal_rule(a, c, n, function);
```

and

```
I1R = trapezoidal_rule(c, b, n, function);
```

With a given defined tolerance, being a small number provided by us, we estimate the difference $|I^{(1L)} + I^{(1R)} - I^{(0)}| < \text{tolerance}$. If

Adaptive methods

```
// Simple recursive function that implements the
// adaptive integration using the trapezoidal rule
const int maxrecursions = 50;
const double tolerance = 1.0E-10;
// Takes as input the integration limits, number of points, func
// and the number of steps
void adaptive_integration(double a, double b, double *Integral, int
if ( steps > maxrecursions){
    cout << "Too many recursive steps, the function varies too
    break;
```

Adaptive methods

```
double c = (a+b)*0.5;
// the whole integral
double I0 = trapezoidal_rule(a, b,n, func);
// the left half
double I1L = trapezoidal_rule(a, c,n, func);
// the right half
double I1R = trapezoidal_rule(c, b,n, func);
if (fabs(I1L+I1R-I0) < tolerance ) integral = I0;
else
{
    adaptive_integration(a, c, integral, int n, ++steps, func)
    adaptive_integration(c, b, integral, int n, ++steps, func)
}
// end function Adaptive integration
```

The variables **Integral** and **steps** should be initialized to zero by the function that calls the adaptive procedure.

Overview of week 44

Monte Carlo methods.

- Monday: Repetition from last week
- Parallelization using MPI
- Introduction to Monte Carlo methods and Monte Carlo integration
- Tuesday:
- Monte Carlo integration and importance sampling
- Discussion of project 4.

Going Parallel: Divide et impera - Divide and Conquer (accordingly after Julius Caesar)

We will first meet the concept of

Task parallelism: the work of a global problem can be divided into a number of independent tasks, which rarely need to synchronize. Monte Carlo simulation or integrations are examples of this. It is almost embarrassingly trivial to parallelize Monte Carlo and numerical integration codes.

We will use MPI=Message Passing Interface. MPI is a message-passing library where all the routines have corresponding C/C++-binding

```
MPI_Command_name
```

and Fortran-binding (routine names are in uppercase, but can also be in lower case)

```
MPI_COMMAND_NAME
```


What is Message Passing Interface (MPI)? Yet another library!

MPI is a library, not a language. It specifies the names, calling sequences and results of functions or subroutines to be called from C or Fortran programs, and the classes and methods that make up the MPI C++ library. The programs that users write in Fortran, C or C++ are compiled with ordinary compilers and linked with the MPI library.

MPI is a specification, not a particular implementation. MPI programs should be able to run on all possible machines and run all MPI implementations without change.

An MPI computation is a collection of processes communicating with messages.

See chapter 5.5 of lecture notes for more details.

MPI

MPI is a library specification for the message passing interface, proposed as a standard.

- independent of hardware;
- not a language or compiler specification;
- not a specific implementation or product. A message passing standard for portability and ease-of-use. Designed for high performance. Insert communication and synchronization functions where necessary.

Demands from the HPC community

In the field of scientific computing, there is an ever-lasting wish to do larger simulations using shorter computer time.

Development of the capacity for single-processor computers can hardly keep up with the pace of scientific computing:

- processor speed
- memory size/speed Solution: parallel computing!

The basic ideas of parallel computing

- Pursuit of shorter computation time and larger simulation size gives rise to parallel computing.
- Multiple processors are involved to solve a global problem.
- The essence is to divide the entire computation evenly among collaborative processors. Divide and conquer.

A rough classification of hardware models

* Conventional single-processor computers can be called SISD (single-instruction-single-data) machines.

- SIMD (single-instruction-multiple-data) machines incorporate the idea of parallel processing, which use a large number of processing units to execute the same instruction on different data.
- Modern parallel computers are so-called MIMD (multiple-instruction-multiple-data) machines and can execute different instruction streams in parallel on different data.

Shared memory and distributed memory

- One way of categorizing modern parallel computers is to look at the memory configuration.
- In shared memory systems the CPUs share the same address space. Any CPU can access any data in the global memory.
- In distributed memory systems each CPU has its own memory. The CPUs are connected by some network and may exchange messages.

Different parallel programming paradigms

- **Task parallelism:** the work of a global problem can be divided into a number of independent tasks, which rarely need to synchronize. Monte Carlo simulation is one example. Integration is another. However this paradigm is of limited use.
- **Data parallelism:** use of multiple threads (e.g. one thread per processor) to dissect loops over arrays etc. This paradigm requires a single memory address space. Communication and synchronization between processors are often hidden, thus easy to program. However, the user surrenders much control to a specialized compiler. Examples of data parallelism are compiler-based parallelization and OpenMP directives.

Today's situation of parallel computing

- Distributed memory is the dominant hardware configuration. There is a large diversity in these machines, from MPP (massively parallel processing) systems to clusters of off-the-shelf PCs, which are very cost-effective.
- Message-passing is a mature programming paradigm and widely accepted. It often provides an efficient match to the hardware. It is primarily used for the distributed memory systems, but can also be used on shared memory systems. In these lectures we consider only message-passing for writing parallel programs.

Overhead present in parallel computing

- **Uneven load balance:** not all the processors can perform useful work at all time.
- **Overhead of synchronization.**
- **Overhead of communication.**
- Extra computation due to parallelization. Due to the above overhead and that certain part of a sequential algorithm cannot be parallelized we may not achieve an optimal parallelization.

Parallelizing a sequential algorithm

- Identify the part(s) of a sequential algorithm that can be executed in parallel. This is the difficult part,
- Distribute the global work and data among P processors.

Process and processor

- We refer to process as a logical unit which executes its own code, in an MIMD style.
- The processor is a physical device on which one or several processes are executed.
- The MPI standard uses the concept process consistently throughout its documentation.

Bindings to MPI routines

MPI is a message-passing library where all the routines have corresponding C/C++-binding

```
MPI_Command_name
```

and Fortran-binding (routine names are in uppercase, but can also be in lower case)

```
MPI_COMMAND_NAME
```

The discussion in these slides focuses on the C++ binding.

The most important/used MPI routines

- `MPI_Init` - initiate an MPI computation
- `MPI_Finalize` - terminate the MPI computation and clean up
- `MPI_Comm_size` - how many processes participate in a given MPI communicator?
- `MPI_Comm_rank` - which one am I? (A number between 0 and size-1.)
- `MPI_Reduce`(`Allreduce`) - Collect data from all nodes and either sum them up in one or all (`Allreduce`). Useful for numerical integration
- `MPI_Send` - send a message to a particular process within an MPI communicator
- `MPI_Recv` - receive a message from a particular process within an MPI communicator

Note the `MPI_COMM_WORLD` declaration

- A group of MPI processes with a name (context).
- Any process is identified by its rank. The rank is only meaningful within a particular communicator.
- By default communicator `MPI_COMM_WORLD` contains all the MPI processes.
- Mechanism to identify subset of processes.
- Promotes modular design of parallel libraries.

The first MPI C/C++ program

Let every process write "Hello world" on the standard output. This is `program2.cpp` of chapter 4.

```
using namespace std;
#include <mpi.h>
#include <iostream>
int main (int nargs, char* args[])
{
    int numprocs, my_rank;
    // MPI initializations
    MPI_Init (&nargs, &args);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &my_rank);
    cout << "Hello world, I have rank " << my_rank << " out of " << numprocs << endl;
    // End MPI
    MPI_Finalize ();
}
```

The Fortran program

```
PROGRAM hello
INCLUDE "mpif.h"
INTEGER:: size, my_rank, ierr

CALL MPI_INIT(ierr)
CALL MPI_COMM_SIZE(MPI_COMM_WORLD, size, ierr)
CALL MPI_COMM_RANK(MPI_COMM_WORLD, my_rank, ierr)
WRITE(*,*)"Hello world, I've rank ",my_rank," out of ",size
CALL MPI_FINALIZE(ierr)

END PROGRAM hello
```

Note 1

The output to screen is not ordered since all processes are trying to write to screen simultaneously. It is then the operating system which opts for an ordering. If we wish to have an organized output, starting from the first process, we may rewrite our program as in the next example (`program3.cpp`), see again chapter 5.7 of lecture notes.

Ordered output with `MPI_Barrier`

```
int main (int nargs, char* args[])
{
    int numprocs, my_rank, i;
    MPI_Init (&nargs, &args);
    MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
    MPI_Comm_rank (MPI_COMM_WORLD, &my_rank);
    for (i = 0; i < numprocs; i++) {}
    MPI_Barrier (MPI_COMM_WORLD);
    if (i == my_rank) {
        cout << "Hello world, I have rank " << my_rank << " out of " << numprocs << endl;
        MPI_Finalize ();
    }
}
```

Note 2

Here we have used the `MPI_Barrier` function to ensure that every process has completed its set of instructions in a particular order. A barrier is a special collective operation that does not allow the processes to continue until all processes in the communicator (here `MPI_COMM_WORLD` have called `MPI_Barrier`. The barriers make sure that all processes have reached the same point in the code. Many of the collective operations like `MPI_ALLREDUCE` to be discussed later, have the same property; viz. no process can exit the operation until all processes have started. However, this is slightly more time-consuming since the processes synchronize between themselves as many times as there are processes. In the next Hello world example we use the send and receive functions in order to have a synchronized action.

Strategies


- Develop codes locally, run with some few processes and test your codes. Do benchmarking, timing and so forth on local nodes, for example your laptop. You can install MPICH2 on your laptop (most new laptops come with dual cores). You can test with one node at the lab.
- When you are convinced that your codes run correctly, you start your production runs on available supercomputers. At UiO we have a new machine among the top 100, Abel.

How do I run MPI on the machines at the lab (MPICH2)

The machines at the lab are all quad-cores

Step 1. Compile with `mpicxx` or `mpic++`

Step 2. Set up collaboration between processes and run

```
mpd --ncpus=4 &  run code with mpiexec -n 4 ./nameofprog
```

Here we declare that we will use 4 processes via the `-ncpus` option and via `-n 4` when running.

Step 3. End with

```
mpdallexit
```

Can I do it on my own PC/laptop?

Of course:

- go to <http://www.mcs.anl.gov/research/projects/mpich2/>
- follow the instructions and install it on your own PC/laptop. It works on Windows, Mac and Linux. For Linux (Ubuntu, Linux Mint), go to synaptic package manager and search for MPI. You will get several options.

Ordered output with `MPI_Recv` and `MPI_Send`

```
.....
int numprocs, my_rank, flag;
MPI_Status status;
MPI_Init (&nargs, &args);
MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
MPI_Comm_rank (MPI_COMM_WORLD, &my_rank);
if (my_rank > 0)
MPI_Recv (&flag, 1, MPI_INT, my_rank-1, 100,
MPI_COMM_WORLD, &status);
cout << "Hello world, I have rank " << my_rank << " out of "
<< numprocs << endl;
if (my_rank < numprocs-1)
MPI_Send (&my_rank, 1, MPI_INT, my_rank+1,
100, MPI_COMM_WORLD);
MPI_Finalize ();
```

Note 3

The basic sending of messages is given by the function `MPI_SEND`, which in C/C++ is defined as

```
int MPI_Send(void *buf, int count,
MPI_Datatype datatype,
int dest, int tag, MPI_Comm comm);
```

This single command allows the passing of any kind of variable, even a large array, to any group of tasks. The variable **buf** is the variable we wish to send while **count** is the number of variables we are passing. If we are passing only a single value, this should be 1. If we transfer an array, it is the overall size of the array. For example, if we want to send a 10 by 10 array, count would be $10 \times 10 = 100$ since we are actually passing 100 values.

Note 4

Once you have sent a message, you must receive it on another task. The function `MPI_RECV` is similar to the send call.

```
int MPI_Recv( void *buf, int count, MPI_Datatype datatype,
             int source,
             int tag, MPI_Comm comm, MPI_Status *status )
```

The arguments that are different from those in `MPI_SEND` are **buf** which is the name of the variable where you will be storing the received data, **source** which replaces the destination in the send command. This is the return ID of the sender.

Finally, we have used `MPI_Status` `status` where one can check if the receive was completed.

The output of this code is the same as the previous example, but now process 0 sends a message to process 1, which forwards it further to process 2, and so forth.

Armed with this wisdom, performed all hello world greetings, we are now ready for serious work.

Integrating π

Examples.

- Go to the program package
- Go to the MPI directory and then chapter 5
- Look at `program6.cpp`.
- This code computes π using the trapezoidal rule.

Integration algos

The trapezoidal rule (`program6.cpp`)

$$I = \int_a^b f(x)dx = h(f(a)/2 + f(a+h) + f(a+2h) + \dots + f(b-h) + f(b)/2)$$

Another very simple approach is the so-called midpoint or rectangle method. In this case the integration area is split in a given number of rectangles with length h and height given by the mid-point value of the function. This gives the following simple rule for approximating an integral

$$I = \int_a^b f(x)dx \approx h \sum_{i=1}^N f(x_{i-1/2}),$$

where $f(x_{i-1/2})$ is the midpoint value of f for a given rectangle. This is used in `program5.cpp`.

MPI_reduce and MPI_Allreduce

The parallel integration MPI instructions are simple if we use the functions `MPI_Reduce` or `MPI_Allreduce`. The first function takes information from all processes and sends the result of the MPI operation to one process only, typically the master node. If we use `MPI_Allreduce`, the result is sent back to all processes, a feature which is useful when all nodes need the value of a joint operation. We limit ourselves to `MPI_Reduce` since it is only one process which will print out the final number of our calculation, The arguments to `MPI_Allreduce` are the same.

MPI_reduce

Call as

```
MPI_reduce( void *senddata, void* resultdata, int count,
            MPI_Datatype datatype, MPI_Op, int root, MPI_Comm comm)
```

The two variables `senddata` and `resultdata` are obvious, besides the fact that one sends the address of the variable or the first element of an array. If they are arrays they need to have the same size. The variable `count` represents the total dimensionality, 1 in case of just one variable, while `MPI_Datatype` defines the type of variable which is sent and received.

The new feature is `MPI_Op`. It defines the type of operation we want to do. In our case, since we are summing the rectangle contributions from every process we define `MPI_Op = MPI_SUM`. If we have an array or matrix we can search for the largest or smallest element by sending either `MPI_MAX` or `MPI_MIN`. If we want the location as well (which array element) we simply transfer `MPI_MAXLOC` or `MPI_MINLOC`. If we want the product we write `MPI_PROD`.

Dissection of example `program6.cpp`

```
// Trapezoidal rule and numerical integration using MPI, example
using namespace std;
#include <mpi.h>
#include <iostream>

// Here we define various functions called by the main program
double int_function(double );
double trapezoidal_rule(double , double , int , double (*)(double))

// Main function begins here
int main (int nargs, char* args[])
{
    int n, local_n, numprocs, my_rank;
    double a, b, h, local_a, local_b, total_sum, local_sum;
    double time_start, time_end, total_time;
```

Dissection of example program6.cpp

```
// MPI initializations
MPI_Init (&nargs, &args);
MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
MPI_Comm_rank (MPI_COMM_WORLD, &my_rank);
time_start = MPI_Wtime();
// Fixed values for a, b and n
a = 0.0 ; b = 1.0; n = 1000;
h = (b-a)/n; // h is the same for all processes
local_n = n/numprocs;
// make sure n > numprocs, else integer division gives zero
// Length of each process' interval of
// integration = local_n*h;
local_a = a + my_rank*local_n*h;
local_b = local_a + local_n*h;
```

Dissection of example program6.cpp

```
total_sum = 0.0;
local_sum = trapezoidal_rule(local_a, local_b, local_n,
                             &int_function);
MPI_Reduce(&local_sum, &total_sum, 1, MPI_DOUBLE,
           MPI_SUM, 0, MPI_COMM_WORLD);
time_end = MPI_Wtime();
total_time = time_end-time_start;
if ( my_rank == 0 ) {
    cout << "Trapezoidal rule = " << total_sum << endl;
    cout << "Time = " << total_time
         << " on number of processors: " << numprocs << endl;

    // End MPI
    MPI_Finalize ();
    return 0;
} // end of main program
```

Dissection of example program6.cpp

We use MPI_reduce to collect data from each process. Note also the use of the function MPI_Wtime. The final functions are

```
// this function defines the function to integrate
double int_function(double x)
{
    double value = 4./(1.+x*x);
    return value;
} // end of function to evaluate
```

Dissection of example program6.cpp

Implementation of the trapezoidal rule.

```
// this function defines the trapezoidal rule
double trapezoidal_rule(double a, double b, int n,
                        double (*func)(double))
{
    double trapez_sum;
    double fa, fb, x, step;
    int j;
    step=(b-a)/((double) n);
    fa=(*func)(a)/2. ;
    fb=(*func)(b)/2. ;
    trapez_sum=0.;
    for (j=1; j <= n-1; j++){
        x=j*step+a;
        trapez_sum+=(*func)(x);

        trapez_sum=(trapez_sum+fb+fa)*step;
    } // end trapezoidal_rule
```

Plan for Monte Carlo Lectures, chapters 11-14 in Lecture notes

- This week: intro, MC integration and probability distribution functions (PDFs)
- Next week: More on integration, PDFs, MC integration and random walks.
- Third week: random walks and statistical physics.
- Fourth week: Statistical physics.
- Fifth week: Most likely quantum Monte Carlo Approximately from this week till the end of November.

Monte Carlo Keywords

Consider it is a numerical experiment

- Be able to generate random variables following a given probability distribution function PDF
- Find a probability distribution function (PDF).
- Sampling rule for accepting a move
- Compute standard deviation and other expectation values
- Techniques for improving errors Enhances algorithmic thinking!

Probability Distribution Functions PDF

	Discrete PDF	continuous PDF
Domain	$\{x_1, x_2, x_3, \dots, x_N\}$	$[a, b]$
probability	$p(x_i)$	$p(x)dx$
Cumulative	$P_i = \sum_{j=1}^i p(x_j)$	$P(x) = \int_a^x p(t)dt$
Positivity	$0 \leq p(x_i) \leq 1$	$p(x) \geq 0$
Positivity	$0 \leq P_i \leq 1$	$0 \leq P(x) \leq 1$
Monotonuous	$P_i \geq P_j$ if $x_i \geq x_j$	$P(x_i) \geq P(x_j)$ if $x_i \geq x_j$
Normalization	$P_N = 1$	$P(b) = 1$

Expectation Values

Discrete PDF

$$E[x^k] = \langle x^k \rangle = \frac{1}{N} \sum_{i=1}^N x_i^k p(x_i),$$

provided that the sums (or integrals) $\sum_{i=1}^N p(x_i)$ converge absolutely (viz, $\sum_{i=1}^N |p(x_i)|$ converges)

Continuous PDF

$$E[x^k] = \langle x^k \rangle = \int_a^b x^k p(x) dx,$$

Function $f(x)$

$$E[f^k] = \langle f^k \rangle = \int_a^b f^k p(x) dx,$$

Variance

Important PDFs

uniform distribution

$$p(x) = \frac{1}{b-a} \Theta(x-a) \Theta(b-x),$$

which gives for $a=0, b=1$ $p(x) = 1$ for $x \in [0, 1]$ and zero else.

exponential distribution

$$p(x) = \alpha e^{-\alpha x},$$

with probability different from zero in $[0, \infty]$

normal distribution (Gaussian)

$$p(x) = \frac{1}{\sqrt{2\pi\sigma^2}} \exp\left(-\frac{(x-\mu)^2}{2\sigma^2}\right)$$

with probability different from zero in $[-\infty, \infty]$

All random number generators use the uniform distribution for $x \in [0, 1]$

Why Monte Carlo?

An example from quantum mechanics: most problems of interest in e.g., atomic, molecular, nuclear and solid state physics consist of a large number of interacting electrons and ions or nucleons. The total number of particles N is usually sufficiently large that an exact solution cannot be found. Typically, the expectation value for a chosen hamiltonian for a system of N particles is

$$\langle H \rangle =$$

$$\frac{\int d\mathbf{R}_1 d\mathbf{R}_2 \dots d\mathbf{R}_N \Psi^*(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N) H(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N) \Psi(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N)}{\int d\mathbf{R}_1 d\mathbf{R}_2 \dots d\mathbf{R}_N \Psi^*(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N) \Psi(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N)}$$

an in general intractable problem.

This integral is actually the starting point in a Variational Monte Carlo calculation. **Gaussian quadrature: Forget it!** given 10 particles and 10 mesh points for each degree of freedom and an ideal 1 petaflops machine (all operations take the same time), how

More on dimensionality

As an example from the nuclear many-body problem, we have Schrödinger's equation as a differential equation

$$\hat{H}\Psi(\mathbf{r}_1, \dots, \mathbf{r}_A, \alpha_1, \dots, \alpha_A) = E\Psi(\mathbf{r}_1, \dots, \mathbf{r}_A, \alpha_1, \dots, \alpha_A)$$

where

$$\mathbf{r}_1, \dots, \mathbf{r}_A,$$

are the coordinates and

$$\alpha_1, \dots, \alpha_A,$$

are sets of relevant quantum numbers such as spin and isospin for a system of A nucleons ($A = N + Z$, N being the number of neutrons and Z the number of protons).

Even more on dimensionality

There are

$$2^A \times \binom{A}{Z}$$

coupled second-order differential equations in $3A$ dimensions.

For a nucleus like ^{10}Be this number is **215040**. This is a truly challenging many-body problem.

But what do we gain by Monte Carlo Integration?

A crude approach consists in setting all weights equal 1, $\omega_i = 1$. With $dx = h = (b - a)/N$ where $b = 1$, $a = 0$ in our case and h is the step size. The integral

$$I = \int_0^1 f(x) dx \approx \frac{1}{N} \sum_{i=1}^N f(x_i),$$

can be rewritten using the concept of the average of the function f for a given PDF $p(x)$ as

$$E[f] = \langle f \rangle = \frac{1}{N} \sum_{i=1}^N f(x_i) p(x_i),$$

and identify $p(x)$ with the uniform distribution, viz $p(x) = 1$ when $x \in [0, 1]$ and zero for all other values of x . The integral is then the average of f over the interval $x \in [0, 1]$

$$I = \int_0^1 f(x) dx \approx E[f] = \langle f \rangle$$

But what do we gain by Monte Carlo Integration?

In addition to the average value $\langle f \rangle$ the other important quantity in a Monte-Carlo calculation is the variance σ^2 and the standard deviation σ . We define first the variance of the integral with f for a uniform distribution in the interval $x \in [0, 1]$ to be

$$\sigma_f^2 = \frac{1}{N} \sum_{i=1}^N (f(x_i) - \langle f \rangle)^2 p(x_i),$$

and inserting the uniform distribution this yields

$$\sigma_f^2 = \frac{1}{N} \sum_{i=1}^N f(x_i)^2 - \left(\frac{1}{N} \sum_{i=1}^N f(x_i) \right)^2,$$

or

$$\sigma_f^2 = E[f^2] - (E[f])^2 = (\langle f^2 \rangle - \langle f \rangle^2).$$

which is nothing but a measure of the extent to which f deviates from its average over the region of integration. The standard

But what do we gain by Monte Carlo Integration?

If we consider the above results for a fixed value of N as a measurement, we could however recalculate the above average and variance for a series of different measurements. If each such measurement produces a set of averages for the integral I denoted $\langle f \rangle_I$, we have for M measurements that the integral is given by

$$\langle I \rangle_M = \frac{1}{M} \sum_{I=1}^M \langle f \rangle_I.$$

If we can consider the probability of correlated events to be zero, we can rewrite the variance of these series of measurements as (equating $M = N$)

$$\sigma_N^2 \approx \frac{1}{N} (\langle f^2 \rangle - \langle f \rangle^2) = \frac{\sigma_f^2}{N}.$$

We note that the standard deviation is proportional with the inverse square root of the number of measurements

But what do we gain by Monte Carlo Integration?

- We saw that the trapezoidal rule carries a truncation error $O(h^2)$, with h the step length.
 - Quadrature rules such as Newton-Cotes have a truncation error which goes like $\sim O(h^k)$, with $k \geq 1$. Recalling that the step size is defined as $h = (b - a)/N$, we have an error which goes like $\sim N^{-k}$.
- Monte Carlo integration is more efficient in higher dimensions. Assume that our integration volume is a hypercube with side L and dimension d . This cube contains hence $N = (L/h)^d$ points and therefore the error in the result scales as $N^{-k/d}$ for the traditional methods.
- The error in the Monte Carlo integration is however independent of d and scales as $\sigma \sim 1/\sqrt{N}$, always!

* Comparing this with traditional methods, shows that Monte Carlo integration is more efficient than an order- k algorithm when $d > 2k$

Some simple examples: Example 1: Particles in a Box

Consider a box divided into two equal halves separated by a wall. At the beginning, time $t = 0$, there are N particles on the left side. A small hole in the wall is then opened and one particle can pass through the hole per unit time.

After some time the system reaches its equilibrium state with equally many particles in both halves, $N/2$. Instead of determining complicated initial conditions for a system of N particles, we model the system by a simple statistical model. In order to simulate this system, which may consist of $N \gg 1$ particles, we assume that all particles in the left half have equal probabilities of going to the right half.

Particles in a Box

We introduce the label n_l to denote the number of particles at every time on the left side, and $n_r = N - n_l$ for those on the right side. The probability for a move to the right during a time step Δt is n_l/N . The algorithm for simulating this problem may then look like as follows

- Choose the number of particles N .
- Make a loop over time, where the maximum time should be larger than the number of particles N .
- For every time step Δt there is a probability n_l/N for a move to the right. Compare this probability with a random number x .
- If $x \leq n_l/N$, decrease the number of particles in the left half by one, i.e., $n_l = n_l - 1$. Else, move a particle from the right half to the left, i.e., $n_l = n_l + 1$. **This is our sampling rule**
- Increase the time by one unit (the external loop). In this case, a Monte Carlo sample corresponds to one time unit Δt .

Particles in a Box

