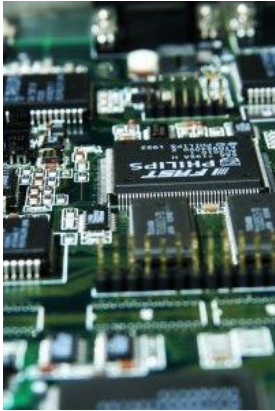


CSE 401

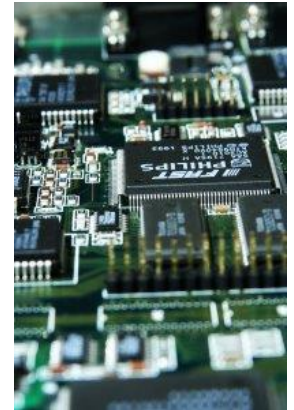
Computer Engineering (2)

هندسة الحاسبات (2)



4th year, Comm. Engineering
Winter 2016

Lecture #6



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Credits to Dr. Ahmed Abdul-Monem Ahmed for the slides

Adminstrivia

- Assignment #2:
 - To be released next week.

Website: <http://hshehata.github.io/courses/zu/cse401/>

Office hours: Monday 11:30am – 12:30pm

Chapter 6. External Memory (*Cont.*)

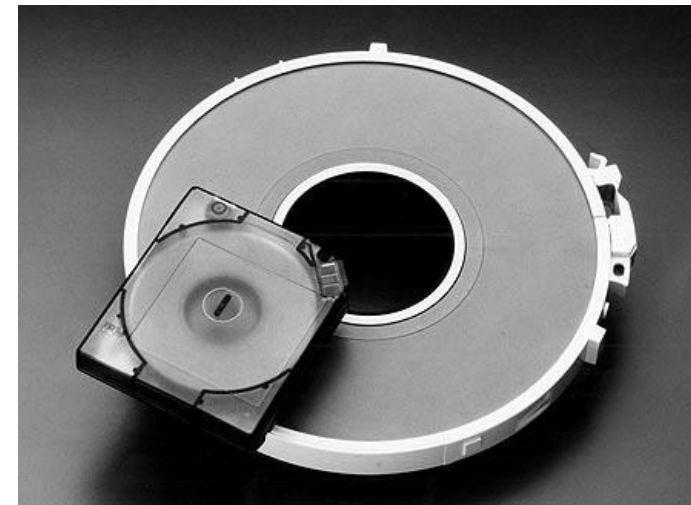
Types of External Memory

- Magnetic Disk
- Redundant Array of Independent Disks (RAID)
- Solid-State Drive (SSD)
- Optical Disk
- Magnetic Tape

Magnetic Tape

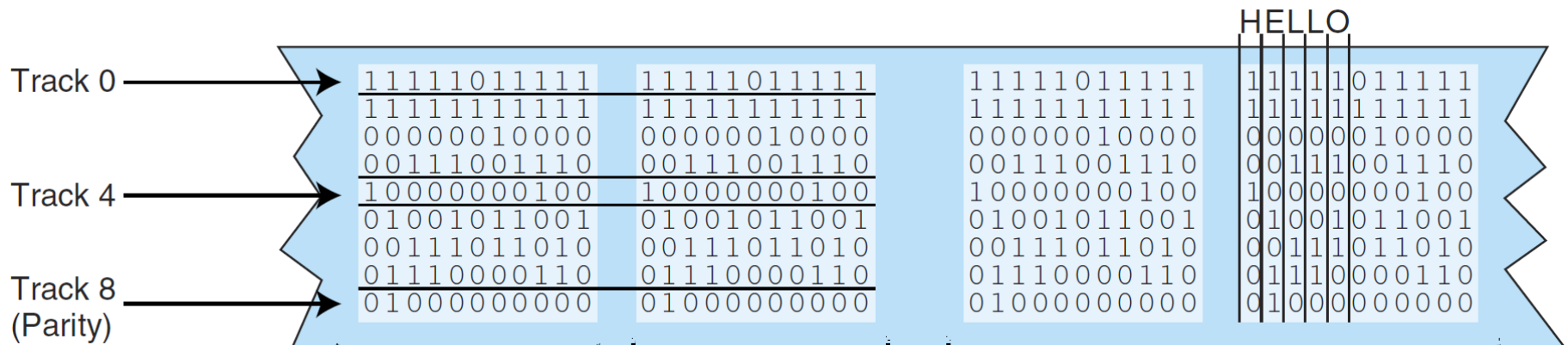


- **Features:**
 - **Oldest** secondary memory (1951).
 - **Slowest-speed** & **lowest-cost** in memory hierarchy.
- **Usage:** **offline storage** (i.e., backup).
- **Access method:** **sequential** → slow!
- **Medium:** flexible polyester tape coated with magnetizable material (metal particles).
- **Tape width:** 0.38cm - 1.27cm.
- **Tape package:**
 - Used to be **open reels**.
 - Now housed in **cartridges**.



