

Introduction to Constraint Programming

Helmut Simonis

ACP Winterschool 2024



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A version of this material was developed as part of the ECLiPSe ELearning course:

<https://eclipseclp.org/ELearning/index.html>.

Support from Cisco Systems and the Silicon Valley Community Foundation is gratefully acknowledged.

Objectives

- Overview of Core Constraint Programming
- Three Main Concepts
 - Constraint Propagation
 - Global Constraints
 - Customizing Search
- Topics will be treated in more detail in later parts of the school
- Based on Examples, not Formal Description

Outline

- Why Constraint Programming?
- Constraint Propagation
- Global Constraints
- Customizing Search

Tutorial Based on ECLiPSe ELearning Course

- Self-study course in constraint programming
- Supported by Cisco Systems and Silicon Valley Community Foundation
- Multi-media format, video lectures, slides, handout etc
- <https://eclipseclp.org/ELearning/index.html>

Also Part of CRT-AI Constraint Week

- Annual one week course on CP and Optimization in Ireland
- Part of national training program for PhD students in AI
- <https://www.crt-ai.ie/>

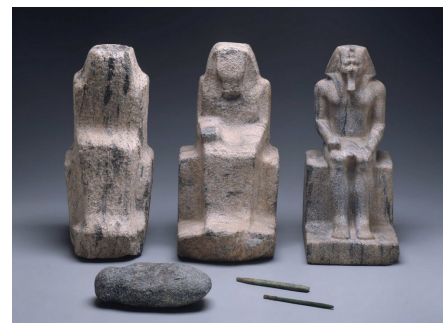
Constraint Programming - in a nutshell

- Declarative description of problems with
 - *Variables* which range over (finite) sets of values
 - *Constraints* over subsets of variables which restrict possible value combinations
 - A *solution* is a value assignment which satisfies all constraints
- Constraint propagation/reasoning
 - Removing inconsistent values for variables
 - Detect failure if constraint can not be satisfied
 - Interaction of constraints via shared variables
 - Incomplete
- Search
 - User controlled assignment of values to variables
 - Each step triggers constraint propagation
- Different domains require/allow different methods

Constraint Programming is Different

- Declarative Programming
 - Concentrate on what you want
 - Not how to get there
 - Program \neq Algorithm
 - Program = Model
- Applied to Combinatorial Problems
 - No complete polynomial algorithms known (exist?)
 - CP less ad-hoc than heuristics
 - Models can evolve

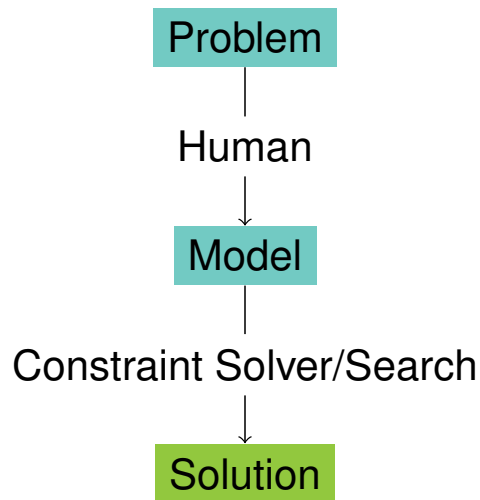
A Subtractive Process



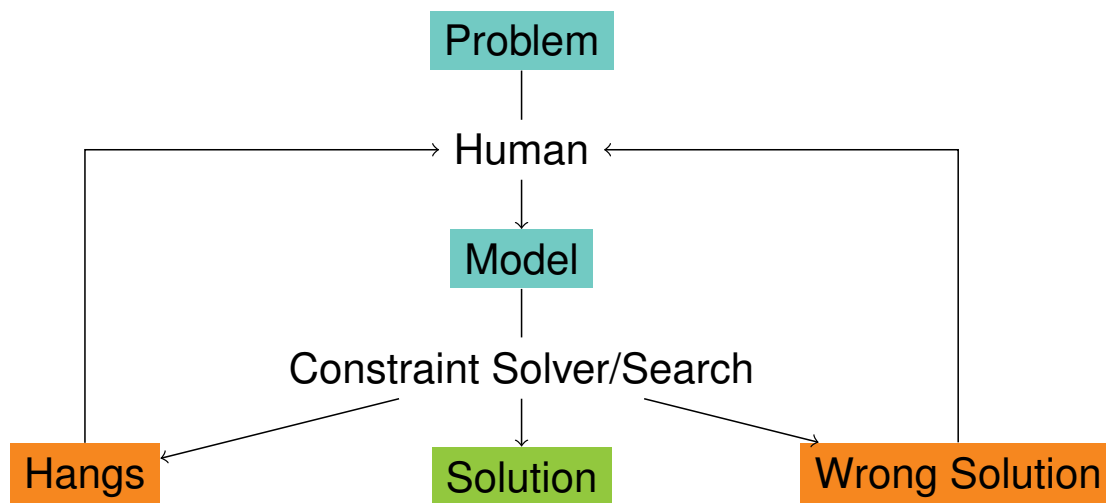
“Oh, bosh, as Mr. Ruskin says. Sculpture, per se, is the simplest thing in the world. All you have to do is to take a big chunk of marble and a hammer and chisel, make up your mind what you are about to create and chip off all the marble you don’t want.”-Paris Gaulois.

Source: <https://quoteinvestigator.com/2014/06/22/chip-away/>

Basic Process



More Realistic



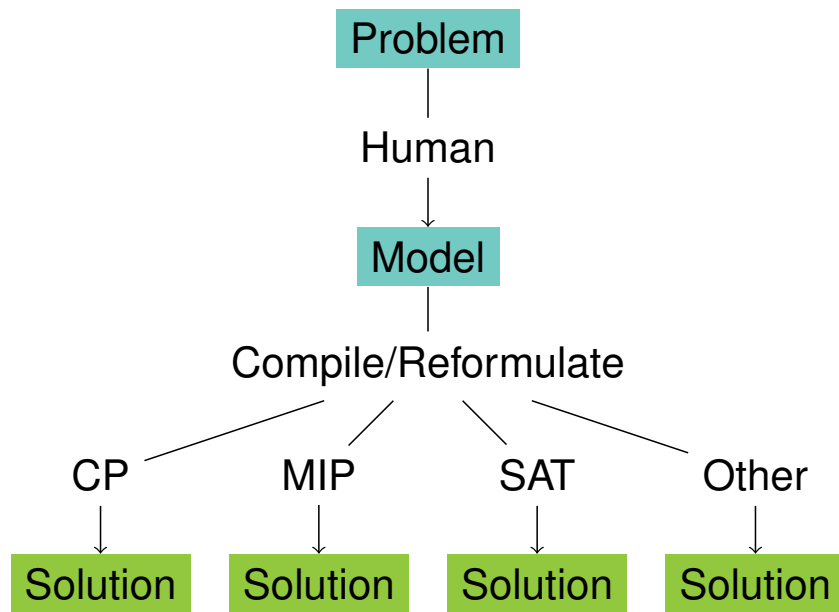
Dual Role of Model

- Allows Human to Express Problem
 - Close to Problem Domain
 - Constraints as Abstractions
- Allows Solver to Execute
 - Variables as Communication Mechanism
 - Constraints as Algorithms

Modelling Frameworks

- MiniZinc (NICTA, Monash University, Australia)
- NumberJack (Insight, Ireland)
- EssencePrime/SavilleRow (UK)
- CPMpy (KU Leuven)
- Allow use of multiple back-end solvers
- Compile model into variants for each solver
- A priori solver independent model(CP, MIP, SAT)

Framework Process



Why use Puzzles as Examples?

- Easy to understand the problem
- Solvable by hand without specialized knowledge
- Possible to compare automated to manual solving process

The puzzle, though inanimate, is presented as a solvable problem without lasting negative consequences, a very low-risk low-reward situation. By being a puzzle, the object is attempting to convince the user that it must be completed.

Source: Every Day Rhetoric

Part I

Basic Constraint Propagation

Example 1: SEND+MORE=MONEY

- Example of Finite Domain Constraint Problem
- Models and Programs
- Constraint Propagation and Search
- Some Basic Constraints: linear arithmetic, alldifferent, disequality
- A Built-in search
- Visualizers for variables, constraints and search

Problem Definition

A Crypt-Arithmetic Puzzle

We begin with the definition of the SEND+MORE=MONEY puzzle. It is often shown in the form of a hand-written addition:

$$\begin{array}{r} \text{S} \text{ E} \text{ N} \text{ D} \\ + \text{M} \text{ O} \text{ R} \text{ E} \\ \hline \text{M} \text{ O} \text{ N} \text{ E} \text{ Y} \end{array}$$

The puzzle was first proposed by Henry Dudeney in the Strand Magazine from 1924.

Rules

- Each character stands for a digit from 0 to 9.
- Numbers are built from digits in the usual, positional notation.
- Repeated occurrence of the same character denote the same digit.
- Different characters denote different digits.
- Numbers do not start with a zero.
- The equation must hold.

$$\begin{array}{r} \text{S} \text{ E} \text{ N} \text{ D} \\ + \text{M} \text{ O} \text{ R} \text{ E} \\ \hline \text{M} \text{ O} \text{ N} \text{ E} \text{ Y} \end{array}$$

Model

- Each character is a variable, which ranges over the values 0 to 9.
- An *alldifferent* constraint between all variables, which states that two different variables must have different values. This is a very common constraint, which we will encounter in many other problems later on.
- Two *disequality constraints* (variable X must be different from value V) stating that the variables at the beginning of a number can not take the value 0.
- An arithmetic *equality constraint* linking all variables with the proper coefficients and stating that the equation must hold.

SEND+MORE=MONEY Models

- ECLiPSe [▶ Show](#)
- MiniZinc [▶ Show](#)
- NumberJack [▶ Show](#)
- CPMpy [▶ Show](#)
- Choco-solver [▶ Show](#)

ECLiPSe Model

```
:- lib(ic).

sendmore(Digits) :-
    Digits = [S,E,N,D,M,O,R,Y],
    Digits :: [0..9],
    alldifferent(Digits),
    S #\= 0,
    M #\= 0,
    1000*S + 100*E + 10*N + D
    + 1000*M + 100*O + 10*R + E
    #= 10000*M + 1000*O + 100*N + 10*E + Y,
    labeling(Digits).
```

► Continue

MiniZinc Model

```
include "alldifferent.mzn";
var 0..9: S;
var 0..9: E;
var 0..9: N;
var 0..9: D;
var 0..9: M;
var 0..9: O;
var 0..9: R;
var 0..9: Y;
constraint S != 0;
constraint M != 0;
constraint
    1000 * S + 100 * E + 10 * N + D
    + 1000 * M + 100 * O + 10 * R + E
    = 10000 * M + 1000 * O + 100 * N + 10 * E + Y;
constraint alldifferent([S,E,N,D,M,O,R,Y]);

solve satisfy;
```

► Continue

NumberJack Model (from <https://github.com/eomahony/Numberjack/>)

```
from Numberjack import *

def get_model():
    model = Model()
    s, m = VarArray(2, 1, 9)
    e, n, d, o, r, y = VarArray(6, 0, 9)
    model.add(
        s*1000 + e*100 + n*10 + d +
        m*1000 + o*100 + r*10 + e ==
        m*10000 + o*1000 + n*100 + e*10 + y)
    model.add(AllDiff((s, e, n, d, m, o, r, y)))
    return s, e, n, d, m, o, r, y, model

def solve(param):
    s, e, n, d, m, o, r, y, model = get_model()
    solver = model.load(param['solver'])
    solver.setVerbosity(param['verbose'])
    solver.solve()
```

► Continue

CPMpy Model (from <https://github.com/CPMpy/>)

```
from cpmPy import *
import numpy as np

s,e,n,d,m,o,r,y = intvar(0,9, shape=8)
model = Model(
    AllDifferent([s,e,n,d,m,o,r,y]),
    (
        sum([s,e,n,d] * np.array([1000, 100, 10, 1])) \
        + sum([m,o,r,e] * np.array([1000, 100, 10, 1])) \
        == sum([m,o,n,e,y] * np.array([10000, 1000, 100, 10, 1])) ),
    s > 0,
    m > 0,
)

model.solve()
```

► Continue

```
Model model = new Model("SEND+MORE=MONEY");
IntVar S = model.intVar("S", 1, 9, false);
IntVar E = model.intVar("E", 0, 9, false);
IntVar N = model.intVar("N", 0, 9, false);
IntVar D = model.intVar("D", 0, 9, false);
IntVar M = model.intVar("M", 1, 9, false);
IntVar O = model.intVar("O", 0, 9, false);
IntVar R = model.intVar("R", 0, 9, false);
IntVar Y = model.intVar("Y", 0, 9, false);

model.allDifferent(new IntVar[]{S, E, N, D, M, O, R, Y}).post();

IntVar[] ALL = new IntVar[]{
    S, E, N, D,
    M, O, R, E,
    M, O, N, E, Y};
int[] COEFFS = new int[]{
    1000, 100, 10, 1,
    1000, 100, 10, 1,
    -10000, -1000, -100, -10, -1};
model.scalar(ALL, COEFFS, "=", 0).post();

Solver solver = model.getSolver();
solver.showStatistics();
solver.showSolutions();
solver.findSolution();
```

