

# Modulator Design for Binary Classification of Poisson Measurements

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#### Introduction

#### What we want to do:

Poisson channel  $\mathbf{y} \sim \mathsf{pois}(\mathbf{A}\mathbf{x}), \mathbf{y} \in \mathbb{Z}_+^m, \mathbf{A} \in \mathbb{R}_+^{m \times n}, \mathbf{x} \in \mathbb{R}_+^n$ 

Two hypotheses  $H_1: \mathbf{x} = \mathbf{u}$ 

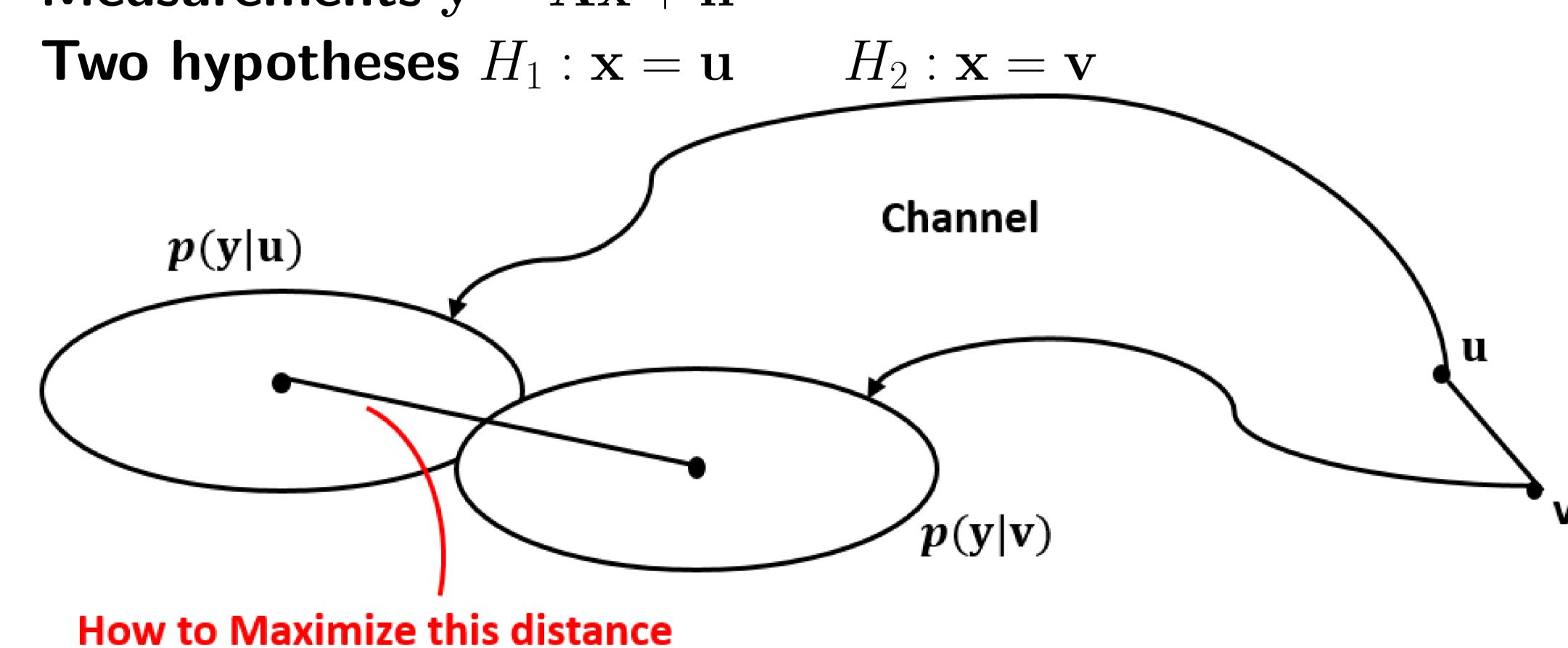
 $H_2: \mathbf{x} = \mathbf{v}$ 

Classification Given y, decide x as u or v without recovering it Goal design an improved system yielding smaller classification error Keyword: Binary Classification, Bayesian Classifier, Diversity, Classification Error

#### Motivation

#### Consider a Gaussian Channel:

Measurements y = Ax + n



**Solution** Rotate  $\mathbf{u} - \mathbf{v}$  to the dominant eigenvector of  $\mathbf{A}^{\dagger} \mathbf{A}$ Back to Poisson case:

