

CSE 321b

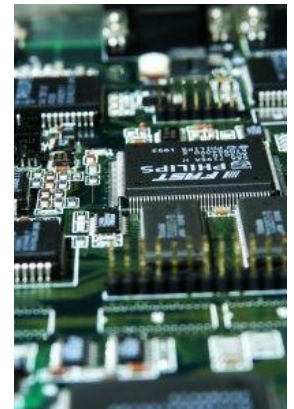
Computer Organization (2)

تنظيم الحاسب (2)



3rd year, Computer Engineering
Winter 2017

Lecture #5



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Dept. of Computer & Systems Engineering

Credits to Dr. Ahmed Abdul-Monem Ahmed for the slides

Adminstrivia

- Assignment #1:
 - Deadline extended again!
 - New due date: **Sunday, Mar. 14, 2017!**

Website: <http://hshehata.github.io/courses/zu/cse321b/>

Office hours: TBA

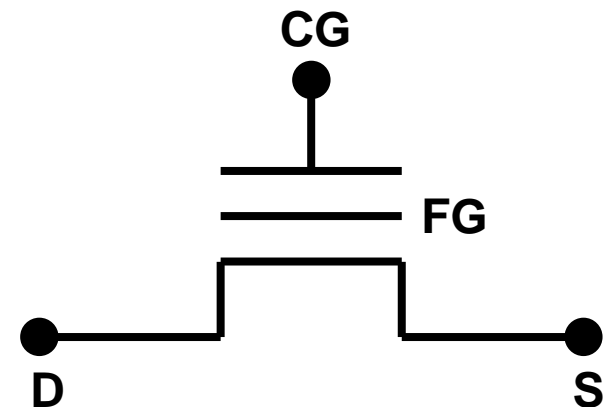
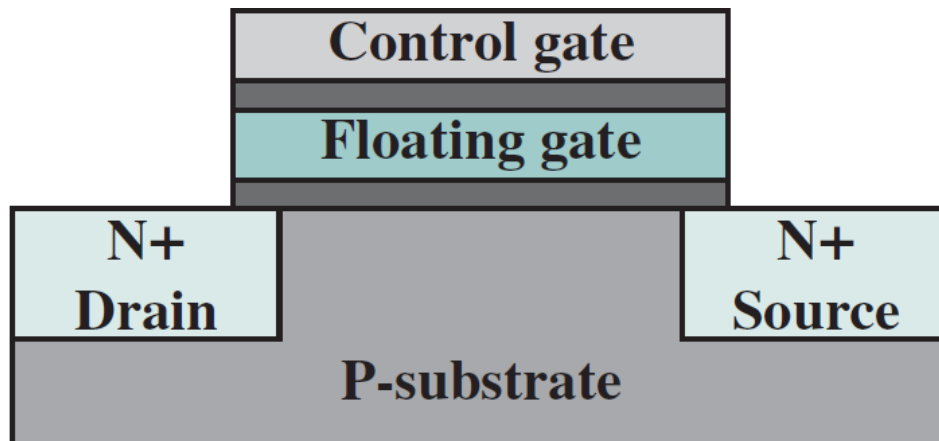
Chapter 6. External Memory (*Cont.*)

Types of External Memory

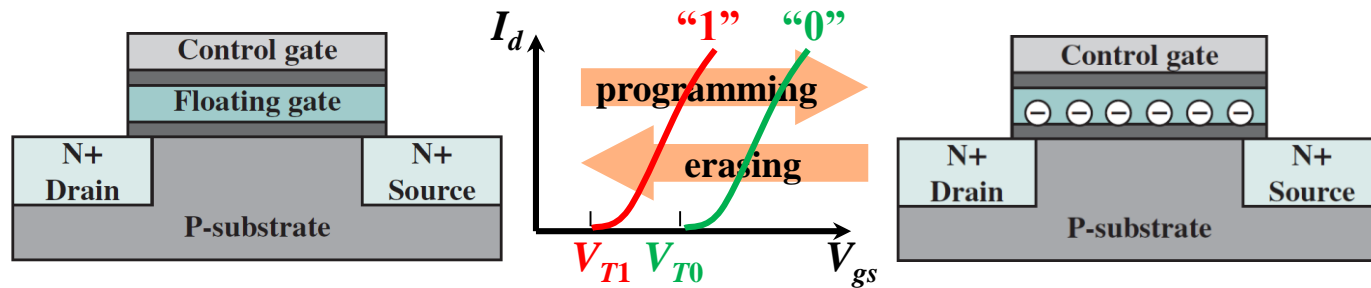
- Magnetic Disk
- Redundant Array of Independent Disks (RAID)
- Optical Disk
- Solid-State Drive (SSD)
- Magnetic Tape

Solid-State Drive (SSD)

- **Purpose:** complement or even replace HDDs!
- The term “solid-state” refers to electronic circuitry built with semiconductors.
- SSD's store data in **flash memory** cells.
- Each flash memory cell is built using a **single** transistor: **floating-gate MOSFET (FG-MOSFET)**.



States of FG-MOSFET



Logic 1

- No electrons trapped on the FG.
- **Smaller** threshold voltage (V_{T1})
 - Forming the channel (i.e., turning transistor on) requires applying a relatively **smaller** voltage to CG.

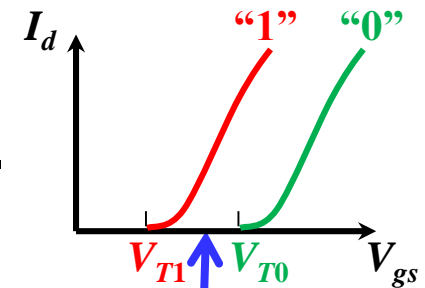
Logic 0

- Electrons trapped on the FG.
- **Higher** threshold voltage (V_{T0})
 - Forming the channel (i.e., turning transistor on) requires applying a slightly **higher** voltage to CG.

Read/Program/Erase Flash Memory Cell

- To read:

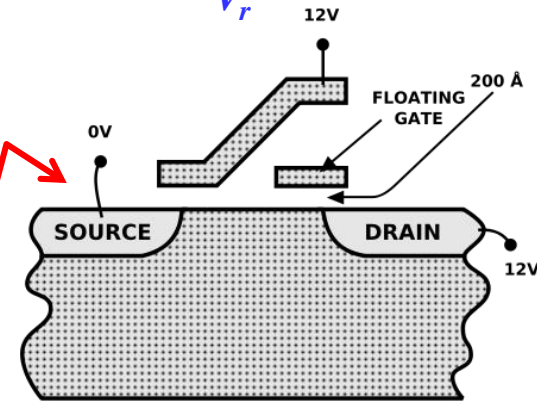
- Apply voltage V_r to CG s.t. $V_{T1} < V_r < V_{T0}$.
- Measure (sense) drain current (I_d).
 - $I_d > 0 \rightarrow$ transistor on \rightarrow logic 1.
 - $I_d = 0 \rightarrow$ transistor off \rightarrow logic 0.



- To program (write "0"):

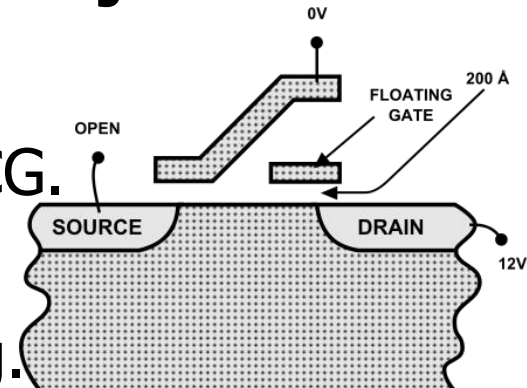
Only Applicable to NOR flash!!

