

# Distance-Regular Graphs

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# Outline

## 1 Distance-Regular Graphs



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# Distance-Regular Graphs



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# Notations

We always assume  $\Gamma = (V, E)$  is a connected graph with diameter  $D$ . For  $x \in X$ ,

$$\Gamma_i(x) := \{y \in X \mid \partial(x, y) = i\}.$$



# Distance-regular graphs

$\Gamma = (V, E)$  is **distance-regular** if and only if for  $i \leq D$ ,

$$c_i := |\Gamma_1(x) \cap \Gamma_{i-1}(y)|,$$

$$a_i := |\Gamma_1(x) \cap \Gamma_i(y)|,$$

$$b_i := |\Gamma_1(x) \cap \Gamma_{i+1}(y)|$$

are **constants** subject to all vertices  $x, y$  with  $\partial(x, y) = i$ .



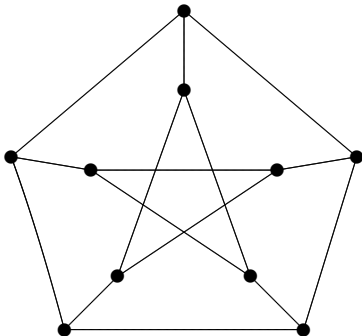
Note that  $a_i + b_i + c_i = b_0$  and  $k := b_0$  is the **valency** of  $\Gamma$ .



# Example

The **Petersen graph** below is a distance-regular with diameter  $D = 2$  and

$$b_0 = 3, c_1 = 1, a_1 = 0, b_1 = 2, c_2 = 1, a_2 = 2, b_2 = 0.$$



## Remark

If  $\Gamma$  is distance-regular with diameter  $D$  and fix a vertex  $x \in V\Gamma$ , then according to the partition

$$\Gamma_0(x) \cup \Gamma_1(x) \cup \cdots \cup \Gamma_D(x)$$

the adjacency matrix of  $\Gamma$  has equitable quotient matrix in the following tridiagonal form

$$\begin{pmatrix} a_0 & b_0 & & & 0 \\ c_1 & a_1 & b_1 & & \\ & \ddots & \ddots & \ddots & \\ & & c_{D-1} & a_{D-1} & b_{D-1} \\ 0 & & & c_D & a_D \end{pmatrix}$$



## Questions

From now on, we assume  $\Gamma = (V, E)$  is a distance-regular graph with diameter  $D$ . For any two vertices  $x, y \in X$  with distance  $\partial(x, y) = k$ , can you compute the **intersection numbers**

$$p_{ij}^k(x, y) := |\Gamma_i(x) \cap \Gamma_j(y)|?$$

Does the numbers  $p_{ij}^k(x, y)$  depend on the choice of vertices  $x, y$ ?





## Remarks on $p_{ij}^k(x, y)$

- ①  $p_{ij}^k(x, y) = 0$  if  $|i + j| < k$  or  $|i - j| > k$ .
- ②  $p_{1i-1}^i = c_i$ ,  $p_{1i}^i = a_i$ , and  $p_{1i+1}^i = b_i$ .
- ③

$$1 = c_1 \leq c_2 \leq \cdots \leq c_D,$$

and

$$b_0 > b_1 \geq b_2 \geq \cdots \geq b_{D-1} > b_D = 0.$$



Compute on  $p_{ii}^0(x, y) = |\Gamma_i(x)|$

$\circ$   
 $x$



- In this case  $x = y$ . Computing the pair  $(u, v)$  of vertices  $u \in \Gamma_{i-1}(x)$ ,  $v \in \Gamma_i(x)$  with  $uv \in R$  in different order, we have

$$p_{i-1 \ i-1}^0(x, y) b_{i-1} = p_i^0(x, y) c_i.$$

- Induction on  $i$  we have

$$p_{ii}^0 = \frac{b_0 b_1 \cdots b_{i-1}}{c_1 c_2 \cdots c_i},$$

independent of the choice of  $x$ .



# Distance matrices

The matrices we are concerned are square matrices with rows and columns indexed by the vertex set  $V\Gamma$ .