► Continue

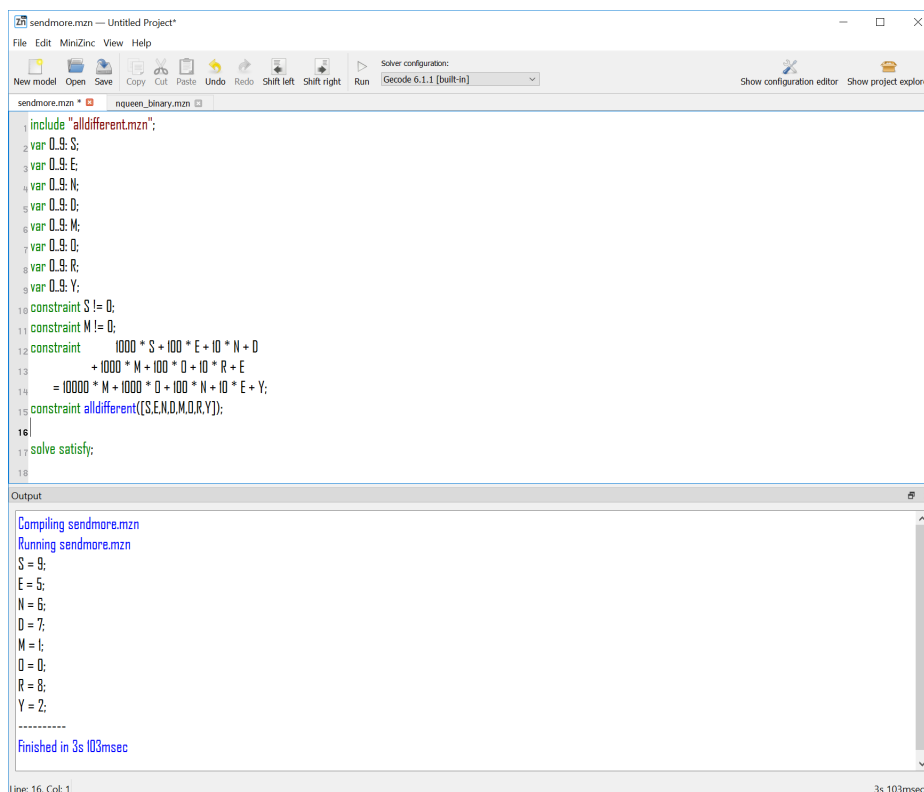
A Note on Syntax

- Some formulations may seem simpler than others
- This largely is an artifact of a very simple problem
- In most models, you do not write down constraints one by one
- You create constraints based on data
- Ease of integration becomes more important than syntax
- Debugging tools for those who need a debugger :-)

Choice of Model

- This is *one* model, not *the* model of the problem
- Many possible alternatives
- Choice often depends on your constraint system
 - Constraints available
 - Reasoning attached to constraints
- Not always clear which is the *best* model
- Often: Not clear what is the *problem*

Running the Program (MiniZinc IDE)



The screenshot shows the MiniZinc IDE interface. The main editor displays the following code:

```
1 include "alldifferent.mzn";
2 var 0..9: S;
3 var 0..9: E;
4 var 0..9: N;
5 var 0..9: D;
6 var 0..9: M;
7 var 0..9: O;
8 var 0..9: R;
9 var 0..9: Y;
10 constraint S != 0;
11 constraint M != 0;
12 constraint 1000 * S + 100 * E + 10 * N + D
13           + 1000 * M + 100 * O + 10 * R + E
14           = 10000 * M + 1000 * O + 100 * N + 10 * E + Y;
15 constraint alldifferent((S,E,N,D,M,O,R,Y));
16
17 solve satisfy;
```

The output window shows the following text:

```
Compiling sendmore.mzn
Running sendmore.mzn
S = 9;
E = 5;
N = 6;
D = 7;
M = 1;
O = 0;
R = 8;
Y = 2;
-----
Finished in 3s 103msec
```

The status bar at the bottom indicates "Line: 16, Col: 1" and "3s 103msec".

Question

- But how did the program come up with this solution?
- We show solution with ECLiPse, other solvers vary slightly

Domain Definition

```
var 0..9: S;  
var 0..9: E;  
var 0..9: N;  
var 0..9: D;  
var 0..9: M;  
var 0..9: O;  
var 0..9: R;  
var 0..9: Y;
```


Domain Visualization

Columns = Values

Rows = Variables

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M			Cells= State							
O										
R										
Y										

Alldifferent Constraint

```
include "alldifferent.mzn";
```

```
constraint alldifferent([S,E,N,D,M,O,R,Y]);
```

- Built-in alldifferent predicate included
- No initial propagation possible
- *Suspends*, waits until variables are changed
- When variable is fixed, remove value from domain of other variables
- *Forward checking*

All different Visualization

Uses the same representation as the domain visualizer

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

Disequality Constraints

```
constraint S != 0;  
constraint M != 0;
```

Remove value from domain

$$S \in \{1..9\}, M \in \{1..9\}$$

Constraints solved, can be removed

Domains after Disequality

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

Equality Constraint

- Normalization of linear terms
 - Single occurrence of variable
 - Positive coefficients
- Propagation

Normalization

$$\begin{array}{rcccccc} & 1000*S+ & 100*E+ & 10*N+ & D & \\ & +1000*M+ & 100*O+ & 10*R+ & E & \\ \hline 10000*M+ & 1000*O+ & 100*N+ & 10*E+ & Y & \\ \text{is transformed into} & & & & & \\ & 1000*S+ & 91*E+ & & D & \\ & & + 10*R & & & \\ \hline 9000*M+ & 900*O+ & 90*N+ & & Y & \end{array}$$

Simplified Equation

$$1000 * S + 91 * E + 10 * R + D = 9000 * M + 900 * O + 90 * N + Y$$

Propagation

$$\underbrace{1000 * S^{1..9} + 91 * E^{0..9} + 10 * R^{0..9} + D^{0..9}}_{1000..9918} = \underbrace{9000 * M^{1..9} + 900 * O^{0..9} + 90 * N^{0..9} + Y^{0..9}}_{9000..89919}$$

Deduction:

$$M = 1, S = 9, O \in \{0..1\}$$

Why? [Skip](#)

Consider lower bound for S

$$\underbrace{1000 * S^{1..9} + 91 * E^{0..9} + 10 * R^{0..9} + D^{0..9}}_{9000..9918} = \underbrace{9000 * M^{1..9} + 900 * O^{0..9} + 90 * N^{0..9} + Y^{0..9}}_{9000..9918}$$

- Lower bound of equation is 9000
- Rest of lhs (left hand side) ($91 * E^{0..9} + 10 * R^{0..9} + D^{0..9}$) is atmost 918
- S must be greater or equal to $\frac{9000-918}{1000} = 8.082$
 - otherwise lower bound of equation not reached by lhs
- S is integer, therefore $S \geq \lceil \frac{9000-918}{1000} \rceil = 9$
- S has upper bound of 9, so $S = 9$

Consider upper bound of M

$$\underbrace{1000 * S^{1..9} + 91 * E^{0..9} + 10 * R^{0..9} + D^{0..9}}_{9000..9918} = \underbrace{9000 * M^{1..9} + 900 * O^{0..9} + 90 * N^{0..9} + Y^{0..9}}_{9000..9918}$$

- Upper bound of equation is 9918
- Rest of rhs (right hand side) $900 * O^{0..9} + 90 * N^{0..9} + Y^{0..9}$ is at least 0
- M must be smaller or equal to $\frac{9918-0}{9000} = 1.102$
- M must be integer, therefore $M \leq \lfloor \frac{9918-0}{9000} \rfloor = 1$
- M has lower bound of 1, so $M = 1$

Consider upper bound of O

$$\underbrace{1000 * S^{1..9} + 91 * E^{0..9} + 10 * R^{0..9} + D^{0..9}}_{9000..9918} = \underbrace{9000 * M^{1..9} + 900 * O^{0..9} + 90 * N^{0..9} + Y^{0..9}}_{9000..9918}$$

- Upper bound of equation is 9918
- Rest of rhs (right hand side) $9000 * 1 + 90 * N^{0..9} + Y^{0..9}$ is at least 9000
- O must be smaller or equal to $\frac{9918-9000}{900} = 1.02$
- O must be integer, therefore $O \leq \lfloor \frac{9918-9000}{900} \rfloor = 1$
- O has lower bound of 0, so $O \in \{0..1\}$

Propagation of equality: Result

	0	1	2	3	4	5	6	7	8	9
S		-	-	-	-	-	-	-	-	☀
E										
N										
D										
M		☀	-	-	-	-	-	-	-	-
O			✕	✕	✕	✕	✕	✕	✕	✕
R										
Y										

Propagation of alldifferent

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

$O = 0, [E, R, D, N, Y] \in \{2..8\}$

Waking the equality constraint

- Triggered by assignment of variables
- *or* update of lower or upper bound

Removal of constants

$$1000 * 9 + 91 * E^{2..8} + 10 * R^{2..8} + D^{2..8} = \\ 9000 * 1 + 900 * 0 + 90 * N^{2..8} + Y^{2..8}$$

$$\mathbf{1000 * 9 + 91 * E^{2..8} + 10 * R^{2..8} + D^{2..8} =} \\ \mathbf{9000 * 1 + 900 * 0 + 90 * N^{2..8} + Y^{2..8}}$$

$$91 * E^{2..8} + 10 * R^{2..8} + D^{2..8} = 90 * N^{2..8} + Y^{2..8}$$

Propagation of equality (Iteration 1)

$$\underbrace{91 * E^{2..8} + 10 * R^{2..8} + D^{2..8}}_{204..816} = \underbrace{90 * N^{2..8} + Y^{2..8}}_{182..728}$$

$$\underbrace{91 * E^{2..8} + 10 * R^{2..8} + D^{2..8}}_{204..728} = 90 * N^{2..8} + Y^{2..8}$$

$$N \geq 3 = \lceil \frac{204 - 8}{90} \rceil, E \leq 7 = \lfloor \frac{728 - 22}{91} \rfloor$$

Propagation of equality (Iteration 2)

$$91 * E^{2..7} + 10 * R^{2..8} + D^{2..8} = 90 * N^{3..8} + Y^{2..8}$$

$$\underbrace{91 * E^{2..7} + 10 * R^{2..8} + D^{2..8}}_{204..725} = \underbrace{90 * N^{3..8} + Y^{2..8}}_{272..728}$$

$$\underbrace{91 * E^{2..7} + 10 * R^{2..8} + D^{2..8}}_{272..725} = 90 * N^{3..8} + Y^{2..8}$$

$$E \geq 3 = \lceil \frac{272 - 88}{91} \rceil$$

Propagation of equality (Iteration 3)

$$91 * E^{3..7} + 10 * R^{2..8} + D^{2..8} = 90 * N^{3..8} + Y^{2..8}$$

$$\underbrace{91 * E^{3..7} + 10 * R^{2..8} + D^{2..8}}_{295..725} = \underbrace{90 * N^{3..8} + Y^{2..8}}_{272..728}$$

$$\underbrace{91 * E^{3..7} + 10 * R^{2..8} + D^{2..8}}_{295..725} = 90 * N^{3..8} + Y^{2..8}$$

$$N \geq 4 = \lceil \frac{295 - 8}{90} \rceil$$

Propagation of equality (Iteration 4)

$$91 * E^{3..7} + 10 * R^{2..8} + D^{2..8} = 90 * N^{4..8} + Y^{2..8}$$

$$\underbrace{91 * E^{3..7} + 10 * R^{2..8} + D^{2..8}}_{295..725} = \underbrace{90 * N^{4..8} + Y^{2..8}}_{362..728}$$

$$\underbrace{91 * E^{3..7} + 10 * R^{2..8} + D^{2..8}}_{362..725} = 90 * N^{4..8} + Y^{2..8}$$

$$E \geq 4 = \lceil \frac{362 - 88}{91} \rceil$$

Propagation of equality (Iteration 5)

$$91 * E^{4..7} + 10 * R^{2..8} + D^{2..8} = 90 * N^{4..8} + Y^{2..8}$$

$$\underbrace{91 * E^{4..7} + 10 * R^{2..8} + D^{2..8}}_{386..725} = \underbrace{90 * N^{4..8} + Y^{2..8}}_{362..728}$$

$$\underbrace{91 * E^{4..7} + 10 * R^{2..8} + D^{2..8}}_{386..725} = 90 * N^{4..8} + Y^{2..8}$$

$$N \geq 5 = \lceil \frac{386 - 8}{90} \rceil$$

Propagation of equality (Iteration 6)

$$91 * E^{4..7} + 10 * R^{2..8} + D^{2..8} = 90 * N^{5..8} + Y^{2..8}$$

$$\underbrace{91 * E^{4..7} + 10 * R^{2..8} + D^{2..8}}_{386..725} = \underbrace{90 * N^{5..8} + Y^{2..8}}_{452..728}$$

$$\underbrace{91 * E^{4..7} + 10 * R^{2..8} + D^{2..8}}_{452..725} = 90 * N^{5..8} + Y^{2..8}$$

$$N \geq 5 = \lceil \frac{452 - 8}{90} \rceil, E \geq 4 = \lceil \frac{452 - 88}{91} \rceil$$

