



# ML and Optimization Intro

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# Quiz

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- What is minimum value of the following function?

$$f(x) = x^2 - 2x + 2$$

$$f(x) = 3 \sin(x) \cos(x) (6x^2 + 3x^3 + x) \tan(x)$$

- How to get the answer?
  - Mathematically
  - Geographically

# What is optimization?

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- Optimization means ‘minimization’ or ‘maximization’ of a (objective) function

$$\min f(x)$$

# Optimization form

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- Unconstrained optimization

- $\arg \max$

$$\operatorname{argmax}_x f(x)$$

- Constrained Optimization

$$\text{minimize } f(x)$$

$$\text{subject to } c_j(x) = 0$$

$$c_k(x) \geq 0$$

- Black-box optimization

- Objective function is unknown

$$y = f(\lambda)$$

# How to solve optimization problems

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- Mathematical methods
  - Simplex
  - Lagrangian method
- Geographically
- Meta Heuristics
- Heuristics
- Differentiation
- Gradient descent
- Optimization by search

# 1. Mathematical methods for constrained optimization

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- Linear Programming
- Integer Programming

# Linear Programming

## ■ Linear Programming

- Object function and constraints are all linear
  - Min. (Max.) a function under constraints.
    - ✓ Material Mixing in production system
    - ✓ Job allocation
    - ✓ Transportation ....

A,B 두 상품을 생산하는데 상품 A는 개당 2원의 이익이 나고, B는 개당 5원의 이익이 발생한다. 상품 A를 생산하는 데 9개의 재료와 3시간 동안 기계를 사용해야 하며, B는 5개의 재료와 4시간의 기계를 사용해야 한다. 이때 재료는 총 300개를 사용할 수 있으며, 기계 가동 시간은 최대 200시간이라고 한다. 또 상품 A는 최소 5개 이상을 생산해야만 한다고 한다. 이때 최대의 이익을 산출해 내는 상품 A와 B의 생산량을 결정하라.

결정 변수: 제품 A의 생산량  $\Rightarrow x_1$   
제품 B의 생산량  $\Rightarrow x_2$

목적함수 제약식	{	$Maximize \quad 2x_1 + 5x_2$	← 이익의 최대화
		$s.t. \quad 3x_1 + 4x_2 \leq 200$	← 기계 가동시간 제약
		$9x_1 + 5x_2 \leq 300$	← 재료 사용량 제약
		$x_1 \geq 5$	← A의 최소 생산량 제약
		$x_2 \geq 0$	← 비음인 해만을 구함

- Minimax 문제

Minimize maximum $\{12x_1 - 21x_2, 17x_1 - 10x_2\}$

$$\begin{aligned} \text{s.t.} \quad & 2x_1 - 7x_2 \geq 12 \\ & 6x_1 + 11x_2 \geq 41 \\ & 9x_1 + 17x_2 \leq 102 \\ & x_i \geq 0, \quad \forall i \end{aligned}$$



Minimize Z

$$\begin{aligned} \text{s.t.} \quad & Z \geq 2x_1 - 21x_2 \\ & Z \geq 17x_1 - 10x_2 \\ & 2x_1 - 7x_2 \geq 12 \\ & 6x_1 + 11x_2 \geq 41 \\ & 9x_1 + 17x_2 \leq 102 \\ & x_i \geq 0, \quad \forall i \end{aligned}$$

- Minimum Absolute Value
- Goal Programming
- ....



# Convex optimization

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- In the constrained optimization,

$$\text{minimize } f(x)$$

$$\text{subject to } c_j(x) = 0$$

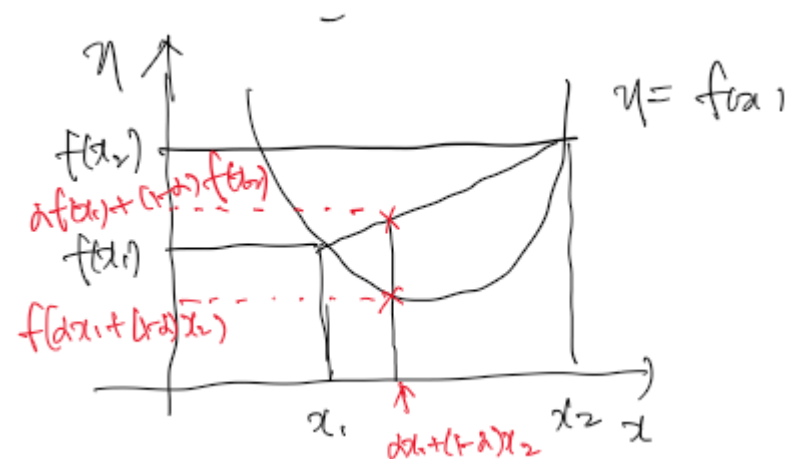
$$c_k(x) \geq 0$$

- ① Objective function is convex and ② feasible set is convex!

# Convex function

- A function  $f(x)$  is convex if

$$f(\alpha x_1 + (1 - \alpha)x_2) \leq \alpha f(x_1) + (1 - \alpha)f(x_2), \quad 0 \leq \alpha \leq 1$$



## Convex set

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- A set  $F$  is convex, if

$$x_1, x_2 \in F, \quad 0 \leq \alpha \leq 1$$

$$\alpha x_1 + (1 - \alpha)x_2 \in F$$

# Convex optimality

- Local optimum ( $x^*$ ) is global optimum      ① Objective function is convex and ② feasible set is convex!

If  $f(x') < f(x^*)$ , (Local optimal exists but there is smaller optimal ...)

$$\alpha x^* + (1-\alpha)x' \in F \quad \text{"F is convex"}$$

$$\underline{f(\alpha x^* + (1-\alpha)x')} \leq \alpha f(x^*) + (1-\alpha)f(x') \quad \text{"f is convex"}$$

$$\begin{aligned} &= f(x^*) + (1-\alpha) \underbrace{(f(x') - f(x^*))}_{< 0} \\ &< f(x^*) \end{aligned}$$

if  $\alpha \rightarrow 1$        $x^*$  주위에  $f(x^*)$  보다 작은 값이 존재!  
Feasible set 안!

# Necessity condition for optimality

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- $f' = 0$ : NC for local optimal
- What will be NC for constrained optimization?
  - Lagrangian!

# Lagrangian

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- Constrained optimization  $\Rightarrow$  Unconstrained optimization

$$\min_{\mathbf{x}} f(\mathbf{x})$$

$$\text{subject to } h_j(\mathbf{x}) = 0, j = 1, \dots, k$$



$$\min_{\mathbf{x}, \boldsymbol{\lambda}} L(\mathbf{x}, \boldsymbol{\lambda}) = \min_{\mathbf{x}, \boldsymbol{\lambda}} L(\mathbf{x}, \lambda_1, \lambda_2, \dots, \lambda_k)$$

$$L(\mathbf{x}, \boldsymbol{\lambda}) = f(\mathbf{x}) + \sum_{j=1}^k \lambda_j h_j(\mathbf{x})$$

- Local optimum  $\mathbf{x}^*$  can be obtained where  $L'(\mathbf{x}, \boldsymbol{\lambda}) = 0$

- Using gradient

$$\nabla_{\mathbf{x}} L(\mathbf{x}, \boldsymbol{\lambda}) = 0$$

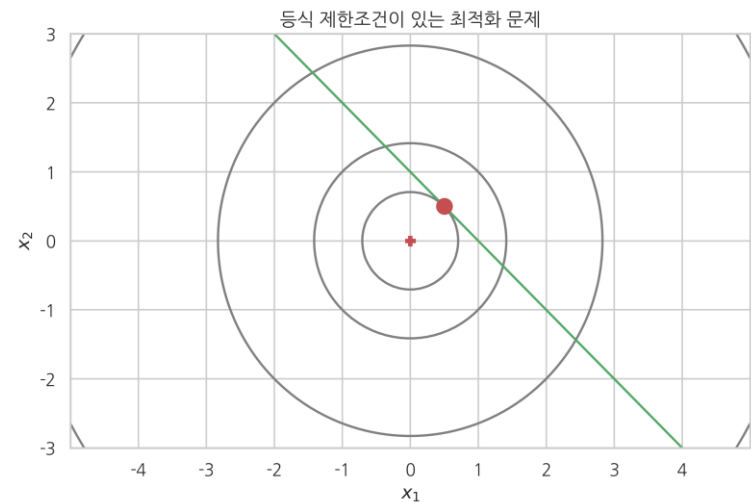
$$\nabla_{\boldsymbol{\lambda}} L(\mathbf{x}, \boldsymbol{\lambda}) = 0 \text{ (or } h_j(\mathbf{x}) = 0, j=1, \dots, k)$$

