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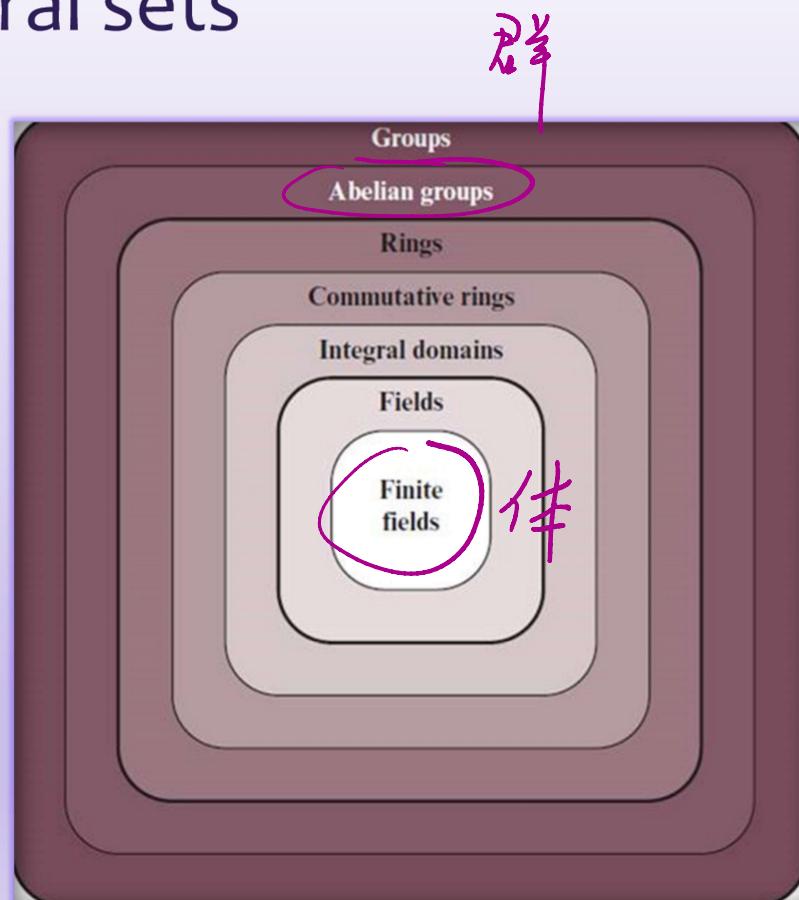
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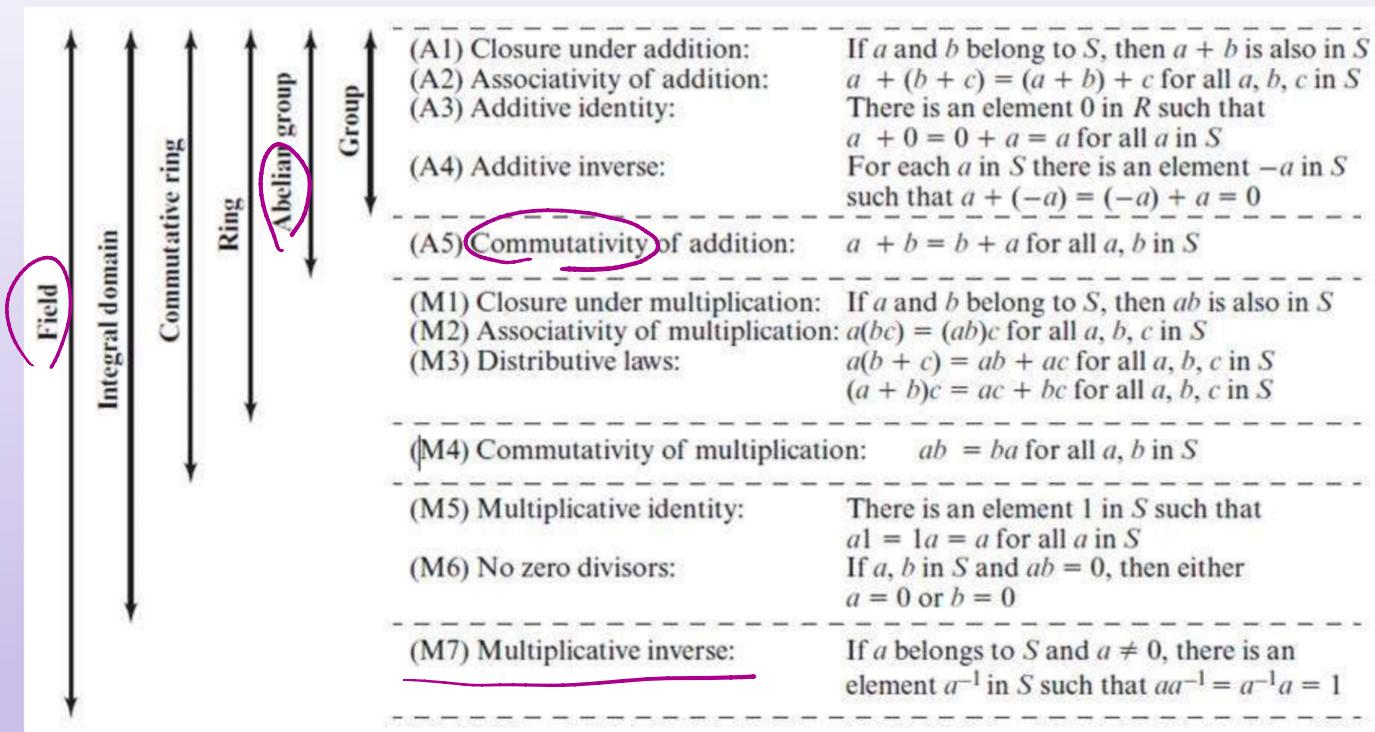