No further propagation at this point

Domains after setup

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

Search

solve satisfy;

- Try to find a feasible solution, choice left to solver
- Naive search strategy shown here
 - Try variable in order given
 - Try values starting from smallest value in domain
 - When failing, backtrack to last open choice
 - *Chronological Backtracking*
 - *Depth First search*

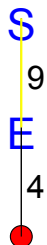
Search Tree Step 1



Variable S already fixed

Step 2, Alternative $E = 4$

Variable $E \in \{4..7\}$, first value tested is 4



Assignment $E = 4$

	0	1	2	3	4	5	6	7	8	9
S										
E					☀	-	-	-		
N										
D										
M										
O										
R										
Y										

Propagation of $E = 4$, equality constraint

$$91 * 4 + 10 * R^{2..8} + D^{2..8} = 90 * N^{5..8} + Y^{2..8}$$

$$\underbrace{91 * 4 + 10 * R^{2..8} + D^{2..8}}_{386..452} = \underbrace{90 * N^{5..8} + Y^{2..8}}_{452..728}$$

$$\underbrace{91 * 4 + 10 * R^{2..8} + D^{2..8}}_{452} = 90 * N^{5..8} + Y^{2..8}$$

$$N = 5, Y = 2, R = 8, D = 8$$

Result of equality propagation

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D			-	-	-	-	-	-		
M										
O										
R			-	-	-	-	-	-		
Y				-	-	-	-	-		

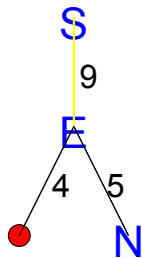
Propagation of alldifferent

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D			-	-	-	-	-	-		
M										
O										
R			-	-	-	-	-	-		
Y				-	-	-	-	-		

Alldifferent fails!

Step 2, Alternative $E = 5$

Return to last open choice, E , and test next value



Assignment $E = 5$

	0	1	2	3	4	5	6	7	8	9
S										
E					-	☀	-	-		
N										
D										
M										
O										
R										
Y										

Propagation of alldifferent

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

$$N \neq 5, N \geq 6$$

Propagation of equality

$$91 * 5 + 10 * R^{2..8} + D^{2..8} = 90 * N^{6..8} + Y^{2..8}$$

$$\underbrace{91 * 5 + 10 * R^{2..8} + D^{2..8}}_{477..543} = \underbrace{90 * N^{6..8} + Y^{2..8}}_{542..728}$$

$$\underbrace{91 * 5 + 10 * R^{2..8} + D^{2..8}}_{542..543} = 90 * N^{6..8} + Y^{2..8}$$

$$N = 6, Y \in \{2, 3\}, R = 8, D \in \{7..8\}$$

Result of equality propagation

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

Propagation of alldifferent

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

$D = 7$

Propagation of equality

$$91 * 5 + 10 * 8 + 7 = 90 * 6 + Y^{2..3}$$

$$\underbrace{91 * 5 + 10 * 8 + 7}_{542} = \underbrace{90 * 6 + Y^{2..3}}_{542..543}$$

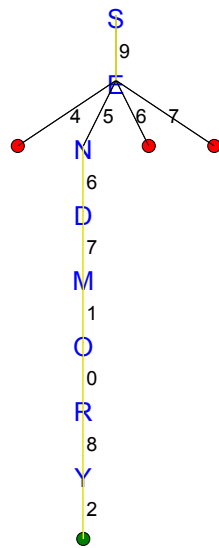
$$\underbrace{91 * 5 + 10 * 8 + 7}_{542} = 90 * 6 + Y^{2..3}$$

$$Y = 2$$

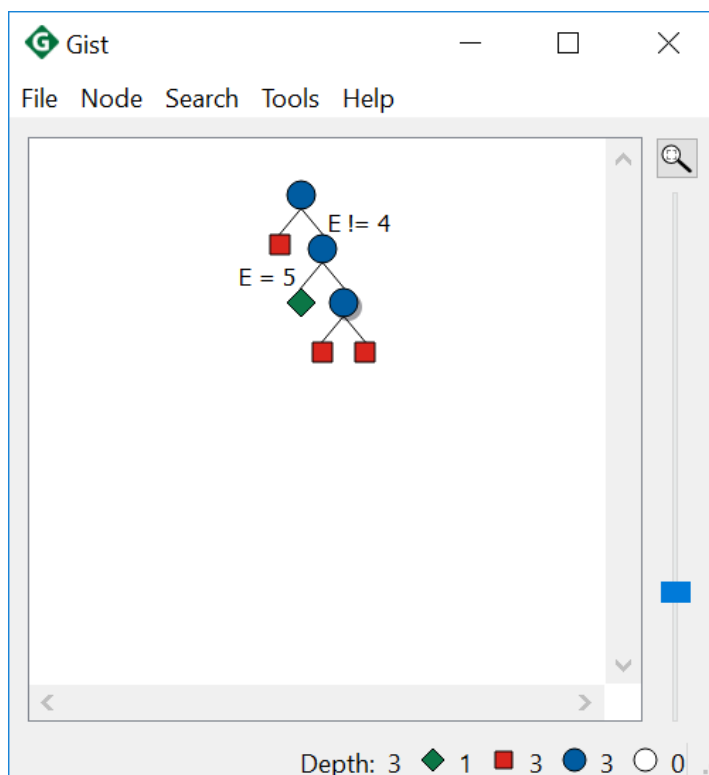
Last propagation step

	0	1	2	3	4	5	6	7	8	9
S										
E										
N										
D										
M										
O										
R										
Y										

Complete Search Tree



Search Tree with Gecode/GIST



Some Differences

- Uses binary branching
 - var equal value, var not equal value
- Solutions in green, failure leafs in red, internal nodes in blue
- By default, shows all failed sub trees collapsed
- By default, uses different search strategy

Solution

$$\begin{array}{r} 9 \ 5 \ 6 \ 7 \\ + \ 1 \ 0 \ 8 \ 5 \\ \hline 1 \ 0 \ 6 \ 5 \ 2 \end{array}$$

Points to Remember

- Constraint models are expressed by variables and constraints.
- Problems can have many different models, which can behave quite differently. Choosing the best model is an art.
- Constraints can take many different forms.
- Propagation deals with the interaction of variables and constraints.
- It removes some values that are inconsistent with a constraint from the domain of a variable.
- Constraints only communicate via shared variables.

Points to Remember

- Propagation usually is not sufficient, search may be required to find a solution.
- Propagation is data driven, and can be quite complex even for small examples.
- The default search uses chronological depth-first backtracking, systematically exploring the complete search space.
- The search choices and propagation are interleaved, after every choice some more propagation may further reduce the problem.

For Puzzle Purists Only

- We did not follow the puzzler ethos!
- We should solve the puzzle without making choices
- Even a case analysis should be avoided
- The puzzle has a single solution, we should be able to deduce the solution
- In the process shown, we are limited by the underlying assumptions
 - Treat each constraint on its own, they interact only by domains of variables
 - We only use the constraints that we stated in the model
- Can we do better than that?

► Skip

Better Reasoning

- Three possible approaches (possibly many more)
 - Full domain reasoning for arithmetic, not just bound reasoning
 - Interaction of *sum* and *alldifferent* constraints
 - Deduced implied constraint

Looking at more than just bounds

- We only considered the smallest and largest values that can be achieved in the sum constraint
- We can do more
 - Can any of the values between be expressed as the sum of the terms
 - Consider holes in the domains, and in the range of possible values for LHS and RHS
- Usually not done in actual solvers for arithmetic constraints
- Easy to do with Dynamic Programming

Consider the interaction of multiple constraints

- Usually ignored, as only interaction is via domains of shared variables
- Here: Sum and *alldifferent* interact
 - When considering the bounds, we cannot assume that each variable takes its smallest/largest value independently
 - Find feasible assignment that minimizes/maximizes the total weight
 - To do this properly, we need some non-trivial reasoning
- Do we do this combined reasoning automatically, or only when prompted by the modeler?

Deduced Implied Constraints

- Look at the partially solved puzzle

$$\begin{array}{r} 9\text{END} \\ +10\text{RE} \\ \hline \end{array}$$

- 10NEY
- In the hundreds position, we have
 $E + 0 + C_{10} = N + 10 * C_{100}$, with C_{10} the 0/1 carry from the tens position
- NB: No carry C_{100} into the thousands, $C_{100} = 0$
- N must be equal to $E + 1$ with $C_{10} = 1$
- If $C_{10} = 0$, then $N = E$, not possible
- We can substitute $N = E + 1$ into our main equation, but keep $N = E + 1$ as well

Expert Mode Reasoning

Starting with

$$91 * E^{4..7} + 10 * R^{2..8} + D^{2..8} = 90 * N^{5..8} + Y^{2..8}$$

we get

$$91 * E^{4..7} + 10 * R^{2..8} + D^{2..8} = 90 * E^{4..7} + 90 + Y^{2..8}$$

Eliminating duplicate occurrences of E

$$\underbrace{E^{4..7} + 10 * R^{2..8} + D^{2..8}}_{26..95} = \underbrace{90 + Y^{2..8}}_{92..98}$$

shared range 92..95 To reach 92, R must be equal to 8, therefore N, E, D, Y must be less than 8

As $N = E + 1$, E must be less than 7

$$E^{4..6} + 10 * 8 + D^{2..7} = 90 + Y^{2..7}$$

Simplification yields

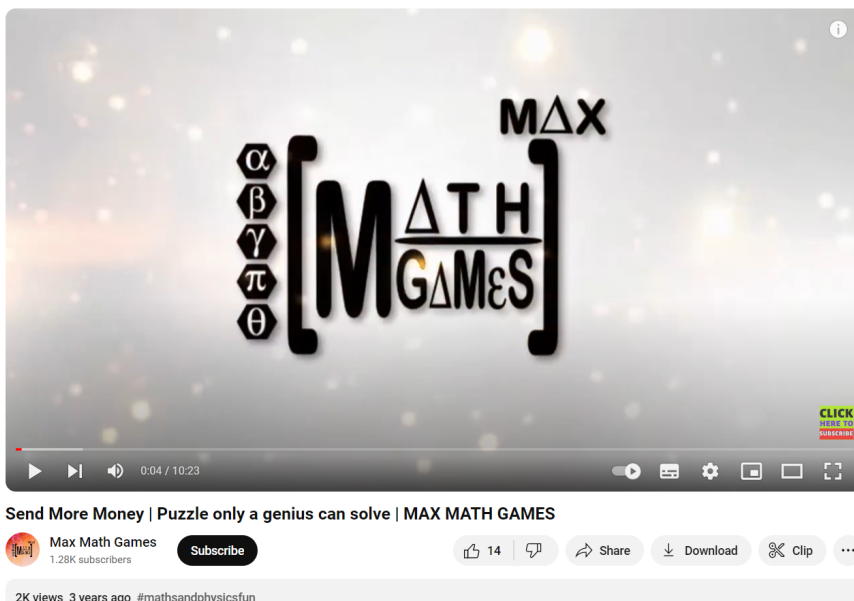
$$\underbrace{E^{4..6} + D^{2..7}}_{6..13} = \underbrace{10 + Y^{2..7}}_{12..17}$$

shared range 12..13 To reach 12 on LHS, E must be greater than 4, R must be less than 5

Expert Mode Summary

- Often there is more propagation that can be done
- Can be difficult/expensive to do
- Balancing
 - How much work it done at each step of search?
 - How many steps of search you need?
- For hard problems, doing all possible propagation may be exponential
- Not aware that any CP system does the full reasoning shown here

This is how people solve the puzzle by hand



- When writing the first version of this puzzle for CHIP (in 1986), we wanted to mimic the way we solve the puzzle by hand

Part II

Global Constraints

Example 2: Sudoku

- Global Constraints
 - Powerful modelling abstractions
 - Non-trivial propagation
 - Different consistency levels
- Example: Sudoku puzzle

Problem Definition

Sudoku

Fill in numbers from 1 to 9 so that each row, column and 3x3 block contain each number exactly once

4			8					

4	2	8	5	6	3	1	7	9
3	5	9	1	7	2	4	6	8
7	6	1	4	8	9	5	3	2
1	4	6	3	9	8	2	5	7
5	9	2	7	4	1	3	8	6
8	3	7	6	2	5	9	4	1
2	7	4	9	5	6	8	1	3
6	8	3	2	1	4	7	9	5
9	1	5	8	3	7	6	2	4

Model

- A variable for each cell, ranging from 1 to 9
- A 9x9 matrix of variables describing the problem
- Preassigned integers for the given hints
- `alldifferent` constraints for each row, column and 3x3 block

- ECLiPSe [▶ Show](#)
- MiniZinc [▶ Show](#)
- NumberJack [▶ Show](#)
- CPMpy [▶ Show](#)
- Choco-solver [▶ Show](#)