- Meaning

$$\nabla L = \nabla f + \lambda \nabla h = 0$$

$$f(x_1, x_2) = x_1^2 + x_2^2$$

$$h(x_1, x_2) = x_1 + x_2 - 1 = 0$$



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## Finding solution using Lagrangian

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- By solving system of equations

$$\nabla_x L(\mathbf{x}, \boldsymbol{\lambda}) = 0$$

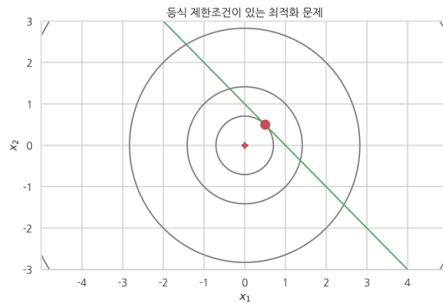
$$\nabla_{\boldsymbol{\lambda}} L(\mathbf{x}, \boldsymbol{\lambda}) = 0 \text{ (or } h_j(\mathbf{x}) = 0, j=1, \dots, k)$$

- Another solution
  - Find  $\arg \min_x L(\mathbf{x}, \boldsymbol{\lambda})$  by  $\nabla_x L(\mathbf{x}, \boldsymbol{\lambda}) = 0$
  - Use the result (find  $L(\boldsymbol{\lambda})$ )
  - Find  $\boldsymbol{\lambda}$  by  $\nabla_{\boldsymbol{\lambda}} L(\boldsymbol{\lambda}) = 0$  (maximize)



$$f(x_1, x_2) = x_1^2 + x_2^2$$

$$h(x_1, x_2) = x_1 + x_2 - 1 = 0$$



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$$\text{ex)} \quad f(x_1, x_2) = x_1^2 + x_2^2$$

$$\textcircled{M1} \quad h(x_1, x_2) = x_1 + x_2 - 1 = 0$$

$$L(x, \lambda) = x_1^2 + x_2^2 + \lambda(x_1 + x_2 - 1)$$

$$\nabla_{x_1} L = 2x_1 + \lambda = 0 \quad x_1 = -\frac{\lambda}{2}$$

$$\nabla_{x_2} L = 2x_2 + \lambda = 0 \quad x_2 = -\frac{\lambda}{2}$$

$$\nabla_{\lambda} L = x_1 + x_2 - 1 = 0 \quad -\lambda - 1 = 0 \quad \lambda = -1 \quad x_1 = \frac{1}{2} \quad x_2 = \frac{1}{2}$$

$$f(x) = \frac{1}{4} + \frac{1}{4} = \frac{1}{2}$$

$$\textcircled{M2} \quad \nabla_x L = 0 \quad 2x_1 + \lambda = 0, \quad 2x_2 + \lambda = 0 \quad x_1 = -\frac{\lambda}{2}$$

$$L(x) = \frac{\lambda^2}{4} + \frac{\lambda^2}{4} + \lambda(-\frac{\lambda}{2} - \frac{\lambda}{2} - 1)$$

$$= \frac{\lambda^2}{2} - \lambda^2 - \lambda$$

$$= -\frac{\lambda^2}{2} - \lambda$$

$$\nabla_{\lambda} L = -\lambda - 1 = 0 \quad \therefore \lambda = -1$$

## With inequality in the condition

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- With inequality

$$\min_{\mathbf{x}} f(\mathbf{x})$$

$$\text{subject to } g_i(\mathbf{x}) \leq 0, i = 1, \dots, m$$

$$h_j(\mathbf{x}) = 0, j = 1, \dots, k$$

$$L(\mathbf{x}, \boldsymbol{\lambda}) = f(\mathbf{x}) + \sum_{i=1}^m \mu_i g_i(\mathbf{x}) + \sum_{j=1}^k \lambda_j h_j(\mathbf{x})$$

- Need additional things: KKT (Karush-Kuhn-Tucker) conditions

- Stationarity:

$$\nabla_x f(\mathbf{x}) + \sum_{i=1}^m \mu_i \nabla_x g_i(\mathbf{x}) + \sum_{j=1}^k \lambda_j \nabla_x h_j(\mathbf{x}) = 0$$

- Primal constraints:

$$g_i(\mathbf{x}) \leq 0, i = 1, \dots, m$$

$$h_j(\mathbf{x}) = 0, j = 1, \dots, k$$

- Dual constraints:

$$\mu_i \geq 0, i = 1, \dots, m$$


- Complementary slackness:

$$\mu_i g_i(\mathbf{x}) = 0, i = 1, \dots, m$$

# How to solve

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- By KKT
- By duality
  - Find  $\arg \min_x L(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda})$  by  $\nabla_{\mathbf{x}} L(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda}) = 0$
  - Find  $\boldsymbol{\lambda}$  by  $\nabla_{\boldsymbol{\mu}, \boldsymbol{\lambda}} \min_x L(\mathbf{x}, \boldsymbol{\mu}, \boldsymbol{\lambda}) = 0$  (maximize)


$$q(\boldsymbol{\mu}, \boldsymbol{\lambda})$$

# Duality

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- Dual problem as a lower bound to the primal

$$f(x) \geq L(x, \mu, \lambda) \geq \min_x L(x, \mu, \lambda) = q(\mu, \lambda)$$



Lower bound

$$f^* = f(x^*) \geq L(x^*, \mu, \lambda)$$

$$q(\mu, \lambda) \leq \max_{\mu \geq 0, \lambda} q(\mu, \lambda) = q^* \leq f^* \text{ (weak duality)}$$

$$q^* = f^* \text{ (strong duality)}$$

- Duality gap

▪ Optimal

- Feasible  $x, \mu, \lambda$  such that  $q=f$

$$q(\mu, \lambda) \leq q^* \leq f^* \leq f(x)$$

▪ If  $q=f$

$$q(\mu, \lambda) = q^*$$

$$f(x) = f^*$$

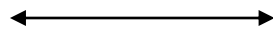
$$\min_x f(x)$$

$$\text{s. t. } g(x) \leq 0$$

$$h(x) = 0$$

$$\max q(\mu, \lambda)$$

$$\text{s. t. } \mu \geq 0$$



# Duality in LP

$$\begin{array}{ll}\text{minimize} & 6x_1 + 4x_2 + 2x_3 \\ \text{subject to} & 4x_1 + 2x_2 + x_3 \geq 5, \\ & x_1 + x_2 \geq 3, \\ & x_2 + x_3 \geq 4, \\ & x_1, x_2, x_3 \geq 0.\end{array}$$

$$6x_1 + 4x_2 + 2x_3 \geq 4x_1 + 2x_2 + x_3 \geq 5$$

$$6x_1 + 4x_2 + 2x_3 \geq (4x_1 + 2x_2 + x_3) + 2 \cdot (x_1 + x_2) \geq 11$$

LB

If we generalize...

$$\begin{array}{ll}\text{minimize} & 6x_1 + 4x_2 + 2x_3 \\ \text{subject to} & 4x_1 + 2x_2 + x_3 \geq 5, \\ & x_1 + x_2 \leq 3, \\ & x_2 + x_3 = 4, \\ & x_1 \geq 0, \\ & x_2 \leq 0, \\ & (x_3 \text{ is free}).\end{array}$$



$$\begin{array}{ll}\text{maximize} & 5y_1 + 3y_2 + 4y_3 \\ \text{subject to} & y_1 \geq 0, \\ & y_2 \leq 0, \\ & (y_3 \text{ is free}), \\ & 4y_1 + y_2 \leq 6, \\ & 2y_1 + y_2 + y_3 \geq 4, \\ & y_1 + y_3 = 2.\end{array}$$

$$6x_1 + 4x_2 + 2x_3 \geq y_1 \cdot (4x_1 + 2x_2 + x_3) + y_2 \cdot (x_1 + x_2) + y_3 \cdot (x_2 + x_3)$$

$$y_1 \cdot (4x_1 + 2x_2 + x_3) + y_2 \cdot (x_1 + x_2) + y_3 \cdot (x_2 + x_3) \geq 5y_1 + 3y_2 + 4y_3$$

$$y_1 \geq 0, \quad y_2 \leq 0, \quad y_3 \text{ free}$$

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▪ Ex

$$\min x_1 + x_2$$

$$\text{s. t. } x_1^2 + x_2^2 \leq 1$$





# Convex Optimization in ML:SVM

- A separating hyperplane can be written as

$$\mathbf{W} \cdot \mathbf{X} + b = 0$$

where  $\mathbf{W} = \{w_1, w_2, \dots, w_n\}$  is a weight vector and  $b$  a scalar (bias)

