

## Solutions

2021年4月22日 上午 08:53

1.

$i$	$r_i$	$\delta_i$	$x_i$	$y_i$
-1	184		1	0
0	57	3	0	1
1	13	4	1	-3
2	5	2	-4	13
3	3	1	9	-29
4	2	1	-13	42
5	1		22	-91
			$x$	$y$

2. By the CRT formula,

$$\begin{aligned} x &= \left[ 2 \cdot (7 \cdot 13) ((7 \cdot 13)^{-1} \bmod 6) \right. \\ &\quad + 4 \cdot (6 \cdot 13) ((6 \cdot 13)^{-1} \bmod 7) \\ &\quad \left. + 1 \cdot (6 \cdot 7) ((6 \cdot 7)^{-1} \bmod 13) \right] \bmod 6 \cdot 7 \cdot 13 \\ &= 326 \end{aligned}$$

3. By the Euclidean algorithm:

$$2x^2 + x + 5 \quad \text{or} \quad a \cdot (2x^2 + x + 5), a \in \mathbb{Z}_7^*$$

$$\rightarrow \begin{cases} 4x^2 + 2x + 3 & a=2 \\ 6x^2 + 3x + 1 & a=3 \\ x^2 + 4x + 6 & a=4 \\ 3x^2 + 5x + 4 & a=5 \\ 5x^2 + 6x + 2 & a=6 \end{cases}$$

4.

$$a = 37_H = 0011\ 0111$$

$$b = 16 = 1010\ 0110$$

$i$	$b_i$	$f(\text{shifted})$	$f(\bmod g(x))$
0		0000 0000	0000 0000
7	1	0011 0111	0011 0111
6	0	0110 1110	0110 1110
5	1	1110 1011	1110 1011
4	0	1110 1010	1110 1010
3	0	1000 1000	1000 1000

T	V	110 0110 1100 110
3	0	1 (00) 1010 1000 0001
2	1	1001 0101 0010 1110
1	1	0110 1011 0110 1011
0	0	1101 0110 1101 0110

D 6

5. Add Roundkey  $\rightarrow$  substitution

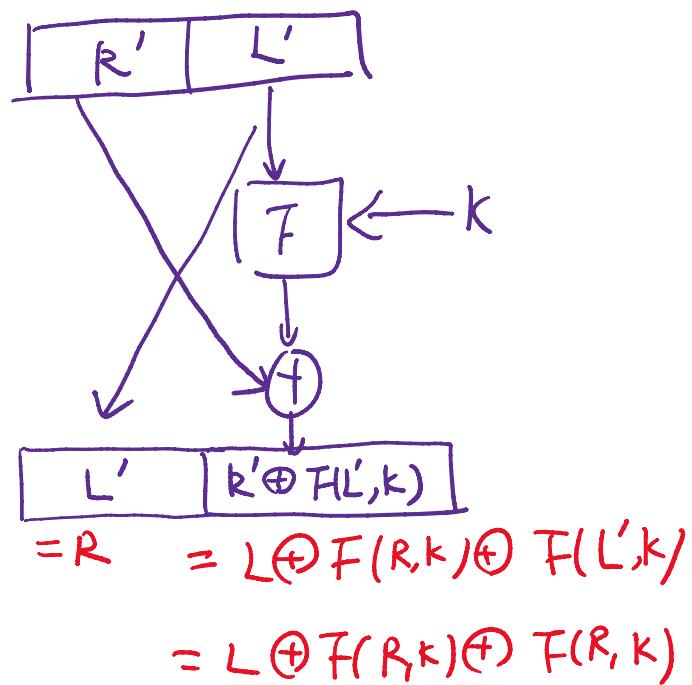
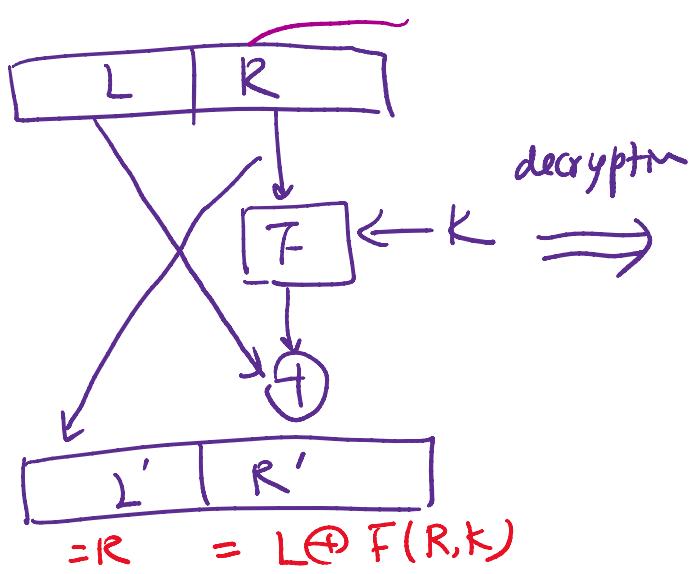
Sub Byte  $\rightarrow$  substitution

