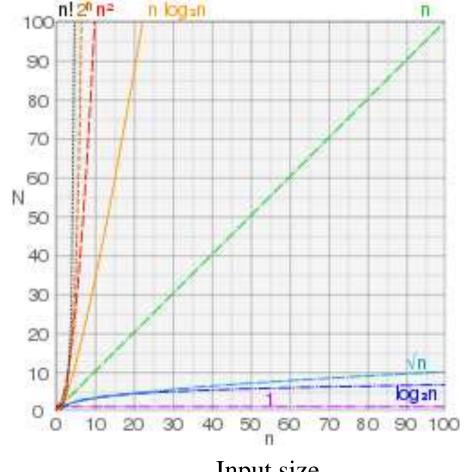
# Time Complexity of Algorithms

# Graphs of functions

#### **#Operations**

- > Show the number of operations *N* versus input size *n* for each function.
- Used in the analysis of algorithms

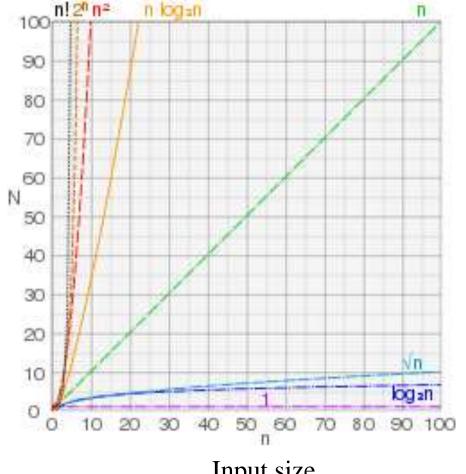


Input size

# Graphs of functions

#### **#Operations**

- > Show the number of operations *N* versus input size *n* for each function.
- Used in the analysis of algorithms



Input size

# Time Complexity

- Time complexity of an <u>algorithm</u> quantifies the amount of time taken by an algorithm to run as a <u>function</u> of the length of the <u>string</u> representing the input.
- The time complexity of an algorithm is commonly expressed using <a href="big O">big O</a> <a href="notation">notation</a>, which excludes coefficients and lower order terms.
- > Asymptotical time complexity: the time complexity of an algorithm as the input size goes to infinity.
- $\triangleright$  e.g.: if the time required by an algorithm on all inputs of size n is at most  $5n^3 + 3n$  for any n (bigger than some  $n_0$ ), the asymptotic time complexity is  $O(n^3)$ .

```
vector a;
int n;
void foo(int n) {
        for (int i = 0; i < n; ++i)
                 a.push_back( i );
void g() {
        cin >> n;
        foo(n);
What is the time complexity of the function foo?
```

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int n;
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void g() {
         cin >> n;
         foo(n);
```

```
vector a;
int n;
void foo(int n) {
         for (int i = 0; i < n; ++i)
                   a.push_back( i );
void g() {
         cin >> n;
         <u>foo</u>(n);
```

```
vector a;
int n;
void foo(int n) {
        for (int i = 0; i < n; ++i)
                a.push_back(i);
void g() {
        cin >> n;
        foo(n);
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What is the time complexity of the function foo?
```

Cost term	Item & Purpose	Times
		1

```
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int n;
void foo(int n) {
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                 a.push_back( i );
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What is the time complexity of the function foo?
```

Cost term	Item & Purpose	Times
	A1	1
	A2	
	A3	
	A4	
	A5	
	A6	
		n*C5

```
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int n;
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        cin >> n;
        foo(n);
What is the time complexity of the function foo?
```

Cost term	Item & Purpose	Times
	Parameter passing and function call	1
	Initialize counter i	
	i < n	n+1
	++i	
	Pass parameter to a.push_back	
	Invoke a.push_back	

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vector a;
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What is the time complexity of the function foo?
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Cost term	Item & Purpose	Times
CO	Parameter passing and function call	1
C1	Initialize counter i	
C2	i < n	
C3	++i	
C4	Pass parameter to a.push_back	
C5	Invoke a.push_back	
С		
= total cost		

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What is the time complexity of the function foo?
```

Cost term	Item & Purpose	Times
C0	Parameter passing and function call	A1
C1	Initialize counter i	A2
C2	i < n	A3
C3	++j	A4
C4	Pass parameter to a.push_back	A5
C5	Invoke a.push_back	<b>A6</b>
C = total cost		n*C5