- Apply high +ve voltage to CG & D.
- Electrons jump from channel through insulating layer onto FG \rightarrow **hot-electron injection.**



- To erase (write "1"):

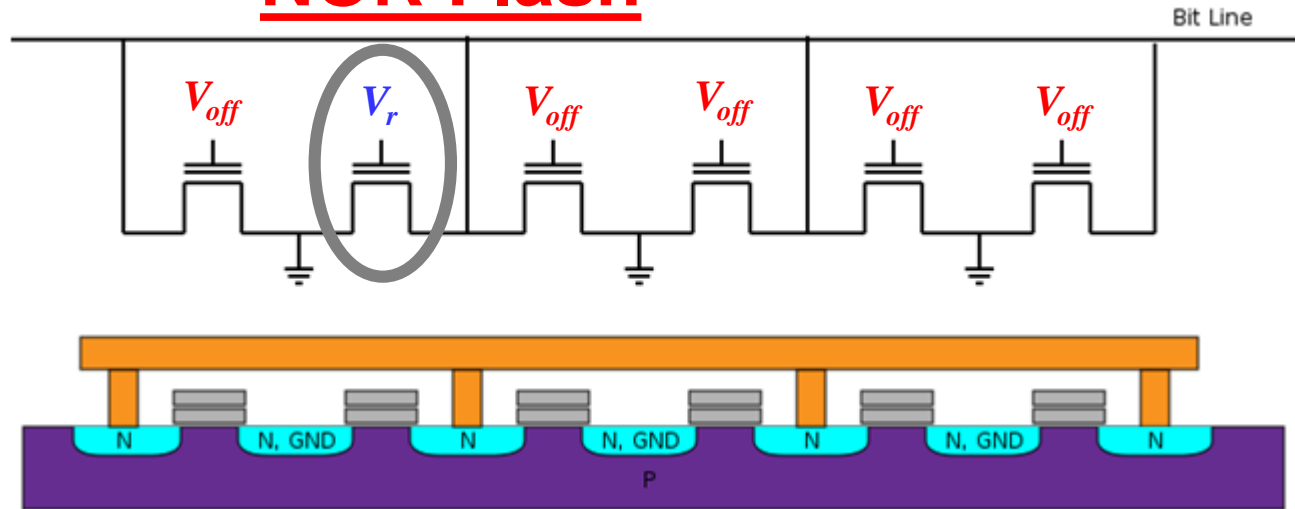
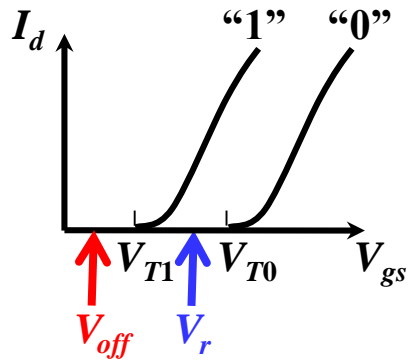
- Apply high voltage diff. between D & CG.
- Electrons pulled off FG through insulating layer to D \rightarrow **quantum tunneling.**



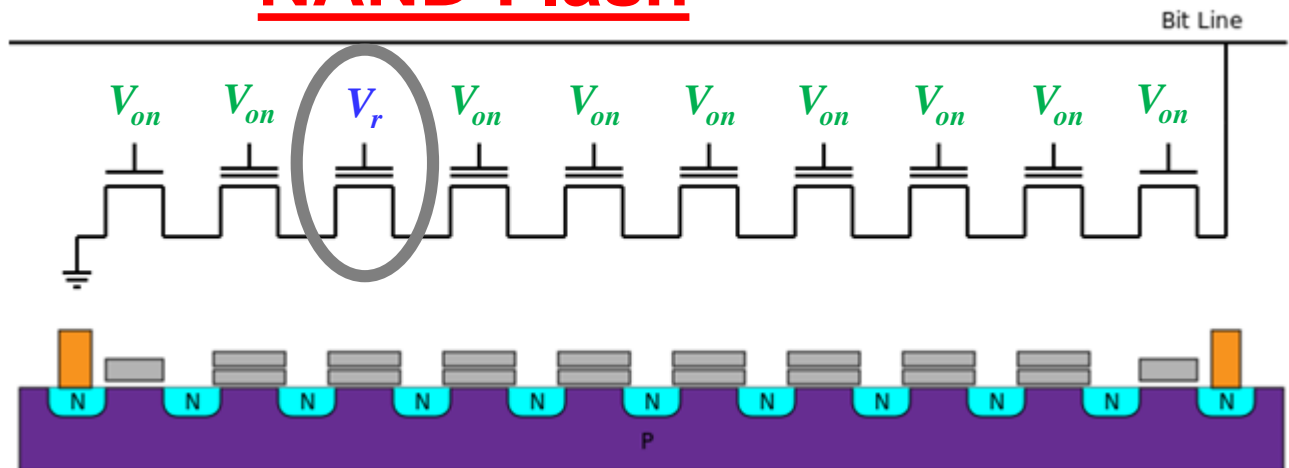
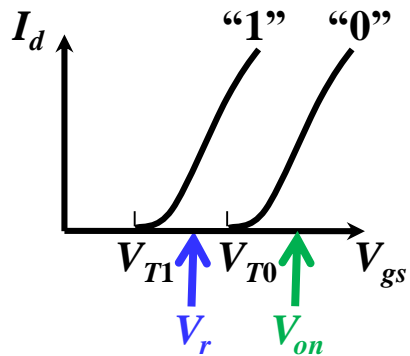
- ## Factor program/erase NAND Flash

NAND/NOR Flash – Reading

NOR Flash



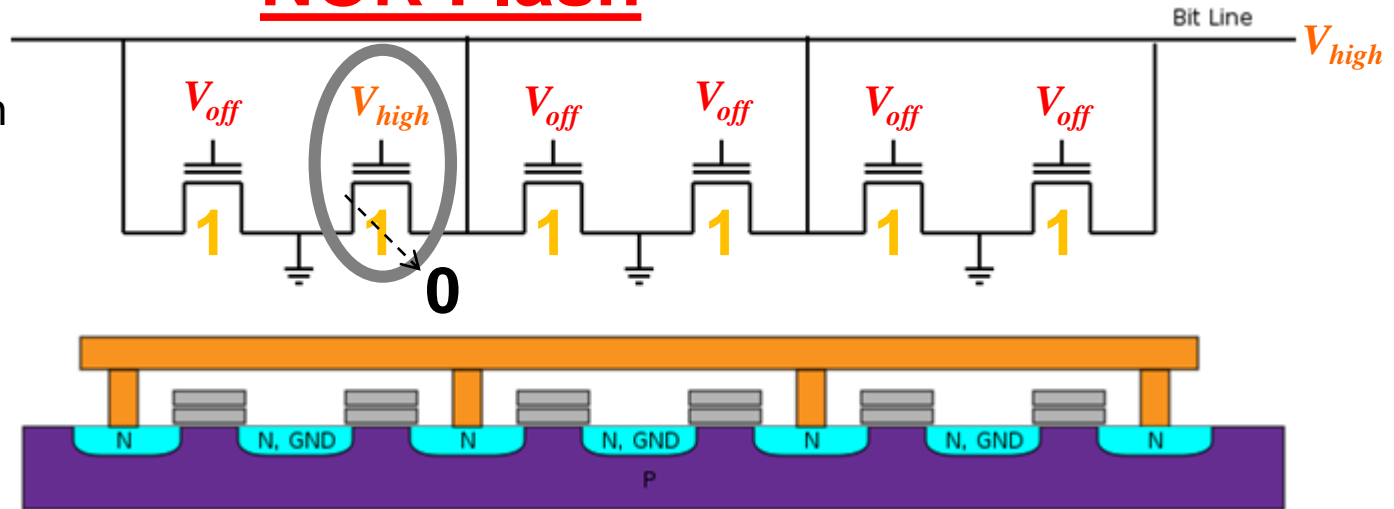
NAND Flash



NAND/NOR Flash – Programming (X → 0)

