CS 314 Operating Systems Lab

Lab Assignment 3

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Part 1:

Q1] Modify the Minix3 source code such that the string “PID swapped in” is printed, whenever a user-level process is brought in by the scheduler.

A] We need to modify the MINIX source code inside the directory /usr/src/minix/servers/sched in this directory the main.c file calls the do_start_scheduling(), do_stop_scheduling(), do_nice() and do_noquantum() functions which is located in the schedule.c file in same directory. Whenever a user level process in brought in by the scheduler the do_start_scheduling() function is called which in turn calls the schedule_process() function in the same file. This function is globally defined with the return type as int. In the schedule_process() function we set the values for the different fields of the schedproc object which is rmp.

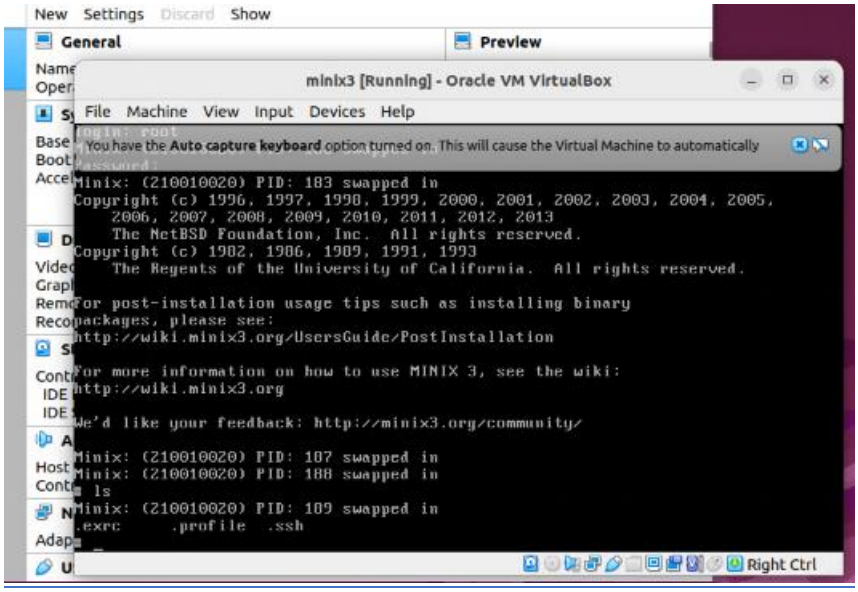
We make the following modification to the file schedule.c in order to print whenever a new process is swapped in:

```
if (rmp->priority >= USER_Q) :

        printf("Minix: (210010020) PID: %d swapped in\n", _ENDPOINT_P(rmp->endpoint));
```

The if statement checks if the priority of the process is greater than equal to the USER_Q(which is defined in minix/include/minix/config.h) .

In order to make the changes we need to run the run.sh file given along with this using command **bash run.sh**. After running the script the following output can be seen here.

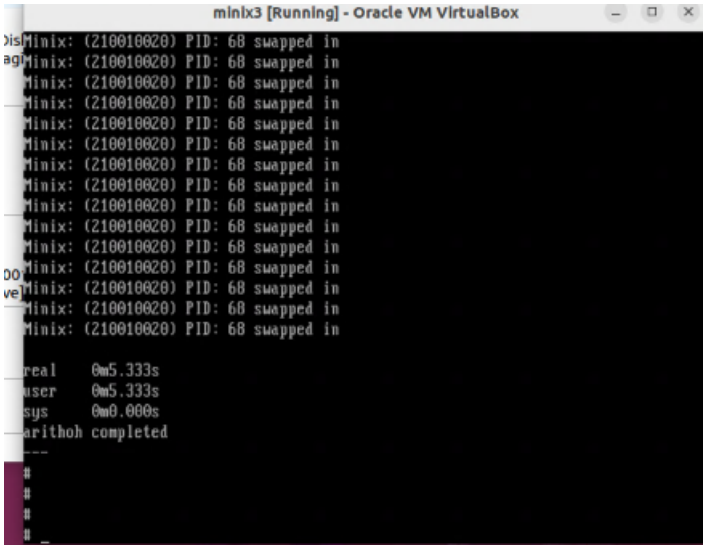


Part2

To classify the processes as I/o Bound or CPU bound we need to examine the **user and sys** times of the processes.