- For each  $i$  let  $A_i$  be a 01-matrix with entries

$$(A_i)_{xy} = \begin{cases} 1, & \text{if } d(x, y) = i; \\ 0, & \text{else.} \end{cases}$$

- $A_i$  is called  **$i$ -th distance matrix**, and  $A = A_1$  is also called the **adjacency matrix** of  $\Gamma$ .
- Note  $A_0 = I$  and  $A_{-1} = A_{D+1} = 0$ .



# Theorem

$$AA_i = b_{i-1}A_{i-1} + a_iA_i + c_{i+1}A_{i+1} \quad 0 \leq i \leq D,$$

where  $c_{D+1} := 1$ .

Proof.

$$(AA_i)_{xy} = \begin{cases} b_{i-1}, & \text{if } \partial(x, y) = i-1; \\ a_i, & \text{if } \partial(x, y) = i; \\ c_{i+1}, & \text{if } \partial(x, y) = i+1. \end{cases}$$



# Distance polynomials

In last page we show

$$AA_i = b_{i-1}A_{i-1} + a_iA_i + c_{i+1}A_{i+1} \quad 0 \leq i \leq D.$$

- Consider polynomials  $f_0(x) := 1$ ,  $f_1(x) := x$  and
- $f_i(x)$  is defined recursively using

$$xf_i(x) = b_{i-1}f_{i-1}(x) + a_if_i(x) + c_{i+1}f_{i+1}(x) \quad 2 \leq i \leq D.$$

- Note that  $A_i = f_i(A)$ ,  $f_{D+1}(A) = A_{D+1} = 0$  (using  $c_{D+1} := 1$ ), and  $f_i(x)$  has degree  $i$ .
- $A$  has  $D + 1$  distinct eigenvalues, which are the roots of  $f_{D+1}(x)$ .
- $f_0(x)$ ,  $f_1(x)$ ,  $\dots$ ,  $f_D(x)$  are called **distance polynomials**.
- Distance polynomials are orthogonal in a kind of inner product to be mentioned later.



# Proposition

Let  $\Gamma$  be a distance-regular graph with diameter  $D$ . If  $\Gamma$  has  $d+1$  distinct eigenvalues then  $d = D$ .

Proof.

Since the adjacency matrix  $A = A(\Gamma)$  is symmetric, the minimal polynomial of  $A(\Gamma)$  has degree  $d+1$ , where  $d \geq D$  has been proved in the first Section. As  $f_{D+1}(A) = 0$  in the last page, we also have  $D \geq d$ . Hence  $D = d$ . □



# Algebra

An **algebra**  $\mathcal{A}$  over  $\mathbb{C}$  is a vector space together with a multiplication  $\cdot : \mathcal{A} \times \mathcal{A} \rightarrow \mathcal{A}$  such that

(i)

$$\begin{cases} \alpha(a \cdot b) = (\alpha a) \cdot b = a \cdot (\alpha b), \\ a \cdot (b + c) = a \cdot b + a \cdot c, \\ (b + c) \cdot a = b \cdot a + c \cdot a, \end{cases} \quad (\text{The map } \cdot \text{ is bilinear});$$

(ii)  $a \cdot (b \cdot c) = (a \cdot b) \cdot c$  (The product is associative);

(iii)  $\exists 1 \in \mathcal{A}, a \cdot 1 = 1 \cdot a = a;$

where  $a, b, c \in \mathcal{A}$  and  $\alpha \in \mathbb{C}$ .

If there is no confusion, we will omit the multiplication symbol  $\cdot$ .



## Example

The set  $\text{Mat}_n(\mathbb{C})$  of all  $n \times n$  matrices over  $\mathbb{C}$  is an algebra under usual addition, scalar multiplication, and multiplication of matrices.





# Bose-Mesner Algebra

Let  $\Gamma$  be a distance-regular graph with diameter  $D$ .

- Consider the algebra  $\mathcal{M} = \langle A \rangle$  generated by the adjacency matrix  $A$  of  $\Gamma$ .
- Then the set

$$\{I = A_0, A = A_1, A_2 = f_2(A), \dots, A_D = f_D(A)\}$$

is a basis of  $\mathcal{M}$ .