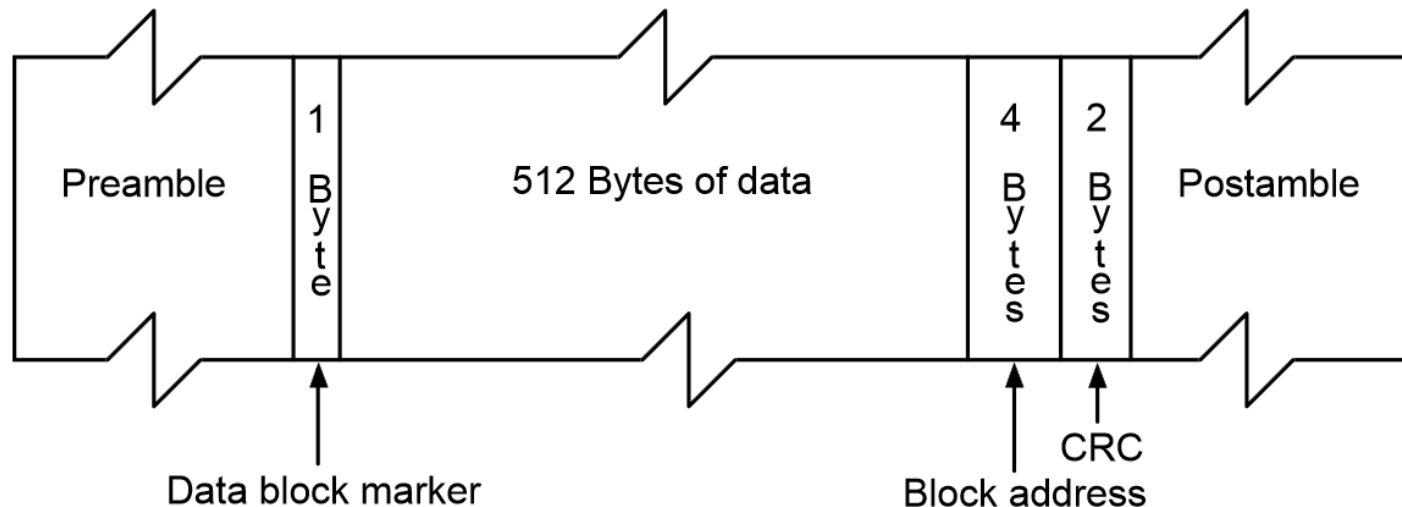
Data Layout and Recording Techniques

- Data are laid out on tape along **parallel tracks** running lengthwise → **linear recording**.
- Tape moves underneath **magnetic head** in drive.
- Head has multiple (8-32) read/write elements → read/write multiple (8-32) tracks simultaneously.
- There are two types of **linear recording**:
 1. **Parallel recording**: Each byte/word is stored on multiple tracks in parallel, .e.g., 9-track tapes.



