

# IME ACM-ICPC Team Notebook

## Contents

### 1 Template + vimrc

1.1	Template	1
1.2	vimrc	1

### 2 Graphs

2.1	DFS	2
2.2	BFS	2
2.3	Zero-One-BFS	2
2.4	Toposort	2
2.5	MST (Kruskal)	3
2.6	MST (Prim)	3
2.7	Shortest Path (Dijkstra)	3
2.8	Shortest Path (SPFA)	4
2.9	Max Flow	4
2.10	Min Cost Max Flow	4
2.11	Max Bipartite Cardinality Matching (Kuhn)	4
2.12	Lowest Common Ancestor	4
2.13	2-SAT	4
2.14	Assignment Problem	5

### 3 Mathematics

3.1	Fast Exponential	5
3.2	Fast Fourier Transform	5
3.3	Mobius Function	6
3.4	Sieve	6

### 4 Strings

4.1	Aho-Corasick	6
4.2	Rabin-Karp	6
4.3	Manacher	7
4.4	Suffix Automaton	7
4.5	Knuth-Morris-Pratt	7
4.6	Z Function	7
4.7	String Hashing	8

### 5 Data Structures

5.1	BIT (Range Update, Point Query)	8
5.2	Centroid Decomposition	8
5.3	Min Queue Stack	8
5.4	Min Queue Deque	8
5.5	Persistent Segment Tree	9
5.6	Segment Tree with Lazy	9
5.7	Segment Tree	9
5.8	Treap	9
5.9	Union Find Simple	10
5.10	Union Find with Rollback	10
5.11	Union Find Partial Persistent	10

### 6 Dynamic Programming

6.1	Convex Hull Trick	10
6.2	Steiner Tree	10

### 7 Geometry

7.1	Convex Hull	11
7.2	Minimum Enclosing Circle	11

### 8 Miscellaneous

8.1	Merge Sort	11
8.2	Parallel Binary Search	11
8.3	Ternary Search	12

8.4	Ternary Search with Double	12
-----	----------------------------	----

### 9 Math Extra

9.1	Combinatorial formulas	12
9.2	Number theory identities	12
9.3	Stirling Numbers of the second kind	12
9.4	Burnside's Lemma	13
9.5	Numerical integration	13

## 1 Template + vimrc

### 1.1 Template

```
#include <bits/stdc++.h>
using namespace std;

mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());

#define st first
#define nd second
#define pb push_back

#ifdef ONLINE_JUDGE
#define db(x) cerr << #x << " == " << x << endl
#define dbs(x) cerr << x << endl
#define _ << " , " <<
#else
#define db(x) ((void)0)
#define dbs(x) ((void)0)
#endif

typedef long long ll;
typedef long double ld;

typedef pair<int, int> pii;
typedef pair<int, pii> piii;
typedef pair<ll, ll> pll;

const int INF = 0x3f3f3f3f, MOD = 1e9+7;
const int N = 1e5+5;

int main() {
    ios_base::sync_with_stdio(0), cin.tie(0);
}
```

### 1.2 vimrc

```
syntax on
set encoding=utf-8
set et ts=2 sw=0 sts=-1 ai nu hls cindent
set noswapfile
set mouse=a
"colorscheme darcula
set t_Co=256
set t_ut=
colorscheme codedark
nnoremap ; :
vnoremap ; :
noremap <c-j> 15gj
noremap <c-k> 15gk
nnoremap <s-k> i<CR><ESC>
inoremap .. <esc>
vnoremap .. <esc>
nnoremap .. <esc>

let g:darcula_colorterm = 0

"hi Normal guibg=NONE ctermbg=NONE

if empty(glob('~/vim/autoload/plug.vim'))
    silent !curl -fLo ~/vim/autoload/plug.vim --create-dirs
    \ https://raw.githubusercontent.com/junegunn/vim-plug/master/plug.vim
    autocmd VimEnter * PlugInstall --sync | source $MYVIMRC
endif

"Plugins
```

```

call plug#begin('~/.vim/plugged')
Plug 'vim-airline/vim-airline-themes'
Plug 'vim-airline/vim-airline'
Plug 'tpope/vim-fugitive'
Plug 'tomasiser/vim-code-dark'
call plug#end()

let g:airline_powerline_fonts = 1
set guifont=Inconsolata\ for\ Powerline:h1
let g:Powerline_symbols = 'fancy'
set encoding=utf-8
set fillchars+=stl:\ ,stlnc:\
set term=xterm-256color
nmap <F8> :tabprev<CR>
nmap <F9> :tabnext<CR>

let g:airline#extensions#tabline#enabled = 1
let g:airline#extensions#branch#enabled=1
"let g:airline_theme='badwolf'
let g:airline_theme='codedark'
let g:airline#extensions#tabline#enabled=1
let g:airline#extensions#tabline#formatter='unique_tail'
let g:airline_powerline_fonts=1
let g:airline_section_a = airline#section#create(['mode', ' ', 'branch'])

```

## 2 Graphs

### 2.1 DFS

```

// Depth First Search O(V+E)

const int N = 1e5 + 5;

int vis[N];
vector<int> adj[N];

void dfs(int x){
    vis[x] = 1;
    for(auto u : adj[x]) if(!vis[u]) dfs(u);
}

```

### 2.2 BFS

```

// Breadth First Search O(V+E)

const int N = 1e5 + 5;

int vis[N];
vector<int> adj[N];
queue<int> q;

void bfs(int x){
    vis[x] = 1;
    q.push(x);
    while(!q.empty()){
        int u = q.front(); q.pop();
        for(auto v : adj[u]) if(!vis[v]) {
            vis[v] = 1, q.push(v);
        }
    }
}

```

### 2.3 Zero-One-BFS

```

// 0-1 BFS - O(V+E)

const int N = 1e5 + 5;

int dist[N];
vector<pii> adj[N];
deque<pii> dq;

void zero_one_bfs (int x){

```

```

    cl(dist, 63);
    dist[x] = 0;
    dq.push_back({x, 0});
    while(!dq.empty()){
        int u = dq.front().st;
        int ud = dq.front().nd;
        dq.pop_front();
        if(dist[u] < ud) continue;
        for(auto x : adj[u]){
            int v = x.st;
            int w = x.nd;
            if(dist[u] + w < dist[v]){
                dist[v] = dist[u] + w;
                if(w) dq.push_back({v, dist[v]});
                else dq.push_front({v, dist[v]});
            }
        }
    }
}

```

### 2.4 Toposort

```

// Kahn - Topological Sort O(V + E)

const int N = 1e5+5;

vector<int> adj[N];
int n, in[N];

// For directed graph: in[x] == 0
// For undirected graph: in[x] <= 1

void kahn(){
    queue<int> q;
    for(int i = 1; i <= n; i++) if(!in[i]) q.push(i);

    while(q.size()){
        int u = q.front(); q.pop();
        for(auto x : adj[u]) if(in[x] and --in[x] == 0) q.push(x);
    }
}