- For 2-D it can be written as

$$w_0 + w_1 x_1 + w_2 x_2 = 0$$

- The hyperplane defining the sides of the margin:

$$H_1: w_0 + w_1 x_1 + w_2 x_2 \geq 1 \quad \text{for } y_i = +1, \text{ and}$$

$$H_2: w_0 + w_1 x_1 + w_2 x_2 \leq -1 \quad \text{for } y_i = -1$$

- Any training tuples that fall on hyperplanes  $H_1$  or  $H_2$  (i.e., the sides defining the margin) are **support vectors**
- This becomes a **constrained (convex) quadratic optimization** problem: Quadratic objective function and linear constraints  $\rightarrow$  *Quadratic Programming (QP)*  $\rightarrow$  Lagrangian multipliers

- Objective function

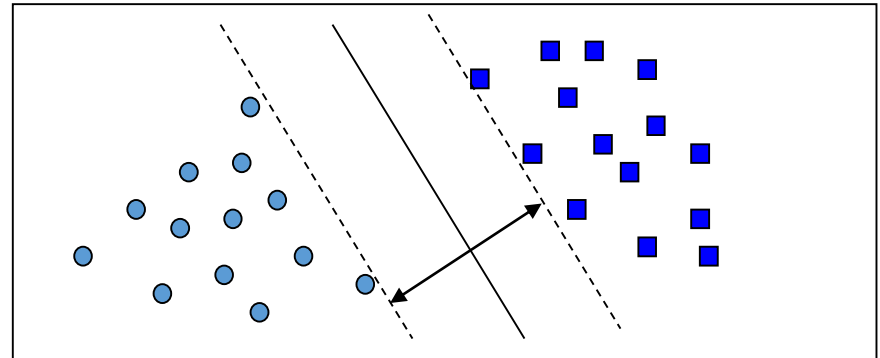
$$b_{i1}: \mathbf{w} \cdot \mathbf{x} + b = 1$$

$$b_{i2}: \mathbf{w} \cdot \mathbf{x} + b = -1$$

$$\mathbf{w} \cdot (\mathbf{x}_1 - \mathbf{x}_2) = 2$$

$$\|\mathbf{w}\| \times d = 2$$

$$d = \frac{2}{\|\mathbf{w}\|}$$



- Constraints

$$\mathbf{w} \cdot \mathbf{x}_i + b \geq 1 \quad (y_i = 1)$$

$$\mathbf{w} \cdot \mathbf{x}_i + b \leq -1 \quad (y_i = -1)$$

$$y_i(\mathbf{w} \cdot \mathbf{x}_i + b) \geq 1, i = 1, 2, \dots, N$$

- SVM Learning

$$\min_{\mathbf{w}} f(\mathbf{w}) = \frac{\|\mathbf{w}\|^2}{2}$$

$$\text{s.t.} \quad y_i(\mathbf{w} \cdot \mathbf{x}_i + b) \geq 1, i = 1, 2, \dots, N$$

- Convex optimization => Lagrangian

$$L_P = \frac{1}{2} \|\mathbf{w}\|^2 - \sum_{i=1}^N \lambda_i (y_i(\mathbf{w} \cdot \mathbf{x}_i + b) - 1)$$

$$\frac{\partial L_P}{\partial \mathbf{w}} = 0 \implies \mathbf{w} = \sum_{i=1}^N \lambda_i y_i \mathbf{x}_i$$

$$\frac{\partial L_P}{\partial b} = 0 \implies \sum_{i=1}^N \lambda_i y_i = 0$$

- Dual Lagrangian

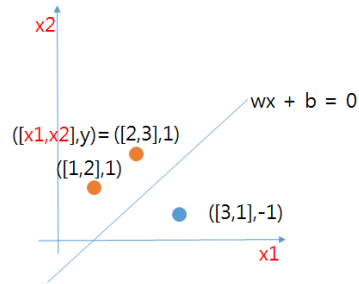
$$L_D = \sum_{i=1}^N \lambda_i - \frac{1}{2} \sum_{i,j} \lambda_i \lambda_j y_i y_j \mathbf{x}_i \cdot \mathbf{x}_j$$

- KKT (Karush-Khun-Tucker) condition for inequality in the constraints: [necessary condition](#)

최적해가 되기 위한 조건

$$\lambda_i \geq 0$$

$$\lambda_i (y_i(\mathbf{w} \cdot \mathbf{x}_i + b) - 1) = 0$$



$$L = \frac{1}{2}w^2 + \sum \alpha_i [y_i(wx + b) - 1]$$

Then differentiate the above equations in terms of w and b

$$\frac{\partial L}{\partial w} = w - \sum \alpha_i y_i x_i = 0$$

$$w = \sum \alpha_i y_i x_i \dots (1)$$

$$\frac{\partial L}{\partial b} = \sum \alpha_i y_i = 0 \dots (2)$$

Apply (1) and (2) into L

$$L = \frac{1}{2}(\sum \alpha_i y_i x_i)^2 - \sum [\alpha_i y_i (\sum \alpha_j y_j x_j) + \alpha_i y_i b - \alpha_i]$$

$$= \frac{1}{2}(\sum \alpha_i y_i x_i)^2 - \sum \alpha_i y_i x_i (\sum \alpha_j y_j x_j) - \sum \alpha_i y_i b + \sum \alpha_i$$

According to equation(2),  $\sum \alpha_i y_i b = 0$

$$= \frac{1}{2}(\sum \alpha_i y_i x_i)^2 - (\sum \alpha_i y_i x_i)(\sum \alpha_j y_j x_j) + \sum \alpha_i$$

$$= -\frac{1}{2}(\sum \alpha_i y_i x_i \sum \alpha_j y_j x_j) + \sum \alpha_i$$

$$= \sum \alpha_i - \frac{1}{2}(\sum \alpha_i y_i x_i \sum \alpha_j y_j x_j)$$

$$= \sum \alpha_i - \frac{1}{2}(\sum \sum \alpha_i \alpha_j y_i y_j x_i x_j)$$

$$x_1 = (1, 2), y_1 = 1$$

$$x_2 = (2, 3), y_2 = 1$$

$$x_3 = (3, 1), y_3 = -1$$

By applying into the equations,

$$L = \sum \alpha_i - \frac{1}{2} \sum \alpha_i y_i x_i (\alpha_1 y_1 x_1 + \alpha_2 y_2 x_2 + \alpha_3 y_3 x_3)$$

$$= \sum \alpha_i - \frac{1}{2} \sum \alpha_i y_i x_i (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3)$$

$$= \sum \alpha_i - \frac{1}{2} (\alpha_1 y_1 x_1 (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3) - \frac{1}{2} (\alpha_2 y_2 x_2 (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3)) - \frac{1}{2} (\alpha_3 y_3 x_3 (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3)))$$

$$= \sum \alpha_i - \frac{1}{2} (\alpha_1 x_1 (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3) - \frac{1}{2} (\alpha_2 x_2 (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3)) + \frac{1}{2} (\alpha_3 x_3 (\alpha_1 x_1 + \alpha_2 x_2 - \alpha_3 x_3)))$$

$$= \sum \alpha_i - \frac{1}{2} (5\alpha_1^2 + 13\alpha_2^2 + 10\alpha_3^2 + 16\alpha_1\alpha_2 - 10\alpha_1\alpha_3 - 18\alpha_2\alpha_3)$$

$$= \alpha_1 + \alpha_2 + \alpha_3 - \frac{1}{2} (5\alpha_1^2 + 13\alpha_2^2 + 10\alpha_3^2 + 16\alpha_1\alpha_2 - 10\alpha_1\alpha_3 - 18\alpha_2\alpha_3)$$

$$\frac{\partial L}{\partial \alpha_1} = 1 + \frac{1}{2} (-10\alpha_1 - 16\alpha_2 + 10\alpha_3) = 0$$

Remember we know that  $\sum \alpha_i y_i = 0$ , so  $\alpha_1 + \alpha_2 = \alpha_3$

$$1 + \frac{1}{2} (-10\alpha_1 - 16\alpha_2 + 10(\alpha_1 + \alpha_2)) = 0$$

$$\alpha_2 = \frac{1}{3}$$

$$\frac{\partial L}{\partial \alpha_3} = 1 - \frac{1}{2} (26\alpha_2 + 16\alpha_1 - 18\alpha_3) = 0$$

$$\frac{\partial L}{\partial \alpha_3} = 1 - (13\alpha_2 + 8\alpha_1 - 9\alpha_3) = 0$$

$$8\alpha_1 - 9\alpha_3 = -\frac{10}{3}$$

We know that  $\alpha_1 + \alpha_2 = \alpha_3$ , that is,  $\alpha_1 + \frac{1}{3} = \alpha_3$