Chapter 5

Finite Fields

Structural sets





Abelian Group

- (G, \circ) is a group, where G is a set of elements and the binary operator \circ has the following properties:
 - Closure
 - $a \circ b$, for $a, b \text{ in } G$
 - Associativity
 - $a \circ (b \circ c) = (a \circ b) \circ c$, for $a, b, c \text{ in } G$
 - Identity
 - There is an element e in G such that $a \circ e = e \circ a = a$, for $a \text{ in } G$
 - Inverse
 - For each a in G , there is an element a^{-1} in G such that $a \circ a^{-1} = a^{-1} \circ a = e$
 - **Commutative:** $a \circ b = b \circ a$, for $a, b \text{ in } G$

Examples

- Additive groups
 - $(R, +)$
 - $(\mathbb{Z}, +)$
 - $(\mathbb{Z}_n, +)$, $\mathbb{Z}_n = \{0, 1, 2, \dots, n-1\}$, “+” is mod n
- Multiplicative groups
 - $(R - \{0\}, \times)$, $\underset{=}^{(?, \times)} \times \rightarrow ? \oplus$
 - $(Q - \{0\}, \times)$
 - (\mathbb{Z}_n^*, \times) , $\mathbb{Z}_n^* = \{a: 1 \leq a < n, \gcd(a, n) = 1\}$, “ \times ” is mod n
- $N_n = \{\pi: \pi \text{ is a permutation over } \{1, 2, \dots, n\}\}$
 - Not abelian $(1\ 2\ 3)(3\ 2\ 1)(2\ 1\ 3) \cdots$
- The operator is omitted if no misunderstanding occurs

Cyclic group

- G is **cyclic** if there is a generator $g \in G$ such that for any $a \in G$, $a = g^k$ for some k
 - g spans all elements of G , that is, $G = \{g^k | k \geq 0\}$
- Notation
 - $g^k = g \circ g \circ \dots \circ g$ (k times)
 - $a^0 = e$: identity
 - $a^{-k} = (a^{-1})^k$
- \mathbb{Z}_n^* is a cyclic group with generators 3 and 5
 - $3^0 = 1, 3^1 = 3, 3^2 = 2, 3^3 = 6, 3^4 = 4, 3^5 = 5$
 - $5^0 = 1, 5^1 = 5, 5^2 = 4, 5^3 = 6, 5^4 = 2, 5^5 = 3$