```

### 2.5 MST (Kruskal)

```

// Kruskal - MST O(ElogE)

// + Union-Find
////////////////////////////////////
int par[N], sz[N];

//Path Compression
int find(int a) { return a == par[a] ? a : par[a] = find(par[a]); }

//Ranking
void unite(int a, int b){
    if(find(a) == find(b)) return;
    a = find(a), b = find(b);
    if(sz[a] < sz[b]) swap(a, b);
    sz[a] += sz[b], par[b] = a;
}

//in main
for(int i = 0; i < N; i++) par[i] = i, sz[i] = 1;
////////////////////////////////////

vector<pii> edges;
// dist, node1, node2

sort(edges.begin(), edges.end());
int cost = 0;
for(auto e : edges) if(find(e.nd.st) != find(e.nd.nd)){
    cost += e.st, unite(e.nd.st, e.nd.nd);
}

```

### 2.6 MST (Prim)

```
// Prim - MST  $O(E \log E)$ 

int cost, vis[N];
vector<pii> adj[N];
priority_queue<pii> pq;

void prim(int s = 1) {
    pq.push({0, s});
    while(!pq.empty()) {
        int ud = -pq.top().st;
        int u = pq.top().nd;
        pq.pop();
        if(vis[u]) continue;
        vis[u] = 1;
        cost += ud;
        for(auto x : adj[u]) {
            int v = x.st;
            int w = x.nd;
            if(!vis[v]) pq.push({-w, v});
        }
    }
}
```

## 2.7 Shortest Path (Dijkstra)

```
// Dijkstra -  $O((V+E) \log E)$ 

int dist[N];
vector<pii> adj[N];
priority_queue<pii> pq;

void dijkstra(int s) {
    cl(dist, 63);
    dist[s] = 0;
    pq.push({0, s});
    while(!pq.empty()) {
        int ud = -pq.top().st;
        int u = pq.top().nd;
        pq.pop();
        if(dist[u] < ud) continue;
        for(auto x : adj[u]) {
            int v = x.st;
            int w = x.nd;
            if(dist[u] + w < dist[v]) {
                dist[v] = dist[u] + w;
                pq.push({-dist[v], v});
            }
        }
    }
}
```

## 2.8 Shortest Path (SPFA)

```
// Shortest Path Faster Algorithm  $O(VE)$ 

int n, inq[N], dist[N];
vector<pair<int, int>> adj[N];

void spfa(int s) {
    memset(dist, 63, sizeof dist);
    queue<int> q;
    dist[s] = 0;
    q.push(s), inq[s] = 1;
    while(q.size()) {
        int u = q.front(); q.pop(); inq[u] = 0;
        for(auto x : adj[u]) {
            int v = x.st;
            int w = x.nd;
            if(dist[u] + w < dist[v]) {
                dist[v] = dist[u] + w;
                if(!inq[v]) q.push(v), inq[v] = 1;
            }
        }
    }
}
```

## 2.9 Max Flow

```
// Dinic -  $O(n^2 * m)$ 
// Max flow

const int N = 1e5 + 5;
const int INF = 0x3f3f3f3f;

struct edge { int v, c, f; };

int n, s, t, h[N], st[N];
vector<edge> eds;
vector<int> g[N];

// directed from u to v with cap(u, v) = c
void add_edge(int u, int v, int c) {
    int k = eds.size();
    eds.push_back({v, c, 0});
    eds.push_back({u, 0, 0});
    g[u].push_back(k);
    g[v].push_back(k+1);
}

int bfs() {
    memset(h, 0, sizeof h);
    h[s] = 1;
    queue<int> q;
    q.push(s);
    while(q.size()) {
        int u = q.front(); q.pop();
        for(auto i : g[u]) {
            int v = eds[i].v;
            if(!h[v] and eds[i].f < eds[i].c)
                h[v] = h[u] + 1, q.push(v);
        }
    }
    return h[t];
}

int dfs(int u, int flow) {
    if(!flow or u == t) return flow;
    for(int &i = st[u]; i < g[u].size(); i++) {
        edge &dir = eds[g[u][i]], &rev = eds[g[u][i]^1];
        int v = dir.v;
        if(h[v] != h[u] + 1) continue;
        int inc = min(flow, dir.c - dir.f);
        inc = dfs(v, inc);
        if(inc) {
            dir.f += inc, rev.f -= inc;
            return inc;
        }
    }
    return 0;
}

int dinic() {
    int flow = 0;
    while(bfs()) {
        memset(st, 0, sizeof st);
        while(int inc = dfs(s, INF)) flow += inc;
    }
    return flow;
}
```

## 2.10 Min Cost Max Flow

```
// Min Cost Max Flow -  $O(n^2 * m^2)$ 

struct edge { int v, f, c, w; };

vector<int> g[N];
vector<edge> eds;
int s, t, inq[N], p[N], dist[N];

void add_edge(int u, int v, int c, int w) {
    int k = eds.size();
    g[u].push_back(k);
    g[v].push_back(k+1);
    eds.push_back({v, 0, c, w});
    eds.push_back({u, 0, 0, -w});
}
```

```

}

int spfa() {
    memset(dist, 63, sizeof dist);
    queue<int> q;
    dist[s] = 0;
    q.push(s), inq[s] = 1;
    while(q.size()) {
        int u = q.front(); q.pop(); inq[u] = 0;
        for(auto i : g[u]) {
            edge dir = eds[i];
            int v = dir.v;
            int w = dir.w;
            if(dir.f < dir.c and dist[u] + w < dist[v]) {
                dist[v] = dist[u] + w;
                p[v] = i;
                if(!inq[v]) q.push(v), inq[v] = 1;
            }
        }
    }

    if(dist[t] == INF) return 0;

    int inc = INF;
    for(int u = t; u != s; u = eds[p[u]^1].v) {
        edge &dir = eds[p[u]];
        inc = min(inc, dir.c - dir.f);
    }

    int aux = 0;
    for(int u = t; u != s; u = eds[p[u]^1].v) {
        edge &dir = eds[p[u]], &rev = eds[p[u]^1];
        dir.f += inc;
        rev.f -= inc;
        aux += inc*dir.w;
    }

    return aux;
}

int mcmf() {
    int cost = 0;
    while(int inc = spfa()) cost += inc;
    return cost;
}

```

## 2.11 Max Bipartite Cardinality Matching (Kuhn)

```

// Khun (Maximum Bipartite Matching) - O(VE)

int n, cnt, vis[N], match[N], ans;

bool find(int x){
    if(vis[x] == cnt) return false;
    vis[x] = cnt;
    for(auto u : adj[x]) if(!match[u] or find(match[u])) return match[u] = x;
    return false;
}

// Maximum Independent Set on bipartite graph
// MIS = V - MATCH

// Minimum Vertex Cover
// MVC = MATCH

// Minimum Path Cover on DAG
// MPC = V - MATCH

// TIP: If you don't know the sides of the bipartite graph,
// run kuhn for all nodes and match = ans/2;

// in main (only for one of the sides)
for(int i = 1; i <= n; i++) ++cnt, ans += find(i);

