ASSIGNMENT-1 AMOGH HUILGOL CS6240- SECTION 02

For each of the versions of your sequential and multithreaded program detailed in B and C, report the minimum, average, and maximum running time observed over the 10 runs. (5 points)

SEQUENCE WITHOUT FIB()

Maximum execution time for 10.0 rounds of program execution is 5.918 seconds

Minimum execution time for 10.0 rounds of program execution is 2.724 seconds

Average execution time for 10.0 rounds of program execution is 3.4140000000000000 seconds

SEQUENCE WITH FIB()

Maximum execution time for 10.0 rounds of program execution is 16.042 seconds

Minimum execution time for 10.0 rounds of program execution is 13.772 seconds

Average execution time for 10.0 rounds of program execution is 14.065800000000001 seconds

No LOCK WITHOUT FIB()

NO_LOCK WITH FIB()

Maximum execution time for 10 rounds of program execution is 11.809 seconds Minimum execution time for 10 rounds of program execution is 11.021 seconds Average execution time for 10 rounds of program execution is 11.3602 seconds

Coarse Lock Without FIB()

Maximum execution time for 10 rounds of program execution is 2.698 seconds

Minimum execution time for 10 rounds of program execution is 1.468 seconds

Average execution time for 10 rounds of program execution is 1.6396000000000002 seconds

Coarse Lock With FIB()

Maximum execution time for 10 rounds of program execution is 28.951 seconds Minimum execution time for 10 rounds of program execution is 24.61 seconds Average execution time for 10 rounds of program execution is 26.6085 seconds

Fine Lock Without FIB()

Maximum execution time for 10.0 rounds of program execution is 2.792 seconds

Minimum execution time for 10.0 rounds of program execution is 1.58 seconds

Average execution time for 10.0 rounds of program execution is 1.8167000000000002 seconds

FINE LOCK WITH FIB()

Maximum execution time for 10.0 rounds of program execution is 13.062 seconds

Minimum execution time for 10.0 rounds of program execution is 10.848 seconds

Average execution time for 10.0 rounds of program execution is 12.0466000000000002 seconds

No Shared WITHOUT FIB()

Maximum execution time for 10 rounds of program execution is 4.385 seconds

Minimum execution time for 10 rounds of program execution is 1.431 seconds

Average execution time for 10 rounds of program execution is 1.798100000000003 seconds

NO SHARED WITH FIB()

1) Report the number of worker threads used and the speedup of the multithreaded versions based on the corresponding average running times. (5 points)

The number of worker threads spawned are 4. the speedup obtained is as follows

NO LOCK:

PROGRAM B: 2.09 PROGRAM C: 1.23

COARSE_LOCK:

PROGRAM B: 2.08 PROGRAM C: 0.528

FINE LOCK:

PROGRAM B: 1.879 PROGRAM C: 1.16

NO SHARED:

PROGRAM B: 1.898 PROGRAM C: 2.07

- 3) Answer the following questions in a brief and concise manner: (4 points each)
- i) Which program version (SEQ, NO-LOCK, COARSE-LOCK, FINE-LOCK, NO-SHARING) would you normally expect to finish fastest and why? Do the experiments confirm your expectation? If not, try to explain the reasons.

Answer) NO LOCKING will run the fastest as compared to the other versions because

- a) Since there are no locks, we can utilize maximum parallelism and no worker will have to wait for resource
- b) All workers will update the resource at the same time. hence there is a chance for values to be inconsistent
- c) The average values confirm with the expectation
- ii) Which program version (SEQ, NO-LOCK, COARSE-LOCK, FINE-LOCK, NO-SHARING) would you normally expect to finish slowest and why? Do the experiments confirm your expectation? If not, try to explain the reasons.
- Answer) a) Sequential is generally expected to finish the slowest since each step is executed one after the other. In parallel approach, program makes use of multiple cores available in the system. Hence parallel programs usually finish faster.
 - b) Yes the experiments confirm with expectation
- iii) Compare the temperature averages returned by each program version. Report if any of them is incorrect or if any of the programs crashed because of concurrent accesses.

Answer) The effects of parallelism are better felt when input size is large . However we can see from our experiments that the average values for NO_LOCK mechanism is wrong or inconsistent . This could occur due to

- a) Program crashing because there is lack of synchronization between threads which leads to one thread trying to read a value in Hashmap which is not present. this results in code throwing null pointer exception
- b) Since there is no lock involved, multiple threads may update a data structure based on old record leading to overwriting of one record over other, this leads to inconsistency or wrong results.
- iv) Compare the running times of SEQ and COARSE-LOCK. Try to explain why one is slower than the other. (Make sure to consider the results of both B and C—this might support or refute a possible hypothesis.)