$$\mathbb{Z}_n^* \rightarrow \text{cyclic}$$

iff $n = 2, 4, p^d, 2p^q$

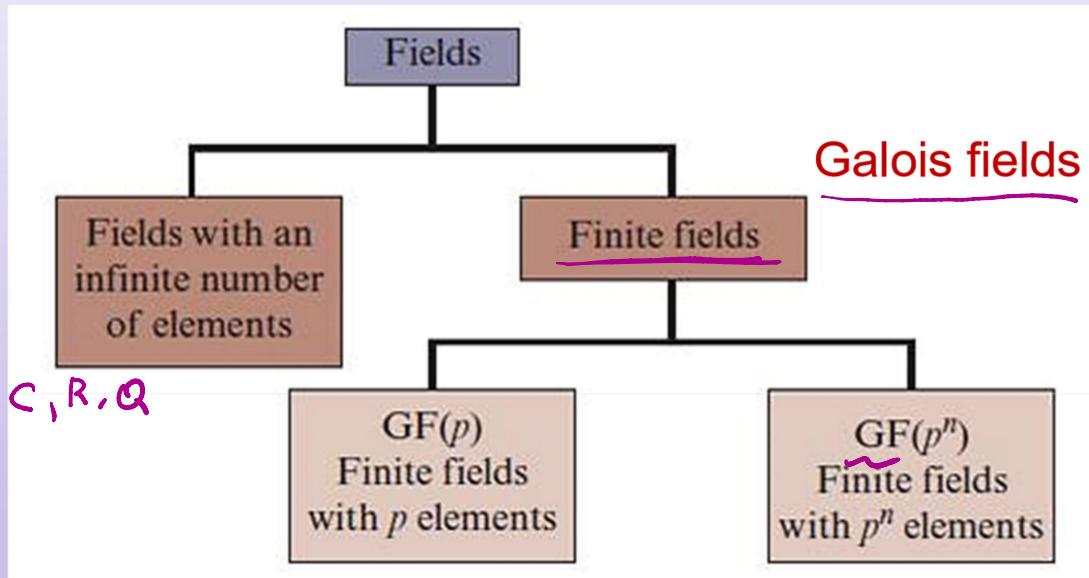
Field: $\{F, +, \times\}$

- $\{F, +\}$ is an additive abelian group
- $\{F - \{0\}, \times\}$ is a multiplicative abelian group
- Distributive laws
 - $a \times (b + c) = a \times b + a \times c$, for a, b, c in F
 - $(a + b) \times c = a \times c + b \times c$, for a, b, c in F
- 0: the identity for $+$
- 1: the identity for \times
- $-a$: the additive inverse of a
- a^{-1} : the multiplicative inverse of a
- $a - b = a + (-b)$
- $a/b = a \times b^{-1}$
- We can do 4 operators (+, -, \times , $/$) over a field

Field: examples

- $\{R, +, \times\}$, where R is the set of reals
- $\{Q, +, \times\}$, where Q is the set of rationals
- $\{Z_p, +, \times\}$, where p is prime and operators are mod p
 - $-a = p - a$ $Z_p = \{0, 1, 2, 3, \dots, p-1\}$
 - $a^{-1}, 1 \leq a < p$
 - Use extended Euclidean algorithm to compute integral (x, y) for
 $xa + yp = 1$
 - $a^{-1} = x \text{ mod } p$
 - Additive identity: 0
 - Multiplicative identity: 1

Field: types



Finite Field: $GF(p^n)$

- Evariste Galois (1811-1832) first studied finite fields
- Finite fields play crucial role in AES and many cryptosystems
- Every finite field F must have p^n elements for some prime p and $n \geq 1$
- For every prime p and $n \geq 1$, there is a finite field of p^n elements
- F of p^n elements may have different forms. Nevertheless, they are all isomorphic
 - Thus, $GF(p^n)$ is the finite field of p^n elements

Finite Field: $GF(p)$, $n = 1$

Finite Field: $GF(p)$, $n = 1$

- $GF(p) = \{Z_p, +, \times\}$, where + and \times under “mod p”
 - The finite field of p elements $Z_p = \{0, 1, 2, \dots, p-1\}$, $\Rightarrow p$ elements
- Example $\{0, 1\}$
 - $GF(2) = \{Z_2, \text{xor}, \text{and}\}$: Boolean algebra
 - $GF(7) = Z_7$

+	0	1	2	3	4	5	6
0	0	1	2	3	4	5	6
1	1	2	3	4	5	6	0
2	2	3	4	5	6	0	1
3	3	4	5	6	0	1	2
4	4	5	6	0	1	2	3
5	5	6	0	1	2	3	4
6	6	0	1	2	3	4	5

\times	0	1	2	3	4	5	6
0	0	0	0	0	0	0	0
1	0	1	2	3	4	5	6
2	0	2	4	6	1	3	5
3	0	3	6	2	5	1	4
4	0	4	1	5	2	6	3
5	0	5	3	1	6	4	2
6	0	6	5	4	3	2	1

