

# Consensus mechanism of HRD chain based on perpetual combinatorial auctions

#### Section one. Introduction

NFT (or heterogeneous token) is an encryption tool that can be used to token unique digital items. NFT is irreplaceable, which means that they can not be replaced by the same goods, and it can be proved that they are scarce, which makes them have intrinsic value.

As NFTs become more and more financialized, they will also need new trading platforms, lending agreements and derivatives. Therefore, the price discovery of NFT will be the next problem to be solved. Efficient price discovery mechanism in the capital market will enable participants to trade more quickly, improve liquidity through token, allow NFT to become collateral without order book, and create rich derivatives with NFT as the target. In other words, price discovery will make the financialization of NFT asset classes possible. Nowadays, there are only several ways to find the price of NFT, such as historical transaction pricing, auction pricing, fragmented pricing and so on. At present, most sellers and buyers in the market prefer to buy NFT through auction and price it. Async.art uses perpetual auctions in their website galleries, and so does superrare. Beeple's auction, with a total turnover of 3.5 million US dollars, proved that although the auction is still in the development stage, it is strong enough. It is worth mentioning that the auction is very helpful for the sale of artworks, because the intrinsic value of NFT assets is often more subjective, and there will be more people in the wait-and-see state. Therefore, we propose the perpetual combinatorial auction model to solve the pricing problem of NFT, which allows the continuous addition of new NFTs in the ore pool. At the same time, we propose the POA (Proof of Aution) mechanism, which transforms the auction value of NFT into the mineral value generated by the block chain mining, and realizes the standardization of NFT assets, so as to solve the two basic problems of difficult circulation and pricing of NFT assets.



# Section two. NFT pricing in perpetual combinatorial auctions

$S_i$	Mine price on the $i$ -th day
$N_{ij}$	The maximum ore output of the $j$ -th $NFT$ single mining on
	the $i-$ th day
$Q_{ij}$	The ore output of the $j$ -th $NFT$ combined mining on the
	i —th day
$V_{ij}$	The value of the $j-$ th $N\!FT$ on the $i-$ th day
$Q_{i,j_1j_\kappa}$	The ore output of these n $NFT$ $j_1,,j_n$ combinations combined
	mining on the $i$ —th day $(j_1 < j_2 < \cdots < j_n)$
$V_{i,j_1j_s}$	The value of these n $NFT$ $j_1,,j_n$ combinations on the $i$ -th
	$day \ (j_1 < j_2 < \dots < j_n)$
$k_{i}$	Number of $NFT$ on the $i$ —th day
$P_{ij}$	The $j-$ th $N\!FT$ auction price on the $i-$ th day
$P_{i,j_1j_s}$	The price of these n $NFT$ $j_1,,j_n$ combinations on the $i$ -th
	$day \ (j_1 < j_2 < \dots < j_n)$
Q	Total ore output per day
r	Fixed daily discount rate

Variation of NFT in ore pool:

The initial price of the j-th NFT digital asset on the i-th day is the auction price  $P_{ij}$ , which allows the NFT digital asset to mine in the public chain. According to the initial price  $P_{ij}$  of the j-th NFT digital asset, the calculation force is anchored, and according to the current mine price  $S_i$ , the output  $N_{ij}$  of the NFT digital asset can be determined. The daily output Q of the mine pool is constant. As time goes on, as shown in the figure above, there are new NFT digital assets



in the mine pool, so the weight of each NFT digital asset is decreasing, and the output  $Q_{ij}$  of the mine is also decreasing. When the output of a NFT digital asset is less than a certain threshold, the NFT digital asset exits the mine pool. The price of the mine  $S_i$  is changing in real time, so the price  $V_{ij}$  of each NFT digital asset is not constant. In the following two sections, the pricing of NFT digital assets with auction portfolio and without auction portfolio under the perpetual combinatorial auction model is given in two cases: without discount and with known fixed discount rate r.