```
// setup of initial conditions
nleft = initial_n_particles;
max_time = 10*initial_n_particles;
idum = -1;
// sampling over number of particles
for( time=0; time <= max_time; time++){
    random_n = ((int) initial_n_particles*ran0(&idum));
    if ( random_n <= nleft){
        nleft -= 1;
    }
    else{
        nleft += 1;
    }
    ofile << setw(15) << time;
    ofile << setw(15) << nleft << endl;
}
```

Example 2: Radioactive Decay

Assume that at the time $t = 0$ we have $N(0)$ nuclei of type X which can decay radioactively. At a time $t > 0$ we are left with $N(t)$ nuclei. With a transition probability ω , which expresses the probability that the system will make a transition to another state during a time step of one second, we have the following first-order differential equation

$$dN(t) = -\omega N(t)dt,$$

whose solution is

$$N(t) = N(0)e^{-\omega t},$$

where we have defined the mean lifetime τ of X as

$$\tau = \frac{1}{\omega}.$$

If a nucleus X decays to a daughter nucleus Y which also can

Radioactive Decay

Probability for a decay of a particle during a time step Δt is

$$\frac{\Delta N(t)}{N(t)\Delta t} = -\lambda$$

λ is inversely proportional to the lifetime

- Choose the number of particles $N(t = 0) = N_0$.
- Make a loop over the number of time steps, with maximum time bigger than the number of particles N_0
- At every time step there is a probability λ for decay. Compare this probability with a random number x .
- If $x \leq \lambda$, reduce the number of particles with one i.e., $N = N - 1$. If not, keep the same number of particles till the next time step. **This is our sampling rule**
- Increase by one the time step (the external loop)

Radioactive Decay

```
idum=-1; // initialise random number generator
// loop over monte carlo cycles
// One monte carlo loop is one sample
for (cycles = 1; cycles <= number_cycles; cycles++){
    n_unstable = initial_n_particles;
    // accumulate the number of particles per time step per trial
    ncumulative[0] += initial_n_particles;
    // loop over each time step
    for (time=1; time <= max_time; time++){
        // for each time step, we check each particle
        particle_limit = n_unstable;
        for ( np = 1; np <= particle_limit; np++) {
            if( ran0(&idum) <= decay_probability) {
                n_unstable=n_unstable-1;
            }
        } // end of loop over particles
        ncumulative[time] += n_unstable;
    } // end of loop over time steps
} // end of loop over MC trials
} // end mc_sampling function
```

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an in general intractable problem.

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or

$$\sigma_f^2 = E[f^2] - (E[f])^2 = (\langle f^2 \rangle - \langle f \rangle^2).$$

which is nothing but a measure of the extent to which f deviates from its average over the region of integration. The standard

But what do we gain by Monte Carlo Integration?

If we consider the above results for a fixed value of N as a measurement, we could however recalculate the above average and variance for a series of different measurements. If each such measurement produces a set of averages for the integral I denoted $\langle f \rangle_I$, we have for M measurements that the integral is given by

$$\langle I \rangle_M = \frac{1}{M} \sum_{I=1}^M \langle f \rangle_I.$$

If we can consider the probability of correlated events to be zero, we can rewrite the variance of these series of measurements as (equating $M = N$)

$$\sigma_N^2 \approx \frac{1}{N} (\langle f^2 \rangle - \langle f \rangle^2) = \frac{\sigma_f^2}{N}.$$

We note that the standard deviation is proportional with the inverse square root of the number of measurements

But what do we gain by Monte Carlo Integration?

- We saw that the trapezoidal rule carries a truncation error $O(h^2)$, with h the step length.
 - Quadrature rules such as Newton-Cotes have a truncation error which goes like $\sim O(h^k)$, with $k \geq 1$. Recalling that the step size is defined as $h = (b - a)/N$, we have an error which goes like $\sim N^{-k}$.
- Monte Carlo integration is more efficient in higher dimensions. Assume that our integration volume is a hypercube with side L and dimension d . This cube contains hence $N = (L/h)^d$ points and therefore the error in the result scales as $N^{-k/d}$ for the traditional methods.
- The error in the Monte Carlo integration is however independent of d and scales as $\sigma \sim 1/\sqrt{N}$, always!

* Comparing this with traditional methods, shows that Monte Carlo integration is more efficient than an order- k algorithm when $d > 2k$

Example: Acceptance-Rejection Method

This is a rather simple and appealing method after von Neumann. Assume that we are looking at an interval $x \in [a, b]$, this being the domain of the PDF $p(x)$. Suppose also that the largest value our distribution function takes in this interval is M , that is

$$p(x) \leq M \quad x \in [a, b].$$

Then we generate a random number x from the uniform distribution for $x \in [a, b]$ and a corresponding number s for the uniform distribution between $[0, M]$. If

$$p(x) \geq s,$$

we accept the new value of x , else we generate again two new random numbers x and s and perform the test in the latter equation again.

Acceptance-Rejection Method

As an example, consider the evaluation of the integral

$$I = \int_0^3 \exp(x) dx.$$

Obviously to derive it analytically is much easier, however the integrand could pose some more difficult challenges. The aim here is simply to show how to implement the acceptance-rejection algorithm. The integral is the area below the curve $f(x) = \exp(x)$. If we uniformly fill the rectangle spanned by $x \in [0, 3]$ and $y \in [0, \exp(3)]$, the fraction below the curve obtained from a uniform distribution, and multiplied by the area of the rectangle, should approximate the chosen integral. It is rather easy to implement this numerically, as shown in the following code.

Acceptance-Rejection Method

```
// Loop over Monte Carlo trials n
integral = 0.;
for ( int i = 1; i <= n; i++){
// Finds a random value for x in the interval [0,3]
x = 3*ran0(&idum);
// Finds y-value between [0,exp(3)]
y = exp(3.0)*ran0(&idum);
// if the value of y at exp(x) is below the curve, we accept
// THIS IS OUR SAMPLING RULE
if ( y < exp(x)) s = s+ 1.0;
// The integral is area enclosed below the line f(x)=exp(x)

// Then we multiply with the area of the rectangle and
// divide by the number of cycles
Integral = 3.*exp(3.)*s/n
```

Monte Carlo Integration

With uniform distribution $p(x) = 1$ for $x \in [0, 1]$ and zero else

$$I = \int_0^1 f(x) dx \approx \frac{1}{N} \sum_{i=1}^N f(x_i),$$

$$I = \int_0^1 f(x) dx \approx E[f] = \langle f \rangle.$$

$$\sigma_f^2 = \frac{1}{N} \sum_{i=1}^N f(x_i)^2 - \left(\frac{1}{N} \sum_{i=1}^N f(x_i) \right)^2,$$

or

$$\sigma_f^2 = E[f^2] - (E[f])^2 = \langle f^2 \rangle - \langle f \rangle^2.$$

Brute Force Algorithm for Monte Carlo Integration

- Choose the number of Monte Carlo samples N .
- Make a loop over N and for every step generate a random number x_i in the interval $x_i \in [0, 1]$ by calling a random number generator.
- Use this number to compute $f(x_i)$.
- Find the contribution to the variance and the mean value for every loop contribution.
- After N samplings, compute the final mean value and the standard deviation

Brute Force Integration

```
// crude mc function to calculate pi
int i, n;
long idum;
double crude_mc, x, sum_sigma, fx, variance;
cout << "Read in the number of Monte-Carlo samples" << endl;
cin >> n;
crude_mc = sum_sigma=0. ; idum=-1 ;
// evaluate the integral with the a crude Monte-Carlo method
for ( i = 1; i <= n; i++){
x=ran0(&idum);
fx=func(x);
crude_mc += fx;
sum_sigma += fx*fx;

crude_mc = crude_mc/((double) n );
sum_sigma = sum_sigma/((double) n );
variance=sum_sigma-crude_mc*crude_mc;
```

Or: another Brute Force Integration

```
// crude mc function to calculate pi
int main()
{
const int n = 1000000;
double x, fx, pi, invers_period, pi2;
int i;
invers_period = 1./RAND_MAX;
srand(time(NULL));
pi = pi2 = 0.;
for (i=0; i<n;i++){
{
x = double(rand())*invers_period;
// This is our sampling rule, all points accepted
fx = 4./(1+x*x);
pi += fx;
pi2 += fx*fx;
}
pi /= n; pi2 = pi2/n - pi*pi;
cout << "pi=" << pi << " sigma^2=" << pi2 << endl;
return 0;
}
```

Brute Force Integration

Note the call to a function which generates random numbers according to the uniform distribution

```
long idum;
idum=-1 ;
....
x=ran0(&idum);
....
```

or

```
...
invers_period = 1./RAND_MAX;
srand(time(NULL));
...
x = double(rand())*invers_period;
```

Algorithm for Monte Carlo Integration, Results

N	I
10 & 3.10263E+00 & 3.98802E-01	
100 & 3.02933E+00 & 4.04822E-01	
1000 & 3.13395E+00 & 4.22881E-01	
10000 & 3.14195E+00 & 4.11195E-01	
100000 & 3.14003E+00 & 4.14114E-01	
1000000	3.14213E+00 & 4.13838E-01
10000000	3.14177E+00 & 4.13523E-01
10^9	3.14162E+00 & 4.13581E-01

We note that as N increases, the integral itself never reaches more than an agreement to the fourth or fifth digit. The variance also oscillates around its exact value $4.13581E-01$. Note well that the variance need not be zero but can, with appropriate redefinitions of the integral be made smaller. A smaller variance yields also a smaller standard deviation. This is the topic of importance sampling.

Transformation of Variables

The starting point is always the uniform distribution

$$p(x)dx = \begin{cases} dx & 0 \leq x \leq 1 \\ 0 & \text{else} \end{cases}$$

with $p(x) = 1$ and satisfying

$$\int_{-\infty}^{\infty} p(x)dx = 1.$$

All random number generators provided in the program library generate numbers in this domain.

When we attempt a transformation to a new variable $x \rightarrow y$ we have to conserve the probability

$$p(y)dy = p(x)dx,$$

which for the uniform distribution implies

Transformation of Variables

Let us assume that $p(y)$ is a PDF different from the uniform PDF $p(x) = 1$ with $x \in [0, 1]$. If we integrate the last expression we arrive at

$$x(y) = \int_0^y p(y')dy',$$

which is nothing but the cumulative distribution of $p(y)$, i.e.,

$$x(y) = P(y) = \int_0^y p(y')dy'.$$

This is an important result which has consequences for eventual improvements over the brute force Monte Carlo.

Example 1, a general Uniform Distribution

Suppose we have the general uniform distribution

$$p(y)dy = \begin{cases} \frac{dy}{b-a} & a \leq y \leq b \\ 0 & \text{else} \end{cases}$$

If we wish to relate this distribution to the one in the interval $x \in [0, 1]$ we have

$$p(y)dy = \frac{dy}{b-a} = dx,$$

and integrating we obtain the cumulative function

$$x(y) = \int_a^y \frac{dy'}{b-a},$$

yielding

$$y = a + (b-a)x,$$

a well-known result!

Example 2, from Uniform to Exponential

Assume that

$$p(y) = e^{-y},$$

which is the exponential distribution, important for the analysis of e.g., radioactive decay. Again, $p(x)$ is given by the uniform distribution with $x \in [0, 1]$, and with the assumption that the probability is conserved we have

$$p(y)dy = e^{-y}dy = dx,$$

which yields after integration

$$x(y) = P(y) = \int_0^y \exp(-y')dy' = 1 - \exp(-y),$$

or

$$y(x) = -\ln(1-x).$$

Example 2, from Uniform to Exponential

This means that if we can factor out $\exp(-y)$ from an integrand we may have

$$I = \int_0^{\infty} F(y)dy = \int_0^{\infty} \exp(-y)G(y)dy$$

which we rewrite as

$$\int_0^{\infty} \exp(-y)G(y)dy = \int_0^{\infty} \frac{dx}{dy}G(y)dy \approx \frac{1}{N} \sum_{i=1}^N G(y(x_i)),$$

where x_i is a random number in the interval $[0, 1]$.

Note that in practical implementations, our random number generators for the uniform distribution never return exactly 0 or 1, but we may come very close. We should thus in principle set $x \in (0, 1)$.

Example 2, from Uniform to Exponential

The algorithm is rather simple. In the function which sets up the integral, we simply need the random number generator for the uniform distribution in order to obtain numbers in the interval $[0,1]$. We obtain y by the taking the logarithm of $(1-x)$. Our calling function which sets up the new random variable y may then include statements like

```
.....
idum=-1;
x=ran0(&idum);
y=-log(1.-x);
.....
```

Example 3

Another function which provides an example for a PDF is

$$p(y)dy = \frac{dy}{(a+by)^n},$$

with $n > 1$. It is normalizable, positive definite, analytically integrable and the integral is invertible, allowing thereby the expression of a new variable in terms of the old one. The integral

$$\int_0^\infty \frac{dy}{(a+by)^n} = \frac{1}{(n-1)ba^{n-1}},$$

gives

$$p(y)dy = \frac{(n-1)ba^{n-1}}{(a+by)^n} dy,$$

which in turn gives the cumulative function

$$P(y) = \int_0^y \frac{(n-1)ba^{n-1}}{(a+bv)^n} dv =$$

Example 4, from Uniform to Normal

For the normal distribution, expressed here as

$$g(x,y) = \exp(-(x^2+y^2)/2) dx dy.$$

it is rather difficult to find an inverse since the cumulative distribution is given by the error function $\text{erf}(x)$.

If we however switch to polar coordinates, we have for x and y

$$r = (x^2 + y^2)^{1/2} \quad \theta = \tan^{-1} \frac{y}{x},$$

resulting in

$$g(r,\theta) = r \exp(-r^2/2) dr d\theta,$$

where the angle θ could be given by a uniform distribution in the region $[0, 2\pi]$. Following example 1 above, this implies simply multiplying random numbers $x \in [0, 1]$ by 2π .

Example 4, from Uniform to Normal

The variable r , defined for $r \in [0, \infty)$ needs to be related to random numbers $x' \in [0, 1]$. To achieve that, we introduce a new variable

$$u = \frac{1}{2}r^2,$$

and define a PDF

$$\exp(-u) du,$$

with $u \in [0, \infty)$. Using the results from example 2, we have that

$$u = -\ln(1-x'),$$

where x' is a random number generated for $x' \in [0, 1]$. With

$$x = r \cos(\theta) = \sqrt{2u} \cos(\theta),$$

and

Example 4, from Uniform to Normal

A function which yields such random numbers for the normal distribution would include statements like

```
.....
idum=-1;
radius=sqrt(-2*ln(1.-ran0(idum)));
theta=2*pi*ran0(idum);
x=radius*cos(theta);
y=radius*sin(theta);
.....
```

Box-Mueller Method for Normal Deviates

```
// random numbers with gaussian distribution
double gaussian_deviate(long * idum)
{
    static int iset = 0;
    static double gset;
    double fac, rsq, v1, v2;
    if ( idum < 0 ) iset = 0;
    if ( iset == 0 ) {
        do {
            v1 = 2.*ran0(idum) - 1.0;
            v2 = 2.*ran0(idum) - 1.0;
            rsq = v1*v1+v2*v2;
        } while (rsq >= 1.0 || rsq == 0.);
        fac = sqrt(-2.*log(rsq)/rsq);
        gset = v1*fac;
        iset = 1;
        return v2*fac;
    } else {
        iset = 0;
        return gset;
    }
}
```

Importance Sampling, chapter 11.4

With the aid of the above variable transformations we address now one of the most widely used approaches to Monte Carlo integration, namely importance sampling.

Let us assume that $p(y)$ is a PDF whose behavior resembles that of a function F defined in a certain interval $[a, b]$. The normalization condition is

$$\int_a^b p(y) dy = 1.$$

We can rewrite our integral as

$$I = \int_a^b F(y) dy = \int_a^b p(y) \frac{F(y)}{p(y)} dy. \quad (76)$$

Importance Sampling

Since random numbers are generated for the uniform distribution $p(x)$ with $x \in [0, 1]$, we need to perform a change of variables $x \rightarrow y$ through

$$x(y) = \int_a^y p(y') dy',$$

where we used

$$p(x) dx = dx = p(y) dy.$$

If we can invert $x(y)$, we find $y(x)$ as well.

Importance Sampling

With this change of variables we can express the integral of Eq. (76) as

$$I = \int_a^b p(y) \frac{F(y)}{p(y)} dy = \int_a^b \frac{F(y(x))}{p(y(x))} dx,$$

meaning that a Monte Carlo evaluation of the above integral gives

$$\int_a^b \frac{F(y(x))}{p(y(x))} dx = \frac{1}{N} \sum_{i=1}^N \frac{F(y(x_i))}{p(y(x_i))}.$$

The advantage of such a change of variables in case $p(y)$ follows closely F is that the integrand becomes smooth and we can sample over relevant values for the integrand. It is however not trivial to find such a function p . The conditions on p which allow us to perform these transformations are

- p is normalizable and positive definite,
- it is analytically integrable and

Importance Sampling

The algorithm for this procedure is

Use the uniform distribution to find the random variable y in the interval $[0, 1]$. $p(x)$ is a user provided PDF.

Evaluate thereafter

$$I = \int_a^b F(x) dx = \int_a^b p(x) \frac{F(x)}{p(x)} dx,$$

by rewriting

$$\int_a^b p(x) \frac{F(x)}{p(x)} dx = \int_a^b \frac{F(x(y))}{p(x(y))} dy,$$

since

$$\frac{dy}{dx} = p(x).$$

Perform then a Monte Carlo sampling for

Demonstration of Importance Sampling

$$I = \int_0^1 F(x) dx = \int_0^1 \frac{1}{1+x^2} dx = \frac{\pi}{4}.$$

We choose the following PDF (which follows closely the function to integrate)

$$p(x) = \frac{1}{3}(4-2x) \quad \int_0^1 p(x) dx = 1,$$

resulting

$$\frac{F(0)}{p(0)} = \frac{F(1)}{p(1)} = \frac{3}{4}.$$

Check that it fulfils the requirements of a PDF. We perform then the change of variables (via the Cumulative function)

$$y(x) = \int_0^x p(x') dx' = \frac{1}{3}x(4-x),$$

or

Simple Code

```
// evaluate the integral with importance sampling
for (int i = 1; i <= n; i++){
    x = ran0(hidm); // random numbers in [0,1]
    y = 2 - sqrt(4-3*x); // new random numbers
    fy = 3*func(y)/(4-2*y); // weighted function
    int_mc += fy;
    sum_sigma += fy*fy;

    int_mc = int_mc/((double) n);
    sum_sigma = sum_sigma/((double) n);
    variance=(sum_sigma-int_mc*int_mc);
}
```

Test Runs and Comparison with Brute Force for

$\pi = 3.14159$

The suffix *cr* stands for the brute force approach while *is* stands for the use of importance sampling. All calculations use `ran0` as function to generate the uniform distribution.

	N
10000 & 3.13395E+00 & 4.22881E-01 & 3.14163E+00 & 6.49921E-03	100000
100000 & 3.14195E+00	1000000
	10000000

However, it is unfair to study one-dimensional integrals with MC methods!

Multidimensional Integrals

When we deal with multidimensional integrals of the form

$$I = \int_0^1 dx_1 \int_0^1 dx_2 \dots \int_0^1 dx_d g(x_1, \dots, x_d),$$

with x_i defined in the interval $[a_i, b_i]$ we would typically need a transformation of variables of the form

$$x_i = a_i + (b_i - a_i)t_i,$$

if we were to use the uniform distribution on the interval $[0, 1]$. In this case, we need a Jacobi determinant

$$\prod_{i=1}^d (b_i - a_i),$$

and to convert the function $g(x_1, \dots, x_d)$ to

Example: 6-dimensional Integral

Consider the following six-dimensional integral

$$\int_{-\infty}^{\infty} dx dy g(x, y),$$

where

$$g(x, y) = \exp(-x^2 - y^2 - (x - y)^2/2),$$

with $d = 6$.

Example: 6-dimensional Integral

We can solve this integral by employing our brute force scheme, or using importance sampling and random variables distributed according to a gaussian PDF. For the latter, if we set the mean value $\mu = 0$ and the standard deviation $\sigma = 1/\sqrt{2}$, we have

$$\frac{1}{\sqrt{\pi}} \exp(-x^2),$$

and through

$$\pi^3 \int \prod_{i=1}^6 \left(\frac{1}{\sqrt{\pi}} \exp(-x_i^2) \right) \exp(-(x - y)^2/2) dx_1 \dots dx_6,$$

we can rewrite our integral as

$$\int f(x_1, \dots, x_d) F(x_1, \dots, x_d) \prod_{i=1}^6 dx_i,$$

where f is the gaussian distribution.

Brute Force I

```

.....
// evaluate the integral without importance sampling
// Loop over Monte Carlo Cycles
for (int i = 1; i <= n; i++){
// x[] contains the random numbers for all dimensions
for (int j = 0; j < 6; j++){
    x[j] = -length*2*length*ran0(&idum);

    fx = brute_force_MC(x);
    int_mc += fx;
    sum_sigma += fx*fx;

    int_mc = int_mc/((double) n);
    sum_sigma = sum_sigma/((double) n);
    variance = sum_sigma - int_mc*int_mc;
}
.....

```

Brute Force II

```

double brute_force_MC(double *x)
{
    double a = 1.; double b = 0.5;
    // evaluate the different terms of the exponential
    double xx = x[0]*x[0] + x[1]*x[1] + x[2]*x[2];
    double yy = x[3]*x[3] + x[4]*x[4] + x[5]*x[5];
    double xy = pow((x[0]-x[3]),2) + pow((x[1]-x[4]),2) + pow((x[2]-x[5]),2);
    return exp(-a*xx - a*yy - b*xy);
}

```

Importance Sampling I

```

.....
// evaluate the integral with importance sampling
for (int i = 1; i <= n; i++){
// x[] contains the random numbers for all dimensions
for (int j = 0; j < 6; j++) {
x[j] = gaussian_deviate(&idum)*sqrt2;

fx=gaussian_MC(x);
int_mc += fx;
sum_sigma += fx*fx;

int_mc = int_mc/((double) n );
sum_sigma = sum_sigma/((double) n );
variance=sum_sigma-int_mc*int_mc;
.....

```

Importance Sampling II

```

// this function defines the integrand to integrate
double gaussian_MC(double *x)
{
double a = 0.5;
// evaluate the different terms of the exponential
double xy=pow((x[0]-x[3]),2)+pow((x[1]-x[4]),2)+pow((x[2]-x[5]),2);
return exp(-a*xy);
} // end function for the integrand

```

Test Runs for six-dimensional Integral

Results for as function of number of Monte Carlo samples N . The exact answer is $I \approx 10.9626$ for the integral. The suffix cr stands for the brute force approach while gd stands for the use of a Gaussian distribution function. All calculations use ran0 as function to generate the uniform distribution.

	N	I_{cr}
10000	& 1.15247E+01	1.09128E+01
100000	& 1.29650E+01 & 1.09522E+01	
	1000000	1.18226E+01 & 1.09673E+01
	10000000	1.04925E+01 & 1.09612E+01

Overview of week 45

Monte Carlo methods.

- Monday: Repetition from last week
- More on importance sampling and multi-dimensional integrals
- Central limit theorem and discussion on variance, covariance and standard deviation
- Tuesday:
 - Central limit theorem and discussion on variance, covariance and standard deviation
- Random number generators (RNG).
- Random walks, Markov chains, Diffusion and Master equation

Monte Carlo Keywords

Consider it is a numerical experiment

- Be able to generate random variables following a given probability distribution function PDF
- Find a probability distribution function (PDF).
- Sampling rule for accepting a move
- Compute standard deviation and other expectation values
- Techniques for improving errors Enhances algorithmic thinking!

Example, six-dimensional integral (project 3 2012)

The task of this project is to integrate in a brute force manner a six-dimensional integral which is used to determine the ground state correlation energy between two electrons in a helium atom. We will employ both Gaussian quadrature and Monte-Carlo integration. We assume that the wave function of each electron can be modelled like the single-particle wave function of an electron in the hydrogen atom. The single-particle wave function for an electron i in the $1s$ state is given in terms of a dimensionless variable (the wave function is not properly normalized)

$$\mathbf{r}_i = x_i \mathbf{e}_x + y_i \mathbf{e}_y + z_i \mathbf{e}_z,$$

as

$$\psi_{1s}(\mathbf{r}_i) = e^{-\alpha r_i},$$

where α is a parameter and

$$\sqrt{2, \dots, 2, \dots, 2}$$

Switch to spherical coordinates

Useful to change to spherical coordinates

$$d\mathbf{r}_1 d\mathbf{r}_2 = r_1^2 dr_1 r_2^2 dr_2 d\cos(\theta_1) d\cos(\theta_2) d\phi_1 d\phi_2,$$

and

$$\frac{1}{r_{12}} = \frac{1}{\sqrt{r_1^2 + r_2^2 - 2r_1 r_2 \cos(\beta)}}$$

with

$$\cos(\beta) = \cos(\theta_1) \cos(\theta_2) + \sin(\theta_1) \sin(\theta_2) \cos(\phi_1 - \phi_2)$$

How do I do importance sampling in spherical coordinates

$r_{1,2} \in [0, \infty)$, here we use the mapping $r_{1,2} = -\ln(1 - ran)$ with $ran \in [0, 1]$, a uniform distribution point.

$\theta_{1,2} \in [0, \pi]$, use mapping $\theta_{1,2} = \pi * ran$ with $ran \in [0, 1]$ a uniform distribution point.

$\phi_{1,2} \in [0, 2\pi]$, use mapping $\phi_{1,2} = 2\pi * ran$ with $ran \in [0, 1]$ a uniform distribution point.

Be careful with the integrand

$$\frac{\exp(-4(r_1 + r_2)) r_1^2 dr_1 r_2^2 dr_2 d\cos(\theta_1) d\cos(\theta_2) d\phi_1 d\phi_2}{\sqrt{r_1^2 + r_2^2 - 2r_1 r_2 \cos(\beta)}}$$

Example, six-dimensional integral (project 3 2012)

The ansatz for the wave function for two electrons is then given by the product of two 1s wave functions as

$$\Psi(\mathbf{r}_1, \mathbf{r}_2) = e^{-\alpha(r_1 + r_2)}.$$

Note that it is not possible to find an analytic solution to Schrödinger's equation for two interacting electrons in the helium atom.

The integral we need to solve is the quantum mechanical expectation value of the correlation energy between two electrons, namely

$$\left\langle \frac{1}{|\mathbf{r}_1 - \mathbf{r}_2|} \right\rangle = \int_{-\infty}^{\infty} d\mathbf{r}_1 d\mathbf{r}_2 e^{-2\alpha(r_1 + r_2)} \frac{1}{|\mathbf{r}_1 - \mathbf{r}_2|} = \frac{5\pi^2}{16^2} = 0.192765711. \quad (77)$$

Note that our wave function is not normalized. There is a normalization factor missing, but for this project we don't need to worry about that

Example, six-dimensional integral (project 3 2012)

- a-b) Use Gaussian quadrature and compute the integral by integrating

for each variable $x_1, y_1, z_1, x_2, y_2, z_2$ from $-\infty$ to ∞ . How many mesh points do you need before the results converges at the level of the fourth leading digit? Hint: the single-particle wave function $e^{-\alpha r_i}$ is more or less zero at $r_i \approx ?$. You can therefore replace the integration limits $-\infty$ and ∞ with $-\Lambda$ and Λ , respectively. You need to check that this approximation is satisfactory.

- c) Compute the same integral but now with brute force Monte Carlo

and compare your results with those from the previous point. Discuss the differences. With brute force we mean that you should use the uniform distribution.

- d) Improve your brute force Monte Carlo calculation by using importance sampling. Hint: use the exponential distribution.

Does the answer depend on the CPU time used and/or the number of mesh points?

Example, six-dimensional integral, Gauss-Legendre

