CSE 321b

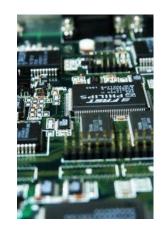
Computer Organization (2)

تنظيم الحاسب (2)



3rd year, Computer Engineering
Winter 2016

Lecture #1



Dr. Hazem Ibrahim Shehata Dept. of Computer & Systems Engineering

Credits to Dr. Ahmed Abdul-Monem Ahmed for the slides

Teaching Staff

- Instructor:
 - —Hazem Ibrahim Shehata
 - —Email: hshehata@uwaterloo.ca
 - —Lectures: Thursday 9:00am-11:30am
 - —Office Hours: TBA
- Teaching Assistant:
 - —Hisham Abdullah
 - —Email: eng hisham22@yahoo.com
 - —Tutorials: TBA
 - —Office Hours: TBA

Course Info

- Course website:
 - —http://www.bit.do/hshehata-courses-zu-cse321b
 - http://www.googledrive.com/host/0B9ExmUsPoGjSdUxvc1BzRndQS1U
- Textbook:
 - —"Computer Organization and Architecture: Designing for Performance", William Stallings, 9th Edition, 2013, www.williamstallings.com/ComputerOrganization

Course Info (Cont.)

• Grading:

Course work	Grade distribution		
Participation	3pt	30	
Assignments	12pt		
Midterm Exam	15pt		
Final Exam	70pt		
Total Points	100		

Course Overview

- Ch. 5: Internal Memory Technology
 - —Semiconductor MM, error correction, ..., etc.
- Ch. 6: External Memory
 - —Magnetic disks, optical disks, magnetic tapes, ... etc.
- Ch. 7: Input / Output
 - —Programmed i/o, interrupt-driven i/o, DMA, ..., etc.
- Ch. 10: Computer Arithmetic
 - —Integer representation, Integer arithmetic, FP representation, FP arithmetic, ..., etc.
- Ch. 14: Processor Structure and Function
 - —Processor organization, register organization, Instruction pipelining, ..., etc.

Ch 5: Internal Memory Technology

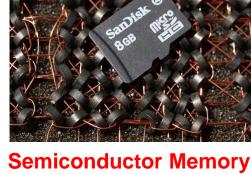
Memory Cell

 Semiconductor memory: electronic memory implemented on a semiconductor-based IC.

- Memory cell: basic element of a semiconductor memory.
 - Holds one bit.
 - Properties



- Can be written into to set the state.
- Can be read to sense the state.

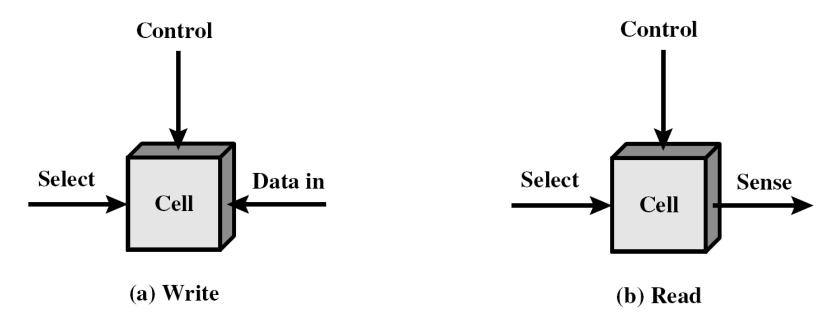


Semiconductor Memory
vs.

Magnetic-core Memory

Conceptual Operation of a Memory Cell

- Three terminals: select, control, data in/sense.
 - Select: select a memory cell for read/write.
 - Control: indicate required operation: read or write.
 - Data in/sense:
 - Read: Output the state of the cell.
 - Write: electrical signal that sets the state to 0 or 1.



