



Merge Insertion

Hybrid Sort

Nathan, Shun Kah, Huai Zhi



A

Implementation

Implementation of
hybrid sort

B

Generation of input data

Generation of
random data sets

C

Analysis of time complexity

Analysing of time
complexity with
different inputs

D

Comparison with original mergesort

Comparing hybrid
sort with original
merge sort



Part A

Algorithm Implementation

Pseudo-code for Merge Sort

```
merge_sort(list):  
    find length of list  
  
    if length of list is equals to 1:  
        return the list  
    else:  
        find mid-point of list  
        merge_sort(left subarray, S)  
        merge_sort(right subarray, S)  
  
    return merge(left subarray, right subarray)
```


Pseudo-code for Hybrid Sort

```
merge_sort_hybrid(list, S):
```

```
    find length of list
```

Different base case

```
    if length of list is smaller than or equals to S:  
        perform insertion sort and return the list
```

```
    else:
```

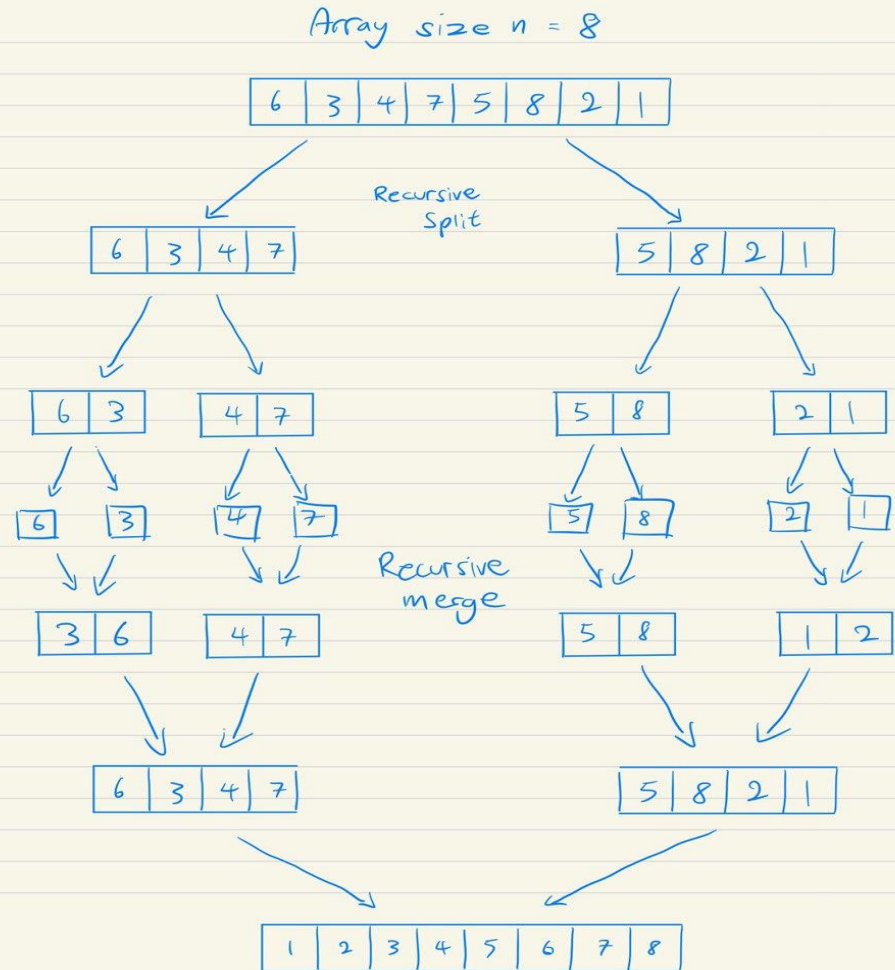
```
        find mid-point of list
```

```
        merge_sort_hybrid(left subarray, S)
```

```
        merge_sort_hybrid(right subarray, S)
```

```
    return merge(left subarray, right subarray)
```

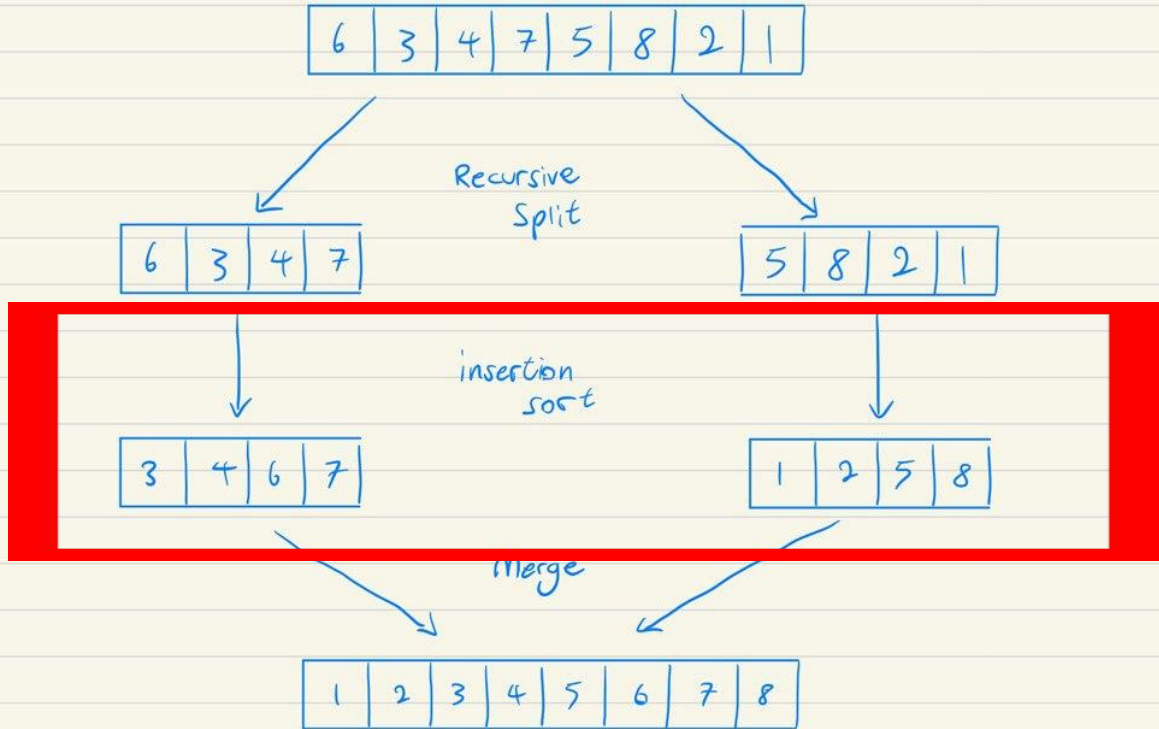
Visualisation (Merge Sort)