Ordinary polynomials: arithmetic

$$\begin{array}{r} x^3 + x^2 \quad + 2 \\ + \quad (x^2 - x + 1) \\ \hline x^3 + 2x^2 - x + 3 \end{array}$$

(a) Addition

$$\begin{array}{r} x^3 + x^2 \quad + 2 \\ - \quad (x^2 - x + 1) \\ \hline x^3 \quad + x + 1 \end{array}$$

(b) Subtraction

$$\begin{array}{r} x^3 + x^2 \quad + 2 \\ \times \quad (x^2 - x + 1) \\ \hline x^3 + x^2 \quad + 2 \\ -x^4 - x^3 \quad - 2x \\ \hline x^5 + x^4 \quad + 2x^2 \\ \hline x^5 \quad + 3x^2 - 2x + 2 \end{array}$$

(c) Multiplication

$$\begin{array}{r} x + 2 \\ \hline x^2 - x + 1 \sqrt{x^3 + x^2 \quad + 2} \\ \hline x^3 - x^2 + x \\ \hline 2x^2 - x + 2 \\ \hline 2x^2 - 2x + 2 \\ \hline x \end{array}$$

(d) Division

Polynomials over $GF(p)$

- A polynomial of degree $n-1$ over $GF(p)$ is of form over $GF(3)$

$$\checkmark f(x) = a_{n-1}x^{n-1} + a_{n-2}x^{n-2} + \cdots + a_1x + a_0$$

$\frac{2x^3+x^2+x+1}{2x+1} \quad) \overline{x^4+x^3+0+0+1}$

$\frac{x^4+2x^3}{x^4+2x^3}$

$\frac{0+2x^3+0+0+1}{2x^3+x^2}$

$\frac{2x^2+0+1}{2x+1}$

$\frac{2x^2+x}{2x+1}$

$\frac{2x+1}{0}$

where each $a_i \in GF(p), 0 \leq i < n$

- Polynomials over $GF(2)$

- $0, 1, x, x + 1, x^2, x^2 + 1, x^2 + x, x^2 + x + 1, x^3, \dots$

- Addition over $GF(2)$

$$(x^2 + 1) + (x^2 + x) = 2x^2 + x + 1 = x + 1$$

- Multiplication over $GF(2)$

$$(x^2 + 1) \times (x^2 + 1) = x^4 + x^2 + x^2 + 1 = x^4 + 1$$

Polynomials over $GF(p)$ mod $m(x)$

- $f(x) \text{ mod } \underline{\underline{m(x)}} = r(x)$, where $\underline{\deg(r(x))} < \underline{\deg(m(x))}$
 - $f(x) = q(x)m(x) + r(x)$
- Example over $GF(2)$
 - $f(x) = x^4 + x^2 + 1$
 - $m(x) = x^3 + x + 1$
 - $f(x) = x \cdot m(x) + (x + 1)$
- Simple way of computing $f(x) \text{ mod } m(x)$
 - Substitute $m(x) = 0$ to $f(x)$ and get $r(x)$
 - $m(x) = 0 \rightarrow x^3 = x + 1$
 - $f(x) \text{ mod } m(x) = x(x + 1) + x^2 + 1 = x + 1$

Irreducible polynomial over $GF(p)$

- $m(x)$ is **irreducible** if $m(x)$ cannot be factored into a product of two polynomials over $GF(p)$ of degree ≥ 1
- Example
 - $x^3 + x + 1$ is irreducible over $GF(2)$
 - $x^3 + x^2 + x + 1 = (x + 1)^3$ is reducible over $GF(2)$
- Let $m(x)$ be degree- n irreducible polynomial over $GF(p)$
 - $f(x)$ over $GF(p)$ with $\deg(f(x)) \leq n - 1$
 - $\gcd(f(x), m(x)) = 1$ for $f(x) \neq 0$
 - $f^{-1}(x) \bmod m(x)$ exists for $f(x) \neq 0$
 - Use extended Euclidean algorithm to find $(a(x), b(x))$ for $a(x)f(x) + b(x)m(x) = 1 = \gcd(f(x), m(x))$
 - $f^{-1}(x) \bmod m(x) = a(x) \bmod m(x)$