Finally we can find  $(\alpha_1, \alpha_2, \alpha_3) = (\frac{1}{3}, \frac{1}{3}, \frac{2}{3})$

By applying  $\alpha$  into  $w = \sum \alpha_i y_i x_i$ , we can get  $w$

$$w = (-1, 1)$$

Our plane is  $y = wx + b$ , therefore  $y = (-a, a)(x_1, x_2) + b = -ax_1 + ax_2 + b$

Apply  $x_1 = (1, 2), y_1 = 1, x_2 = (3, 1), y_3 = -1$  and get  $a = \frac{2}{3}, b = \frac{1}{3}$

- For test case  $\mathbf{z}$ ,

$$f(\mathbf{z}) = \text{sign}(\mathbf{w} \cdot \mathbf{z} + b) = \text{sign}\left(\sum_{i=1}^N \lambda_i y_i x_i \cdot \mathbf{z} + b\right)$$

- If  $f=1$ ,  $\mathbf{z}$  will be classified as positive

# Why Is SVM Effective on High Dimensional Data?

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- The **complexity** of trained classifier is characterized by the # of support vectors rather than the dimensionality of the data
- The **support vectors** are the essential or critical training examples —they lie closest to the decision boundary (MMH)
- If all other training examples are removed and the training is repeated, the same separating hyperplane would be found
- The number of support vectors found can be used to compute an (upper) bound on the expected error rate of the SVM classifier, which is independent of the data dimensionality
- Thus, an SVM with a small number of support vectors can have good generalization, even when the dimensionality of the data is high

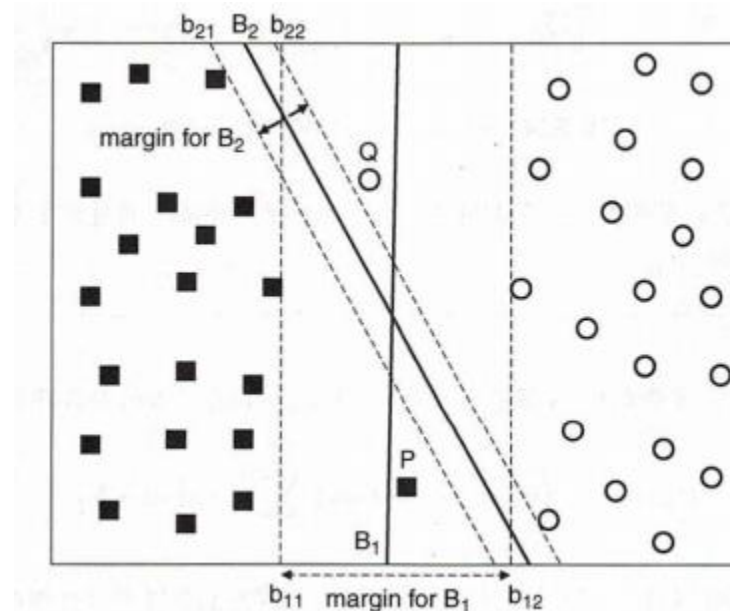
# Soft Margin

$$\mathbf{w} \cdot \mathbf{x}_i + b \geq 1 - \xi_i \quad (y_i = 1)$$

$$\mathbf{w} \cdot \mathbf{x}_i + b \leq -1 + \xi_i \quad (y_i = -1)$$

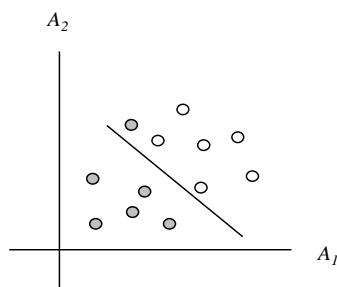
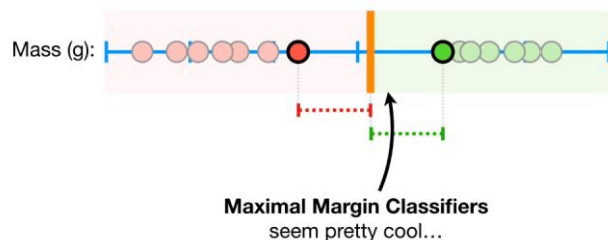
$$\min_{\mathbf{w}} f(\mathbf{w}) = \frac{\|\mathbf{w}\|^2}{2} + C \left( \sum_{i=1}^N \xi_i \right)^k$$

$$L_p = \frac{1}{2} \|\mathbf{w}\|^2 - C \sum_{i=1}^N \xi_i - \sum_{i=1}^N \lambda_i \{ \mathbf{y}_i (\mathbf{w} \cdot \mathbf{x}_i + b) - 1 + \xi_i \} - \sum_{i=1}^N \mu_i \xi_i$$



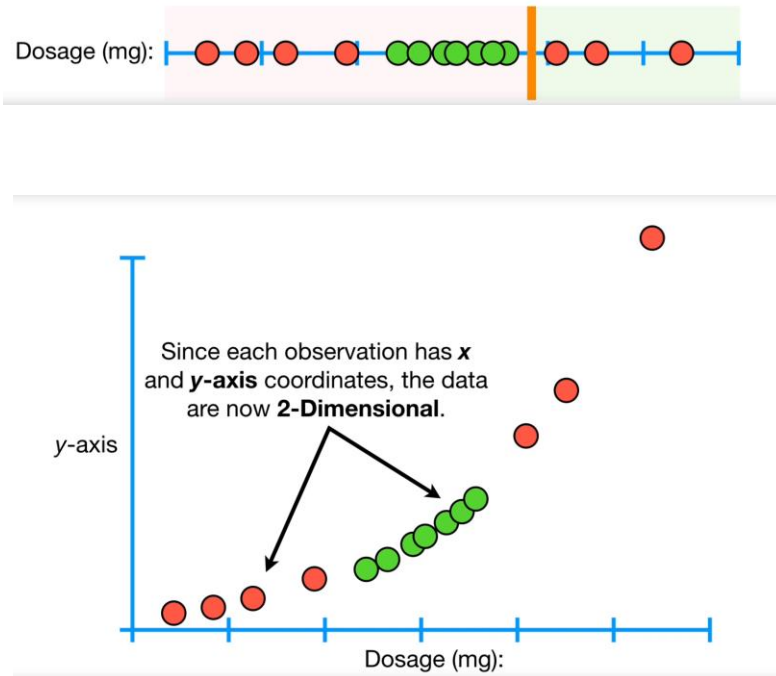
# SVM—Linearly Inseparable

- Transform the original input data into a higher dimensional space



Example 6.8 Nonlinear transformation of original input data into a higher dimensional space. Consider the following example. A 3D input vector  $\mathbf{X} = (x_1, x_2, x_3)$  is mapped into a 6D space  $\mathbf{Z}$  using the mappings  $\phi_1(\mathbf{X}) = x_1, \phi_2(\mathbf{X}) = x_2, \phi_3(\mathbf{X}) = x_3, \phi_4(\mathbf{X}) = (x_1)^2, \phi_5(\mathbf{X}) = x_1x_2$ , and  $\phi_6(\mathbf{X}) = x_1x_3$ . A decision hyperplane in the new space is  $d(\mathbf{Z}) = \mathbf{WZ} + b$ , where  $\mathbf{W}$  and  $\mathbf{Z}$  are vectors. This is linear. We solve for  $\mathbf{W}$  and  $b$  and then substitute back so that we see that the linear decision hyperplane in the new ( $\mathbf{Z}$ ) space corresponds to a nonlinear second order polynomial in the original 3-D input space,

$$\begin{aligned} d(\mathbf{Z}) &= w_1x_1 + w_2x_2 + w_3x_3 + w_4(x_1)^2 + w_5x_1x_2 + w_6x_1x_3 + b \\ &= w_1z_1 + w_2z_2 + w_3z_3 + w_4z_4 + w_5z_5 + w_6z_6 + b \end{aligned}$$



- Search for a linear separating hyperplane in the new space



- Kernel: Higher degree mapping, inner product together

$$\begin{aligned}\max L(\alpha_i) &= \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j t_i t_j \Phi(x_i)^T \Phi(x_j) \\ &= \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i=1}^n \sum_{j=1}^n \alpha_i \alpha_j t_i t_j K(x_i, x_j)\end{aligned}$$

$$K(x_i, x_j) = \Phi(x_i)^T \Phi(x_j)$$

- Kernel trick:  $\Phi(x_i)^T, \Phi(x_j)$ 를 함수로 구현

$\Phi(x) = Ax$  :  $m$ 차 feature를  $n$ 차로  
 $n \times 1$     $n \times m$     $m \times 1$

$$K(x_i, x_j) = \underbrace{\Phi(x_i)^T}_{1 \times n} \underbrace{\Phi(x_j)}_{n \times 1} = x_i^T A^T A x_j \Rightarrow 1 \times 1 \text{ (scalar)}$$

- 변환을 해서 내적을 하는것이 아닌, 변환된 곳에서의 내적을 위처럼 표현이 가능하다.
  - 데이터 -> 변환 -> 내적 순서대로 진행하는것이 아닌, 변환을 건너뛰고 바로 내적을 이끌어 내는 과정이 kernel이다.
  - 즉, Kernel은 “**높은 차원에서의 내적을 계산하는 함수**” 이다.