```
double *x = new double [N];
double *w = new double [N];
// set up the mesh points and weights
gauleg(a,b,x,w, N);

// evaluate the integral with the Gauss-Legendre method
// Note that we initialize the sum
double int_gauss = 0.;
for (int i=0; i<N; i++){
  for (int j = 0; j<N; j++){
    for (int k = 0; k<N; k++){
      for (int l = 0; l<N; l++){
        for (int m = 0; m<N; m++){
          for (int n = 0; n<N; n++){
            int_gauss += w[i]*w[j]*w[k]*w[l]*w[m]*w[n]
              *int_function(x[i],x[j],x[k],x[l],x[m],x[n]);
          }}}}}
}
```

Example, six-dimensional integral, Gauss-Legendre

```
// this function defines the function to integrate
double int_function(double x1, double y1, double z1,
  double x2, double y2, double z2)
{
  double alpha = 2.;
  // evaluate the different terms of the exponential
  double exp1 = -2*alpha*sqrt(x1*x1+y1*y1+z1*z1);
  double exp2 = -2*alpha*sqrt(x2*x2+y2*y2+z2*z2);
  double deno = sqrt(pow((x1-x2),2)+pow((y1-y2),2)+pow((z1-z2),2));
  if(deno < pow(10,-6)) { return 0; }
  else return exp(exp1+exp2)/deno;
} // end of function to evaluate
```

Example, six-dimensional integral, brute force MC

```
double int_mc = 0.; double variance = 0.;
double sum_sigma = 0.; long idum = -1;
double length = 1.5; // we fix the max size of the box to L=3
double volume = pow((2*length), 6.);

// evaluate the integral with importance sampling
for (int i = 1; i <= n; i++){
    // x[] contains the random numbers for all dimensions
    for (int j = 0; j < 6; j++){
        // Maps U[0,1] to U[-L,L]
        x[j] = -length + 2*length*ran0(&idum);

        fx = brute_force_MC(x);
        int_mc += fx;
        sum_sigma += fx*fx;
    }

    int_mc = int_mc/((double) n);
    sum_sigma = sum_sigma/((double) n);
    variance = sum_sigma - int_mc*int_mc;
    ....
}
```

Example, six-dimensional integral, brute force MC

```
double brute_force_MC(double *x)
{
    double alpha = 2.;
    // evaluate the different terms of the exponential
    double exp1 = -2*alpha*sqrt(x[0]*x[0]+x[1]*x[1]+x[2]*x[2]);
    double exp2 = -2*alpha*sqrt(x[3]*x[3]+x[4]*x[4]+x[5]*x[5]);
    double deno = sqrt(pow((x[0]-x[3]), 2)
        + pow((x[1]-x[4]), 2) + pow((x[2]-x[5]), 2));
    double value = exp(exp1+exp2)/deno;
    return value;
} // end function for the integrand
```

Example, six-dimensional integral, importance sampling

```
double int_mc = 0.; double variance = 0.;
double sum_sigma = 0.; long idum = -1;
// The 'volume' contains 4 jacobideterminants(pi, pi, 2pi, 2pi)
// and a scaling factor 1/16
double volume = 4*pow(acos(-1.), 4)*1./16;
// evaluate the integral with importance sampling
for (int i = 1; i <= n; i++){
    for (int j = 0; j < 2; j++){
        y = ran0(&idum);
        x[j] = -0.25*log(1.-y);
    }
    for (int j = 2; j < 4; j++){
        x[j] = 2*acos(-1.)*ran0(&idum);
    }
    for (int j = 4; j < 6; j++){
        x[j] = acos(-1.)*ran0(&idum);
    }
    fx = integrand_MC(x);
    ....
}
```

Example, six-dimensional integral, importance sampling

```
// this function defines the integrand to integrate
double integrand_MC(double *x)
{
    double num = x[0]*x[0]*x[1]*x[1]*sin(x[4])*sin(x[5]);
    double deno = sqrt(x[0]*x[0]+x[1]*x[1]-2*x[0]*x[1]*
        (sin(x[4])*sin(x[5])*cos(x[2]-x[3])+cos(x[4])*cos(x[5]))));
    return num/deno;
} // end function for the integrand
```

Test Runs and Comparison with Brute Force and Gauss-Legendre

The suffix *br* stands for the brute force approach while *is* stands for the use of importance sampling.

	N	I_{br}	σ_{br}	time
1E6 & 0.19238	3.85124E-4	0.6	0.19176 & 1.01515E-4	1.4
10E6	0.18607		1.18053E-4	6
100E6	0.18846		4.37163E-4	57
1000E6	0.18843		1.35879E-4	58

Gauss-Legendre results:

	N	time(s)	I_n	$ I_n - I $
20 & 31	0.18047	& 1.123E-2		
30 & 354	0.18501	& 7.76E-3		
40 & 1999	0.18653	& 6.24E-3		
50 & 7578	0.18722	& 5.54E-3		

Variance, covariance, errors etc, chapter 11.2

Statistical analysis.

- Monte Carlo simulations can be treated as *computer experiments*
- The results can be analysed with the same statistics tools we would use in analysing laboratory experiments
- As in all other experiments, we are looking for expectation values and an estimate of how accurate they are, i.e., the error

Variance, covariance, errors etc

Statistical analysis.

- As in other experiments, Monte Carlo experiments have two classes of errors:
- Statistical errors
- Systematic errors
- Statistical errors can be estimated using standard tools from statistics
- Systematic errors are method specific and must be treated differently from case to case.

Variance, covariance, errors etc

A *stochastic process* is a process that produces sequentially a chain of values:

$$\{x_1, x_2, \dots, x_k, \dots\}.$$

We will call these values our *measurements* and the entire set as our measured *sample*. The action of measuring all the elements of a sample we will call a stochastic *experiment* (since, operationally, they are often associated with results of empirical observation of some physical or mathematical phenomena; precisely an experiment). We assume that these values are distributed according to some PDF $p_X(x)$, where X is just the formal symbol for the stochastic variable whose PDF is $p_X(x)$. Instead of trying to determine the full distribution p we are often only interested in finding the few lowest moments, like the mean μ_X and the variance σ_X^2 .

Variance, covariance, errors etc

The *probability distribution function (PDF)* is a function $p(x)$ on the domain which, in the discrete case, gives us the probability or relative frequency with which these values of X occur:

$$p(x) = \text{Prob}(X = x)$$

In the continuous case, the PDF does not directly depict the actual probability. Instead we define the probability for the stochastic variable to assume any value on an infinitesimal interval around x to be $p(x)dx$. The continuous function $p(x)$ then gives us the *density* of the probability rather than the probability itself. The probability for a stochastic variable to assume any value on a non-infinitesimal interval $[a, b]$ is then just the integral:

$$\text{Prob}(a \leq X \leq b) = \int_a^b p(x)dx$$

Qualitatively speaking, a stochastic variable represents the values of numbers chosen as if by chance from some specified PDF so that

Variance, covariance, errors etc

Also of interest to us is the *cumulative probability distribution function (CDF)*, $P(x)$, which is just the probability for a stochastic variable X to assume any value less than x :

$$P(x) = \text{Prob}(X \leq x) = \int_{-\infty}^x p(x')dx'$$

The relation between a CDF and its corresponding PDF is then:

$$p(x) = \frac{d}{dx}P(x)$$

Variance, covariance, errors etc

A particularly useful class of special expectation values are the *moments*. The n -th moment of the PDF p is defined as follows:

$$\langle x^n \rangle \equiv \int x^n p(x) dx$$

The zero-th moment $\langle 1 \rangle$ is just the normalization condition of p . The first moment, $\langle x \rangle$, is called the *mean* of p and often denoted by the letter μ :

$$\langle x \rangle = \mu \equiv \int x p(x) dx$$

Variance, covariance, errors etc

A special version of the moments is the set of *central moments*, the n -th central moment defined as:

$$\langle (x - \langle x \rangle)^n \rangle \equiv \int (x - \langle x \rangle)^n p(x) dx$$

The zero-th and first central moments are both trivial, equal 1 and 0, respectively. But the second central moment, known as the *variance* of p , is of particular interest. For the stochastic variable X , the variance is denoted as σ_X^2 or $\text{Var}(X)$:

$$\begin{aligned} \sigma_X^2 &= \text{Var}(X) = \langle (x - \langle x \rangle)^2 \rangle = \int (x - \langle x \rangle)^2 p(x) dx \\ &= \int (x^2 - 2x\langle x \rangle + \langle x \rangle^2) p(x) dx \\ &= \langle x^2 \rangle - 2\langle x \rangle \langle x \rangle + \langle x \rangle^2 \\ &= \langle x^2 \rangle - \langle x \rangle^2 \end{aligned}$$

Variance, covariance, errors etc

Another important quantity is the so called covariance, a variant of the above defined variance. Consider again the set $\{X_i\}$ of n stochastic variables (not necessarily uncorrelated) with the multivariate PDF $P(x_1, \dots, x_n)$. The *covariance* of two of the stochastic variables, X_i and X_j , is defined as follows:

$$\begin{aligned}\text{Cov}(X_i, X_j) &\equiv \langle (x_i - \langle x_i \rangle)(x_j - \langle x_j \rangle) \rangle \\ &= \int \dots \int (x_i - \langle x_i \rangle)(x_j - \langle x_j \rangle) P(x_1, \dots, x_n) dx_1 \dots dx_n\end{aligned}\quad (78)$$

with

$$\langle x_i \rangle = \int \dots \int x_i P(x_1, \dots, x_n) dx_1 \dots dx_n$$

Variance, covariance, errors etc

If we consider the above covariance as a matrix $C_{ij} = \text{Cov}(X_i, X_j)$, then the diagonal elements are just the familiar variances, $C_{ii} = \text{Cov}(X_i, X_i) = \text{Var}(X_i)$. It turns out that all the off-diagonal elements are zero if the stochastic variables are uncorrelated. This is easy to show, keeping in mind the linearity of the expectation value. Consider the stochastic variables X_i and X_j , ($i \neq j$):

$$\begin{aligned}\text{Cov}(X_i, X_j) &= \langle (x_i - \langle x_i \rangle)(x_j - \langle x_j \rangle) \rangle \\ &= \langle x_i x_j - x_i \langle x_j \rangle - \langle x_i \rangle x_j + \langle x_i \rangle \langle x_j \rangle \rangle \\ &= \langle x_i x_j \rangle - \langle x_i \rangle \langle x_j \rangle - \langle x_i \rangle \langle x_j \rangle + \langle x_i \rangle \langle x_j \rangle \\ &= \langle x_i x_j \rangle - \langle x_i \rangle \langle x_j \rangle - \langle x_i \rangle \langle x_j \rangle + \langle x_i \rangle \langle x_j \rangle \\ &= \langle x_i x_j \rangle - \langle x_i \rangle \langle x_j \rangle\end{aligned}$$

Variance, covariance, errors etc

If X_i and X_j are independent, we get $\langle x_i x_j \rangle = \langle x_i \rangle \langle x_j \rangle$, resulting in $\text{Cov}(X_i, X_j) = 0$ ($i \neq j$).

Also useful for us is the covariance of linear combinations of stochastic variables. Let $\{X_i\}$ and $\{Y_j\}$ be two sets of stochastic variables. Let also $\{a_i\}$ and $\{b_j\}$ be two sets of scalars. Consider the linear combination:

$$U = \sum_i a_i X_i \quad V = \sum_j b_j Y_j$$

By the linearity of the expectation value

$$\text{Cov}(U, V) = \sum_{i,j} a_i b_j \text{Cov}(X_i, Y_j)$$

Variance, covariance, errors etc

Now, since the variance is just $\text{Var}(X_i) = \text{Cov}(X_i, X_i)$, we get the variance of the linear combination $U = \sum_i a_i X_i$:

$$\text{Var}(U) = \sum_{i,j} a_i a_j \text{Cov}(X_i, X_j)$$

And in the special case when the stochastic variables are uncorrelated, the off-diagonal elements of the covariance are as we know zero, resulting in:

$$\text{Var}(U) = \sum_i a_i^2 \text{Cov}(X_i, X_i) = \sum_i a_i^2 \text{Var}(X_i)$$

$$\text{Var}\left(\sum_i a_i X_i\right) = \sum_i a_i^2 \text{Var}(X_i)$$

which will become very useful in our study of the error in the mean value of a set of measurements.

Variance, covariance, errors etc

In practical situations a sample is always of finite size. Let that size be n . The expectation value of a sample, the *sample mean*, is then defined as follows:

$$\bar{x}_n \equiv \frac{1}{n} \sum_{k=1}^n x_k$$

The *sample variance* is:

$$\text{Var}(x) \equiv \frac{1}{n} \sum_{k=1}^n (x_k - \bar{x}_n)^2$$

its square root being the *standard deviation of the sample*. The *sample covariance* is:

$$\text{Cov}(x) \equiv \frac{1}{n} \sum_{kl} (x_k - \bar{x}_n)(x_l - \bar{x}_n)$$

Variance, covariance, errors etc

Note that the sample variance is the sample covariance without the cross terms. In a similar manner as the covariance in eq. (78) is a measure of the correlation between two stochastic variables, the above defined sample covariance is a measure of the sequential correlation between succeeding measurements of a sample.

These quantities, being known experimental values, differ significantly from and must not be confused with the similarly named quantities for stochastic variables, mean μ_X , variance $\text{Var}(X)$ and covariance $\text{Cov}(X, Y)$.

Variance, covariance, errors etc

The law of large numbers states that as the size of our sample grows to infinity, the sample mean approaches the true mean μ_X of the chosen PDF:

$$\lim_{n \rightarrow \infty} \bar{x}_n = \mu_X$$

The sample mean \bar{x}_n works therefore as an estimate of the true mean μ_X .

What we need to find out is how good an approximation \bar{x}_n is to μ_X . In any stochastic measurement, an estimated mean is of no use to us without a measure of its error. A quantity that tells us how well we can reproduce it in another experiment. We are therefore interested in the PDF of the sample mean itself. Its standard deviation will be a measure of the spread of sample means, and we will simply call it the *error* of the sample mean, or just sample error, and denote it by err_X . In practice, we will only be able to produce an *estimate* of the sample error since the exact value would require the knowledge of the true PDFs behind, which

Variance, covariance, errors etc

The straight forward brute force way of estimating the sample error is simply by producing a number of samples, and treating the mean of each as a measurement. The standard deviation of these means will then be an estimate of the original sample error. If we are unable to produce more than one sample, we can split it up sequentially into smaller ones, treating each in the same way as above. This procedure is known as *blocking* and will be given more attention shortly. At this point it is worth while exploring more indirect methods of estimation that will help us understand some important underlying principles of correlational effects.

Variance, covariance, errors etc

Let us first take a look at what happens to the sample error as the size of the sample grows. In a sample, each of the measurements x_i can be associated with its own stochastic variable X_i . The stochastic variable \bar{X}_n for the sample mean \bar{x}_n is then just a linear combination, already familiar to us:

$$\bar{X}_n = \frac{1}{n} \sum_{i=1}^n X_i$$

All the coefficients are just equal $1/n$. The PDF of \bar{X}_n , denoted by $p_{\bar{X}_n}(x)$ is the desired PDF of the sample means.

Variance, covariance, errors etc

The probability density of obtaining a sample mean \bar{x}_n is the product of probabilities of obtaining arbitrary values x_1, x_2, \dots, x_n with the constraint that the mean of the set $\{x_i\}$ is \bar{x}_n :

$$p_{\bar{X}_n}(x) = \int p_X(x_1) \cdots \int p_X(x_n) \delta\left(x - \frac{x_1 + x_2 + \cdots + x_n}{n}\right) dx_n \cdots dx_1$$

And in particular we are interested in its variance $\text{Var}(\bar{X}_n)$.

Variance, covariance, errors etc

It is generally not possible to express $p_{\bar{X}_n}(x)$ in a closed form given an arbitrary PDF p_X and a number n . But for the limit $n \rightarrow \infty$ it is possible to make an approximation. The very important result is called *the central limit theorem*. It tells us that as n goes to infinity, $p_{\bar{X}_n}(x)$ approaches a Gaussian distribution whose mean and variance equal the true mean and variance, μ_X and σ_X^2 , respectively:

$$\lim_{n \rightarrow \infty} p_{\bar{X}_n}(x) = \left(\frac{n}{2\pi \text{Var}(X)}\right)^{1/2} e^{-\frac{n(x-\mu_X)^2}{2\text{Var}(X)}}$$

Variance, covariance, errors etc

The desired variance $\text{Var}(\bar{X}_n)$, i.e. the sample error squared err_X^2 , is given by:

$$\text{err}_X^2 = \text{Var}(\bar{X}_n) = \frac{1}{n^2} \sum_{ij} \text{Cov}(X_i, X_j) \quad (79)$$

We see now that in order to calculate the exact error of the sample with the above expression, we would need the true means μ_{X_i} of the stochastic variables X_i . To calculate these requires that we know the true multivariate PDF of all the X_i . But this PDF is unknown to us, we have only got the measurements of one sample. The best we can do is to let the sample itself be an estimate of the PDF of each of the X_i , estimating all properties of X_i through the measurements of the sample.

Variance, covariance, errors etc

Our estimate of μ_{X_i} is then the sample mean \bar{x} itself, in accordance with the central limit theorem:

$$\mu_{X_i} = \langle x_i \rangle \approx \frac{1}{n} \sum_{k=1}^n x_k = \bar{x}$$

Using \bar{x} in place of μ_{X_i} we can give an *estimate* of the covariance in eq. (79):

$$\begin{aligned} \text{Cov}(X_i, X_j) &= \langle (x_i - \langle x_i \rangle)(x_j - \langle x_j \rangle) \rangle \approx \langle (x_i - \bar{x})(x_j - \bar{x}) \rangle \\ &\approx \frac{1}{n} \sum_i \left(\frac{1}{n} \sum_k \right. \\ &= \frac{1}{n} \text{Cov}(x) \end{aligned}$$

Variance, covariance, errors etc

By the same procedure we can use the sample variance as an estimate of the variance of any of the stochastic variables X_i :

$$\begin{aligned} \text{Var}(X_i) &= \langle x_i - \langle x_i \rangle \rangle \approx \langle x_i - \bar{x}_n \rangle \\ &\approx \frac{1}{n} \sum_{k=1}^n (x_k - \bar{x}_n) \\ &= \text{Var}(x) \end{aligned} \quad (80)$$

Now we can calculate an estimate of the error err_X of the sample mean \bar{x}_n :

$$\begin{aligned} \text{err}_X^2 &= \frac{1}{n^2} \sum_{ij} \text{Cov}(X_i, X_j) \\ &\approx \frac{1}{n^2} \sum_{ij} \frac{1}{n} \text{Cov}(x) = \frac{1}{n^2} n^2 \frac{1}{n} \text{Cov}(x) \end{aligned}$$

Variance, covariance, errors etc

In the special case that the measurements of the sample are uncorrelated (equivalently the stochastic variables X_i are uncorrelated) we have that the off-diagonal elements of the covariance are zero. This gives the following estimate of the sample error:

$$\begin{aligned} \text{err}_X^2 &= \frac{1}{n^2} \sum_{ij} \text{Cov}(X_i, X_j) = \frac{1}{n^2} \sum_i \text{Var}(X_i) \\ &\approx \frac{1}{n^2} \sum_i \text{Var}(x) \\ &= \frac{1}{n} \text{Var}(x) \end{aligned} \quad (82)$$

where in the second step we have used eq. (80). The error of the sample is then just its standard deviation divided by the square root of the number of measurements the sample contains. This is a very useful formula which is easy to compute. It acts as a first approximation to the error, but in numerical experiments, we

Variance, covariance, errors etc

For computational purposes one usually splits up the estimate of err_X^2 , given by eq. (81), into two parts:

$$\begin{aligned} \text{err}_X^2 &= \frac{1}{n} \text{Var}(x) + \frac{1}{n} (\text{Cov}(x) - \text{Var}(x)) \\ &= \frac{1}{n^2} \sum_{k=1}^n (x_k - \bar{x}_n)^2 + \frac{2}{n^2} \sum_{k < l} (x_k - \bar{x}_n)(x_l - \bar{x}_n) \end{aligned} \quad (83)$$

The first term is the same as the error in the uncorrelated case, Eq. (82). This means that the second term accounts for the error correction due to correlation between the measurements. For uncorrelated measurements this second term is zero.

Variance, covariance, errors etc

Computationally the uncorrelated first term is much easier to treat efficiently than the second.

$$\text{Var}(x) = \frac{1}{n} \sum_{k=1}^n (x_k - \bar{x}_n)^2 = \left(\frac{1}{n} \sum_{k=1}^n x_k^2 \right) - \bar{x}_n^2$$

We just accumulate separately the values x^2 and x for every measurement x we receive. The correlation term, though, has to be calculated at the end of the experiment since we need all the measurements to calculate the cross terms. Therefore, all measurements have to be stored throughout the experiment.

Random Numbers

Most used are so-called 'Linear congruential'

$$N_i = (aN_{i-1} + c) \text{MOD}(M),$$

and to find a number in $x \in [0, 1]$

$$x_i = N_i / M$$

M is called the period and should be as big as possible. The start value is N_0 and is called the seed.

- The random variables should result in the uniform distribution
- No correlations between numbers (zero covariance)
- As big as possible period M
- Fast algo

Random Numbers

The problem with such generators is that their outputs are periodic; they will start to repeat themselves with a period that is at most M . If however the parameters a and c are badly chosen, the period may be even shorter.

Consider the following example

$$N_i = (6N_{i-1} + 7) \text{MOD}(5),$$

with a seed $N_0 = 2$. This generator produces the sequence 4, 1, 3, 0, 2, 4, 1, 3, 0, 2, ... i.e., a sequence with period 5. However, increasing M may not guarantee a larger period as the following example shows

$$N_i = (27N_{i-1} + 11) \text{MOD}(54),$$

which still, with $N_0 = 2$, results in 11, 38, 11, 38, 11, 38, ... a period of just 2.

Random Numbers

Typical periods for the random generators provided in the program library are of the order of $\sim 10^9$ or larger. Other random number generators which have become increasingly popular are so-called shift-register generators. In these generators each successive number depends on many preceding values (rather than the last values as in the linear congruential generator). For example, you could make a shift register generator whose l th number is the sum of the $l-i$ th and $l-j$ th values with modulo M ,

$$N_l = (aN_{l-i} + cN_{l-j}) \text{MOD}(M).$$

Such a generator again produces a sequence of pseudorandom numbers but this time with a period much larger than M . It is also possible to construct more elaborate algorithms by including more than two past terms in the sum of each iteration. One example is the generator of Marsaglia and Zaman (Computers in Physics **8** (1994) 117) which consists of two congruential relations

$$N_l = (N_{l-3} - N_{l-1}) \text{MOD}(2^{31} - 69),$$

Random Numbers

Using modular addition, we could use the bitwise exclusive-OR (\oplus) operation so that

$$N_l = (N_{l-i} \oplus N_{l-j})$$

where the bitwise action of \oplus means that if $N_{l-i} = N_{l-j}$ the result is 0 whereas if $N_{l-i} \neq N_{l-j}$ the result is 1. As an example, consider the case where $N_{l-i} = 6$ and $N_{l-j} = 11$. The first one has a bit representation (using 4 bits only) which reads 0110 whereas the second number is 1011. Employing the \oplus operator yields 1101, or $2^3 + 2^2 + 2^0 = 13$.

In Fortran90, the bitwise \oplus operation is coded through the intrinsic function IEXOR(m, n) where m and n are the input numbers, while in C it is given by $m \wedge n$.

Random Numbers

The function *ran0* implements

$$N_l = (aN_{l-1}) \text{MOD}(M).$$

Note that $c = 0$ and that it cannot be initialized with $N_0 = 0$. However, since a and N_{l-1} are integers and their multiplication could become greater than the standard 32 bit integer, there is a trick via Schrage's algorithm which approximates the multiplication of large integers through the factorization

$$M = aq + r,$$

where we have defined

$$q = \lfloor M/a \rfloor,$$

and

$$r = M \text{ MOD } a.$$

Random Numbers

To see how this works we note first that

$$(aN_{l-1}) \text{MOD}(M) = (aN_{l-1} - \lfloor N_{l-1}/q \rfloor M) \text{MOD}(M),$$

since we can add or subtract any integer multiple of M from aN_{l-1} . The last term $\lfloor N_{l-1}/q \rfloor M \text{MOD}(M)$ is zero since the integer division $\lfloor N_{l-1}/q \rfloor$ just yields a constant which is multiplied with M . Rewrite as

$$(aN_{l-1}) \text{MOD}(M) = (aN_{l-1} - \lfloor N_{l-1}/q \rfloor (aq + r)) \text{MOD}(M),$$

Random Numbers

It gives

$$(aN_{l-1}) \text{MOD}(M) = (aN_{l-1} - \lfloor N_{l-1}/q \rfloor q - \lfloor N_{l-1}/q \rfloor r) \text{MOD}(M),$$

yielding

$$(aN_{l-1}) \text{MOD}(M) = (aN_{l-1} \text{MOD}(q) - \lfloor N_{l-1}/q \rfloor r) \text{MOD}(M).$$

- $\lfloor N_{l-1}/q \rfloor r$ is always smaller or equal $N_{l-1}(r/q)$ and with $r < q$ we obtain always a number smaller than N_{l-1} , which is smaller than M .
- $N_{l-1} \text{MOD}(q)$ is between zero and $q - 1$ then $a(N_{l-1} \text{MOD}(q)) < aq$.
- Our definition of $q = \lfloor M/a \rfloor$ ensures that this term is also smaller than M meaning that both terms fit into a 32-bit signed integer. None of these two terms can be negative, but their difference could.

Random Numbers

