# Lecture 08 – Multiobjective Optimisation

**Optimisation CT5141** 

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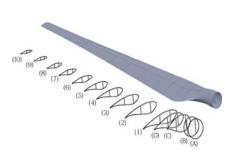


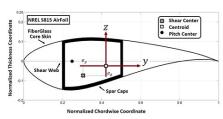
### **Overview**

- Applications
- Multi-objective optimisation algorithms

## Wind turbine engineering

In a wind turbine, what **shape** should the sails be?

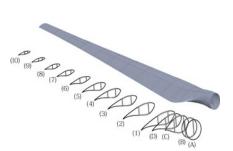


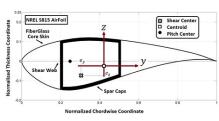


Sheibani and Akbari

## Wind turbine engineering

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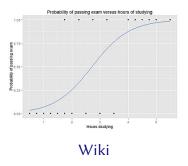


#### Sheibani and Akbari

We program a finite-element simulation to estimate the **annual energy production** for a given shape. We **maximise**.

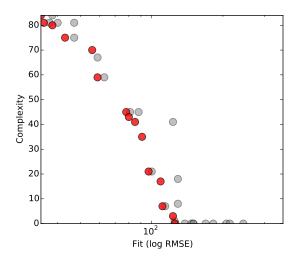
We also want to **minimise** construction cost.

## Regression with regularisation



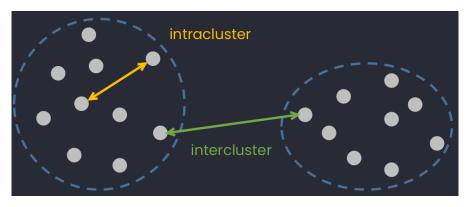
We want to **minimise** the error on the training data, but also **minimise** model complexity, i.e. use **regularisation**.

# **Genetic Programming symbolic regression**



Minimise regression error, minimise equation complexity.

# Clustering



https://dinhanhthi.com/metrics-for-clustering/

We want to **minimise** the sum of intra-cluster distances and **maximise** the sum of between-cluster distances.

## **Print Shop job selection**

- Busy print shop, many jobs arrive each morning
- Every job j gives some profit (value) v<sub>j</sub>, and requires time (weight) w<sub>j</sub>
- We can only work  $w_{\text{max}}$  hours
- Which jobs should we choose to max profit while working less than time limit?
- This is a **knapsack problem**

## Multi-objective Unit Commitment

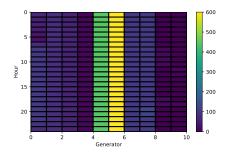
In a small country there are 10 generators of four types: Hydroelectric, Solid fuel, Gas, Solar. Each generator has a lower and upper bound on its production per hour (in MW/h).

		•	•	•	
	type	LB (MW)	UB (MW)	c (EUR/MW)	p (tons CO2/MW)
A	hydro	10	100	1.4	0.024
В	hydro	10	80	1.4	0.024
C	hydro	10	60	1.4	0.024
D	hydro	1	10	1.4	0.024
Ε	solid	100	900	4.4	0.82
F	solid	100	600	4.4	0.82
G	solid	10	100	4.4	0.82
Η	gas	100	400	9.1	0.49
I	solar	0	70	6.6	0
J	solar	0	20	6.6	0

## Minimising cost

Minimise  $\sum_{i,j} c_j X_{i,j}$ 

Subject to meeting demand, *X* within lower/upper bounds, etc.



Cost: EUR151518

■ Emissions: 23810 tons CO2

## Alternative: minimising emissions

Minimise  $\sum_{i,j} c_j X_{i,j}$ 

■ Cost: EUR151518

■ Emissions: 23810 tons CO2

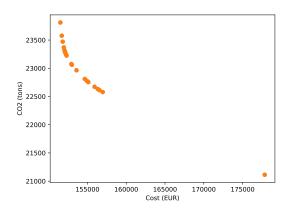
Minimise  $\sum_{i,j} p_j X_{i,j}$ 

■ Cost: EUR177846

■ Emissions: 21113 tons CO2

## A spectrum of solutions

We have considered **only cost**, and then **only emissions**. We really want both! Remember, lower is better for both objectives:



## **Conflicting objectives**

Multi-objective problems arise when we want to maximise or minimise **multiple objectives** 

In particular: multiple conflicting objectives

## Regression

- $\blacksquare$  Maximise  $R^2$
- Minimise RMSE

This are **not conflicting objectives** 

We **don't need** multi-objective optimisation.