## 1. NFT pricing without discounting

### (1) There is no auction portfolio

$$\label{eq:Theorem 1.1:} \text{ $Q_{ij}$} = \begin{cases} \frac{N_{ij}}{k_i} Q & \text{, } 1 \leq j \leq k_i \\ \sum\limits_{j=1}^{k_i} N_{ij} & \text{.} \\ 0 & \text{, } k_i < j \end{cases} .$$

Proof: When  $1 \le j \le k_i$ , the proportion of the j-th in the total NFT on the i-th day is

$$\frac{P_{ij}}{\sum_{j=1}^{k_i} P_{ij}} = \frac{\frac{P_{ij}}{S_i}}{\sum_{j=1}^{k_i} P_{ij}} = \frac{N_{ij}}{\sum_{j=1}^{k_i} N_{ij}}.$$

When  $k_{\scriptscriptstyle i} < j$  ,  $Q_{\scriptscriptstyle ij} = 0$  ,

so

$$Q_{ij} = \begin{cases} \frac{N_{ij}}{\sum_{j=1}^{k_i} N_{ij}} Q & , & 1 \le j \le k_i \\ \sum_{j=1}^{k_i} N_{ij} & . & . \\ 0 & , & k_i < j \end{cases}$$

$$\text{Theorem 1.2:} \quad V_{ij} = \begin{cases} \frac{N_{ij}}{k_i} QS_i & \text{,} \quad 1 \leq j \leq k_i \\ \sum\limits_{j=1}^{k_i} N_{ij} \\ 0 & \text{,} \quad k_i < j \end{cases}$$



Proof: We know from the theorem 1.1

$$Q_{ij} = \begin{cases} \frac{N_{ij}}{\sum_{j=1}^{k_i} N_{ij}} Q & , & 1 \le j \le k_i \\ \sum_{j=1}^{k_i} N_{ij} & , & \\ 0 & , & k_i < j \end{cases}$$

SO

$$V_{ij} = Q_{ij}S_i = egin{cases} rac{N_{ij}}{k_i} QS_i & , & 1 < j \le k_i \ \sum_{j=1}^{k_i} N_{ij} & & . \ 0 & , & k_i < j \end{cases}$$

$$\text{Theorem 1.3:} \quad V_j = \sum_{n=1}^{\infty} \sum_{i=n}^{\infty} \frac{N_{ij}}{\sum_{i=1}^{k_i} N_{ij}} QS_i \mathbf{1}_{(k_{n-1},k_n)}(j) \text{ , where } \mathbf{1}_{(k_{n-1},k_n)}(j) = \begin{cases} 1 \text{ , } & j \in (k_{n-1},k_n) \\ 0 \text{ , } & j \notin (k_{n-1},k_n) \end{cases} \text{, } k_0 = 0$$

Proof: We know from the theorem 1.2

$$V_{ij} = \begin{cases} \frac{N_{ij}}{k_i} QS_i &, & 1 \leq j \leq k_i \\ \sum_{j=1}^{k_i} N_{ij} & & \\ 0 &, & k_i < j \end{cases}$$

when  $j \in (k_{n-1}, k_n]$ ,

$$V_j = \sum_{i=1}^{\infty} V_{ij} = \sum_{i=n}^{\infty} V_{ij},$$

when  $j \notin (k_{n-1}, k_n]$ ,

$$V_i = 0$$
,

so

$$V_{j} = \sum_{n=1}^{\infty} \sum_{i=n}^{\infty} V_{ij} 1_{(k_{n-1},k_{n}]}(j) = \sum_{n=1}^{\infty} \sum_{i=n}^{\infty} \frac{N_{ij}}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j),$$

where 
$$1_{(k_{n-1},k_n]}(j) = \begin{cases} 1, & j \in (k_{n-1},k_n] \\ 0, & j \notin (k_{n-1},k_n] \end{cases}$$

Theorem 1.4:  $\lim_{i \to \infty} N_{ij} = 0$ .