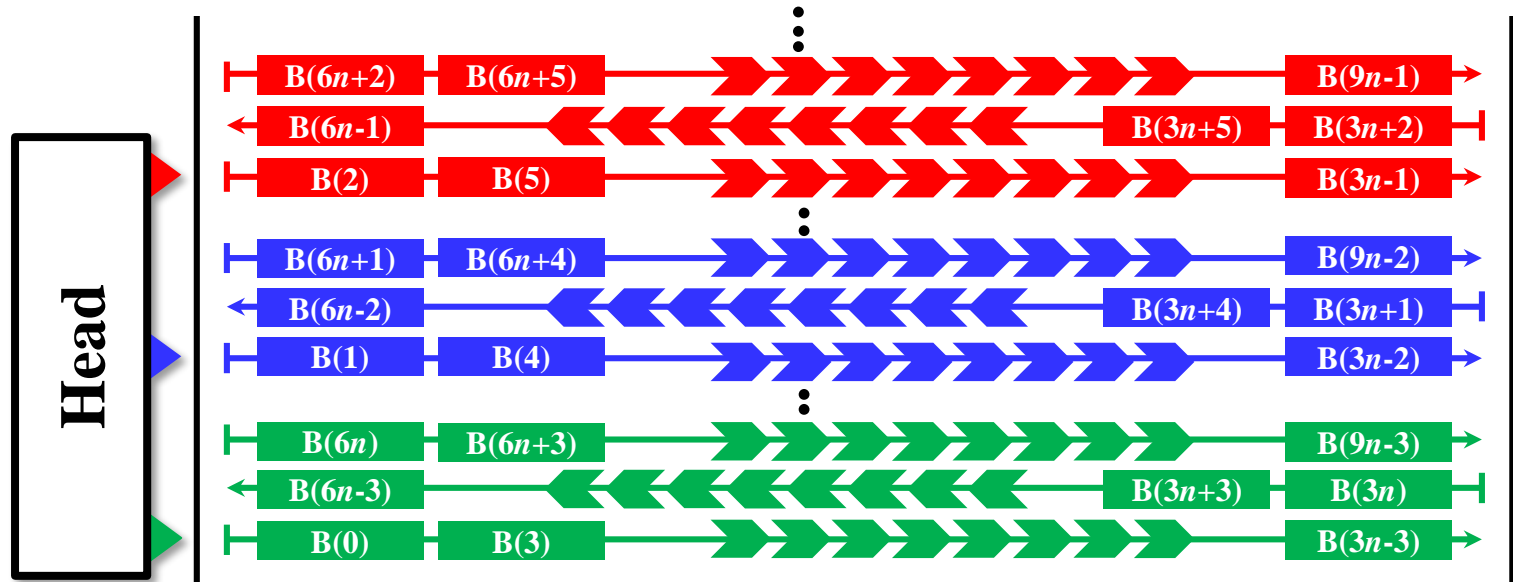
Serial Recording

- ... types of **linear recording** ... (cont.):
- 2. **Serial Recording**: data laid out as sequence of bits along each track, similar to magnetic disks.
 - Data read/written in contiguous blocks called **records**.
 - Records are separated by inter-record gaps.
 - Tape is formatted to assist in locating records.



Serpentine Recording

- Typical recording technique in serial tapes.
 - First set of bits is recorded on tracks in **forward direction**.
 - When tape reaches end, rd/wr elements are repositioned to new tracks & recording is resumed in **opposite direction**.
 - Process continues, back and forth, until tape is full.
 - Data still recorded **serially along individual tracks**, but blocks in sequence are stored in adjacent tracks.
 - **Adv.:** more tracks than head/elements → increasing capacity.



Linear Tape-Open (LTO)

- Dominant magnetic tape technology today.
- Developed late 1990s as an **open source** alternative to proprietary tape systems.

| | LTO-1 | LTO-2 | LTO-3 | LTO-4 | LTO-5 | LTO-6 | LTO-7 | LTO-8 |
|--------------------------|---------|---------|----------|----------|----------|----------|----------|-----------|
| Release date | 2000 | 2003 | 2005 | 2007 | 2010 | 2012 | 2015 | TBA |
| Compressed capacity | 200 GB | 400 GB | 800 GB | 1600 GB | 3000 GB | 6250 GB | 16 TB | 32 TB |
| Compressed transfer rate | 40 MB/s | 80 MB/s | 160 MB/s | 240 MB/s | 280 MB/s | 400 MB/s | 788 MB/s | 1.18 GB/s |
| Linear density (bits/mm) | 4880 | 7398 | 9638 | 13250 | 15142 | 15143 | | |
| Tape tracks | 384 | 512 | 704 | 896 | 1280 | 2176 | 3584 | |
| Tape length (m) | 609 | 609 | 680 | 820 | 846 | 846 | 960 | |
| Tape width (cm) | 1.27 | 1.27 | 1.27 | 1.27 | 1.27 | 1.27 | 1.27 | |
| Write elements | 8 | 8 | 16 | 16 | 16 | 16 | 32 | |
| WORM? | No | No | Yes | Yes | Yes | Yes | Yes | Yes |
| Encryption Capable? | No | No | No | Yes | Yes | Yes | Yes | Yes |
| Partitioning? | No | No | No | No | Yes | Yes | Yes | Yes |

Chapter 7. Input / Output

Outline

- External Devices
 - Types
 - Structure
- I/O Modules
 - Function
 - Structure
- I/O Techniques
 - Programmed I/O
 - Interrupt-Driven I/O
 - Direct Memory Access
- I/O Channels & Processors