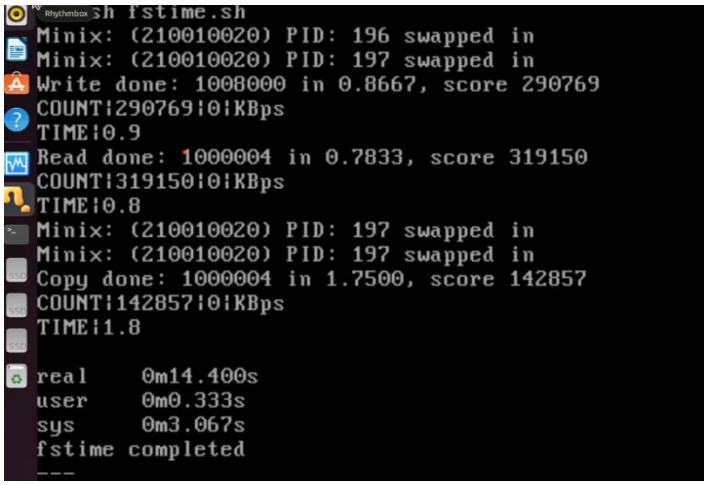
- CPU-Bound Process:
  - A process is considered CPU-bound if it spends a significant portion of its time using the CPU.
  - If the user time is significantly higher than the sys time, it suggests that the process is consuming a lot of CPU time.
- I/O-Bound Process:
  - A process is considered I/O-bound if it spends a significant portion of its time waiting for I/O operations (e.g., disk I/O, network I/O).
  - If the sys time is relatively higher than the user time, it might indicate that the process is spending more time in system mode, possibly waiting for I/O operations to complete.

1. arithoh.sh:



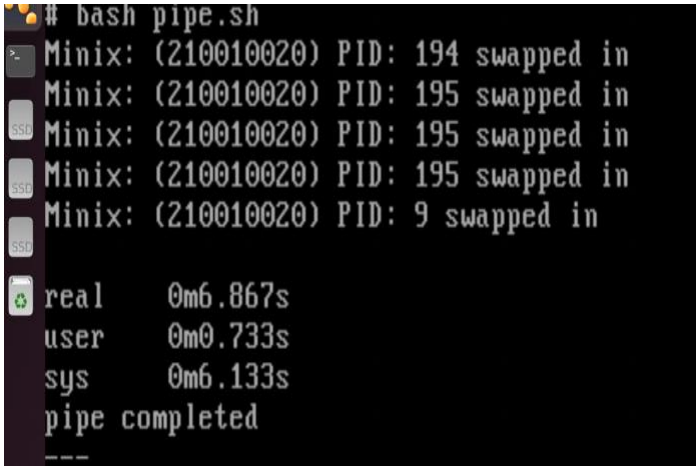
- The process with pid 68 runs when arithoh.sh script is executed.
- The process spends 5.33s in user mode and 0.00s in kernel mode, total wall clock time spent = 5.33s.
- This is CPU Bound Process.

2. fstime.sh



- The process with pid 196 runs when fstime.sh script is executed.
- The process spends 0.333s in user mode and 3.067s in kernel mode, total wall clock time spent = 14.400s.
- This is I/O Bound Process as it waits for the Input for long time.