 $\text{Proof: Because } N_i = \frac{P_{ij}}{S_i} \text{ , and because } P_{ij} \text{ is a bounded quantity and } \lim_{l \to \infty} S_i = +\infty,$ 

SO

$$\lim_{i \to \infty} N_{ij} = 0.$$



Theorem 1.5:  $\forall \varepsilon > 0, \{i \mid N_{ij} < \varepsilon\} \neq \emptyset$ 

Proof: We know from the theorem 1.4 and the definition of limit,  $\forall \varepsilon > 0$ ,  $\exists I = I(j, \varepsilon)$ , when i > I, we have

$$N_{ii} < \varepsilon$$
,

so

$$\{i \mid N_{ij} < \varepsilon\} \neq \emptyset.$$

Theorem 1.6: Set threshold  $\varepsilon > 0$ , when  $N_{ij} < \varepsilon$ , it is considered that the j-th NFT exits the ore pool at this time, let  $\min \left\{ i \mid N_{ij} < \varepsilon \right\} = T(j,\varepsilon)$ , then

$$V_{j} = \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,c)} \frac{N_{ij}}{\sum_{i=1}^{k_{j}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n})}(j)$$

where  $k_0 = 0$ .

Proof: When  $N_{ij} < \epsilon$ , it is considered that the j-th NFT exits the ore pool at this time,

here  $N_{ij} = 0$ .

when  $i > T(j, \varepsilon)$ , we have

$$\begin{split} V_{j} &= \sum_{i=n}^{\infty} V_{ij} = \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} V_{ij} 1_{(k_{n-1},k_{n}]}(j) + \sum_{n=1}^{\infty} \sum_{i=T(j,\varepsilon)+1}^{\infty} V_{ij} 1_{(k_{n-1},k_{n}]}(j) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} \frac{N_{ij}}{\sum_{i=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) + \sum_{n=1}^{\infty} \sum_{i=T(j,\varepsilon)-1}^{\infty} \frac{N_{ij}}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} \frac{N_{ij}}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) + \sum_{n=1}^{\infty} \sum_{i=T(j,\varepsilon)-1}^{\infty} \frac{0}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} \frac{N_{ij}}{\sum_{i=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) \end{split}$$

### (2) There is an auction portfolio

$$\text{Theorem 2.1:} \quad Q_{i,j_1\dots j_n} = \begin{cases} \frac{N_{i,j_1\dots j_n}}{N_{i,j_1\dots j_n}} Q & \text{,} & 1 \leq j_n \leq k_i \\ N_{i,j_1\dots j_n} + \sum_{\substack{j=1 \\ j \neq j_1\dots j_n}}^{k_i} N_{ij} & \\ 0 & \text{,} & j_n > k_i \end{cases}$$

Proof: When  $1 \le j_n \le k_i$ , and  $j_1 ... j_n$  Combinatorial auction,  $P_{i,j_1...j_n} < \sum_{k=1}^n P_{ij_k}$ , so the proportion of



corresponding  $\mathit{NFT}_{j_1 \dots j_\kappa}$  in total ore production Q on day i is

$$\frac{P_{i,j_{1}...j_{n}}}{p_{i,j_{1}...j_{n}} + \sum\limits_{\substack{j=1\\j \neq j_{1}...j_{n}}}^{k_{i}} p_{ij}} = \frac{\frac{P_{i,j_{1}...j_{n}}}{S_{i}}}{p_{i,j_{1}...j_{n}} + \sum\limits_{\substack{j=1\\j \neq j_{1}...j_{n}}}^{k_{i}} p_{ij}} = \frac{N_{i,j_{1}...j_{n}}}{N_{i,j_{1}...j_{n}} + \sum\limits_{\substack{j=1\\j \neq j_{1}...j_{n}}}^{k_{i}} N_{ij}}$$

 $\text{ when } j_n > k_i \text{ , } j_1 \cdots j_n \text{ does not exist at this time , so} \frac{P_{i,j_1 \dots j_n}}{p_{i,j_1 \dots j_n} + \sum\limits_{\substack{j=1 \\ j \neq h \dots j_n}}^{k_i} p_{ij}} = 0$ 

$$\text{Hence } Q_{i,j_1\dots j_n} = \begin{cases} \frac{N_{l,j_1\dots j_n}}{N_{l,j_1\dots j_n}} Q & \text{, } 1 \leq j_n \leq k_i \\ N_{i,j_1\dots j_n} + \sum\limits_{\substack{j=1 \\ j \neq j_1\dots j_n}}^{k_i} N_{ij} & \\ 0 & \text{, } j_n > k_i \end{cases}$$