Shift Row  $\rightarrow$  permutation

Mix Column  $\rightarrow$  substitution

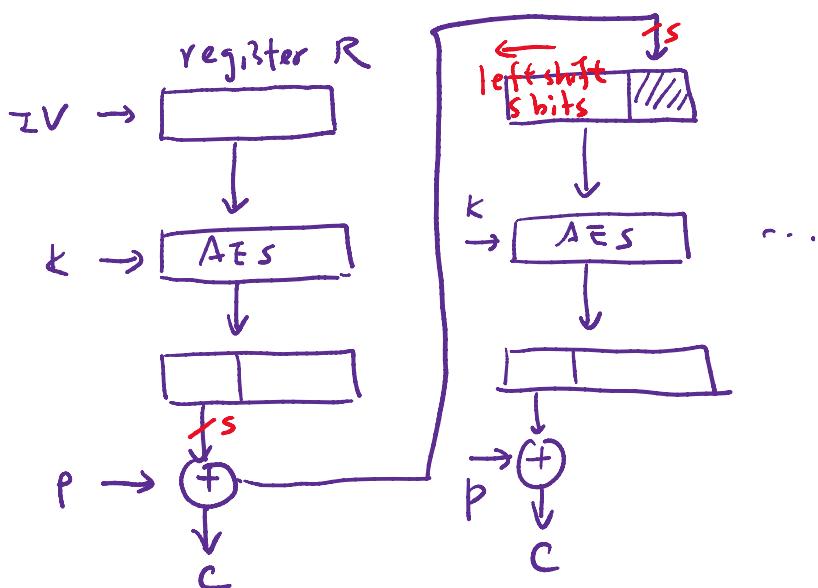
Explanation can be found in class slides

6.

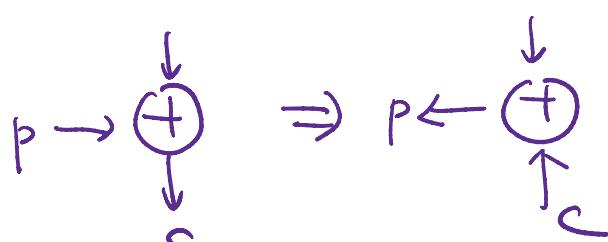


Thus, no matter what  $F$  is,  $L$  can be recovered.

7. (a)



Decryption: the same except that



(b) Since an erroneous ciphertext of  $s$  bytes would stay in register R for  $\lceil \frac{16}{s} \rceil$  blocks in decryption, the next  $\lceil \frac{16}{s} \rceil$  blocks of decryption are wrong.

So,  $\underbrace{P_2, P_3, P_4, P_5, P_6, P_7, P_8, P_9, P_{10}, P_{11}}_{\text{and}}$

$\underbrace{P_{83}, P_{84}, P_{85}, P_{86}, P_{87}, P_{88}, P_{89}}$  are decrypted incorrectly.

8. (a)  $\varphi(n) = (11-1)(13-1) = 120$

$$d = e^{-1} \pmod{\varphi(n)} = 7^{-1} \pmod{120} = 103$$

(b) (i) in "mod 11" space.

$$C' = C \pmod{11} = 8$$

$$d' = d \pmod{11-1} = 3$$

$$m' = (C')^{d'} \pmod{11} = 8^3 \pmod{11} = 6$$

(ii) in "mod 13" space

$$C'' = C \pmod{13} = 8$$

$$d'' = d \pmod{13-1} = 7$$

$$m'' = (C'')^{d''} \pmod{13} = 8^7 \pmod{13} = 5$$

Then, by CRT

$$m = (6 \cdot 13 \cdot (13^{-1} \pmod{11}) + 5 \cdot 11 (11^{-1} \pmod{13})) \pmod{143}$$

$$= (6 \cdot 13 \cdot 6 + 5 \cdot 11 \cdot 6) \pmod{143}$$

$$= 83$$

(c) Case I:  $M \perp n$ ,

$$C^d \pmod{n} = M^{ed} \pmod{n} = M^{k \cdot \varphi(n)} \cdot M \pmod{n}$$

$$= ((M^{\varphi(n)} \pmod{n})^k \cdot M) \pmod{n} \quad \text{By Euler's theorem:}$$

$$= (1^k \cdot M) \pmod{n}$$

$$\text{since } M^{\varphi(n)} \pmod{n} = 1$$

since  $M \perp n$ .

Case II:  $\gcd(M, n) \neq 1$

Assume that  $M = aP$ ,  $0 \leq a < p-1$

Let  $x = C^d \pmod{n}$ .

We have

$$\dots \rightarrow m \pmod{n} = n = M \pmod{a} = r_1$$

We have

$$(i) x \bmod p = 0 = M \bmod p = r_1$$

$$(ii) x \bmod q = M^{k(p-1)(q-1)+1} \bmod q$$

$$= (M^{q-1} \bmod q)^{k(p-1)} \cdot M \bmod q \quad \text{since } M \perp q.$$

$$= 1^{k(p-1)} \cdot M \bmod q \quad \begin{matrix} \swarrow \\ M^{q-1} \bmod q = 1 \end{matrix}$$

$$= M \bmod q = r_2 \quad \text{by Fermat's little theorem}$$

By CRT, there is a unique solution  $x_0$ ,  $0 \leq x_0 \leq n-1$

for the equations:  $x \bmod p = r_1 = 0$ ,  $x \bmod q = r_2 = M \bmod q$ .

Since  $M \bmod p = r_1$  and  $M \bmod q = r_2$ ,  $M$  is a solution in  $(0, n-1]$

Thus,  $x = M$ .