3. pipe.sh

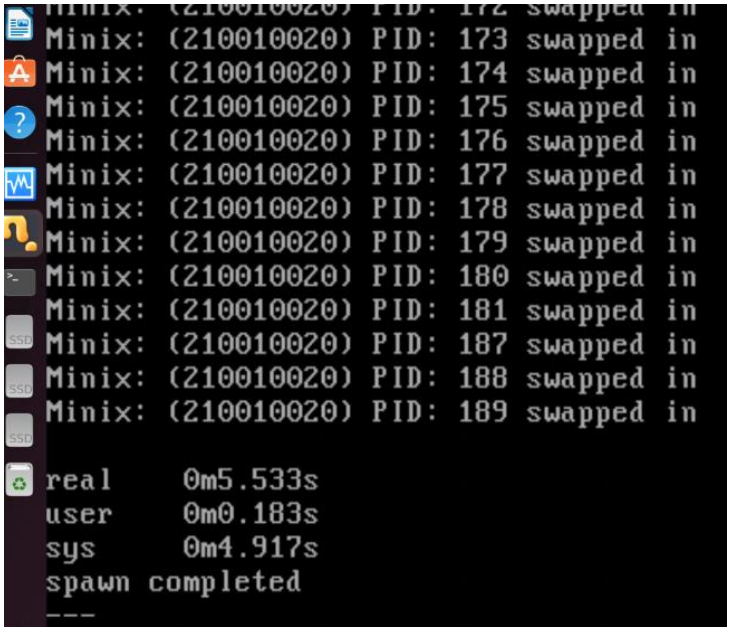


- The process with pid 194,195 and 9 runs when pipe.sh script is executed.
- The process spends 0.733s in user mode and 6.133s in kernel mode, total wall clock time spent = 6.867.

From the code, it appears to involve both CPU-bound and user-bound aspects.

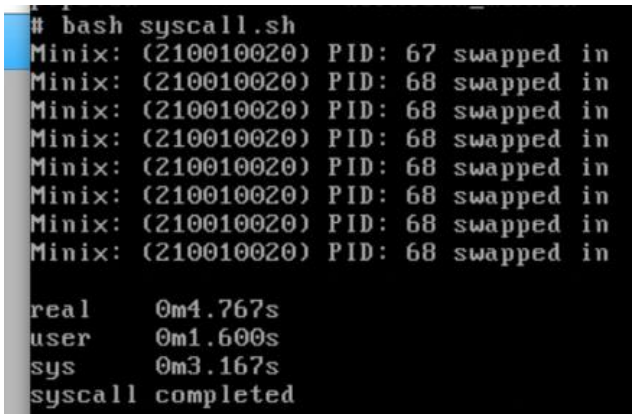
The CPU is actively engaged in the loop iterations, and the handling of interrupts indicates potential user-bound characteristics due to interactions with the operating system.

4. **spawn.sh:**



- The process with pid 172,173,174,175,176,177,178,179,180,181,187,188 ,189 runs when spawn.sh script is executed.
- The process spends 0.733s in user mode and 6.133s in kernel mode, total wall clock time spent = 6.867.
- Based on the analysis, the code is more likely to be CPU-bound. The main workload involves creating and terminating processes in a loop, which is a CPU-intensive operation. The absence of explicit I/O operations in the loop supports the classification as CPU-bound.