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What is the time complexity of the function foo?
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Cost term	Item & Purpose	Times
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C1	Initialize counter i	1
C2	i < n	n+1
C3	++j	n
C4	Pass parameter to a.push_back	n
C5	Invoke a.push_back	n
C = total cost	C0+C1 +(n+1)*C2+n*C3+n*C4 +n*C5 = C0+C1+C2 + n*(C6) = O(n)	

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$$C6 = C2 + C3 + C4 + C5$$

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C = total cost	C0+C1 +(n+1)*C2+n*C3+n*C4 +n*C5 = C0+C1+C2 + n*(C6) = O(n*C6) = O(n)	

$$C6 = C2 + C3 + C4 + C5$$

#### **Intuitive idea**

Assume that a.push\_back takes k units of time to execute and k is a fixed value. Also, there is an extra amount of time to set the for-loop, which is c.

Then it will take c + k \* n units of time.

#### Thus:

Time complexity of the algorithm

$$= c + k*n = O(c + k*n) = O(n).$$

As n -> infinity, c is ignored. k is a constant, it is eliminated.

Cost term	Item & Purpose	Times
C0	Parameter passing and function call	1
C1	Initialize counter i	1
C2	i < n	n+1
C3	++i	n
C4	Pass parameter to a.push_back	n
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C = total cost	C0+C1 +(n+1)*C2+n*C3+n*C4 +n*C5 = $C0+C1+C2 + n*(C6)$ = $O(n*C6) = O(n)$	

$$C6 = C2 + C3 + C4 + C5$$

# Time Complexity

Big-O: Computational complexity: refer to the upper bound for the asymptotic computational complexity of an algorithm or a problem, which is usually written in terms of the big-O notation, e.g.,  $O(n^3)$ .

Big Omega  $\Omega(n)$ : Other types of (asymptotic) computational complexity estimates are lower bounds ("Big Omega" notation; e.g.,  $\Omega(n)$ ) and asymptotically tight estimates.

When the asymptotic upper and lower bounds coincide (written using the "big Theta"; e.g.,  $\Theta(n \log n)$ ).

asymptotic: approaching a value or curve arbitrary closely

# Time Complexity

Time complexity can be one of the followings:

```
>O(1)
                                 constant time
>O(log n)
>O(n)
                                 linear
\trianglerightO(n log n)
>O(n^2)
                                quadratic
>O(n^3)
>0(2^n)
                                exponential
\trianglerightO(n<sup>log n</sup>), and more
```

- Assume k is a constant. n is varying (depends on input size). Use big-O notation to indicate the following functions:
- 1. 5
- 2. 5 + k
- 3. k + k \* k
- 4. n\*n+n
- 5.  $n^*(k^*n+k^5)$
- 6. n\*k + n\*log n

• Assume k is a constant. n is varying (depends on input size). Use big-O notation to indicate the following functions:

- 1. 5
- 2. 5 + k
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- 4. n\*n+n
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- 6. n\*k + n\*log n

- = A1
- = A2
- = A3
- = A4
- = A5
- = A6

 Assume k is a constant. n is varying (depends on input size). Use big-O notation to indicate the following functions:

1. 
$$5 = O(1)$$

2. 
$$5 + k = O(1)$$

3. 
$$k + k *k = O(1)$$

4. 
$$n*n+n = O(n^2)$$

5. 
$$n^*(k^*n+k^5) = O(n^2)$$

6. 
$$n*k + n * log n$$
 = O(n log n)

```
bool find(int key, const vector<int> &a) {
         bool flg = false;
         for (int i = 0; i < a.size(); ++i) {
                  if (a[i] == key) {
                           flg = true; break;
         return flg;
```

```
bool find(int key, const vector<int> &a) { // C0
        bool flg = false;
                                             // C1
        for (int i = 0; i < a.size(); ++i) {
                  if (a[i] == key) {
                           flg = true; break;
         return flg;
n = a.size()
```

```
bool find(int key, const vector<int> &a) { // C0
        bool flg = false;
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```