```

## 2.12 Lowest Common Ancestor

```

// Lowest Common Ancestor - <O(nlog n), O(log n)>

const int N = 1e6;
const int M = 25; //m = log N

```

```

int anc[M][N], h[N], rt;
vector<int> adj[N];

void dfs(int x = rt, int p = -1, int ht = 0){
    anc[0][x] = p, h[x] = ht;
    for(auto v : adj[x]) if(v != p) dfs(v, x, ht+1);
}

void build(){
    dfs(), anc[0][rt] = rt;

    for(int j = 1; j < M; j++){
        for(int i = 1; i <= n; i++) // 1-indexed
            anc[j][i] = anc[j-1][anc[j-1][i]];
    }

    int lca(int u, int v){
        if(h[u] < h[v]) swap(u, v);
        for(int i = M-1; i >= 0; i--) if(h[u] - (1<<i) >= h[v]) u = anc[i][u];

        if(u == v) return u;
        for(int i = M-1; i >= 0; i--) if(anc[i][u] != anc[i][v])
            u = anc[i][u], v = anc[i][v];

        return anc[0][u];
    }
}

```

## 2.13 2-SAT

```

// 2-SAT - O(V+E)

int n, vis[2*N], ord[2*N], ordn, cnt, cmp[2*N], val[N];
vector<int> adj[2*N], adjt[2*N];

// for a variable u with idx i
// u is 2*i and !u is 2*i+1
// (a v b) == !a -> b ^ !b -> a

int v(int x) { return 2*x; }
int nv(int x) { return 2*x+1; }

// add a -> b
void add(int a, int b) {
    adj[a].push_back(b);
    adj[b^1].push_back(a^1);
    adjt[b].push_back(a);
    adjt[a^1].push_back(b^1);
}

// add clause (a v b)
void add_or(int a, int b){
    adj[a^1].push_back(b);
    adj[b^1].push_back(a);
    adjt[b].push_back(a^1);
    adjt[a].push_back(b^1);
}

void dfs(int x){
    vis[x] = 1;
    for(auto v : adj[x]) if(!vis[v]) dfs(v);
    ord[ordn++] = x;
}

void dfst(int x){
    cmp[x] = cnt, vis[x] = 0;
    for(auto v : adjt[x]) if(vis[v]) dfst(v);
}

bool run2sat(){
    for(int i = 1; i <= n; i++) {
        if(!vis[v(i)]) dfs(v(i));
        if(!vis[nv(i)]) dfs(nv(i));
    }
    for(int i = ordn-1; i >= 0; i--)
        if(vis[ord[i]]) cnt++, dfst(ord[i]);
    for(int i = 1; i <= n; i++){
        if(cmp[v(i)] == cmp[nv(i)]) return false;
        val[i] = cmp[v(i)] > cmp[nv(i)];
    }
    return true;
}

```

## 2.14 Assignment Problem

```
// Hungarian - O(m*n^2)
// Assignment Problem

int n, m;
int pu[N], pv[N], cost[N][M];
int pairV[N], way[M], minv[M], used[M];

void hungarian() {
    for(int i = 1, j0 = 0; i <= n; i++) {
        pairV[0] = i;
        memset(minv, 63, sizeof minv);
        memset(used, 0, sizeof used);
        do {
            used[j0] = 1;
            int i0 = pairV[j0], delta = INF, j1;
            for(int j = 1; j <= m; j++) {
                if(used[j]) continue;
                int cur = cost[i0][j] - pu[i0] - pv[j];
                if(cur < minv[j]) minv[j] = cur, way[j] = j0;
                if(minv[j] < delta) delta = minv[j], j1 = j;
            }

            for(int j = 0; j <= m; j++) {
                if(used[j]) pu[pairV[j]] += delta, pv[j] -= delta;
                else minv[j] -= delta;
            }
            j0 = j1;
        } while(pairV[j0]);

        do {
            int j1 = way[j0];
            pairV[j0] = pairV[j1];
            j0 = j1;
        } while(j0);
    }
}

// in main
// for(int j = 1; j <= m; j++)
//     if(pairV[j]) ans += cost[pairV[j]][j];
//
```

## 3 Mathematics

### 3.1 Fast Exponential

```
// Fast Exponential - O(log b)
ll fexp (ll b, ll e, ll mod) {
    ll ans = 1;
    while (e) {
        if (e&1) ans = (ans*b) % mod;
        b = (b*b) % mod, e >>= 1;
    }
    return ans;
}
```