**Idea** Design a modulator  ${f R}$  that "rotates"  ${f u}-{f v}$  appropriately designed system  $y \sim pois(ARx)$ 

## Problem Setup

#### Error of optimal classifier

Hypothesis testing:

$$H_1: \mathbf{y} \sim \mathsf{pois}(\mathbf{ARu})$$
 $H_2: \mathbf{y} \sim \mathsf{pois}(\mathbf{ARv})$ 

Bayesian(optimal) classifier:

$$\frac{\mathsf{pois}(\mathbf{y}|\mathbf{A}\mathbf{R}\mathbf{u})}{\mathsf{pois}(\mathbf{y}|\mathbf{A}\mathbf{R}\mathbf{v})} \overset{H_1}{\geqslant} 1$$

Classification error:

$$P_{e} = \sum_{\mathbf{y}} \min\{\mathsf{pois}(\mathbf{y}|\mathbf{A}\mathbf{R}\mathbf{u}), \mathsf{pois}(\mathbf{y}|\mathbf{A}\mathbf{R}\mathbf{v})\}$$

$$\leq \exp\{-\frac{1}{2}\sum_{i}\left(\sqrt{(\mathbf{A}\mathbf{R}\mathbf{u})_{i}} - \sqrt{(\mathbf{A}\mathbf{R}\mathbf{v})_{i}}\right)^{2}\}$$

$$\triangleq P^{UB}$$

#### Formulation

$$\max_{\mathbf{R}} \sum_{i} \left( \sqrt{(\mathbf{A}\mathbf{R}\mathbf{u})_{i}} - \sqrt{(\mathbf{A}\mathbf{R}\mathbf{v})_{i}} \right)^{2}$$
s.t.  $0 \le R_{i,j} \le 1$  (P0)

#### A relaxed problem

#### Difficulty of problem P0

Convex maximization: optima occurs at vertex Enumerating: infeasible, complexity  $O(2^{n^2})$ 

A relaxed problem

Assmue  $(\mathbf{ARu})_i + (\mathbf{ARv})_i \leq C, \ \forall i = 1, \dots, m$ 

Relaxed error upper bound

$$\tilde{P}_e^{UB} = \exp\left\{-\frac{1}{4C} \|\mathbf{A}\mathbf{R}\mathbf{u} - \mathbf{A}\mathbf{R}\mathbf{v}\|^2\right\}$$

Relaxed problem

$$\max_{\mathbf{R}} \|\mathbf{A}\mathbf{R}\mathbf{u} - \mathbf{A}\mathbf{R}\mathbf{v}\|^2$$
s.t.  $0 \le R_{i,j} \le 1$  (P1)

Proposition Denote  $\mathbf{r} = \text{vec}(\mathbf{R}^{\top})$ ,  $\boldsymbol{\gamma} = \text{vec}(\mathbf{R})$ ,  $\mathbf{d} = \mathbf{u} - \mathbf{v}$  $\|\mathbf{A}\mathbf{R}\mathbf{u} - \mathbf{A}\mathbf{R}\mathbf{v}\|^2 = \mathbf{r}^{\top} [(\mathbf{A}^{\top}\mathbf{A}) \otimes (\mathbf{d}\mathbf{d}^{\top})] \mathbf{r} = \boldsymbol{\gamma}^{\top} [(\mathbf{d}\mathbf{d}^{\top}) \otimes (\mathbf{A}^{\top}\mathbf{A})] \boldsymbol{\gamma}$ 

P1 as a fixed rank 0-1 quadratic maximization problem Search the vertices of a zonotope, complexity  $O(n^{2rank(\mathbf{D})})$ P1 as a mincut problem

Let  $\rho = 2\mathbf{r} - \mathbf{1}$ , then  $\rho_i = \pm 1$ 

$$\mathbf{r}^{\top}\mathbf{D}\mathbf{r} = \frac{1}{4} \underbrace{\begin{bmatrix} \mathbf{1}\boldsymbol{\rho}^{\top} \end{bmatrix}}_{\mathbf{s}^{\top}} \underbrace{\begin{bmatrix} \mathbf{1}^{\top}\mathbf{D}\mathbf{1}\mathbf{1}^{\top}\mathbf{D} \\ \mathbf{D}\mathbf{1} & \mathbf{D} \end{bmatrix}}_{\mathbf{G}} \begin{bmatrix} 1 \\ \boldsymbol{\rho} \end{bmatrix} = -\frac{1}{2} \sum_{i,j:s_i=-s_j} G_{i,j} + \text{const.}$$