ECLiPSe Sudoku Model (from <https://eclipseclp.org/>)

```
:- lib(ic).
:- import alldifferent/1 from ic_global.

top :-
    problem(Board),
    print_board(Board),
    sudoku(Board),
    labeling(Board),
    print_board(Board).

sudoku(Board) :-
    dim(Board, [N,N]),
    Board :: 1..N,
    ( for(I,1,N), param(Board) do
        alldifferent(Board[I,*]),
        alldifferent(Board[*,I])
    ),
    NN is integer(sqrt(N)),
    ( multifor([I,J],1,N,NN), param(Board,NN) do
        alldifferent(concat(Board[I..I+NN-1, J..J+NN-1]))
    ).

print_board(Board) :-
    dim(Board, [N,N]),
    ( for(I,1,N), param(Board,N) do
        ( for(J,1,N), param(Board,I) do
            X is Board[I,J],
            ( var(X) -> write(" _") ; printf(" %2d", [X]) )
        ), nl
    ), nl.
```

ECLiPSe Data Definition

```
problem([ (
  [] (4, _, 8, _, _, _, _, _),
  [] (_, _, _, 1, 7, _, _, _),
  [] (_, _, _, 8, _, _, 3, 2),
  [] (_, _, 6, _, _, 8, 2, 5, _),
  [] (_, 9, _, _, _, _, 8, _),
  [] (_, 3, 7, 6, _, _, 9, _, _),
  [] (2, 7, _, _, 5, _, _, _),
  [] (_, _, _, 1, 4, _, _, _),
  [] (_, _, _, _, _, 6, _, 4) ) ).
```

► Continue

MiniZinc Sudoku Model

```
int: s;
int: n=s*s;
array[1..n,1..n] of var 1..n: puzzle;

include "sudoku.dzn";

predicate alldifferent(array[int] of var int: x) =
  forall(i,j in index_set(x) where i < j)
    (x[i] != x[j]);

constraint forall(i in 1..n)
  (alldifferent([puzzle[i,j] | j in 1..n]));
constraint forall(j in 1..n)
  (alldifferent([ puzzle[i,j] | i in 1..n]));
constraint forall(i,j in 1..s)
  (alldifferent([puzzle[s*(i-1)+p, s*(j-1)+q] |
    p,q in 1..s]));
solve satisfy;
```

MiniZinc Output

```
output [ "sudoku:\\n" ] ++
[ show(puzzle[i,j]) ++
  if j = n then
    if i mod s = 0 /\ i < n then "\\n\\n"
    else "\\n"
    endif
  else
    if j mod s = 0 then " "
    else " "
    endif
  endif
| i,j in 1..n ];
```

MiniZinc Data File (sudoku.dzn)

```
s=3;
puzzle=[|
  4, _, 8, _, _, _, _, _|
  _, _, _, 1, 7, _, _, _|
  _, _, _, _, 8, _, 3, 2|
  _, _, 6, _, _, 8, 2, 5, _|
  _, 9, _, _, _, _, 8, _|
  _, 3, 7, 6, _, _, 9, _|
  2, 7, _, _, 5, _, _, _|
  _, _, _, _, 1, 4, _, _|
  _, _, _, _, _, 6, _, 4|
|];
```

► Continue

NumberJack Sudoku Model

```
from Numberjack import *

def get_model(N, clues):
    grid = Matrix(N*N, N*N, 1, N*N)

    sudoku = Model([AllDiff(row) for row in grid.row],
                   [AllDiff(col) for col in grid.col],
                   [AllDiff(grid[x:x + N, y:y + N]) for x in range(0, N*N, N)
                                                            for y in range(0, N * N, N)],
                   [(x == int(v)) for x, v in
                    zip(grid.flat, "".join(open(clues)).split()) if v != '*']
                  ])
    return grid, sudoku

def solve(param):
    N = param['N']
    clues = param['file']

    grid, sudoku = get_model(N, clues)

    solver = sudoku.load(param['solver'])
    solver.setVerbosity(param['verbose'])
    solver.setTimeLimit(param['tcutoff'])

    solver.solve()
```

NumberJack Data File

```
4 * 8 * * * * *
* * * 1 7 * * *
* * * * 8 * * 3 2
* * 6 * * 8 2 5 *
* 9 * * * * 8 *
* 3 7 6 * * 9 *
2 7 * * 5 * *
* * * * 1 4 *
* * * * * 6 * 4
```

► Continue

CPMpy Sudoku Model (from <https://github.com/CPMpy/>)

```
import numpy as np
from cpmPy import *

# Variables
puzzle = intvar(1,9, shape=given.shape, name="puzzle")

model = Model(
    # Constraints on values (cells that are not empty)
    puzzle[given!=e] == given[given!=e], # numpy's indexing, vectorized equality
    # Constraints on rows and columns
    [AllDifferent(row) for row in puzzle],
    [AllDifferent(col) for col in puzzle.T], # numpy's Transpose
)

# Constraints on blocks
for i in range(0,9, 3):
    for j in range(0,9, 3):
        model += AllDifferent(puzzle[i:i+3, j:j+3]) # python's indexing

model.solve()
```

CPMpy Data Definition

```
e = 0 # value for empty cells
given = np.array([
    [4, e, 8, e, e, e, e, e, e],
    [e, e, e, 1, 7, e, e, e, e],
    [e, e, e, e, 8, e, e, 3, 2],
    [e, e, 6, e, e, 8, 2, 5, e],
    [e, 9, e, e, e, e, e, 8, e],
    [e, 3, 7, 6, e, e, 9, e, e],
    [2, 7, e, e, 5, e, e, e, e],
    [e, e, e, e, 1, 4, e, e, e],
    [e, e, e, e, e, e, 6, e, 4]
])
```

► Continue

Choco-solver Sudoku Model

```
Model model = new Model("Sudoku");
int blockSize = 3;
int m = blockSize*blockSize;

IntVar[][] vars = new IntVar[m][m];
for(int i=0;i<m;i++){
    for(int j=0;j<m;j++){
        vars[i][j] = model.intVar("X"+i+" "+j, 1, m);
        if (data[i][j]>0) {
            model.arithm(vars[i][j],"=",data[i][j]).post();
        }
    }
}
for(int i=0;i<m;i++){
    model.allDifferent(row(i,m,vars)).post();
    model.allDifferent(column(i,m,vars)).post();
}
for(int i=0;i<m;i+=blockSize){
    for(int j=0;j<m;j+=blockSize){
        model.allDifferent(block(i,j,blockSize,vars)).post();
    }
}
Solver solver = model.getSolver();
solver.solve();
```

Choco-solver Data

```
int[][] data = new int[m][m]{
    {4, 0, 8, 0, 0, 0, 0, 0, 0},
    {0, 0, 0, 1, 7, 0, 0, 0, 0},
    {0, 0, 0, 0, 8, 0, 0, 3, 2},
    {0, 0, 6, 0, 0, 8, 2, 5, 0},
    {0, 9, 0, 0, 0, 0, 0, 8, 0},
    {0, 3, 7, 6, 0, 0, 9, 0, 0},
    {2, 7, 0, 0, 5, 0, 0, 0, 0},
    {0, 0, 0, 0, 1, 4, 0, 0, 0},
    {0, 0, 0, 0, 0, 0, 6, 0, 4}
};
```

Choco-solver Utilities

```

IntVar[] row(int row, int size, IntVar[][] array){
    return array[row];
}

IntVar[] column(int col,int size,IntVar[][] array){
    IntVar[] column = new IntVar[size];
    for(int i=0; i<size; i++){
        column[i] = array[i][col];
    }
    return column;
}

IntVar[] block(int x,int y,int blockSize,IntVar[][] array){
    IntVar[] block = new IntVar[blockSize*blockSize];
    int k=0;
    for(int i=0;i<blockSize;i++){
        for(int j=0;j<blockSize;j++){
            block[k++] = array[x+i][y+j];
        }
    }
    return block;
}

```

▶ Continue

Domain Visualizer

- Problem shown as matrix
- Each cell corresponds to a variable
- Instantiated: Shows integer value (large)
- Uninstantiated: Shows values in domain

[illegible]

Initial State (Forward Checking)

4	1 2 3 4 5 6 7 8 9	8	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9
1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 7	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9
1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	8	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	3	2	
1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	6	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	8	2	5		
1 2 3 4 5 6 7 8 9	9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	8	1 2 3 4 5 6 7 8 9	
1 2 3 4 5 6 7 8 9	3	7	6	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	
2	7	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	5	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	
1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1	4	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	
1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	1 2 3 4 5 6 7 8 9	6	1 2 3 4 5 6 7 8 9	4	

Propagation Steps (Forward Checking)

4	1 2 5 6	8	2 3 5 9	2 3 6 9	2 3 5 6 9 7	1 5	1 6 7 9	1 5 6 7 9
3 6 9	2 5 6	3 5 9	1	7	2 3 5 6 9	4 5 8	4 6 9	5 6 8 9
1 6 7 9	1 5 6	1 5 9	4 5 9	8	5 6 9 7	4 5	3	2
1	4	6	3 7 9	3 9	8	2	5	3
5	9	2	4 7	3 4 6	1 3 7 6	1 3 4 7	8	1 3 6
8	3	7	6 4	2 5	1 2 5	9	1 2 4	1 5
2	7	1 3 4 9	3 8 9	5	1 3 6 9	1 3 4 8	1 4 6 9	1 3 6 8 9
3 6 7 9	2 5 6 8	3 5 9	2 3 5 9	1	4	5 7 8	3 6 9 7 8 9	3 5 6 7 8 9
3 7 9	1 2 5 8	1 5 9	2 3 5 9	2 3 5 9	1 2 3 5 9	6	1 2 7 9	4

After Setup (Forward Checking)

4	^{1 2} 5 6	8	^{2 3} 5 9	³ 6 9	^{2 3 1} 6 9 7	¹ 6 9	^{5 6} 7 9
³ 6 9	² 5 6	³ 5 9	1	7	^{2 3} 6 9	^{4 5} 8	⁶ 5 6 8 9
¹ 6 9	¹ 5 6	^{4 5} 9	8	¹ 6 9	^{4 5} 7	3	2
1	4	6	³ 7 9	³ 9	8	2	5
5	9	2	⁴ 7	³ 4	^{1 3} 7	³ 8	³ 6 7
8	3	7	6	2	5	9	4
2	7	^{1 3} 4 9	³ 8 9	5	^{3 1} 6 9	^{3 1} 8 9	³ 8 9
³ 6 9	^{5 6} 8	⁵ 9 7	^{2 3} 8 9	1	4	⁵ 7 8	² 7 9
^{3 1} 9 8	¹ 5 9	³ 7 8	^{2 3} 9	³ 7 9	^{2 3} 6	^{1 2} 7 9	4

Can we do better?

- The alldifferent constraint is missing propagation
 - How can we do more propagation?
 - Do we know when we derive all possible information from the constraint?
- Constraints only interact by changing domains of variables

A Simpler Example

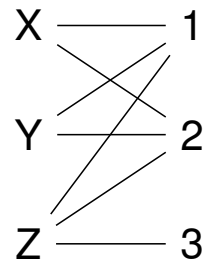
```
include "alldifferent.mzn";  
  
var 1..2:X;  
var 1..2:Y;  
var 1..3:Z;  
  
constraint alldifferent([X,Y,Z]);  
  
solve satisfy;
```

Using Forward Checking

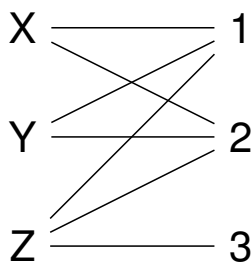
- No variable is assigned
- No reduction of domains
- But, values 1 and 2 can be removed from Z
- This means that Z is assigned to 3

Visualization of alldifferent as Graph

- Show problem as graph with two types of nodes
 - Variables on the left
 - Values on the right
- If value is in domain of variable, show link between them
- This is called a *bipartite* graph



A Simpler Example



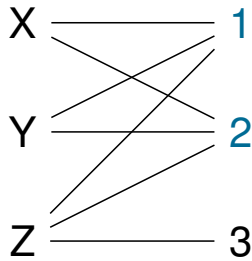
Value Graph for

```
var 1..2:X;
```

```
var 1..2:Y;
```

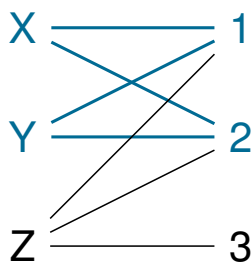
```
var 1..3:Z;
```

A Simpler Example



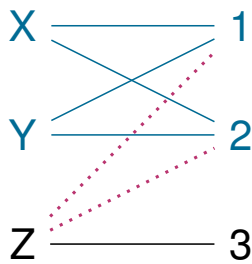
Check interval $[1,2]$

A Simpler Example



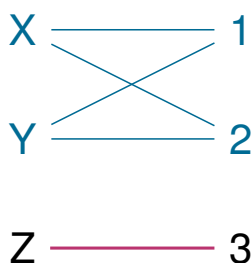
- Find variables completely contained in interval
- There are two: X and Y
- This uses up the capacity of the interval