```

/* ran0() is an "Minimal" random number generator of Park and Miller
** Set or reset the input value
** idum to any integer value (except the unlikely value MASK)
** to initialize the sequence; idum must not be altered between
** calls for successive deviates in a sequence.
** The function returns a uniform deviate between 0.0 and 1.0.
*/
double ran0(long &idum)
{
    const int a = 16807, m = 2147483647, q = 127773;
    const int r = 2836, MASK = 123459876;
    const double am = 1./m;
    long k;
    double ans;
    idum ^= MASK;
    k = (*idum)/q;
    idum = a*(idum - k*q) - r*k;
    // add m if negative difference
    if(idum < 0) idum += m;
    ans = am*(idum);
    idum ^= MASK;
    return ans;
} // End: function ran0()

```

Random Numbers

Important tests of random numbers are the standard deviation σ and the mean $\mu = \langle x \rangle$.

For the uniform distribution we have

$$\langle x^k \rangle = \int_0^1 dx p(x) x^k = \int_0^1 dx x^k = \frac{1}{k+1},$$

since $p(x) = 1$. The mean value μ is then

$$\mu = \langle x \rangle = \frac{1}{2}$$

while the standard deviation is

$$\sigma = \sqrt{\langle x^2 \rangle - \mu^2} = \frac{1}{\sqrt{12}} = 0.2886.$$

Random Numbers

Number of x -values for various intervals generated by 4 random number generators, their corresponding mean values and standard deviations. All calculations have been initialized with the variable $idum = -1$.

x -bin	ran0	ran1	ran2	ran3
0.0-0.1	1013	991	938	1047
0.1-0.2	1002	1009	1040	1030
0.2-0.3	989	999	1030	993
0.3-0.4	939	960	1023	937
0.4-0.5	1038	1001	1002	992
0.5-0.6	1037	1047	1009 & 1009	
0.6-0.7	1005	989	1003	989
0.7-0.8	986	962	985	954
0.8-0.9	1000	1027	1009	1023
0.9-1.0	991	1015	961	1026
μ	0.4997	0.5018	0.4992 & 0.4990	
σ	0.2882	0.2892	0.2861	0.2915

Overview of week 46

Monte Carlo methods chapter 12.

- Monday: Repetition from last week
- Brownian motion and Markov chains, chapters 12.2 and 12.3 of lecture notes
- Tuesday:
 - More on Markov chains (chapter 12.4 and 12.5)
 - The Metropolis algorithm (chapter 12.5) and simple implementations of it
- Discussion of project 5, there are three versions, available from Tuesday the 12th. Next week we will continue to discuss the various projects. For the quantum mechanical version we will also discuss how to perform quantum mechanical calculations. We will also discuss how to parallelize the diffusion equation.

Why Markov Chains?

- We want to study a physical system which evolves towards equilibrium, from given initial conditions.
- We start with a PDF $w(x_0, t_0)$ and we want to understand how it evolves with time.
- We want to reach a situation where after a given number of time steps we obtain a steady state. This means that the system reaches its most likely state (equilibrium situation)
- Our PDF is normally a multidimensional object whose normalization constant is impossible to find.
- Analytical calculations from $w(x, t)$ are not possible.
- To sample directly from $w(x, t)$ is not possible/difficult.
- The transition probability W is also not known.
- How can we establish that we have reached a steady state? Sounds impossible! Use Markov chain Monte Carlo

Brownian Motion and Markov Processes

A Markov process is a random walk with a selected probability for making a move. The new move is independent of the previous history of the system. The Markov process is used repeatedly in Monte Carlo simulations in order to generate new random states. The reason for choosing a Markov process is that when it is run for a long enough time starting with a random state, we will eventually reach the most likely state of the system. In thermodynamics, this means that after a certain number of Markov processes we reach an equilibrium distribution.

This mimics the way a real system reaches its most likely state at a given temperature of the surroundings.

Brownian Motion and Markov Processes

To reach this distribution, the Markov process needs to obey two important conditions, that of **ergodicity** and **detailed balance**. These conditions impose then constraints on our algorithms for accepting or rejecting new random states. The Metropolis algorithm discussed here abides to both these constraints. The Metropolis algorithm is widely used in Monte Carlo simulations and the understanding of it rests within the interpretation of random walks and Markov processes.

Brownian Motion and Markov Processes

In a random walk one defines a mathematical entity called a **walker**, whose attributes completely define the state of the system in question. The state of the system can refer to any physical quantities, from the vibrational state of a molecule specified by a set of quantum numbers, to the brands of coffee in your favourite supermarket. The walker moves in an appropriate state space by a combination of deterministic and random displacements from its previous position.

This sequence of steps forms a **chain**.

Sequence of ingredients

- We want to study a physical system which evolves towards equilibrium, from given initial conditions.
- Markov chains are intimately linked with the physical process of diffusion. Proof in lecture notes.
- From a Markov chain we can then derive the conditions for detailed balance and ergodicity. These are the conditions needed for obtaining a steady state.
- The widely used algorithm for doing this is the so-called Metropolis algorithm, in its refined form the Metropolis-Hastings algorithm.

Applications: almost every field

Financial engineering, see for example Patriarca *et al*, Physica **340**, page 334 (2004).

Neuroscience, see for example Lipinski, Physics Medical Biology **35**, page 441 (1990) or Farnell and Gibson, Journal of Computational Physics **208**, page 253 (2005)

Tons of applications in physics

and chemistry

and biology, medicine

Nobel prize in economy to Black and Scholes

$$\frac{\partial V}{\partial t} + \frac{1}{2}\sigma^2 S^2 \frac{\partial^2 V}{\partial S^2} + rS \frac{\partial V}{\partial S} - rV = 0.$$

The Black and Scholes equation is a partial differential equation, which describes the price of the option over time. The derivation is based on Brownian motion (Langevin and Fokker-Planck, 12.6)

A simple Example

The obvious case is that of a random walker on a one-, or two- or three-dimensional lattice (dubbed coordinate space hereafter)

Consider a system whose energy is defined by the orientation of single spins. Consider the state i , with given energy E_i represented by the following N spins

$\uparrow \uparrow \uparrow \dots \uparrow \downarrow \uparrow \dots \uparrow \downarrow$
 $1 \ 2 \ 3 \ \dots \ k-1 \ k \ k+1 \ \dots \ N-1 \ N$

We may be interested in the transition with one single spinflip to a new state j with energy E_j

$\uparrow \uparrow \uparrow \dots \uparrow \uparrow \uparrow \dots \uparrow \downarrow$
 $1 \ 2 \ 3 \ \dots \ k-1 \ k \ k+1 \ \dots \ N-1 \ N$

This change from one microstate i (or spin configuration) to another microstate j is the **configuration space** analogue to a random walk on a lattice. Instead of jumping from one place to another in space, we 'jump' from one microstate to another.

Diffusion from Markov Chain

From experiment there are strong indications that the flux of particles $j(x, t)$, viz., the number of particles passing x at a time t is proportional to the gradient of $w(x, t)$. This proportionality is expressed mathematically through

$$j(x, t) = -D \frac{\partial w(x, t)}{\partial x}, \quad (84)$$

where D is the so-called diffusion constant, with dimensionality length² per time. If the number of particles is conserved, we have the continuity equation

$$\frac{\partial j(x, t)}{\partial x} = -\frac{\partial w(x, t)}{\partial t}, \quad (85)$$

which leads to

$$\frac{\partial w(x, t)}{\partial t} = D \frac{\partial^2 w(x, t)}{\partial x^2}, \quad (86)$$

which is the diffusion equation in one dimension. Solved as a

Diffusion from Markov Chain

With the probability distribution function $w(x, t)dx$ we can compute expectation values such as the mean distance

$$\langle x(t) \rangle = \int_{-\infty}^{\infty} xw(x, t)dx, \quad (87)$$

or

$$\langle x^2(t) \rangle = \int_{-\infty}^{\infty} x^2w(x, t)dx, \quad (88)$$

which allows for the computation of the variance $\sigma^2 = \langle x^2(t) \rangle - \langle x(t) \rangle^2$. Note well that these expectation values are time-dependent. In a similar way we can also define expectation values of functions $f(x, t)$ as

$$\langle f(x, t) \rangle = \int_{-\infty}^{\infty} f(x, t)w(x, t)dx. \quad (89)$$

Diffusion from Markov Chain

Since $w(x, t)$ is now treated as a PDF, it needs to obey the same criteria as discussed in the previous chapter. However, the normalization condition

$$\int_{-\infty}^{\infty} w(x, t)dx = 1 \quad (90)$$

imposes significant constraints on $w(x, t)$. These are

$$w(x = \pm\infty, t) = 0 \quad \frac{\partial^n w(x, t)}{\partial x^n} \Big|_{x=\pm\infty} = 0, \quad (91)$$

implying that when we study the time-derivative $\partial \langle x(t) \rangle / \partial t$, we obtain after integration by parts and using Eq. (86)

$$\frac{\partial \langle x \rangle}{\partial t} = \int_{-\infty}^{\infty} x \frac{\partial w(x, t)}{\partial t} dx = D \int_{-\infty}^{\infty} x \frac{\partial^2 w(x, t)}{\partial x^2} dx, \quad (92)$$

leading to

Diffusion from Markov Chain

This means in turn that $\langle x \rangle$ is independent of time. If we choose the initial position $x(t=0) = 0$, the average displacement $\langle x \rangle = 0$. If we link this discussion to a random walk in one dimension with equal probability of jumping to the left or right and with an initial position $x = 0$, then our probability distribution remains centered around $\langle x \rangle = 0$ as function of time. However, the variance is not necessarily 0. Consider first

$$\frac{\partial \langle x^2 \rangle}{\partial t} = Dx^2 \frac{\partial w(x, t)}{\partial x} \Big|_{x=\pm\infty} - 2D \int_{-\infty}^{\infty} x \frac{\partial w(x, t)}{\partial x} dx, \quad (95)$$

where we have performed an integration by parts as we did for $\frac{\partial \langle x \rangle}{\partial t}$. A further integration by parts results in

$$\frac{\partial \langle x^2 \rangle}{\partial t} = -Dxw(x, t) \Big|_{x=\pm\infty} + 2D \int_{-\infty}^{\infty} w(x, t)dx = 2D, \quad (96)$$

leading to

Random walks, chapter 12.2

Consider now a random walker in one dimension, with probability R of moving to the right and L for moving to the left. At $t = 0$ we place the walker at $x = 0$. The walker can then jump, with the above probabilities, either to the left or to the right for each time step. Note that in principle we could also have the possibility that the walker remains in the same position. This is not implemented in this example. Every step has length $\Delta x = l$. Time is discretized and we have a jump either to the left or to the right at every time step.

Random walks

Let us now assume that we have equal probabilities for jumping to the left or to the right, i.e., $L = R = 1/2$. The average displacement after n time steps is

$$\langle x(n) \rangle = \sum_i^n \Delta x_i = 0 \quad \Delta x_i = \pm l,$$

since we have an equal probability of jumping either to the left or to right. The value of $\langle x(n)^2 \rangle$ is

$$\langle x(n)^2 \rangle = \left(\sum_i^n \Delta x_i \right)^2 = \sum_i^n \Delta x_i^2 + \sum_{i \neq j}^n \Delta x_i \Delta x_j = l^2 n.$$

For many enough steps the non-diagonal contribution is

$$\sum_{i \neq j}^N \Delta x_i \Delta x_j = 0,$$

Random walks

The variance is then

$$\langle x(n)^2 \rangle - \langle x(n) \rangle^2 = l^2 n.$$

It is also rather straightforward to compute the variance for $L \neq R$. The result is

$$\langle x(n)^2 \rangle - \langle x(n) \rangle^2 = 4LRl^2 n.$$

The variable n represents the number of time steps. If we define $n = t/\Delta t$, we can then couple the variance result from a random walk in one dimension with the variance from diffusion by defining the diffusion constant as

$$D = \frac{l^2}{\Delta t}.$$

Diffusion from Markov Chain

When solving partial differential equations such as the diffusion equation numerically, the derivatives are always discretized. We can rewrite the time derivative as

$$\frac{\partial w(x, t)}{\partial t} \approx \frac{w(i, n+1) - w(i, n)}{\Delta t}, \quad (100)$$

whereas the gradient is approximated as

$$D \frac{\partial^2 w(x, t)}{\partial x^2} \approx D \frac{w(i+1, n) + w(i-1, n) - 2w(i, n)}{(\Delta x)^2}, \quad (101)$$

resulting in the discretized diffusion equation

$$\frac{w(i, n+1) - w(i, n)}{\Delta t} = D \frac{w(i+1, n) + w(i-1, n) - 2w(i, n)}{(\Delta x)^2},$$

where n represents a given time step and i a step in the x -direction.

Diffusion from Markov Chain

A Markov process allows in principle a microscopic description of Brownian motion. As with the random walk, we consider a particle which moves along the x -axis in the form of a series of jumps with step length $\Delta x = l$. Time and space are discretized and the subsequent moves are statistically independent, i.e., the new move depends only on the previous step and not on the results from earlier trials. We start at a position $x = jl = j\Delta x$ and move to a new position $x = i\Delta x$ during a step $\Delta t = \epsilon$, where $i \geq 0$ and $j \geq 0$ are integers. The original probability distribution function (PDF) of the particles is given by $w_i(t=0)$ where i refers to a specific position on a grid, with $i=0$ representing $x=0$. The function $w_i(t=0)$ is now the discretized version of $w(x, t)$. We can regard the discretized PDF as a vector.

Diffusion from Markov Chain

For the Markov process we have a transition probability from a position $x = jl$ to a position $x = il$ given by

$$W_{ij}(\epsilon) = W(il - jl, \epsilon) = \begin{cases} \frac{1}{2} & |i - j| = 1 \\ 0 & \text{else} \end{cases}$$

We call W_{ij} for the transition probability and we can represent it, see below, as a matrix. Our new PDF $w_i(t = \epsilon)$ is now related to the PDF at $t = 0$ through the relation

$$w_i(t = \epsilon) = W(j \rightarrow i)w_j(t = 0).$$

Diffusion from Markov Chain

This equation represents the discretized time-development of an original PDF. Since both W and w represent probabilities, they have to be normalized, i.e., we require that at each time step we have

$$\sum_i w_i(t) = 1,$$

and

$$\sum_j W(j \rightarrow i) = 1.$$

The further constraints are $0 \leq W_{ij} \leq 1$ and $0 \leq w_j \leq 1$. Note that the probability for remaining at the same place is in general not necessarily equal zero. In our Markov process we allow only for jumps to the left or to the right.

Diffusion from Markov Chain

The time development of our initial PDF can now be represented through the action of the transition probability matrix applied n times. At a time $t_n = n\epsilon$ our initial distribution has developed into

$$w_i(t_n) = \sum_j W_{ij}(t_n)w_j(0),$$

and defining

$$W(il - jl, n\epsilon) = (W^n(\epsilon))_{ij}$$

we obtain

$$w_i(n\epsilon) = \sum_j (W^n(\epsilon))_{ij}w_j(0),$$

or in matrix form

$$w(n\epsilon) = \hat{W}^n(\epsilon)\hat{w}(0). \quad (102)$$

Brownian Motion and Markov Processes

We wish to study the time-development of a PDF after a given number of time steps. We define our PDF by the function $w(t)$. In addition we define a transition probability W . The time development of our PDF $w(t)$, after one time-step from $t = 0$ is given by

$$w_i(t = \epsilon) = W(j \rightarrow i)w_j(t = 0).$$

Normally we don't know the form of W !! This equation represents the discretized time-development of an original PDF. We can rewrite this as a

$$w_i(t = \epsilon) = W_{ij}w_j(t = 0).$$

with the transition matrix W for a random walk left or right (cannot stay in the same position) given by

$$W_{ij}(\epsilon) = W(il - jl, \epsilon) = \begin{cases} \frac{1}{2} & |i - j| = 1 \\ 0 & \text{else} \end{cases}$$

Brownian Motion and Markov Processes

Both W and w represent probabilities and they have to be normalized, meaning that at each time step we have

$$\sum_i w_i(t) = 1,$$

and

$$\sum_j W(j \rightarrow i) = 1.$$

Further constraints are $0 \leq W_{ij} \leq 1$ and $0 \leq w_j \leq 1$. We can thus write the action of W as

$$w_i(t+1) = \sum_j W_{ij} w_j(t),$$

or as vector-matrix relation

$$\mathbf{w}(t+1) = \mathbf{W}\mathbf{w}(t)$$

Brownian Motion and Markov Processes, a simple Example

Consider the simple 3×3 matrix \hat{W}

$$\hat{W} = \begin{pmatrix} 1/4 & 1/8 & 2/3 \\ 3/4 & 5/8 & 0 \\ 0 & 1/4 & 1/3 \end{pmatrix},$$

and we choose our initial state as

$$\hat{w}(t=0) = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}.$$

The first iteration is

$$w_i(t=\epsilon) = W(j \rightarrow i) w_j(t=0),$$

resulting in

$$\hat{w}(t=\epsilon) = \begin{pmatrix} 1/4 \\ 3/4 \\ 0 \end{pmatrix}.$$

Brownian Motion and Markov Processes, a simple Example

The next iteration results in

$$w_i(t=2\epsilon) = W(j \rightarrow i) w_j(t=\epsilon),$$

resulting in

$$\hat{w}(t=2\epsilon) = \begin{pmatrix} 5/32 \\ 21/32 \\ 6/32 \end{pmatrix}.$$

Note that the vector \hat{w} is always normalized to 1. We find the steady state of the system by solving the linear set of equations

$$\mathbf{w}(t=\infty) = \mathbf{W}\mathbf{w}(t=\infty).$$

Brownian Motion and Markov Processes, a simple Example

This linear set of equations reads

$$\begin{aligned} W_{11}w_1(t=\infty) + W_{12}w_2(t=\infty) + W_{13}w_3(t=\infty) &= w_1(t=\infty) \\ W_{21}w_1(t=\infty) + W_{22}w_2(t=\infty) + W_{23}w_3(t=\infty) &= w_2(t=\infty) \\ W_{31}w_1(t=\infty) + W_{32}w_2(t=\infty) + W_{33}w_3(t=\infty) &= w_3(t=\infty) \end{aligned} \quad (103)$$

with the constraint that

$$\sum_i w_i(t=\infty) = 1,$$

yielding as solution

$$\hat{w}(t=\infty) = \begin{pmatrix} 4/15 \\ 8/15 \\ 3/15 \end{pmatrix}.$$

Brownian Motion and Markov Processes, a simple Example

Convergence of the simple example

Iteration	w_1	w_2	w_3
0 & 1.00000	0.00000	0.00000	
1 & 0.25000	0.75000	0.00000	
2 & 0.15625	0.62625	0.18750	
3 & 0.24609	0.52734	0.22656	
4	0.27848	0.51416	0.20736
5	0.27213	0.53021	0.19766
6	0.26608	0.53548	0.19844
7	0.26575	0.53424	0.20002
8	0.26656	0.53321	0.20023
9	0.26678	0.53318	0.20005
10	0.26671	0.53332	0.19998
11	0.26666	0.53335	0.20000
12	0.26666	0.53334	0.20000
13	0.26667	0.53333	0.20000
$\hat{w}(t=\infty)$	0.26667	0.53333	0.20000

In a Markov chain Monte Carlo w is normally given, we need

Brownian Motion and Markov Processes, what is happening?

We have after t -steps

$$\hat{\mathbf{w}}(t) = \hat{\mathbf{W}}^t \hat{\mathbf{w}}(0),$$

with $\hat{\mathbf{w}}(0)$ the distribution at $t=0$ and $\hat{\mathbf{W}}$ representing the transition probability matrix. We can always expand $\hat{\mathbf{w}}(0)$ in terms of the right eigenvectors $\hat{\mathbf{v}}_i$ of $\hat{\mathbf{W}}$ as

$$\hat{\mathbf{w}}(0) = \sum_i \alpha_i \hat{\mathbf{v}}_i,$$

resulting in

$$\hat{\mathbf{w}}(t) = \hat{\mathbf{W}}^t \hat{\mathbf{w}}(0) = \hat{\mathbf{W}}^t \sum_i \alpha_i \hat{\mathbf{v}}_i = \sum_i \lambda_i^t \alpha_i \hat{\mathbf{v}}_i,$$

with λ_i the i^{th} eigenvalue corresponding to the eigenvector $\hat{\mathbf{v}}_i$.

Brownian Motion and Markov Processes, what is happening?

If we assume that λ_0 is the largest eigenvector we see that in the limit $t \rightarrow \infty$, $\hat{\mathbf{w}}(t)$ becomes proportional to the corresponding eigenvector $\hat{\mathbf{v}}_0$. This is our steady state or final distribution.

In our discussion below in connection with the entropy of a system and later statistical physics and quantum physics applications, we will relate these properties to correlation functions such as the time-correlation function.

That will allow us to define the so-called *equilibration time*, viz the time needed for the system to reach its most likely state. From that state and on we can compute contributions to various statistical variables.

Brownian Motion and Markov Processes, what is happening?

We anticipate parts of the discussion on statistical physics.

We can relate this property to an observable like the mean magnetization of say a magnetic material. With the probability $\hat{\mathbf{w}}(t)$ we can write the mean magnetization as

$$\langle \mathcal{M}(t) \rangle = \sum_{\mu} \hat{\mathbf{w}}(t)_{\mu} \mathcal{M}_{\mu},$$

or as the scalar of a vector product

$$\langle \mathcal{M}(t) \rangle = \hat{\mathbf{w}}(t) \mathbf{m},$$

with \mathbf{m} being the vector whose elements are the values of \mathcal{M}_{μ} in its various microstates μ .

Recall our definition of an expectation value with a discrete PDF $p(x_i)$:

Brownian Motion and Markov Processes, what is happening?

We rewrite the last relation as

$$\langle \mathcal{M}(t) \rangle = \hat{\mathbf{w}}(t) \mathbf{m} = \sum_i \lambda_i^t \alpha_i \hat{\mathbf{v}}_i \mathbf{m}_i.$$

If we define $m_i = \hat{\mathbf{v}}_i \mathbf{m}_i$ as the expectation value of \mathcal{M} in the i^{th} eigenstate we can rewrite the last equation as

$$\langle \mathcal{M}(t) \rangle = \sum_i \lambda_i^t \alpha_i m_i.$$

Since we have that in the limit $t \rightarrow \infty$ the mean magnetization is dominated by the largest eigenvalue λ_0 , we can rewrite the last equation as

$$\langle \mathcal{M}(t) \rangle = \langle \mathcal{M}(\infty) \rangle + \sum_{i \neq 0} \lambda_i^t \alpha_i m_i.$$

Brownian Motion and Markov Processes, what is happening?

We define the quantity

$$\tau_i = -\frac{1}{\log \lambda_i},$$

and rewrite the last expectation value as

$$\langle \mathcal{M}(t) \rangle = \langle \mathcal{M}(\infty) \rangle + \sum_{i \neq 0} \alpha_i m_i e^{-t/\tau_i}.$$

The quantities τ_i are the correlation times for the system. They control also the time-correlation functions.

The longest correlation time is obviously given by the second largest eigenvalue λ_1 , which normally defines the correlation time discussed above. For large times, this is the only correlation time that survives. If higher eigenvalues of the transition matrix are well separated from λ_1 and we simulate long enough, τ_1 may well define the correlation time. In other cases we may not be able to

Entropy and Equilibrium, section 12.4

The definition of the entropy S (as a dimensionless quantity here) is

$$S = -\sum_i w_i \ln(w_i),$$

where w_i is the probability of finding our system in a state i . For our one-dimensional random walk it represents the probability for being at position $i = i\Delta x$ after a given number of time steps. Assume now that we have N random walkers at $i = 0$ and $t = 0$ and let these random walkers diffuse as function of time.

Entropy and Equilibrium, section 12.4

We compute then the probability distribution for N walkers after a given number of steps i along x and time steps j . We can then compute an entropy S_j for a given number of time steps by summing over all probabilities i . The code used to compute these results is in programs/chapter12/program4.cpp. Here we have used 100 walkers on a lattice of length from $L = -50$ to $L = 50$ employing periodic boundary conditions meaning that if a walker reaches the point $x = L$ it is shifted to $x = -L$ and if $x = -L$ it is shifted to $x = L$.

Entropy

```
// loop over all time steps
for (int step=1; step <= time_steps; step++){
    // move all walkers with periodic boundary conditions
    for (int walks = 1; walks <= walkers; walks++){
        if (ran0(hidum) <= move_probability) {
            if ( x[walks] +1 > length) {
                x[walks] = -length;
            }
            else{
                x[walks] += 1;
            }
        }
    }
}
```

Entropy

