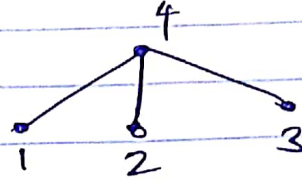


Seidel

Normal access

$$A \rightarrow \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$



$A \rightarrow$ adjacency matrix

$i, j \rightarrow$ vertices

$D_{i,j} \rightarrow$ shortest distance between vertices i & j

$$A[i][1] = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

$$A[i][2] = \begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

and so on

lu

$$\begin{bmatrix} A_{11} & A_{12} & A_{13} \\ A_{21} & A_{22} & A_{23} \\ A_{31} & A_{32} & A_{33} \end{bmatrix} = \begin{bmatrix} U_{11} & U_{12} & U_{13} \\ 0 & U_{22} & U_{23} \\ 0 & 0 & U_{33} \end{bmatrix} \begin{bmatrix} L_{11} & 0 & 0 \\ L_{21} & L_{22} & 0 \\ L_{31} & L_{32} & L_{33} \end{bmatrix}$$

in our algo,

~~for~~ $k=0$
 $j=1$
 $i=1$

$$\Rightarrow a[0][1] = a[0][1] / a[0][0]$$
$$a[1][1] = a[1][1] - (a[1][0] \times a[0][1])$$

Tiled access

The matrix is split into
 $(N/32)$ small matrices

and the same operation is performed.

where lbv is the lower triangle vertices & ubv is upper triangle vertices.

Smaller access sizes than normal access as $(N/32)$ small matrices are traversed to form the actual matrix.