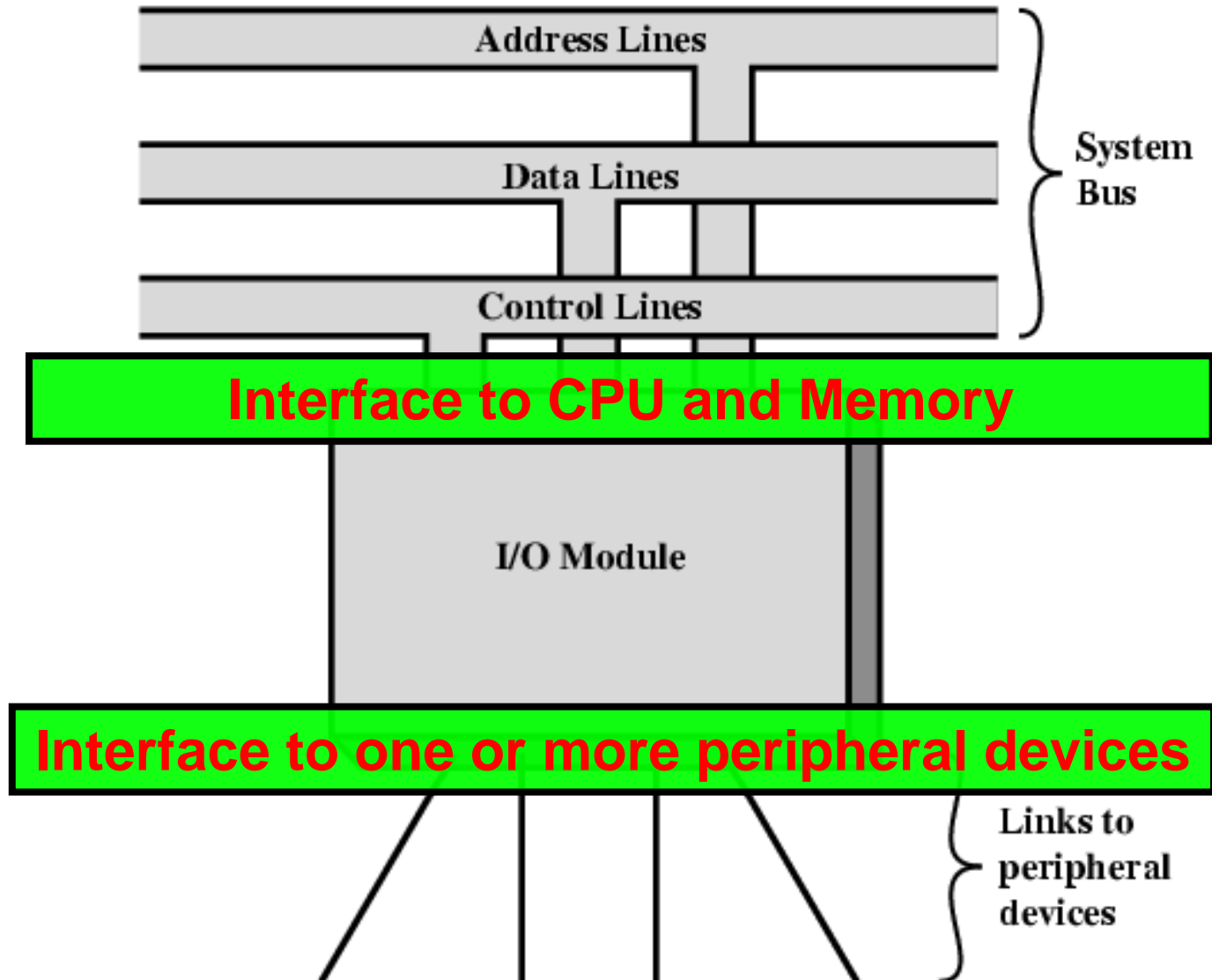
Terms

- Essential Computer Units
 - CPU and Memory
- Peripheral (or External or I/O) devices
 - Any device attached to a computer in order to increase its functionality.
 - Input: keyboard, mouse, scanner, ... etc.
 - Output: printers, speakers, ... etc.
 - Input and output: hard disk, modem, ... etc.
- I/O (Input/Output) Operations
 - Transfer of data to/from computer from/to peripheral device (done by program, operation, or device).
 - Input: from a device to the computer
 - Output: from the computer to a device.

Input/Output Problems

- There is a wide variety of peripherals!
 - Different methods of operation (H/W).
 - Delivering different amounts of data.
 - At different speeds (which are also different from CPU and memory).
 - In different formats (e.g., word length).
- **Conclusion:** Hard to connect such variety of different devices directly to same Bus!!
- **Solution:** **I/O Module**

I/O Module



Types of Peripherals

- Human readable
 - Screen, printer, keyboard, ... etc.
- Machine readable
 - Magnetic disk, tape, ... etc.
- Communication
 - Modem, Network Interface Card (NIC), Wireless Network Adaptor, ... etc.

Peripheral (External) Device

- **Control Signals**

- Send data to module, receive data from module, send status, position disk head.

- **Status Signals**

- READY, NOT READY, ...