## SVM: Different Kernel functions

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- Instead of computing the dot product on the transformed data, it is math. equivalent to applying a kernel function  $K(\mathbf{X}_i, \mathbf{X}_j)$  to the original data, i.e.,  $K(\mathbf{X}_i, \mathbf{X}_j) = \Phi(\mathbf{X}_i) \cdot \Phi(\mathbf{X}_j)$
- Typical Kernel Functions

Polynomial kernel of degree  $h$  :  $K(\mathbf{X}_i, \mathbf{X}_j) = (\mathbf{X}_i \cdot \mathbf{X}_j + 1)^h$

Gaussian radial basis function kernel :  $K(\mathbf{X}_i, \mathbf{X}_j) = e^{-\|\mathbf{X}_i - \mathbf{X}_j\|^2 / 2\sigma^2}$

Sigmoid kernel :  $K(\mathbf{X}_i, \mathbf{X}_j) = \tanh(\kappa \mathbf{X}_i \cdot \mathbf{X}_j - \delta)$

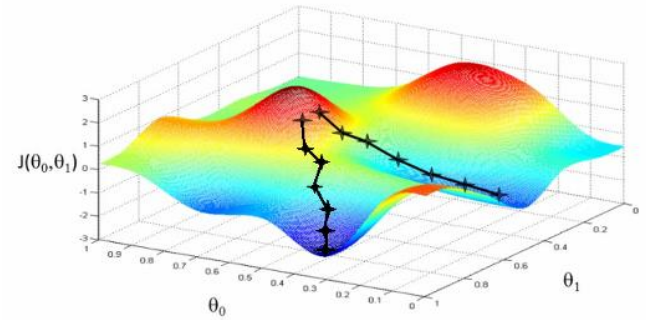
- SVM can also be used for classifying multiple ( $> 2$ ) classes and for regression analysis (with additional parameters)

## 2. Using differentiation for unconstrained optimization

- Minimum (Maximum) value is on the point that

$$f'(x) = 0$$

- Finding  $x$  such that  $f'(x) = 0$

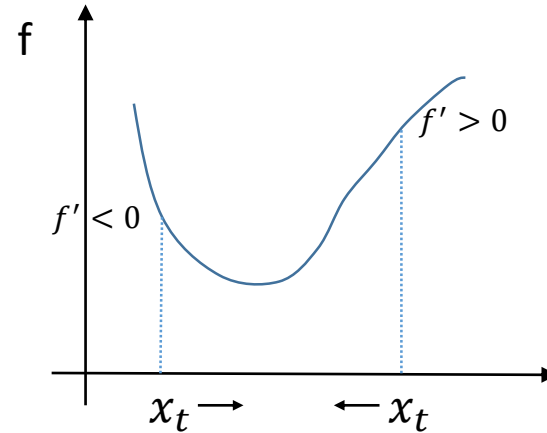


출처: <http://blog.datumbox.com/tuning-the-learning-rate-in-gradient-descent/>

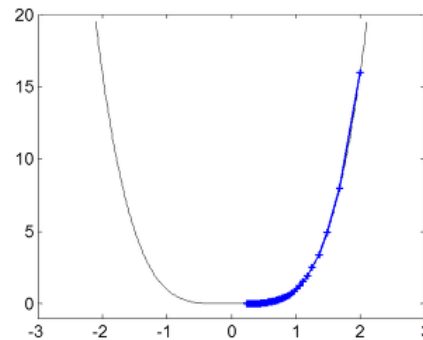
# Gradient descent

- Update  $x$  value starting from  $x_0$

$$x_{t+1} = x_t - \lambda f'(x_t)$$



$$f(x)=x^4$$



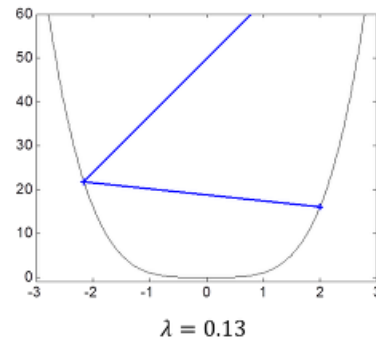
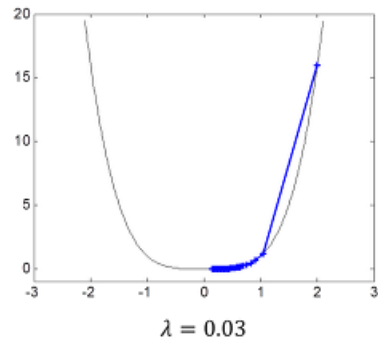
$x_0 = 2.0$   
 $x_1 = x_0 - \lambda f'(x_0) = 1.6800$   
 $x_2 = x_1 - \lambda f'(x_1) = 1.4903$   
 $x_3 = x_2 - \lambda f'(x_2) = 1.3579$   
 $x_4 = x_3 - \lambda f'(x_3) = 1.2578$   
 $x_5 = x_4 - \lambda f'(x_4) = 1.1782$   
 $x_6 = x_5 - \lambda f'(x_5) = 1.1128$   
 $x_7 = x_6 - \lambda f'(x_6) = 1.0576$   
 $x_8 = x_7 - \lambda f'(x_7) = 1.0103$   
 $x_9 = x_8 - \lambda f'(x_8) = 0.9691$   
...

<https://darkpgmr.tistory.com/149>

- Problems

- Even after 200 times of updating, it is not close to  $x=0$

- If we use larger value of  $\lambda$



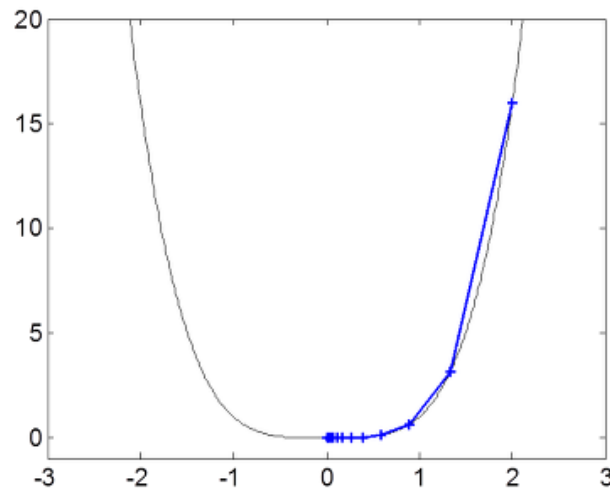
## Using secondary derivatives: Newton's method

- Update using the equation: 보다 빠르게 수렴

$$x_{t+1} = x_t - \frac{f'(x_t)}{f''(x_t)}$$

- Problems: 변곡점에서는 정의되지 않음

$$f''(x_t)=0$$



$x_0 = 2.0$   
 $x_1 = 1.3333$   
 $x_2 = 0.8889$   
 $x_3 = 0.5926$   
 $x_4 = 0.3951$   
 $x_5 = 0.2634$   
 $x_6 = 0.1756$   
 $x_7 = 0.1171$   
 $x_8 = 0.0780$   
 $x_9 = 0.0520$   
...