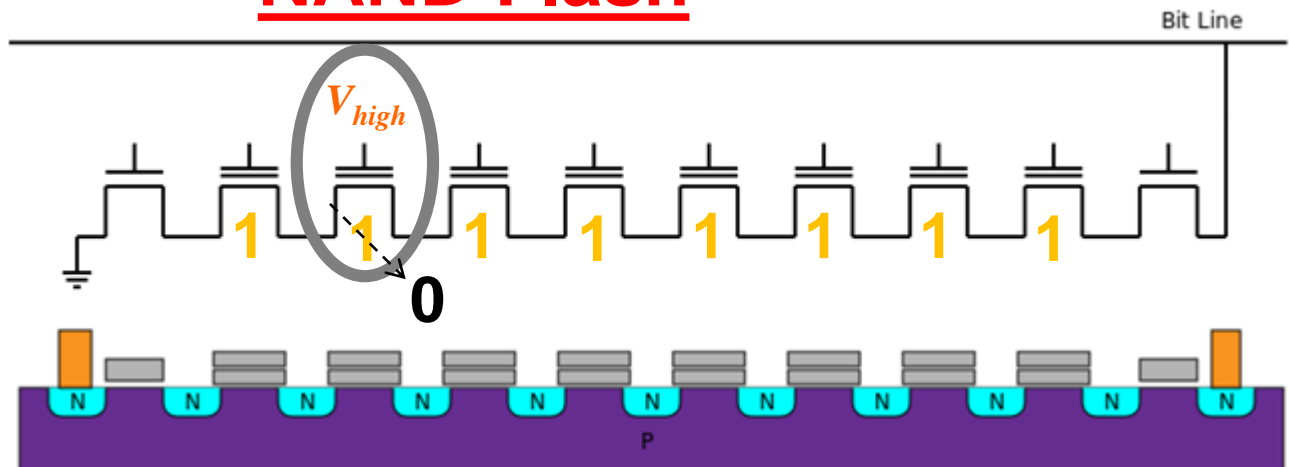
NOR Flash

- Technique:
 - Hot electron injection.
- Explanation:
 - Electrons jump from channel to FG.



NAND Flash

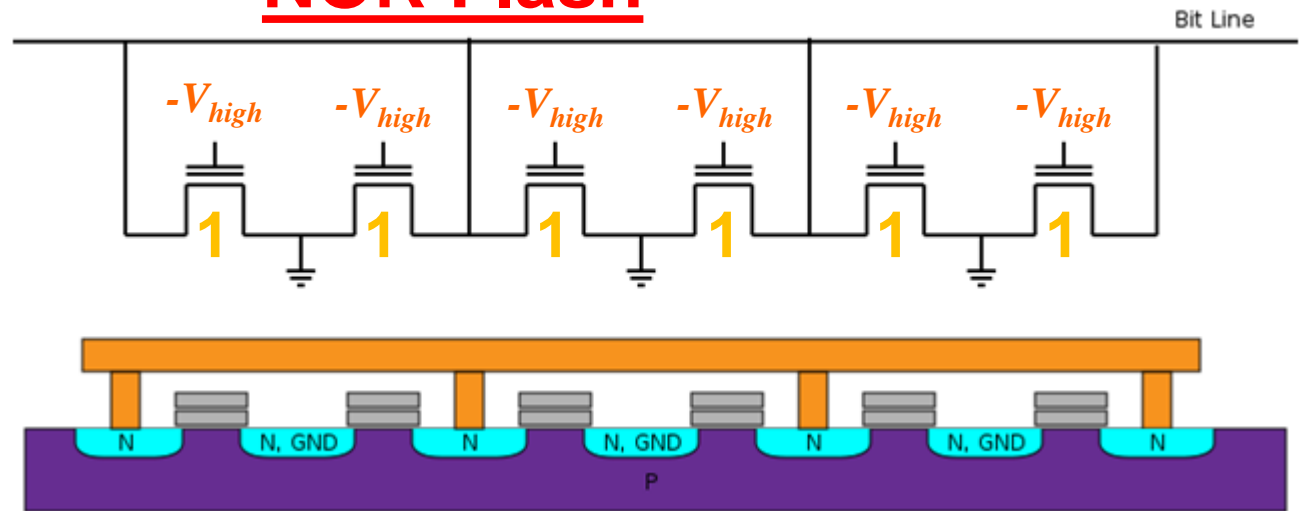
- Technique:
 - Quantum tunneling (injection).
- Explanation:
 - Electrons tunnel from B to FG.



NAND/NOR Flash – Block Erasure ($X \rightarrow 1$)

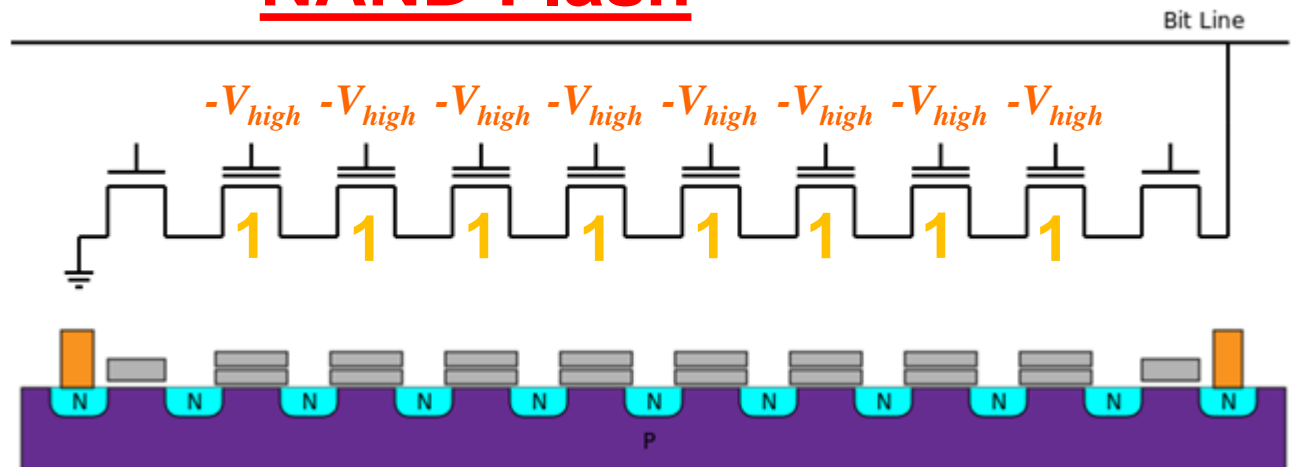
NOR Flash

- Technique:
 - Quantum tunneling (ejection).
- Explanation:
 - Electrons tunnel from FG to S/B.



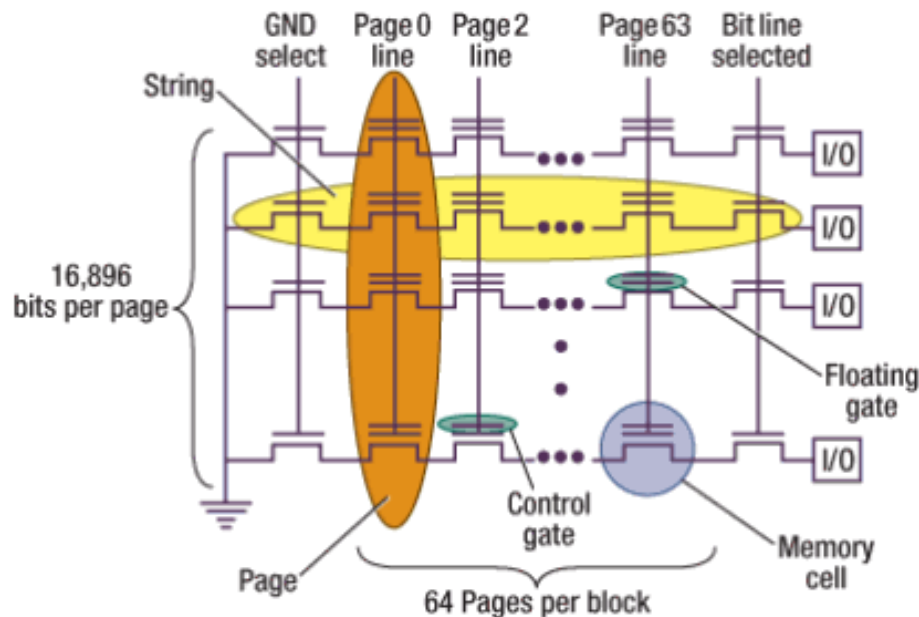
NAND Flash

- Technique:
 - Quantum tunneling (ejection).
- Explanation:
 - Electrons tunnel from FG to B.