```
        else {
            if ( x[walks] -1 < -length) {
                x[walks] = length;
            }
            else{
                x[walks] -= 1;
            }
        } // end of loop over walks
    } // end of loop over trials
}
```

Entropy

```
// at the final time step we compute the probability
// by counting the number of walkers at every position
for ( int i = -length; i <= length; i++){
    int count = 0;
    for( int j = 1; j <= walkers; j++){
        if ( x[j] == i) {
            count += 1;
        }
    }
    probability[i+length] = count;
}
```

Entropy

```
// Writes the results to screen
void output(int length, int time_steps, int walkers, int *probability)
{
    double entropy, histogram;
    // find norm of probability
    double norm = 1.0/walkers;
    // compute the entropy
    entropy = 0.; histogram = 0.;
    for( int i = -length; i <= length; i++){
        histogram = (double) probability[i+length]*norm;
        if ( histogram > 0.0) {
            entropy -= histogram*log(histogram);
        }
    }
    // then write entropy to file
}
```

Entropy

At small time steps the entropy is very small, reflecting the fact that we have an ordered state. As time elapses, the random walkers spread out in space (here in one dimension) and the entropy increases as there are more states, that is positions accessible to the system. We say that the system shows an increased degree of disorder. After several time steps, we see that the entropy reaches a constant value, a situation called a steady state. This signals that the system has reached its equilibrium situation and that the random walkers spread out to occupy all possible available states. At equilibrium it means thus that all states are equally probable and this is not baked into any dynamical equations such as Newton's law of motion.

Entropy

It occurs because the system is allowed to explore all possibilities. An important hypothesis is the ergodic hypothesis which states that in equilibrium all available states of a closed system have equal probability. This hypothesis states also that if we are able to simulate long enough, then one should be able to trace through all possible paths in the space of available states to reach the equilibrium situation. Our Markov process should be able to reach any state of the system from any other state if we run for long enough.

Detailed Balance

In a Markov Monte Carlo w is normally given, we need to find W ! But we need to find which distribution we obtain when the steady state has been achieved.

Markov process with transition probability from a state j to another state i

$$\sum_j W(j \rightarrow i) = 1$$

Note that the probability for remaining at the same place is not necessarily equal zero.

PDF w_i at time $t = n\epsilon$

$$w_i(t) = \sum_j W(j \rightarrow i)^n w_j(t=0)$$

$$\sum_i w_i(t) = 1$$

Detailed Balance

Detailed balance condition

$$\sum_i W(j \rightarrow i) w_j = \sum_i W(i \rightarrow j) w_i$$

Ensures that it is the correct distribution which is achieved when equilibrium is reached.

When a Markov process reaches equilibrium we have

$$\mathbf{w}(t = \infty) = \mathbf{W}\mathbf{w}(t = \infty)$$

General condition at equilibrium

$$W(j \rightarrow i) w_j = W(i \rightarrow j) w_i$$

which is the detailed balance condition. Proof is simple.

Detailed Balance

To derive the conditions for equilibrium, we start from the so-called Master equation, which relates the temporal development of a PDF $w_i(t)$. The equation is given

$$\frac{dw_i(t)}{dt} = \sum_j [W(j \rightarrow i) w_j - W(i \rightarrow j) w_i],$$

which simply states that the rate at which the systems moves from a state j to a final state i (the first term on the right-hand side of the last equation) is balanced by the rate at which the systems undergoes transitions from the state i to a state j (the second term). If we have reached the so-called steady state, then the temporal dependence is zero. This means that in equilibrium we have

$$\frac{dw_i(t)}{dt} = 0.$$

Ergodicity

It should be possible for any Markov process to reach every possible state of the system from any starting point if the simulations is carried out for a long enough time.

Any state in a Boltzmann distribution has a probability different from zero and if such a state cannot be reached from a given starting point, then the system is not ergodic.

Example: Boltzmann Distribution

At equilibrium detailed balance gives

$$\frac{W(j \rightarrow i)}{W(i \rightarrow j)} = \frac{w_i}{w_j}$$

Boltzmann distribution

$$\frac{w_i}{w_j} = \exp(-\beta(E_i - E_j))$$

Selection Rule

In general

$$W(i \rightarrow j) = g(i \rightarrow j) A(i \rightarrow j)$$

where g is a selection probability while A is the probability for accepting a move. It is also called the acceptance ratio.

With detailed balance this gives

$$\frac{g(j \rightarrow i) A(j \rightarrow i)}{g(i \rightarrow j) A(i \rightarrow j)} = \exp(-\beta(E_i - E_j))$$

Metropolis Algorithm

For a system which follows the Boltzmann distribution the Metropolis algorithm reads

$$A(j \rightarrow i) = \begin{cases} \exp(-\beta(E_i - E_j)) & E_i - E_j > 0 \\ 1 & \text{else} \end{cases}$$

This algorithm satisfies the condition for detailed balance and ergodicity.

Implementation

- Establish an initial energy E_b
- Do a random change of this initial state by e.g., flipping an individual spin. This new state has energy E_t . Compute then $\Delta E = E_t - E_b$
- If $\Delta E \leq 0$ accept the new configuration.
- If $\Delta E > 0$, compute $w = e^{-(\beta \Delta E)}$.
- Compare w with a random number r . If $r \leq w$ accept, else keep the old configuration.
- Compute the terms in the sums $\sum A_s P_s$.
- Repeat the above steps in order to have a large enough number of microstates
- For a given number of MC cycles, compute then expectation values.

Test of the Metropolis Algorithm

Want to show that the Metropolis algorithm generates the Boltzmann distribution

$$P(\beta) = \frac{e^{-\beta E}}{Z},$$

with $\beta = 1/kT$ being the inverse temperature, E is the energy of the system and Z is the partition function. The only functions you will need are those to generate random numbers.

We are going to study one single particle in equilibrium with its surroundings, the latter modeled via a large heat bath with temperature T .

The model used to describe this particle is that of an ideal gas in **one** dimension and with velocity $-v$ or v . We are interested in finding $P(v)dv$, which expresses the probability for finding the system with a given velocity $v \in [v, v + dv]$. The energy for this one-dimensional system is

$$E = \frac{1}{2}kT = \frac{1}{2}v^2$$

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Test of the Metropolis Algorithm, closed form results

The partition function of the system of interest is:

$$Z = \int_{-\infty}^{+\infty} e^{-\beta v^2/2} dv = \sqrt{2\pi\beta}^{-1/2}$$

The mean velocity

$$\langle v \rangle = \int_{-\infty}^{+\infty} v e^{-\beta v^2/2} dv = 0$$

The expressions for $\langle E \rangle$ and σ_E assume the following form:

$$\langle E \rangle = \int_{-\infty}^{+\infty} \frac{v^2}{2} e^{-\beta v^2/2} dv = -\frac{1}{Z} \frac{\partial Z}{\partial \beta} = \frac{1}{2} \beta^{-1} = \frac{1}{2} T$$

$$\langle E^2 \rangle = \int_{-\infty}^{+\infty} \frac{v^4}{4} e^{-\beta v^2/2} dv = \frac{1}{Z} \frac{\partial^2 Z}{\partial \beta^2} = \frac{3}{4} \beta^{-2} = \frac{3}{4} T^2$$

and

Test of the Metropolis Algorithm

```
for( montecarlo_cycles=1; Max_cycles; montecarlo_cycles++) {
    ...
    // change speed as function of delta v
    v_change = (2*rani(kidum) - 1) * delta_v;
    v_new = v_old + v_change;
    // energy change
    delta_E = 0.5*(v_new*v_new - v_old*v_old) ;
    .....
    // Metropolis algorithm begins here
    if ( rani(kidum) <= exp(-beta*delta_E) ) {
        accept_step = accept_step + 1 ;
        v_old = v_new ;
    }

    // thereafter we must fill in P[N] as a function of
    // the new speed
    // upgrade mean velocity, energy and variance
}
```


Test of the Metropolis Algorithm

Analytical vs numerical results. $T = 4$, 10^8 MC tries, $\Delta v = 0.2$

Observable & Analytical value & Numerical value	
$\langle v \rangle$	0.00000
$\langle E \rangle$	2.00000 & 1.99855
σ_E	8.00000 & 8.06669

Code for Metropolis test

```
v_current = v0;

// start simulation
ofile.open("evamc.dat");
for (tries = 1; tries <= MC; tries++){

    v_change = (2.*ran0(&idum) - 1.) * dv;
    v_trial = v_current + v_change;

    // evaluate dE
    delta_E = 0.5 * ( v_trial * v_trial - v_current * v_current );
```

Code for Metropolis test

```
// Metropolis test
if (delta_E <= 0) {
    acceptance++; v_current = v_trial;
} else if (ran0(&idum) <= exp( -beta * delta_E )){
    acceptance++; v_current = v_trial;

// check if velocity value lies within given limits
if (abs(v_current) > v_max) {
    cout<<"Velocity out of range."; exit(1);
}
```

Code for Metropolis test

```
// save event in P array
address = (int) floor( v_current / dv ) + N/2 + 1;
P[address]++;

// update mean velocity, mean energy and energy variance values
mean_v += v_current;
mean_E += 0.5 * v_current * v_current;
E_variance += 0.25 * v_current * v_current * v_current * v_curr
```

Code for Metropolis test

```
// initialize model parameters
beta = 1./T; v_max = 10. * sqrt( T);
// calculate amount of P-array elements
N = 2 * (int)(v_max/dv) + 1;
// initialize P-array
P = new int [N];

for (int i=0; i < N; i++) P[i] = 0;

mean_v = 0.; mean_E = 0.; E_variance = 0.;
acceptance = 0;

} // initialize
```

Overview of week 47

Summary and exam discussion.

- Monday:
 - Brief repetition from last week and discussion of projects
 - Metropolis algorithm and Variational Monte Carlo
 - Emphasis on the Monte Carlo project (integration and Variational Monte Carlo)
- Tuesday:
 - Discussion of projects and parallelization (all projects)
 - Discussion of Monte Carlo simulation of Markov chains (linked with the diffusion project)

Pros and Cons of Quantum Monte Carlo

- Is physically intuitive.
- Allows one to study systems with many degrees of freedom. Diffusion Monte Carlo (DMC) and Green's function Monte Carlo (GFMC) yield in principle the exact solution to Schrödinger's equation.
- Variational Monte Carlo (VMC) is easy to implement but needs a reliable trial wave function, can be difficult to obtain.
- DMC/GFMC for fermions (spin with half-integer values, electrons, baryons, neutrinos, quarks) has a sign problem. Nature prefers an anti-symmetric wave function. PDF in this case given distribution of random walkers ($p \geq 0$).
- The solution has a statistical error, which can be large.
- There is a limit for how large systems one can study, DMC needs a huge number of random walkers in order to achieve stable results.
- Obtain only the lowest-lying states with a given symmetry. Can get excited states.

Where and why do we use Monte Carlo Methods in Quantum Physics

- Quantum systems with many particles at finite temperature: Path Integral Monte Carlo with applications to dense matter and quantum liquids (phase transitions from normal fluid to superfluid). Strong correlations.
- Bose-Einstein condensation of dilute gases, method transition from non-linear PDE to Diffusion Monte Carlo as density increases.
- Light atoms, molecules, solids and nuclei.
- Lattice Quantum-Chromo Dynamics. Impossible to solve without MC calculations.
- Simulations of systems in solid state physics, from semiconductors to spin systems. Many electrons active and possibly strong correlations. Chapter 13 on statistical physics simulations will not be discussed this year.

Quantum Monte Carlo and Schrödinger's equation

For one-body problems (one dimension)

$$-\frac{\hbar^2}{2m}\nabla^2\Psi(x,t) + V(x,t)\Psi(x,t) = i\hbar\frac{\partial\Psi(x,t)}{\partial t},$$

$$P(x,t) = \Psi(x,t)^*\Psi(x,t)$$

$$P(x,t)dx = \Psi(x,t)^*\Psi(x,t)dx$$

Interpretation: probability of finding the system in a region between x and $x + dx$. Always real

$$\Psi(x,t) = R(x,t) + iI(x,t)$$

yielding

$$\Psi(x,t)^*\Psi(x,t) = (R - iI)(R + iI) = R^2 + I^2$$

Quantum Monte Carlo and Schrödinger's equation

Petit digression

Choose $\tau = i\hbar/t$.

The time-dependent (1-dim) Schrödinger equation becomes then

$$\frac{\partial\Psi(x,\tau)}{\partial\tau} = \frac{\hbar^2}{2m}\frac{\partial^2\Psi(x,\tau)}{\partial x^2} - V(x,\tau)\Psi(x,\tau).$$

With $V = 0$ we have a diffusion equation in complex time with diffusion constant

$$D = \frac{\hbar^2}{2m}.$$

Used in diffusion Monte Carlo calculations. Topic for FYS4411, Computational Physics II

Quantum Monte Carlo and Schrödinger's equation

Conditions which Ψ has to satisfy:

Normalization

$$\int_{-\infty}^{\infty} P(x,t)dx = \int_{-\infty}^{\infty} \Psi(x,t)^*\Psi(x,t)dx = 1$$

meaning that

$$\int_{-\infty}^{\infty} \Psi(x,t)^*\Psi(x,t)dx < \infty$$

And

- $\Psi(x,t)$ and $\partial\Psi(x,t)/\partial x$ must be finite,
- $\Psi(x,t)$ and $\partial\Psi(x,t)/\partial x$ must be continuous,
- $\Psi(x,t)$ and $\partial\Psi(x,t)/\partial x$ must be single valued square integrable functions.

First Postulate

Any physical quantity $A(\vec{r},\vec{p})$ which depends on position \vec{r} and momentum \vec{p} has a corresponding quantum mechanical operator by replacing $\vec{p} \rightarrow i\hbar\vec{\nabla}$, yielding the quantum mechanical operator

$$\hat{A} = A(\vec{r}, -i\hbar\vec{\nabla}).$$

Quantity & Classical definition & QM operator	
Position	\vec{r}
Momentum	\vec{p}
Orbital momentum	$\vec{L} = \vec{r} \times \vec{p}$
Kinetic energy	$T = (\vec{p})^2/2m$
Total energy	$H = (\vec{p}^2/2m) + V(\vec{r})$

Second Postulate

The only possible outcome of an ideal measurement of the physical quantity A are the eigenvalues of the corresponding quantum mechanical operator $\hat{\mathbf{A}}$.

$$\hat{\mathbf{A}}\psi_\nu = a_\nu\psi_\nu,$$

resulting in the eigenvalues a_1, a_2, a_3, \dots as the only outcomes of a measurement. The corresponding eigenstates $\psi_1, \psi_2, \psi_3, \dots$ contain all relevant information about the system.

Third Postulate

Assume Φ is a linear combination of the eigenfunctions ψ_ν for $\hat{\mathbf{A}}$,

$$\Phi = c_1\psi_1 + c_2\psi_2 + \dots = \sum_\nu c_\nu\psi_\nu.$$

The eigenfunctions are orthogonal and we get

$$c_\nu = \int (\Phi)^*\psi_\nu d\tau.$$

From this we can formulate the third postulate: When the eigenfunction is Φ , the probability of obtaining the value a_ν as the outcome of a measurement of the physical quantity A is given by $|c_\nu|^2$ and ψ_ν is an eigenfunction of $\hat{\mathbf{A}}$ with eigenvalue a_ν .

Third Postulate

As a consequence one can show that: when a quantal system is in the state Φ , the mean value or expectation value of a physical quantity $A(\vec{r}, \vec{p})$ is given by

$$\langle A \rangle = \int (\Phi)^*\hat{\mathbf{A}}(\vec{r}, -i\hbar\vec{\nabla})\Phi d\tau.$$

We have assumed that Φ has been normalized, viz., $\int (\Phi)^*\Phi d\tau = 1$. Else

$$\langle A \rangle = \frac{\int (\Phi)^*\hat{\mathbf{A}}\Phi d\tau}{\int (\Phi)^*\Phi d\tau}.$$

Fourth Postulate

The time development of a quantal system is given by

$$i\hbar\frac{\partial\Psi}{\partial t} = \hat{\mathbf{H}}\Psi,$$

with $\hat{\mathbf{H}}$ the quantal Hamiltonian operator for the system.

Quantum Monte Carlo

Most quantum mechanical problems of interest in e.g., atomic, molecular, nuclear and solid state physics consist of a large number of interacting electrons and ions or nucleons. The total number of particles N is usually sufficiently large that an exact solution cannot be found. Typically, the expectation value for a chosen hamiltonian for a system of N particles is

$$\langle H \rangle =$$

$$\frac{\int d\mathbf{R}_1 d\mathbf{R}_2 \dots d\mathbf{R}_N \Psi^*(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N) H(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N) \Psi(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N)}{\int d\mathbf{R}_1 d\mathbf{R}_2 \dots d\mathbf{R}_N \Psi^*(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N) \Psi(\mathbf{R}_1, \mathbf{R}_2, \dots, \mathbf{R}_N)}$$

an in general intractable problem.

Quantum Monte Carlo

Given a hamiltonian H and a trial wave function Ψ_T , the variational principle states that the expectation value of $\langle H \rangle$, defined through

$$E[H] = \langle H \rangle = \frac{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) H(\mathbf{R}) \Psi_T(\mathbf{R})}{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) \Psi_T(\mathbf{R})},$$

is an upper bound to the ground state energy E_0 of the hamiltonian H , that is

$$E_0 \leq \langle H \rangle.$$

In general, the integrals involved in the calculation of various expectation values are multi-dimensional ones. Traditional integration methods such as the Gauss-Legendre will not be adequate for say the computation of the energy of a many-body system.

Quantum Monte Carlo

The trial wave function can be expanded in the eigenstates of the hamiltonian since they form a complete set, viz.,

$$\Psi_T(\mathbf{R}) = \sum_i a_i \Psi_i(\mathbf{R}),$$

and assuming the set of eigenfunctions to be normalized one obtains

$$\frac{\sum_{nm} a_m^* a_n \int d\mathbf{R} \Psi_m^*(\mathbf{R}) H(\mathbf{R}) \Psi_n(\mathbf{R})}{\sum_{nm} a_m^* a_n \int d\mathbf{R} \Psi_m^*(\mathbf{R}) \Psi_n(\mathbf{R})} = \frac{\sum_n a_n^2 E_n}{\sum_n a_n^2} \geq E_0,$$

where we used that $H(\mathbf{R}) \Psi_n(\mathbf{R}) = E_n \Psi_n(\mathbf{R})$. In general, the integrals involved in the calculation of various expectation values are multi-dimensional ones. The variational principle yields the lowest state of a given symmetry.

Quantum Monte Carlo

In most cases, a wave function has only small values in large parts of configuration space, and a straightforward procedure which uses homogeneously distributed random points in configuration space will most likely lead to poor results. This may suggest that some kind of importance sampling combined with e.g., the Metropolis algorithm may be a more efficient way of obtaining the ground state energy. The hope is then that those regions of configurations space where the wave function assumes appreciable values are sampled more efficiently.

Quantum Monte Carlo

The tedious part in a VMC calculation is the search for the variational minimum. A good knowledge of the system is required in order to carry out reasonable VMC calculations. This is not always the case, and often VMC calculations serve rather as the starting point for so-called diffusion Monte Carlo calculations (DMC). DMC is a way of solving exactly the many-body Schrödinger equation by means of a stochastic procedure. A good guess on the binding energy and its wave function is however necessary. A carefully performed VMC calculation can aid in this context.

Quantum Monte Carlo

Construct first a trial wave function $\psi_T^\alpha(\mathbf{R})$, for a many-body system consisting of N particles located at positions $\mathbf{R} = (\mathbf{R}_1, \dots, \mathbf{R}_N)$. The trial wave function depends on α variational parameters $\alpha = (\alpha_1, \dots, \alpha_N)$.

Then we evaluate the expectation value of the hamiltonian H

$$E[H] = \langle H \rangle = \frac{\int d\mathbf{R} \Psi_T^\alpha(\mathbf{R}) H(\mathbf{R}) \Psi_T^\alpha(\mathbf{R})}{\int d\mathbf{R} \Psi_T^\alpha(\mathbf{R}) \Psi_T^\alpha(\mathbf{R})}.$$

Thereafter we vary α according to some minimization algorithm and return to the first step.

Quantum Monte Carlo

Choose a trial wave function $\psi_T(\mathbf{R})$.

$$P(\mathbf{R}) = \frac{|\psi_T(\mathbf{R})|^2}{\int |\psi_T(\mathbf{R})|^2 d\mathbf{R}}.$$

This is our new probability distribution function (PDF). The approximation to the expectation value of the Hamiltonian is now

$$E[H] \approx \frac{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) H(\mathbf{R}) \Psi_T(\mathbf{R})}{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) \Psi_T(\mathbf{R})}.$$

Define a new quantity

$$E_L(\mathbf{R}) = \frac{1}{\psi_T(\mathbf{R})} H \psi_T(\mathbf{R}),$$

called the local energy, which, together with our trial PDF yields

$$E[H] \approx \langle H \rangle \approx \int P(\mathbf{R}) E_L(\mathbf{R}) d\mathbf{R} \approx \frac{1}{N} \sum_{i=1}^N E_L(\mathbf{R}_i)$$

Quantum Monte Carlo

Algo:

- Initialisation: Fix the number of Monte Carlo steps. Choose an initial \mathbf{R} and variational parameters α and calculate $|\psi_T^\alpha(\mathbf{R})|^2$.
- Initialise the energy and the variance and start the Monte Carlo calculation (thermalize)
 - 1 Calculate a trial position $\mathbf{R}_p = \mathbf{R} + r * \text{step}$ where r is a random variable $r \in [0, 1]$.
 - 2 Metropolis algorithm to accept or reject this move $w = P(\mathbf{R}_p)/P(\mathbf{R})$.
 - 3 If the step is accepted, then we set $\mathbf{R} = \mathbf{R}_p$. Update averages
- Finish and compute final averages.

Observe that the jumping in space is governed by the variable *step*. Called brute-force sampling. Need importance sampling to get more relevant sampling.

Quantum Monte Carlo

The radial Schrodinger equation for the hydrogen atom can be written as

$$-\frac{\hbar^2}{2m} \frac{\partial^2 u(r)}{\partial r^2} - \left(\frac{ke^2}{r} - \frac{\hbar^2 l(l+1)}{2mr^2} \right) u(r) = Eu(r),$$

or with dimensionless variables

$$-\frac{1}{2} \frac{\partial^2 u(\rho)}{\partial \rho^2} - \frac{u(\rho)}{\rho} + \frac{l(l+1)}{2\rho^2} u(\rho) - \lambda u(\rho) = 0,$$

with the hamiltonian

$$H = -\frac{1}{2} \frac{\partial^2}{\partial \rho^2} - \frac{1}{\rho} + \frac{l(l+1)}{2\rho^2}.$$

Use variational parameter α in the trial wave function

$$u_T^\alpha(\rho) = \alpha \rho e^{-\alpha \rho}.$$

Quantum Monte Carlo

Inserting this wave function into the expression for the local energy E_L gives

$$E_L(\rho) = -\frac{1}{\rho} - \frac{\alpha}{2} \left(\alpha - \frac{2}{\rho} \right).$$

α	$\langle H \rangle$	σ^2	σ
7.00000E-01	-4.57759E-01 & 4.51201E-02 & 6.71715E-04		
8.00000E-01	-4.81461E-01 & 3.05736E-02 & 5.52934E-04		
9.00000E-01	-4.95899E-01 & 8.20497E-03 & 2.86443E-04		
1.00000E-00	-5.00000E-01 & 0.00000E+00 & 0.00000E+00		
1.10000E+00	-4.93738E-01 & 1.16989E-02 & 3.42036E-04		
1.20000E+00	-4.75563E-01 & 8.85899E-02 & 9.41222E-04		
1.30000E+00	-4.54341E-01 & 1.45171E-01 & 1.20487E-03		

Quantum Monte Carlo

We note that at $\alpha = 1$ we obtain the exact result, and the variance is zero, as it should. The reason is that we then have the exact wave function, and the action of the hamiltonian on the wave function

$$H\psi = \text{constant} \times \psi,$$

yields just a constant. The integral which defines various expectation values involving moments of the hamiltonian becomes then

$$\langle H^n \rangle = \frac{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) H^n(\mathbf{R}) \Psi_T(\mathbf{R})}{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) \Psi_T(\mathbf{R})} = \text{constant} \times \frac{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) \Psi_T(\mathbf{R})}{\int d\mathbf{R} \Psi_T^*(\mathbf{R}) \Psi_T(\mathbf{R})} = \text{constant}$$

This gives an important information: the exact wave function leads to zero variance! Variation is then performed by minimizing both the energy and the variance.

Quantum Monte Carlo

The helium atom consists of two electrons and a nucleus with charge $Z = 2$. The contribution to the potential energy due to the attraction from the nucleus is

$$-\frac{2ke^2}{r_1} - \frac{2ke^2}{r_2},$$

and if we add the repulsion arising from the two interacting electrons, we obtain the potential energy

$$V(r_1, r_2) = -\frac{2ke^2}{r_1} - \frac{2ke^2}{r_2} + \frac{ke^2}{r_{12}},$$

with the electrons separated at a distance $r_{12} = |\mathbf{r}_1 - \mathbf{r}_2|$.

Quantum Monte Carlo

The hamiltonian becomes then

$$\hat{H} = -\frac{\hbar^2 \nabla_1^2}{2m} - \frac{\hbar^2 \nabla_2^2}{2m} - \frac{2ke^2}{r_1} - \frac{2ke^2}{r_2} + \frac{ke^2}{r_{12}},$$

and Schrödingers equation reads

$$\hat{H}\psi = E\psi.$$

All observables are evaluated with respect to the probability distribution

$$P(\mathbf{R}) = \frac{|\psi_T(\mathbf{R})|^2}{\int |\psi_T(\mathbf{R})|^2 d\mathbf{R}}.$$

generated by the trial wave function. The trial wave function must approximate an exact eigenstate in order that accurate results are to be obtained. Improved trial wave functions also improve the importance sampling, reducing the cost of obtaining a certain statistical accuracy.

Quantum Monte Carlo

Choice of trial wave function for Helium: Assume $r_1 \rightarrow 0$.

$$E_L(\mathbf{R}) = \frac{1}{\psi_T(\mathbf{R})} H \psi_T(\mathbf{R}) = \frac{1}{\psi_T(\mathbf{R})} \left(-\frac{1}{2} \nabla_1^2 - \frac{Z}{r_1} \right) \psi_T(\mathbf{R}) + \text{finite terms}.$$

$$E_L(R) = \frac{1}{\mathcal{R}_T(r_1)} \left(-\frac{1}{2} \frac{d^2}{dr_1^2} - \frac{1}{r_1} \frac{d}{dr_1} - \frac{Z}{r_1} \right) \mathcal{R}_T(r_1) + \text{finite terms}$$

For small values of r_1 , the terms which dominate are

$$\lim_{r_1 \rightarrow 0} E_L(R) = \frac{1}{\mathcal{R}_T(r_1)} \left(-\frac{1}{r_1} \frac{d}{dr_1} - \frac{Z}{r_1} \right) \mathcal{R}_T(r_1),$$

since the second derivative does not diverge due to the finiteness of Ψ at the origin.

Quantum Monte Carlo

This results in

$$\frac{1}{\mathcal{R}_T(r_1)} \frac{d\mathcal{R}_T(r_1)}{dr_1} = -Z,$$

and

$$\mathcal{R}_T(r_1) \propto e^{-Zr_1}.$$

A similar condition applies to electron 2 as well. For orbital momenta $l > 0$ we have

$$\frac{1}{\mathcal{R}_T(r)} \frac{d\mathcal{R}_T(r)}{dr} = -\frac{Z}{l+1}.$$

Similarlry, studying the case $r_{12} \rightarrow 0$ we can write a possible trial wave function as

$$\psi_T(\mathbf{R}) = e^{-\alpha(r_1+r_2)} e^{\beta r_{12}}.$$

VMC code for helium, vmc_para.cpp

