

1.

A. Virtual address format

13	12	11	10	9	8	7	6	5	4	3	2	1	0
0	0	0	0	1	0	0	1	1	1	1	1	0	0

B. Address translation

Parameter	Value
VPN	0x09
TLB index	0x1
TLB tag	0x02
TLB hit?(Y/N)	N
Page fault?(Y/N)	N
PPN	0x17

C. Physical address format

11	10	9	8	7	6	5	4	3	2	1	0
0	1	0	1	1	1	1	1	1	1	0	0

D. Physical memory reference

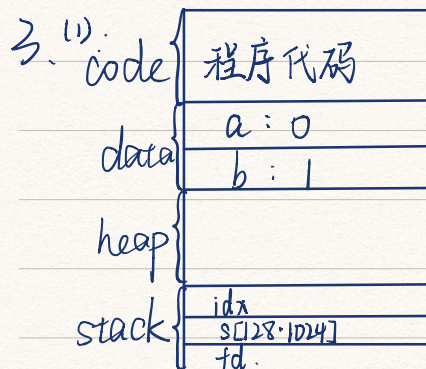
Parameter	Value
Byte offset	0x0
Cache index	0x0F
Cache tag	0x17
Cache hit?(Y/N)	N
Cache byte returned	—

2. $\text{page size} = 2048 \text{ B} = 2^{11} \text{ B}$

$64 - 11 = 53$, 故共有 2^{53} 个 page

一页中可以装 $2048 \div 4 = 2^9$ 个页表项

而 $(2^9)^6 > 2^{53}$, 所以至少需要 6 级 levels 才能将所有页面放在表中.



12) ① $S = 0xFFC00 = 1111\ 1111\ 1100\ 0000\ 0000$

segment = 11, 位于 stack 段

Offset = 11 1111 1100 0000 0000 = 255k

VA 中 stack 段占 $1MB/4 = 256KB$

\therefore 实际的负偏移量 = $256k - 255k = 1k$.

$\therefore PA = 1MB - 1k = 1023k$

②. $VA = 258K = 0100\ 0000\ 1000\ 0000\ 0000$

segment = 01, 位于 data 段,

\therefore Offset = 00 0000 1000 0000 0000 = 2k

\therefore data 段 size 为 2KB, 所以内存访问越界.

③. $VA = 514K = 1000\ 0000\ 1000\ 0000\ 0000$

segment = 10, 位于 heap 段

\therefore Offset = 1000 0000 0000 = 2k < 8k

$\therefore PA = 7MB + 2KB = 7170k$