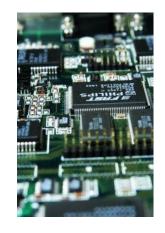
CSE 321a

Computer Organization (1) (1) تنظیم الحاسبات



3rd year, Computer Engineering Fall 2017

Lecture #3



Dr. Hazem Ibrahim Shehata Dept. of Computer & Systems Engineering

Credits to Dr. Ahmed Abdul-Monem Ahmed for the slides

Administrivia

- Office Hours:
 - Sunday 1:00pm 2:00pm
- Tutorials:
 - Day/Time: **Tuesday 8:30am 10:00am**.
- Assignment #1:
 - Released: by the end of this week.
 - Due on: by the end of next week.
 - Work in groups of two.
- Academic integrity:
 - Copying assignments will not be tolerated.
 - Involved groups get 0 points!!

Website: http://hshehata.github.io/courses/zu/cse321a

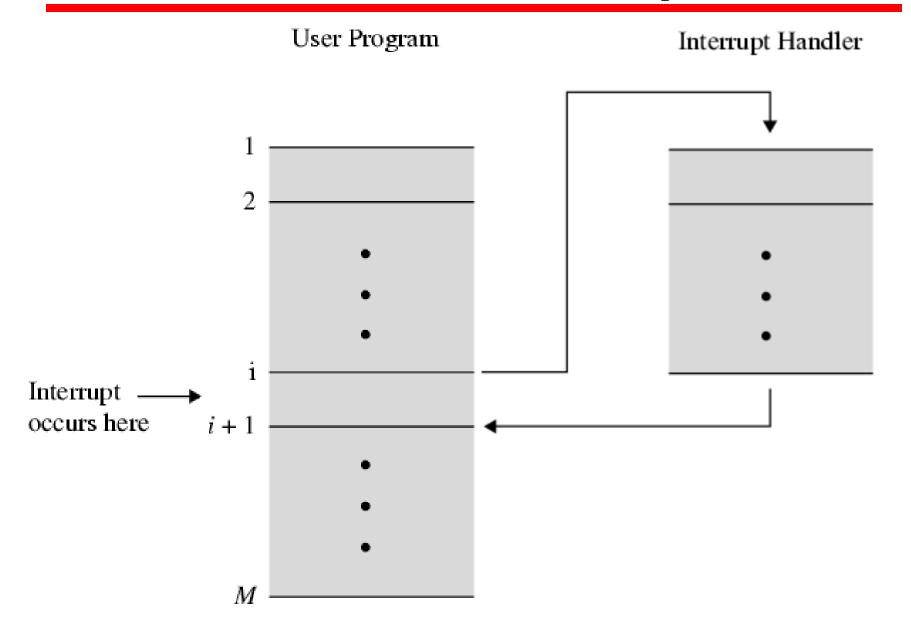
Office hours: Sunday 1:00pm – 2:00pm

Chapter 3. A Top-Level View of Computer Function and Interconnection (Cont.)

Overview

- Computer Structure (Components):
 - -CPU, Memory, I/O.
- Computer Function:
 - —Instruction Fetch/Execute.
 - —Interrupts.
- Interconnection Structure:
 - —Transfer between mem. & proc., i/o & proc. ... etc.
- Bus Interconnection:
 - —Bus structure: address, data, and control lines.
 - —Multiple-bus hierarchies.
 - —Elements of bus design.