```
// Here we define global variables used in various functions
// These can be changed by reading from file the different parameters
int dimension = 3; // three-dimensional system
int charge = 2; // we fix the charge to be that of the helium atom
int my_rank, numprocs; // these are the parameters used by MPI
double step_length = 1.0; // define which node and how many
int number_particles = 2; // we fix the brute force jump to 1 Bohr
// we fix also the number of electrons
```

VMC code for helium, vmc_para.cpp, main part

```
// MPI initializations
MPI_Init (&argc, &argv);
MPI_Comm_size (MPI_COMM_WORLD, &numprocs);
MPI_Comm_rank (MPI_COMM_WORLD, &my_rank);
time_start = MPI_Wtime();

if (my_rank == 0 && argc <= 2) {
    cout << "Bad Usage: " << argv[0] << endl;
    " read also output file on same line" << endl;
    exit(1);
}

if (my_rank == 0 && argc > 2) {
    outfilename=argv[1];
    ofile.open(outfilename);
}
```

VMC code for helium, vmc_para.cpp, main part

```
// Setting output file name for this rank:
ostringstream ost;
ost << "blocks_rank" << my_rank << ".dat";
// Open file for writing:
blockofile.open(ost.str().c_str(), ios::out | ios::binary);

total_cumulative_e = new double[max_variations+1];
total_cumulative_e2 = new double[max_variations+1];
cumulative_e = new double[max_variations+1];
cumulative_e2 = new double[max_variations+1];

// initialize the arrays by zeroing them
for( i=1; i <= max_variations; i++){
    cumulative_e[i] = cumulative_e2[i] = 0.0;
    total_cumulative_e[i] = total_cumulative_e2[i] = 0.0;
}
```

VMC code for helium, vmc_para.cpp, main part

```
// broadcast the total number of variations
MPI_Bcast (&max_variations, 1, MPI_INT, 0, MPI_COMM_WORLD);
MPI_Bcast (&number_cycles, 1, MPI_INT, 0, MPI_COMM_WORLD);
total_number_cycles = number_cycles*numprocs;
// array to store all energies for last variation of alpha
all_energies = new double[number_cycles+1];
// Do the mc sampling and accumulate data with MPI_Reduce
mc_sampling(max_variations, number_cycles, cumulative_e,
            cumulative_e2, all_energies);
// Collect data in total averages
for( i=1; i <= max_variations; i++){
    MPI_Reduce(&cumulative_e[i], &total_cumulative_e[i], 1, MPI_DOUBLE,
              MPI_SUM, 0, MPI_COMM_WORLD);
    MPI_Reduce(&cumulative_e2[i], &total_cumulative_e2[i], 1, MPI_DOUBLE,
              MPI_SUM, 0, MPI_COMM_WORLD);
}
```

VMC code for helium, vmc_para.cpp, main part

```
blockofile.write((char*)(all_energies+1),
                number_cycles*sizeof(double));
blockofile.close();
delete [] total_cumulative_e; delete [] total_cumulative_e2;
delete [] cumulative_e; delete [] cumulative_e2; delete [] all_energies;
// End MPI
MPI_Finalize ();
return 0;
} // end of main function
```

VMC code for helium, vmc_para.cpp, sampling

```
alpha = 0.5*charge;
// every node has its own seed for the random numbers
idum = -1*my_rank;
// allocate matrices which contain the position of the particles
r_old = (double **)matrix(number_particles,dimension,sizeof(double))
r_new = (double **)matrix(number_particles,dimension,sizeof(double))
for (i = 0; i < number_particles; i++) {
    for ( j=0; j < dimension; j++) {
        r_old[i][j] = r_new[i][j] = 0;
    }
}
// loop over variational parameters
```

VMC code for helium, vmc_para.cpp, sampling

```
for (variate=1; variate <= max_variations; variate++){
    // initialisations of variational parameters and energies
    alpha += 0.1;
    energy = energy2 = 0; accept = 0; delta_e=0;
    // initial trial position, note calling with alpha
    for (i = 0; i < number_particles; i++) {
        for ( j=0; j < dimension; j++) {
            r_old[i][j] = step_length*(ran2(&idum)-0.5);
        }
    }
    wfold = wave_function(r_old, alpha);
```

VMC code for helium, vmc_para.cpp, sampling

```
// loop over monte carlo cycles
for (cycles = 1; cycles <= number_cycles; cycles++){
    // new position
    for (i = 0; i < number_particles; i++) {
        for ( j=0; j < dimension; j++) {
            r_new[i][j] = r_old[i][j]+step_length*(ran2(&idum)-0.5);
        }
    }
    // for the other particles we need to set the position to the old
    // we move only one particle at the time
    for (k = 0; k < number_particles; k++) {
        if (k != i) {
            for ( j=0; j < dimension; j++) {
                r_new[k][j] = r_old[k][j];
            }
        }
    }
}
```

VMC code for helium, vmc_para.cpp, sampling

```
    wfnew = wave_function(r_new, alpha);
    // The Metropolis test is performed by moving one particle at the time
    if(ran2(&idum) <= wfnew*wfnew/wfold/wfold) {
        for ( j=0; j < dimension; j++) {
            r_old[i][j]=r_new[i][j];
        }
        wfold = wfnew;
    } // end of loop over particles
```

VMC code for helium, vmc_para.cpp, sampling

```
// compute local energy
delta_e = local_energy(r_old, alpha, wfold);
// save all energies on last variate
if(variate==max_variations){
    all_energies[cycles] = delta_e;
}
// update energies
energy += delta_e;
energy2 += delta_e*delta_e;
} // end of loop over MC trials
// update the energy average and its squared
cumulative_e[variate] = energy;
cumulative_e2[variate] = energy2;
} // end of loop over variational steps
```

VMC code for helium, vmc_para.cpp, wave function

```
// Function to compute the squared wave function, simplest form
double wave_function(double **r, double alpha)
{
    int i, j, k;
    double wf, argument, r_single_particle, r_12;

    argument = wf = 0;
    for (i = 0; i < number_particles; i++) {
        r_single_particle = 0;
        for (j = 0; j < dimension; j++) {
            r_single_particle += r[i][j]*r[i][j];
        }
        argument += sqrt(r_single_particle);
    }
    wf = exp(-argument*alpha);
    return wf;
}
```

VMC code for helium, vmc_para.cpp, local energy

```
// Function to calculate the local energy with num derivative
double local_energy(double **r, double alpha, double wfold)
{
    int i, j, k;
    double e_local, wfminus, wfplus, e_kinetic, e_potential, r_12,
           r_single_particle;
    double **r_plus, **r_minus;
```

VMC code for helium, vmc_para.cpp, local energy

```
// allocate matrices which contain the position of the particles
// the function matrix is defined in the program library
r_plus = (double **)matrix(number_particles, dimension, sizeof(double));
r_minus = (double **)matrix(number_particles, dimension, sizeof(double));
for (i = 0; i < number_particles; i++) {
    for (j = 0; j < dimension; j++) {
        r_plus[i][j] = r_minus[i][j] = r[i][j];
```

VMC code for helium, vmc_para.cpp, local energy

```
// compute the kinetic energy
e_kinetic = 0;
for (i = 0; i < number_particles; i++) {
    for (j = 0; j < dimension; j++) {
        r_plus[i][j] = r[i][j] + h;
        r_minus[i][j] = r[i][j] - h;
        wfminus = wave_function(r_minus, alpha);
        wfplus = wave_function(r_plus, alpha);
        e_kinetic += (wfminus + wfplus - 2 * wfold);
        r_plus[i][j] = r[i][j];
        r_minus[i][j] = r[i][j];
    }
}

// include electron mass and hbar squared and divide by wave function
e_kinetic = 0.5 * h2 * e_kinetic / wfold;
```

VMC code for helium, vmc_para.cpp, local energy

```
// compute the potential energy
e_potential = 0;
// contribution from electron-proton potential
for (i = 0; i < number_particles; i++) {
    r_single_particle = 0;
    for (j = 0; j < dimension; j++) {
        r_single_particle += r[i][j] * r[i][j];
    }
    e_potential += charge / sqrt(r_single_particle);
}
```

VMC code for helium, vmc_para.cpp, local energy

```
// contribution from electron-electron potential
for (i = 0; i < number_particles - 1; i++) {
    for (j = i + 1; j < number_particles; j++) {
        r_12 = 0;
        for (k = 0; k < dimension; k++) {
            r_12 += (r[i][k] - r[j][k]) * (r[i][k] - r[j][k]);
        }
        e_potential += 1 / sqrt(r_12);
    }
}
```

Structuring the code

During the development of our code we need to make several checks. It is also very instructive to compute a closed form expression for the local energy. Since our wave function is rather simple it is straightforward to find an analytic expressions. Consider first the case of the simple helium function

$$\Psi_T(\mathbf{r}_1, \mathbf{r}_2) = e^{-\alpha(r_1 + r_2)}$$

The local energy is for this case

$$E_{L1} = (\alpha - Z) \left(\frac{1}{r_1} + \frac{1}{r_2} \right) + \frac{1}{r_{12}} - \alpha^2$$

which gives an expectation value for the local energy given by

$$\langle E_{L1} \rangle = \alpha^2 - 2\alpha \left(Z - \frac{5}{16} \right)$$

Structuring the code

With analytic formulae we can speed up the computation of the correlation. In our case we write it as

$$\Psi_C = \exp \left\{ \sum_{i < j} \frac{ar_{ij}}{1 + \beta r_{ij}} \right\},$$

which means that the gradient needed for the local energy can be calculated analytically. This will speed up your code since the computation of the correlation part and the Slater determinant are the most time consuming parts in your code.

We will refer to this correlation function as Ψ_C or the *linear Padé-Jastrow*.

Structuring the code

We can test this by computing the local energy for our helium wave function

$$\psi_T(\mathbf{r}_1, \mathbf{r}_2) = \exp(-\alpha(r_1 + r_2)) \exp\left(\frac{r_{12}}{2(1 + \beta r_{12})}\right),$$

with α and β as variational parameters.

The local energy is for this case

$$E_{L2} = E_{L1} + \frac{1}{2(1 + \beta r_{12})^2} \left\{ \frac{\alpha(r_1 + r_2)}{r_{12}} \left(1 - \frac{\mathbf{r}_1 \mathbf{r}_2}{r_1 r_2}\right) - \frac{1}{2(1 + \beta r_{12})^2} - \frac{2}{r_{12}} + \right.$$

Getting a closed form expression for the local energy means that you don't need to compute derivatives numerically.

Structuring the code, simple task

- Make another copy of your code.
- Implement the closed form expression for the local energy
- Compile the new and old codes with the -pg option for profiling.
- Run both codes and profile them afterwards using `gprof <executable> > outprofile`
- Study the time usage in the file outprofile

Structuring the code, simple task

- Use also better compiler options like `c++ -O3`
- Can speed up considerably your code

Overview of week 48

Summary and exam discussion.

- Monday:
- Brief repetition from last week and discussion of projects
- Project discussions and structure of report
- Summary and exam discussion

No lecture on Tuesday. Monday is the last lecture this semester.

The report

What should it contain? A possible structure.

- An introduction where you explain the rationale for the physics case and what you have done. At the end of the introduction you should give a brief summary of the structure of the report
- Theoretical models and technicalities. This is the methods section.
- Results and discussion
- Conclusions and perspectives
- Appendix with extra material, include your program!!
- Bibliography

The report

What should I focus on? Introduction.

You don't need to answer all questions in a chronological order. When you write the introduction you could focus on the following aspects

- Motivate the reader, the first part of the introduction gives always a motivation and tries to give the overarching ideas
- What I have done
- The structure of the report, how it is organized etc

The report

What should I focus on? Methods sections.

- Describe the methods and algorithms
- You need to explain how you implemented the methods and also say something about the structure of your algorithm and present some parts of your code
- You should plug in some calculations to demonstrate your code, such as selected runs used to validate and verify your results. The latter is extremely important!! A reader needs to understand that your code reproduces selected benchmarks and reproduces previous results, either numerical and/or well-known closed form expressions.

The report

What should I focus on? Results.

- Present your results
- Give a critical discussion of your work and place it in the correct context.
- Relate your work to other calculations/studies
- An eventual reader should be able to reproduce your calculations if she/he wants to do so. All input variables should be properly explained.
- Make sure that figures and tables should contain enough information in their captions, axis labels etc so that an eventual reader can gain a first impression of your work by studying figures and tables only.

The report

What should I focus on? Conclusions.

- State your main findings and interpretations
- Try as far as possible to present perspectives for future work
- Try to discuss the pros and cons of the methods and possible improvements

The report

What should I focus on? additional material.

- Additional calculations used to validate the codes
- Selected calculations, these can be listed with few comments
- Listing of the code if you feel this is necessary You can consider moving parts of the material from the methods section to the appendix. You can also place additional material on your webpage.

The report

What should I focus on? References.

- Give always references to material you base your work on, either scientific articles/reports or books.
- *Wikipedia is not accepted as a scientific reference.* Under no circumstances.
- Refer to articles as: name(s) of author(s), journal, volume (boldfaced), page and year in parenthesis.
- Refer to books as: name(s) of author(s), title of book, publisher, place and year, eventual page numbers

Exam FYS3150/FYS4150

Pensum/syllabus.

Go to <http://www.uio.no/studier/emner/matnat/fys/FYS3150/h13/>

and click on syllabus.
For the written exam, you can bring with you 2 pages (written on both sides) with own notes. See also the exams from last years.

Exam FYS3150/FYS4150

Topics.

- Linear algebra and eigenvalue problems. (Lecture notes chapters 6.1-6.5 and 7.1-7.5 and projects 1 and 2).
- Numerical integration, standard methods and Monte Carlo methods (Lecture notes chapters 5.1-5.5 and 11 and project 5).
- Ordinary differential equations (Lecture notes chapter 8 and projects 3 and 5)
- Partial differential equations (Lecture notes chapter 10 and projects 4 and 5)
- Monte Carlo methods in physics (Lecture notes chapters 11, 12 and 14, project 5)

Exam FYS3150/FYS4150

Linear algebra and eigenvalue problems chapters 6.1-6.5 and 7.1-7.5.

- Know Gaussian elimination and LU decomposition (project 1)
- How to solve linear equations (project 1)
- How to obtain the inverse and the determinant of a real symmetric matrix
- Cubic spline
- Tridiagonal matrix decomposition (project 1)
- Householder's tridiagonalization technique and finding eigenvalues based on this
- Jacobi's method for finding eigenvalues (project 2)

Exam FYS3150/FYS4150

Numerical integration standard methods and Monte Carlo methods (5.1-5.5 and 11).

- Trapezoidal, rectangle and Simpson's rules
- Gaussian quadrature, emphasis on Legendre polynomials, but you need to know about other polynomials as well (project 5).
- Brute force Monte Carlo integration (project 5)
- Random numbers (simplest algo, ran0) and probability distribution functions, expectation values
- Improved Monte Carlo integration and importance sampling (project 5).

Exam FYS3150/FYS4150

Monte Carlo methods in physics (12 and 14).

- Random walks and Markov chains and relation with diffusion equation (project 5)
- Metropolis algorithm, detailed balance and ergodicity (project 5)
- Applications to quantum mechanical systems (project 5)

Exam FYS3150/FYS4150

Ordinary differential equations (Chapter 8).

- Euler's method and improved Euler's method, truncation errors (project 3)
- Runge Kutta methods, 2nd and 4th order, truncation errors (project 3)
- Leap-frog and Verlet algorithm (project 5)
- How to implement a second-order differential equation, both linear and non-linear. How to make your equations dimensionless.

Exam FYS3150/FYS4150

Partial differential equations chapter 10.

- Set up diffusion, Poisson and wave equations up to 2 spatial dimensions and time
- Set up the mathematical model and algorithms for these equations, with boundary and initial conditions. The stability conditions for the diffusion equation.
- Explicit, implicit and Crank-Nicolson schemes, and how to solve them. Remember that they result in triangular matrices (project 4).
- Diffusion equation in two dimensions (project 5)
- How to compute the Laplacian in Poisson's equation (project 5).
- How to solve the wave equation in one and two dimensions using an explicit scheme.

x

Other courses in Computational Science at UiO

Bachelor/Master/PhD Courses.

- INF-MAT4350 Numerical linear algebra
- MAT-INF3300/3310/4300/4310, PDEs and Sobolev spaces I and II
- INF-MAT3360 Partial differential equation
- INF3380 Parallel programming for scientific problems
- INF5620 Numerical methods for PDEs, finite element method
- FYS4411 Computational physics II, computational quantum mechanics.
- FYS4460: Computational statistical mechanics

AND, GOOD LUCK TO YOU ALL!



Version control, Git and dropbox

Why version control?

- It allows you to keep track of different change made
- It allows people you collaborate with to see the recent changes
- It becomes an excellent logbook and allows people to verify your results
- It can also be used to build up a large directory of codes, with validation examples as well Git is a very popular version control software, and easy to use. To install on ubuntu just write `sudo apt-get install git-core`.

This applies to Dropbox as well.

Version control, Git

When you want to start tracking a project just go to the project's directory and type

```
git init
```

This creates a subdirectory .git that contains all of your necessary repository files. At this point nothing in your project is tracked yet. To start, there are a few commands you need in the beginning. As an example

```
git add *.cpp *.hpp
git add ANOTHERFILE
git commit -m "This is first setup of my superduper project"
```

using

```
git status
```

Tells you about the modifications made. You can also standard unix commands like `mv`, `rm` etc

```
git mv file_from file_to
```

Version control, Git

If you want to get a copy of an existing Git repository for example, a project you would like to contribute to, the command you need is

```
git clone
```

If you are familiar with other VCS systems such as Subversion, you will notice that the command is clone and not checkout. This is an important distinction. Git receives a copy of nearly all data that the server has. Every version of every file for the history of the project is pulled down when you run `git clone`. In fact, if your server disk gets corrupted, you can use any of the clones on any client to set the server back to the state it was in when it was cloned.

```
git status
```

Optimization and profiling, useful for project 3

Till now we have not paid much attention to speed and possible optimization possibilities inherent in the various compilers. We have compiled and linked as

```
c++ -c mycode.cpp
c++ -o mycode.exe mycode.o
```

This is what we call a flat compiler option and should be used when we develop the code. It produces normally a very large and slow code when translated to machine instructions. We use this option for debugging and for establishing the correct program output because every operation is done precisely as the user specified it.

It is instructive to look up the compiler manual for further instructions

```
man c++ > out_to_file
```

Optimization and profiling

We have additional compiler options for optimization. These may include procedure inlining where performance may be improved, moving constants inside loops outside the loop, identify potential parallelism, include automatic vectorization or replace a division with a reciprocal and a multiplication if this speeds up the code.

```
c++ -O3 -c mycode.cpp
c++ -O3 -o mycode.exe mycode.o
```

This is the recommended option. **But you must check that you get the same results as previously.**

Optimization and profiling

It is also useful to profile your program under the development stage. You would then compile with (Mac and unix/linux)

```
c++ -pg -O3 -c mycode.cpp
c++ -pg -O3 -o mycode.exe mycode.o
```

After you have run the code you can obtain the profiling information via

```
gprof mycode.exe > out_to_profile
```

When you have profiled properly your code, you must take out this option as it increases your CPU expenditure.

Optimization and profiling

Other hints

- avoid if tests or call to functions inside loops, if possible.
- avoid multiplication with constants inside loops if possible.

Bad code

```
for i = 1:n a(i) = b(i) + c*d e = g(k) end
```

Better code

```
temp = c*d for i = 1:n a(i) = b(i) + temp end e = g(k)
```

What do we have then?

- Several functions under the program link, see the file OOCodes.tar.gz
- Matrix and array manipulations (similar to Blitz++/Armadillo)
- Random numbers and numerical integration
- Functions to convert arrays from C++ to Numpy and viceversa
- Numerical derivatives and differential equation solvers. These files can serve as a help when you want to start to write your own classes. We will discuss some of these classes during the course.

But what should I do else????? Some hints.

If your needs (common in most problems) include handling of large arrays and linear algebra problem, I would not recommend to write your own vector-matrix or more general array handling class. It is easy to make error.

- Old-fashioned allocation of arrays and explicit handling of all loops in for example matrix-matrix multiplication.
- You can use what we have developed for Blitz++ or the Array class (but more limited than Blitz++)
- or (recommended) you can use armadillo, a great C++ library for handling arrays and doing linear algebra.
- Armadillo provides a user friendly interface to lapack and blas functions. Here an example of using the Blas function DGEMM for matrix-matrix multiplication.
- After having installed armadillo, compile with **c++ -O3 -o test.x test.cpp -lblas -lamadillo -llapack.**

Matrix-matrix multiplication

```
#include <cstdlib>
#include <ios>
#include <iostream>
#include <armadillo>
using namespace std;
using namespace arma;

/*Because fortran files don't have any header files,
 *we need to declare the functions ourself.*/
extern "C"
{
    void dgemm_(char*, char*, int*, int*, int*, double*,
               double*, int*, double*, int*, double*, double*, int*);
}
```

Matrix-matrix multiplication

```
int main(int argc, char** argv)
{
    //Dimensions
    int n = atoi(argv[1]);
    int m = n;
    int p = m;

    /*Create random matrices
     * (note that older versions of armadillo uses "rand" instead of "randu")
     */
    srand(time(NULL));
    mat A(n, p);
    A.randu();
}
```

Matrix-matrix multiplication

```
// Pretty print, and pretty save, are as easy as the two following
// cout << A << endl;
// A.save("A.mat", raw_ascii);
mat A_trans = trans(A);
mat B(p, m);
B.randu();
mat C(n, m);
// cout << B << endl;
// B.save("B.mat", raw_ascii);
```

Matrix-matrix multiplication

```
// ARMADILLO TEST
cout << "Starting armadillo multiplication\n";
//Simple wall_clock timer is a part of armadillo.
wall_clock timer;
timer.tic();
C = A*B;
double num_sec = timer.toc();
cout << "-- Finished in " << num_sec << " seconds.\n\n";
```

Matrix-matrix multiplication

```
C = zeros<mat>(n, m);
cout << "Starting blas multiplication.\n";

char trans = 'N';
double alpha = 1.0;
double beta = 0.0;
int _numRowA = A.n_rows;
int _numColA = A.n_cols;
int _numRowB = B.n_rows;
int _numColB = B.n_cols;
int _numRowC = C.n_rows;
int _numColC = C.n_cols;
int lda = (A.n_rows >= A.n_cols) ? A.n_rows : A.n_cols;
int ldb = (B.n_rows >= B.n_cols) ? B.n_rows : B.n_cols;
int ldc = (C.n_rows >= C.n_cols) ? C.n_rows : C.n_cols;
```

Matrix-matrix multiplication, calling DGEMM