A Simpler Example



No other variable can use that interval

A Simpler Example



Only one value left in domain of Z,
this can be assigned

Idea (Hall Intervals)

- Take each interval of possible values, say size N
- Find all K variables whose domain is completely contained in interval
- If $K > N$ then the constraint is infeasible
- If $K = N$ then no other variable can use that interval
- Remove values from such variables if their bounds change
- If $K < N$ do nothing
- Re-check whenever domain bounds change

Implementation

- Problem: Too many intervals ($O(n^2)$) to consider
- Solution:
 - Check only those intervals which update bounds
 - Enumerate intervals incrementally
 - Starting from lowest(highest) value
 - Using sorted list of variables
- Complexity: $O(n \log(n))$ in standard implementations
- Important: Only looks at min/max bounds of variables

Bounds Consistency

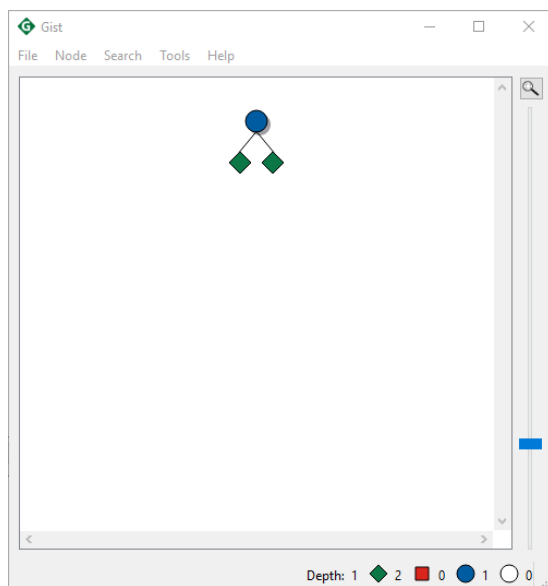
Definition

A constraint achieves *bounds consistency*, if for the lower and upper bound of every variable, it is possible to find values for all other variables between their lower and upper bounds which satisfy the constraint.

Annotation: :: bounds

```
include "alldifferent.mzn";  
  
var 1..2:X;  
var 1..2:Y;  
var 1..3:Z;  
  
constraint alldifferent([X,Y,Z]) :: bounds;  
solve satisfy;
```

Running with Gecode Gist



All Solutions

The screenshot shows the Gecode Gist Console window. It has a menu bar with 'Clear' and 'Stay on top'. The console displays the following state for the root node:

```
X = [1..2];  
Y = [1..2];  
Z = 3;
```

Node Inspector (Root)

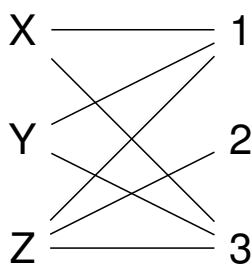
Can we do even better?

- Bounds consistency only considers min/max bounds
- Ignores “holes” in domain
- Sometimes we can improve propagation looking at those holes

Another Simple Example

```
include "alldifferent.mzn";  
  
var {1,3}:X; % note enumerated domain  
var {1,3}:Y;  
var 1..3:Z; % note domain as interval  
  
% annotated constraint  
constraint alldifferent([X,Y,Z]) :: bounds;  
solve satisfy;
```

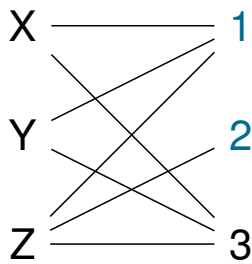
Another Simple Example



Value Graph for

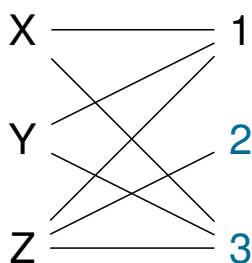
```
var {1,3}:X;  
var {1,3}:Y;  
var 1..3:Z;
```

Another Simple Example



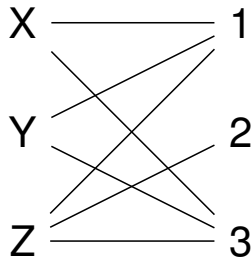
- Check interval $[1,2]$
- No domain of a variable completely contained in interval
- No propagation

Another Simple Example



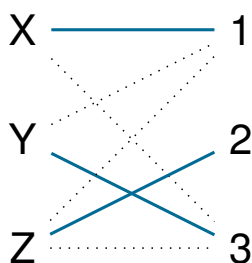
- Check interval $[2,3]$
- No domain of a variable completely contained in interval
- No propagation

Another Simple Example



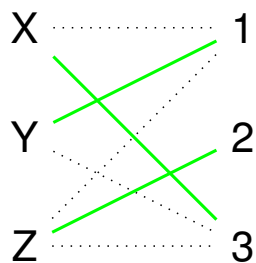
But, more propagation is possible,
there are only two solutions

Another Simple Example



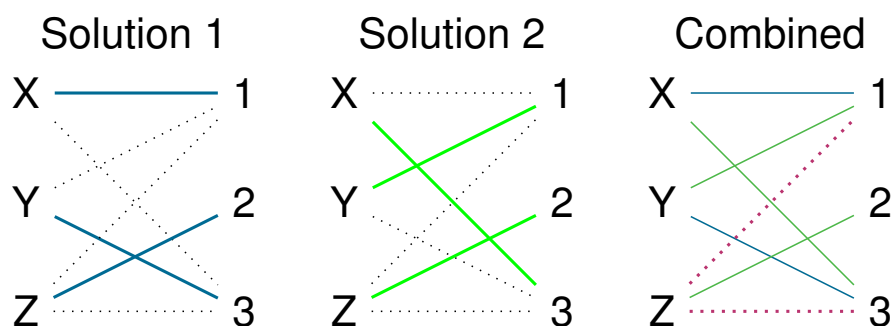
Solution 1: assignment in blue

Another Simple Example



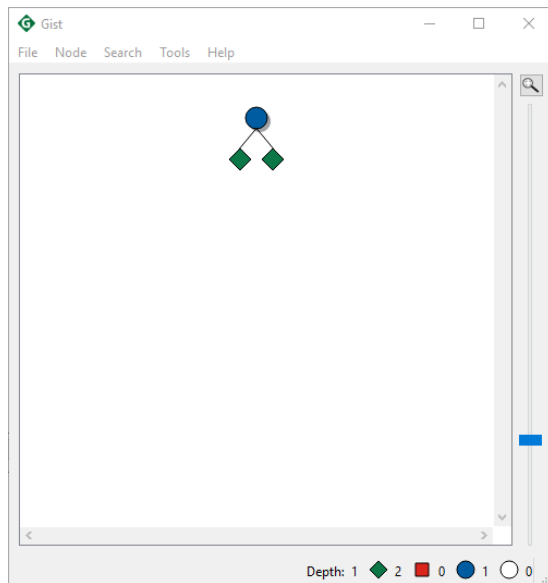
Solution 2: assignment in green

Another Simple Example

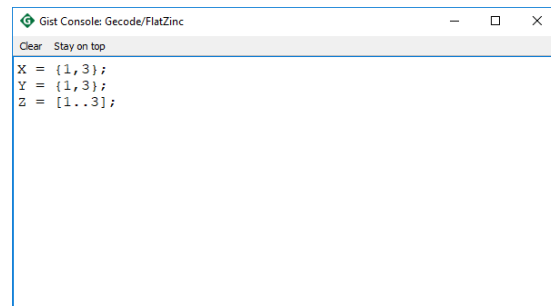


Combining solutions shows that $Z=1$ and $Z=3$ are not possible. Can we deduce this without enumerating solutions?

Bounds Consistency with Gecode Gist: No Propagation



All Solutions



Node Inspector (Root)

Solutions and Maximal Matchings

- A *Matching* is subset of edges which do not coincide in any node
- No matching can have more edges than number of variables
- Every solution corresponds to a *maximal matching* and vice versa
- If a link does not belong to some maximal matching, then it can be removed

Implementation

- Possible to compute all links which belong to some matching
 - Without enumerating all of them!
- Enough to compute **one** maximal matching
- Requires algorithm for *strongly connected components*
- Extra work required if more values than variables
- All links (values in domains) which are not supported can be removed
- Complexity: $O(n^{1.5}d)$

Domain Consistency

Definition

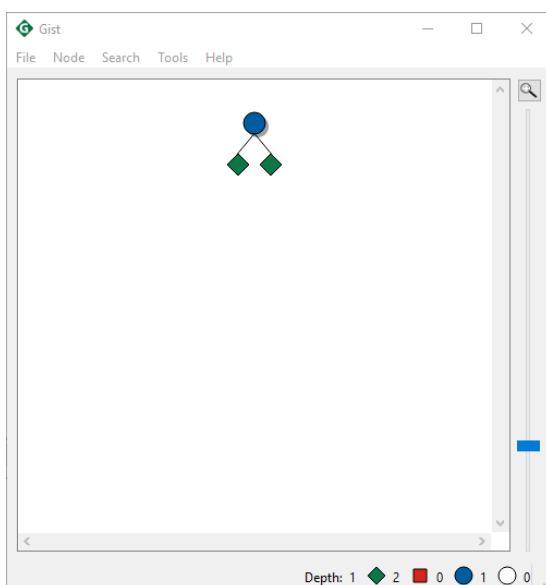
A constraint achieves *domain consistency*, if for every variable and for every value in its domain, it is possible to find values in the domains of all other variables which satisfy the constraint.

- Also called *generalized arc consistency (GAC)*
- or *hyper arc consistency*

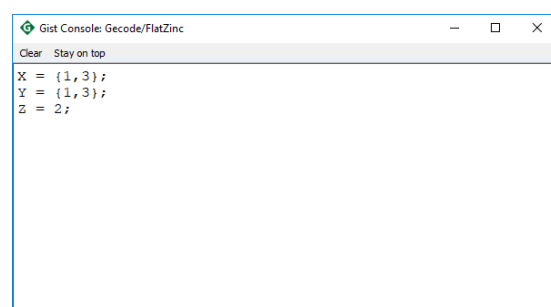
Simple Example Revisited

```
include "alldifferent.mzn";  
  
var {1,3}:X; % note enumerated domain  
var {1,3}:Y;  
var 1..3:Z; % note domain as interval  
  
% note different annotation  
constraint alldifferent([X,Y,Z]) :: domain;  
solve satisfy;
```

Domain Consistency with Gecode Gist: Propagation



All Solutions



Node Inspector (Root)

Can we still do better?

- NO! This extracts all information from this one constraint
- We could perhaps improve speed, but not propagation
- But possible to use different model
- Or model interaction of multiple constraints

Should all constraints achieve domain consistency?

- Domain consistency is usually more expensive than bounds consistency
 - Overkill for simple problems
 - Nice to have choices
- For some constraints achieving domain consistency is NP-hard
 - We have to live with more restricted propagation

Modified MiniZinc Program

```
int: s;
int: n=s*s;
array[1..n,1..n] of var 1..n: puzzle;

include "sudoku.dzn";

include "alldifferent.mzn";

constraint forall(i in 1..n)
    (alldifferent([puzzle[i,j] | j in 1..n])::domain);
constraint forall(j in 1..n)
    (alldifferent([ puzzle[i,j] | i in 1..n])::domain);
constraint forall(i,j in 1..s)
    (alldifferent([puzzle[s*(i-1)+p, s*(j-1)+q] |
                    p,q in 1..s])::domain);

solve satisfy;
```

Modified Choco-solver Sudoku Model

```
Model model = new Model("Sudoku");
int blockSize = 3;
int m = blockSize*blockSize;

IntVar[][] vars = new IntVar[m][m];
for(int i=0;i<m;i++){
    for(int j=0;j<m;j++){
        vars[i][j] = model.intVar("X"+i+" "+j, 1, m);
        if (data[i][j]>0) {
            model.arithm(vars[i][j],"=",data[i][j]).post();
        }
    }
}

// Consistency level AC: domain consistency, BC: bounds consistency, default: mix
for(int i=0;i<m;i++){
    model.allDifferent(row(i,m,vars),AC).post();
    model.allDifferent(column(i,m,vars),AC).post();
}
for(int i=0;i<m;i+=blockSize){
    for(int j=0;j<m;j+=blockSize){
        model.allDifferent(block(i,j,blockSize,vars),AC).post();
    }
}

Solver solver = model.getSolver();
solver.solve();
```

Initial State (Domain Consistency)

4	<div>1 2 3 4 5 6 7 8 9</div>	8	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	1	7	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	8	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	3	2	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	6	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	8	2	5	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	9	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	8	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	3	7	6	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	9	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
2	7	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	5	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	1	4	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>
<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	<div>1 2 3 4 5 6 7 8 9</div>	6	<div>1 2 3 4 5 6 7 8 9</div>	4	<div>1 2 3 4 5 6 7 8 9</div>

Propagation Steps (Domain Consistency)

4	2	8	5	6	3	1	<div>1 7</div>	<div>1 6</div>	<div>5 6</div>
<div>3 6</div>	5	<div>5 9</div>	1	7	2	4	6	8	
7	6	1	4	8	9	5	3	2	
1	4	6	<div>3 7</div>	<div>3 9</div>	8	2	5	<div>3 7</div>	
5	9	2	<div>3 7</div>	4	1	<div>1 4</div>	8	6	
8	3	7	6	2	5	9	4	1	
2	7	4	<div>3 8</div>	5	6	8	1	<div>1 6</div>	<div>3 8</div>
6	8	<div>5 9</div>	2	1	4	<div>5 7</div>	<div>2 6</div>	5	
<div>3 9</div>	1	5	8	<div>2 9</div>	7	6	2	4	

After Setup (Domain Consistency)

4	2	8	5	6	3	1		
	5		1	7	2	4	6	8
7	6	1	4	8	9	5	3	2
1	4	6			8	2	5	
5	9	2		4	1		8	6
8	3	7	6	2	5	9	4	1
2	7	4		5	6	8	1	
6	8		2	1	4			5
	1	5	8		7	6	2	4

Comparison

Forward Checking

4		8						
	5		1	7				
7	6	1	4	8	9	5	3	2
1	4	6			8	2	5	
5	9	2		4	1		8	6
8	3	7	6	2	5	9	4	1
2	7	4		5	6	8	1	
6	8		2	1	4			5
	1	5	8		7	6	2	4

Bounds Consistency

4		8	5	6				
	5		1	7				
7	6	1	4	8	9	5	3	2
1	4	6			8	2	5	
5	9	2		4	1		8	6
8	3	7	6	2	5	9	4	1
2	7	4		5	6	8	1	
6	8		2	1	4			5
	1	5	8		7	6	2	4

Domain Consistency

4	2	8	5	6	3	1		
	5		1	7	2	4	6	8
7	6	1	4	8	9	5	3	2
1	4	6			8	2	5	
5	9	2		4	1		8	6
8	3	7	6	2	5	9	4	1
2	7	4		5	6	8	1	
6	8		2	1	4			5
	1	5	8		7	6	2	4

Typical?