Nontrivial if there exists any negative edge,  $G_{i,j} < 0$ 

# Even all edges are nonnegative, complexity is at least $O(n^4 \log^3 n)$

# Fast Algorithm

Decompose the objective w.r.t.  ${f R}$ 's rows:

$$\mathbf{r}^{\mathsf{T}}\mathbf{D}\mathbf{r} = \mathbf{r}_{1}^{\mathsf{T}}[\|\mathbf{a}_{1}\|^{2}\mathbf{d}\mathbf{d}^{\mathsf{T}}]\mathbf{r}_{1} + (\mathbf{r}_{2}^{\mathsf{T}}[\|\mathbf{a}_{2}\|^{2}\mathbf{d}\mathbf{d}^{\mathsf{T}}]\mathbf{r}_{2} + 2[(\mathbf{a}_{1}^{\mathsf{T}}\mathbf{a}_{2})\mathbf{r}_{1}^{\mathsf{T}}\mathbf{d}\mathbf{d}^{\mathsf{T}}]\mathbf{r}_{2}) + \dots$$

#### **Algorithm 1** Fast Algorithm for problem P1

```
Input: Two hypothetical inputs: \mathbf{u}, \mathbf{v}. Sensing matrix \mathbf{A} = [\mathbf{a}_1, \dots, \mathbf{a}_n]
Output: binary matrix \mathbf{R} = [\mathbf{r}_1, \dots, \mathbf{r}_n]^{\top}
     \mathbf{d} = \mathbf{u} - \mathbf{v}
     if \mathbf{1}^{\mathsf{T}}\mathbf{d} > 0 then
         \mathbf{r}_1 \leftarrow \frac{1}{2}[\operatorname{sign}(\mathbf{d}) + 1]
         \mathbf{r}_1 \leftarrow \frac{1}{2}[\operatorname{sign}(-\mathbf{d}) + 1]
     end if
    for k=2,\ldots,n do
         if (\|\mathbf{a}_k\|^2 \mathbf{1}^\top + 2 \sum_{i=1}^{k-1} (\mathbf{a}_i^\top \mathbf{a}_k) \mathbf{r}_i^\top) \mathbf{d} > 0 then
               \mathbf{r}_k \leftarrow \frac{1}{2}[\operatorname{sign}(\mathbf{d}) + 1]
          \mathbf{else}
              \mathbf{r}_k \leftarrow \frac{1}{2}[\operatorname{sign}(-\mathbf{d}) + 1]
          end if
     end for
```

#### Complexity $O(n^3)$

### **Extension: Adding Flux-preserving Constraint**

$$\max_{\mathbf{R}} \|\mathbf{A}\mathbf{R}\mathbf{u} - \mathbf{A}\mathbf{R}\mathbf{v}\|^2$$
s.t.  $0 \le R_{i,j} \le 1, \sum_{i} R_{i,j} \le 1$  (P2)

Decompose the objective w.r.t.  $\mathbf{R}$ 's columns:

$$\boldsymbol{\gamma}^{\mathsf{T}} \mathbf{E} \boldsymbol{\gamma} = d_1^2 \boldsymbol{\gamma}_1^{\mathsf{T}} \mathbf{A}^{\mathsf{T}} \mathbf{A} \boldsymbol{\gamma}_1 + (d_2^2 \boldsymbol{\gamma}_2^{\mathsf{T}} \mathbf{A}^{\mathsf{T}} \mathbf{A} \boldsymbol{\gamma}_2 + 2(d_1 d_2) \boldsymbol{\gamma}_1^{\mathsf{T}} (\mathbf{A}^{\mathsf{T}} \mathbf{A}) \boldsymbol{\gamma}_2) + \dots$$

Algorithm 2 Fast Algorithm for problem P2

**Input:** Two hypothetical inputs:  $\mathbf{u}, \mathbf{v}$ . Sensing matrix  $\mathbf{A} = [\mathbf{a}_1, \dots, \mathbf{a}_n]$ Output: binary matrix  $\mathbf{R} = [\gamma_1, \dots, \gamma_n]$ 

 $\mathbf{d} = \mathbf{u} - \mathbf{v}$ 

 $\gamma_{1,i} = 1 \text{ where } i = \arg\max_i \|\mathbf{a}_i\|^2$ 

for  $i=2,\ldots,n$  do

 $\gamma_i = \mathbf{e}_k$  where  $\mathbf{e}_k$  is the k-th natural basis such that

$$k = \arg \max_{k} d_i^2 \mathbf{e}_k^{\top} \mathbf{A}^{\top} \mathbf{A} \mathbf{e}_k + 2 \left[ \sum_{j=1}^{i-1} (d_j d_i) \boldsymbol{\gamma}_j^{\top} (\mathbf{A}^{\top} \mathbf{A}) \right] \mathbf{e}_k$$