```
dgemm_(&trans, &trans, &_numRowA, &_numColB, &_numColA, &alpha,
      A.memptr(), &lda, B.memptr(), &ldb, &beta, C.memptr(), &ldc);
```

Overview of week 44

Monte Carlo methods chapter 14.

- Monday: Repetition from last week
- Statistical physics and the Metropolis algorithm (covered by chapter 13)
- Discussion of project 4.
- Wednesday:
- Statistical physics and the Metropolis algorithm (covered by chapter 13)
- Discussion of project 4.
- Introduction to differential equations
- Differential equations, general properties.
- Runge-Kutta methods Chapter 14 on quantum mechanical Monte Carlo is not part of this year's curriculum.

Most Common Ensembles in Statistical Physics

	Microcanonical	Canonical
Exchange of heat with the environment	no	yes
Exchange of particles with the environment	no	no
Thermodynamical parameters	V, \mathcal{M}, D E N	V, \mathcal{M}, D T N
Potential	Entropy	Helmholtz
Energy	Internal	Internal & Internal & Enthalpy

Microcanonical Ensemble

Entropy

$$S = k_B \ln \Omega \quad (104)$$

$$dS = \frac{1}{T} dE + \frac{p}{T} dV - \frac{\mu}{T} dN \quad (105)$$

Temperature

$$\frac{1}{k_B T} = \left(\frac{\partial \ln \Omega}{\partial E} \right)_{N, V} \quad (106)$$

Pressure

$$\frac{p}{k_B T} = \left(\frac{\partial \ln \Omega}{\partial V} \right)_{N, E} \quad (107)$$

Chemical potential

Canonical Ensemble

Helmholtz Free Energy

$$F = -k_B T \ln Z \quad (109)$$

$$dF = -SdT - pdV + \mu dN \quad (110)$$

Entropy

$$S = k_B \ln Z + k_B T \left(\frac{\partial \ln Z}{\partial T} \right)_{N, V} \quad (111)$$

Pressure

$$p = k_B T \left(\frac{\partial \ln Z}{\partial V} \right)_{N, T} \quad (112)$$

Chemical Potential

Grand Canonical Ensemble

Potential

$$pV = k_B T \ln \Xi \quad (115)$$

$$d(pV) = SdT + Nd\mu + pdV \quad (116)$$

Entropy

$$S = k_B \ln \Xi + k_B T \left(\frac{\partial \ln \Xi}{\partial T} \right)_{V, \mu} \quad (117)$$

Particles

$$N = k_B T \left(\frac{\partial \ln \Xi}{\partial \mu} \right)_{V, T} \quad (118)$$

Pressure

Pressure Canonical Ensemble

Gibbs Free Energy

$$G = -k_B T \ln \Delta \quad (120)$$

$$dG = -SdT + Vdp + \mu dN \quad (121)$$

Entropy

$$S = k_B \ln \Delta + k_B T \left(\frac{\partial \ln \Delta}{\partial T} \right)_{p, N} \quad (122)$$

Volume

$$V = -k_B T \left(\frac{\partial \ln \Delta}{\partial p} \right)_{N, T} \quad (123)$$

Chemical potential

Expectation Values

At a given temperature we have the probability distribution

$$P_i(\beta) = \frac{e^{-\beta E_i}}{Z}$$

with $\beta = 1/kT$ being the inverse temperature, k the Boltzmann constant, E_i is the energy of a state i while Z is the partition function for the canonical ensemble defined as

$$Z = \sum_{i=1}^M e^{-\beta E_i},$$

where the sum extends over all states M . P_i expresses the probability of finding the system in a given configuration i .

Expectation Values

For a system described by the canonical ensemble, the energy is an expectation value since we allow energy to be exchanged with the surroundings (a heat bath with temperature T). This expectation value, the mean energy, can be calculated using the probability distribution P_i as

$$\langle E \rangle = \sum_{i=1}^M E_i P_i(\beta) = \frac{1}{Z} \sum_{i=1}^M E_i e^{-\beta E_i},$$

with a corresponding variance defined as

$$\sigma_E^2 = \langle E^2 \rangle - \langle E \rangle^2 = \frac{1}{Z} \sum_{i=1}^M E_i^2 e^{-\beta E_i} - \left(\frac{1}{Z} \sum_{i=1}^M E_i e^{-\beta E_i} \right)^2.$$

If we divide the latter quantity with kT^2 we obtain the specific heat at constant volume

$$C_V = \frac{1}{kT^2} (\langle E^2 \rangle - \langle E \rangle^2).$$

Expectation Values

We can also write

$$\langle E \rangle = -\frac{\partial \ln Z}{\partial \beta}.$$

The specific heat is

$$C_V = \frac{1}{kT^2} \frac{\partial^2 \ln Z}{\partial \beta^2}$$

These expressions link a physical quantity (in thermodynamics) with the microphysics given by the partition function. Statistical physics is the field where one relates microscopic quantities to observables at finite temperature.

Expectation Values

$$\langle \mathcal{M} \rangle = \sum_i^M \mathcal{M}_i P_i(\beta) = \frac{1}{Z} \sum_i^M \mathcal{M}_i e^{-\beta E_i},$$

and the corresponding variance

$$\sigma_{\mathcal{M}}^2 = \langle \mathcal{M}^2 \rangle - \langle \mathcal{M} \rangle^2 = \frac{1}{Z} \sum_{i=1}^M \mathcal{M}_i^2 e^{-\beta E_i} - \left(\frac{1}{Z} \sum_{i=1}^M \mathcal{M}_i e^{-\beta E_i} \right)^2.$$

This quantity defines also the susceptibility χ

$$\chi = \frac{1}{kT} (\langle \mathcal{M}^2 \rangle - \langle \mathcal{M} \rangle^2).$$

Phase Transitions

NOTE: Helmholtz free energy and canonical ensemble

$$F = \langle E \rangle - TS = -kT \ln Z$$

meaning $\ln Z = -F/kT = -\beta F$ and

$$\langle E \rangle = -\frac{\partial \ln Z}{\partial \beta} = \frac{\partial(\beta F)}{\partial \beta}.$$

and

$$C_V = -\frac{1}{kT^2} \frac{\partial^2(\beta F)}{\partial \beta^2}.$$

We can relate observables to various derivatives of the partition function and the free energy. When a given derivative of the free energy or the partition function is discontinuous or diverges (logarithmic divergence for the heat capacity from the Ising model) we talk of a phase transition of order of the derivative.

Phase Transitions

- An important quantity is the correlation length. The correlation length defines the length scale at which the overall properties of a material start to differ from its bulk properties. It is the distance over which the fluctuations of the microscopic degrees of freedom (for example the position of atoms) are significantly correlated with each other. Usually it is of the order of few interatomic spacings for a solid.
- The correlation length ξ depends however on external conditions such as pressure and temperature.
- A phase transition is marked by abrupt macroscopic changes as external parameters are changed, such as an increase of temperature.
- The point where a phase transition takes place is called a critical point.

Two Scenarios for Phase Transitions

- First order/discontinuous phase transitions: Two or more states on either side of the critical point also coexist exactly at the critical point. As we pass through the critical point we observe a discontinuous behavior of thermodynamical functions. The correlation length is normally finite at the critical point. Phenomena such as hysteresis occur, viz. there is a continuation of state below the critical point into one above the critical point. This continuation is metastable so that the system may take a macroscopically long time to readjust. Classical example, melting of ice.
- Second order or continuous transitions: The correlation length diverges at the critical point, fluctuations are correlated over all distance scales, which forces the system to be in a unique critical phase. The two phases on either side of the critical point become identical. Smooth behavior of first derivatives of the partition function, while second derivatives diverge. Strong correlations make a perturbative treatment impossible. Renormalization group theory.

Examples of Phase Transitions

System	
Liquid-gas & Condensation/evaporation Binary liquid	m
Quantum liquid & Normal fluid/superfluid Liquid-solid & Melting/crystallisation & Reciprocal lattice vector Magnetic solid & Ferromagnetic	Antif Ar
Dielectric solid & Ferroelectric	

Ising Model

The model we will employ in our studies of phase transitions at finite temperature for magnetic systems is the so-called Ising model. In its simplest form the energy is expressed as

$$E = -J \sum_{\langle kl \rangle} s_k s_l - B \sum_k s_k,$$

with $s_k = \pm 1$, N is the total number of spins, J is a coupling constant expressing the strength of the interaction between neighboring spins and B is an external magnetic field interacting with the magnetic moment set up by the spins. The symbol $\langle kl \rangle$ indicates that we sum over nearest neighbors only.

Ising Model

Notice that for $J > 0$ it is energetically favorable for neighboring spins to be aligned. This feature leads to, at low enough temperatures, to a cooperative phenomenon called spontaneous magnetization. That is, through interactions between nearest neighbors, a given magnetic moment can influence the alignment of spins that are separated from the given spin by a macroscopic distance. These long range correlations between spins are associated with a long-range order in which the lattice has a net magnetization in the absence of a magnetic field. This phase is normally called the ferromagnetic phase. With $J < 0$, we have a so-called antiferromagnetic case. At a critical temperature we have a phase transition to a disordered phase, a so-called paramagnetic phase.

Treatment of Boundaries

With two spins, since each spin takes two values only, it means that in total we have $2^2 = 4$ possible arrangements of the two spins. These four possibilities are

$$1 = \uparrow\uparrow \quad 2 = \uparrow\downarrow \quad 3 = \downarrow\uparrow \quad 4 = \downarrow\downarrow$$

What is the energy of each of these configurations?

For small systems, the way we treat the ends matters. In the first case we employ what is called free ends. For the one-dimensional case, the energy is then written as a sum over a single index

$$E_i = -J \sum_{j=1}^{N-1} s_j s_{j+1},$$

If we label the first spin as s_1 and the second as s_2 we obtain the following expression for the energy

Treatment of Boundaries

The calculation of the energy for the one-dimensional lattice with free ends for one specific spin-configuration can easily be implemented in the following lines

```
for ( j=1; j < N; j++) {
    energy += spin[j]*spin[j+1];
}
```

where the vector $spin[]$ contains the spin value $s_k = \pm 1$. For the specific state E_1 , we have chosen all spins up. The energy of this configuration becomes then

$$E_1 = E_{\uparrow\uparrow} = -J.$$

The other configurations give

$$E_2 = E_{\uparrow\downarrow} = +J,$$

$$E_3 = E_{\downarrow\uparrow} = +J,$$

and

Treatment of Boundaries

We can also choose so-called periodic boundary conditions. This means that if $i = N$, we set the spin number to $i = 1$. In this case the energy for the one-dimensional lattice reads

$$E_i = -J \sum_{j=1}^N s_j s_{j+1},$$

and we obtain the following expression for the two-spin case

$$E = -J(s_1 s_2 + s_2 s_1).$$

If we choose to use periodic boundary conditions we can code the above expression as

```
jm=N;
for ( j=1; j <=N ; j++) {
    energy += spin[j]*spin[jm];
    jm = j ;
}
```

Treatment of Boundaries

State	Energy (FE)	Energy (PBC)	Magnetization
1 = ↑↑	-J	-2J	2
2 = ↑↓	J	2J	0
3 = ↓↑	J	2J	0
4 = ↓↓	-J	-2J	-2

Number spins up	Degeneracy	Energy (FE)	Energy (PBC)	Magne
2	1	-J	-2J	
1	2	J	2J	
0	1	-J	-2J	

Treatment of Boundaries

It is worth noting that for small dimensions of the lattice, the energy differs depending on whether we use periodic boundary conditions or free ends. This means also that the partition functions will be different, as discussed below. In the thermodynamic limit however, $N \rightarrow \infty$, the final results do not depend on the kind of boundary conditions we choose. The magnetization is however the same, defined as

$$\mathcal{M}_i = \sum_{j=1}^N s_j,$$

where we sum over all spins for a given configuration i .

Treatment of Boundaries

In a similar way, we could enumerate the number of states for a two-dimensional system consisting of two spins, i.e., a 2×2 Ising model on a square lattice with *periodic boundary conditions*. In this case we have a total of $2^4 = 16$ states. Some examples of configurations with their respective energies are listed here

$$E = -8J \quad \begin{array}{cc} \uparrow & \uparrow \\ \uparrow & \uparrow \end{array} \quad E = 0 \quad \begin{array}{cc} \uparrow & \uparrow \\ \uparrow & \downarrow \end{array} \quad E = 0 \quad \begin{array}{cc} \downarrow & \downarrow \\ \uparrow & \downarrow \end{array}$$

Treatment of Boundaries

We can group these configurations according to their total energy and magnetization.

Number spins up	Degeneracy	Energy & Magnetization
4	1	-8J 4
3	4	0 2
2	4	0 0
2	2	8J 0
1	4	0 -2
0	1	-8J -4

Modelling the Ising Model

The code uses periodic boundary conditions with energy

$$E_i = -J \sum_{j=1}^N s_j s_{j+1},$$

In our case we have as the Monte Carlo sampling function the probability for finding the system in a state s given by

$$P_s = \frac{e^{-(\beta E_s)}}{Z},$$

with energy E_s , $\beta = 1/kT$ and Z is a normalization constant which defines the partition function in the canonical ensemble

$$Z(\beta) = \sum_s e^{-(\beta E_s)}$$

This is difficult to compute since we need all states. In a calculation of the Ising model in two dimensions, the number of configurations is given by 2^N with $N = L \times L$ the number of spins

Metropolis Algorithm

- Establish an initial state with energy E_b by positioning yourself at a random position in the lattice
- Change the initial configuration by flipping e.g., one spin only. Compute the energy of this trial state E_t .
- Calculate $\Delta E = E_t - E_b$. The number of values ΔE is limited to five for the Ising model in two dimensions, see the discussion below.
- If $\Delta E \leq 0$ we accept the new configuration, meaning that the energy is lowered and we are hopefully moving towards the energy minimum at a given temperature. Go to step 7.
- If $\Delta E > 0$, calculate $w = e^{-(\beta \Delta E)}$.
- Compare w with a random number r . If $r \leq w$ then accept the new configuration, else we keep the old configuration and its values.
- The next step is to update various expectations values.
- The steps (2)-(7) are then repeated in order to obtain a sufficiently good representation of states.

Modelling the Ising Model

In the calculation of the energy difference from one spin configuration to the other, we will limit the change to the flipping of one spin only. For the Ising model in two dimensions it means that there will only be a limited set of values for ΔE . Actually, there are only five possible values. To see this, select first a random spin position x, y and assume that this spin and its nearest neighbors are all pointing up. The energy for this configuration is $E = -4J$. Now we flip this spin as shown below. The energy of the new configuration is $E = 4J$, yielding $\Delta E = 8J$.

$$E = -4J \quad \begin{array}{c} \uparrow \\ \uparrow \uparrow \uparrow \\ \uparrow \end{array} \implies E = 4J \quad \begin{array}{c} \uparrow \\ \uparrow \downarrow \uparrow \\ \uparrow \end{array}$$

Modelling the Ising Model

The four other possibilities are as follows

$$E = -2J \quad \begin{array}{c} \uparrow \\ \downarrow \uparrow \uparrow \\ \uparrow \end{array} \implies E = 2J \quad \begin{array}{c} \uparrow \\ \downarrow \downarrow \uparrow \\ \uparrow \end{array}$$

with $\Delta E = 4J$,

$$E = 0 \quad \begin{array}{c} \uparrow \\ \downarrow \uparrow \uparrow \\ \downarrow \end{array} \implies E = 0 \quad \begin{array}{c} \uparrow \\ \downarrow \downarrow \uparrow \\ \downarrow \end{array}$$

with $\Delta E = 0$

Modelling the Ising Model

$$E = 2J \quad \begin{array}{c} \downarrow \\ \downarrow \uparrow \uparrow \\ \downarrow \end{array} \implies E = -2J \quad \begin{array}{c} \downarrow \\ \downarrow \downarrow \uparrow \\ \downarrow \end{array}$$

with $\Delta E = -4J$ and finally

$$E = 4J \quad \begin{array}{c} \downarrow \\ \downarrow \uparrow \downarrow \\ \downarrow \end{array} \implies E = -4J \quad \begin{array}{c} \downarrow \\ \downarrow \downarrow \downarrow \\ \downarrow \end{array}$$

with $\Delta E = -8J$. This means in turn that we could construct an array which contains all values of $e^{\beta \Delta E}$ before doing the Metropolis sampling. Else, we would have to evaluate the exponential at each Monte Carlo sampling.

The loop over T in main

```
for ( double temp = initial_temp; temp <= final_temp; temp+=temp)
// initialise energy and magnetization
E = M = 0.;
// setup array for possible energy changes
for( int de =-8; de <= 8; de++) w[de+8] = 0;
for( int de =-8; de <= 8; de+=4) w[de+8] = exp(-de/temp);
// initialise array for expectation values
for( int i = 0; i < 5; i++) average[i] = 0.;
initialize(n_spins, temp, spin_matrix, E, M);
// start Monte Carlo computation
for (int cycles = 1; cycles <= mcs; cycles++){
    Metropolis(n_spins, idum, spin_matrix, E, M, w);
    // update expectation values
    average[0] += E;    average[1] += E*M;
    average[2] += M;    average[3] += M*M; average[4] += fabs(M);

// print results
output(n_spins, mcs, temp, average);
```

The Initialise function

```
void initialize(int n_spins, double temp, int **spin_matrix,
double& E, double& M)
{
    // setup spin matrix and initial magnetization
    for(int y = 0; y < n_spins; y++) {
        for (int x = 0; x < n_spins; x++){
            spin_matrix[y][x] = 1; // spin orientation for the ground state
            M += (double) spin_matrix[y][x];
        }
    }
    // setup initial energy
    for(int y = 0; y < n_spins; y++) {
        for (int x = 0; x < n_spins; x++){
            E -= (double) spin_matrix[y][x]*
(spin_matrix[periodic(y,n_spins,-1)][x] +
spin_matrix[y][periodic(x,n_spins,-1)]);
        }
    }
} // end function initialize
```

The periodic function

A compact way of dealing with periodic boundary conditions is given as follows:

```
// inline function for periodic boundary conditions
inline int periodic(int i, int limit, int add) {
    return (i+limit+add) % (limit);
}
```

with the following example from the function initialise

```
E -= (double) spin_matrix[y][x]*
(spin_matrix[periodic(y,n_spins,-1)][x] +
spin_matrix[y][periodic(x,n_spins,-1)]);
```

Alternative way for periodic boundary conditions

A more pedagogical way is given by the (here Fortran as example) program

```
DO y = 1,lattice_y
  DO x = 1,lattice_x
    right = x+1 ; IF(x == lattice_x ) right = 1
    left = x-1 ; IF(x == 1 ) left = lattice_x
    up = y+1 ; IF(y == lattice_y ) up = 1
    down = y-1 ; IF(y == 1 ) down = lattice_y
    energy=energy - spin_matrix(x,y)*(spin_matrix(right,y)+&
      spin_matrix(left,y)+spin_matrix(x,up)+ &
      spin_matrix(x,down) )
    magnetization = magnetization + spin_matrix(x,y)
  ENDDO
ENDDO
energy = energy*0.5
```

Computing ΔE and ΔM

The energy difference between a state E_1 and a state E_2 with zero magnetic field is

$$\Delta E = E_2 - E_1 = J \sum_{\langle kl \rangle} s_k^1 s_l^1 - J \sum_{\langle kl \rangle} s_k^2 s_l^2,$$

which we can rewrite as

$$\Delta E = -J \sum_{\langle kl \rangle} s_k^2 (s_l^2 - s_l^1),$$

where the sum now runs only over the nearest neighbors k of the spin. Since the spin to be flipped takes only two values, $s_l^1 = \pm 1$ and $s_l^2 = \pm 1$, it means that if $s_l^1 = 1$, then $s_l^2 = -1$ and if $s_l^1 = -1$, then $s_l^2 = 1$. The other spins keep their values, meaning that $s_k^1 = s_k^2$. If $s_l^1 = 1$ we must have $s_l^1 - s_l^2 = 2$, and if $s_l^1 = -1$ we must have $s_l^1 - s_l^2 = -2$. From these results we see that the energy difference can be coded efficiently as

The Metropolis function

```
// loop over all spins
for(int y=0; y < n_spins; y++) {
  for (int x= 0; x < n_spins; x++){
    int ix = (int) (ran1(&idum)*(double)n_spins); // RANDOM SPIN
    int iy = (int) (ran1(&idum)*(double)n_spins); // RANDOM SPIN
    int deltaE = 2*spin_matrix[iy][ix]*
      (spin_matrix[iy][periodic(ix,n_spins,-1)]+
      spin_matrix[periodic(iy,n_spins,-1)][ix] +
      spin_matrix[iy][periodic(ix,n_spins,1)] +
      spin_matrix[periodic(iy,n_spins,1)][ix]);
    if ( ran1(&idum) <= w[deltaE*8] ) {
      spin_matrix[iy][ix] += -1; // flip one spin and accept new spin co
      M += (double) 2*spin_matrix[iy][ix];
      E += (double) deltaE;
    }
  }
}
```

Expectation Values

```
double norm = 1/((double) (mcs)); // divided by total number of cycles
double Eaverage = average[0]*norm;
double E2average = average[1]*norm;
double Maverage = average[2]*norm;
double M2average = average[3]*norm;
double Mabsaverage = average[4]*norm;
// all expectation values are per spin, divide by 1/n_spins/n_spins
double Evariance = (E2average- Eaverage*Eaverage)/n_spins/n_spins
double Mvariance = (M2average - Mabsaverage*Mabsaverage)/n_spins/n_spins
ofile << setiosflags(ios::showpoint | ios::uppercase);
ofile << setw(15) << setprecision(8) << temp;
ofile << setw(15) << setprecision(8) << Eaverage/n_spins/n_spins;
ofile << setw(15) << setprecision(8) << Evariance/temp/temp;
ofile << setw(15) << setprecision(8) << Maverage/n_spins/n_spins;
ofile << setw(15) << setprecision(8) << Mvariance/temp;
ofile << setw(15) << setprecision(8) << Mabsaverage/n_spins/n_spins;
```

Code example for the parallel two-dimensional diff equation

```
int Jacobi_P(int mynode, int numnodes, int N, double **A, double *x)
{
  int i,j,k,i_global;
  int maxit = 100000;
  int rows_local,local_offset,last_rows_local,*count,*displacements;
  double sum1,sum2,*xold;
  double error_sum_local, error_sum_global;
  MPI_Status status;

  rows_local = (int) floor((double)N/numnodes);
  local_offset = mynode*rows_local;
  if(mynode == (numnodes-1))
    rows_local = N - rows_local*(numnodes-1);
```

Code example for the parallel two-dimensional diff equation

```

/*Distribute the Matrix and R.H.S. among the processors */
if(mynode == 0){
    for(i=1;i<numnodes-1;i++){
        for(j=0;j<rows_local;j++){
            MPI_Send(A[i*rows_local+j],N,MPI_DOUBLE,i,j,MPI_COMM_WORLD);
            MPI_Send(b+i*rows_local,rows_local,MPI_DOUBLE,i,rows_local,
                MPI_COMM_WORLD);

            last_rows_local = N-rows_local*(numnodes-1);
            for(j=0;j<last_rows_local;j++){
                MPI_Send(A[(numnodes-1)*rows_local+j],N,MPI_DOUBLE,numnodes-1,
                    MPI_COMM_WORLD);
                MPI_Send(b+(numnodes-1)*rows_local,last_rows_local,MPI_DOUBLE,numnodes-1,
                    last_rows_local,MPI_COMM_WORLD);
            }
        }
    }
}

```

Code example for the parallel two-dimensional diff equation

```

else{
    A = CreateMatrix(rows_local,N);
    x = new double[rows_local];
    b = new double[rows_local];
    for(i=0;i<rows_local;i++){
        MPI_Recv(A[i],N,MPI_DOUBLE,0,i,MPI_COMM_WORLD,&status);
        MPI_Recv(b,rows_local,MPI_DOUBLE,0,rows_local,MPI_COMM_WORLD,&status);
    }
}

```

Code example for the parallel two-dimensional diff equation

```

xold = new double[N];
count = new int[numnodes];
displacements = new int[numnodes];
//set initial guess to all 1.0
for(i=0; i<N; i++){
    xold[i] = 1.0;

    for(i=0;i<numnodes;i++){
        count[i] = (int) floor((double)N/numnodes);
        displacements[i] = i*count[i];

        count[numnodes-1] = N - ((int)floor((double)N/numnodes))*numnodes;
    }
}

```

Code example for the parallel two-dimensional diff equation

```

for(k=0; k<maxit; k++){
    error_sum_local = 0.0;
    for(i = 0; i<rows_local; i++){
        i_global = local_offset+i;
        sum1 = 0.0; sum2 = 0.0;
        for(j=0; j < i_global; j++){
            sum1 = sum1 + A[i][j]*xold[j];
            for(j=i_global+1; j < N; j++){
                sum2 = sum2 + A[i][j]*xold[j];
            }

            x[i] = (-sum1 - sum2 + b[i])/A[i][i_global];
            error_sum_local += (x[i]-xold[i_global])*(x[i]-xold[i_global]);
        }
    }
}

```

Code example for the parallel two-dimensional diff equation

```

MPI_Allreduce(&error_sum_local,&error_sum_global,1,MPI_DOUBLE,
    MPI_SUM,MPI_COMM_WORLD);
MPI_Allgather(x,rows_local,MPI_DOUBLE,xold,count,displacements,
    MPI_DOUBLE,MPI_COMM_WORLD);

```

Code example for the parallel two-dimensional diff equation

```

if(sqrt(error_sum_global)<abstol){
    if(mynode == 0){
        for(i=0;i<N;i++){
            x[i] = xold[i];
        }
    }
    else{
        DestroyMatrix(A,rows_local,N);
        delete[] x;
        delete[] b;

        delete[] xold;
        delete[] count;
        delete[] displacements;
        return k;
    }
}

```

Code example for the parallel two-dimensional diff equation