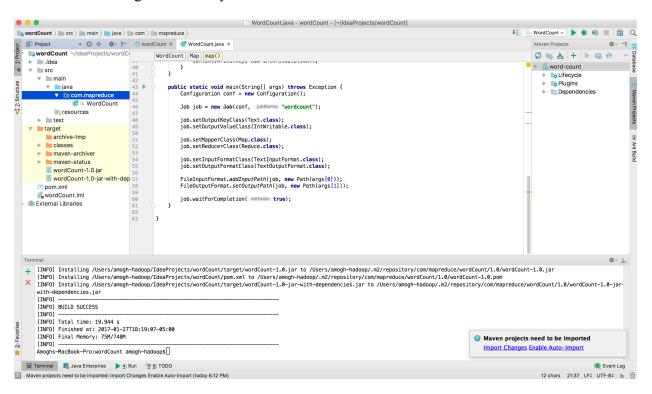
Answer) Sequential is generally slower than coarse lock mechanism. However we have seen that when delay fib(17) is introduced, sequential program performs faster, the possible reasons could be:

- a) When inputs are small, the tasks are performed quickly with less transfer of locks between threads. Hence parallelism is utilized to speed up the computation
- b) However as the inputs increase, the wait time for threads increases as they have to wait more often for other thread to release locks. This defeats the improvement in performance we obtain by running a program parallel
- v) How does the higher computation cost in part C (additional Fibonacci computation) affect the difference between COARSE-LOCK and FINE-LOCK? Try to explain the reason.

Answer) In COARSE_LOCK, the lock is held over the entire data structures and this ensures that no other thread is trying to update the hashmap, thus it can be guaranteed that no two Fibonacci computations will happen in parallel. Thus all delays will happen in sequential way. However for FINE_LOCK there is a chance that some Fibonacci computations may happen in parallel since each thread holds lock over the variable which it updates (not the whole hashmap). Hence with higher computation cost, the fine lock performs better than coarse lock

WORD COUNT LOCAL EXECUTION

Screenshot showing the directory structure

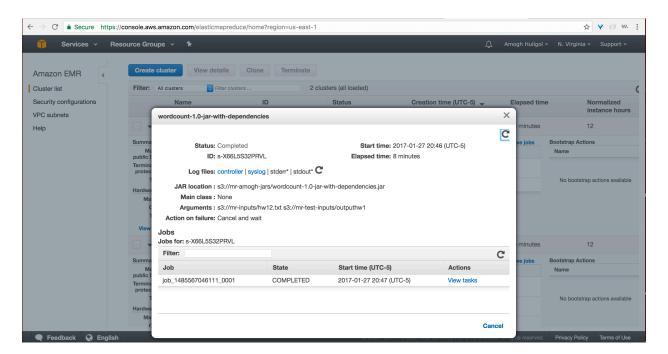


Screen shot showing last 20 lines of Hadoop execution in IDE

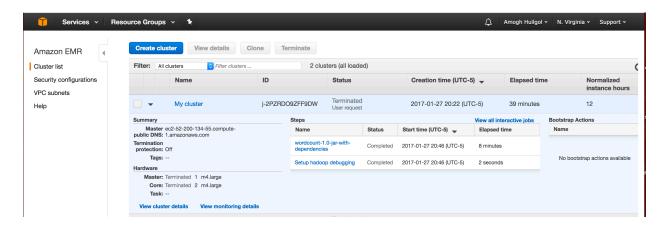
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2017-01-27 17:30:40,343 INFO [pool-3-thread-1] outpt 2017-01-27 17:30:40,343 INFO [pool-3-thread-1] mapre 2017-01-27 17:30:40,343 INFO [pool-3-thread-1] mapre 2017-01-27 17:30:40,344 INFO [pool-3-thread-1] mapre 2017-01-27 17:30:40,344 INFO [pool-3-thread-1] mapre 2017-01-27 17:30:40,344 INFO [pool-3-thread-1] mapre 2017-01-27 17:30:40,444 INFO [main] mapreduce.job (j 2017-01-27 17:30:40,441 INFO [main] mapreduce.job (j 2017-01-30:40,441 INFO [mai
                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                                        [pool-3-thread-1] output.FileOutputCommitter (FileOutputCommitter.java:commitTask(535)) - Saved output of task 'attempt_local1771957700_0001_r_000000_0' to fi' [pool-3-thread-1] mapred.LocalJobRunner (LocalJobRunner.java:statusUpdate(591)) - reduce > reduce | redu
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EXECUTION ON AWS

Screenshot showing successful execution of word count in AWS-EMR



Screenshot showing cluster details



CONTROLLER.txt

SYSLOG.txt