Finite field: $GF(\underline{p^n})/m(x), n \geq 2$

Finite field: $GF(p^n)/\underline{m(x)}$, $n \geq 2$

- $GF(p^n)$ is the set of polynomials over $GF(p)$ of degree at most $n - 1$
- $m(x)$ is an irreducible degree- n monic polynomial over $GF(p)$
- Coefficient operations are over $GF(p)$
- Multiplicative operations are "mod $m(x)$ "
- Additive identity: 0
- Additive inverse: $-f(x)$
- Multiplicative identity: 1
- Multiplicative inverse: $f(x)^{-1} \text{ mod } m(x)$
- Closure, inverse, associative, commutative, and distributive rules are satisfied

Finite field $GF(2^3)/\underline{x^3 + x + 1}$

- $S = \{0, 1, x, x + 1, x^2, x^2 + 1, x^2 + x, x^2 + x + 1\}$
- $m(x) = x^3 + x + 1$ irreducible over $GF(2)$

- $S = \{0, 1, x, x+1, x^2, x^2+1, x^2+x, x^2+x+1\}$

- $m(x) = \underline{x^3 + x + 1}$: irreducible over $GF(2)$

- Example

- $(x^2 + 1) + (x + 1) = x^2 + x$

- $(x^2 + 1) \times (x + 1) \bmod m(x) = x^2$

- $-(x^2 + x) = x^2 + x$

$$x^3 + x^2 + x^2 + x = \underline{x^3 + x}$$

- $(x^2 + x)^{-1} \bmod m(x) = \underline{x + 1}$

- $(x^2)^{-1} \bmod m(x) = x^2 + x + 1$

- $GF(2^3)/\underline{x^3 + x + 1}$ and $GF(2^3)/\underline{x^3 + 1}$ are isomorphic since $x^3 + x^2 + 1$ is also irreducible over $GF(2)$

$\hat{x} \times \textcolor{red}{1}$	$+$	000	001	010	011	100	101	110	111
000	0	1	x	$x+1$	x^2	x^2+1	x^2+x	x^2+x+1	
001	1	0	$x+1$	x	x^2+1	x^2	x^2+x+1	x^2+x	
010	x	x	$x+1$	0	x^2+x	x^2+x+1	x^2	x^2+1	
011	$x+1$	$x+1$	x	1	x^2+x+1	x^2+x	x^2+1	x^2	
100	x^2	x^2	x^2+x	x^2+x+1	0	1	x	$x+1$	
101	x^2+1	x^2+1	x^2+x+1	x^2+x	1	0	$x+1$	x	
110	x^2+x	x^2+x+1	x^2	x^2+x+1	x	$x+1$	0	1	

100	x^2	x^2	$x^2 + 1$	$x^2 + x$	$x^2 + x + 1$	0	1	x	$x + 1$
101	$x^2 + 1$	$x^2 + 1$	x^2	$x^2 + x + 1$	$x^2 + x$	1	0	$x + 1$	x
110	$x^2 + x$	$x^2 + x$	$x^2 + x + 1$	x^2	$x^2 + 1$	x	$x + 1$	0	1
111	$x^2 + x + 1$	$x^2 + x + 1$	$x^2 + x$	$x^2 + 1$	x^2	$x + 1$	x	1	0

		000	001	010	011	100	101	110	111
	\times	0	1	x	$x + 1$	x^2	$x^2 + 1$	$x^2 + x$	$x^2 + x + 1$
000	0	0	0	0	0	0	0	0	0
001	1	0	1	x	$x + 1$	x^2	$x^2 + 1$	$x^2 + x$	$x^2 + x + 1$
010	x	0	x	x^2	$x^2 + x$	$x + 1$	1	$x^2 + x + 1$	$x^2 + 1$
011	$x + 1$	0	$x + 1$	$x^2 + x$	$x^2 + 1$	$x^2 + x + 1$	x^2	1	x
100	x^2	0	x^2	$x + 1$	$x^2 + x + 1$	$x^2 + x$	x	$x^2 + 1$	1
101	$x^2 + 1$	0	$x^2 + 1$	1	x^2	x	$x^2 + x + 1$	$x + 1$	$x^2 + x$
110	$x^2 + x$	0	$x^2 + x$	$x^2 + x + 1$	1	$x^2 + 1$	$x + 1$	x	x^2
111	$x^2 + x + 1$	0	$x^2 + x + 1$	$x^2 + 1$	x	1	$x^2 + x$	x^2	$x + 1$