Semiconductor Memory Types

Volatile

- —Random Access Memory (RAM)
 - Dynamic RAM (DRAM)
 - Static RAM (SRAM)

Non-volatile

- —Read-Only Memory (ROM)
- —Programmable ROM (PROM)
- —Erasable Programmable ROM (EPROM)
- —Electrically Erasable Programmable ROM (EEPROM)
- —Flash Memory

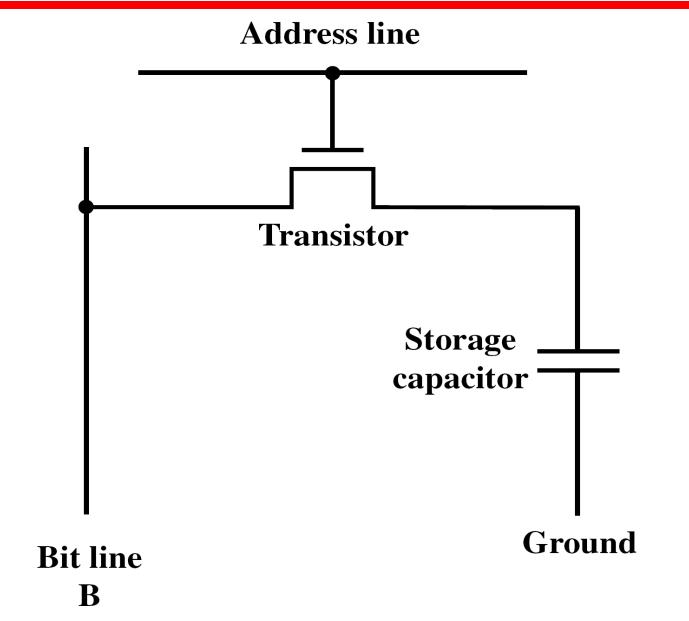
RAM

- Individual words of memory are directly accessed through wired-in addressing logic.
- Misnamed!! All semiconductor memories are random access!!
- Read/Write: by electrical signals.
- Volatile: must be provided with a constant power supply → temporary storage.
- Dynamic or static.

Dynamic RAM (DRAM)

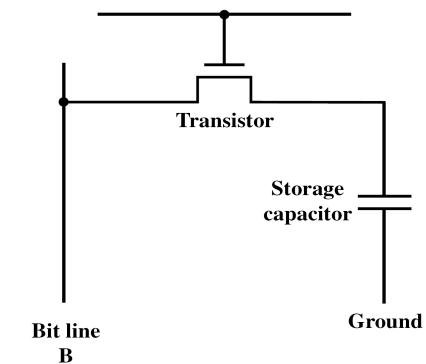
- Bits are stored as charge on capacitors.
 - Charge \rightarrow 1, no charge \rightarrow 0.
- Capacitors discharge → DRAM needs periodic charge refreshing even when powered.
- Analog device: capacitor can store any charge value within a range → a threshold value is used.
- Pros
 - Simpler construction.
 - Smaller per bit.
 - Less expensive.
- Cons
 - Need refresh circuits.
 - Slower.
- Most common usage: main memory.

DRAM Structure



DRAM Operation

- Address line active when bit read or written.
 - —Transistor switch closed (current flows). Address line
- Write
 - —Voltage to bit line
 - High for 1 low for 0.
 - —Then signal address line
 - Transfers charge to capacitor.
- Read
 - —Address line selected
 - transistor turns on.



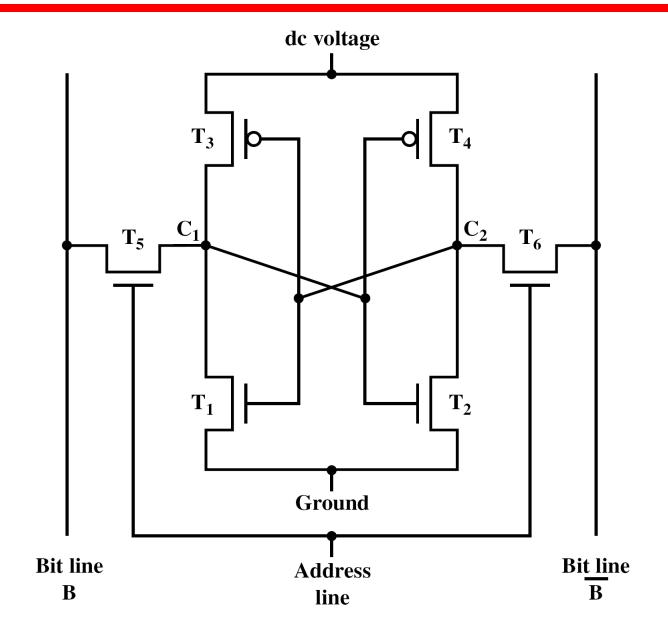
- —Charge from capacitor fed via bit line to a sense amplifier
 - Compares with threshold/reference value to determine 0 or 1.
- Readout discharges capacitor

 charge must be restored

Static RAM (SRAM)