- **Buffer**

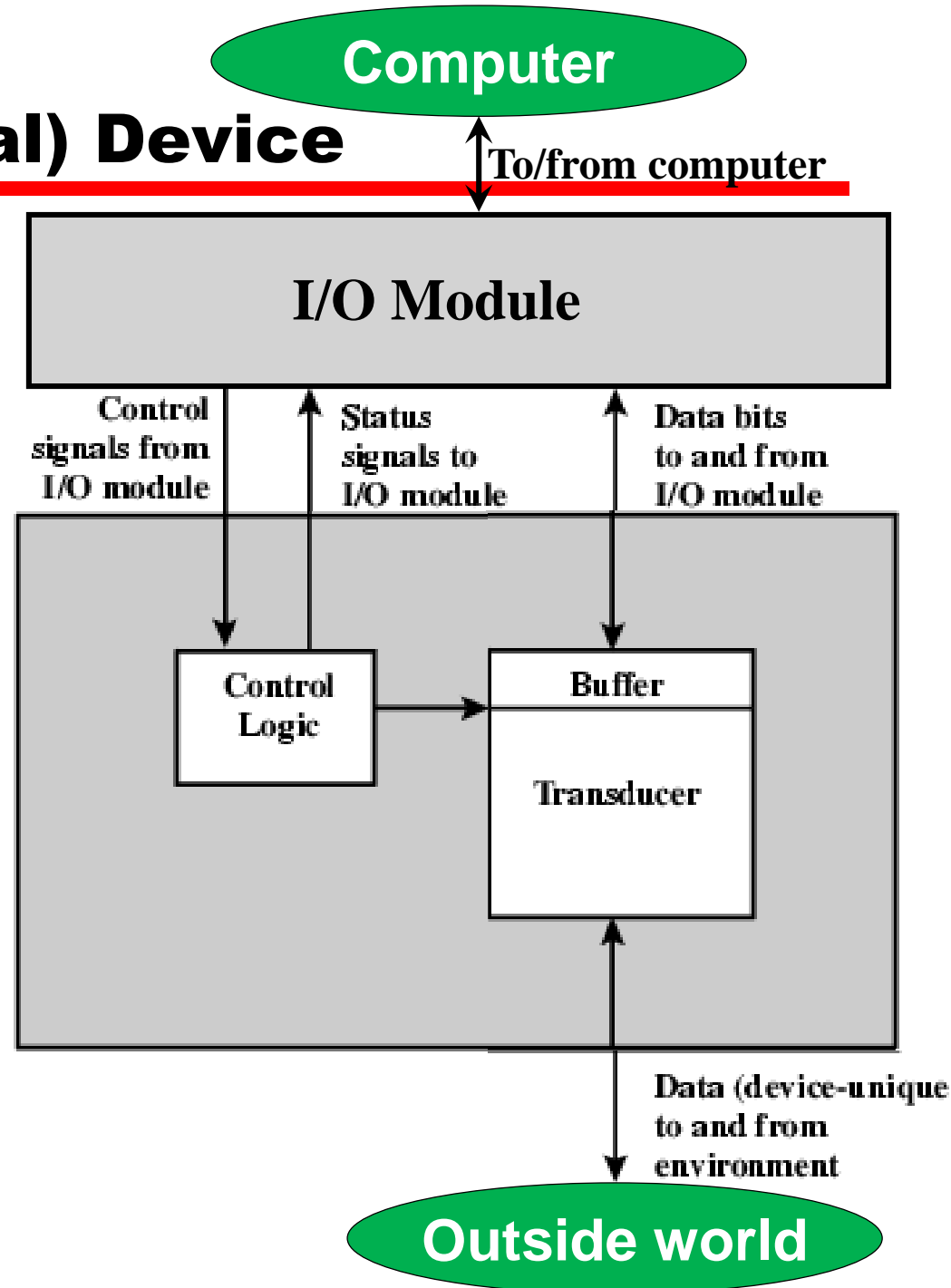
- Temporarily hold data being transferred. Size: x bytes → x Kbytes!!

- **Transducer**

- Converts energy: electrical ↔ other.

- **Control logic**

- Controls operation.



Examples: Keyboard/Monitor, and Disk Drive

- **Keyboard** (input)
 - A key is pressed.
 - Transducer translates signal into ASCII.
 - ASCII is transmitted to I/O module in the computer.
 - Text can be stored as ASCII in the computer.
- **Monitor** (output)
 - Computer sends ASCII to I/O module. I/O module sends ASCII to external device (monitor).
 - Transducer at the monitor sends electronic signals to display the character.
- **Hard Disk Drive** (input/output)
 - Head moves in and out across disk surface.
 - Transducer converts magnetic patterns to/from bits.

Functions of I/O Module

1. Control & Timing.
2. CPU Communication.
3. Device Communication.
4. Data Buffering.
5. Error Detection.

1. Control & Timing

- I/O includes **control & timing** requirement to coordinate the flow of traffic between internal resources (CPU, MM, ...) and external devices.
- Ex.: Transfer data from input device to CPU:
 1. CPU checks I/O module device status.
 2. I/O module returns status.
 3. If ready, CPU requests data transfer (command to I/O module).
 4. I/O module gets data from external device.
 5. I/O module transfers data to CPU.
- If transfer goes through bus, each CPU/module interaction may involve 1+ bus arbitrations.