- $\mathcal{M}$  is called the **Bose-Mesner Algebra** of  $\Gamma$ .



# Theorem

$$p_{ij}^k(x, y) = p_{ij}^k,$$

independent the choice  $x, y$ .

Proof.

Note that  $A_i A_j \in \mathcal{M}$ . Hence

$$A_i A_j = \sum_{k=0}^D p_{ij}^k A_k$$

for some constant  $p_{ij}^k \in \mathbb{R}$ . For  $x, y \in X$  with  $\partial(x, y) = k$  consider the  $xy$  entry of both sides of the above equality, we find  $p_{ij}^k(x, y) = p_{ij}^k$ , independent the choice  $x, y$ . □



# Strongly regular graphs

A distance-regular graph of diameter 2 is called a **strongly regular graph**.

Strongly regular graphs are studied earlier than the study of distance-regular graphs. They have relations to designs, finite geometries, graph eigenvalues, etc.



# Distance-transitive graphs

Let  $\Gamma = (V, E)$  be a graph.

- A map  $\phi : V \rightarrow V$  is an **automorphism** of  $\Gamma$  if  $\phi$  is bijection and

$$uv \in E \Leftrightarrow \phi(u)\phi(v) \in E.$$

- $\Gamma$  is **vertex-transitive (VT)** if for any two vertices  $u, v$  there exists an automorphism  $\phi$  of  $\Gamma$  such that  $\phi(u) = (v)$ .
- $\Gamma$  is **distance-transitive (DT)** if for any two pairs  $(u, v), (x, y)$  of vertices with  $\partial(u, v) = \partial(x, y)$ , there exists an automorphism  $\phi$  of  $\Gamma$  such that  $x = \phi(u)$  and  $y = \phi(v)$ .
- Note that a distance transitive graph is distance-regular.



## Recall the Hamming graphs $H(D, q)$

- Let  $X$  be a finite field of  $q$  elements,  $V = X^D$ , and

$$E = \{uv \mid u, v \in V \text{ differ in exact one coordinate}\}.$$

- $\Gamma = (V, E)$  is called the **Hamming graph**  $H(D, q)$ .
- Note that  $H(1, q) = K_q$ ,  $H(2, 2)$  is a square and  $H(3, 2)$  is a cube.
- For vertices  $x$  and  $y$  in  $H(D, q)$ ,  $\partial(x, y) = i$  if and only if  $x$  and  $y$  differ in exactly  $i$  coordinates.
- Hence  $H(D, q)$  is  $D(q-1)$ -regular with order  $q^D$  and diameter  $D$ .



# Theorem

The Hamming graph  $H(D, q)$  is distance-regular with intersection numbers

$$a_i = i(q - 2), \quad b_i = (D - i)(q - 1), \quad c_i = i.$$

Proof.

For two vertices  $x$  and  $y$  in  $H(D, q)$  at distance  $\partial(x, y) = i$ ,  $x$  and  $y$  differ at exactly  $i$  coordinates. Hence for  $z \in \Gamma_i(x) \cap \Gamma_1(y)$ , the unique coordinate that  $y$  and  $z$  differ is in one of the above  $i$  coordinates and at this coordinate  $z$  has one of the  $q - 2$  values which also differs to that of both  $x$  and  $y$ . This proves  $a_i = i(q - 2)$ . The proof of  $c_i$  is similar. Use  $a_i + b_i + c_i = b_0 = D(q - 1)$  to find  $b_i$ . □



# Coding Theory

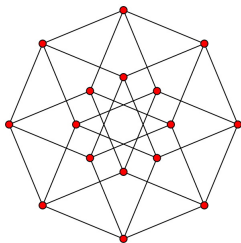
In  $H(D, 2)$  for each positive integer  $c$  find  $C \subseteq X$  with  $|C| = c$  to maximize

$$\text{Min}\{\partial(x, y) \mid x, y \in C, x \neq y\}.$$

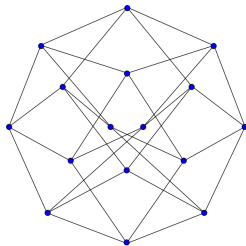
P. Delsarte studied coding theory in any distance-regular graphs in his 1973 PH. D. thesis "An Algebraic Approach to the Association Schemes of Coding Theory".