Visualisation (Hybrid Sort)

Insertion Sort instead of
splitting the subarray
further, when subarray
size = S

Array size $n = 8$, Cutoff $S = 4$



Code Implementation for Hybrid Sort

Array : [0]

```
def merge_sort_hybrid(list, S, key_comparisons):  
    length = len(list)  
  
    if length <= S:  # base case  
        return insertion_sort(list, key_comparisons)  
  
    mid = length // 2  
  
    left = merge_sort_hybrid(list[:mid], S, key_comparisons)[0]  
    right = merge_sort_hybrid(list[mid:], S, key_comparisons)[0]  
  
    return merge(left, right, key_comparisons)
```


Code Implementation for Hybrid Sort

```
def insertion_sort(list, key_comparisons):  
    for i in range(1, len(list)):  
        current = list[i]  
        for j in range(i, 0, -1):  
            key_comparisons[0] += 1  
            if list[j] < list[j-1]:  
                list[j] = list[j-1]  
                list[j-1] = current  
            else:  
                break  
  
    return list, key_comparisons
```

```
def merge(left, right, key_comparisons):  
    output = []  
    i = j = 0  
  
    while i < len(left) and j < len(right):  
        key_comparisons[0] += 1  
        if left[i] < right[j]:  
            output.append(left[i])  
            i += 1  
        else:  
            output.append(right[j])  
            j += 1  
  
    output.extend(left[i:])  
    output.extend(right[j:])  
  
    return output, key_comparisons
```

Sample Output

```
def main():  
    unsorted = [99, 0, 5, 20, 123, 0, -1, 72, 21, 22, 13, 8, 7, 67, 29, 1, 2, 4]  
    S = 5  
    comparisons = merge_sort_hybrid(unsorted, S,[0])  
    print(unsorted)  
    print(comparisons)
```

```
main()
```

```
[99, 0, 5, 20, 123, 0, -1, 72, 21, 22, 13, 8, 7, 67, 29, 1, 2, 4]  
([-1, 0, 0, 1, 2, 4, 5, 7, 8, 13, 20, 21, 22, 29, 67, 72, 99, 123], [58])
```

Tuple: (Sorted Array , Number of Key Comps)



Part B

Generating input datasets

```
from numpy import random
```

```
def rng():  
    inputs = []  
    inputs.append(random.randint(1000, size=1000))  
    for i in range(1, 5):  
        size = 1000 * pow(10, i)  
        inputs.append(random.randint(size, size=int(size/2)))  
        inputs.append(random.randint(size, size=size))  
  
    return inputs
```

```
inputs = rng()  
for i in inputs:  
    print(len(i))
```

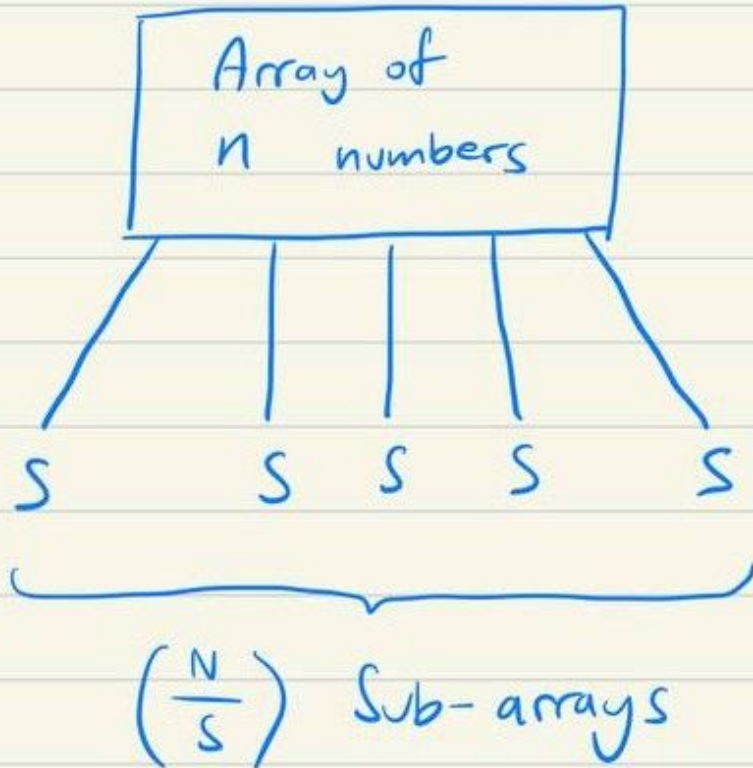
1k, 5k, 10k, 50k, 100k, 500k, 1m, 5m, 10m



Part C

Time complexity analysis

Theoretical Analysis of Time Complexity (Explanation)