- This does not always happen
- Sometimes, two methods produce same amount of propagation
- Possible to predict in certain special cases
- In general, tradeoff between speed and propagation
- Not always fastest to remove inconsistent values early
- But often required to find a solution at all

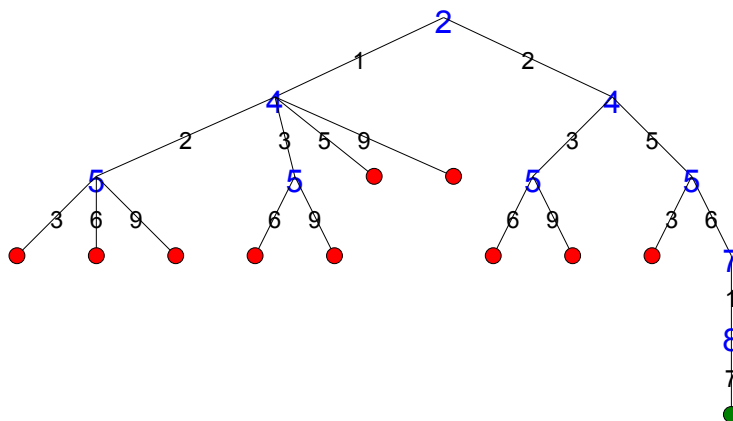
Simple search routine

- Enumerate variables in given order
- Try values starting from smallest one in domain
- Complete, chronological backtracking
- Advantage: Results can be compared with each other
- Disadvantage: Usually not a very good strategy

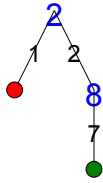
Asking for Naive Search in MiniZinc

```
solve :: int_search(
    puzzle,
    input_order,
    indomain_min)
satisfy;
```

Search Tree (Forward Checking)



Search Tree (Bounds Consistency)



Search Tree (Domain Consistency)



Trading Propagation Against Search

- If we perform more propagation, search is more constrained
- Fewer values left, fewer alternatives to explore in search
- Best compromise is not obvious
- But can be learned from examples or during search
- Annotations are optional
 - Some MiniZinc back-end solvers do the search they want, not the one you specify
 - Some solvers simply do not work in a way that these search annotations apply

Are there other Global Constraints?

- alldifferent is the most commonly used constraint
- Propagation methods can be explained
- But there are many more

Global Constraint Catalog

- <https://sofdem.github.io/gccat/>
- Description of 354 global constraints, 2800 pages
- Not all of them are widely used
- Detailed, meta-data description of constraints in Prolog

Families of Global Constraints

- Value Counting
 - alldifferent, global cardinality
- Scheduling
 - cumulative
- Properties of Sequences
 - sequence, no_valley
- Graph Properties
 - circuit, tree

Common Algorithmic Techniques

- Bi-Partite Matchning
- Flow Based Algorithms
- Automata
- Task Intervals
- Reduced Cost Filtering
- Decomposition

Part III

Customizing Search

What we want to introduce

- Importance of search strategy, constraints alone are not enough
- Two schools of thought
 - Black-box solver, solver decides by itself
 - Human control over process
- Dynamic variable ordering exploits information from propagation
- Variable and value choice
- Hard to find strategy which works all the time
- Different way of improving stability of search routine

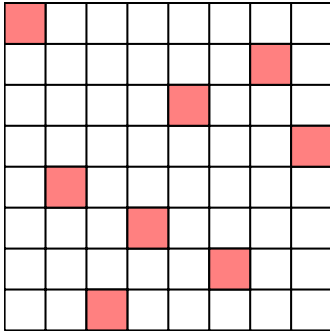
Example Problem

- N-Queens puzzle
- Rather weak constraint propagation
- Many solutions, limited number of symmetries
- Easy to scale problem size

Problem Definition

8-Queens

Place 8 queens on an 8×8 chessboard so that no queen attacks another. A queen attacks all cells in horizontal, vertical and diagonal direction. Generalizes to boards of size $N \times N$.



Solution for board size 8×8

Basic Model

- Cell based Model
 - A 0/1 variable for each cell to say if it is occupied or not
 - Constraints on rows, columns and diagonals to enforce no-attack
 - N^2 variables, $6N - 2$ constraints
- Column (Row) based Model
 - A 1..N variable for each column, stating position of queen in the column
 - Based on observation that each column must contain exactly one queen
 - N variables, $N^2/2$ binary constraints

assign $[X_1, X_2, \dots, X_N]$

s.t.

$$\forall 1 \leq i \leq N: X_i \in 1..N$$

$$\forall 1 \leq i < j \leq N: X_i \neq X_j$$

$$\forall 1 \leq i < j \leq N: X_i + j \neq X_j + i$$

$$\forall 1 \leq i < j \leq N: X_i + i \neq X_j + j$$

Nqueens Models

- ECLiPSe [▶ Show](#)
- MiniZinc [▶ Show](#)
- NumberJack [▶ Show](#)
- CPMpy [▶ Show](#)
- Choco-solver [▶ Show](#)

ECLiPSe N-Queens Model

```
:- lib(lists).
:- lib(ic).
:- lib(ic_search).

top:-
    queens(8,Board),
    search(Board, 0, input_order, indomain, complete).

queens(N, Board) :-
    length(Board, N),
    Board :: 1..N,
    ( fromto(Board, [Q1|Cols], Cols, []) do
        ( foreach(Q2, Cols), param(Q1), count(Dist,1,_) do
            noattack(Q1, Q2, Dist)
        )
    ).

noattack(Q1,Q2,Dist) :-
    Q2 #\= Q1,
    Q2 - Q1 #\= Dist,
    Q1 - Q2 #\= Dist.
```

► Continue

MiniZinc N-Queens Model

```
int: n=8;
array[1..n] of var 1..n: queens;
constraint
    forall(i, j in 1..n where i < j) (
        queens[i] != queens[j] /\
        queens[i] + i != queens[j] + j /\
        queens[i] - i != queens[j] - j
    )
;
solve :: int_search(
    queens,
    input_order,
    indomain_min)
satisfy;
```

► Continue

NumberJack N-Queens Model

```
from Numberjack import *

def get_model(N):
    queens = VarArray(N, N)
    model = Model(
        AllDiff(queens),
        AllDiff([queens[i] + i for i in range(N)]),
        AllDiff([queens[i] - i for i in range(N)])
    )
    return queens, model

def solve(param):
    queens, model = get_model(param['N'])
    solver = model.load(param['solver'])
    solver.setHeuristic(param['heuristic'], param['value'])
    solver.setVerbosity(param['verbose'])
    solver.setTimeLimit(param['tcutoff'])
    solver.solve()
```

► Continue

CPMpy N-Queens Model

```
def nqueens_naive(n=8):
    queens = IntVar(1,n, shape=n)

    model = Model()
    for i in range(n):
        for j in range(i):
            model += [queens[i] != queens[j],
                      queens[i] + i != queens[j] + j,
                      queens[i] - i != queens[j] - j,
                      ]
```

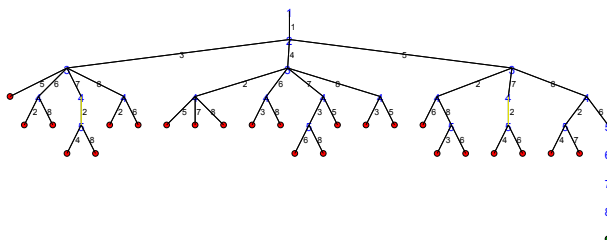
► Continue

Choco-solver N-Queens Program

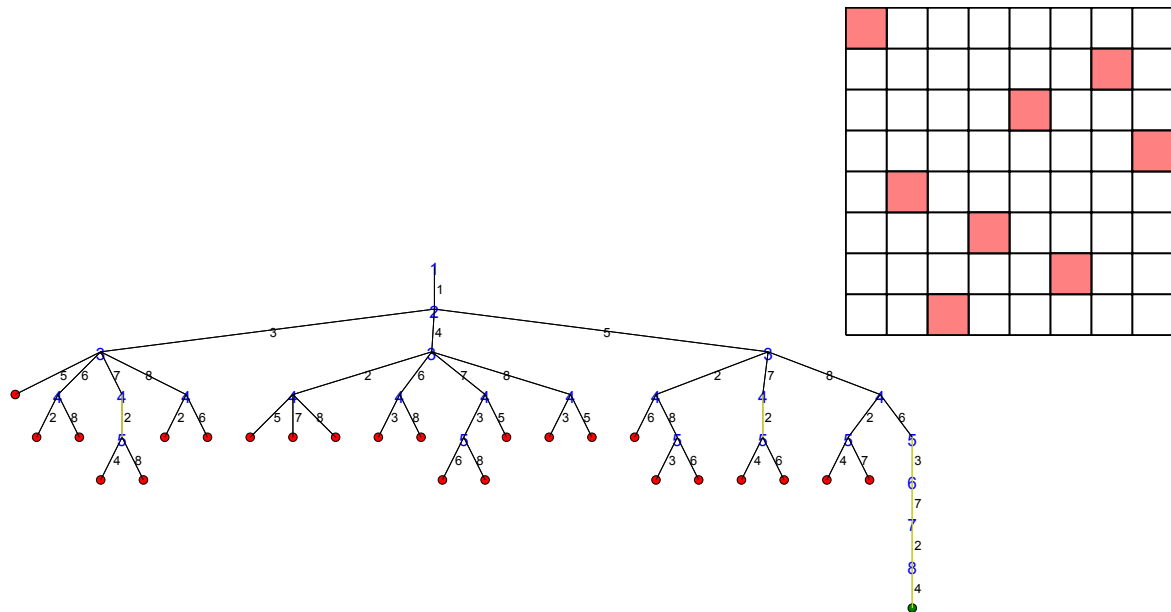
```
int n = 8;
Model model = new Model(n + "-queens problem");
IntVar[] vars = new IntVar[n];
for(int q = 0; q < n; q++){
    vars[q] = model.intVar("Q_"+q, 1, n);
}
for(int i = 0; i < n-1; i++){
    for(int j = i + 1; j < n; j++){
        model.arithm(vars[i], "!=" ,vars[j]).post();
        model.arithm(vars[i], "!=" , vars[j], "-", j - i).post();
        model.arithm(vars[i], "!=" , vars[j], "+", j - i).post();
    }
}
Solution solution = model.getSolver().findSolution();
if(solution != null){
    System.out.println(solution.toString());
}
```