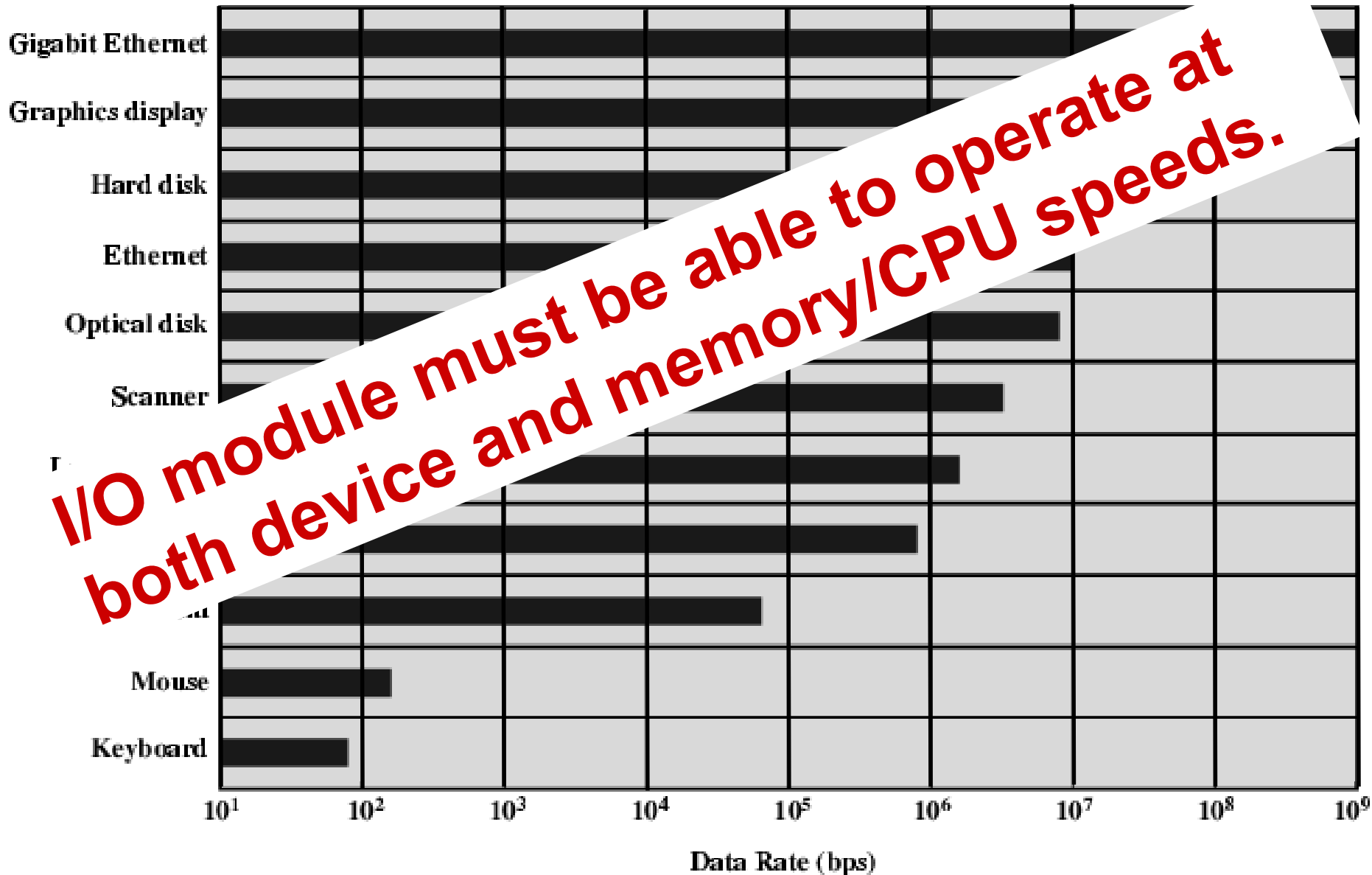
2. CPU Communication

- CPU communication involves the following:
 - Command decoding
 - Module accepts commands from CPU on control lines.
 - Command parameters can be sent over data line.
 - e.g., SEEK track in a disk drive: SEEK command sent on control lines and track # sent on data lines.
 - Address recognition
 - One unique address for each peripheral it controls.
 - Data exchange
 - Between CPU and device over the data bus.
 - Status reporting
 - BUSY, READY, or some error conditions.

3. Device Communication

- I/O module must also be able to do device communication:
 - Commands
 - Status information
 - Data