Transfer of Control via Interrupts



Multiple Interrupts

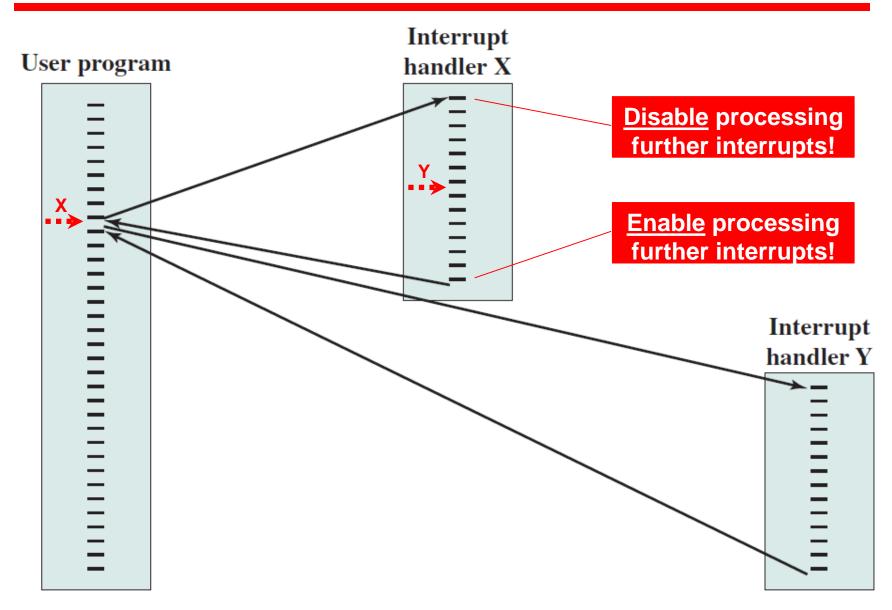
Disable interrupts

- Processor will ignore further interrupts while processing one interrupt.
- —Interrupts remain pending and are checked after first interrupt has been processed.
- —Interrupts handled in **sequence** as they occur.

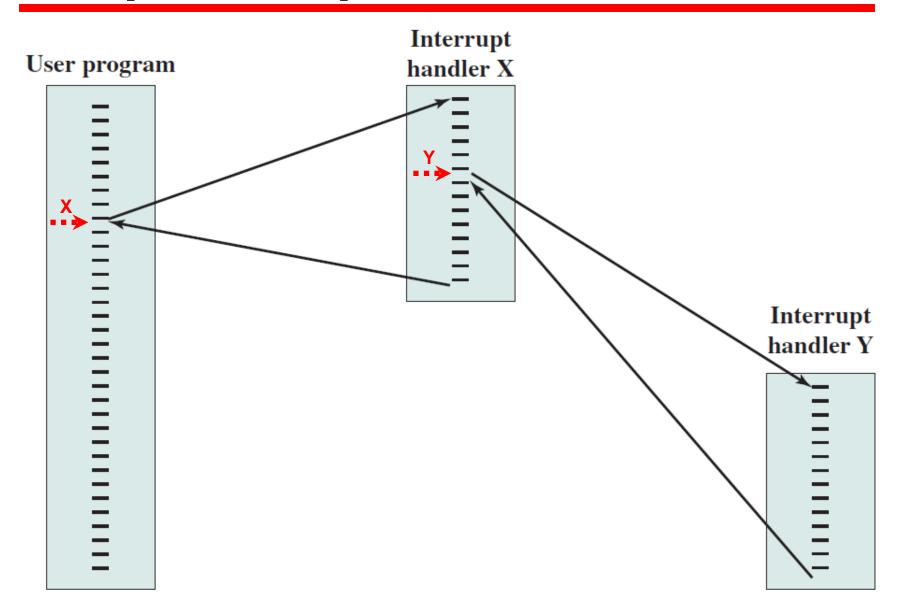
Define priorities

- —Low priority interrupts can be interrupted by higher priority interrupts.
- —When higher priority interrupt has been processed, processor returns to previous interrupt.

Multiple Interrupts - Sequential



Multiple Interrupts – Nested

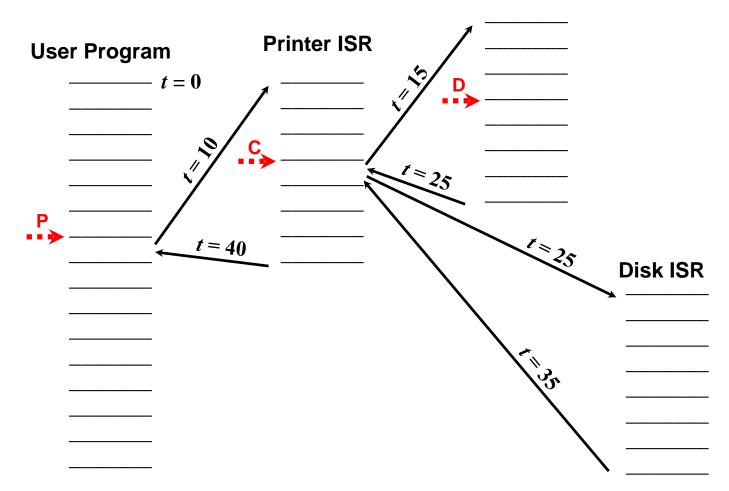


Multiple Interrupts – Nested (example)