```
(1, n+1, n), (C2, C3, C4)
```

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bool find(int key, const vector<int> &a) { // C0
       bool flg = false;
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                                                   (1, n+1, n), (C2, C3, C4)
       for (int i = 0; i < a.size(); ++i) {
                                                   n, C5
                if (a[i] == key) {
                        flg = true; break;
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bool find(int key, const vector<int> &a) { // C0
       bool flg = false;
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       for (int i = 0; i < a.size(); ++i) { (1, n+1, n), (C2, C3, C4)
                                                 n, C5
               if ( a[i] == key) {
                       flg = true; break; 🛑
                                                  1, C6
       return flg;
n = a.size()
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bool find(int key, const vector<int> &a) { // C0
       bool flg = false;
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       for (int i = 0; i < a.size(); ++i) { (1, n+1, n), (C2, C3, C4)
                                                  n, C5
               if ( a[i] == key) {
                       flg = true; break; 🛑
                                                   1, C7
       return flg;
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bool find(int key, const vector<int> &a) { // C0
       bool flg = false;
                                        // C1
                                                  (1, n+1, n), (C2, C3, C4)
       for (int i = 0; i < a.size(); ++i) {
                                                   n, C5
               if ( a[i] == key) {
                       flg = true; break; 🛑
                                                    1, C7
       return flg;
O(n), where n is the size of a.
```

```
bool find(int key, const vector<int> &a) {
        bool flg = false;
        for (int i = 0; i < a.size(); ++i) {
                if (a[i] == key) {
                        flg = true; break;
        return flg;
Assume there are m keys. We need to find
whether they appear in a. What is the
time complexity?
```

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                        flg = true; break;
        return flg;
Assume there are m keys. We need to find
whether they appear in a. What is the
time complexity? m*O(n) = O(m n)
```

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bool find(int key, const vector<int> &a) {
        bool flg = false;
        for (int i = 0; i < a.size(); ++i) {
                if (a[i] == key) {
                         flg = true; break;
        return flg;
Assume that there are m keys to find. If m
is a constant, what is the time complexity?
```

```
bool find(int key, const vector<int> &a) {
        bool flg = false;
        for (int i = 0; i < a.size(); ++i) {
                if (a[i] == key) {
                        flg = true; break;
        return flg;
Assume that there are m keys to find. If m
is a constant, what is the time complexity?
m*O(n) = O(m n) = O(n)
```

## Big-O notation

That is, f(x) = O(g(x)) if and only if there exists a positive real number M and a real number  $x_0$  such that

$$| f(x) | \le M | g(x) |$$
, for all  $x >= x_0$ .

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, for all  $x >= x_0$ .

- 1. 5 = O(1)
- 2. 5 + k = O(1)
- 3. k + k \*k = O(1)
- 4.  $x*x+x = O(x^2)$
- 5.  $x*(k*x+k^{5000000}) = O(x^2)$
- 6. x\*k + x \* log x = O(x log x)

k is a constant.

# Supplemental Materials

#### Insert an element to a linked list

```
void insert( Node *node ) {
    node->next = head;
    if (!head) tail = node;
    head = node;
}
```

What is the time complexity?

## Append an element to a linked list

```
void append( Node *node ) {
      if (tail) {
             tail->next = node;
             tail = node;
      } else {
             head = tail = node;
             node->next = nullptr;
What is the time complexity?
```

### Find an element in a linked list

```
bool find ( int key ) {
      Node *t = head;
      while (t) {
             if (t->key == key) return true;
             t = t->next;
      return false;
What is the time complexity?
```

### Bubble sort

```
void bubbleSort( int a[], int n) {
   while (true) {
       bool swap_flg = false;
       for (int i = 0; i < n-1; ++i) {
               if (a[i] > a[i+1]) {
                       swap(a[i], a[i+1]);
                       swap_flg = true;
       if ( !swap_flg ) break;
What is the time complexity?
```

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### Bubble sort

What is the time complexity?

```
void bubbleSort( int a[ ], int n) {
   while (true) {
        bool swap_flg = false;
       for (int i = 0; i < n-1; ++i) {
               if (a[i] >= a[i+1]) {
                                              // Can we use >=?
                       swap(a[i], a[i+1]);
                       swap_flg = true;
       if ( !swap_flg ) break;
```

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### Big-O notation

Let f and g be two functions defined on some subset of the real numbers. One writes

$$f(x) = O(g(x))$$
, as  $x \rightarrow \infty$ 

if and only if there is a positive constant M such that for all sufficiently large values of x, the absolute value of f(x) is at most M multiplied by the absolute value of g(x).