# Using random sampling

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- Sampling base inference

# Black-box optimization

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- Bayesian optimization



# Conclusions

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- ML for opt.
- Opt. for ML

## references

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- [https://developers.google.com/optimization/lp/lp\\_example#python\\_7](https://developers.google.com/optimization/lp/lp_example#python_7)
- <https://darkpgmr.tistory.com/149>

# Geometric way

$$\text{Maximize } 2x_1 + 5x_2$$

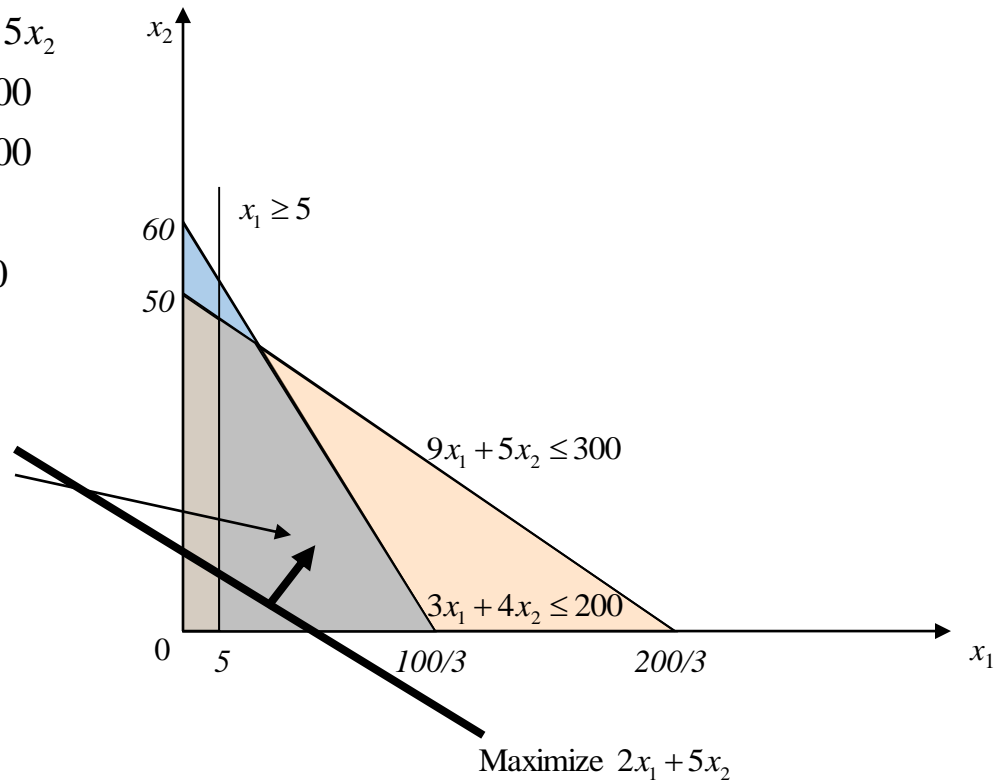
$$\text{s.t. } 3x_1 + 4x_2 \leq 200$$

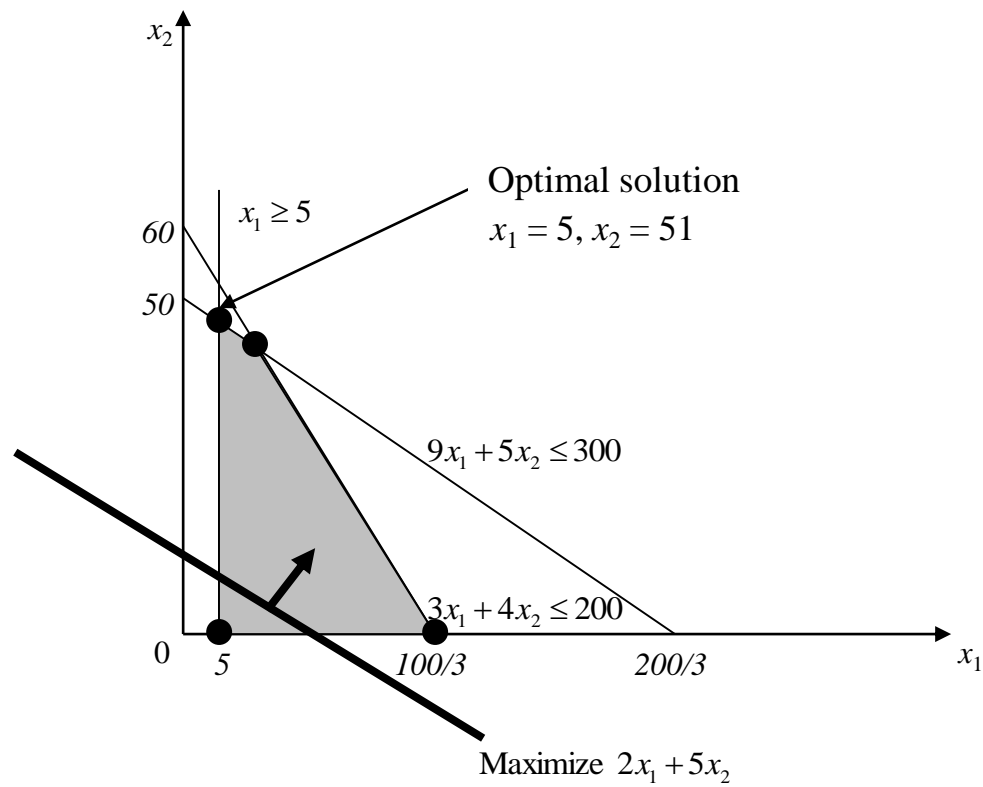
$$9x_1 + 5x_2 \leq 300$$

$$x_1 \geq 5$$

$$x_2 \geq 0$$

Feasible solution space





[Ex] Production plan for maximization of profit

A factory produces three products: A, B, and C. One unit of product A requires 2.4 minutes of rolling time and 5.0 minutes of assembly process. Profit is \$600. One unit of product B requires 3.0 minutes of rolling time, 2.5 minutes of welding process, and \$700 of profit is generated. One unit of product C requires 2.0 minutes of rolling time, 1.5 minutes of welding time, 2.5 minutes of assembly time, and generates a profit of \$500.

The production time of the rolling process is 1,200 minutes per week, the welding process is 600 minutes per week, and the assembly process can be 1,500 minutes per week.

What is the production of products A, B, and C that can generate the greatest profit?

$$\begin{array}{ll} \text{Maximize } 600x_1 + 700x_2 + 500x_3 & \longleftarrow \text{Profit maximization} \\ \text{s.t. } 2.4x_1 + 3.0x_2 + 2.0x_3 \leq 1200 & \longleftarrow \text{Rolling time constraints} \\ 0.0x_1 + 2.5x_2 + 1.5x_3 \leq 600 & \longleftarrow \text{Welding time constraints} \\ 5.0x_1 + 0.0x_2 + 2.5x_3 \leq 1500 & \longleftarrow \text{Assembly time constraints} \\ x_1, x_2, x_3 \geq 0 & \longleftarrow \text{Non-negative} \end{array}$$

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■ Geometric approach?

- Solutions are located at Vertexes
- Within 2 or 3 dimension
  - Not easy for more than 3-dimensional problems
  - Not easy for a problem with a large number of constraints

=> Need algorithmic approach

- Simplex Method
  - Dantzig
  - Exploring the optimal solution among the intersections of constraints.
- Karmarkar Method

# Simplex Method

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- Simplex method(單體法)
  - Using simultaneous linear equations
    - Matrix operation: Gauss-Jordan elimination
  - Easy to understand, easy to use
  - Find the opt. sol. among vertexes in a feasible space
    - Initial basic solutions => [repetitive improvement] => optimal solution

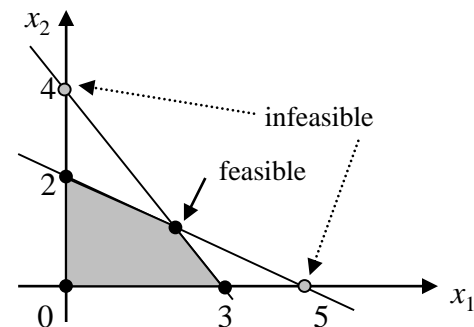
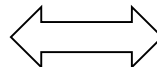
$$\left\{ \begin{array}{ll} \text{Maximize} & 12x_1 + 15x_2 \\ \text{s.t.} & 4x_1 + 3x_2 \leq 12 \\ & 2x_1 + 5x_2 \leq 10 \\ & x_1, x_2 \geq 0 \end{array} \right\} \xrightarrow{\text{To equality}} \left\{ \begin{array}{ll} \text{Maximize} & 12x_1 + 15x_2 \\ \text{s.t.} & 4x_1 + 3x_2 + x_3 = 12 \\ & 2x_1 + 5x_2 + x_4 = 10 \\ & x_1, x_2, x_3, x_4 \geq 0 \end{array} \right\}$$

System of equations

$$\left\{ \begin{array}{ll} 4x_1 + 3x_2 + x_3 & = 12 \\ 2x_1 + 5x_2 + x_4 & = 10 \end{array} \right\} \text{ Infinite number of solutions}$$

6 solutions by pairs of solutions ( ${}_4C_2$ )

구분	해 $(x_1, x_2, x_3, x_4)$
$(x_1, x_2)$	$(15/7, 8/7, 0, 0)$
$(x_1, x_3)$	$(5, 0, -8, 0)$
$(x_1, x_4)$	$(3, 0, 0, 4)$
$(x_2, x_3)$	$(0, 2, 6, 0)$
$(x_2, x_4)$	$(0, 4, 0, -10)$
$(x_3, x_4)$	$(0, 0, 12, 10)$





Object function  $\xrightarrow{\text{Maximize } 12x_1 + 15x_2}$

A single solution that maximizes the O-function => optimal solution

Basis	Sol. $(x_1, x_2, x_3, x_4)$	Values of O-function	
$(x_1, x_2)$	$(15/7, 8/7, 0, 0)$	$300/7$	Infeasible
$(x_1, x_3)$	$(5, 0, -8, 0)$	60	Optimal solution
$(x_1, x_4)$	$(3, 0, 0, 4)$	36	
$(x_2, x_3)$	$(0, 2, 6, 0)$	30	
$(x_2, x_4)$	$(0, 4, 0, -10)$	60	
$(x_3, x_4)$	$(0, 0, 12, 10)$	0	Infeasible

A solution using simply a system of equations of constraints should find  ${}_nC_m$  solutions if it is a problem with m constraints consisting of n determinants. Depending on the values of n and m, the search space increases exponentially.