Insertion Sort:

$$(N/S) * O(S) = O(N)$$

[BEST CASE]

$$(N/S) * O(S^2) = O(NS)$$

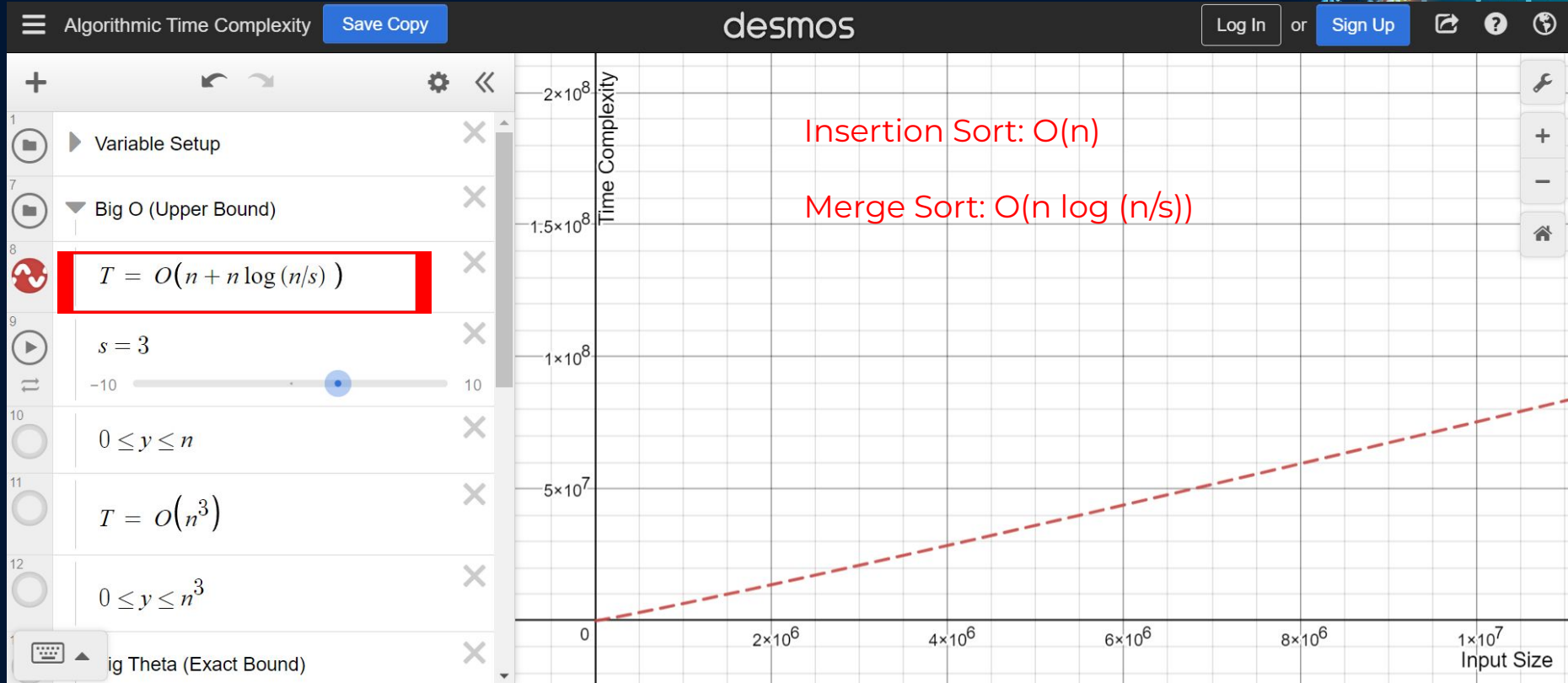
[WORST/AVG]

Merge Sort:

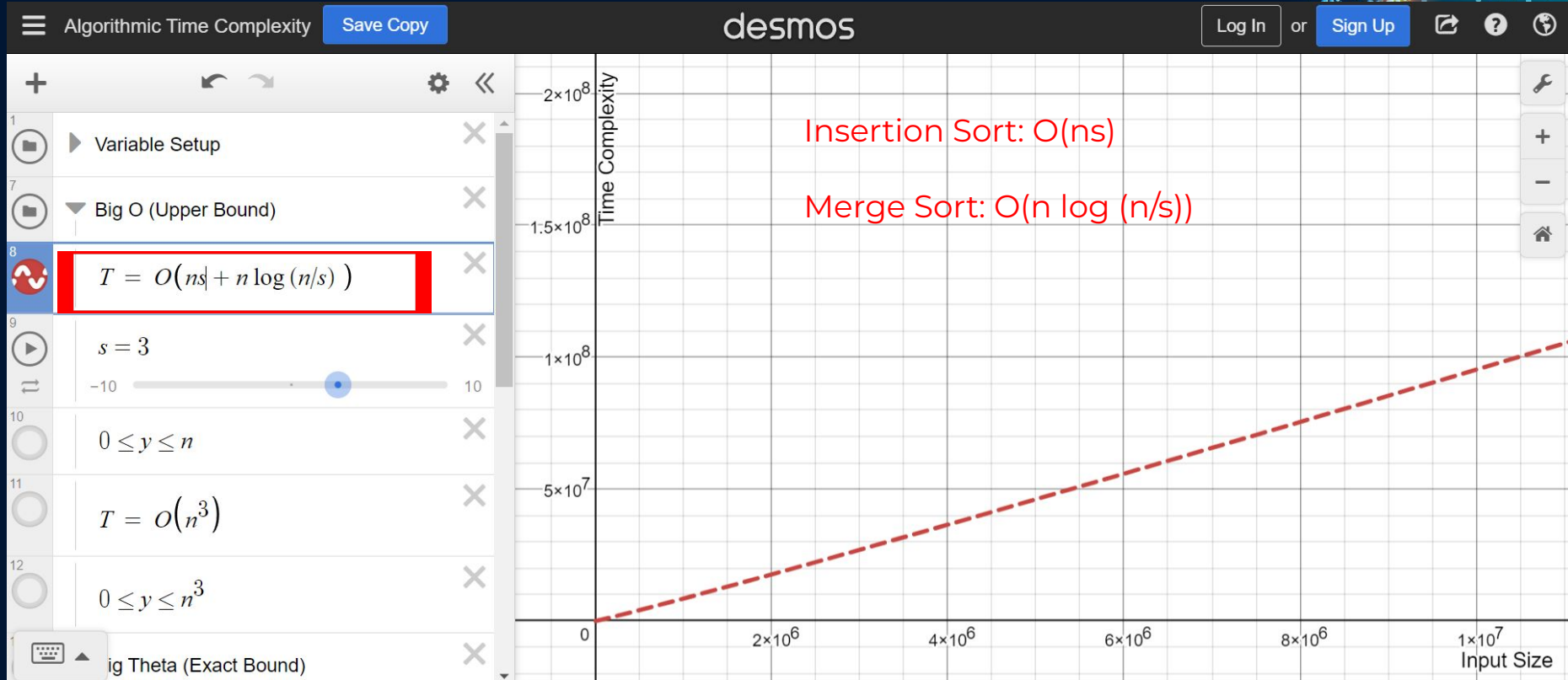
$$O(N \log (N/S))$$

[ALL CASES]

