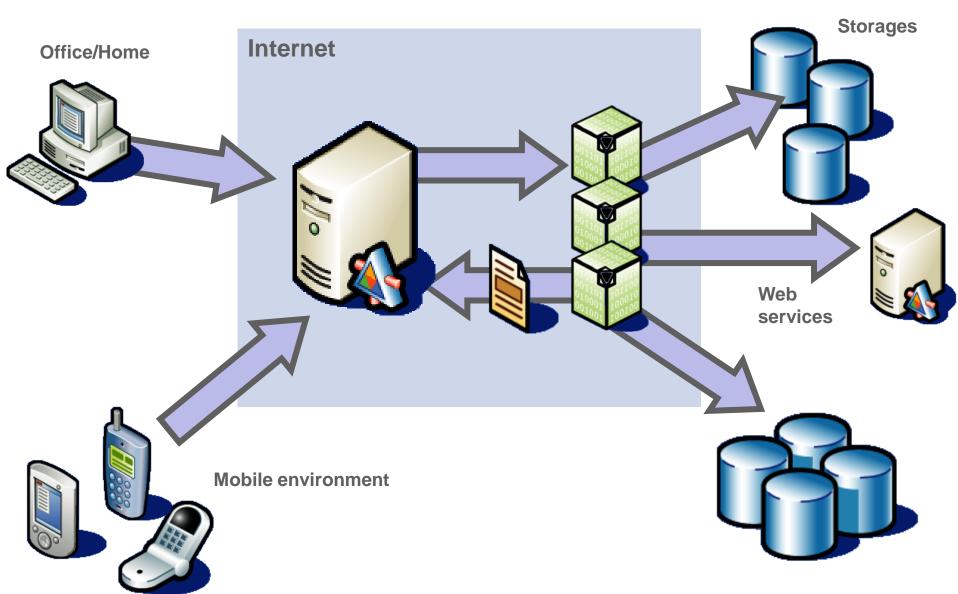


Text Books and Grading

- Major text: A. Silberschatz, P.B. Galvin, and G. Gagne, "Operating System Concepts"
- Grading
 - Mid-term: 30%
 - Final term: 30%
 - 4 Projects: 30%
 - Your grade will be limited below C+ if you do not submit any projects.
 - Source code cheating: grades will be deducted.
 - No co-work!!!!!
 - Attendance: 10%

The Internet of Things(IoT)



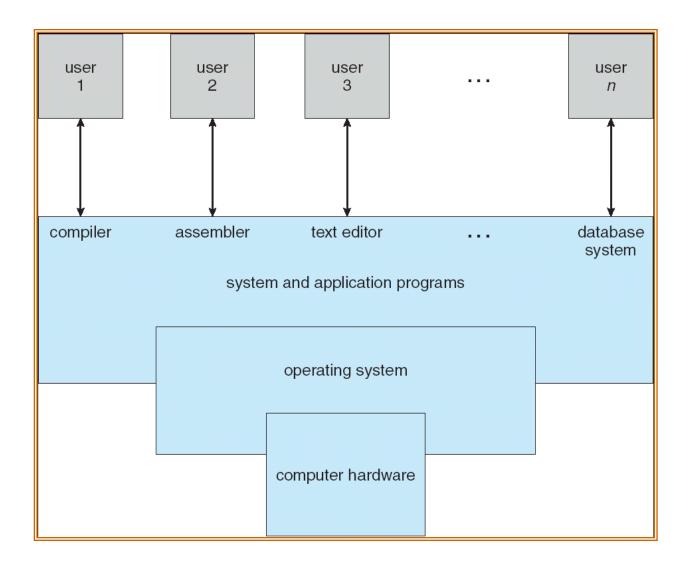
Modern Operating Systems



- Operating systems for servers and desktops
- Mobile operating systems
- Embedded systems

- Operating systems for web servers and multimedia systems

Four Components of a Computer System



Computer System Structure

- Computer system can be divided into four components
 - Hardware provides basic computing resources
 - CPU, memory, I/O devices
 - Operating system
 - Controls and coordinates use of hardware among various applications and users
 - System and application programs define the ways in which the system resources are used to solve the computing problems of the users
 - Word processors, compilers, web browsers, database systems, video games
 - Users
 - People, machines, other computers

What if we didn't have an Operating System?