4. Data Buffering (Speed Mismatch)

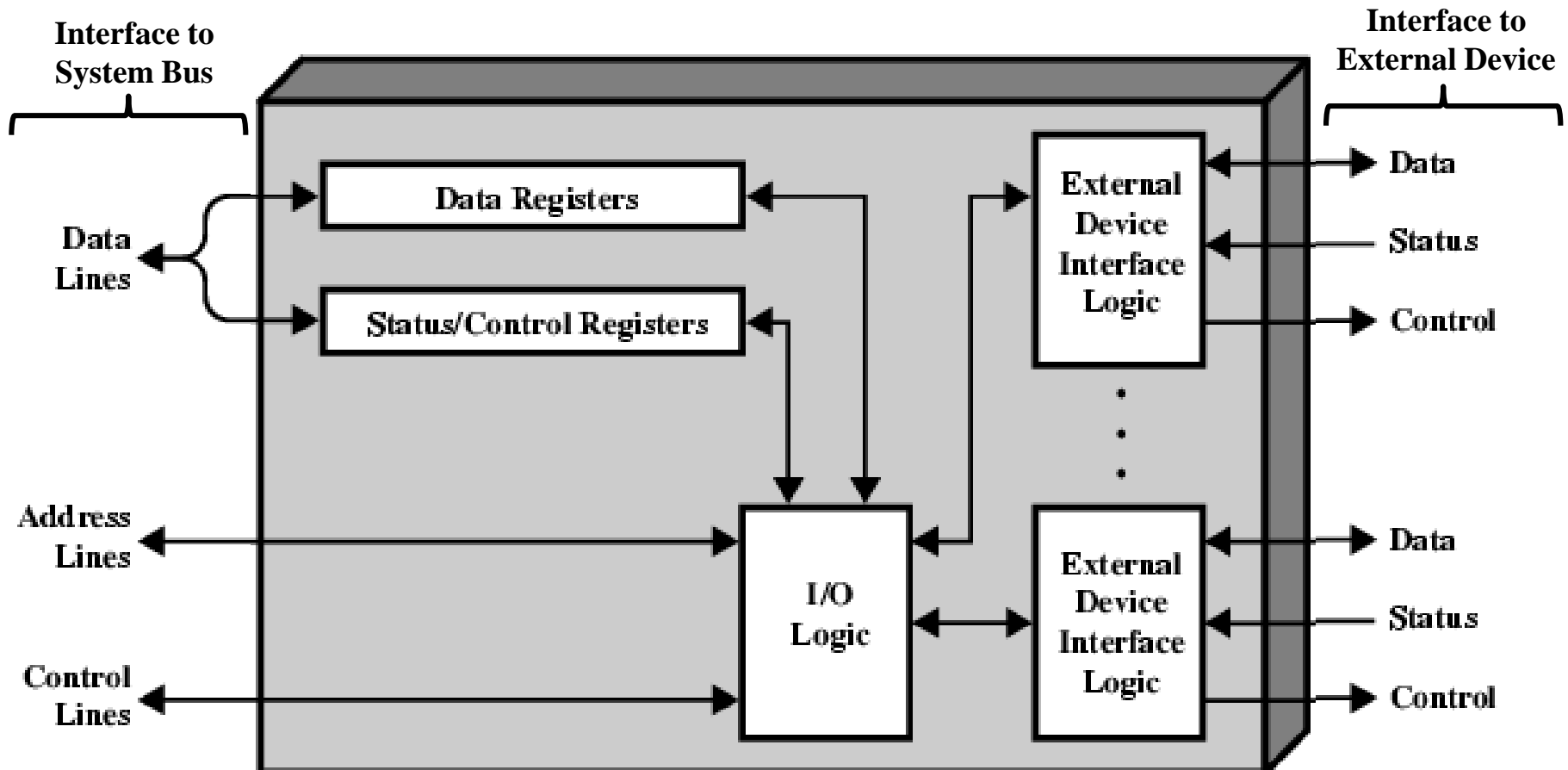


5. Error Detection

- Mechanical and electrical malfunctions
 - Report to CPU.
 - e.g., paper jam, bad disk sector/track.
- Unintentional changes to transmitted bit pattern
 - Detected using error-detecting codes.
 - e.g., parity bit (ASCII).

I/O Module Structure

- St./Ctrl. registers: hold device status or accept control info from CPU.
- CPU issues commands to I/O module via control lines.
- Some control lines are also used by I/O module for bus arbitration.
- Module controlling more than 1 device has a set of unique addresses.



I/O Module Design Decisions

- I/O module lets CPU view a wide range of devices in a simple-minded way.
- Module may **hide device properties** from CPU:
 - Quite **complex module design**.
 - Simple CPU commands** (e.g., render object).
 - Referred to as **I/O channel (I/O processor)**.
 - Common in mainframes.
- Module may **reveal device properties** to CPU:
 - Relatively **simple module design**.
 - Detailed CPU commands** (e.g., rewind tape).
 - Referred to as **I/O controller (device controller)**.
 - Common in microcomputers.