C(i): Theoretical Analysis of Time Complexity (best case)



C(i): Theoretical Analysis of Time Complexity (worst/avg case)



C(ii): Theoretical Analysis of Time Complexity

As $S \uparrow$

Insertion Sort \uparrow

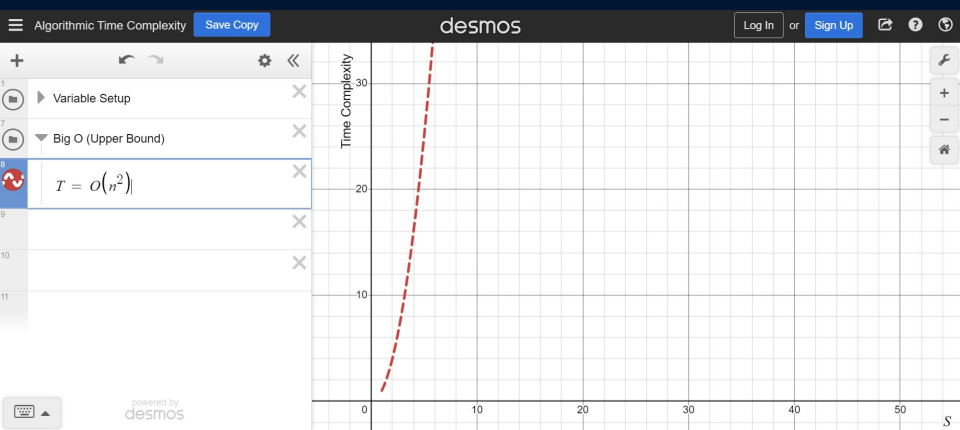
Merge Sort \downarrow

Time complexity will resemble
Insertion Sort

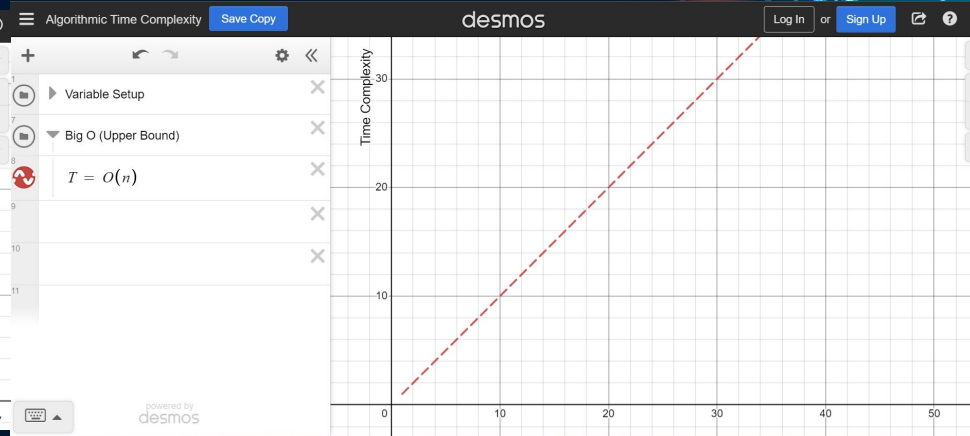


C(ii): Theoretical Analysis of Time Complexity

Worst/ Avg case

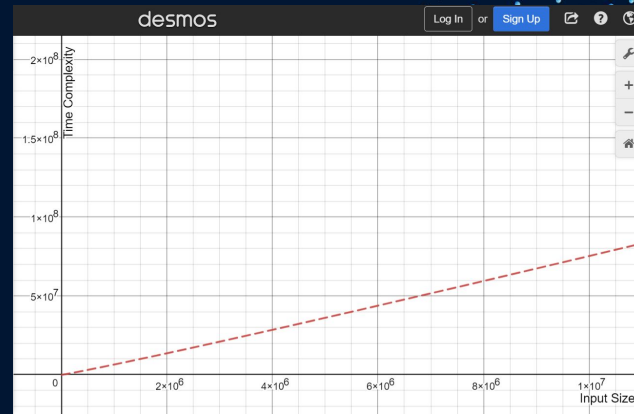
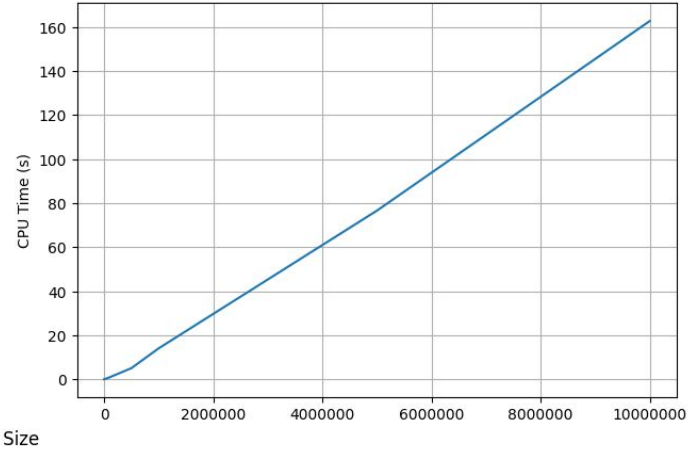
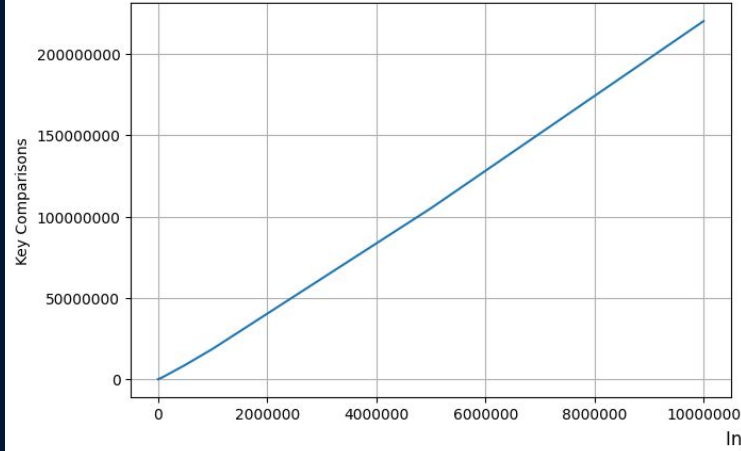