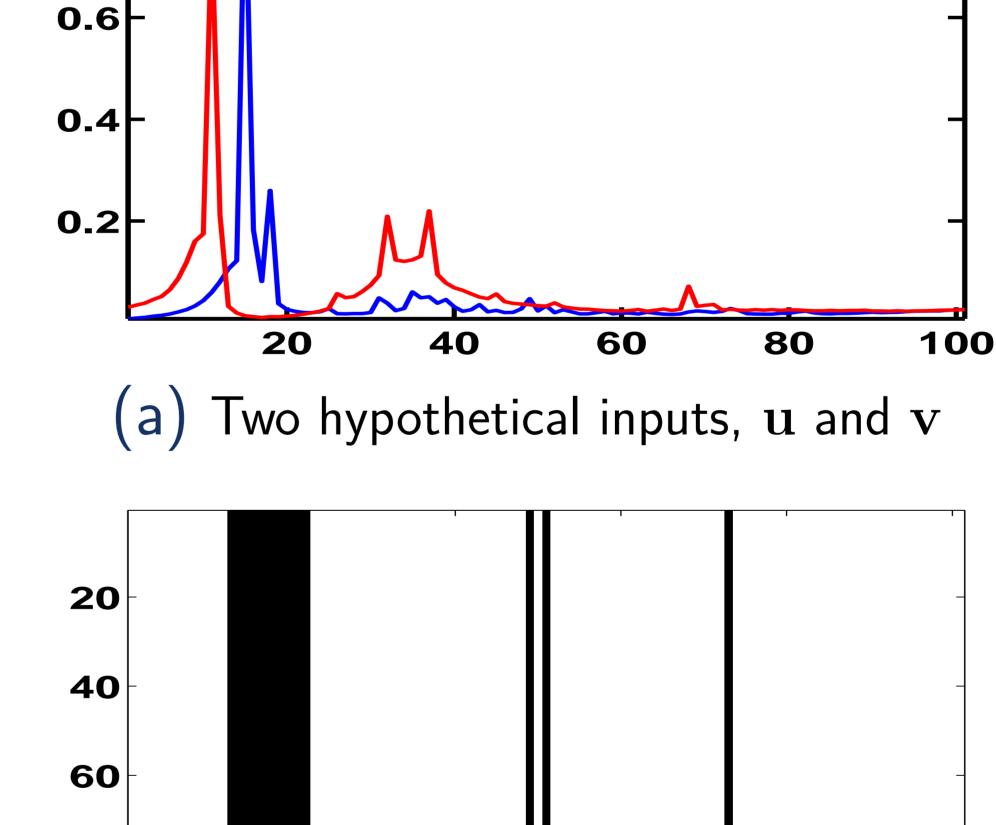
Example

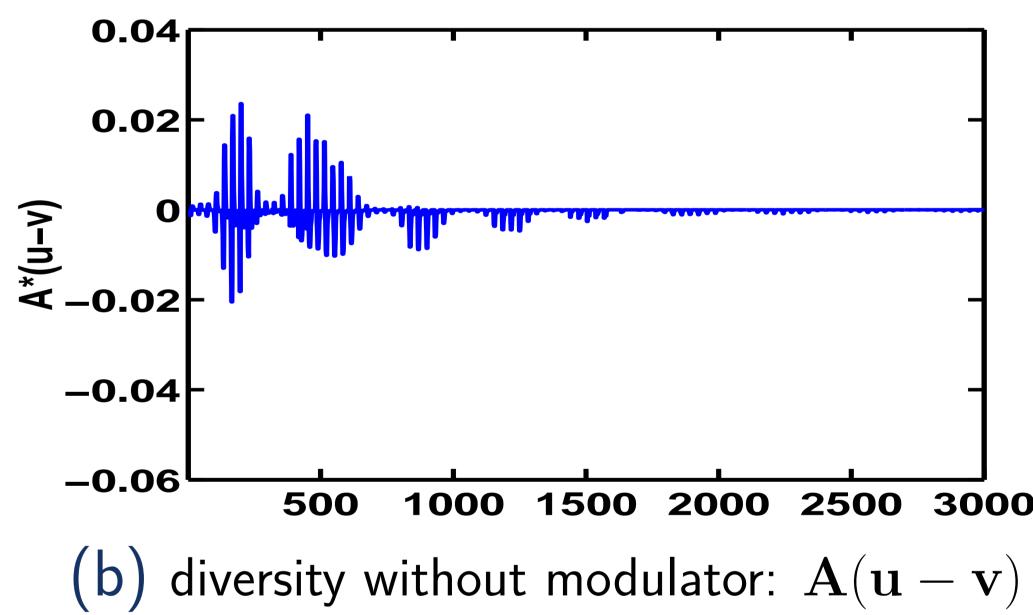
