



## Tentamen 28 Maart 2019, antwoorden

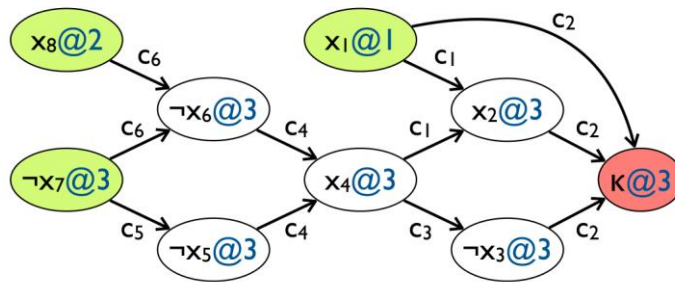
Knowledge Representation (Vrije Universiteit Amsterdam)

## Examination for Knowledge Representation, 28 March 2009

### 1. Propositional logic Logic (10pt)

- 0
- Somewhere between 1 and  $2^k$  (or  $(2^k)-1$  if F is satisfiable but not valid)
- 0 (inconsistent = unsatisfiable)
- $2^k$

### 2. SAT solving: clause Learning (10pt)



(some students applied pure-literal simplification, and then there is no conflict. For those students the question is excluded from the exam)

### 3. SAT solving: GSAT (15pt)

```

procedure GSAT(Sigma)
  for i := 1 to MAX-TRIES ; These are the restarts
    T := random(Sigma) ; random assignment
    for j := 1 to MAX-FLIPS ; To ensure termination
      if T satisfies Sigma then return T
      else T := T with variable flipped to maximize
        number of satisfied clauses
    ; It doesn't matter if the number does
    ; not increase. This are the sideways flips
  end
end

```

- True (each variable always has a value),  
alternatively:  
False, because not each variable gets *changed* at each iteration
- False
- False (it is sound, but not complete)
- True
- To avoid exploring forever parts of the search space that do not contain a solution
- To make sure you explore the whole local search space before jumping elsewhere (and miss a nearby solution)
- To void wandering for too long in parts of the search space that do not contain a solution

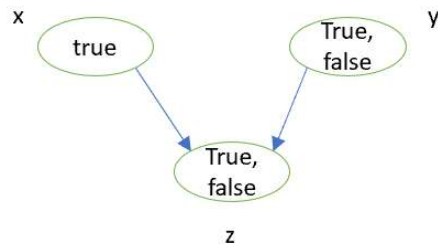
### 4. Description Logic (10pt)

- If you know at least one actor then all your friends are envious
- Happy parents are precisely those parents all of whose children are either doctors or have a child who is a doctor

### 5. Description Logic (10pt)

- {Alice,Bob}
- {Bob}
- {Clair}
- {Clair}
- {}

## 6. CSP (10pt)



- For every value of x (true) there is some value of z that is consistent with  $z = x \wedge y$  (in fact, both values of z might be consistent with  $x = \text{true}$ )

$Y = \text{true}$  is consistent with both  $z = \text{true}$  and  $z = \text{false}$

$Y = \text{false}$  is consistent with  $z = \text{false}$

So the problem is directionally arc-consistent

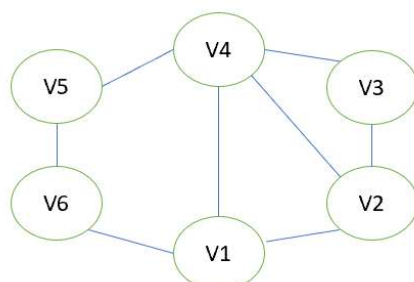
Vice versa:  $z = \text{true}$  is consistent with  $x = \text{true}$  and with  $y = \text{true}$ ,  
 $z = \text{false}$  is obviously consistent with  $y = \text{false}$ , as well as with  $x = \text{true}$   
 So the problem is also bi-directionally arc-consistent

The pair  $\{x = \text{true}, y = \text{true}\}$  is consistent with  $z = \text{true}$   
 The pair  $\{x = \text{true}, y = \text{false}\}$  is consistent with  $z = \text{false}$   
 The pair  $\{x = \text{true}, z = \text{true}\}$  is consistent with  $y = \text{true}$   
 The pair  $\{x = \text{true}, z = \text{false}\}$  is consistent with  $y = \text{false}$

So, any pair of consistent values can be extended to a consistent third value in the constraint, hence the problem is also hyper-arc consistent

- Yes. Simply give each variable the domain  $\{\text{true}, \text{false}\}$ , the constraint says that from each clause, at least one literal has to be true.
- No, CSPs can also be over very different domains (e.g. real numbers) and can have very different constraints (eg inequalities).

## 7. CSP (10pt)



- All variables initially have domain  $\{1, 2, 3, 4\}$

The least constraining values will be for V3, V5, V6, so because of the tie-breaking constraint we choose  $V3 = 1$ . This causes the reductions  $V4 = \{2, 3, 4\}$  and  $V2 = \{2, 3, 4\}$ .

The next least constraining values is now V5, and the tie-breaker causes us to choose V5=1, causing V6={2,3,4}. We can now choose V1=1 without causing any further value constraints. At this point, only the  $v2 \neq v4$  constraint is unsatisfied. The tie-breaker chooses V2=2.

This gives the solution

V1=1,  
V2={2},  
V3=1,  
V4={3,4}  
V5=1,  
V6={2,3,4}

- b. V4 is the most constraining variable, so pick V4={1}. This causes V1={2,3,4}, V2={2,3,4}, V3={2,3,4} and V5={2,3,4}. Now V1 and V2 are most constraining, so pick V1=2. This causes V2={3,4} and V6={1,3,4}. Now V2 is most constraining, so pick V2={3}. This causes V3={2,4}. Now V3, V5 and V6 are most constraining, so pick V3={2} and then V5={2} leading to the solution:

V1=2  
V2=3  
V3=2  
V4=1  
V5=2  
V6=1,3,4

## 8. QR (10pt)

**Influence:** The magnitude of the causing quantity determines the derivative of the affected quantity.

**Proportionality:** Directed propagation of derivative information from the causing quantity to the affected quantity.

## 9. QR (15pt)

### a. Sim2 (5 states)

b.

State	Q1:M	Q1:∂	Q2:M	Q2:∂
1	V1	-	P2	-
2	V1	-	V2	-
3	0	0	0	0
4	V1	-	0	0
5	0	0	V2	-

First a point-value that changes (from S1 to S2), from S2 there are 3 options (Q1 goes to 0 (S5) or Q2 goes to 0 (S4), or both (S3)), after that from S4 and S5 the other quantity also goes to 0.