Best case



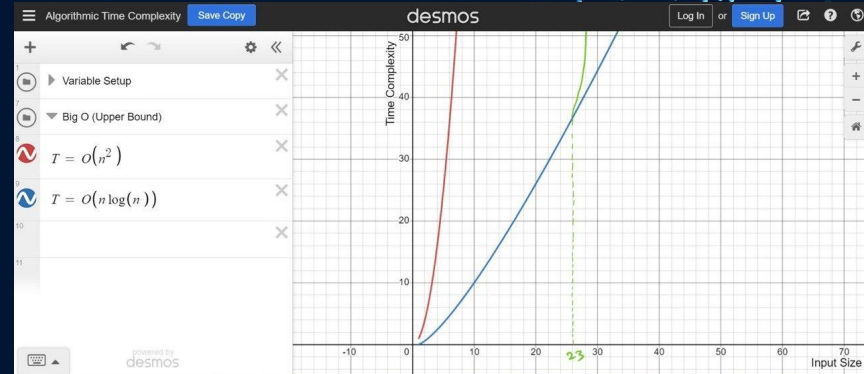
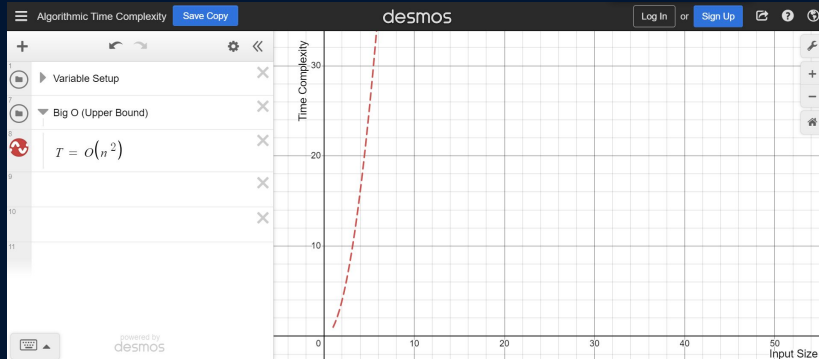
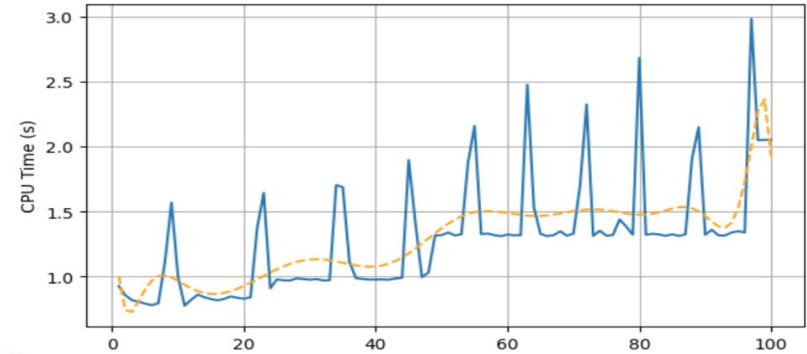
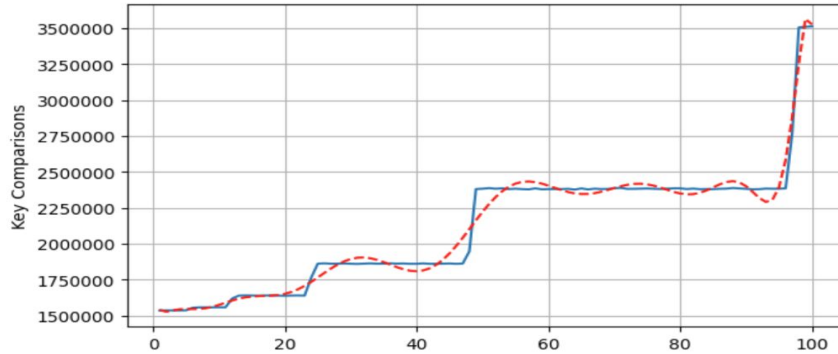
C(i): Empirical vs theoretical

Merge + Insertion Sort



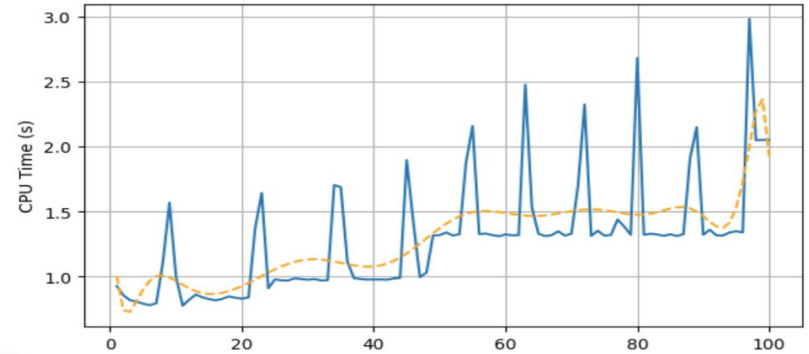
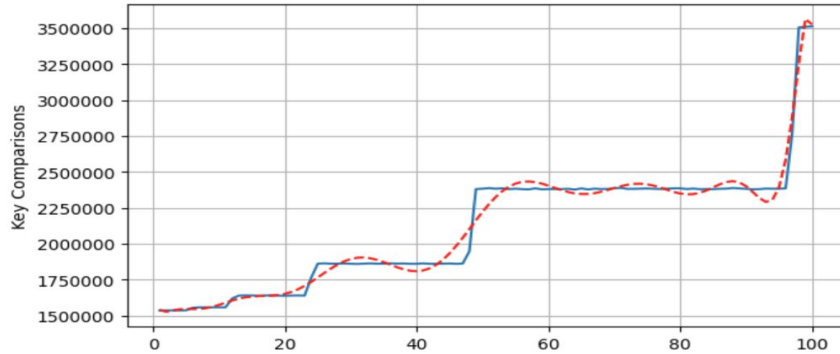
C(ii): Empirical vs theoretical