### 3.2 Fast Fourier Transform

```
// FFT - Polynomial Multiplication in O(n log n)
// Made by tourist

// p(x)^k -> remeber to use fast exponentiation
// mod multiplication -> every coefficient in range [0, mod-1]
// be careful with overflow!

namespace fft {
    typedef double dbl;

    struct num {
        dbl x, y;
        num() { x = y = 0; }
        num(dbl x, dbl y) : x(x), y(y) {}
    };
}
```

```
};

inline num operator+ (num a, num b) { return num(a.x + b.x, a.y + b.y); }
inline num operator- (num a, num b) { return num(a.x - b.x, a.y - b.y); }
inline num operator* (num a, num b) { return num(a.x * b.x - a.y * b.y, a.x * b.y + a.y * b.x); }
inline num conj(num a) { return num(a.x, -a.y); }

int base = 1;
vector<num> roots = {{0, 0}, {1, 0}};
vector<int> rev = {0, 1};

const dbl PI = acos(-1.0);

void ensure_base(int nbase) {
    if(nbase <= base) return;

    rev.resize(1 << nbase);
    for(int i=0; i < (1 << nbase); i++) {
        rev[i] = (rev[i] >> 1) >> 1 + ((i & 1) << (nbase - 1));
    }
    roots.resize(1 << nbase);

    while(base < nbase) {
        dbl angle = 2*PI / (1 << (base + 1));
        for(int i = 1 << (base - 1); i < (1 << base); i++) {
            roots[i << 1] = roots[i];
            dbl angle_i = angle * (2 * i + 1 - (1 << base));
            roots[(i << 1) + 1] = num(cos(angle_i), sin(angle_i));
        }
        base++;
    }
}

void fft(vector<num> &a, int n = -1) {
    if(n == -1) {
        n = a.size();
    }
    assert((n & (n-1)) == 0);
    int zeros = __builtin_ctz(n);
    ensure_base(zeros);
    int shift = base - zeros;
    for(int i = 0; i < n; i++) {
        if(i < (rev[i] >> shift)) {
            swap(a[i], a[rev[i] >> shift]);
        }
    }
    for(int k = 1; k < n; k <= 1) {
        for(int i = 0; i < n; i += 2 * k) {
            for(int j = 0; j < k; j++) {
                num z = a[i+j+k] * roots[j+k];
                a[i+j+k] = a[i+j] - z;
                a[i+j] = a[i+j] + z;
            }
        }
    }
}

vector<num> fa, fb;
vector<int> multiply(vector<int> &a, vector<int> &b) {
    int need = a.size() + b.size() - 1;
    int nbase = 0;
    while((1 << nbase) < need) nbase++;
    ensure_base(nbase);
    int sz = 1 << nbase;
    if(sz > (int) fa.size()) {
        fa.resize(sz);
    }
    for(int i = 0; i < sz; i++) {
        int x = (i < (int) a.size() ? a[i] : 0);
        int y = (i < (int) b.size() ? b[i] : 0);
        fa[i] = num(x, y);
    }
    fft(fa, sz);
    num r(0, -0.25 / sz);
    for(int i = 0; i <= (sz >> 1); i++) {
        int j = (sz - i) & (sz - 1);
        num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
        if(i != j) {
            fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r;
        }
        fa[i] = z;
    }
    fft(fa, sz);
    vector<int> res(need);
    for(int i = 0; i < need; i++) {
        res[i] = fa[i].x + 0.5;
    }
    return res;
}
```

```

vector<int> multiply_mod(vector<int> &a, vector<int> &b, int m, int eq = 0) {
    int need = a.size() + b.size() - 1;
    int nbase = 0;
    while ((1 << nbase) < need) nbase++;
    ensure_base(nbase);
    int sz = 1 << nbase;
    if (sz > (int) fa.size()) {
        fa.resize(sz);
    }
    for (int i = 0; i < (int) a.size(); i++) {
        int x = (a[i] % m + m) % m;
        fa[i] = num(x & ((1 << 15) - 1), x >> 15);
    }
    fill(fa.begin() + a.size(), fa.begin() + sz, num {0, 0});
    fft(fa, sz);
    if (sz > (int) fb.size()) {
        fb.resize(sz);
    }
    if (eq) {
        copy(fa.begin(), fa.begin() + sz, fb.begin());
    } else {
        for (int i = 0; i < (int) b.size(); i++) {
            int x = (b[i] % m + m) % m;
            fb[i] = num(x & ((1 << 15) - 1), x >> 15);
        }
        fill(fb.begin() + b.size(), fb.begin() + sz, num {0, 0});
        fft(fb, sz);
    }
    dbl ratio = 0.25 / sz;
    num r2(0, -1);
    num r3(ratio, 0);
    num r4(0, -ratio);
    num r5(0, 1);
    for (int i = 0; i <= (sz >> 1); i++) {
        int j = (sz - i) & (sz - 1);
        num a1 = (fa[i] + conj(fa[j]));
        num a2 = (fa[i] - conj(fa[j])) * r2;
        num b1 = (fb[i] + conj(fb[j])) * r3;
        num b2 = (fb[i] - conj(fb[j])) * r4;
        if (i != j) {
            num c1 = (fa[j] + conj(fa[i]));
            num c2 = (fa[j] - conj(fa[i])) * r2;
            num d1 = (fb[j] + conj(fb[i])) * r3;
            num d2 = (fb[j] - conj(fb[i])) * r4;
            fa[i] = c1 * d1 + c2 * d2 * r5;
            fb[i] = c1 * d2 + c2 * d1;
        }
        fa[j] = a1 * b1 + a2 * b2 * r5;
        fb[j] = a1 * b2 + a2 * b1;
    }
    fft(fa, sz);
    fft(fb, sz);
    vector<int> res(need);
    for (int i = 0; i < need; i++) {
        long long aa = fa[i].x + 0.5;
        long long bb = fb[i].x + 0.5;
        long long cc = fa[i].y + 0.5;
        res[i] = (aa + ((bb % m) << 15) + ((cc % m) << 30)) % m;
    }
    return res;
}

vector<int> square_mod(vector<int> &a, int m) {
    return multiply_mod(a, a, m, 1);
}
}

```

### 3.3 Mobius Function

```

// Mobius Function
// u(1) = 1
// u(p) = -1
// u(p^k) = 0, k >= 2
// u(a*b) = u(a)*u(b), for a,b co-primes

// Sum for d|n of u(d) = [n == 1]

// Calculate Mobius all integers - O(n log(log n))

int cmp[N], mob[N];
void mobius() {
    for(int i = 1; i < N; i++) mob[i] = 1;
    for(ll i = 2; i < N; i++) if(!cmp[i]) {
        for(ll j = i; j < N; j += i) cmp[j] = 1, mob[j] *= -1;
    }
}

```

```

    for(ll j = i+i; j < N; j += i+i) mob[j] = 0;
}

// Calculate Mobius for 1 integer - O(sqrt(n))
int mobius(int n) {
    if(n == 1) return 1;
    int p = 0;
    for(int i = 2; i+i <= n; i++) {
        if(n % i == 0) {
            n = n/i, p++;
            if(n % i == 0) return 0;
        }
    }
    if(n > 1) p++;
    return p%1 ? -1 : 1;
}

```

### 3.4 Sieve

```

// Sieve - O(n log(log n))

int cmp[N];
vector<int> p;

void sieve() {
    for(ll i = 2; i < N; i++) if(!cmp[i]) {
        p.push_back(i);
        for(ll j = i+i; j < N; j += i) cmp[j] = 1;
    }
}

```

## 4 Strings

### 4.1 Aho-Corasick

```

// Aho-Corasick - <O(sum(m)), O(n + #matches)>

// number of nodes
const int N = 1e5 + 5;
// number of strings
const int S = 2e3 + 5;

int nxt[N][26], ch[N], p[N], f[N], szt = 1, cnt[N];
bitset<S> elem[N];

void add(string &pt, int x) {
    int u = 1;
    for(auto c : pt) {
        int j = c - 'a';
        if(!nxt[u][j]) {
            szt++;
            ch[szt] = j;
            p[szt] = u;
            nxt[u][j] = szt;
        }
        u = nxt[u][j];
    }
    cnt[u]++;
    elem[u].set(x);
}

void build() {
    queue<int> q;
    for(int i = 0; i < 26; i++) {
        nxt[0][i] = 1;
        if(nxt[1][i]) q.push(nxt[1][i]);
    }
    while(q.size()) {
        int v = q.front(); q.pop();
        int u = f[p[v]];
        while(u and !nxt[u][ch[v]]) u = f[u];
        f[v] = nxt[u][ch[v]];
        cnt[v] += cnt[f[v]];
        elem[v] |= elem[f[v]];
        for(int i = 0; i < 26; i++) {
            if(nxt[v][i]) q.push(nxt[v][i]);
        }
        /* Pre-Computation of next states
    }
}

```

```

    else {
        int ax = f[v];
        while(ax and !nxt[ax][i]) ax = f[ax];
        nxt[v][i] = nxt[ax][i];
    }
}
*/
}
}

int match(string &s) {
    int s = 0, u = 1;
    bitset<S> ans;
    for(auto c : s) {
        int j = c - 'a';
        while(u and !nxt[u][j]) u = f[u];
        u = nxt[u][j];
        s += cnt[u];
        ans |= elem[u];
    }
    return s;
}
}