$$\text{Theorem 2.2:} \quad V_{i,j_{1}\dots j_{n}} = \begin{cases} \frac{N_{i,j_{1}\dots j_{n}}}{N_{i,j_{1}\dots j_{n}} + \sum\limits_{\substack{j=1\\j \neq j_{1}\dots j_{n}}}^{k_{i}} N_{ij}} QS_{i} & , & 1 \leq j_{n} \leq k_{i} \\ \\ 0 & , & j_{n} > k_{i} \end{cases}$$

Proof: We know from the theorem 2.1

$$Q_{i,j_1\dots j_n} = \begin{cases} \frac{N_{i,j_1\dots j_n}}{N_{i,j_1\dots j_n}} & Q & , & 1 \leq j_n \leq k_i \\ N_{i,j_1\dots j_n} + \sum_{\substack{j=1 \\ j \neq j_1\dots j_n}}^{k_i} N_{ij} & & \\ 0 & & , & j_n > k_i \end{cases}$$

$$\text{Hence } V_{i,j_1\ldots j_n} = Q_{i,j_1\ldots j_n}S_i = \begin{cases} \frac{N_{i,j_1\ldots j_n}}{N_{i,j_1\ldots j_n}}QS_i &, & 1 \leq j_n \leq k_i \\ N_{i,j_1\ldots j_n} + \sum\limits_{\substack{j=1\\j \neq j_1\ldots j_n}}^{k_i}N_{ij} & \\ 0 & , & j_n > k_i \end{cases}$$

Theorem 2.3:

$$V_{j_1...j_n} = \sum_{m=1}^{\infty} \sum_{i=m}^{\infty} \frac{N_{i,j_1...j_n}}{N_{i,j_1...j_n} + \sum_{\substack{j=1 \ j \neq j_1...j_n}}^{k_i} N_{ij}} QS_i 1_{(k_{m-1},k_m]}(j_n)$$

where  $k_0 = 0$ .

Proof: We know from the theorem 2.2



$$V_{i,j_1\dots j_n} = \begin{cases} \frac{N_{i,j_1\dots j_n}}{N_{i,j_1\dots j_n}} QS_i &, \quad 1 \leq j_n \leq k_i \\ N_{i,j_1\dots j_n} + \sum_{\substack{j=1 \\ j \neq j_1\dots j_n}}^{k_i} N_{ij} & \\ 0 &, \quad j_n > k_i \end{cases}$$

SO

$$V_{j_1...j_n} = \sum_{m=1}^{\infty} \sum_{i=m}^{\infty} \frac{N_{i,j_1...j_n}}{N_{i,j_1...j_n} + \sum_{\substack{j=1 \ j \neq j_1...j_n}}^{k_i} N_{ij}} QS_i 1_{(k_{m-1},k_{m})}(j_n)$$

Theorem 2.4:  $\lim_{i\to\infty} N_{i,j_1\dots j_n}=0$ 

Proof: Because

$$N_{i,j_1...j_n} = \frac{P_{i,j_1...j_n}}{S_i}$$

And because  $P_{i,j_1...j_s}$  is a bounded quantity and

$$\lim_{i\to\infty} S_i = +\infty \quad \to \quad \lim_{i\to\infty} \frac{1}{S_i} = 0,$$

so

$$\lim_{l o \infty} N_{i,j_1...j_n} = \lim_{l o \infty} rac{P_{l,j_1...j_n}}{S_i} = 0.$$

Theorem 2.5:  $\forall \epsilon > 0, \{i \mid N_{i,j_1...j_s} < \epsilon\} \neq \emptyset.$ 

Proof: We know from the theorem 1.3 and the definition of limit,  $\forall \varepsilon > 0$ ,  $\exists I = I(j_1...j_n, \varepsilon)$ , when i > I, we have

$$N_{ii} < \varepsilon$$
,

so

$$\{i \mid N_{i,j_1...j_n} < \varepsilon\} \neq \emptyset.$$

Theorem 2.6: Set threshold  $\varepsilon > 0$ , let  $\min \left\{ i \mid N_{i,j_1...j_n} < \varepsilon \right\} = T(j_1...j_n,\varepsilon)$ , when  $N_{i,j_1...j_n} < \varepsilon$ , it is considered that the  $j_1,...,j_n$  NFT exits the ore pool at this time, then

where  $k_0 = 0$ .