#### **Overview**

- Applications
- 2 Multi-objective optimisation algorithms

## Types of algorithms

#### Weighting schemes:

- Choose a weight for each objective
- Then apply single-objective optimisation

#### Simple MOO algorithms:

- Pareto archive
- Random-objective tournament selection

#### Complex MOO algorithms:

NSGA2 and similar

## Weighting schemes

Suppose we have two objectives, maximise  $f_1$  and maximise  $f_2$ .

Define some weights  $w_1$  and  $w_2$  such that  $w_1 + w_2 = 1$ . Define our objective as a weighted sum:

$$f(x) = w_1 f_1(x) + w_2 f_2(x)$$

For any value of  $w_1$ ,  $w_2$ , this is now a **single objective** we can maximise.

## Weighting plus grid search

- **11** For many evenly-spaced values of  $w_1$ ,
- **2** Define  $w_2 = 1 w_1$
- Define the weighted objective and optimise
- Take the Pareto front across the resulting solutions' cost and emissions values.

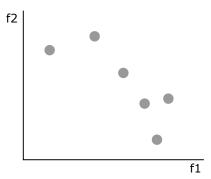
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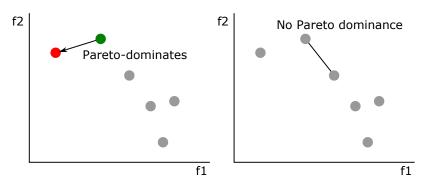
What is the **Pareto front**?

## Plotting our options

For any potential solution x, let's calculate  $f_1(x)$  and  $f_2(x)$ . (Suppose higher is better, for both.)



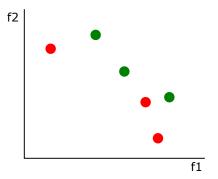
## Pareto dominance



We say that an option  $x_1$  **Pareto-dominates** an option  $x_2$  if  $x_1$  is **better than**  $x_2$  on at least one  $f_i$ , and **at least equal to**  $x_2$  on all  $f_i$ .

An option is **Pareto-dominated** if some other option Pareto-dominates it.

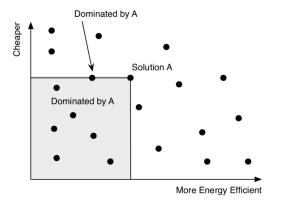
## Pareto front



This allows us to discard dominated options, those shown in red.

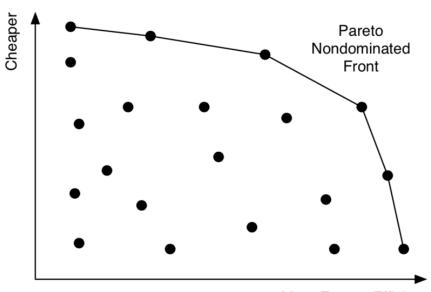
The remainder in green are called the **Pareto front**, the **efficient front**, and similar names. But we have no way to choose between them!

### Pareto Dominance



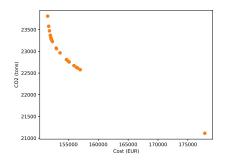
This diagram and later ones are from Luke, Essentials of Metaheuristics, Chapter 7

## **Pareto Front**



More Energy Efficient

## **Multi-objective Unit Commitment**



Remember, all of these solutions are Pareto-optimal.

Having obtained this Pareto front, we are finished. A decision-maker now has to **decide**. Sometimes they'll aim for a **knee** in the plot, or an extreme, or they'll introduce new criteria or other issues.

# Some simple algorithms

#### **Elitism in MOO**

In MOO, the "elite" is just the Pareto front of all points considered so far. It seems natural to **never discard an elite individual**. This leads to a simple algorithm.

## Pareto Archive algorithm

- Create an empty archive, create one individual and add it to the archive.
- **2** Select random individual(s) from the archive.
- **3** Generate new individual by mutation and/or crossover.
- Add the new individual to the archive. If that individual Pareto-dominates any individuals in the set, then **they** are discarded. If it is Pareto-dominated **by** any individual in the set, then **it** is discarded instead.
- **5** Go to 2.

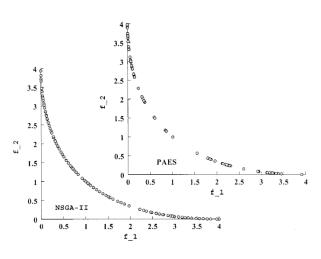
In this algorithm we don't have **generations**. We don't have a fixed population size – in fact the set is likely to grow over time.