=> There is a need for a method that does not search for infeasible solutions, but only within the feasible, but steadily increases the value of the objective function.

## Simplex method

Maximize  $4x_1 + 5x_2 + 8x_3 + 11x_4$

s.t.  $x_1 + x_2 + x_3 + x_4 \leq 15$

$3x_1 + 5x_2 + 10x_3 + 15x_4 \leq 105$

$7x_1 + 5x_2 + 3x_3 + 2x_4 \leq 120$

$x_j \geq 0, \forall j$

Min. Coef. = -11

Init. Sol.

$$\begin{array}{rcl}
 z - 4x_1 - 5x_2 - 8x_3 - 11x_4 & = & 0 \\
 x_1 + x_2 + x_3 + x_4 + (x_5) & = & 15 \quad \leftarrow 15/1 = 15 \\
 3x_1 + 5x_2 + 10x_3 + 15x_4 + (x_6) & = & 105 \quad \leftarrow 105/15 = 7 \\
 7x_1 + 5x_2 + 3x_3 + 2x_4 + (x_7) & = & 120 \quad \leftarrow 120/2 = 60 \\
 x_j \geq 0, \forall j
 \end{array}$$

Adding slack var.

pivoting

Min. ratio

$$\begin{array}{rcl}
 z - \frac{9}{5}x_1 - \frac{4}{3}x_2 - \frac{2}{3}x_3 + \frac{11}{15}x_6 & = & 77 \\
 \frac{4}{5}x_1 + \frac{2}{3}x_2 + \frac{1}{3}x_3 - \frac{1}{15}x_6 & = & 8 \\
 \frac{1}{5}x_1 + \frac{1}{3}x_2 + \frac{2}{3}x_3 + \frac{1}{15}x_6 & = & 7 \\
 \frac{33}{5}x_1 + \frac{13}{3}x_2 + \frac{5}{3}x_3 - \frac{2}{15}x_6 + (x_7) & = & 106 \\
 x_j \geq 0, \forall j
 \end{array}$$

A new solution (0, 0, 0, 7, 8, 0, 106)

A new value of O-function 77

$$\begin{array}{rcl}
 z & \frac{9}{5}x_1 - \frac{4}{3}x_2 - \frac{2}{3}x_3 & + \frac{11}{15}x_6 = 77 \\
 & \frac{4}{5}x_1 + \frac{2}{3}x_2 + \frac{1}{3}x_3 & + x_5 - \frac{1}{15}x_6 = 8 \\
 & \frac{1}{5}x_1 + \frac{1}{3}x_2 + \frac{2}{3}x_3 + x_4 & + \frac{1}{15}x_6 = 7 \\
 & \frac{33}{5}x_1 + \frac{13}{3}x_2 + \frac{5}{3}x_3 & - \frac{2}{15}x_6 + x_7 = 106 \\
 & x_j \geq 0, \forall j
 \end{array}$$

Pivoting

All the coef.  $\geq 0$

$$\begin{array}{rcl}
 z & + \frac{1}{6}x_2 + \frac{1}{12}x_3 & + \frac{9}{4}x_5 + \frac{7}{12}x_6 = 95 \\
 \textcircled{x_1} & + \frac{5}{6}x_2 + \frac{5}{12}x_3 & + \frac{5}{4}x_5 - \frac{1}{12}x_6 = 10 \\
 & \frac{1}{6}x_2 + \frac{7}{12}x_3 + \textcircled{x_4} - \frac{1}{4}x_5 + \frac{1}{12}x_6 & = 5 \\
 & -\frac{7}{6}x_2 - \frac{13}{12}x_3 - \frac{33}{4}x_5 + \frac{5}{12}x_6 + \textcircled{x_7} & = 40 \\
 & x_j \geq 0, \forall j
 \end{array}$$

Opt. Sol.  $x^* = (10, 0, 0, 5, 0, 0, 40)$

Opt. Val.  $z^* = 95$

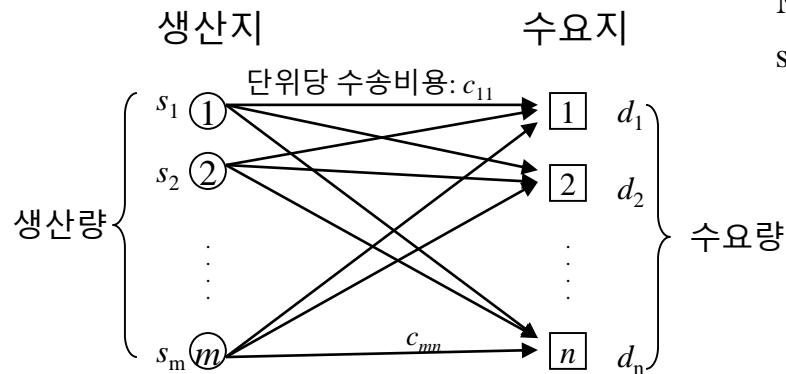
## 단체표의 이용 [요약정리]

$z$	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$	
1	-4	-5	-8	-11	0	0	0	0
0	1	1	1	1	1	0	0	15
0	3	5	10	15	0	1	0	105
0	7	5	3	2	0	0	1	120
1	-9/5	-4/3	-2/3	0	0	11/15	0	77
0	4/5	2/3	1/3	0	1	-1/15	0	8
0	1/5	1/3	2/3	1	0	1/15	0	7
0	33/5	13/3	5/3	0	0	-2/15	1	106
1	0	1/6	1/12	0	9/4	7/12	0	95
0	1	5/6	5/12	0	5/4	-1/12	0	10
0	0	1/6	7/12	1	-1/4	1/12	0	5
0	0	-7/6	-13/12	0	-33/4	5/12	1	40

# Other Linear Problems

## ■ Transportation Problem

$m$ 개의 생산지로부터  $n$ 개의 수요지로 수요량을 만족시키면서, 최소의 비용으로 전달하는 문제



$$\text{Minimize } c_{11}x_{11} + c_{12}x_{12} + \dots + c_{mn}x_{mn}$$

$$\begin{aligned} \text{s.t. } & x_{11} + x_{12} + \dots + x_{1n} \leq s_1 \\ & \dots \\ & x_{m1} + x_{m2} + \dots + x_{mn} \leq s_m \\ & x_{11} + x_{21} + \dots + x_{m1} \geq d_1 \\ & \dots \\ & x_{1n} + x_{2n} + \dots + x_{mn} \geq d_n \\ & x_{ij} \geq 0, \quad \forall i, j \end{aligned}$$

# Optimizing transportation

A 건설회사에서 3곳의 야산으로부터 모래를 운반하여 4곳의 아파트 부지에 공급한다. 모래의 운반과 관련한 비용 및 생산량과 수요량이 다음의 행렬에 정리되어 있다. 최소의 운반 비용을 얻을 수 있는 수송 경로를 구하여라.