```

## 4.2 Rabin-Karp

```

// Rabin-Karp (String Matching + Hashing)
const int MOD = 1e9+9;
const int B = 313;
char s[N], p[N];
int n, m; // n = strlen(s), m = strlen(p)

// Chance of collision for k generated values and N possible hash values
// e^(-k*(k-1)/2*N)

int rabin() {
    if(n < m) return 0;
    ull hp = 0, hs = 0, E = 1, oc = 0;
    for(int i = 0; i < m; i++) {
        hp = ((hp*B)%MOD + p[i])%MOD;
        hs = ((hs*B)%MOD + s[i])%MOD;
        E = (E*B)%MOD;
    }

    if(hs == hp) oc++; //match at 0
    for(int i = m; i < n; i++) {
        hs = ((hs*B)%MOD + s[i])%MOD;
        hs = (hs - s[i-m]*E%MOD + MOD)%MOD;
        if(hs == hp) oc++; //match at i-m+1
    }
    return oc;
}
}

```

## 4.3 Manacher

```

// Manacher - O(n)

// d1 -> odd : size = 2*d1[i] - 1
// d2 -> even : size = 2*d2[i]

vector<int> d1, d2;

void manacher(string &s) {
    int n = s.size();
    d1.resize(n), d2.resize(n);
    for(int i = 0, l1 = 0, l2 = 0, r1 = -1, r2 = -1; i < n; i++) {
        d1[i] = 1;
        if(i <= r1) d1[i] = min(d1[r1+l1-i], r1-i+1);
        if(i <= r2) d2[i] = min(d2[r2+l2-i-1], r2-i);
        while(i - d1[i] >= 0 and i + d1[i] < n and
            s[i - d1[i]] == s[i + d1[i]]) d1[i]++;
        while(i - d2[i] >= 0 and i + d2[i] + 1 < n and
            s[i - d2[i]] == s[i + d2[i] + 1]) d2[i]++;
        if(i + d1[i] - 1 > r1) l1 = i - d1[i] + 1, r1 = i + d1[i] - 1;
        if(i + d2[i] > r2) l2 = i - d2[i] + 1, r2 = i + d2[i];
    }
}
}

```

## 4.4 Suffix Automaton

```

#include <bits/stdc++.h>
using namespace std;

const int N = 1e5 + 5;

int len[2*N], link[2*N], nxt[2*N][26], szt = 1, last = 1;

void add(string &s) {
    for(auto c : s) {
        int cur = ++szt, p = last, j = c - 'a';
        len[cur] = len[last] + 1;
        while(p and !nxt[p][j]) {
            nxt[p][j] = cur;
            p = link[p];
        }
        if(!p) link[cur] = 1;
        else {
            int q = nxt[p][j];
            if(len[p] + 1 == len[q]) link[cur] = q;
            else {
                int clone = ++szt;
                len[clone] = len[p] + 1;
                link[clone] = link[q];
                for(int i = 0; i < 26; i++)
                    if(nxt[q][i]) nxt[clone][i] = nxt[q][i];
                while(p and nxt[p][j] == q) {
                    nxt[p][j] = clone;
                    p = link[p];
                }
                link[q] = link[cur] = clone;
            }
        }
        last = cur;
    }
}

int main() {
}

```

## 4.5 Knuth-Morris-Pratt

```

// KMP - O(n + m)

// max size pattern
const int N = 1e5 + 5;

int lps[N], cont;

void prekmp(string &p) {
    for (int i = 1, j = 0; i < p.size(); i++) {
        while (j > 0 and p[j] != p[i]) j = lps[j-1];
        if (p[j] == p[i]) j++;
        lps[i] = j;
    }
}

void kmp(string &s, string &p) {
    for (int i = 0, j = 0; i < s.size(); i++) {
        while (j > 0 and p[j] != s[i]) j = lps[j-1];
        if (p[j] == s[i]) j++;
        if (j == p.size()) {
            // match i-j+1
            cont++;
            j = lps[j-1];
        }
    }
}
}

```

## 4.6 Z Function

```

// Z-Function - O(n)

```

```
vector<int> z(string s){
    vector<int> z(s.size());
    for(int i = 1, l = 0, r = 0, n = s.size(); i < n; i++){
        if(i <= r) z[i] = min(z[i-1], r - i + 1);
        while(i + z[i] < n and s[z[i]] == s[z[i] + i]) z[i]++;
        if(i + z[i] - 1 > r) l = i, r = i + z[i] - 1;
    }
    return z;
}
```

## 4.7 String Hashing

```
// String Hashing
// Rabin Karp - O(n + m)

// max size txt + 1
const int N = 1e6 + 5;

// lowercase letters p = 31
// uppercase letters p = 53
// any character p = 313

const int MOD = 1e9+9;
ull h[N], p[N] = {1, 31};
int cnt;

void build(string &s) {
    for(int i = 1; i <= s.size(); i++) {
        h[i] = ((p[i]*h[i-1]) % MOD + s[i-1]) % MOD;
        p[i] = (p[i]*p[i-1]) % MOD;
    }
}

// 1-indexed
ull fhash(int l, int r) {
    return (h[r] - ((h[l-1]*p[r-l+1]) % MOD) + MOD) % MOD;
}

ull shash(string &pt) {
    ull h = 0;
    for(int i = 0; i < pt.size(); i++)
        h = ((h*p[1]) % MOD + pt[i]) % MOD;
    return h;
}

void rabin_karp(string &s, string &pt) {
    build(s);
    ull hp = shash(pt);
    for(int i = 0, m = pt.size(); i + m <= s.size(); i++) {
        if(fhash(i+1, i+m) == hp) {
            // match at i
            cnt++;
        }
    }
}
```

## 5 Data Structures

### 5.1 BIT (Range Update, Point Query)

```
// Binary Indexed Tree
// Range Update and Point Query
// Update - O(log n)
// Query - O(log n)

int bit[N];

void add(int p, int v) {
    for (p+=2; p<N; p+=p&-p) bit[p] += v;
}

void update(int l, int r, int val){
    add(l, val);
    if(r != N) add(r+1, -val);
}

int query(int p) {
```

```
int r = 0;
for (p+=2; p; p-=p&-p) r += bit[p];
return r;
}
```

### 5.2 Centroid Decomposition

```
// Centroid Decomposition - O(n log n)
int n, m, sz[N], forb[N], par[N];

void dfs(int u, int p) {
    sz[u] = 1;
    for(auto v : adj[u]) {
        if(v != p and !forb[v]) {
            dfs(v, u);
            sz[u] += sz[v];
        }
    }

    int cent(int u, int p, int amt) {
        for(auto v : adj[u]) {
            if(v == p or forb[v]) continue;
            if(sz[v] > amt/2) return cent(v, u, amt);
        }
        return u;
    }

    void decomp(int u, int p) {
        dfs(u, -1);
        int cen = cent(u, -1, sz[u]);
        forb[cen] = 1;
        if(p != -1) par[cen] = p;

        for(auto v : adj[u])
            if(!forb[v]) decomp(v, cen);
    }

    // in main
    // decomp(1, -1);
}
```