Interrupt priorities: Printer (P): 2, Disk (D): 4, Communication line (C): 5

ISR execution: 10 units, Interrupt scenario: P @ t=10, C @ t=15, D @ t=20

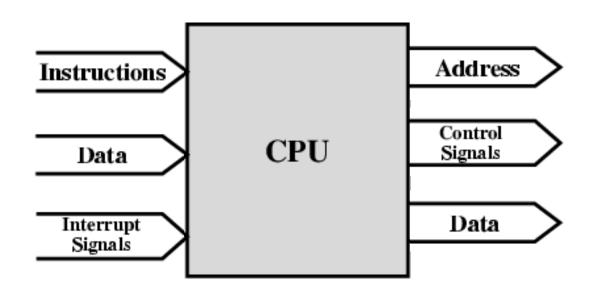
Communication ISR



Interconnection Structures

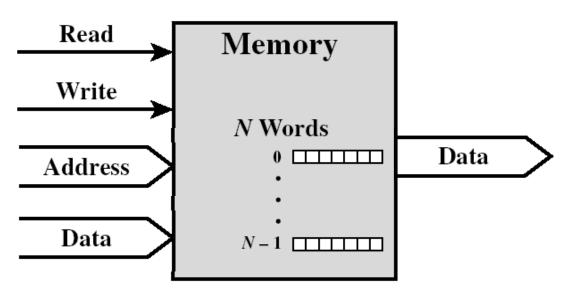
- All computer components/modules must be connected.
- Collection of paths connecting various modules is called: Interconnection Structures
- Different type of connection is required for different type of module.
 - —CPU
 - —Memory
 - —Input/Output

CPU Connection



- Reads instructions and data
- Writes out data (after processing)
- Sends control signals to other units
- Receives (& acts on) interrupts

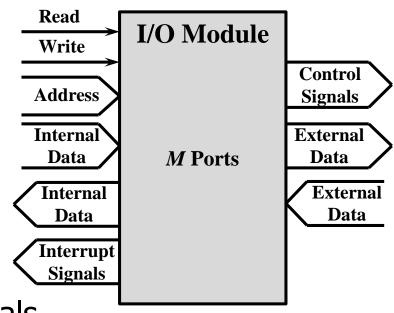
Memory Connection



- Receives control signals.
 - —Read
 - —Write
- Receives addresses (of locations).
- Sends data (read operation).
- Receives data (write operation).

I/O Connection

- Similar to memory from computer's viewpoint.
- Receives control signals from computer.
 - e.g., read and write.
- Receives addresses from computer
 - e.g., port no. to identify peripheral.
- Output
 - Receives data from computer
 - Sends data to peripheral
- Input
 - Receives data from peripheral
 - Sends data to computer
- Sends control signals to peripherals.
 - I/o commands (e.g., spin disk).
- Sends interrupt signals to computer



Types of Transfers

- CPU to memory.
- Memory to CPU.
- CPU to I/O.
- I/O to CPU.
- I/O to memory.
- Memory to I/O.

Direct Memory Access (DMA)

Bus Interconnection

There are a number of possible interconnection systems.

Single and multiple Bus structures are most

common.

Examples of a bus:

—Internal: PCI, SATA

—External: USB

—Internal/external: SCSI





What is a Bus? (1)

- A communication pathway connecting two or more devices.
- Usually by broadcast.
- Consists of multiple lines (each carries 1 bit)
 - —e.g. 32 lines are needed to transfer 32 bits of data in parallel.
- Bus lines typically:
 - —carry address, data, and/or control information
 - —supply power (VCC, GND, ... etc.)
- Power lines may not be shown.

What is a Bus? (2)

Data lines

- Carries data: Remember that there is no difference between "data" and "instruction" at this level!
- Width is a key determinant of performance: 8, 16, 32, 64 bit.