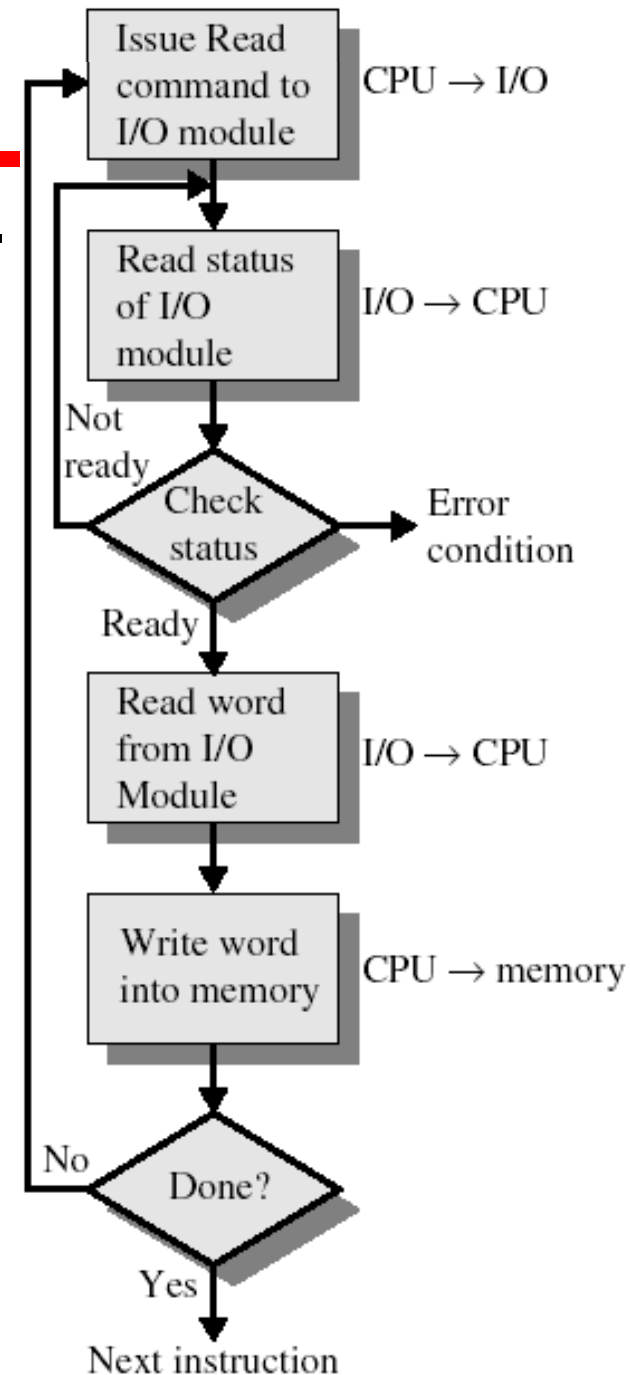
Input Output Techniques

- Programmed I/O.
- Interrupt-driven I/O.
- Direct Memory Access (DMA).

| | No Interrupts | Use of Interrupts |
|------------------------------------|----------------|----------------------------|
| I/O-to-memory transfer through CPU | Programmed I/O | Interrupt-driven I/O |
| Direct I/O-to-memory transfer | | Direct Memory Access (DMA) |

Programmed I/O

- CPU (program) has direct control over I/O.
 - Issuing commands
 - Sensing status
 - Transferring data
- CPU issues a command to I/O module.
- CPU checks status bits periodically.
 - This process is called: **polling**.
- I/O module performs operation.
- I/O module sets status bits.
- CPU transfers data: device \leftrightarrow memory.
- Notes:
 - I/O module does not inform CPU directly.
 - I/O module **does not interrupt** CPU.
 - CPU waits for I/O operation to complete.
 - **Disadvantage: Wasting CPU time!**



I/O Commands vs. I/O Instructions

- I/O Command

- Signal**: issued by (or Sent from) CPU to I/O module.
- Types:
 - Control: activate device & tell it what to do (e.g., rewind tape).
 - Test: check status (e.g., is power on? is error detected?)
 - Read/Write: transfer data CPU \leftrightarrow buffer \leftrightarrow peripheral

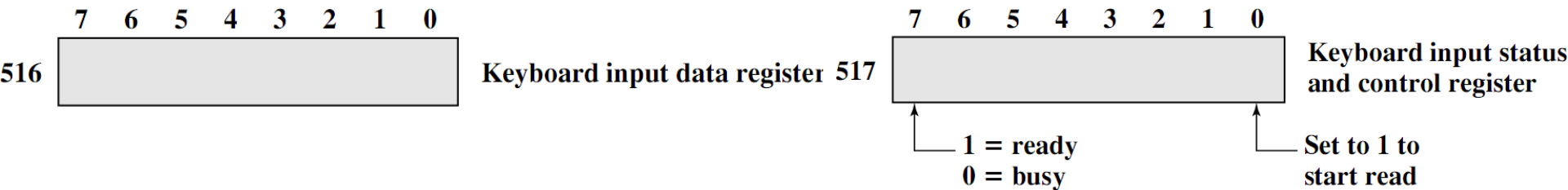
- I/O instruction

- Step in program**: fetched from MM & executed by CPU.
 - To execute an I/O instruction: (1) CPU issues address of I/O module & device, (2) CPU issues an I/O command.
 - Instruction form depends on how devices are addressed.
- In programmed I/O, there is a **one-to-one mapping** between I/O instructions and I/O commands.

Addressing Techniques of I/O Devices

- Two ways to assign addresses to I/O devices:
 1. Memory-mapped I/O
 - Devices & memory **share** same address space.
 - I/O looks just like memory read/write.
 - No **special instructions** for I/O → “load”, “store”, ... etc.
 - Bus has **one Read & one Write** control line.
 - **Pros**: Large selection of memory access instructions.
 - **Cons**: Valuable memory address space is used up!
 2. Isolated I/O
 - **Separate** address space for devices.
 - **Special instructions** for I/O → “in”, “out”, “test”, ... etc.
 - Need two Rd & two Wr control lines.
 - **Pros**: efficient use of memory address space.
 - **Cons**: Not so many I/O instructions.

I/O Mapping - Example



| ADDRESS | INSTRUCTION | OPERAND | COMMENT |
|---------|--------------------|---------|------------------------|
| 200 | Load AC | "1" | Load accumulator |
| | Store AC | 517 | Initiate keyboard read |
| 202 | Load AC | 517 | Get status byte |
| | Branch if Sign = 0 | 202 | Loop until ready |
| | Load AC | 516 | Load data byte |

(a) Memory-mapped I/O

| ADDRESS | INSTRUCTION | OPERAND | COMMENT |
|---------|------------------|---------|------------------------|
| 200 | Load I/O | 5 | Initiate keyboard read |
| 201 | Test I/O | 5 | Check for completion |
| | Branch Not Ready | 201 | Loop until complete |
| | In | 5 | Load data byte |

(b) Isolated I/O

Reading Material

- Stallings, Chapter 6:
 - Pages 215-217
- Stallings, Chapter 7:
 - Pages 222-232