- Source Code ⇒ Compiler ⇒ Object Code ⇒ Hardware
- How do you get object code onto the hardware?
- How do you print out the answer?

What is Operating System?

- Operating system goals
 - Execute user programs
 - Provide libraries of standard services and make solving user problems easier.
 - Make the computer system convenient to use.
 - Use the computer hardware in an efficient manner.

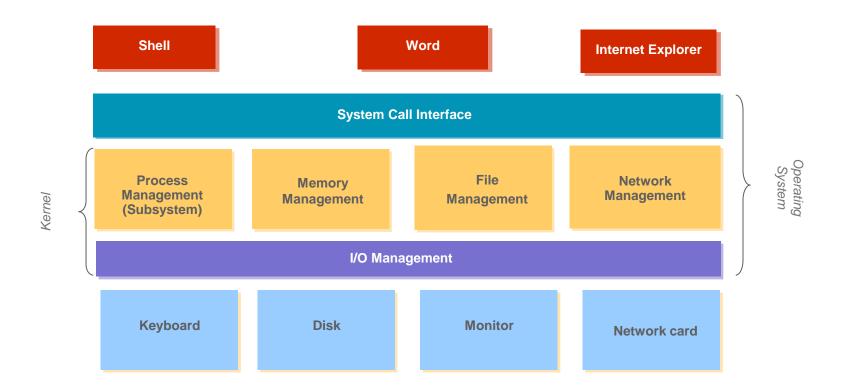
Definition

- A program that acts as an intermediary between a user of a computer and the computer hardware.
- Common functions of controlling and allocating resources are brought together into one piece of software
- The one program running at all times on the computer, usually called kernel
- No universally accepted definition
 - "Everything a vendor ships when you order an operating system" is good approximation, but varies wildly

Operating System Roles

- User view
 - To maximize works (play) that users are performing. In this case, the operating system is designed mostly for ease of use.
- System view
 - OS is a resource allocator
 - Manages all resources
 - Decides between conflicting requests for efficient and fair resource use
 - Process management, memory management, etc.
 - OS is a control program
 - Controls execution of programs to prevent errors and improper use of the computer
 - Process management, memory management, etc

Operating System Structure



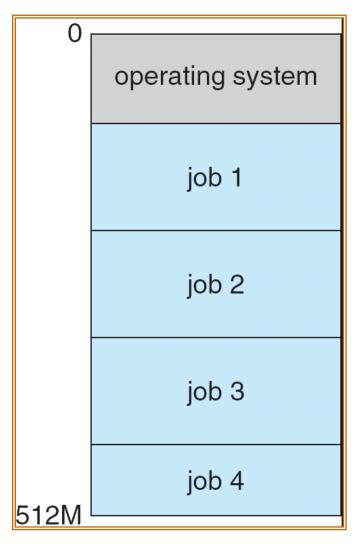
Operating System Structure (Cont'd)

- Multiprogramming/Multitasking
 - Single user or job cannot keep CPU and I/O devices busy at all times
 - A subset of total jobs in system is kept in memory
 - When one job has to wait (for I/O for example), OS switches to another job

Subsystems required to support multiprogramming

- program execution in memory ⇒ process
 - Process management
- If several jobs ready to run at the same time ⇒ CPU scheduling
 - Process management
- If processes don't fit in memory, swapping moves them in and out to run
 - Memory management
- Virtual memory allows execution of processes not completely in memory
 - Memory management

Memory Layout for Multiprogrammed System



Process Management

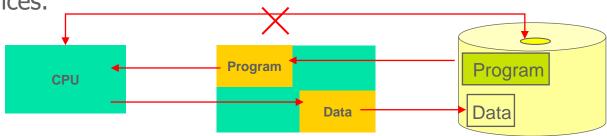
- What is a process?
 - a program in execution.
 - Process has several resources: CPU, memory, I/O, files, data
- Process termination requires reclaim of any reusable resources
- Typically system has many processes, some user, some operating system running concurrently on one or more CPUs
 - Concurrency by multiplexing the CPUs among the processes / threads

Process management is responsible for the following activities

- Creating and deleting both user and system processes
- Suspending and resuming processes
- Providing mechanisms for process synchronization
- Providing mechanisms for process communication
- Providing mechanisms for deadlock handling

Memory Management

- The usage of main memory
 - A repository of quickly accessible data that is shared among the CPU and I/O devices.