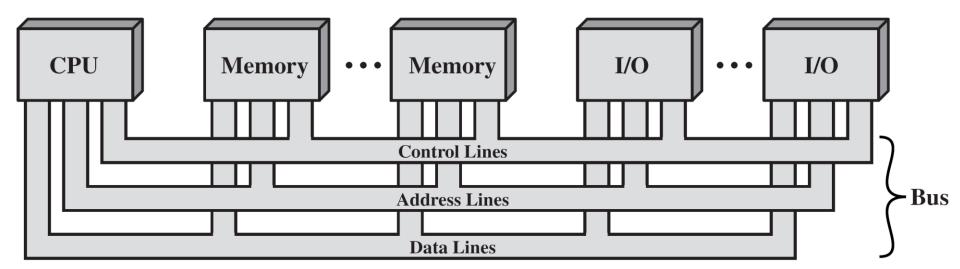
Address lines

- Identify the source or destination of data.
 - e.g. CPU needs to read an instruction (data) from a given location in memory.
- Bus width determines maximum memory capacity of system.
 - e.g. 8080 has 16 bit address bus giving 64k address space.

Control lines

 Control and timing information: memory read/write signal, interrupt request, clock signals

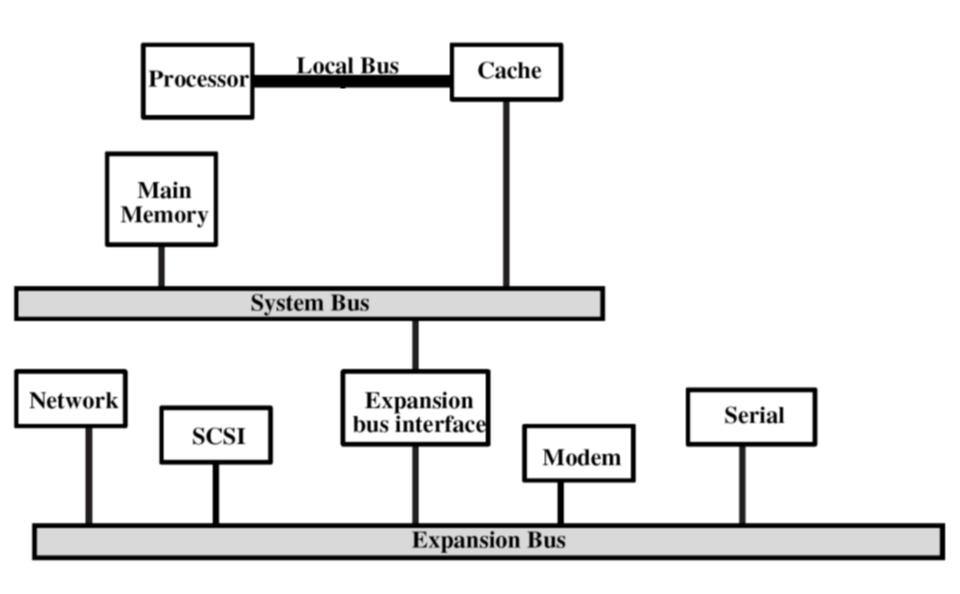
Bus Interconnection Scheme



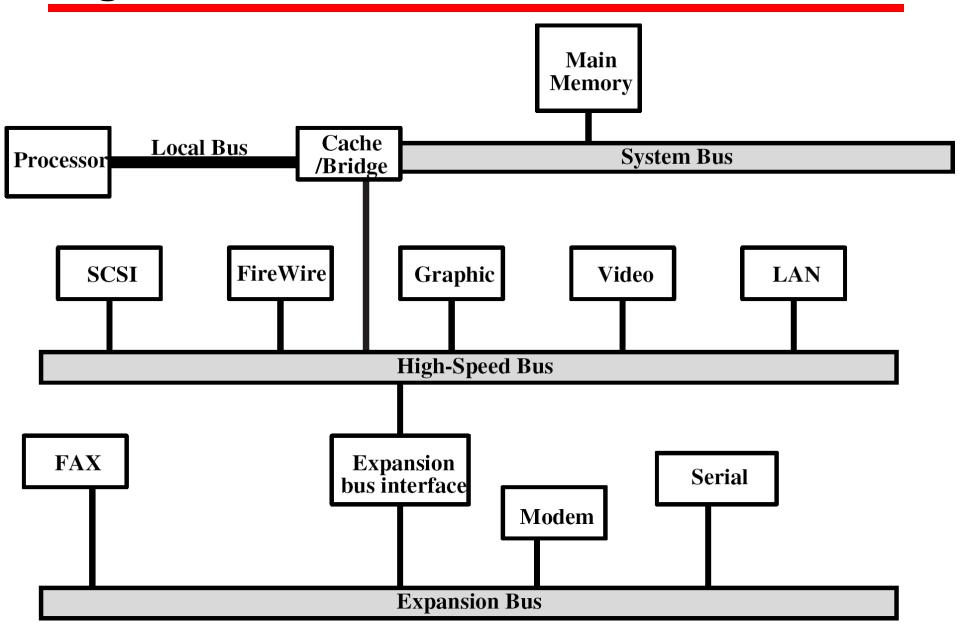
Single Bus Problems

- If great number of devices are connected to bus, performance will suffer. Why?
 - —Greater bus length → greater propagation delay → more time required for device coordination
 - Aggregate data transfer demand approaches bus capacity → bus speed becomes a bottleneck.
- To overcome this problem:
 - —Most systems use multiple buses laid out hierarchically.

Traditional Bus Architecture



High Performance Bus