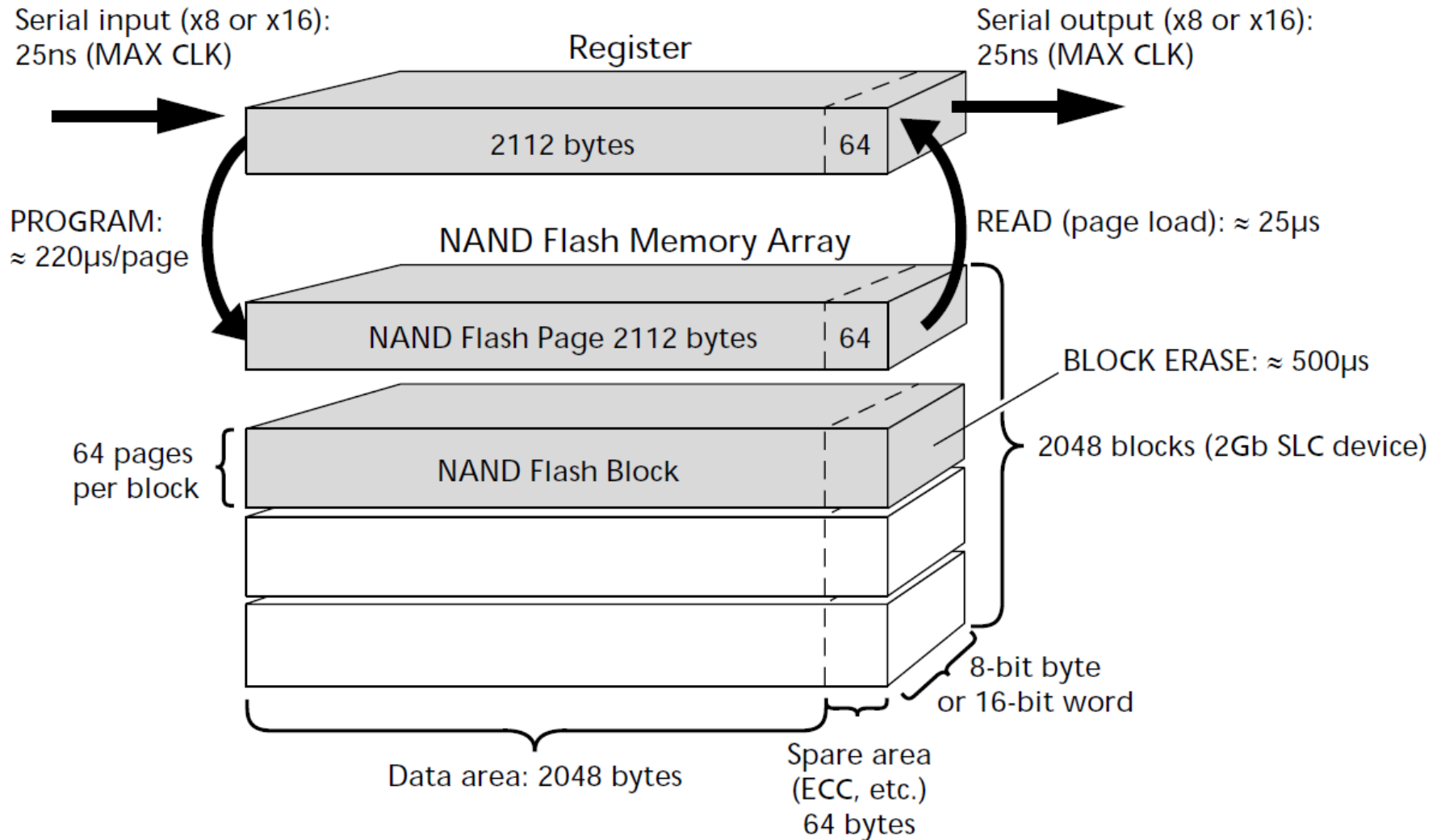
NAND Flash Block Organization

- NAND flash memory is organized as a set of independent **blocks**. Each block has a set of **pages**.
- **Blocks** are the smallest **erasable** units.
- **Pages** are the smallest **programmable** units.
 - Partial pages can be programmed in some devices!



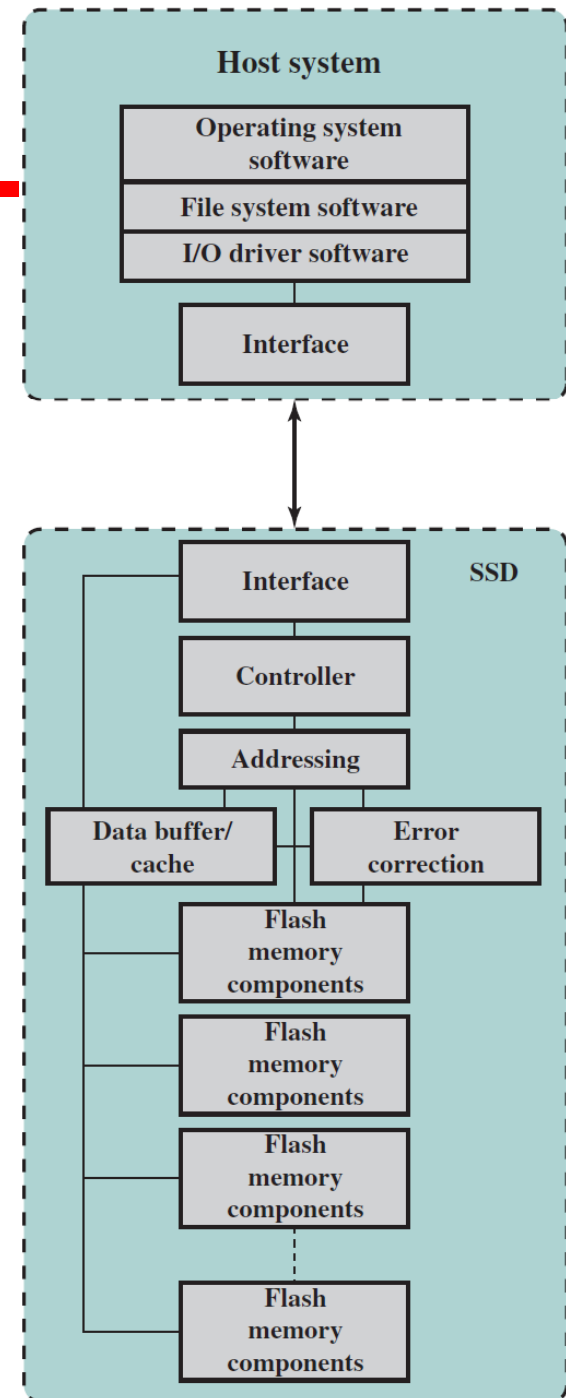
**** The numbers of bits-per-page and pages-per-block here are typical for a 2Gb NAND flash device ****

NAND Flash Memory Architecture








SSD Architecture

- **Controller**: provides SSD device level interfacing & firmware execution.
- **Addressing**: logic to select one of the flash memory components.
- **Data buffer/cache**: High-speed RAM for speed matching and increasing data throughput.
- **Error correction**: Logic for error detection and correction.
- **Flash memory components**: Individual NAND flash chips.



SSD vs. HDD

- SSDs have the following advantages over HDDs:
 - Higher input/output operations per second (IOPS)
 - Longer lifespan: no mechanical wear.
 - Lower power consumption.
 - Quieter and cooler running capabilities.
 - Lower access times & latency rates: >10x faster.