► Continue

Default Strategy



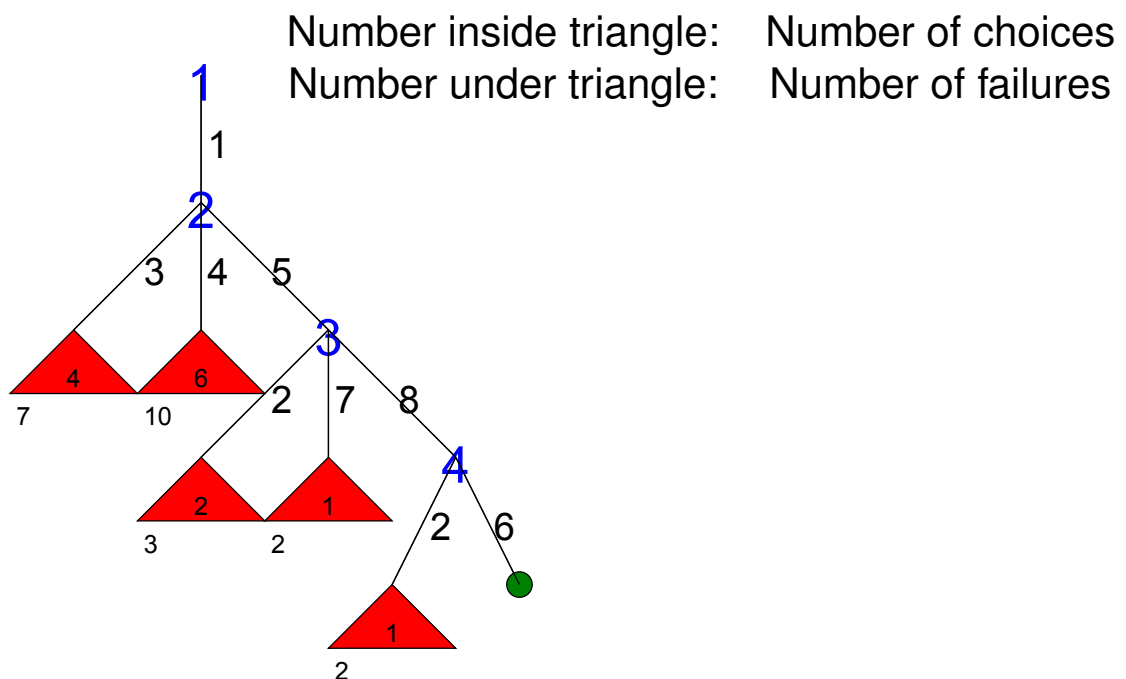
First Solution



Observations

- Even for small problem size, tree can become large
- Not interested in all details
- Ignore all automatically fixed variables
- For more compact representation abstract failed sub-trees

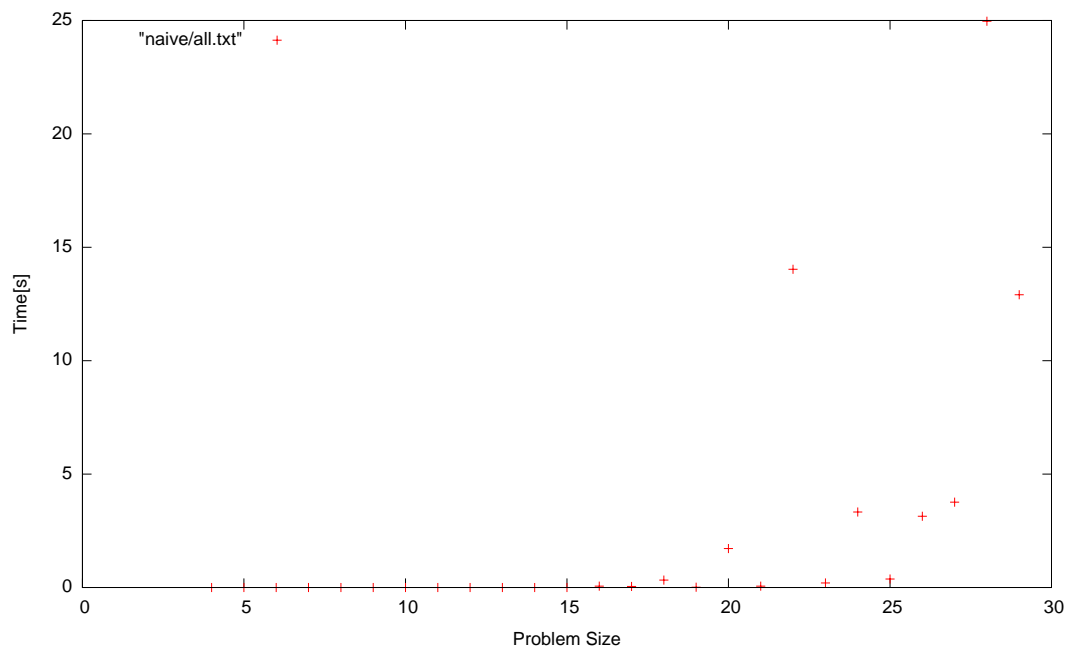
Compact Representation



Exploring other board sizes

- How stable is the model?
- Try all sizes from 4 to 100
- Timeout of 100 seconds

Naive Strategy, Problem Sizes 4-100



Observations

- Time very reasonable up to size 20
- Sizes 20-30 times very variable
- Not just linked to problem size
- No size greater than 30 solved within timeout

Possible Improvements

- Better constraint reasoning
 - Remodelling problem with 3 `alldifferent` constraints
 - Global reasoning as described before
- Better control of search
 - Static vs. dynamic variable ordering
 - Better value choice
 - Not using complete depth-first chronological backtracking

Static vs. Dynamic Variable Ordering

- Heuristic Static Ordering
 - Sort variables before search based on heuristic
 - Most important decisions
 - Smallest initial domain
- Dynamic variable ordering
 - Use information from constraint propagation
 - Different orders in different parts of search tree
 - Use all information available

First Fail strategy

- Dynamic variable ordering
- At each step, select variable with smallest domain
- Idea: If there is a solution, better chance of finding it
- Idea: If there is no solution, smaller number of alternatives
- Needs tie-breaking method

Search Strategy Choices

- Minizinc [▶ Show](#)
- Choco-solver [▶ Show](#)

Modified MiniZinc Program

```
int: n=8;
array[1..n] of var 1..n: queens;
constraint
    forall(i, j in 1..n where i < j) (
        queens[i] != queens[j] /\
        queens[i] + i != queens[j] + j /\
        queens[i] - i != queens[j] - j
    )
;
solve :: int_search(
    queens,
    first_fail,
    indomain_min)
satisfy;
```

Variable Choice (MiniZinc)

- Determines the order in which variables are assigned
- `input_order` assign variables in static order given
- `smallest` assign variable with smallest value in domain first
- `first_fail` select variable with smallest domain first
- `dom_w_deg` consider ratio of domain size and failure count
- Others, including programmed selection for specific solvers

Value Choice (MiniZinc)

- Determines the order in which values are tested for selected variables
- `indomain_min` Start with smallest value, on backtracking try next larger value
- `indomain_median` Start with value closest to middle of domain
- `indomain_random` Choose values in random order
- `indomain_split` Split domain into two intervals

► Continue

Modified Choco-solver Model

```
int n = 8;

Model model = new Model(n + "-queens problem");
IntVar[] vars = model.intVarArray("Q", n, 1, n, false);
IntVar[] diag1 = IntStream.range(0, n).
    mapToObj(i -> vars[i].sub(i).intVar()).
    toArray(IntVar[]::new);
IntVar[] diag2 = IntStream.range(0, n).
    mapToObj(i -> vars[i].add(i).intVar()).
    toArray(IntVar[]::new);

model.post (
    model.allDifferent(vars),
    model.allDifferent(diag1),
    model.allDifferent(diag2)
);

Solver solver = model.getSolver();
solver.showStatistics();
solver.setSearch(Search.domOverWDegSearch(vars));
Solution solution = solver.findSolution();

if (solution != null) {
    System.out.println(solution.toString());
}
```

VariableSelector Choice (Choco-solver)

- Determines the order in which variables are assigned
- `InputOrder` assign variables in static order given
- `Smallest` assign variable with smallest value in domain first
- `FirstFail` select variable with smallest domain first
- `DomOverWDeg` consider ratio of domain size and failure count
- `ActivityBased` dynamic, based on dynamic observed behaviour
- `ImpactBased` dynamic, based on dynamic observed behaviour

IntValueSelector Choice (Choco-solver)

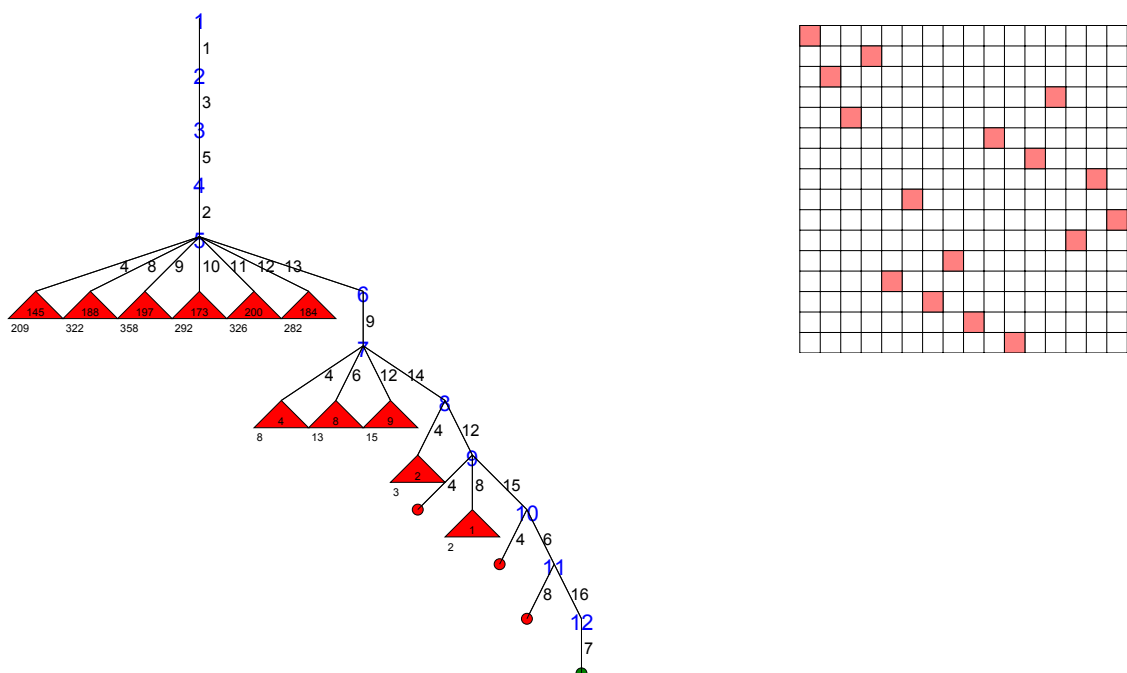
- Determines the order in which values are tested for selected variables
- `IntDomainMin` Start with smallest value, on backtracking try next larger value
- `IntDomainMiddle` Start with value closest to middle of domain
- `IntDomainRandom` Choose values in random order
- `IntDomainRandomBound` Randomly choose between smallest and largest value

► Continue

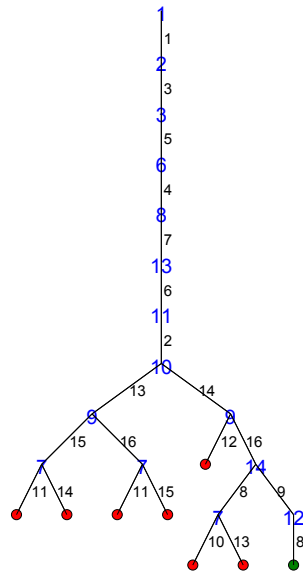
Comparison

- Board size 16x16
- Naive (Input Order) Strategy
- First Fail variable selection

Naive (Input Order) Strategy (Size 16)

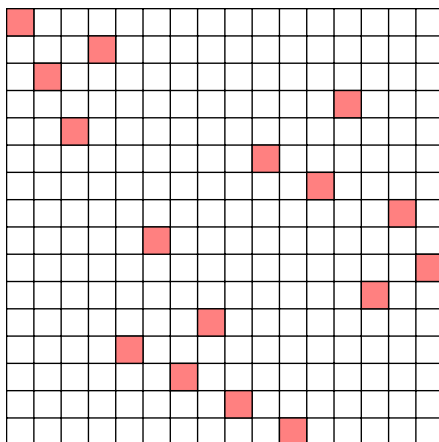


FirstFail Strategy (Size 16)

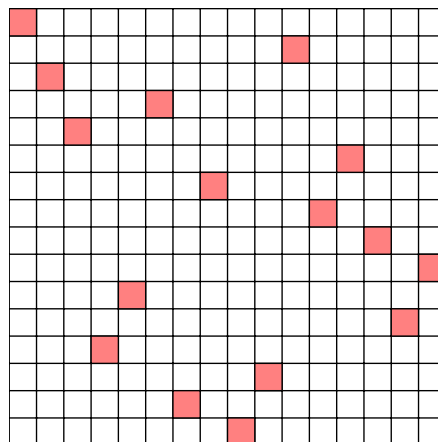


Comparing Solutions

Naive

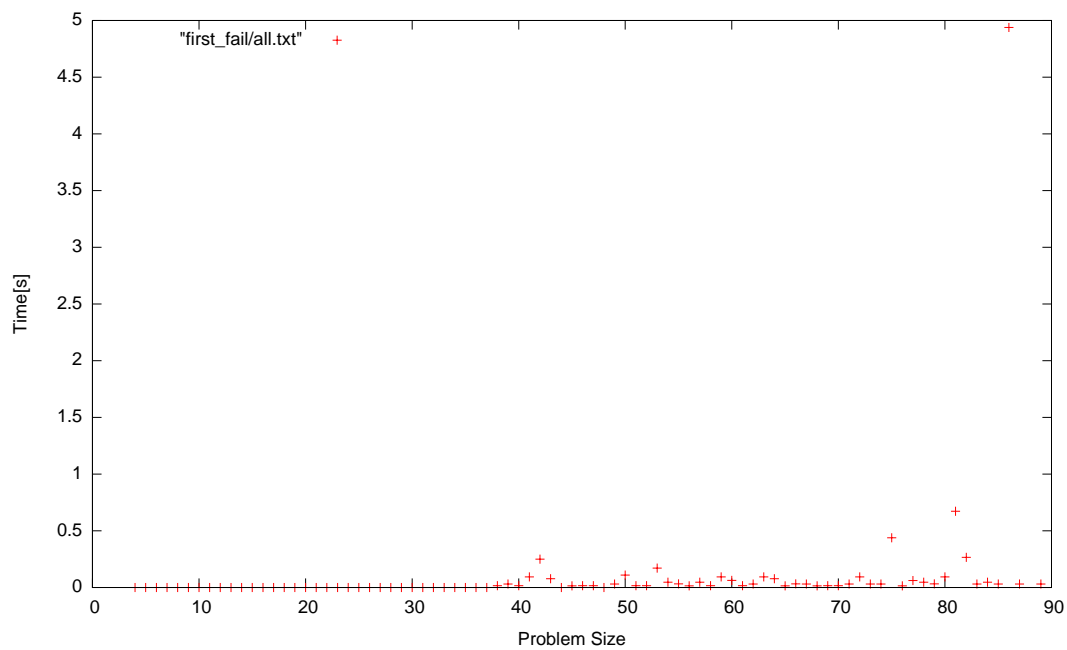


First Fail



Solutions are different!