Elements of Bus Design

- Bus Type
- Method of Arbitration
- Timing
- Data Transfer Type

Bus Types

- Dedicated
 - -Functional: separate data & address lines.
 - Physical: multiple buses connecting different subsets of modules.
- Multiplexed
 - —Shared lines.
 - —Address valid or data valid control line.
 - —Advantages:
 - fewer lines.
 - —Disadvantages
 - More complex control.
 - Lower performance.

Bus Arbitration (I)

- More than one module controlling the bus.
 - —e.g. CPU and DMA controller.
- Only one module may control bus at a time.
- Need a special h/w for arbitration.
- Arbitration may be centralized or distributed.

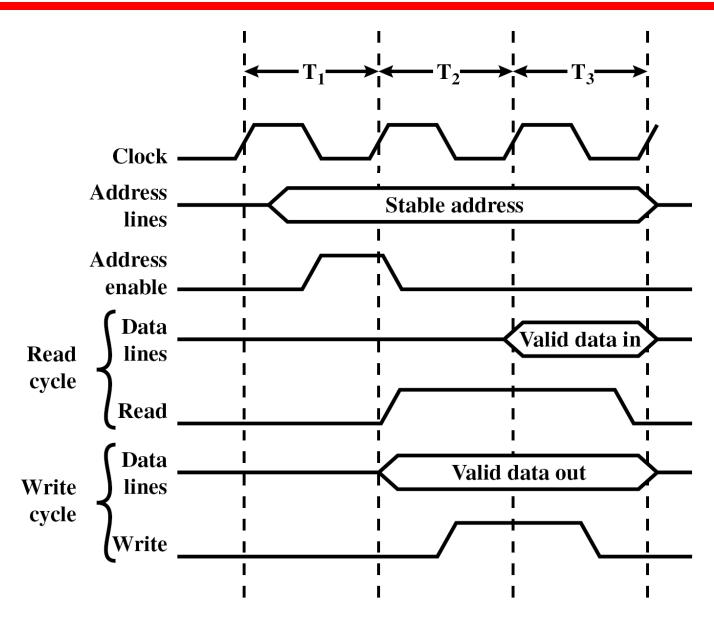
Bus Arbitration (II)

- Centralized Arbitration
 - —Single hardware device controlling bus access
 - Bus controller or arbiter.
 - —May be part of CPU or separate.
- Distributed Arbitration
 - —Each module may claim the bus.
 - —Control logic on all modules.