```
cerr << "Jacobi: Maximum Number of Iterations Reached Without Co
if(mynode == 0){
    for(i=0;i<N;i++)
        x[i] = xold[i];
    else{
        DestroyMatrix(A,rows_local,N);
        delete[] x;
        delete[] b;

        delete[] xold;
        delete[] count;
        delete[] displacements;

    return maxit;
```

How do I use the titan.uio.no cluster?

- Computational Physics requires High Performance Computing (HPC) resources
- USIT and the Research Computing Services (RCS) provides HPC resources and HPC support
- Resources: titan.uio.no
- Support: 14 people
- Contact: hpl@usit.uio.no

Titan

Hardware.

- 304 dual-cpu quad-core SUN X2200 Opteron nodes (total 2432 cores), 2.2 Ghz, and 8 - 16 GB RAM and 250 - 1000 GB disk on each node
- 3 eight-cpu quad-core Sun X4600 AMD Opteron nodes (total 96 cores), 2.5 Ghz, and 128, 128 and 256 GB memory, respectively
- Infiniband interconnect
- Heterogenous cluster!

Titan

Software.

- Batch system: SLURM and MAUI
- Message Passing Interface (MPI):
- OpenMPI
- Scampi
- MPICH2
- Compilers: GCC, Intel, Portland and Pathscale
- Optimized math libraries and scientific applications
- All you need may be found under /site
- Available software:
http://www.hpc.uio.no/index.php/Titan_software

Getting started

Batch systems.

- A batch system controls the use of the cluster resources
- Submits the job to the right resource
- Monitors the job while executing
- Restarts the job in case of failure
- Takes care of priorities and queues to control execution order of unrelated jobs

Sun Grid Engine.

- SGE is the batch system used on Titan
- Jobs are executed either interactively or through job scripts
- Useful commands: showq, qlogin, sbatch
- http://hpc.uio.no/index.php/Titan_User_Guide

Getting started

Modules.

- Different compilers, MPI-versions and applications need different sets of user environment variables
- The modules package lets you load and remove the different variable sets
- Useful commands:
- List available modules: module avail
- Load module: module load <environment>
- Unload module: module unload <environment>
- Currently loaded: module list
- http://hpc.uio.no/index.php/Titan_User_Guide

Example

Interactively.

The job script

job.sge.

Example

Submitting.

Example

Checking execution.

Tips and admonitions

Tips.

- Titan FAQ: <http://www.hpc.uio.no/index.php/FAQ>
- man-pages, e.g. `man sbatch`
- Ask us

Admonitions.

- Remember to exit from qlogin-sessions; the resource is reserved for you until you exit
- Don't run jobs on login-nodes; these are only for compiling and editing files

Projects: quantum mechanics

The expectation value of the kinetic energy expressed in atomic units for electron i is

$$K_i = -\frac{1}{2} \frac{\nabla_i^2 \Psi}{\Psi}.$$

$$\frac{\nabla^2 \Psi}{\Psi} = \frac{\nabla^2 \Psi_D}{\Psi_D} + \frac{\nabla^2 \Psi_J}{\Psi_J} + 2 \frac{\nabla \Psi_D}{\Psi_D} \cdot \frac{\nabla \Psi_J}{\Psi_J} \quad (125)$$

Quantum mechanical project

We define the correlated function as

$$\Psi_J = \prod_{i<j} g(r_{ij}) = \prod_{i<j} g(r_{ij}) = \prod_{i=1}^N \prod_{j=i+1}^N g(r_{ij}),$$

with $r_{ij} = |\mathbf{r}_i - \mathbf{r}_j| = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2 + (z_i - z_j)^2}$ for three dimensions and $r_{ij} = |\mathbf{r}_i - \mathbf{r}_j| = \sqrt{(x_i - x_j)^2 + (y_i - y_j)^2}$ for two dimensions.

In our particular case we have

$$\Psi_J = \prod_{i<j} g(r_{ij}) = \exp \left\{ \sum_{i<j} f(r_{ij}) \right\} = \exp \left\{ \sum_{i<j} \frac{a r_{ij}}{1 + \beta r_{ij}} \right\},$$

Quantum mechanical project

The second derivative of the Jastrow factor divided by the Jastrow factor (the way it enters the kinetic energy) is

$$\left[\frac{\nabla^2 \Psi_J}{\Psi_J} \right]_x = 2 \sum_{k=1}^N \sum_{i=1}^{k-1} \frac{\partial^2 g_{ik}}{\partial x_k^2} + \sum_{k=1}^N \left(\sum_{i=1}^{k-1} \frac{\partial g_{ik}}{\partial x_k} - \sum_{i=k+1}^N \frac{\partial g_{ki}}{\partial x_i} \right)^2$$

But we have a simple form for the function, namely

$$\Psi_J = \prod_{i<j} \exp f(r_{ij}) = \exp \left\{ \sum_{i<j} \frac{a r_{ij}}{1 + \beta r_{ij}} \right\},$$

and it is easy to see that for particle k we have

$$\frac{\nabla_k^2 \Psi_J}{\Psi_J} = \sum_{ij \neq k} \frac{(\mathbf{r}_k - \mathbf{r}_i)(\mathbf{r}_k - \mathbf{r}_j)}{r_{ki} r_{kj}} f'(r_{ki}) f'(r_{kj}) + \sum_{j \neq k} \left(f''(r_{kj}) + \frac{2}{r_{kj}} f'(r_{kj}) \right)$$

Jastrow gradient

We have

$$\Psi_J = \prod_{i<j} g(r_{ij}) = \exp \left\{ \sum_{i<j} \frac{a r_{ij}}{1 + \beta r_{ij}} \right\},$$

the gradient needed for the local energy is easy to compute. We get for particle k

$$\frac{\nabla_k \Psi_J}{\Psi_J} = \sum_{j \neq k} \frac{\mathbf{r}_{kj}}{r_{kj}} \frac{a}{(1 + \beta r_{kj})^2},$$

which is rather easy to code. Remember to sum over all particles when you compute the local energy.

Two-dimensional wave equation, Tsunami model

We assume that we can approximate the coastline with a quadratic grid. As boundary condition at the coastline we will employ

$$\frac{\partial u}{\partial n} = \nabla u \cdot \mathbf{n} = 0,$$

where $\partial u / \partial n$ is the derivative in the direction normal to the boundary.

Here you must pay particular attention to the endpoints.

Two-dimensional wave equation

We are going to model the impact of an earthquake on sea water. This is normally modelled via an elevation of the sea bottom. We will assume that the movement of the sea bottom is very rapid compared with the period of the propagating waves. This means that we can approximate the bottom elevation with an initial surface elevation. The initial conditions are then given by (with L the length of the grid)

$$u(x, y, 0) = f(x, y) \quad x, y \in (0, L),$$

and

$$\partial u / \partial t|_{t=0} = 0 \quad x, y \in (0, L).$$

We will approximate the initial elevation with the function

$$f(x, y) = A_0 \exp \left(- \left[\frac{x - x_c}{\sigma_x} \right]^2 - \left[\frac{y - y_c}{\sigma_y} \right]^2 \right),$$

Exam FYS3150 place FV329, Lab

Place and duration.

- 13/12-17/12.
- duration: ~ 45 min
- ca 20-25 min for discussion of the project, your presentation, reproduction of results, test runs etc
- 20-25 min for questions from one of the five topics listed below.
- Check final exam list by the end of this week, see webpage. Changes must be communicated to me at mhjensen@fys.uio.no.

Exam FYS3150

How to prepare your Talk.

- You can use slides, ps, pdf, ppt etc files (projector and slide projector at room FV329). Latex example on webpage. Choose among one of the two versions of project 5.
- 10-15 mins (5-15 slides) for your presentation, rest of 10-15 mins questions and test of code from classfronter You should discuss
- The mathematical model and the physics
- Your algorithm and how you implemented it, with selected results
- How you dealt with eventual comments and your corrections of these
- Any improvement you can think of

Week 44, October 26-30

Ising model phase transitions and project 4.

- Monday: Repetition from last week
- Discussion of project 4, results near the phase transition
- Definition of correlation time and equilibration time (optional part of project 4).
- Wednesday:
- End Monte Carlo discussion and summary of chapters 8-10. Chapter 11 on quantum Monte Carlo will not be discussed and is not part of the curriculum.

Mean Field Theory and the Ising Model

In studies of phase transitions we are interested in minimizing the free energy by varying the average magnetization, which is the order parameter (disappears at T_C).

In mean field theory the local magnetization is treated as a constant, all effects from fluctuations are neglected. A way to achieve this is to rewrite by adding and subtracting the mean magnetization $\langle s \rangle$

$$s_i s_j = (s_i - \langle s \rangle + \langle s \rangle)(s_j - \langle s \rangle + \langle s \rangle) \approx \langle s \rangle^2 + \langle s \rangle(s_i - \langle s \rangle) + \langle s \rangle(s_j - \langle s \rangle),$$

where we have ignored terms of the order $(s_i - \langle s \rangle)(s_j - \langle s \rangle)$, which leads to correlations between neighbouring spins. In mean field theory we ignore correlations.

Mean Field Theory and the Ising Model

This means that we can rewrite the Hamiltonian

$$E = -J \sum_{\langle ij \rangle} s_i s_j - B \sum_i s_i,$$

as

$$E = -J \sum_{\langle ij \rangle} (\langle s \rangle^2 + \langle s \rangle(s_i - \langle s \rangle) + \langle s \rangle(s_j - \langle s \rangle)) - B \sum_i s_i,$$

resulting in

$$E = -(B + zJ\langle s \rangle) \sum_i s_i + zJ\langle s \rangle^2,$$

with z the number of nearest neighbours for a given site i . We can define an effective field which all spins see, namely

$$B_{\text{eff}} = (B + zJ\langle s \rangle).$$

Mean Field Theory and the Ising Model

How do we get $\langle s \rangle$?

Here we use the canonical ensemble. The partition function reads in this case

$$Z = e^{-NzJ\langle s \rangle^2/kT} (2 \cosh(B_{\text{eff}}/kT))^N,$$

with a free energy

$$F = -kT \ln Z = -NkT \ln(2) + NzJ\langle s \rangle^2 - NkT \ln(\cosh(B_{\text{eff}}/kT))$$

and minimizing F wrt $\langle s \rangle$ we arrive at

$$\langle s \rangle = \tanh(2 \cosh(B_{\text{eff}}/kT)).$$

Connecting to Landau Theory

Close to the phase transition we expect $\langle s \rangle$ to become small and eventually vanish. We can then expand F in powers of $\langle s \rangle$ as

$$F = -NkT \ln(2) + NzJ\langle s \rangle^2 - NkT \ln(\cosh(B_{\text{eff}}/kT))$$

and using $\langle M \rangle = N\langle s \rangle$ we can rewrite as

$$F = F_0 - B\langle M \rangle + \frac{1}{2}a\langle M \rangle^2 + \frac{1}{4}b\langle M \rangle^4 + \dots$$

Connecting to Landau Theory

Let $\langle M \rangle = m$ and

$$F = F_0 + \frac{1}{2}am^2 + \frac{1}{4}bm^4 + \frac{1}{6}cm^6$$

F has a minimum at equilibrium $F'(m) = 0$ and $F''(m) > 0$

$$F'(m) = 0 = m(a + bm^2 + cm^4),$$

and if we assume that m is real we have two solutions

$$m = 0,$$

or

$$m^2 = \frac{b}{2c} \left(-1 \pm \sqrt{1 - 4ac/b^2} \right)$$

Second Order Phase Transition

Can describe both first and second-order phase transitions. Here we consider the second case. Assume $b > 0$ and $a \ll 1$ small since we want to study a perturbation around $m = 0$. We reach the critical point when $a = 0$.

$$m^2 = \frac{b}{2c} \left((-1 \pm \sqrt{1 - 4ac/b^2}) \right) \approx -a/b$$

Assume that

$$a(T) = \alpha(T - T_C),$$

with $\alpha > 0$ and T_C the critical temperature where the magnetization vanishes. If a is negative we have two solutions

$$m = \pm \sqrt{-a/b} = \pm \sqrt{\frac{\alpha(T_C - T)}{b}}$$

m evolves continuously to the critical temperature where $F = 0$ for $T \leq T_C$ (see separate graph).

Entropy and Specific Heat

We can now compute the entropy

$$S = - \left(\frac{\partial F}{\partial T} \right)$$

For $T \geq T_C$ we have $m = 0$ and

$$S = - \left(\frac{\partial F_0}{\partial T} \right)$$

and for $T \leq T_C$

$$S = - \left(\frac{\partial F_0}{\partial T} \right) - \alpha^2(T_C - T)/2b,$$

and we see that there is a smooth crossover at T_C .

Entropy and Specific Heat

We can now compute the specific heat

$$C_V = T \left(\frac{\partial S}{\partial T} \right)$$

and T_C we get a discontinuity of

$$\Delta C_V = -\alpha^2/2b,$$

signalling a second-order phase transition. Landau theory gives irrespective of dimension critical exponents

$$m \sim (T_C - T)^\beta,$$

and

$$C_V \sim (T_C - T)^\alpha,$$

with $\beta = 1/2$ and $\alpha = 1$. It predicts a phase transition for one dimension as well. For the Ising model there is no phase transition

Scaling Results

Near T_C we can characterize the behavior of many physical quantities by a power law behavior. As an example, the mean magnetization is given by

$$\langle M(T) \rangle \sim (T - T_C)^\beta, \quad (126)$$

where β is a so-called critical exponent. A similar relation applies to the heat capacity

$$C_V(T) \sim |T_C - T|^{-\alpha}, \quad (127)$$

the susceptibility

$$\chi(T) \sim |T_C - T|^{-\gamma}. \quad (128)$$

and the correlation length

$$\xi(T) \sim |T_C - T|^{-\nu}. \quad (129)$$

$\alpha = 0, \beta = 1/2, \gamma = 7/4$ and $\nu = 1$. Below we will derive these

Scaling Results

Through finite size scaling relations it is possible to relate the behavior at finite lattices with the results for an infinitely large lattice. The critical temperature scales then as

$$T_C(L) - T_C(L = \infty) \sim aL^{-1/\nu}, \quad (130)$$

$$\langle M(T) \rangle \sim (T - T_C)^\beta \rightarrow L^{-\beta/\nu}, \quad (131)$$

$$C_V(T) \sim |T_C - T|^{-\gamma} \rightarrow L^{\gamma/\nu}, \quad (132)$$

and

$$\chi(T) \sim |T_C - T|^{-\alpha} \rightarrow L^{\alpha/\nu}. \quad (133)$$

We can compute the slope of the curves for M , C_V and χ as function of lattice sites and extract the exponent ν .

Analytic Results: two-dimensional Ising model

The analytic expression for the Ising model in two dimensions was obtained in 1944 by the Norwegian chemist Lars Onsager (Nobel prize in chemistry). The exact partition function for N spins in two dimensions and with zero magnetic field \mathcal{H} is given by

$$Z_N = \left[2 \cosh(\beta J) e^I \right]^N,$$

with

$$I = \frac{1}{2\pi} \int_0^\pi d\phi \ln \left[\frac{1}{2} \left(1 + (1 - \kappa^2 \sin^2 \phi)^{1/2} \right) \right],$$

and

$$\kappa = 2 \sinh(2\beta J) / \cosh^2(2\beta J).$$

Analytic Results: two-dimensional Ising model

The resulting energy is given by

$$\langle E \rangle = -J \coth(2\beta J) \left[1 + \frac{2}{\pi} (2 \tanh^2(2\beta J) - 1) K_1(q) \right],$$

with $q = 2 \sinh(2\beta J) / \cosh^2(2\beta J)$ and the complete elliptic integral of the first kind

$$K_1(q) = \int_0^{\pi/2} \frac{d\phi}{\sqrt{1 - q^2 \sin^2 \phi}}.$$

Analytic Results: two-dimensional Ising model

Differentiating once more with respect to temperature we obtain the specific heat given by

$$C_V = \frac{4k_B}{\pi} (\beta J \coth(2\beta J))^2$$

$$\left\{ K_1(q) - K_2(q) - (1 - \tanh^2(2\beta J)) \left[\frac{\pi}{2} + (2 \tanh^2(2\beta J) - 1) K_1(q) \right] \right\},$$

with

$$K_2(q) = \int_0^{\pi/2} d\phi \sqrt{1 - q^2 \sin^2 \phi}.$$

is the complete elliptic integral of the second kind. Near the critical temperature T_C the specific heat behaves as

$$C_V \approx -\frac{2}{k_B \pi} \left(\frac{J}{T_C} \right)^2 \ln \left| 1 - \frac{T}{T_C} \right| + \text{const.}$$

Analytic Results: two-dimensional Ising model

In theories of critical phenomena one has that

$$C_V \sim \left| 1 - \frac{T}{T_C} \right|^{-\alpha},$$

and Onsager's result is a special case of this power law behavior. The limiting form of the function

$$\lim_{\alpha \rightarrow 0} \frac{1}{\alpha} (Y^{-\alpha} - 1) = -\ln Y,$$

meaning that the analytic result is a special case of the power law singularity with $\alpha = 0$.

Analytic Results: two-dimensional Ising model

One can also show that the mean magnetization per spin is

$$\left[1 - \frac{(1 - \tanh^2(\beta J))^4}{16 \tanh^4(\beta J)} \right]^{1/8}$$

for $T < T_C$ and 0 for $T > T_C$. The behavior is thus as $T \rightarrow T_C$ from below

$$M(T) \sim (T_C - T)^{1/8}$$

The susceptibility behaves as

$$\chi(T) \sim |T_C - T|^{-7/4}$$

Time Auto-correlation Function

The so-called time-displacement autocorrelation $\phi(t)$ for the magnetization is given by

$$\phi(t) = \int dt' [\mathcal{M}(t') - \langle \mathcal{M} \rangle] [\mathcal{M}(t' + t) - \langle \mathcal{M} \rangle],$$

which can be rewritten as

$$\phi(t) = \int dt' [\mathcal{M}(t') \mathcal{M}(t' + t) - \langle \mathcal{M} \rangle^2],$$

where $\langle \mathcal{M} \rangle$ is the average value of the magnetization and $\mathcal{M}(t)$ its instantaneous value. We can discretize this function as follows, where we used our set of computed values $\mathcal{M}(t)$ for a set of discretized times (our Monte Carlo cycles corresponding to a sweep over the lattice)

$$\phi(t) = \frac{1}{t_{\max} - t} \sum_{t'=0}^{t_{\max}-t} \mathcal{M}(t') \mathcal{M}(t'+t) - \frac{1}{t_{\max} - t} \sum_{t'=0}^{t_{\max}-t} \mathcal{M}(t') \times \frac{1}{t_{\max}}$$

Correlation Time, how to compute

```
// define m-value at each cycle within loop over cycles
m_matrix[cycles] = Maverage/(n_spins**2)/cycles
// Then compute phi(i) as (in pseudocode)
for i = 1, mcs
  r = 1.0/(mcs-1)
  s = 0.0; v = 0.0; p = 0.0
  for j = 1, mcs-1
    p = p + m_matrix(j)*m_matrix(j+i)
    s = s + m_matrix(j)
    v = v + m_matrix(j+i)
  end for
  phi(i) = r*p-r*s+v
end for
```

Time Auto-correlation Function

One should be careful with times close to t_{\max} , the upper limit of the sums becomes small and we end up integrating over a rather small time interval. This means that the statistical error in $\phi(t)$ due to the random nature of the fluctuations in $\mathcal{M}(t)$ can become large.

One should therefore choose $t \ll t_{\max}$.

Note also that we could replace the magnetization with the mean energy, or any other expectation values of interest.

The time-correlation function for the magnetization gives a measure of the correlation between the magnetization at a time t' and a time $t' + t$. If we multiply the magnetizations at these two different times, we will get a positive contribution if the magnetizations are fluctuating in the same direction, or a negative value if they fluctuate in the opposite direction. If we then integrate over time, or use the discretized version of, the time correlation function $\phi(t)$ should take a non-zero value if the fluctuations are correlated, else it should gradually go to zero. For

Time Auto-correlation Function

We can derive the correlation time by observing that our Metropolis algorithm is based on a random walk in the space of all possible spin configurations. Our probability distribution function $\hat{\mathbf{w}}(t)$ after a given number of time steps t could be written as

$$\hat{\mathbf{w}}(t) = \hat{\mathbf{W}}^t \hat{\mathbf{w}}(0),$$

with $\hat{\mathbf{w}}(0)$ the distribution at $t = 0$ and $\hat{\mathbf{W}}$ representing the transition probability matrix. We can always expand $\hat{\mathbf{w}}(0)$ in terms of the right eigenvectors of $\hat{\mathbf{W}}$ as

$$\hat{\mathbf{w}}(0) = \sum_i \alpha_i \hat{\mathbf{v}}_i,$$

resulting in

$$\hat{\mathbf{w}}(t) = \hat{\mathbf{W}}^t \hat{\mathbf{w}}(0) = \hat{\mathbf{W}}^t \sum_i \alpha_i \hat{\mathbf{v}}_i = \sum_i \lambda_i^t \alpha_i \hat{\mathbf{v}}_i,$$

with λ_i the i^{th} eigenvalue corresponding to the eigenvector $\hat{\mathbf{v}}_i$.

Time Auto-correlation Function

If we assume that λ_0 is the largest eigenvalue we see that in the limit $t \rightarrow \infty$, $\hat{\mathbf{w}}(t)$ becomes proportional to the corresponding eigenvector $\hat{\mathbf{v}}_0$. This is our steady state or final distribution.

We can relate this property to an observable like the mean magnetization. With the probability $\hat{\mathbf{w}}(t)$ (which in our case is the Boltzmann distribution) we can write the mean magnetization as

$$\langle \mathcal{M}(t) \rangle = \sum_{\mu} \hat{\mathbf{w}}(t)_{\mu} \mathcal{M}_{\mu},$$

or as the scalar of a vector product

$$\langle \mathcal{M}(t) \rangle = \hat{\mathbf{w}}(t) \mathbf{m},$$

with \mathbf{m} being the vector whose elements are the values of \mathcal{M}_{μ} in its various microstates μ .

Time Auto-correlation Function

We rewrite this relation as

$$\langle \mathcal{M}(t) \rangle = \hat{\mathbf{w}}(t) \mathbf{m} = \sum_i \lambda_i^t \alpha_i \hat{\mathbf{v}}_i \mathbf{m}_i.$$

If we define $m_i = \hat{\mathbf{v}}_i \mathbf{m}_i$ as the expectation value of \mathcal{M} in the i^{th} eigenstate we can rewrite the last equation as

$$\langle \mathcal{M}(t) \rangle = \sum_i \lambda_i^t \alpha_i m_i.$$

Since we have that in the limit $t \rightarrow \infty$ the mean magnetization is dominated by the the largest eigenvalue λ_0 , we can rewrite the last equation as

$$\langle \mathcal{M}(t) \rangle = \langle \mathcal{M}(\infty) \rangle + \sum_{i \neq 0} \lambda_i^t \alpha_i m_i.$$

We define the quantity

Time Auto-correlation Function

The quantities τ_i are the correlation times for the system. They control also the auto-correlation function discussed above. The longest correlation time is obviously given by the second largest eigenvalue τ_1 , which normally defines the correlation time discussed above. For large times, this is the only correlation time that survives. If higher eigenvalues of the transition matrix are well separated from λ_1 and we simulate long enough, τ_1 may well define the correlation time. In other cases we may not be able to extract a reliable result for τ_1 . Coming back to the time correlation function $\phi(t)$ we can present a more general definition in terms of the mean magnetizations $\langle \mathcal{M}(t) \rangle$. Recalling that the mean value is equal to $\langle \mathcal{M}(\infty) \rangle$ we arrive at the expectation values

$$\phi(t) = \langle \mathcal{M}(0) - \mathcal{M}(\infty) \rangle \langle \mathcal{M}(t) - \mathcal{M}(\infty) \rangle,$$

resulting in

$$\phi(t) = \sum m_i \alpha_i m_j \alpha_j e^{-t/\tau_i},$$

Correlation Time

If the correlation function decays exponentially

$$\phi(t) \sim \exp(-t/\tau)$$

then the exponential correlation time can be computed as the average

$$\tau_{\text{exp}} = -\left\langle \frac{t}{\log \left| \frac{\phi(t)}{\phi(0)} \right|} \right\rangle.$$

If the decay is exponential, then

$$\int_0^\infty dt \phi(t) = \int_0^\infty dt \phi(0) \exp(-t/\tau) = \tau \phi(0),$$

which suggests another measure of correlation

$$\tau_{\text{int}} = \sum_k \frac{\phi(k)}{\phi(0)},$$

Time Auto-correlation Function

Here we mention that one can show, using scaling relations, that at the critical temperature the correlation time τ relates to the lattice size L as

$$\tau \sim L^{d+z},$$

with d the dimensionality of the system. For the Metropolis algorithm based on a single spin-flip process, Nightingale and Blöte obtained $z = 2.1665 \pm 0.0012$. This is a rather high value, meaning that our algorithm is not the best choice when studying properties of the Ising model near T_C .

We can understand this behavior by studying the development of the two-dimensional Ising model as function of temperature.

Time Auto-correlation Function

Cooling the system down to the critical temperature we observe clusters pervading larger areas of the lattice. The reason for the large correlation time (and the parameter z) for the single-spin flip Metropolis algorithm is the development of these large domains or clusters with all spins pointing in one direction. It is quite difficult for the algorithm to flip over one of these large domains because it has to do it spin by spin, with each move having a high probability of being rejected due to the ferromagnetic interaction between spins. Since all spins point in the same direction, the chance of performing the flip

$$E = -4J \begin{array}{c} \uparrow \\ \uparrow \uparrow \uparrow \\ \uparrow \end{array} \implies E = 4J \begin{array}{c} \uparrow \\ \uparrow \downarrow \uparrow \\ \uparrow \end{array}$$

leads to an energy difference of $\Delta E = 8J$. Using the exact critical temperature $k_B T_C/J \approx 2.269$, we obtain a probability $\exp-(8/2.269) = 0.029429$ which is rather small. The increase in large correlation times due to increasing lattices can be diminished

Better Algorithm needed

Monte Carlo simulations close to a phase transition are affected by critical slowing down. In the 2-D Ising system, the correlation length becomes very large, and the correlation time, which measures the number of steps between independent Monte Carlo configurations behaves like

$$\tau \sim \xi^z,$$

with $z \approx 2.1$ for the Metropolis algorithm. The exponent z is called the dynamic critical exponent. The maximum possible value for ξ in a finite system of $N = L \times L$ spins is $\xi \sim L$, because ξ cannot be larger than the lattice size! This implies that $\tau \sim L^z \approx N$. This makes simulations difficult because the Metropolis algorithm time scales like N , so the time to generate independent Metropolis configurations scales like

$$N\tau \sim N^2 = L^4.$$

If the lattice size

Better Algorithm needed

There is a simple physical argument which helps understand why $z = 2$. The Metropolis algorithm is a local algorithm, i.e., one spin is tested and flipped at a time. Near T_C the system develops large domains of correlated spins which are difficult to break up. So the most likely change in configuration is the movement of a whole domain of spins. But one Metropolis sweep of the lattice can move a domain at most by approximately one lattice spacing in each time step. This motion is stochastic, i.e., like a random walk. The distance traveled in a random walk scales like $\sqrt{\text{time}}$, so to move a domain a distance of order ξ takes $\tau \sim \xi^2$ Monte Carlo steps. This argument suggests that the way to speed up a Monte Carlo simulation near T_C is to use a non-local algorithm.