 $\text{Proof:} \quad \text{when } N_{i,j_1,\dots,j_n} < \epsilon \text{ , it is considered that the } \quad j_1,\dots,j_n \quad NFT \quad \text{exits the ore pool at this time , }$ 

here  $N_{i,j_1\dots j_n}=0$  ,  $% i=1,\dots,j_n,$  when  $i>T(j_1\dots j_n,\epsilon)$  , we have



$$\begin{split} V_{j_{1}\dots j_{n}} &= \sum_{i=n}^{\infty} V_{i,j_{1}\dots j_{n}} = \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} V_{i,j_{1}\dots j_{n}} 1_{(k_{n-1},k_{n}]}(j_{n}) + \sum_{n=1}^{\infty} \sum_{i=T(j_{1}\dots j_{n},\varepsilon)+1}^{\infty} V_{i,j_{1}\dots j_{n}} 1_{(k_{n-1},k_{n}]}(j_{n}) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} \frac{N_{i,j_{1}\dots j_{n}}}{\sum_{j=1}^{k_{i}} N_{i,j_{1}\dots j_{n}}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) + \sum_{n=1}^{\infty} \sum_{i=T(j_{1}\dots j_{n},\varepsilon)+1}^{\infty} \frac{N_{ij}}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} \frac{N_{i,j_{1}\dots j_{n}}}{\sum_{j=1}^{k_{i}} N_{i,j_{1}\dots j_{n}}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) + \sum_{n=1}^{\infty} \sum_{i=T(j_{1}\dots j_{n},\varepsilon)+1}^{\infty} \frac{0}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} \frac{N_{i,j_{1}\dots j_{n}}}{\sum_{j=1}^{k_{i}} N_{i,j_{1}\dots j_{n}}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) \end{split}$$

# 2. NFT pricing with discount

## (1) There is no auction portfolio

$$\text{Theorem 3.1:} \quad V_{ij} = \begin{cases} \frac{N_{ij}}{k_i} QS_i e^{-ir} &, \quad 1 \leq j \leq k_i \\ \sum_{j=1}^{k} N_{ij} & \\ 0 &, \quad k_i < j \end{cases}$$

Proof: When  $1 \le j \le k_i$  , the proportion of the j —th in the total NFT on the i —th day is

$$\frac{P_{ij}}{\sum_{j=1}^{k_i} P_{ij}} = \frac{\frac{P_{ij}}{S_i}}{\sum_{j=1}^{k_i} P_{ij}} = \frac{N_{ij}}{\sum_{j=1}^{k_i} N_{ij}}.$$

When  $k_{\scriptscriptstyle i} < j$  ,  $V_{\scriptscriptstyle ij} = 0$  ,

so

$$V_{ij} = \begin{cases} \frac{N_{ij}}{k_i} QS_i e^{-ir} &, & 1 < j \le k_i \\ \sum_{j=1}^{k_i} N_{ij} &, & k_i < j \end{cases}$$



$$\text{Theorem 3.2:} \quad V_j = \sum_{n=1}^{\infty} \sum_{i=n}^{\infty} \frac{N_{ij}}{\sum_{i=1}^{k_i} N_{ij}} QS_i \mathbf{1}_{(k_{n-1},k_n]}(j) e^{-ir} \text{ , where } \mathbf{1}_{(k_{n-1},k_n]}(j) = \begin{cases} 1 &, & j \in (k_{n-1},k_n] \\ 0 &, & j \not\in (k_{n-1},k_n] \end{cases} \text{, } k_0 = 0$$