n = 100000

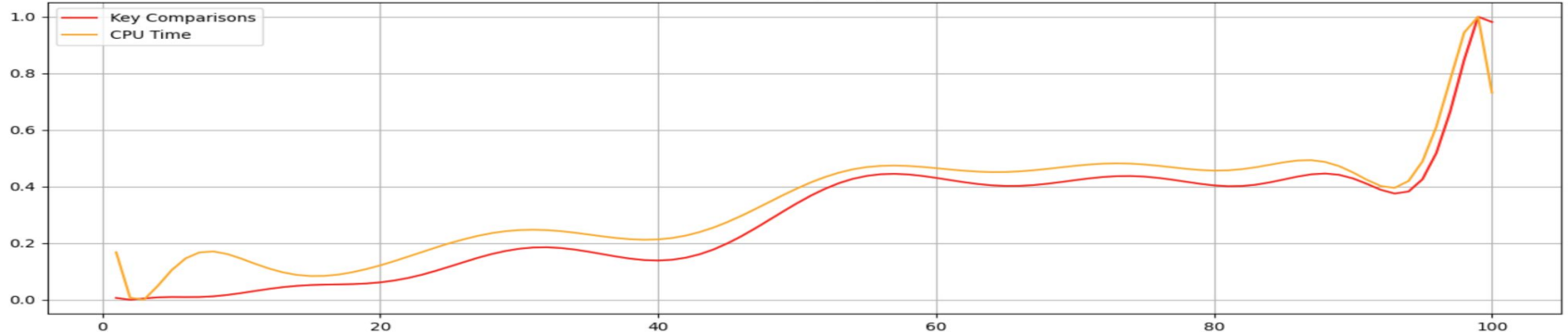


C(ii): Empirical vs theoretical

n = 100000



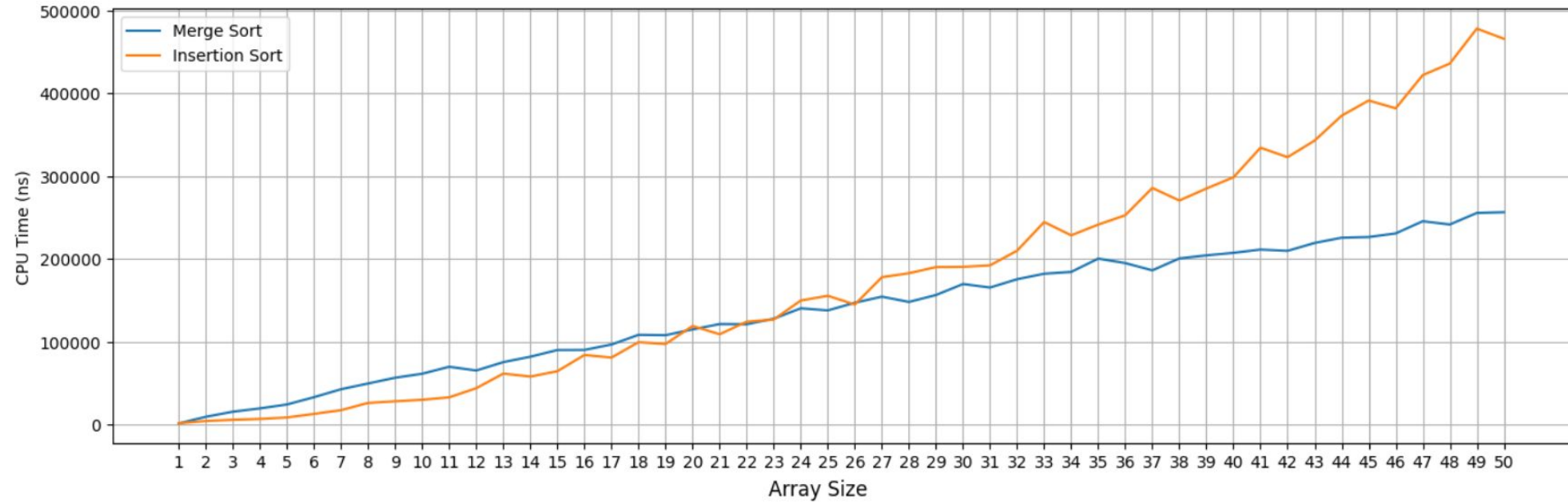
Normalised Trend for Key Comp and CPU Time



Correlation coefficient: 0.97959

C(iii): Finding Optimal S Value

Merge & Insertion Sort Comparison



C(iii): Finding Optimal S-value

Based on the intersection between the insertion and merge sort algorithm, where the hybrid sort will work best as insertion sort before the optimal S-value, and will work best as merge sort after the optimal S-value. The optimal S-value should be the intersection, which is 23.