### 5.3 Min Queue Stack

```
// Min Queue
// Operations in O(1)

struct MinQueue{
    vector<pii> s1, s2;

    int size(){ return s1.size() + s2.size(); }

    int mini(){
        if(s1.empty() or s2.empty()) return s1.empty() ? s2.back().nd : s1.back().nd;
        return min(s1.back().nd, s2.back().nd);
    }

    void add(int x){
        int mn = s1.empty() ? x : min(s1.back().nd, x);
        s1.pb({x, mn});
    }

    void rem(){
        if(s2.empty()){
            while(s1.size()){
                int x = s1.back().st; s1.pop_back();
                int mn = s2.empty() ? x : min(x, s2.back().nd);
                s2.pb({x, mn});
            }
            s2.pop_back();
        }
    }
};
```

### 5.4 Min Queue Deque



```
// Min Queue - Operations in O(1)

struct MinQueue {
    int sum = 0, l = 1, r = 0;
    deque<pair<int, int>> d;

    void reset(){
        sum = 0, l = 1, r = 0;
        d.clear();
    }

    void push(int x){
        x -= sum;
        while(!d.empty() and d.back().first >= x) d.pop_back();
        d.push_back({x, ++r});
    }

    void add(int x) { sum += x; }
    void pop() { if(!d.empty() and d.front().second == l++) d.pop_front(); }
    int min() { return d.front().first + sum; }
    int size() { return r - l + 1; }
};
```

## 5.5 Persistent Segment Tree

```
//Persistent Segment Tree
//Update and Query - O(log n)
//Space - O(n log n)

//M -> n log n

int nxt;
int lc[M], rc[M], st[M];
vector<int> root;

int update(int n, int l, int r, int idx, int v){
    if(l == r) { st[nxt] = v; return nxt++; }
    int mid = (l+r)/2;
    int u = nxt++;
    if(i <= mid){
        lc[u] = update(lc[n], l, mid, idx, v);
        rc[u] = rc[n];
    }else{
        rc[u] = update(rc[n], mid, r, idx, v);
        lc[u] = lc[n];
    }
    st[u] = st[lc[u]] + st[rc[u]]; // RMQ -> min/max, RSQ -> +
    return u;
}

int query(int n, int l, int r, int i, int j){
    if(i > r or l > j) return 0;
    if(i <= l and j >= r) return st[n];
    return query(lc[n], l, (l+r)/2, i, j) + query(rc[n], (l+r)/2 + 1, r, i, j); // RMQ -> min/max, RSQ -> +
}
```

## 5.6 Segment Tree with Lazy

```
//Segment Tree (Range Query and Point Update)
//Update and Query - O(log n)
//Space - O(n)

int n, v[N], lz[4*N], st[4*N];

// n - size of the array (up to N)
void build(int p = 1, int l = 1, int r = n){
    if(l == r) { st[p] = v[l]; return; }
    build(2*p, l, (l+r)/2), build(2*p + 1, (l+r)/2 + 1, r);
    st[p] = st[2*p] + st[2*p + 1];
}

void push(int p, int l, int r){
    if(lz[p]){
        st[p] = lz[p]*(r-l+1);
        if(l != r) lz[2*p] = lz[2*p + 1] = lz[p];
        lz[p] = 0;
    }
}

int query(int i, int j, int p = 1, int l = 1, int r = n){
```

```
    push(p, l, r);
    if(l > j or r < i) return 0;
    if(l >= i and j >= r) return st[p];
    return query(i, j, 2*p, l, (l+r)/2) + query(i, j, 2*p + 1, (l+r)/2 + 1, r);
}

void update(int i, int j, int v, int p = 1, int l = 1, int r = n){
    push(p, l, r);
    if(l > j or r < i) return;
    if(l >= i and j >= r) { lz[p] = v, push(p, l, r); return; }
    update(i, j, v, 2*p, l, (l+r)/2), update(i, j, v, 2*p+1, (l+r)/2+1, r);
    st[p] = st[2*p] + st[2*p + 1];
}
```

## 5.7 Segment Tree

```
//Segment Tree (Range Query and Point Update)
//NOT the Lazy Propagation version
//Update and Query - O(log n)
//Space - O(n)

int n, v[N], st[4*N];

// n - size of the array (up to N)
// You could do point update in all values of v, instead of using build
void build(int p = 1, int l = 1, int r = n){
    if(l == r) { st[p] = v[l]; return; }
    build(2*p, l, (l+r)/2), build(2*p + 1, (l+r)/2 + 1, r);
    st[p] = st[2*p] + st[2*p + 1];
}

int query(int i, int j, int p = 1, int l = 1, int r = n){
    if(l >= i and j >= r) return st[p];
    if(l > j or r < i) return 0;
    return query(i, j, 2*p, l, (l+r)/2) + query(i, j, 2*p + 1, (l+r)/2 + 1, r);
}

void update(int idx, int v, int p = 1, int l = 1, int r = n){
    if(l == r) { st[p] = v; return; }
    if(idx <= (l+r)/2) update(idx, v, 2*p, l, (l+r)/2);
    else update(idx, v, 2*p + 1, (l+r)/2 + 1, r);
    st[p] = st[2*p] + st[2*p + 1];
}
```

## 5.8 Treap

```
// Treap
// Operations in O(log n)

mt19937_64 llrand(random_device{}());

struct node{
    int val, cnt;
    node *r, *l;
    ll pri;
    node(int x) : val(x), cnt(1), pri(llrand()), l(0), r(0) {}
};

struct treap{
    node *root;

    int cnt(node *t) {return t ? t->cnt : 0;}

    void update(node *t){
        if(!t) return;
        t->cnt = cnt(t->l) + cnt(t->r) + 1;
    }

    node *merge(node *l, node *r){
        if(!l and !r) return nullptr;
        if(!l or !r) return l ? l : r;
        node *t;
        if(l->pri > r->pri) t = l, t->r = merge(l->r, r);
        else t = r, t->l = merge(l, r->l);
        update(t);
        return t;
    }

    pair<node*, node*> split(node *t, int pos){
        if(!t) return {0, 0};
```

```

    if(cnt(t->l) < pos){
        auto x = split(t->r, pos - cnt(t->l) - 1);
        t->r = x.st;
        update(t);
        return {t, x.nd};
    }

    auto x = split(t->l, pos);
    t->l = x.nd;
    update(t);
    return {x.st, t};
}
};

```

## 5.9 Union Find Simple

```

//Union-Find
//Union and Find - O(alpha n)

int par[N], sz[N];

//Path Compression
int find(int a) { return a == par[a] ? a : par[a] = find(par[a]); }

//Ranking
void unite(int a, int b){
    if(find(a) == find(b)) return;
    a = find(a), b = find(b);
    if(sz[a] < sz[b]) swap(a, b);
    sz[a] += sz[b], par[b] = a;
}

//in main
for(int i = 0; i < N; i++) par[i] = i, sz[i] = 1;

