hw2 Writeup

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RSA

• Overview:

In this challenge, the steps I took are listed as follows:

```
1. use binary search to find out the original p where n = p * q1 * q2
```

- 2. calculate d where $d = e^{(-1)} \mod (p-1)(q1-1)(q2-1)$
- 3. decrypt the cyphertext c with the private key (n,d)

• Observation:

- getPrime(n) is a function from the package pycryptodome which generates a random n-bit prime number.
- o for every p we pick, q1 and q2 can be deduced

```
p = getPrime(512)
q1 = next_prime(2 * p)  # q1 depends on p
q2 = next_prime(3 * q1)  # q2 depends on q1
```

o since the domain of p is [2^511, 2^512 - 1], bruteforcing this entire domain can be really slow. Hence, we can employ binary search to find out the correct p.

• Exploitation:

1 of 3 10/27/20, 3:34 PM

```
left = 2 ** 511
right = 2 ** 512 - 1
while left <= right:
    mid = left + int((right - left) / 2)
    p = mid
    q1 = next_prime(2 * p)
    q2 = next_prime(3 * q1)
    n = p * q1 * q2
    if n > real_n:
        right = mid - 1
    elif n < real n:
        left = mid + 1
    else: # n == real_n
        d = inverse(e, (p - 1) * (q1 - 1) * (q2 - 1))
        m = pow(real_c, int(d), n)
        print(long_to_bytes(m))
        break
```

• Flag:

flag: FLAG{Ew9xeANumjDr6bXemHsh}
exploit: please see the attachment

LSB (LSB Oracle Attack)

Overview:

- 1. Modify the exploit: https://github.com/oalieno/Crypto-Course/blob/master/RSA/LSB-Oracle/solve2.py (https://github.com/oalieno/Crypto-Course/blob/master/RSA/LSB-Oracle/solve2.py)
- 2. Change all base2-related stuff to base3

• Flag:

flag: FLAG{nF9Px2LtlNh5fJiq3QtG}

exploit:

```
#!/usr/bin/env python3
from pwn import *
from Crypto.Util.number import *
r = remote('140.112.31.97', 30001)
n = int(r.recvline().split(b' = ')[1])
c = int(r.recvline().split(b' = ')[1])
e = 65537
inv = inverse(3, n)
inve = pow(inv, e, n)
flag, x = 0, 0
for i in range(1024):
    r.sendline(str(c))
    m = int(r.recvline().split(b' = ')[1])
    bit = (m - x) \% 3
    x = inv * (x + bit) % n
    flag += bit * pow(3, i)
    print(long_to_bytes(flag))
    c = (c * inve) % n
print(long_to_bytes(flag))
```

3 of 3 10/27/20, 3:34 PM