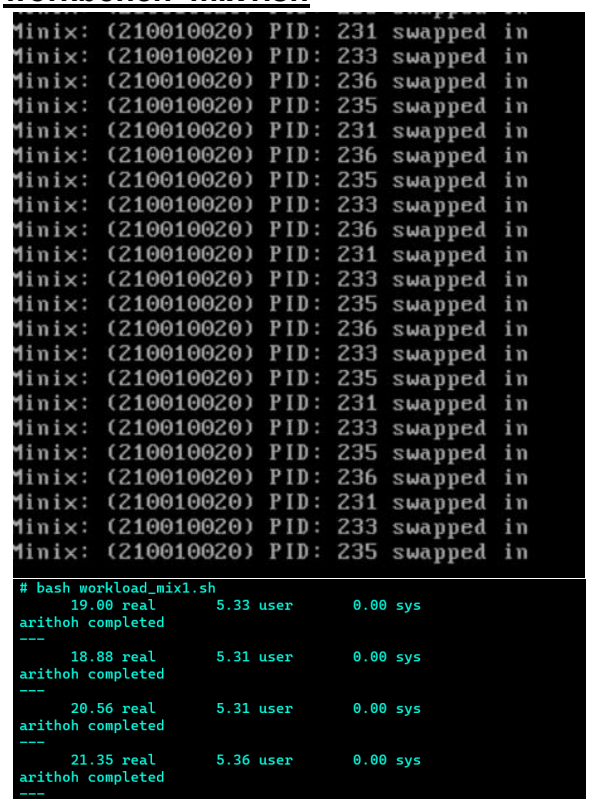
5. **syscall.sh**



- The process with pid 67 and 68 runs when syscall.sh script is executed.
- The process spends 1.600s in user mode and 3.167s in kernel mode, total wall clock time spent = 4.767s.

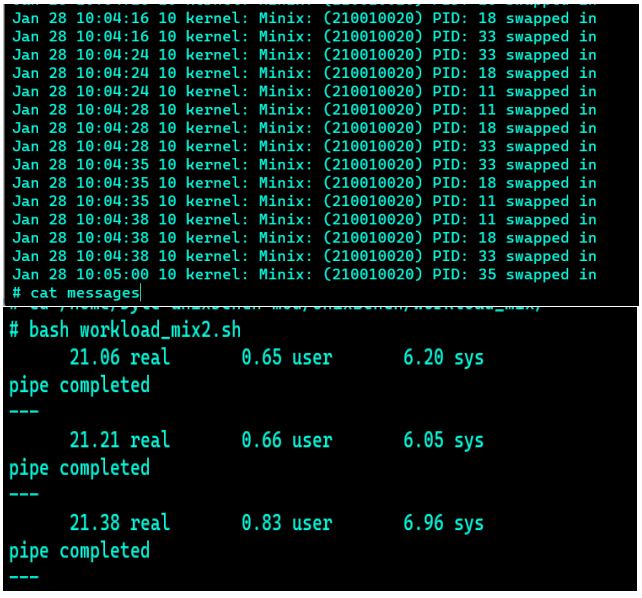
This is The code exhibits a mix of CPU-bound and user-bound characteristics, depending on the specific calls:"mix" includes a mix of CPU-bound and user-bound operations whereas close" and "getpid" are more user-bound due to the nature of the system calls. "exec" involves process creation and execution, which includes both CPU-bound and user-bound aspects.

6. workbench_mix1.sh



- As seen above arithoh.sh is CPU Bound process.
- Four instances of arithoh.sh run initiated run one after the other until their time slice is over.
- Processes 231,233,236 and 235 correspond to 4 instances of arithoh.sh.

7. workload_mix2.sh



- Each instance of pipe.sh is a I/O bound process as explained earlier.
- When executed process waits for the input meanwhile others who have the input use CPU in round robin manner.

8. workload_mix3.sh

```
Jan 28 11:22:47 10 kernel: Minix: (210010020) PID: 55 swapped in
Jan 28 11:22:48 10 kernel: Minix: (210010020) PID: 55 swapped in
Jan 28 11:22:48 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:48 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:48 10 kernel: Minix: (210010020) PID: 55 swapped in
Jan 28 11:22:48 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:49 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:49 10 kernel: Minix: (210010020) PID: 55 swapped in
Jan 28 11:22:49 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:49 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:49 10 kernel: Minix: (210010020) PID: 55 swapped in
Jan 28 11:22:49 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:50 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:50 10 kernel: Minix: (210010020) PID: 55 swapped in
Jan 28 11:22:50 10 kernel: Minix: (210010020) PID: 62 swapped in
Jan 28 11:22:51 10 last message repeated 6 times
Jan 28 11:23:01 10 kernel: Minix: (210010020) PID: 64 swapped in
```

```
# bash workload_mix3.sh
       7.51 real       0.43 user       6.41 sys
pipe completed
---
      22.98 real       5.23 user       0.00 sys
arithoh completed
---
      13.23 real       5.25 user       0.00 sys
arithoh completed
```

- When arithoh.h and pipe are initiated, one is CPU bound and other is I/O bound.
- While the I/O bound waits for an input from the user the cpu bound process uses the CPU.
- After pipe completes its execution it leaves.