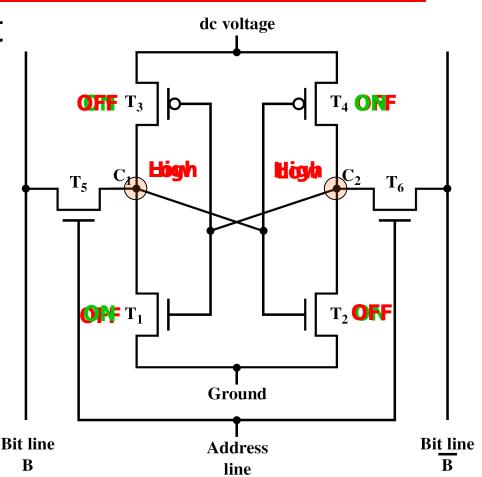
- Bits stored as on/off switches.
- Digital device: uses flip-flops.
- No charges to leak.
- No refreshing needed.
- Pros
 - —Does not need refresh circuits.
 - —Faster.
- Cons
 - —More complex construction.
 - —Larger per bit.
 - —More expensive.
- Most common usage: cache memory.

SRAM Structure



SRAM Operation

- Transistor arrangement gives stable logic state.
- C₁ and C₂: diff. states
- State 1
 - $-C_1$ high, C_2 low
 - $-T_2$ on, T_4 off
 - $-T_1$ off, T_3 on
- State 0
 - C₁ low, C₂ high
 - $-T_2$ off, T_4 on
 - $-T_1$ on, T_3 off



- Address line transistors T₅ T₆ are switches.
- Write apply value to B & compliment to B.
- Read value is on line B.

DRAM vs SRAM

- Both volatile
 - —Power needed to preserve data.
- Dynamic cell
 - —Simpler to build, smaller.
 - —More dense: more cells per unit area.
 - -Less expensive.
 - —Needs refreshment.
 - —Fixed cost of refreshment circuitry → use large memory units to benefit from the small cell cost.
 - —Used in **main memory**.
- Static cell
 - —Faster
 - —Used in **cache memory**.

Read Only Memory (ROM)

- Permanent storage that cannot be changed.
 - —Nonvolatile.
- Can read stored data, cannot write new data.
- Written during fabrication

 - —No room for error. One bit error → throw the whole batch of ROMs.
- Why useful?
 - Data or program is permanently in main memory and need never be loaded from a secondary storage device.
- Applications
 - —Microprogramming.
 - —Library subroutines.
 - —System programs (BIOS).

Programmable ROM (PROM)

- Nonvolatile.
- Can be written into only once.
- Writing (or programming):
 - —Performed electrically using a special equipment.
 - Writing one → do nothing! (all cells store one by default).
 - Writing zero → blow a fuse (or melt an anti-fuse) in the cell.
 - —Performed by supplier or customer (after fabrication).
- Useful when a small number of ROMs with a particular memory content is needed.
- Flexible and convenient.
- ROM is good for high-volume production.

Erasable Programmable ROM (EPROM)

- Nonvolatile.
- Read-mostly memory: read operations are far more than write operations.
- Read and written electrically.
 - —Before a write operation, all cells must be optically erased to the same initial state.
- Erasure
 - —Done optically by exposure of the packaged chip to ultraviolet radiation.
 - —Takes up to 20 minutes.
 - —Can be done repeatedly.
- One transistor per bit → dense.
- More expensive than PROM, but can do multiple updates.

Electrically Erasable Programmable ROM (EEPROM)

- Nonvolatile.
- Read-mostly memory.
- Can be written into without erasing prior contents
 - —Only the byte/bytes addressed are updated.
- Write operation takes longer than read.
- Flexible: updatable in place using ordinaly bus lines.
- More expensive than EPROM.
- Less dense than EPROM: fewer bits per chip.

