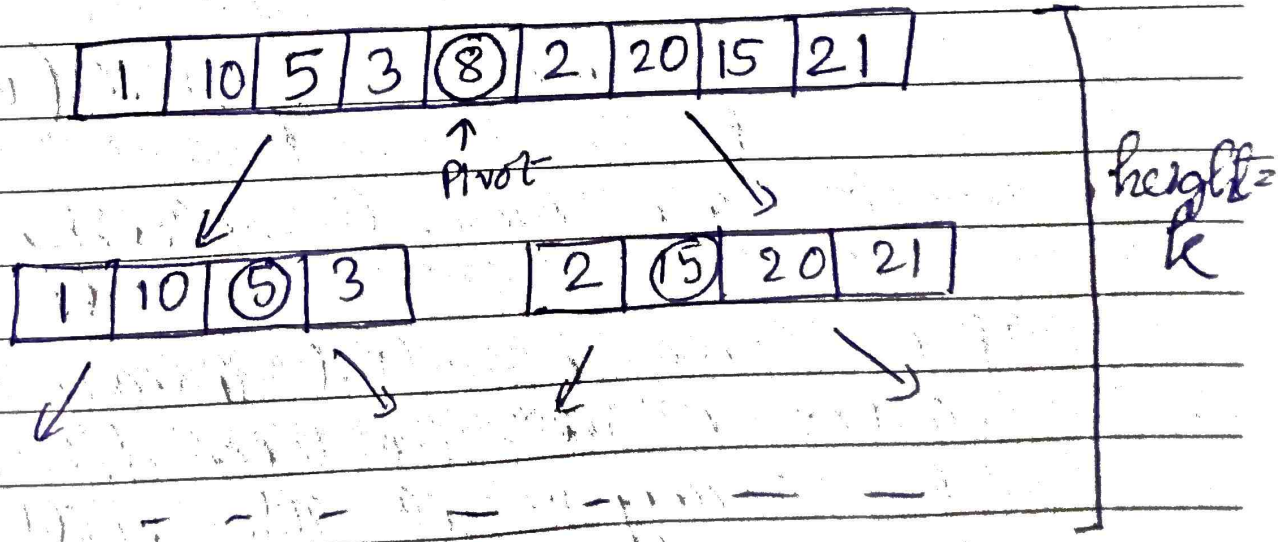
EC B19035

Quick Sort Analysis

Steps for quick sort:
→ partition the array & get a pivot.
→ Recursively call left & right part of pivot to sort itself

- **Best-Case Time Complexity** = $O(n \log n)$

If we assume that the pivot is always the median of the array then,



If the height of recursive tree is k & the array goes on to decrease its range by half, so we have to divide the array $n/2^k$ times & until only one element is left.

August 2015

21 Friday

(233 - 132) Wk 34

08:00

So, $\frac{n}{2^k} = 1$

09:00

$n = 2^k$

10:00

$\boxed{k = \log_2 n}$

11:00

compare & pivot

12:00

So, no. of times we do ~~partition~~ is = n

13:00

& we have $\log_2 n$ times division of array.

14:00

∴ Time Complexity = $O(n \log n)$

15:00

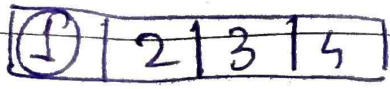
- Worst Case Time Complexity = $O(n^2)$

16:00

If we have a sorted array & the pivot is at the beginning of the array or at the end of the array. Then,

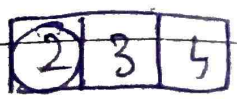
17:00

18:00



n times

19:00



n-1 "

Eve.

n-2 "

n-3 "

- Politeness costs nothing and gains everything

We know,

$$n + (n-1) + (n-2) + (n-3) + \cdots$$

$$> \frac{n(n+1)}{2}$$

$$= \frac{n^2+n}{2}$$

Here, the Time Complexity comes out to be $O(n^2)$.

- To Optimise the Worst Case of Quick Sort

  - Select Middle element as the pivot.

23 Sunday

  - Select Random element as pivot.

August 2015

24 Monday

(236 - 129) Wk 35

August

Su	Mo	Tu	We	Th	Fr	Sa
30	31	•	•	•	•	
2	3	4	5	6	7	
9	10	11	12	13	14	
16	17	18	19	20	21	
23	24	25	26	27	28	

08:00

Space Complexity

09:00

Quick Sort is a recursive algorithm so it doesn't require any extra space, but it will use Stack.

10:00

- In Worst case, quick sort may take stack size of n

11:00

- In Best case, quick sort may take stack size of $\log n$

12:00

So, it can vary from

13:00

$\log n \rightarrow n$

14:00

15:00

16:00

17:00

18:00

19:00

Eve.

- Power comes from sing