Proof: We know from the theorem 3.1

$$V_{ij} = \begin{cases} \frac{N_{ij}}{k_i} Q S_i e^{-ir} &, & 1 \le j \le k_i \\ \sum_{j=1}^{k_i} N_{ij} & & \\ 0 &, & k_i < j \end{cases}$$

when  $j \in (k_{n-1}, k_n]$ ,

$$V_j = \sum_{i=1}^{\infty} V_{ij} = \sum_{i=n}^{\infty} V_{ij},$$

when  $j \notin (k_{n-1}, k_n]$ ,

$$V_{i} = 0,$$

so

$$V_{j} = \sum_{n=1}^{\infty} \sum_{i=n}^{\infty} V_{ij} 1_{(k_{n-1},k_{n}]}(j) = \sum_{n=1}^{\infty} \sum_{i=n}^{\infty} \frac{N_{ij}}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) e^{-ir},$$

where  $1_{(k_{n-1},k_n]}(j) = \begin{cases} 1, & j \in (k_{n-1},k_n] \\ 0, & j \notin (k_{n-1},k_n] \end{cases}$ 

Theorem 3.3:  $\lim_{l\to\infty} N_{ij} = 0$ .

Proof: Because  $N_i=rac{P_{ij}}{S_i}$  , and because  $P_{ij}$  is a bounded quantity and  $\lim_{i o\infty}S_i=+\infty$  ,

$$\lim_{i \to \infty} N_{ij} = 0.$$

Theorem 3.4:  $\forall \varepsilon > 0, \{i \mid N_{ij} < \varepsilon\} \neq \emptyset.$ 

Proof: We know from the theorem 3.3 and the definition of  $[\lim_{t\to 0}] I = I(j,\epsilon)$ , when i>I, we have

$$N_{ij} < \varepsilon$$
,

so

$$\{i \mid N_{ij} < \varepsilon\} \neq \emptyset.$$

Theorem 3.5: Set threshold  $\varepsilon > 0$ , when  $N_{ij} < \varepsilon$ , it is considered that the j-th NFT exits the ore pool at this time, let  $\min \left\{ i \mid N_{ij} < \varepsilon \right\} = T(j,\varepsilon)$ , then



$$V_{j} = \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,e)} \frac{N_{ij}}{\sum_{i=1}^{k_{j}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n})}(j)e^{-ir}$$

where  $k_0 = 0$ .

Proof: When  $N_{ij} < \epsilon$  , it is considered that the j-th NFT exits the ore pool at this time,

here  $N_{ij} = 0$ .

when  $i > T(j, \varepsilon)$ , we have

$$\begin{split} V_{j} &= \sum_{i=n}^{\infty} V_{ij} = \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} V_{ij} 1_{(k_{n-1},k_{n}]}(j) + \sum_{n=1}^{\infty} \sum_{i=T(j,\varepsilon)+1}^{\infty} V_{ij} 1_{(k_{n-1},k_{n}]}(j) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} \frac{N_{ij}}{\sum_{k_{i}}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) e^{-ir} + \sum_{n=1}^{\infty} \sum_{i=T(j,\varepsilon)+1}^{\infty} \frac{N_{ij}}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) e^{-ir} \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} \frac{N_{ij}}{\sum_{i=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) e^{-ir} + \sum_{n=1}^{\infty} \sum_{i=T(j,\varepsilon)+1}^{\infty} \frac{0}{\sum_{j=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) e^{-ir} \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j,\varepsilon)} \frac{N_{ij}}{\sum_{i=1}^{k_{i}} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j) e^{-ir} \,. \end{split}$$

## (2) There is an auction portfolio

$$\text{Theorem 4.1:} \quad V_{i,j_1\dots j_n} = \begin{cases} \frac{N_{i,j_1\dots j_n}}{N_{i,j_1\dots j_n}} QS_i e^{-ir} & \text{,} \quad 1 \leq j_n \leq k_i \\ N_{i,j_1\dots j_n} + \sum_{\substack{j=1 \\ j \neq j_1\dots j_n}}^{k_i} N_{ij} & \\ 0 & \text{,} \quad j_n > k_i \end{cases}$$