Flash Memory

- Nonvolatile.
- Read-mostly memory.
- First introduced in the mid-1980s.
- Intermediate between EPROM and EEPROM in cost and functionality
 - —Like EEPROM, it uses electrical erasing technology.
 - Entire flash memory can be erased in a few seconds
 → much faster than EPROM.
 - —Only a block of memory can be erased.
 - —No byte-level erasure.
 - —Like EPROM, one transistor per bit → higher density than EEPROM.

Semiconductor Memory Types - Summary

Memory Type	Category	Erasure	Write Mechanism	Volatility
Random-access memory (RAM)	Read-write memory	Electrically, byte-level	Electrically	Volatile
Read-only memory (ROM)	Read-only	Not possible	Masks	
Programmable ROM (PROM)	memory			Nonvolatile
Erasable PROM (EPROM)		UV light, chip-level		
Electrically Erasable PROM (EEPROM)	Read-mostly memory	Electrically, byte-level	Electrically	
Flash memory		Electrically, block-level		

Chip Logic

- Semiconductor memory comes in packaged chips.
- Each mem. chip contains an array of memory cells.
- Design issue: number of bits of data that maybe read/written at a time.
 - One extreme: 1-word-per-chip organization: physical arrangement of cells in the array is the same as logical arrangement of words in memory.
 - EX.: 1Mx16 memory = **One** 1Mx16 chip.
 - —Other extreme: 1-bit-per-chip organization: data is read/written 1 bit at a time.
 - Ex.: 1Mx16 memory = sixteen 1Mx1 chips s.t. chip #1 holds bit #1 of each word, chip #2 holds bit #2 of each word, and so on.

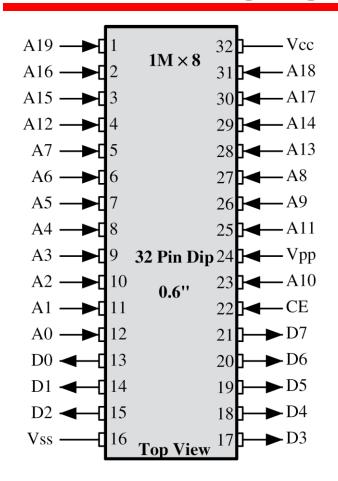
RAS CAS WE OE Typical 16 Mb DRAM (4Mx4) Timing and Control Refresh MUX Counter Row Row Memory array De-Address $(2048 \times 2048 \times 4)$ coder Buffer Data Input Column Buffer A10 Address Refresh circuitry Buffer Data Output Buffer Column Decoder

- Logically, 4 square arrays of 2048 x 2048 elements.
- Each horizontal line connects to the Select terminal of each cell in its row.
- Each vertical line connects to the Data in/Sense terminal of each cell in its column.
- Reduces number of address pins
 - Multiplex row address and column address
 - 11 pins to address (2¹¹=2048)
 - Adding one more pin doubles range of values so x4 capacity

Refreshing

- Refresh circuit included on chip.
- Disable chip.
- Count through rows.
- Data is read out and written back into the same location → each cell is refreshed.
- Takes time.
- Slows down apparent performance.

Chip Packaging



_		
Vcc — 1	24	Vss
D1 ←► [2	2 23] ←► D4
D2 ←► [3	3 22]← ▶D3
WE → []4	1 21	J ←CAS
RAS → [5	5 20	J ←OE
NC → [6	5 19	→ A9
A10 → [7	7 18	→ A8
A0 → [8	3 17	→ A7
A1 → [9) 16	1 ←A6
A2 → []1	10 15	→ A5
A3 → [1	14	→ A4
Vcc — 1	13	Vss