# Two cospectral graphs



The 4-cube.



The Hoffman graph.

Copy from [http://en.wikipedia.org/wiki/Hoffman\\_graph](http://en.wikipedia.org/wiki/Hoffman_graph)

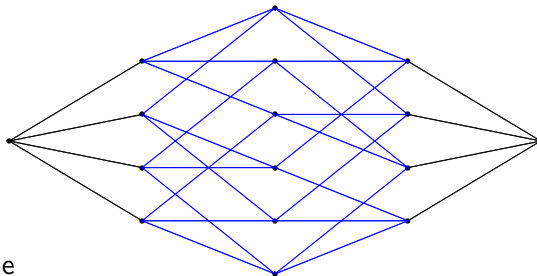
In the Hoffman graph, every vertex has distance at most 3 to the southwestern vertex. So Hoffman graph is not distance-transitive.



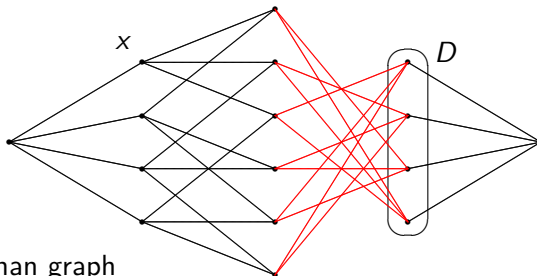
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# Another drawing of Hoffman graph



The 4-cube



The Hoffman graph



## Example (Johnson graphs $J(n, D)$ , $2D \leq n$ )

- Set  $[n] = \{1, 2, \dots, n\}$ ,

$$V = \binom{[n]}{D}, \quad E = \{uv \mid u, v \in V, |u \cap v| = D - 1\}.$$

- Then  $\Gamma = (V, E)$  is called the **Johnson graph**  $J(n, D)$ .
- $J(n, D)$  is a distance-regular graph with diameter  $D$  (Exercise).



# Design theory

- A  $t$ -( $n, D, \lambda$ ) **design** is a pair  $(P, \mathcal{B})$ , where  $P = [n]$  and  $\mathcal{B}$  is a family of  $D$ -subsets of  $P$  such that any  $t$  elements in  $P$  are contained in exactly  $\lambda$   $D$ -subsets in  $\mathcal{B}$ .
- Note that with  $P = [n]$  and  $\mathcal{B} = \binom{[n]}{D}$ , we have a  $t$ -( $n, D, \binom{n-t}{D-t}$ ) design.
- Any  $t$ -( $n, D, \lambda$ )  $(P, \mathcal{B})$  for  $2D \leq n$  can be viewed as a subgraph of  $J(n, D)$  induced on  $\mathcal{B}$  with some nice properties to be specified (a design on  $J(n, D)$ ).

## Question

What is the related design on arbitrary distance-regular graph. (P. Delsarte's thesis)



# Grassmann graph $J_q(n, D)$ , $2D \leq n$

- Let  $F_q$  be a finite field of  $q$  elements and  $V = \binom{F_q^n}{D}$  be the set of  $D$ -dimensional subspaces of  $F_q^n$ , and

$$E = \{uv \mid u, v \in V, \dim(u \cap v) = D - 1\}.$$

- Then  $\Gamma = (V, E)$  is called the **Grassmann graph**  $J_q(n, D)$ , or the  **$q$ -analog of Johnson graph**.
- $J_q(n, D)$  is a distance-regular graph with diameter  $D$  (Exercise).



# A class of non-VT DRG

Invent. math. 162, 189–193 (2005)  
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*Inventiones  
mathematicae*

## A new family of distance-regular graphs with unbounded diameter

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**Abstract.** We construct distance-regular graphs with the same – classical – parameters as the Grassmann graphs on the  $e$ -dimensional subspaces of a  $(2e + 1)$ -dimensional space over an arbitrary finite field. This provides the first known family of non-vertex-transitive distance-regular graphs with unbounded diameter.



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