Part D

Comparison with original merge sort

Pseudo-code for Merge Sort

```
merge_sort(list):  
    find length of list  
  
    if length of list is equals to 1:  
        return the list  
    else:  
        find mid-point of list  
        merge_sort(left subarray, S)  
        merge_sort(right subarray, S)  
  
    return merge(left subarray, right subarray)
```

Merge Sort Implementation

```
def merge_sort(list, key_comparisons):  
    length = len(list)  
  
    if length == 1: # base case  
        return list, key_comparisons  
  
    mid = length // 2  
  
    left = merge_sort(list[:mid], key_comparisons)[0]  
    right = merge_sort(list[mid:], key_comparisons)[0]  
  
    return merge1(left, right, key_comparisons)
```

```
def merge1(left, right, key_comparisons):  
    output = []  
    i = j = 0  
  
    while i < len(left) and j < len(right):  
        key_comparisons[0] += 1  
        if left[i] < right[j]:  
            output.append(left[i])  
            i += 1  
        else:  
            output.append(right[j])  
            j += 1  
  
    output.extend(left[i:])  
    output.extend(right[j:])  
  
    return output, key_comparisons
```

Program for comparison

```
import numpy as np
import matplotlib.pyplot as plt
from time import process_time, process_time_ns

# unsorted = [99, 0, 5, 20, 123, 0, -1, 72, 21, 22, 13, 8, 7, 67, 29, 1, 2, 4]
input = random.randint(1000, size=10000000)
# print(input)
merge_start = process_time_ns()
comparisons = merge_sort(input, [0])
merge_stop = process_time_ns()
# print(merge_stop, merge_start)
merge_time = merge_stop - merge_start
print("Merge Sort Time:", merge_time, "ns")

S = 23
hybrid_start = process_time_ns()
sortResult = merge_sort_hybrid(input, S, [0])
hybrid_stop = process_time_ns()
hybrid_time = hybrid_stop - hybrid_start
print("Hybrid Sort Time:", hybrid_time, "ns")
```

Logs

1. Key Comparisons
2. CPU Time

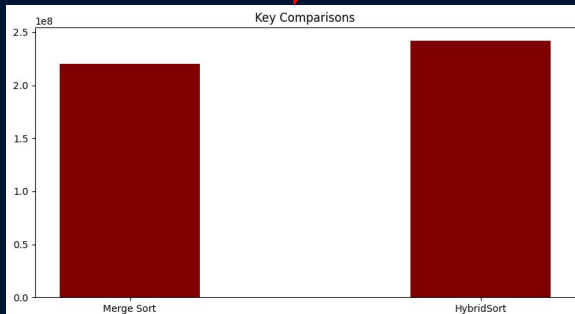
Bar Graph code

```
# creating the dataset
data = {'Merge Sort':comparisons[1][0], 'HybridSort':sortResult[1][0]}
courses = list(data.keys())
values = list(data.values())

fig = plt.figure(figsize = (10, 5))

# creating the bar plot
plt.bar(courses, values, color = 'maroon',
        width = 0.4)

plt.xlabel("")
plt.ylabel("")
plt.title("Key Comparisons")
plt.show()
```