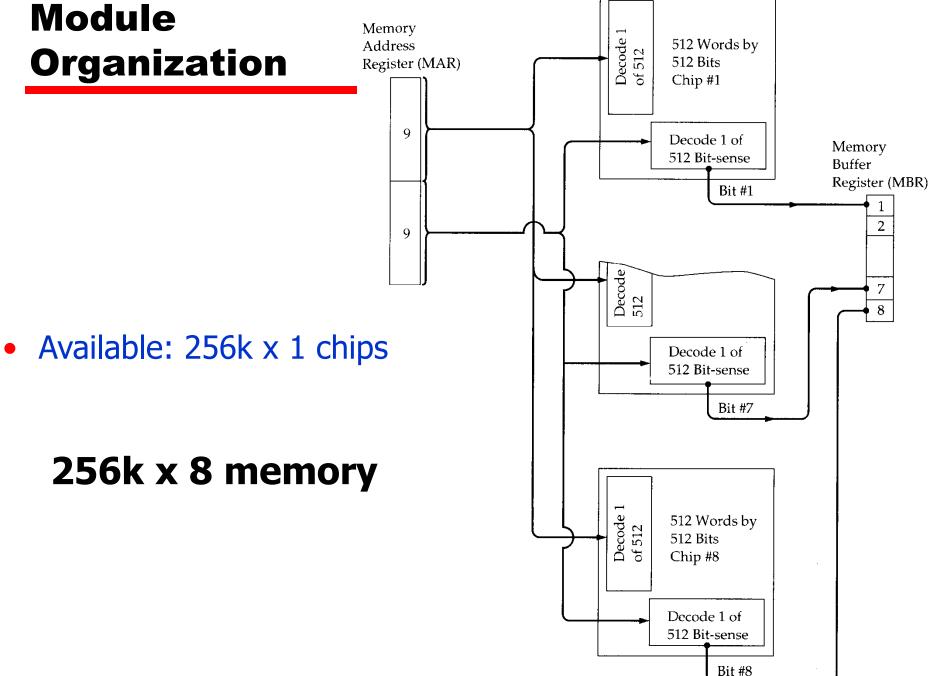
(a) 8 Mbit EPROM

- 8-Mbit EPROM chip, 1M x 8. One-word-per-chip package.
- Address: A0-A19, Data: D0-D7 Vcc; power, Vss: ground, CE: chip enable, Vpp: programming voltage.

(b) 16 Mbit DRAM

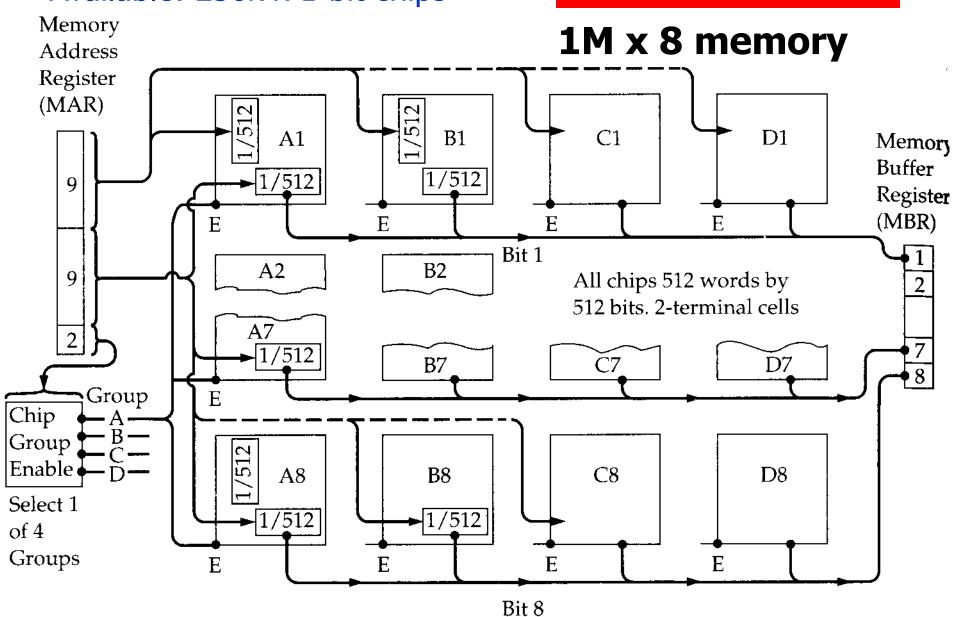
- 16-Mbit DRAM, 4M x 4.
- Updatable → data pins in/out.
- OE: Output Enable
- NC: No Connect → even # of pins

Module Organization



Module Organization (2)

Available: 256k x 1-bit chips



Reading Material

- Stallings, Chapter 5:
 - —Pages 159 169