	NAND Flash Drives	Disk Drives
I/O per second (sustained)	Read: 45,000  Write: 15,000	300
Throughput (MB/s)	Read: 200+  Write: 100+	up to 80
Access Time (ms)	0.1 	4–10
Storage capacity (year: 2014)	up to 2 TB	up to 8 TB 
Cost per capacity (year: 2014)	\$0.45/GB	\$0.05/GB 

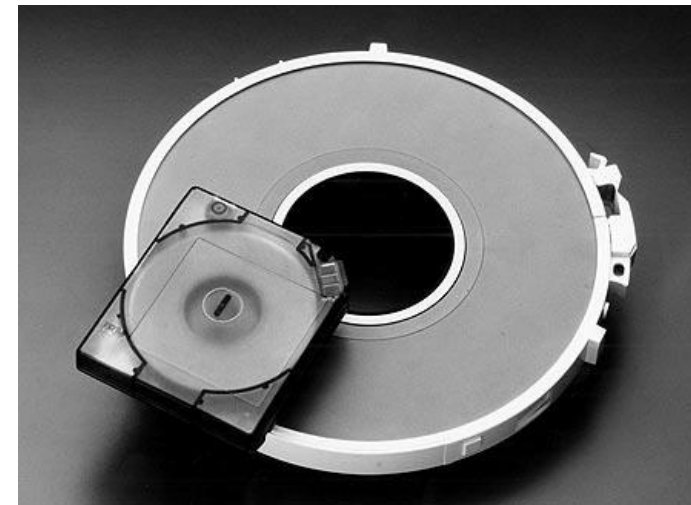
Types of External Memory

- Magnetic Disk
- Redundant Array of Independent Disks (RAID)
- Optical Disk
- Solid-State Drive (SSD)
- Magnetic Tape

Magnetic Tape

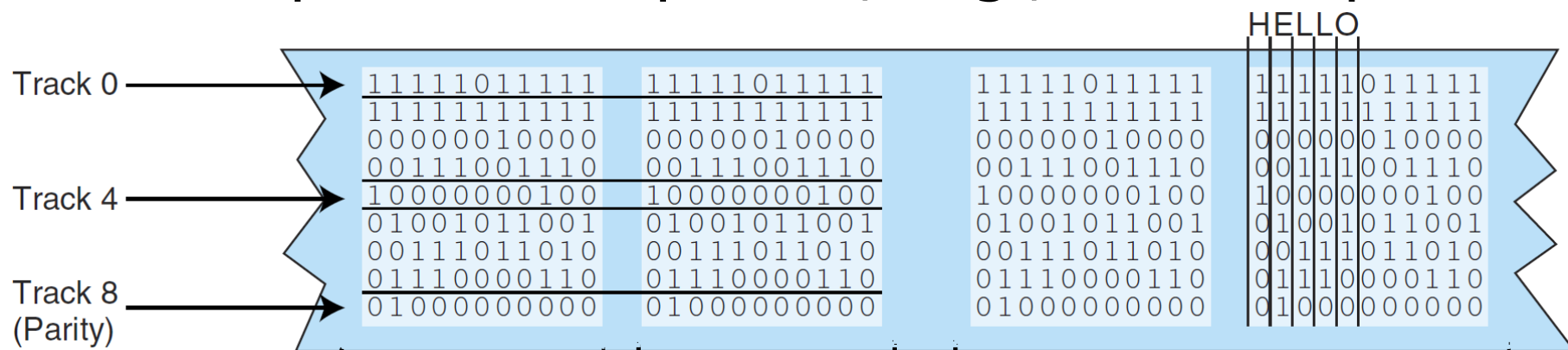


- **Features:**
 - **Oldest** secondary memory (1951).
 - **Slowest-speed** & **lowest-cost** in memory hierarchy.
- **Usage:** **offline storage** (i.e., backup).
- **Access method:** **sequential** → slow!
- **Medium:** flexible polyester tape coated with magnetizable material (metal particles).
- **Tape width:** 0.38cm - 1.27cm.
- **Tape package:**
 - Used to be **open reels**.
 - Now housed in **cartridges**.



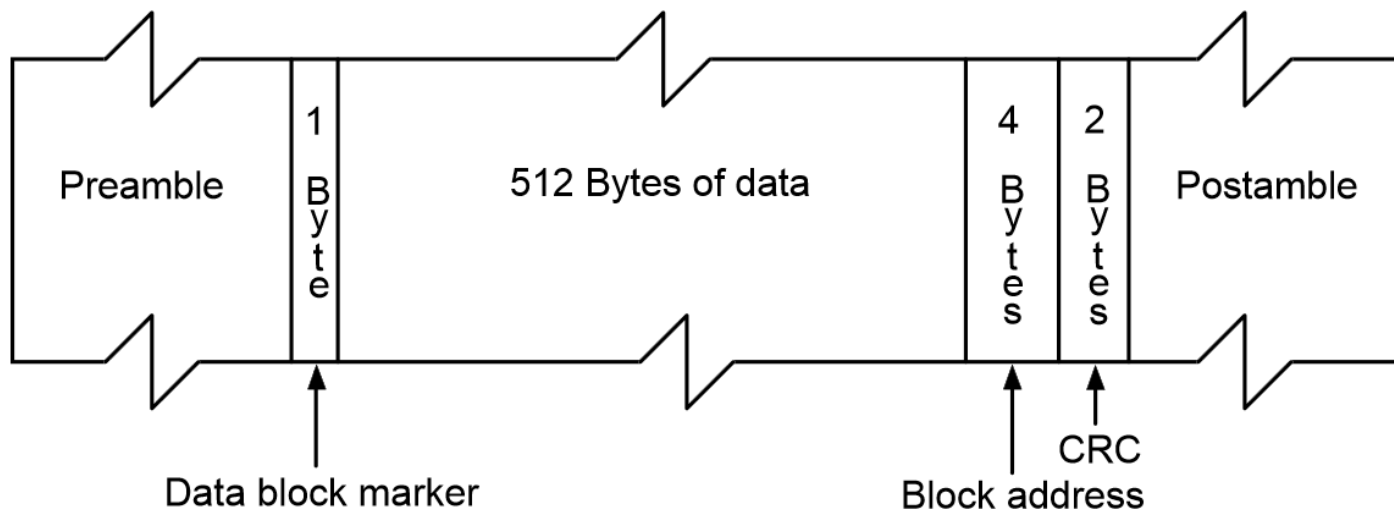
Data Layout and Recording Techniques

- Data are laid out on tape along **parallel tracks** running lengthwise → **linear recording**.
- Tape moves underneath **magnetic head** in drive.
- Head has multiple (8-32) read/write elements → read/write multiple (8-32) tracks simultaneously.
- There are two types of **linear recording**:
 1. **Parallel recording**: Each byte/word is stored on multiple tracks in parallel, .e.g., 9-track tapes.