야산 \ 아파트 부지	d <sub>1</sub>	d <sub>2</sub>	d <sub>3</sub>	d <sub>4</sub>	공급량
s <sub>1</sub>	2	3	11	7	6
s <sub>2</sub>	1	0	6	1	1
s <sub>3</sub>	5	8	15	9	10
수요량	7	5	3	2	합: 17

단위: 100만원/톤

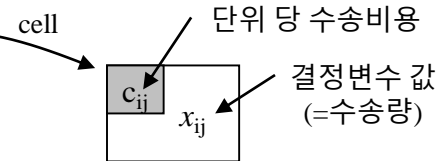
$$\text{Minimize } 2x_{11} + 3x_{12} + 11x_{13} + 7x_{14} + x_{21} + 6x_{23} + x_{24} + 5x_{31} + 8x_{32} + 15x_{33} + 9x_{34}$$

$$\begin{aligned} \text{s.t. } & x_{11} + x_{12} + x_{13} + x_{14} \leq 6 \\ & x_{21} + x_{22} + x_{23} + x_{24} \leq 1 \\ & x_{31} + x_{32} + x_{33} + x_{34} \leq 10 \\ & x_{11} + x_{21} + x_{31} \geq 7 \\ & x_{12} + x_{22} + x_{32} \geq 5 \\ & x_{13} + x_{23} + x_{33} \geq 3 \\ & x_{14} + x_{24} + x_{34} \geq 2 \\ & x_{ij} \geq 0, \quad \forall i, j \end{aligned}$$

- Step 1) 초기해를 여러 가지 방법을 활용하여 설정한다.  
 Step 2) 초기해를 개선시킬 수 있는 수송비용의 음환(negative cycle)을 찾는다.  
 Step 3) 더 이상 해를 개선시킬 수 있는 음환을 찾을 수 없을 때가 최적이다.

1) 최소가법을 이용하여 설정된 초기해

공급지 \ 수요지	d <sub>1</sub>	d <sub>2</sub>	d <sub>3</sub>	d <sub>4</sub>	공급량
s <sub>1</sub>	2	3	11	7	6
s <sub>2</sub>	1	0	6	1	1
s <sub>3</sub>	5	8	15	9	10
수요량	7	5	3	2	총합: 17



$$\text{총비용} = 2 \cdot 6 + 0 \cdot 1 + 5 \cdot 1 + 8 \cdot 4 + 15 \cdot 3 + 9 \cdot 2 = 112$$

2) 음환을 찾는다.

$$\begin{aligned} x_{12} \uparrow &= +3 - 2 + 5 - 8 = -2 < 0 \\ x_{13} \uparrow &= +11 - 2 + 5 - 15 = -1 < 0 \\ x_{14} \uparrow &= +7 - 2 + 5 - 9 = 1 > 0 \\ x_{21} \uparrow &= +1 - 5 + 8 - 0 = 4 > 0 \\ x_{23} \uparrow &= +6 - 0 + 8 - 15 = -1 < 0 \\ x_{24} \uparrow &= +1 - 0 + 8 - 9 = 0 \end{aligned}$$

$x_{12}$  진입

3) 개선된 해

공급지 \ 수요지	d <sub>1</sub>	d <sub>2</sub>	d <sub>3</sub>	d <sub>4</sub>	공급량
s <sub>1</sub>	2	3	11	7	6
s <sub>2</sub>	1	0	6	1	1
s <sub>3</sub>	5	8	15	9	10
수요량	7	5	3	2	총합: 17

$$\text{총비용} = 2 \cdot 2 + 3 \cdot 4 + 0 \cdot 1 + 5 \cdot 5 + 15 \cdot 3 + 9 \cdot 2 = 104$$

# Diet Problem

## ■ Diet

여러 가지 영양분을 지닌 음식들로부터 필수 영양분을 최소의 음식값으로 섭취하는 문제

[예] 주위에서 흔히 볼 수 있는 음식 재료에 포함된 영양분이 다음 표에 정리되어 있다. 최소의 비용으로 식단을 마련해 보고자 한다. 단 하루의 식단에서 쌀은 20포, 쇠고기는 1근, 우유는 2통, 계란은 3개, 배추는 3단을 넘지 않기로 한다.

영양 \ 재료	쌀(포)	쇠고기(근)	우유(통)	계란(12 개)	배추(단)	1 일 필요량
열량(Kcal)	340	1080	362	1040	17	2200
단백질(g)	6.5	167	19	78	1.3	70
비타민(I.U)	0	97	758	7080	255	5000
철분(mg)	0.4	11	0.3	13	0.3	12.5
탄수화물(g)	52	30	25	0	5	
콜레스테롤(u)	0	22	11	120	0	
값(원)	75	1640	370	550	110	

$$\text{Minimize } 75x_1 + 1640x_2 + 370x_3 + 550x_4 + 110x_5$$

$$\text{s.t. } 340x_1 + 1080x_2 + 362x_3 + 1040x_4 + 17x_5 \geq 2200$$

$$6.5x_1 + 167x_2 + 19x_3 + 78x_4 + 1.3x_5 \geq 70$$

$$97x_2 + 758x_3 + 7080x_4 + 255x_5 \geq 5000$$

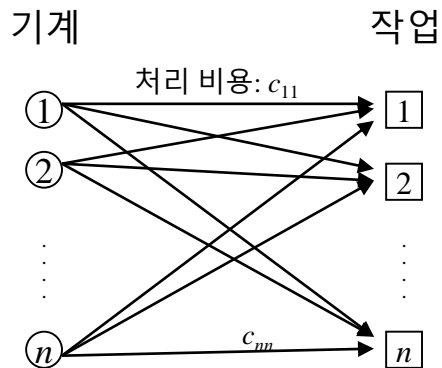
$$0.4x_1 + 11x_2 + 0.3x_3 + 13x_4 + 0.3x_5 \geq 12.5$$

$$x_1 \leq 20, \quad x_2 \leq 1, \quad x_3 \leq 2, \quad x_4 \leq 0.25, \quad x_5 \leq 3$$

$$x_{ij} \geq 0, \quad \forall i, j$$

# Assignment Problem

배정 문제(Assignment Problem):  $n$ 개의 작업을  $n$ 개의 기계에 각기 하나씩 최소의 비용이 되도록 할당하는 문제  
공급량과 수요량이 각기 1씩 발생하는 특별한 경우의 수송문제



만약 기계  $i$ 가 작업  $j$ 에 할당되면  $x_{ij} = 1$   
아니면  $x_{ij} = 0$

$$\text{Minimize } \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij}$$

$$\begin{aligned} \text{s.t. } & \sum_{j=1}^n x_{ij} = 1, \quad i = 1, 2, \dots, n \\ & \sum_{i=1}^n x_{ij} = 1, \quad j = 1, 2, \dots, n \\ & x_{ij} \geq 0, \quad \forall i, j \end{aligned}$$



# Knapsack Problem

## ■ Knapsack Problem

- 한정된 배낭의 용량에 맞게 각각의 용량을 가지는 물건을 최대의 효용을 얻도록 채우는 문제
- 대표적인 정수계획법 문제 중의 하나

[예] 갑돌이는 등산을 계획하고 있는데, 가면서 먹을 음식을 결정해야만 한다. 배당에는 총 1.6 kg까지만 음식을 담기로 결정했다고 한다. 각각의 음식의 무게와 그 음식을 가져 감으로써 얻을 수 있는 만족도가 다음과 같을 때, 가장 큰 만족도를 얻을 수 있는 음식의 조합을 결정하시오.

물건	고기	쌀	라면	과일	빵	
만족도	20	48	14	18	20	배낭의 무게
무게(100g)	8	6	2	3	2	16

$$\text{Maximize } 20x_1 + 48x_2 + 14x_3 + 18x_4 + 20x_5$$

$$\text{s.t.} \quad 8x_1 + 6x_2 + 2x_3 + 3x_4 + 2x_5 \leq 16$$

$$x_i = 1 \text{ or } 0, \forall i$$

## ORTools `!pip install ortools`

```
from ortools.linear_solver import pywraplp

def LinearProgrammingExample():
    """Linear programming sample."""
    # Instantiate a Glop solver, naming it LinearExample.
    solver = pywraplp.Solver.CreateSolver('GLOP')

    # Create the two variables and let them take on any non-negative value.
    x1 = solver.NumVar(0, solver.infinity(), 'x1')
    x2 = solver.NumVar(0, solver.infinity(), 'x2')
    print('Number of variables =', solver.NumVariables())

    # Constraint 0: 4x1 + 3x2 <= 12.
    solver.Add(4*x1 + 3 * x2 <= 12)

    # Constraint 1: 2x1 + 5x2 <= 10.
    solver.Add(2*x1 + 5*x2 <= 10)

    print('Number of constraints =', solver.NumConstraints())

    # Objective function: 12x1 + 15x2.
    solver.Maximize(12 * x1 + 15 * x2)

    # Solve the system.
    status = solver.Solve()

    # print_solution
    if status == pywraplp.Solver.OPTIMAL:
        print('Solution:')
        print('Objective value =', solver.Objective().Value())
        print('x1 =', x1.solution_value())
        print('x2 =', x2.solution_value())
    else:
        print('The problem does not have an optimal solution.')

LinearProgrammingExample()
```

# Home work

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- If  $f$  is convex, prove that  $q$  is concave