FirstFail, Problem Sizes 4-100



Observations

- This is much better
- But some sizes are much harder
- Timeout for sizes 88, 91, 93, 97, 98, 99

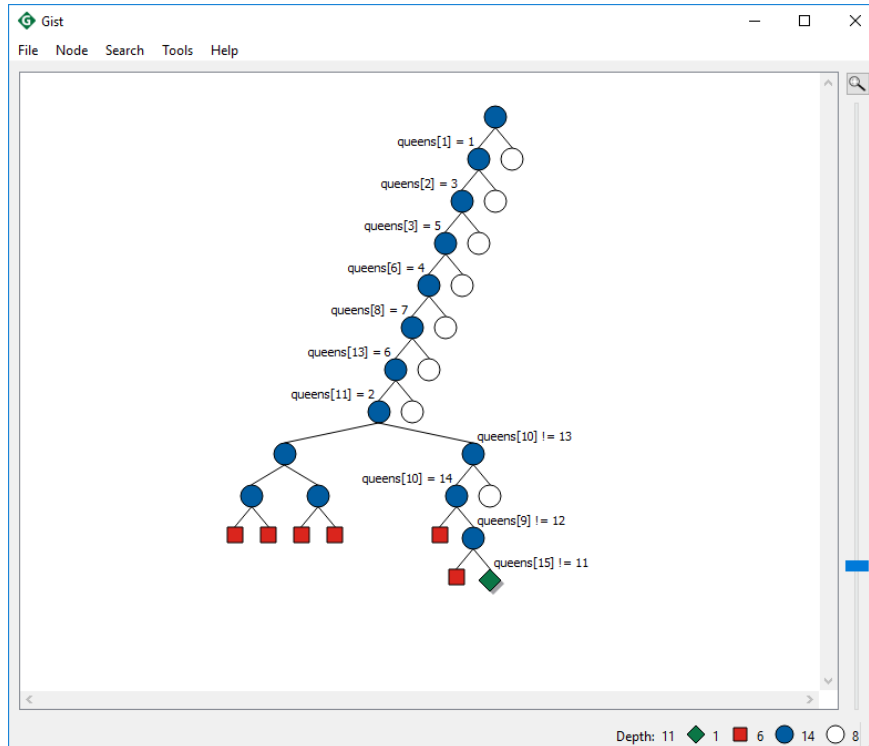
More Reactive Variable Selection

- Domain size is important, but other information is useful as well
- Dom/Weighted Degree: better results in many situations
- Weight Degree: count how often variable has been involved in failure
- Focus on more complicated part of problem
- Changes during search, learns from past performance
- Option **dom_w_deg**

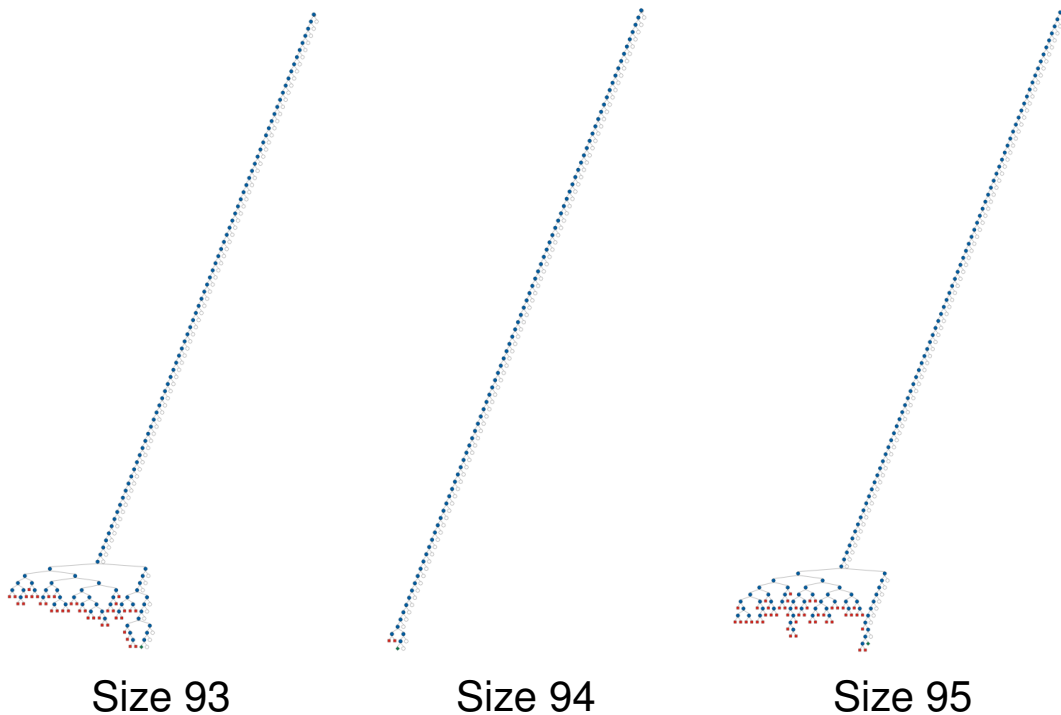
Weighted Degree Variable Selection

```
int: n=8;
array[1..n] of var 1..n: queens;
constraint
    forall(i, j in 1..n where i < j) (
        queens[i] != queens[j] /\
        queens[i] + i != queens[j] + j /\
        queens[i] - i != queens[j] - j
    )
;
solve :: int_search(
    queens,
    dom_w_deg,
    indomain_random)
satisfy;
```

Result for size 16 with Gecode-Gist



Sample Results for Larger Sizes



Approach 1: Heuristic Portfolios

- Try multiple strategies for the same problem
- With multi-core CPUs, run them in parallel
- Only one needs to be successful for each problem

Approach 2: Restart with Randomization

- Only spend limited number of backtracks for a search attempt
- When this limit is exceeded, restart at beginning
- Requires randomization to explore new search branch
- Randomize variable choice by random tie break
- Randomize value choice by shuffling values
- Needs strategy when to restart

Random Variable Choice and Restarts

```
int: n=8;
array[1..n] of var 1..n: queens;
constraint
    forall(i, j in 1..n where i < j) (
        queens[i] != queens[j] /\
        queens[i] + i != queens[j] + j /\
        queens[i] - i != queens[j] - j
    )
;
solve :: int_search(
    queens,
    dom_w_deg,
    indomain_random)
    :: random_linear(100)
    satisfy;
```

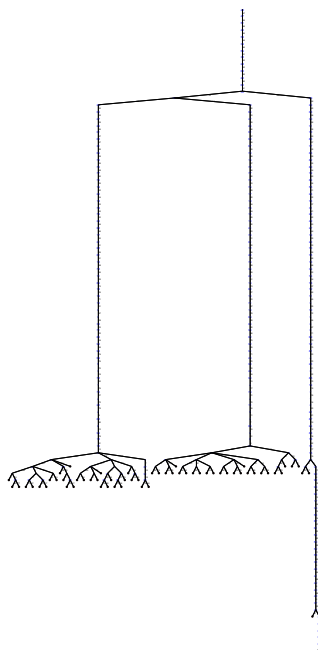
Approach 3: Partial Search

- Abandon depth-first, chronological backtracking
- Don't get locked into a failed sub-tree
- A wrong decision at a level is not detected, and we have to explore the complete subtree below to undo that wrong choice
- Explore more of the search tree
- Spend time in promising parts of tree

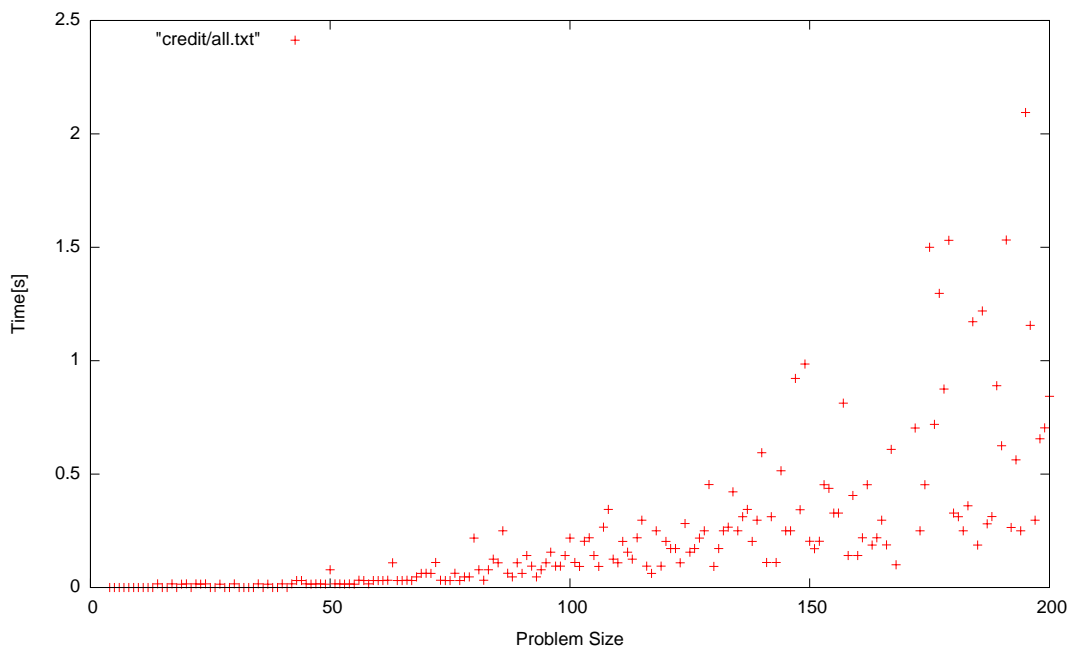
Example: Credit Search

- Not available in all solvers
- Explore top of tree completely, based on credit
- Start with fixed amount of credit
- Each node consumes one credit unit
- Split remaining credit amongst children
- When credit runs out, start bounded backtrack search
- Each branch can use only K backtracks
- If this limit is exceeded, jump to unexplored top of tree

Credit, Search Tree Problem Size 94



Credit, Problem Sizes 4-200



Points to Remember

- Choice of search can have huge impact on performance
- Dynamic variable selection can lead to large reduction of search space
- Packaged search can do a lot, but programming search adds even more
- Depth-first chronological backtracking not always best choice
- How to control this explosion of search alternatives?

Part IV

What is missing?

Many Specialized Topics

- How to design efficient core engine
- Hybrids with LP/MIP tools
- Hybrids with SAT
- Symmetry breaking
- Use of MDD/BDD to encode sets of solutions
- High level modelling tools
- Debugging/visualization

Reformulation

- Just because the user has modelled it this way, it doesn't mean we have to solve it that way
 - Replace some constraint(s) by other, equivalent constraints
 - Because we don't have that constraint in our system
 - For performance

Learning

- While solving the problem we can learn how to strengthen the model/search
 - Understand which constraints/method contribute to propagation and change schedule
 - Learn no-good constraints by explaining failure
 - Adapt search strategy based on search experience

More Learning Resources

- Survey of Methods, Resources, and Formats for Teaching Constraint Programming
 - by Tejas Santanam, Helmut Simonis
 - <https://doi.org/10.48550/arXiv.2403.12717>
 - Based on survey of community for WTCP 2023
 - <https://hsimonis.github.io/WTCP2023/>