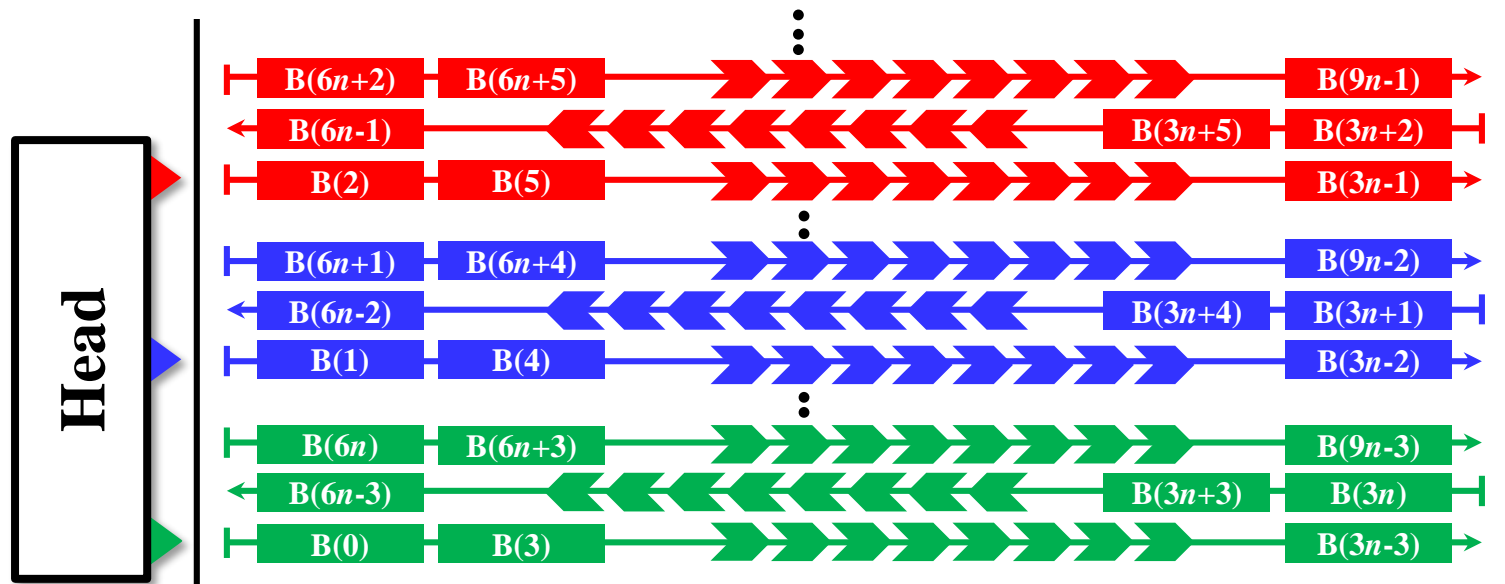
Serial Recording

- ... types of **linear recording** ... (cont.):
- 2. **Serial Recording**: data laid out as sequence of bits along each track, similar to magnetic disks.
 - Data read/written in contiguous blocks called **records**.
 - Records are separated by inter-record gaps.
 - Tape is formatted to assist in locating records.



Serpentine Recording

- Typical recording technique in serial tapes.
 - First set of bits is recorded on tracks in **forward direction**.
 - When tape reaches end, rd/wr elements are repositioned to new tracks & recording is resumed in **opposite direction**.
 - Process continues, back and forth, until tape is full.
 - Data still recorded **serially along individual tracks**, but blocks in sequence are stored in adjacent tracks.
 - **Adv.:** more tracks than head/elements → increasing capacity.



Linear Tape-Open (LTO)

- Dominant magnetic tape technology today.
- Developed late 1990s as an **open source** alternative to proprietary tape systems.

	LTO-1	LTO-2	LTO-3	LTO-4	LTO-5	LTO-6	LTO-7	LTO-8
Release date	2000	2003	2005	2007	2010	2012	2015	TBA
Compressed capacity	200 GB	400 GB	800 GB	1600 GB	3000 GB	6250 GB	16 TB	32 TB
Compressed transfer rate	40 MB/s	80 MB/s	160 MB/s	240 MB/s	280 MB/s	400 MB/s	788 MB/s	1.18 GB/s
Linear density (bits/mm)	4880	7398	9638	13250	15142	15143		
Tape tracks	384	512	704	896	1280	2176	3584	
Tape length (m)	609	609	680	820	846	846	960	
Tape width (cm)	1.27	1.27	1.27	1.27	1.27	1.27	1.27	
Write elements	8	8	16	16	16	16	32	
WORM?	No	No	Yes	Yes	Yes	Yes	Yes	Yes
Encryption Capable?	No	No	No	Yes	Yes	Yes	Yes	Yes
Partitioning?	No	No	No	No	Yes	Yes	Yes	Yes

Chapter 7. Input / Output

Outline

- External Devices
 - Types
 - Structure
- I/O Modules
 - Function
 - Structure
- I/O Techniques
 - Programmed I/O
 - Interrupt-Driven I/O
 - Direct Memory Access
- I/O Channels & Processors

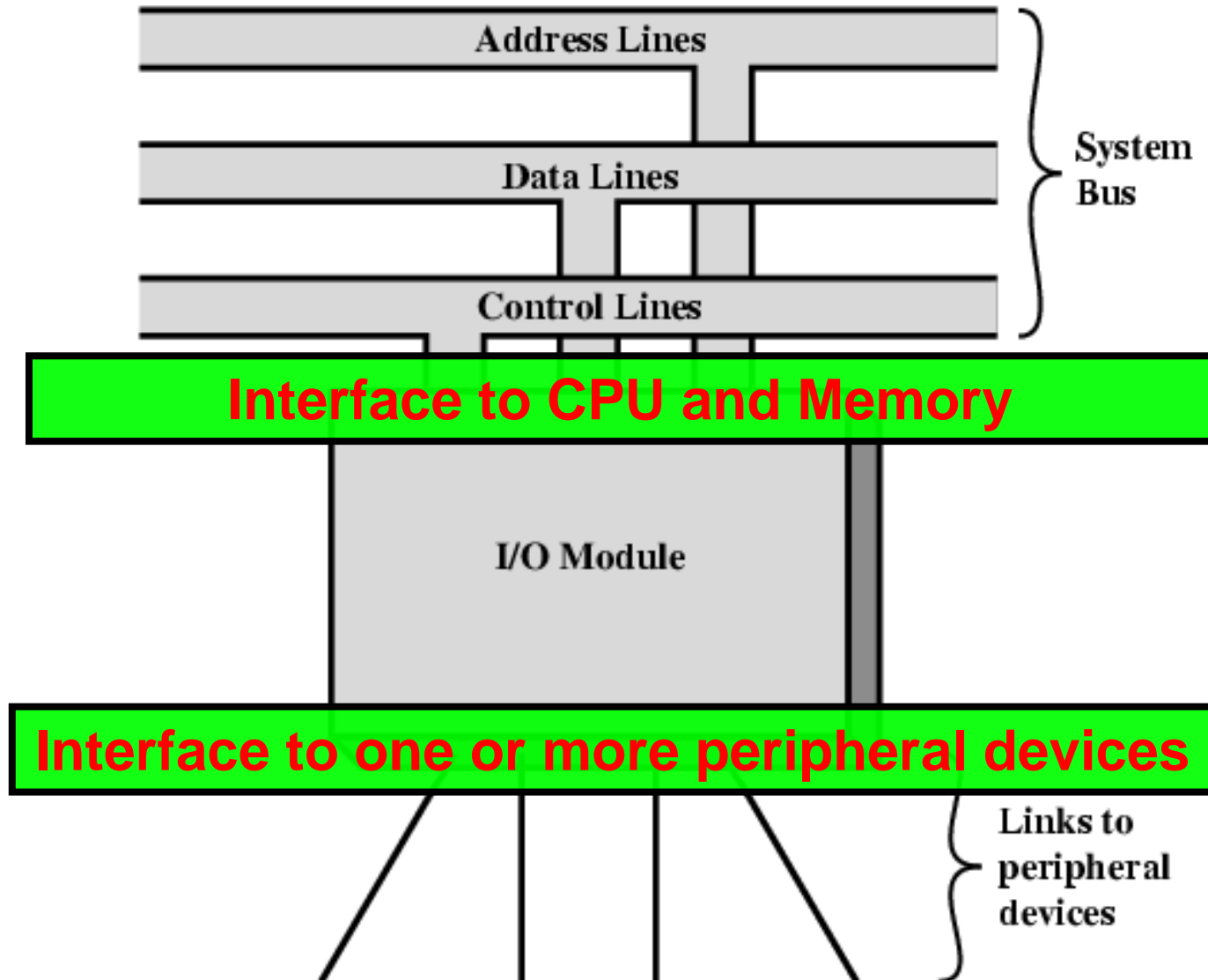
Terms

- Essential Computer Units
 - CPU and Memory
- Peripheral (or External or I/O) devices
 - Any device attached to a computer in order to increase its functionality.
 - Input: keyboard, mouse, scanner, ... etc.
 - Output: printers, speakers, ... etc.
 - Input and output: hard disk, modem, ... etc.
- I/O (Input/Output) Operations
 - Transfer of data to/from computer from/to peripheral device (done by program, operation, or device).
 - Input: from a device to the computer
 - Output: from the computer to a device.