## Random-objective tournament selection

- This uses a **normal population-based algorithm**, but:
- At each selection event, we choose **one of our objectives at** random and use it to determine the tournament winner
- Variant known as lexicographic selection algorithm
- See Luke, Algorithm 95.

## **Drawbacks**

The main drawback of these simpler algorithms (Pareto Archive algorithm and random-objective tournament selection): they fail to fully "spread" the population over the Pareto front.

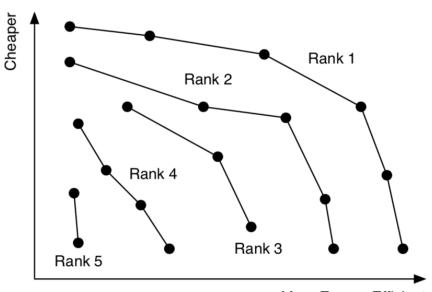


Deb et al.

#### **NSGA2**

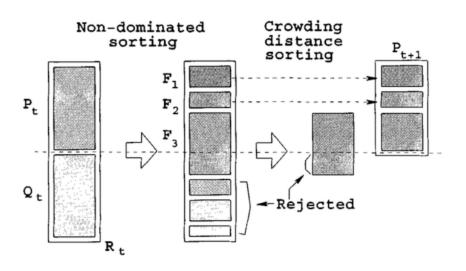
- Non-dominated sorting genetic algorithm 2 (Deb et al., in Blackboard)
- Over 5000 citations.

# Sorting



More Energy Efficient

## **NSGA2** main loop



Deb et al.

## **NSGA2**

#### Main loop:

- Create children from current population
- Merge with population
- Select according to below procedure
- Discard the rest

## **NSGA2**

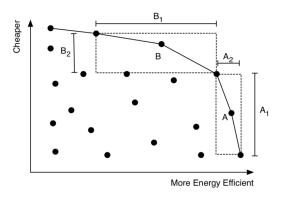
#### Main loop:

- Create children from current population
- Merge with population
- Select according to below procedure
- Discard the rest

#### Selection:

- Select all of the Pareto front
- Then remove the Pareto front and re-calculate the next Pareto front, and select all of this
- Repeat (remove, re-calculate, select) until the next Pareto front would make our population over-full
- From that Pareto front, just select the most **sparse**.

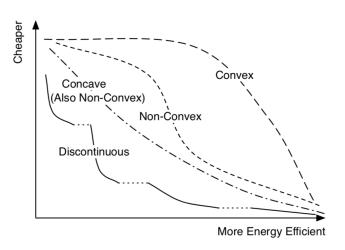
## **Sparsity**



Individual B has sparsity B1+B2. An extreme individual has sparsity  $\infty$ . See Luke, Algorithm 102 for details.

The second main idea in NSGA2 is: we would like to achieve a good **spread** of individuals along a wide Pareto front. We select the individuals with highest sparsity, based on how close they are (in objective space) to other individuals in the same front.

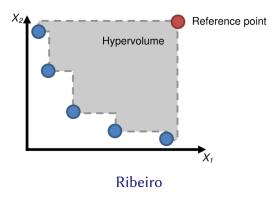
# Types of fronts



## How good is our solution?

In single-objective optimisation we can just report the best value of the objective. But in MOO, we have to report a whole Pareto front. If we are comparing two algorithms, or judging progress over time, that is tricky. How do we tell whether one Pareto front is better than another?

# Hypervolume indicator



Choose some **reference point** worse than all possible solutions on all objectives and calculate this volume: the larger the better.

**DEAP** usage

## How good is our solution?

#### **Principles:**

- Reward Pareto Front for good extreme values
- Reward for diversity across the front (as opposed to clumping)
- If Pareto front PF1 has all the same items as PF2 plus some extra, then PF1 is better.

## Recap

- Multiple conflicting objectives
- Pareto dominance
- Basic approaches
  - Weighting schemes, grid search
  - Pareto Archive
  - Random tournament
- NSGA2, with selection based on:
  - Ranking of Pareto fronts
  - Sparsity
- Hypervolume indicator to measure algorithm success.

## Reading

Luke, Essentials, Chapter on Multiobjective optimisation.

Deb, Multi-objective Optimisation Using Evolutionary Algorithms: An Introduction (in Bb).

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- Multi-objective optimisation algorithms