- Memory management has an important effect on
 - CPU utilization and computer response to users
- Memory management activities
 - Keeping track of which parts of memory are currently being used and by whom
 - Automatically freeing memory space occupied by a process when it is terminated (Garbage collection)
 - Deciding which processes and data to move into and out of memory
 - Allocating and deallocating memory space as needed

File and Storage Management

- File-System management
 - Abstracts physical properties to logical storage unit file
 - Files usually organized into directories
 - Access control on most systems to determine who can access what
- Activities in file management
 - Creating and deleting files and directories
 - Primitives to manipulate files and directories
 - Mapping files onto secondary storage
- Storage Management
 - Disks and SSDs are used to store data that does not fit in main memory or data that must be kept for a "long" period of time.
 - Entire speed of computer operation depends on disk subsystem and its algorithms
- Activities in storage management
 - Free-space management
 - Storage allocation
 - Disk scheduling

I/O Subsystem

- One purpose of OS is to hide peculiarities of hardware devices from the user
- I/O subsystem responsible for
 - Caching (storing parts of data in faster storage for performance)
 - General device-driver interface
 - Drivers for specific hardware devices

Protection and Security

- **Protection** any mechanism for controlling access of processes or users to resources defined by the OS
- Security defense of the system against internal and external attacks
 - Huge range, including denial-of-service, worms, viruses, identity theft, theft of service

Dual-mode Operation

- Dual-mode operation allows OS to protect itself and other system components
 - User mode and kernel mode
- Hardware provides at least two modes:
 - "Kernel" mode (or "supervisor" or "protected")
 - "User" mode: Normal programs executed
- Some instructions/ops prohibited in user mode:
 - Example: cannot modify page tables in user mode
 - Attempt to modify ⇒ Exception generated

Dual-mode Operation (Cont'd)

Exceptions/Interrupts

 Any event that disturbs the normal execution of the processor and focuses the processor into execution of special instructions.

Exceptions

internal events such as events generated by the execution of processor instructions.

Examples of exceptions

- Alignment exceptions: processors read/writes at the beginning of an odd memory.
- Arithmetic operations: division by zero.
- TRAP

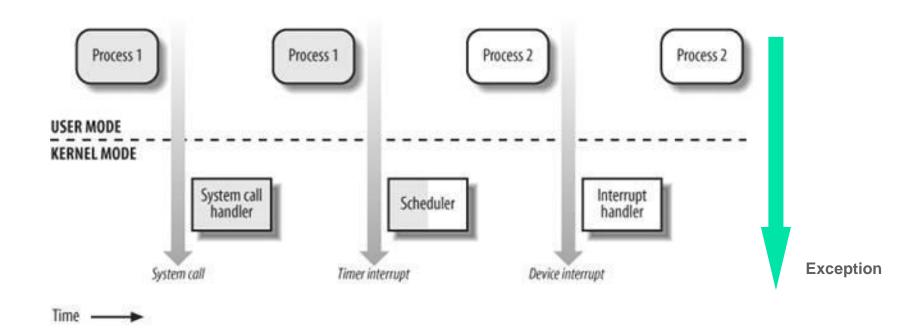
Interrupts

- Electric signal raised by external events, which do not relate to the execution of processor instructions.
- The source of external events are typically external hardware devices.

Examples of interrupts

- The reset button on embedded boards.
- Electrical signals generated on devices.

Transition from User to Kernel Mode



Transition from User to Kernel Mode (Cont'd)

- Implementation
 - Mode bit provided by hardware
 - Provides ability to distinguish when system is running user code or kernel code
 - Some instructions designated as privileged, only executable in kernel mode
 - System call changes mode to kernel, return from call resets it to user