Input/Output Problems

- There is a wide variety of peripherals!
 - Different methods of operation (H/W).
 - Delivering different amounts of data.
 - At different speeds (which are also different from CPU and memory).
 - In different formats (e.g., word length).
- **Conclusion:** Hard to connect such variety of different devices directly to same Bus!!
- **Solution:** **I/O Module**

I/O Module



Types of Peripherals

- Human readable
 - Screen, printer, keyboard, ... etc.
- Machine readable
 - Magnetic disk, tape, ... etc.
- Communication
 - Modem, Network Interface Card (NIC), Wireless Network Adaptor, ... etc.

Peripheral (External) Device

- **Control Signals**

- Send data to module, receive data from module, send status, position disk head.

- **Status Signals**

- READY, NOT READY, ...

- **Buffer**

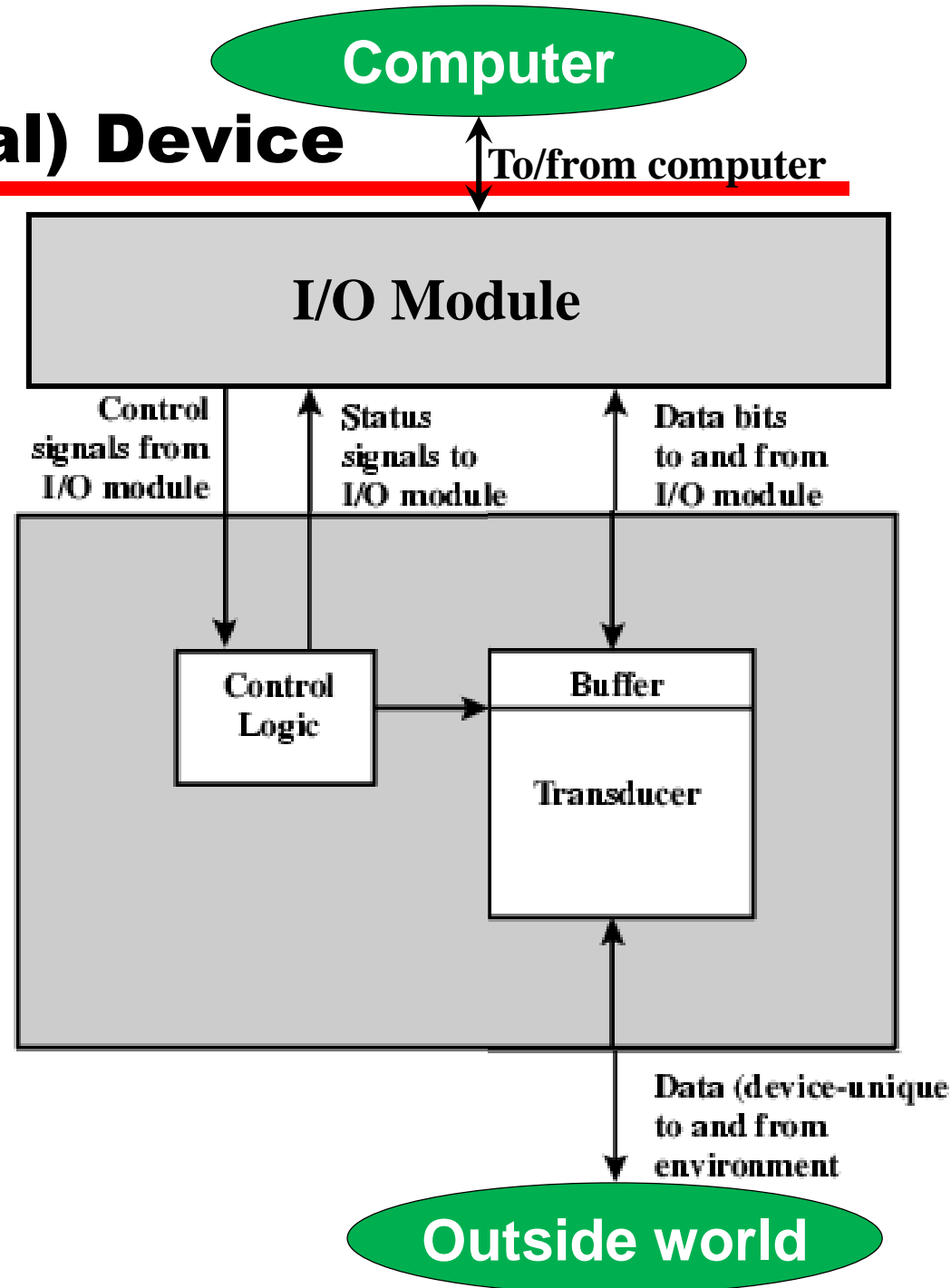
- Temporarily hold data being transferred. Size: x bytes → x Kbytes!!

- **Transducer**

- Converts energy: electrical ↔ other.

- **Control logic**

- Controls operation.



Examples: Keyboard/Monitor, and Disk Drive

- **Keyboard** (input)
 - A key is pressed.
 - Transducer translates signal into ASCII.
 - ASCII is transmitted to I/O module in the computer.
 - Text can be stored as ASCII in the computer.
- **Monitor** (output)
 - Computer sends ASCII to I/O module. I/O module sends ASCII to external device (monitor).
 - Transducer at the monitor sends electronic signals to display the character.
- **Hard Disk Drive** (input/output)
 - Head moves in and out across disk surface.
 - Transducer converts magnetic patterns to/from bits.

Functions of I/O Module

1. Control & Timing.
2. CPU Communication.
3. Device Communication.
4. Data Buffering.
5. Error Detection.

1. Control & Timing

- I/O includes **control & timing** requirement to coordinate the flow of traffic between internal resources (CPU, MM, ...) and external devices.
- Ex.: Transfer data from input device to CPU:
 1. CPU checks I/O module device status.
 2. I/O module returns status.
 3. If ready, CPU requests data transfer (command to I/O module).
 4. I/O module gets data from external device.
 5. I/O module transfers data to CPU.
- If transfer goes through bus, each CPU/module interaction may involve 1+ bus arbitrations.