```

## 5.10 Union Find with Rollback

```

//Union-Find with Rollback
//Union and find - O(log n)

int par[N], sz[N];
vector<pii> old_par, old_sz;

int find(int a) { return par[a] == a ? a : find(par[a]); }

void unite(int a, int b){
    if(find(a) == find(b)) return;
    a = find(a), b = find(b);
    if(sz[a] < sz[b]) swap(a, b);
    old_par.pb({b, par[b]});
    old_par.pb({a, sz[a]});
    sz[a] += sz[b], par[b] = a;
}

void roolback(){
    par[old_par.top().st] = old_par.top().nd;
    sz[old_sz.top().st] = old_sz.top().nd;
    old_par.pop();
    old_sz.pop();
}

//in main
for(int i = 0; i < N; i++) par[i] = i, sz[i] = 1;

```

## 5.11 Union Find Partial Persistent

```

//Union-Find with Partial Persistence
//Union and Find - O(log n)

int t, par[N], sz[N], his[N];

int find(int a, int t){
    if(par[a] == a) return a;
    if(his[a] > t) return a;
    return find(par[a], t);
}

```

```

void unite(int a, int b){
    if(find(a, t) == find(b, t)) return;
    a = find(a, t), b = find(b, t), t++;
    if(sz[a] < sz[b]) swap(a, b);
    sz[a] += sz[b], par[b] = a, his[b] = t;
}

//in main
for(int i = 0; i < N; i++) par[i] = i, sz[i] = 1, his[i] = 0;

```

## 6 Dynamic Programming

### 6.1 Convex Hull Trick

```

// Convex Hull Trick

typedef long long type;
struct line { type b, m; };

int nh, pos;
line hull[N];

bool check(line s, line t, line u){
    //attention for overflow
    return 1d (u.b - t.b)/(t.m - u.m) > 1d (t.b - s.b)/(s.m - t.m);
}

void update(line s){
    if(nh == 1 and hull[nh-1].b == s.b) nh--;
    if(nh > 0 and s.m >= hull[nh-1].m) return;
    while(nh >= 2 and !check(hull[nh-2], hull[nh-1], s)) nh--;
    pos = min(pos, nh);
    hull[nh++] = s;
}

type eval(int id, type x) { return hull[id].b + hull[id].m*x; }

// Linear Query
// queries always move to the right
/*
type query(type x){
    while(pos+1 < nh and eval(pos, x) > eval(pos+1, x)) pos++;
    return eval(pos, x);
}
*/

type query(type x){
    int l = 0, r = nh-1, mid;
    while(r - l > 5){
        mid = (l+r)/2;
        if(eval(mid+1, x) < eval(mid, x)) l = mid;
        else r = mid+1;
    }
    type mn = LINF;
    for(int i = l; i <= r; i++) mn = min(mn, eval(i, x));
    return mn;
}

```

### 6.2 Steiner Tree

```

// Steiner-Tree O(2^t*n^2 + n*3^t + APSP)

// N - number of nodes
// T - number of terminals

int n, t, dp[N][1 << T], dist[N][N];

int steiner_tree() {
    // + APSP
    memset(dp, 63, sizeof dp);
    for(int i = 1; i <= n; i++) dp[i][1 << (i-1)] = 0;

    for(int msk = 0; msk < (1 << t); msk++) {
        for(int i = 1; i <= n; i++) {
            for(int ss = msk; ss > 0; ss = (ss - 1) & msk)
                dp[i][msk] = min(dp[i][msk], dp[i][ss] + dp[i][msk - ss]);
        }
    }
}

```

```

    if(dp[i][msk] != INF)
        for(int j = 1; j <= n; j++)
            dp[j][msk] = min(dp[j][msk], dp[i][msk] + dist[i][j]);
}

int mn = INF;
for(int i = 1; i <= n; i++) mn = min(mn, dp[i][(1 << t) - 1]);
return mn;
}

```

## 7 Geometry

### 7.1 Convex Hull

```

// Convex Hull Graham's Scan Algorithm O(nlogn)

struct Point{
    int x, y;
    Point(int x = 0, int y = 0):x(x), y(y) {}
    Point operator-(Point p){ return Point(x - p.x, y - p.y); }
    Point operator+(Point p){ return Point(x + p.x, y + p.y); }
    int operator%(Point p){ return x*p.y - y*p.x; }
    int operator*(Point p){ return x*p.x + y*p.y; }
};

vector<Point> pts, hull;
Point ori;
//ori -> the highest leftmost point

bool cmp(Point a, Point b){
    if((b - ori) % (a - ori) > 0) return false;
    if((b - ori) % (a - ori) == 0 and (b - ori) * (b - ori) < (a - ori) * (a - ori)) return false;
    return true;
}

void convex_hull (vector<Point>& pts){
    sort(pts.begin() + 1, pts.end(), cmp);
    hull.pb(pts[0]);
    hull.pb(pts[1]);
    for(int i = 2; i < pts.size(); i++){
        while((hull[hull.size()-1] - hull[hull.size()-2]) % (pts[i] - hull[hull.size()-1]) <= 0) hull.
            pop_back();
        hull.pb(pts[i]);
    }
}

```

### 7.2 Minimum Enclosing Circle

```

#include <bits/stdc++.h>
using namespace std;

const double EPS = 1e-9;
const double PI = acos(-1.);

struct point{
    double x, y;
    point(): x(0.0), y(0.0) {}
    point(double x, double y): x(x), y(y) {}
    point operator +(point b){ return point(x+b.x, y+b.y); }
    point operator -(point b){ return point(x-b.x, y-b.y); }
    point operator +(double k) { return point(x+k, y+k); }
    point operator /(double k) { return point(x/k, y/k); }
    double operator %(point b) { return x*b.y - y*b.x; }
};

double dist(point p1, point p2){ return hypot(p1.x - p2.x, p1.y - p2.y); }

point rot90cw(point p1){ return point(p1.y, -p1.x); }

struct circle{
    point c;
    double r;
    circle() { c = point(); r = 0; }
    circle(point c, double r) : c(c), r(r) {}
    bool contains(point p) { return dist(c, p) <= r + EPS; }
};

```

```

circle circumcircle(point a, point b, point c) {
    point u = rot90cw(b-a);
    point v = rot90cw(c-a);
    point n = (c-b)/2;
    point ans = ((a+c)/2 + (v*((u%n)/(v%u))));
    return circle(ans, dist(ans, a));
}

```

```

// Welzl - Minimum Enclosing Circle O(n)
circle minimumCircle(vector<point> p){
    random_shuffle(p.begin(), p.end());
    circle C = circle(p[0], 0.0);
    for(int i = 0; i < p.size(); i++){
        if(C.contains(p[i])) continue;
        C = circle(p[i], 0.0);
        for(int j = 0; j < i; j++){
            if(C.contains(p[j])) continue;
            C = circle((p[j] + p[i])/2, dist(p[j], p[i])/2);
            for(int k = 0; k < j; k++){
                if(C.contains(p[k])) continue;
                C = circumcircle(p[j], p[i], p[k]);
            }
        }
    }
    return C;
}