 $\text{Proof: When } 1 \leq j_n \leq k_i \text{ , and } j_1...j_n \text{ Combinatorial auction, } P_{i,j_1...j_n} < \sum_{k=1}^n P_{j_k} \text{ , so the proportion of corresponding } NFT_{j_1...j_n} \text{ in total ore production } Q \text{ on day } i \text{ is }$ 

$$\frac{P_{i,j_{1}...j_{n}}}{p_{i,j_{1}...j_{n}} + \sum_{\substack{j=1\\j \neq j_{1}...j_{n}}}^{k_{i}} p_{ij}} = \frac{\frac{P_{i,j_{1}...j_{n}}}{S_{i}}}{p_{i,j_{1}...j_{n}} + \sum_{\substack{j=1\\j \neq j_{1}...j_{n}}}^{k_{i}} p_{ij}} = \frac{N_{i,j_{1}...j_{n}}}{N_{i,j_{1}...j_{n}} + \sum_{\substack{j=1\\j \neq j_{1}...j_{n}}}^{k_{i}} N_{ij}}$$

when  $j_n > k_i$  ,  $j_1 \cdots j_n$  does not exist at this time, so  $\frac{P_{i,j_1 \cdots j_n}}{p_{i,j_1 \cdots j_n} + \sum\limits_{\substack{j=1 \ i \neq k, \dots, i}}^{k_i} p_{ij}} = 0.$ 



$$\text{Hence} \quad V_{i,j_1\dots j_n} = \begin{cases} \frac{N_{i,j_1\dots j_n}}{N_{i,j_1\dots j_n}} QS_i e^{-ir} &, \quad 1 \leq j_n \leq k_i \\ N_{i,j_1\dots j_n} + \sum\limits_{\substack{j=1\\j \neq j_1\dots j_n}}^{k_j} N_{ij} &\\ 0 &, \quad j_n > k_i \end{cases}$$

Theorem 4.2:

$$V_{j_1...j_n} = \sum_{m=1}^{\infty} \sum_{i=m}^{\infty} \frac{N_{i,j_1...j_n}}{N_{i,j_1...j_n} + \sum\limits_{\substack{j=1 \ j 
eq j_1...j_n}}^{k_i} N_{ij}} QS_i \, 1_{(k_{m-1},k_m)} (j_n) e^{-ir}$$

where  $k_0 = 0$ .

Proof: We know from the theorem 4.1

$$V_{i,j_{1}\dots j_{n}} = \begin{cases} \frac{N_{i,j_{1}\dots j_{n}}}{N_{i,j_{1}\dots j_{n}}} + \sum_{\substack{j=1\\j\neq j_{1}\dots j_{n}}}^{k_{i}} N_{ij}} QS_{i}e^{-ir} &, & 1 \leq j_{n} \leq k_{i} \\ \\ 0 &, & j_{n} > k_{i} \end{cases}$$

so

$$V_{j_1...j_n} = \sum_{m=1}^{\infty} \sum_{i=m}^{\infty} rac{N_{i,j_1...j_n}}{N_{i,j_1...j_n} + \sum\limits_{\substack{j=1\j
 j,1...j_n}}^{k_i} N_{ij}} QS_i \, 1_{(k_{m-1},k_m)}(j_n) e^{-ir}$$

Theorem 4.3:  $\lim_{i\to\infty} N_{i,j_1...j_n} = 0$ 

Proof: Because

$$N_{i,j_1...j_n} = \frac{P_{i,j_1...j_n}}{S_i}$$

And because  $P_{i,j_1\dots j_s}$   $P_{ij}$  is a bounded quantity and

$$\lim_{i\to\infty} S_i = +\infty \quad \to \quad \lim_{i\to\infty} \frac{1}{S_i} = 0,$$

so

$$\lim_{i\to\infty} N_{i,j_1\dots j_n} = \lim_{i\to\infty} \frac{P_{i,j_1\dots j_n}}{S_i} = 0.$$

Theorem 4.4:  $\forall \varepsilon > 0, \left\{ i \mid N_{i,j_1...j_n} < \varepsilon \right\} \neq \emptyset$ 

