# Part I

## **Encodings**

Using the minimal  $\lambda$ -calculus language we get

- √ functions
- ✓ local binding
- √ booleans
- ✓ numbers

... and recursive functions?

### Factorial in Plait

local binds both in the body expression and in the binding expression

### Factorial in Plait

letrec has the shape of let but the binding structure of local

### Factorial in Plait

Doesn't work, because let binds fac only in the body

Overall goal: Implement letrec as syntactic sugar for Curly

```
{letrec {[name rhs]}
  name}
```

Step I: Implement fac in Plait without letrec

Step 2: Isolate the rhs

**Step 3:** Surrounding as a parse transformation for Curly

This is Difficult...





# Part 2

Overall goal: Implement letrec as syntactic sugar for Curly

```
{letrec {[name rhs]}
  name}
```

Step I: Implement fac in Plait without letrec

Step 2: Isolate the rhs

**Step 3:** Surrounding as a parse transformation for Curly

At the point that we call fac, obviously we have a binding for fac...

... so pass it as an argument!

Wrap this to get fac back...

# Part 3

Overall goal: Implement letrec as syntactic sugar for Curly

```
{letrec {[name rhs]}
  name}
```

Step I: Implement fac in Plait without letrec

```
Step 2: Isolate the rhs
```

**Step 3:** Surrounding as a parse transformation for Curly

But Curly has only single-argument functions...

```
(let ([fac
       (let ([facX
               (lambda (facX)
                 (let ([fac (facX facX)])
                    ; Exactly like original fac:
                    (lambda (n)
                      (if (zero? n)
                          (* n (fac (- n 1)))))))))
          (facX facX))])
  (fac 10))
Oops! — this is an infinite loop
We used to evaluate (facX facX) only when n is non-zero
                                      Delay (facX facX)...
```

```
(define mk-rec
  (lambda (body-proc)
    (let ([fX
           (lambda (fX)
             (let ([f (lambda (x)
                         ((fX fX) x))])
                (body-proc
                f)))])
      (fX fX))))
(let ([fac
       (mk-rec
        (lambda (fac)
          ; Exactly like original fac:
          (lambda (n)
            (if (zero? n)
                1
                 (* n (fac (- n 1)))))))))
  (fac 10))
```

### **Fibonnaci**

### Sum

# Part 4

Overall goal: Implement letrec as syntactic sugar for Curly

```
{letrec {[name rhs]}
  name}
```

Step I: Implement fac in Plait without letrec

Step 2: Isolate the rhs

**Step 3:** Surrounding as a parse transformation for Curly

```
{letrec {[fac {lambda {n}}
                            {if0 n
                                  1
                                  {* n
                                     {fac {- n 1}}}}}}
           {fac 10}}
could be parsed the same as
              {let {[fac
                     {mk-rec
                      {lambda {fac}
                        {lambda {n}
                          {if0 n
                               1
                               {* n
                                  {fac {- n 1}}}}}}}]}
                {fac 10}}
```

```
{letrec {[fac {lambda {n}}
                             {if0 n
                                   1
                                   {* n
                                       {fac {- n 1}}}}}]
            {fac 10}}
                                       mk-rec = {lambda {body-proc}
could be parsed the same as
                                                  {let {[fX
                                                         {lambda {fX}
              {let {[fac
                                                           {let {[f {lambda {x}}
                      {mk-rec
                                                                      {{fX fX} x}}}
                       {lambda {fac}
                                                             {body-proc f}}}]}
                         {lambda {n}
                                                    {fX fX}}}
                           {if0 n
                                1
                                {* n
                                   {fac {- n 1}}}}}}}}
                {fac 10}}
```

```
{letrec {[fac {lambda {n}}
                             {if0 n
                                   1
                                   {* n
                                       {fac {- n 1}}}}}]
            {fac 10}}
                                       mk-rec = {lambda {body-proc}
could be parsed the same as
                                                  {{lambda {fx} {fX fX}}
                                                   {lambda {fX}
              {let {[fac
                                                     {{lambda {f} {body-proc f}}}
                      {mk-rec
                                                      {lambda {x}
                       {lambda {fac}
                                                        {{fX fX} x}}}}
                         {lambda {n}
                           {if0 n
                                1
                                {* n
                                   {fac {- n 1}}}}}}}}
                {fac 10}}
```

```
{letrec {[name rhs]}
    body}
could be parsed the same as
 {let {[name {mk-rec {lambda {name} rhs}}]}
    body}
which is really
 {{lambda {name} body}
   {mk-rec {lambda {name} rhs}}}
which, writing out mk-rec, is really
 {{lambda {name} body}
  {{lambda {body-proc}}
    {let {[fX {fun {fX}}
              {let {[f {lambda {x}}
                     {{fX fX} x}}]}
               {body-proc f}}}]}
     {fX fX}}}
   {lambda {name} rhs}}}
```

# Part 5

## The Big Picture



#### Y Combinator

#### mk-rec is better known as the Y combinator

```
{lambda {body-proc}
  {{lambda {fx} {fX fX}}
    {lambda {fX}
        {{lambda {f}} {body-proc f}}}
        {lambda {x}
        {{fX fX} x}}}}
```

#### Y Combinator

mk-rec is better known as the Y combinator

```
Y = (λ (g)
{(λ (fx) {fX fX})
(λ (fX)
{(λ (f) {g f})
(λ (x)
{fX fX fX} x})})
```

a.k.a. the **fixpoint operator** 

```
{Y (lambda (f<sub>in</sub>) f<sub>out</sub>)}
```

### Y Combinator

mk-rec is better known as the Y combinator

```
Y = (λ (g)
{(λ (fx) {fX fX})
(λ (fX)
{(λ (f) {g f})
(λ (x)
{fX fX fX} x})})
```

See also The Why of Y (Gabriel) or The Little Schemer (Friedman & Felleisen)

# Part 6

### Example with Quasiquote Escapes