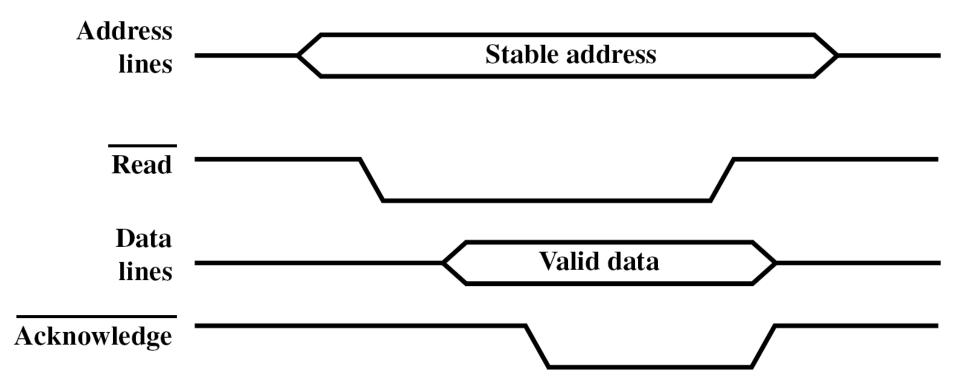
Timing

- Coordination of events on bus.
- Synchronous
 - —Events determined by clock signals.
 - —Clock line belongs to the control bus.
 - —A single 1-0 is a bus cycle.
 - —All devices can read clock line.
 - —Usually sync on leading edge.
 - —Usually a single cycle for an event.
- Asynchronous
 - Occurrence of one event on a bus follows and depends on the occurrence of a previous events.

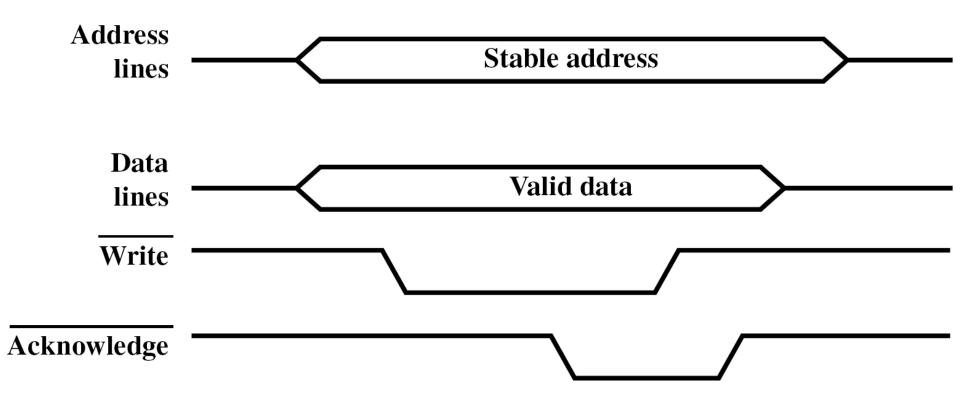
Synchronous Timing Diagram



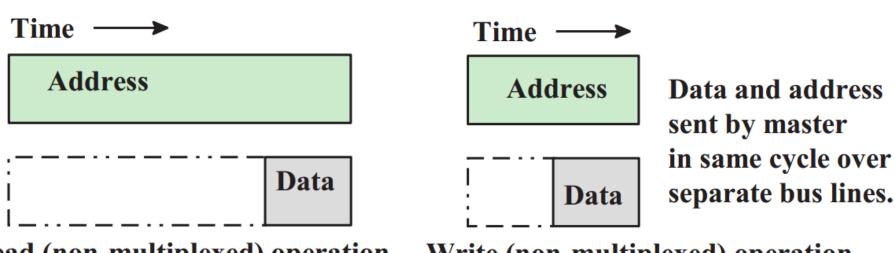
Asynchronous Timing Diagram - Read



Asynchronous Timing Diagram - Write

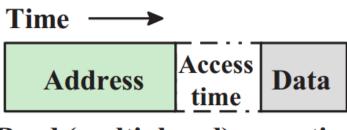


Data Transfer Type (I)



Read (non-multiplexed) operation

Write (non-multiplexed) operation



Read (multiplexed) operation

Time → Data
(1st cycle) (2nd cycle)

Write (multiplexed) operation

Data Transfer Type (II)

Time>				
Address		Data read	Data write	
Read-modify-write operation				
Time				
Address	Data		Data	
	write	L	read	
Read-after-write operation				
Time				
Address	Data	Data	Data	
11441 033	Dutu	Dutu	Dutu	L

Block data transfer

Reading Material

- Stallings, Chapter 3:
 - —Pages 80 93