Proof: We know from the theorem 4.3 and the definition of limit ,  $\forall \varepsilon>0,\ \exists I=I(j_1...j_n,\varepsilon)$  , when i>I , we have

$$N_{ij} < \varepsilon_i$$

so

$$\{i \mid N_{i,j_1...j_n} < \varepsilon\} \neq \emptyset.$$

Theorem 4.5: Set threshold  $\varepsilon > 0$ , let  $\min \left\{ i \mid N_{i,j_1...j_n} < \varepsilon \right\} = T(j_1...j_n,\varepsilon)$ , when  $N_{i,j_1...j_n} < \varepsilon$ , it is



considered that the  $j_1,...,j_n$  NFT exits the ore pool at this time, then

$$V_{j_1...j_n} = \sum_{m=1}^{\infty} \sum_{i=m}^{T(j_1...j_n,\varepsilon)} rac{N_{i,j_1...j_n}}{N_{i,j_1...j_n} + \sum\limits_{\substack{j=1\j
eq j_1...j_n}}^{k_j}} QS_i \, 1_{(k_{m-1},k_m]}(j_n) e^{-ir}$$

where  $k_0 = 0$ .

Proof: when  $N_{i,j_1...j_n} < \epsilon$  , it is considered that the  $j_1,...,j_n$  NFT exits the ore pool at this time,

here  $N_{i,j_1...j_n}=0$  , when  $i>T(j_1...j_n,\varepsilon)$  , we have

$$\begin{split} V_{j_{1}\dots j_{n}} &= \sum_{i=n}^{\infty} V_{i,j_{1}\dots j_{n}} = \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} V_{i,j_{1}\dots j_{n}} 1_{(k_{n-1},k_{n}]}(j_{n}) + \sum_{n=1}^{\infty} \sum_{i=T(j_{1}\dots j_{n},\varepsilon)+1}^{\infty} V_{i,j_{1}\dots j_{n}} 1_{(k_{n-1},k_{n}]}(j_{n}) \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} \frac{N_{i,j_{1}\dots j_{n}}}{\sum_{j=1}^{k} N_{i,j_{1}\dots j_{n}}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) e^{-ir} + \sum_{n=1}^{\infty} \sum_{i=T(j_{1}\dots j_{n},\varepsilon)+1}^{\infty} \frac{N_{ij}}{\sum_{j=1}^{k} N_{ij}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) e^{-ir} \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} \frac{N_{i,j_{1}\dots j_{n}}}{\sum_{j=1}^{k} N_{i,j_{1}\dots j_{n}}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) e^{-ir} \\ &= \sum_{n=1}^{\infty} \sum_{i=n}^{T(j_{1}\dots j_{n},\varepsilon)} \frac{N_{i,j_{1}\dots j_{n}}}{\sum_{j=1}^{k} N_{i,j_{1}\dots j_{n}}} QS_{i} 1_{(k_{n-1},k_{n}]}(j_{n}) e^{-ir}. \end{split}$$

## Section three. POA consensus mechanism(Proof of Aution)

In the HRD pool, the NFT asset value is used as the anchor of pool calculation force, the amount of HRD produced per day is constant, which is halved with the block height. Each NFT asset



is priced in the perpetual combinatorial auction, and the weight  $\frac{N_{ij}}{\sum\limits_{j=1}^{k_i}N_{ij}}$  (1  $\leq j \leq k_i$ ) in the HRD

pool is obtained from the price. As long as the mortgage is successful, HRD  $Q_{ij} (= \frac{N_{ij}}{\sum\limits_{j=1}^{k_i} N_{ij}} Q)$  can

be produced according to the weight. In this way, NFT assets can be standardized through the output HDR, and users can replace NFT assets through HRD.

## Section four. Conclusion

In this paper, under the model of perpetual combinatorial auction, we give the NFT pricing formula with and without auction portfolio without discount. Furthermore, we give the NFT pricing formula with and with auction portfolio with discount, it solves the problem of NFT assets pricing. We further abstract the consensus mechanism of POA to realize the standardized pricing of NFT assets with HDR. Thus, the two problems of NFT assets circulation and pricing are well solved.