Finite field $GF(3^2)/\underline{\underline{(\underline{x^2+1})}}$

- $S = \{0, 1, 2, x, \underline{x + 1}, \underline{x + 2}, \underline{2x}, \underline{2x + 1}, \underline{2x + 2}\}$
- $m(x) = \underline{x^2} + 1$: irreducible over $GF(3)$
- Example
 - $(x + 1) + (x + 2) = 2x$
 - $(x + 1) \times (x + 2) \bmod m(x) = 1$
 - $-(x + 2) = 2x + 1$
 - $(x + 1)^{-1} \bmod m(x) = x + 2$
 - $(2x)^{-1} \bmod m(x) = x$

$+$	0	1	2	x	$x+1$	$x+2$	$2x$	$2x+1$	$2x+2$
0	0	1	2	x	$x+1$	$x+2$	$2x$	$2x+1$	$2x+2$
1	1	2	0	$x+1$	$x+2$	x	$2x+1$	$2x+2$	$2x$
2	2	0	1	$x+2$	x	$x+1$	$2x+2$	$2x$	$2x+1$
x	x	$x+1$	$x+2$	$2x$	$2x+1$	$2x+2$	0	1	2
$x+1$	$x+1$	$x+2$	x	$2x+1$	$2x+2$	$2x$	1	2	0
$x+2$	$x+2$	x	$x+1$	$2x+2$	$2x$	$2x+1$	2	0	1
$2x$	$2x$	$2x+1$	$2x+2$	0	1	2	x	$x+1$	$x+2$
$2x+1$	$2x+1$	$2x+2$	$2x$	1	2	0	$x+1$	$x+2$	x
$2x+2$	$2x+2$	$2x$	$2x+1$	2	0	1	$x+2$	x	$x+1$

\times	0	1	2	x	$x+1$	$x+2$	$2x$	$2x+1$	$2x+2$
0	0	0	0	0	0	0	0	0	0
1	0	1	2	x	$x+1$	$x+2$	$2x$	$2x+1$	$2x+2$
2	0	2	1	$2x$	$2x+2$	$2x+1$	x	$x+2$	$x+1$
x	0	x	$2x$	2	$x+2$	$2x+2$	1	$x+1$	$2x+1$
$x+1$	0	$x+1$	$2x+2$	$x+2$	2	1	$2x+1$	2	x
$x+2$	0	$x+2$	$2x+1$	$2x+2$	1	x	$x+1$	$2x$	2
$2x$	0	$2x$	x	1	$2x+1$	$x+1$	2	$2x+2$	$x+2$
$2x+1$	0	$2x+1$	$x+2$	$x+1$	2	$2x$	$2x+2$	x	1
$2x+2$	0	$2x+2$	$x+1$	$2x+1$	x	2	$x+2$	1	$2x$

$GF(2^n)/m(x)$: computaton

- Represent polynomial $f(x) = \sum_{i=0}^{n-1} b_i x^i$ as binary string $b_{n-1}b_{n-2}\cdots b_1b_0$
- Addition: bitwise XOR (no need to carry)

		000	001	010	011	100	101	110	111	
		+	0	1	2	3	4	5	6	7
000	0	0	1	2	3	4	5	6	7	
001	1	1	0	3	2	5	4	7	6	
010	2	2	3	0	1	6	7	4	5	
011	3	3	2	1	0	7	6	5	4	
100	4	4	5	6	7	0	1	2	3	
101	5	5	4	7	6	1	0	3	2	
110	6	6	7	4	5	2	3	0	1	
111	7	7	6	5	4	3	2	1	0	

- Multiplication: Shift-XOR

- long-hand multiplication, but from high bit to low bit

- Example: $\underline{1001} \times \underline{1101}$ multiplier

- Let $1101 = b_3 b_2 b_1 b_0$

- Initiation: $f = \underline{\underline{0000}} \underline{\underline{0000}}$

- $b_3 = 1 \rightarrow f = \text{shift-left}(f) \oplus 1001 = 0000 \ 1001$

- $b_2 = 1 \rightarrow f = \text{shift-left}(f) \oplus 1001 = 0001 \ 1011$

- $b_1 = 0 \rightarrow f = \text{shift-left}(f) = 0011 \ 0110$

- $b_0 = 1 \rightarrow f = \text{shift-left}(f) \oplus 1001 = 0110 \ 0101$

$$\begin{array}{r} 1 \ 0 \ 0 \ 1 \text{ (multiplicand)} \\ \times 1 \ 1 \ 0 \ 1 \text{ (multiplier)} \\ \hline 1 \ 0 \ 0 \ 1 \\ 1 \ 0 \ 0 \ 1 \\ 0 \ 0 \ 0 \ 0 \\ 1 \ 0 \ 0 \ 1 \\ \hline 1 \ 1 \ 0 \ 0 \ 1 \ 0 \ 1 \end{array}$$