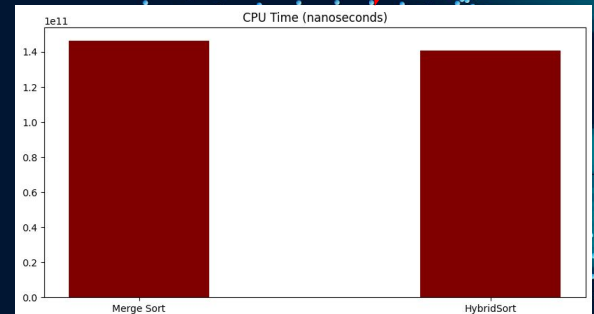


```
# creating the dataset
data = {'Merge Sort':merge_time, 'HybridSort':hybrid_time}
courses = list(data.keys())
values = list(data.values())

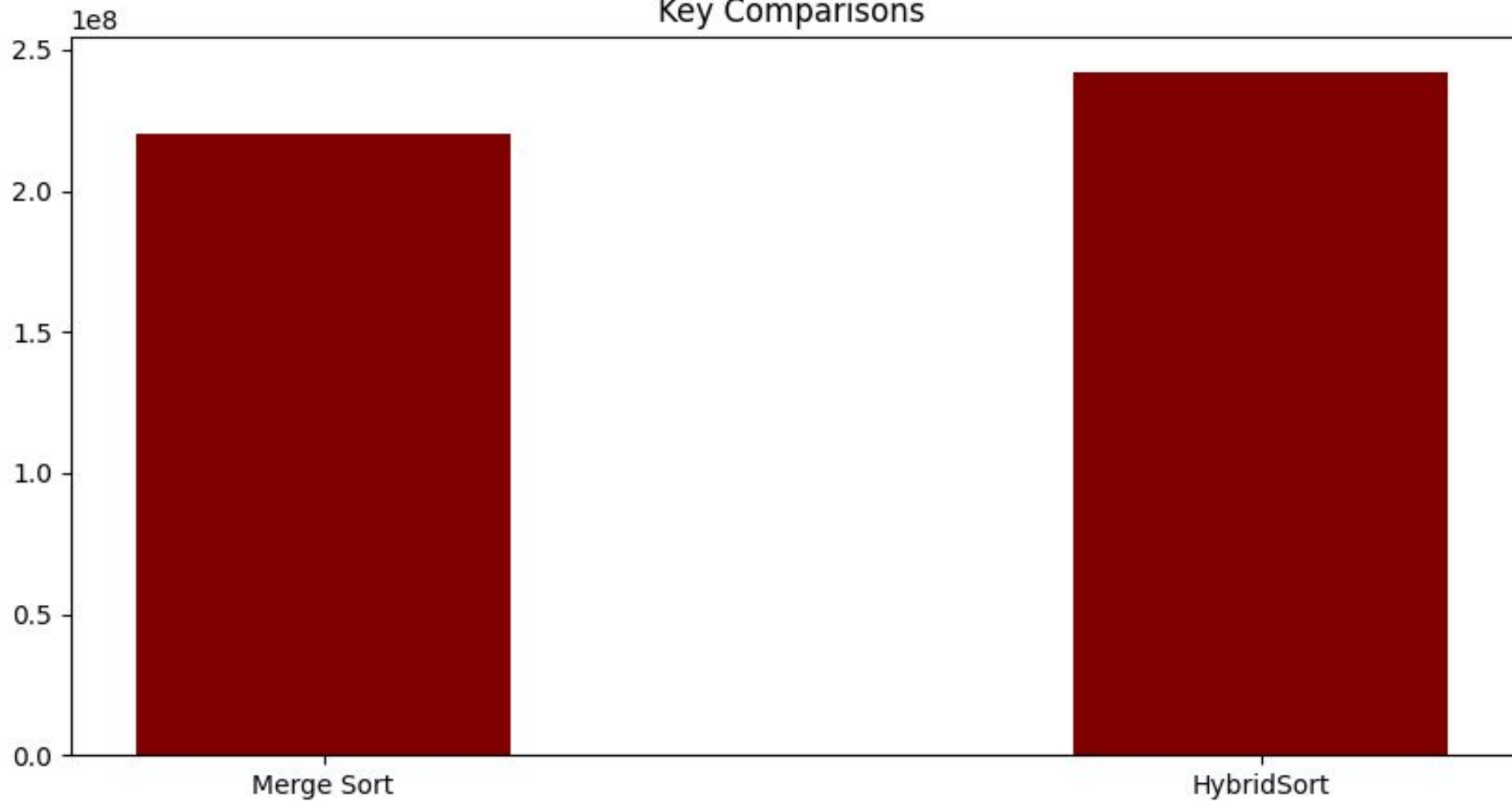
fig = plt.figure(figsize = (10, 5))

# creating the bar plot
plt.bar(courses, values, color = 'maroon',
        width = 0.4)

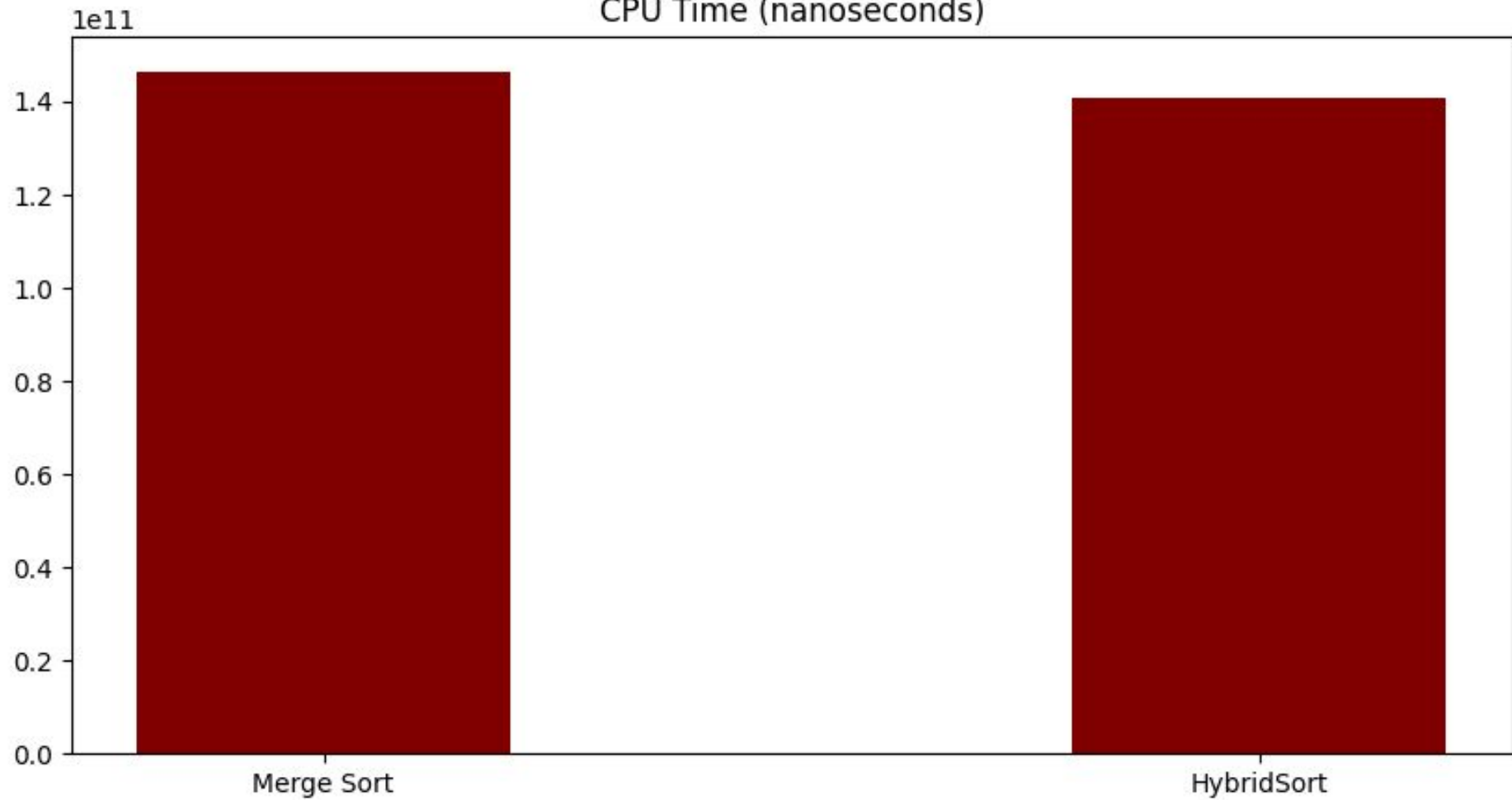
plt.xlabel("")
plt.ylabel("")
plt.title("CPU Time (nanoseconds)")
plt.show()
```



Key Comparisons



CPU Time (nanoseconds)



Interpretation

Since insertion sort typically uses more key comparison than merge sort, when the hybrid sort starts to follow the pattern of insertion sort, its key comparison will be more than if merge sort is used.

Insertion sort worst key comparison

$$\sum_{i=1}^{n-1} i = \frac{(n-1)n}{2}$$

Merge sort worst key comparison

at most (worst case) $n - 1$

Interpretation

Insertion sort is faster on smaller array of numbers as it skips the sorted values. Thus when the hybrid model is used on the smaller array of numbers and insertion sort is used, the cpu time will be faster than the merge sort model as merge sort will still go through the sorted values.



The background features a dark teal color with several wavy, dotted lines in a lighter teal shade. A bright orange and yellow glow emanates from the center-left, creating a lens flare effect.

Thank you!