```

## 8 Miscellaneous

### 8.1 Merge Sort

```

// Merge-sort with inversion count - O(nlog n)

int n, inv;
vector<int> v, ans;

void mergesort(int l, int r, vector<int> &v){
    if(l == r) return;
    int mid = (l+r)/2;
    mergesort(l, mid, v), mergesort(mid+1, r, v);
    int i = l, j = mid + 1, k = l;
    while(i <= mid or j <= r){
        if(i <= mid and (j > r or v[i] <= v[j])) ans[k++] = v[i++];
        else ans[k++] = v[j++], inv += j-k;
    }
    for(int i = l; i <= r; i++) v[i] = ans[i];
}

//in main
ans.resize(v.size());

```

### 8.2 Parallel Binary Search

```

//Parallel Binary Search - O(nlog^2 n + q)

struct Query{ int l, r, i, ans; };
vector<Query> req;

void pbs2(int L, int R, vector<Query> Q){
    int mid = (L+R)/2;
    // mid = (L+R+1)/2 if different from simple upper/lower bound
    if(Q.empty()) return;
    while(i < mid) {
        i++;
        //add value to data structure
    }
    while(i > mid) {
        //remove value to data structure
        i--;
    }

    if(L == R) {
        for(auto q : Q) req[q.i].ans = L;
        return;
    }
}

```

```

vector<Query> vl, vr;
for (auto q : Q) {
    if ( /* cond */ ) vl.pb(q);
    else vr.pb(q);
}
pbs(L, mid, vl);
pbs(mid + 1, R, vr);
}

```

## 8.3 Ternary Search

```

// Ternary Search - O(log n)

int ternary_search_min (int l = 1, int r = N) {
    int mid, mn = INF, idx;
    while (r - l > 5) {
        mid = (l+r)/2;
        if (func(mid + 1) < func(mid)) l = mid;
        else r = mid + 1;
    }
    for (int i = l; i <= r; i++) if (func(i) < mn) mn = func(i), idx = i;
    return idx;
}

int ternary_search_max (int l = 1, int r = N) {
    int mid, mx = -INF, idx;
    while (r - l > 5) {
        mid = (l+r)/2;
        if (func(mid + 1) > func(mid)) l = mid;
        else r = mid + 1;
    }
    for (int i = l; i <= r; i++) if (func(i) > mx) mx = func(i), idx = i;
    return idx;
}

```

## 8.4 Ternary Search with Double

```

// Ternary Search - O(log n)

double ternary_search_min (double l, double r) {
    double mid1, mid2;
    for (int i = 0; i < 100; i++) {
        mid1 = (r-l)/3 + l;
        mid2 = mid1 + (r-l)/3;
        if (func(mid2) < func(mid1)) l = mid1;
        else r = mid2;
    }
    return l;
}

double ternary_search_max (double l, double r) {
    double mid1, mid2;
    for (int i = 0; i < 100; i++) {
        mid1 = (r-l)/3 + l;
        mid2 = mid1 + (r-l)/3;
        if (func(mid2) > func(mid1)) l = mid1;
        else r = mid2;
    }
    return l;
}

```

# 9 Math Extra

## 9.1 Combinatorial formulas

$$\begin{aligned}
 \sum_{k=0}^n k^2 &= n(n+1)(2n+1)/6 \\
 \sum_{k=0}^n k^3 &= n^2(n+1)^2/4 \\
 \sum_{k=0}^n k^4 &= (6n^5 + 15n^4 + 10n^3 - n)/30 \\
 \sum_{k=0}^n k^5 &= (2n^6 + 6n^5 + 5n^4 - n^2)/12
 \end{aligned}$$

$$\begin{aligned}
 \sum_{k=0}^n x^k &= (x^{n+1} - 1)/(x - 1) \\
 \sum_{k=0}^n kx^k &= (x - (n+1)x^{n+1} + nx^{n+2})/(x - 1)^2 \\
 \binom{n}{k} &= \frac{n!}{(n-k)!k!} \\
 \binom{n}{k} &= \binom{n-1}{k} + \binom{n-1}{k-1} \\
 \binom{n}{k} &= \frac{n}{n-k} \binom{n-1}{k} \\
 \binom{n}{k} &= \frac{n-k+1}{k} \binom{n}{k-1} \\
 \binom{n+1}{k} &= \frac{n+1}{n-k+1} \binom{n}{k} \\
 \binom{n}{k+1} &= \frac{n-k}{k+1} \binom{n}{k} \\
 \sum_{k=1}^n k \binom{n}{k} &= n2^{n-1} \\
 \sum_{k=1}^n k^2 \binom{n}{k} &= (n + n^2)2^{n-2} \\
 \binom{n}{r} &= \sum_{k=0}^r \binom{n}{k} \binom{n}{r-k} \\
 \binom{n}{k} &= \prod_{i=1}^k \frac{n-k+i}{i}
 \end{aligned}$$

## 9.2 Number theory identities

**Lucas' Theorem:** For non-negative integers  $m$  and  $n$  and a prime  $p$ ,

$$\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p},$$

where

$$m = m_k p^k + m_{k-1} p^{k-1} + \cdots + m_1 p + m_0$$

is the base  $p$  representation of  $m$ , and similarly for  $n$ .

## 9.3 Stirling Numbers of the second kind

Number of ways to partition a set of  $n$  numbers into  $k$  non-empty subsets.

$$\left\{ \begin{matrix} n \\ k \end{matrix} \right\} = \frac{1}{k!} \sum_{j=0}^k (-1)^{(k-j)} \binom{k}{j} j^n$$

Recurrence relation:

$$\left\{ \begin{matrix} 0 \\ 0 \end{matrix} \right\} = 1$$

$$\left\{ \begin{matrix} n \\ 0 \end{matrix} \right\} = \left\{ \begin{matrix} 0 \\ n \end{matrix} \right\} = 1$$

$$\left\{ \begin{matrix} n+1 \\ k \end{matrix} \right\} = k \left\{ \begin{matrix} n \\ k \end{matrix} \right\} + \left\{ \begin{matrix} n \\ k-1 \end{matrix} \right\}$$

## 9.4 Burnside's Lemma

Let  $G$  be a finite group that acts on a set  $X$ . For each  $g$  in  $G$  let  $X^g$  denote the set of elements in  $X$  that are fixed by  $g$ , which means  $X^g = \{x \in X | g(x) = x\}$ . Burnside's lemma asserts the following formula for the number of orbits, denoted  $|X/G|$ :

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

## 9.5 Numerical integration

RK4: to integrate  $\dot{y} = f(t, y)$  with  $y_0 = y(t_0)$ , compute

$$k_1 = f(t_n, y_n)$$

$$k_2 = f(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_1)$$

$$k_3 = f(t_n + \frac{h}{2}, y_n + \frac{h}{2}k_2)$$

$$k_4 = f(t_n + h, y_n + hk_3)$$

$$y_{n+1} = y_n + \frac{h}{6}(k_1 + 2k_2 + 2k_3 + k_4)$$

	C	A	N	Assunto	Descricao	Diff
A						
B						
C						
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