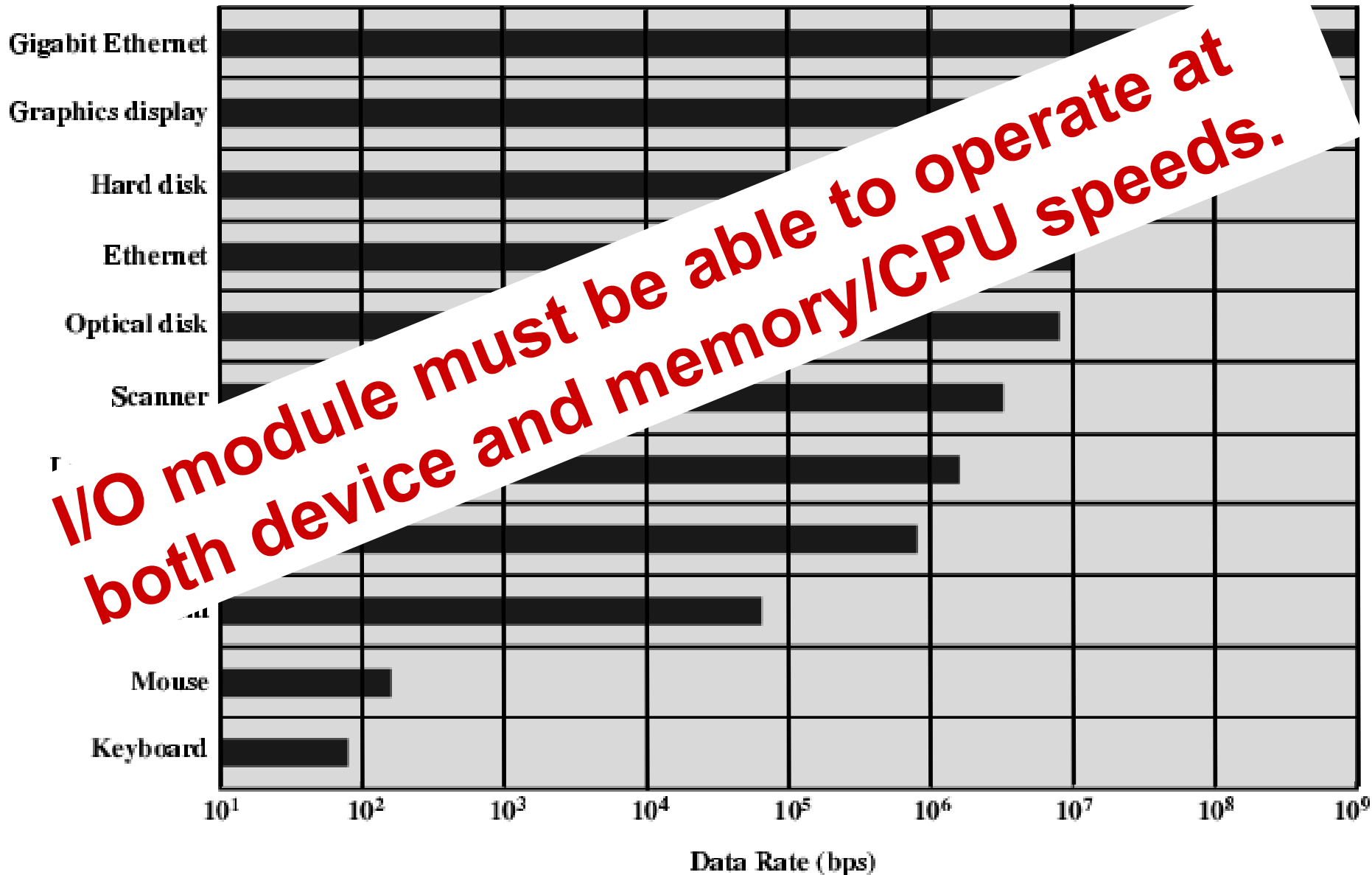
2. CPU Communication

- CPU communication involves the following:
 - Command decoding
 - Module accepts commands from CPU on control lines.
 - Command parameters can be sent over data line.
 - e.g., SEEK track in a disk drive: SEEK command sent on control lines and track # sent on data lines.
 - Address recognition
 - One unique address for each peripheral it controls.
 - Data exchange
 - Between CPU and device over the data bus.
 - Status reporting
 - BUSY, READY, or some error conditions.

3. Device Communication

- I/O module must also be able to do device communication:
 - Commands
 - Status information
 - Data

4. Data Buffering (Speed Mismatch)

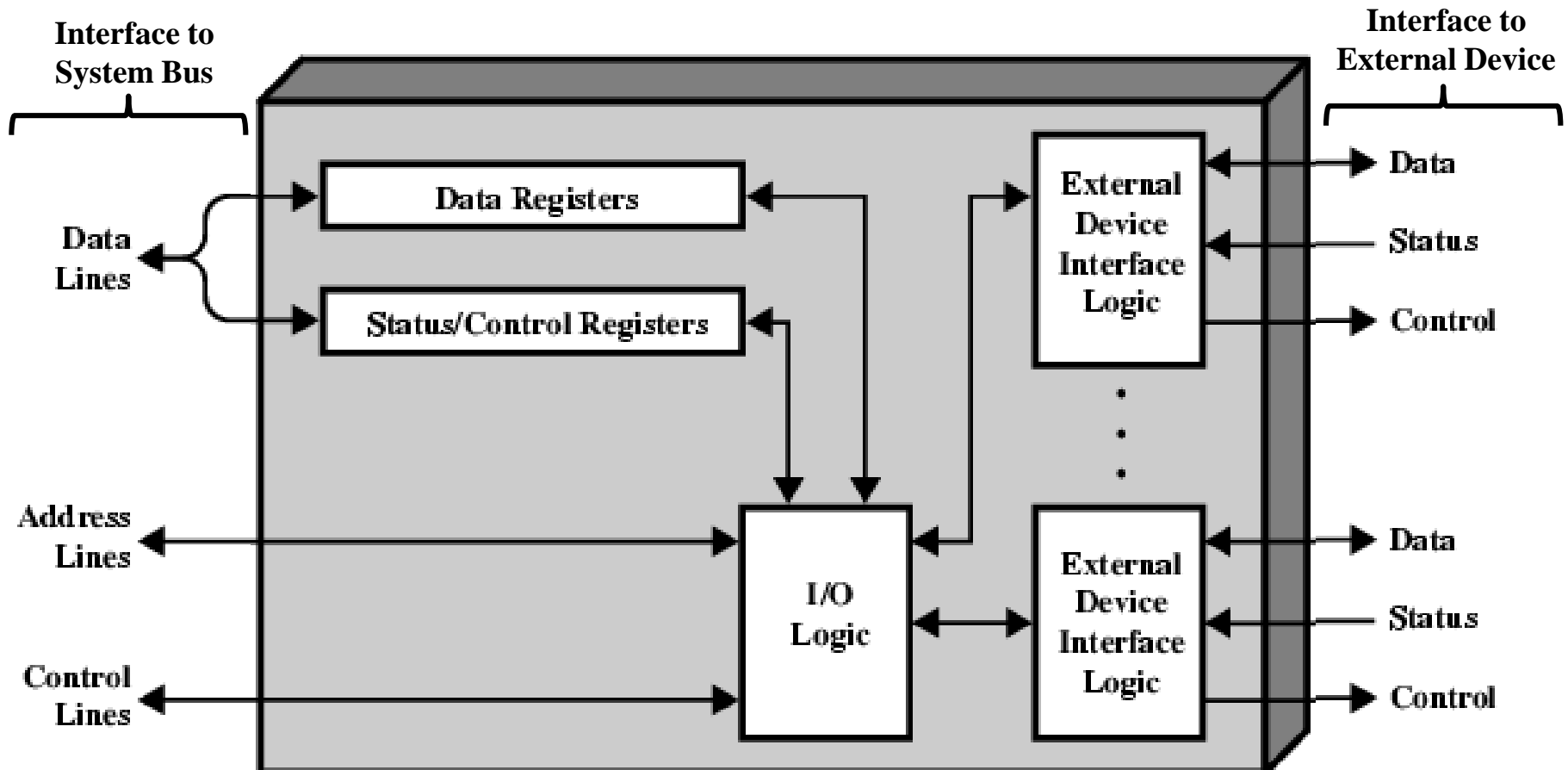


5. Error Detection

- Mechanical and electrical malfunctions
 - Report to CPU.
 - e.g., paper jam, bad disk sector/track.
- Unintentional changes to transmitted bit pattern
 - Detected using error-detecting codes.
 - e.g., parity bit (ASCII).

I/O Module Structure

- St./Ctrl. registers: hold device status or accept control info from CPU.
- CPU issues commands to I/O module via control lines.
- Some control lines are also used by I/O module for bus arbitration.
- Module controlling more than 1 device has a set of unique addresses.



I/O Module Design Decisions

- I/O module lets CPU view a wide range of devices in a simple-minded way.
- Module may **hide device properties** from CPU:
 - Quite **complex module design**.
 - Simple CPU commands** (e.g., render object).
 - Referred to as **I/O channel (I/O processor)**.
 - Common in mainframes.
- Module may **reveal device properties** to CPU:
 - Relatively **simple module design**.
 - Detailed CPU commands** (e.g., rewind tape).
 - Referred to as **I/O controller (device controller)**.
 - Common in microcomputers.

Reading Material

- Stallings, Chapter 6:
 - Pages 205 – 209
 - Pages 215 – 217
- Stallings, Chapter 7:
 - Pages 222 – 228