- Modular multiplication

- Shift-XOR-Mod: like Shift-XOR, do modulo whenever necessary
- $m(x) = x^8 + x^4 + x^3 + x + 1$ (binary: 1 0001 1011)
- $a \times b = 3F \times 86 = 0011\ 1111 \times 1000\ 0110$

i	b_i	f: shift-XOR	f: mod g(x) \rightarrow bitwise XOR
Initial			0000 0000
7	1	0011 1111	0011 1111
6	0	0111 1110	0111 1110
5	0	1111 1100	1111 1100
4	0	1 1111 1000	1110 0011
3	0	1 1100 0110	1101 1101
2	1	1 1000 0101	1001 1110
1	1	1 0000 0011	0001 1000
0	0	0011 0000	0011 0000

Computation: table lookup

- $GF(2^8) / x^8 + x^4 + x^3 + x + 1$
- Build a table for $a(x)b(x) \bmod x^8 + x^4 + x^3 + x + 1$
- Table size: $2^8 \times 2^8 \times 2^8$ bits = 2^{16} bytes = 64K bytes

(p^{n₁n₂})

Field: $GF(\underline{p}^{n_1 \times n_2}) / \underline{m_1}(x), \underline{m_2}(y)$

FIELD: $GF(\underline{p^{n_1}})/\underline{m_1(x)}, \underline{m_2(y)}$

- Consider degree- (n_2-1) polynomials of y over $GF(p^{n_1})/m_1(x)$
- Example
 - $p = 2, n_1 = 3, n_2 = 4$
 - $GF(p^{n_1})/m_1(x) = GF(2^3)/x^3 + x + 1$
 - A polynomial of the field is like : $(x+1)\underline{y^3} + (x^2)\underline{y^2} + \underline{1}$
- Let $m_2(y)$ be an irreducible degree- n_2 polynomial with coefficients over $GF(p^{n_1})/m_1(x)$
- $GF(p^{n_1 \times n_2})/m_1(x), m_2(y)$
 - The element set consists of all degree- (n_2-1) polynomials (of y) with coefficients over $GF(p^{n_1})/m_1(x)$
 - Coefficients are operated over $GF(p^{n_1})/m_1(x)$

Example: $GF(2^{3 \times 4})/m_1(x), m_2(y)$

- $GF(2^{3 \times 4})/x^3 + x + 1, y^4 + (x^2 + 1)y^2 + (x + 1)$
- $m_2(y) = y^4 + (x^2 + 1)y^2 + (x + 1)$ is irreducible over field

- $GF(2^{5+4})/x^5 + x + 1, y^4 + (x^2 + 1)y^2 + (x + 1)$
- $m_2(y) = \underbrace{y^4 + (x^2 + 1)y^2 + (x + 1)}_{\text{is irreducible over field } GF(2^3)/x^3 + x + 1} = 0 \Rightarrow y^4 = (x^2 + 1)y^2 + (x + 1)$
- Multiplication

$$\begin{aligned}
 & [(x + 1)y^3 + xy^2 + 1] \times [y + (x^2 + 1)] \bmod m_2(y) \\
 &= (x + 1)y^4 + [(x + 1)(x^2 + 1) + x]y^3 + [x(x^2 + 1)]y^2 \\
 &\quad + y + (x^2 + 1) \bmod m_2(y) \\
 &= (x + 1)[(x^2 + 1)y^2 + (x + 1)] + (x^2 + x)y^3 + y^2 \\
 &\quad + y + (x^2 + 1) \bmod m_2(y) \\
 &= (x^2 + x)y^3 + (x^2 + 1)y^2 + y + y \cdot \frac{(x+1)(x^2+1)+1}{x+1} \\
 &= \underline{\underline{x+1}}
 \end{aligned}$$