# Complexity $O(n^3)$

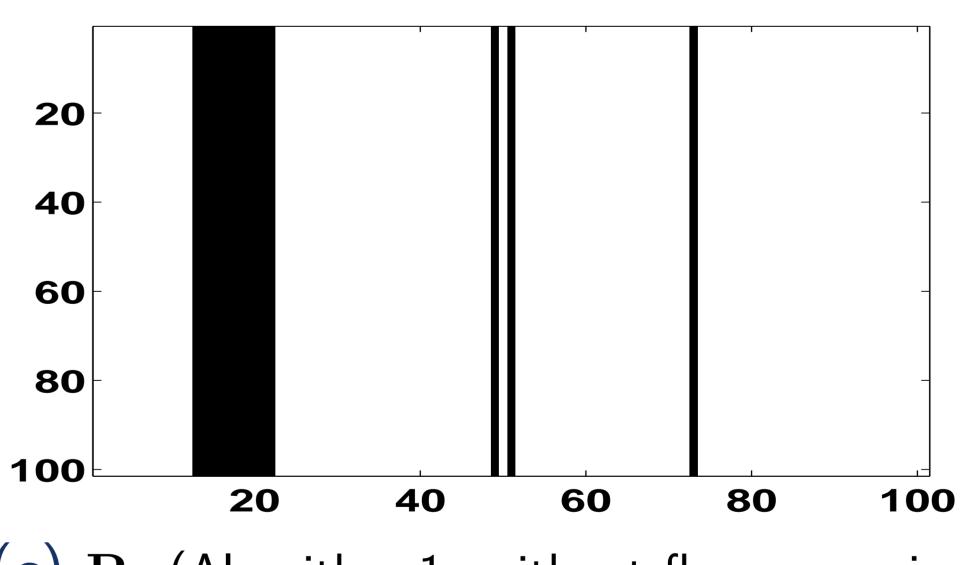
-HDPE |

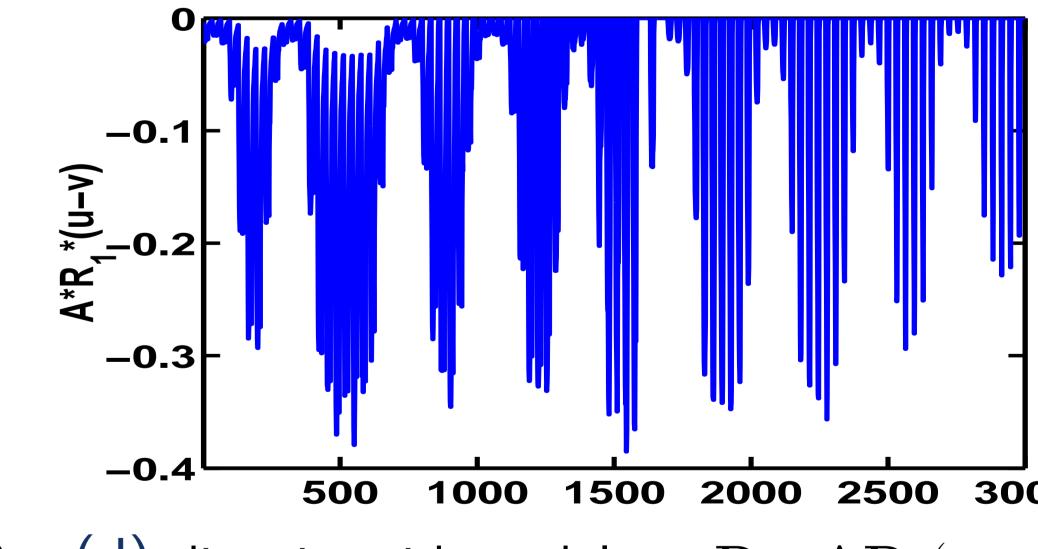
Teflon

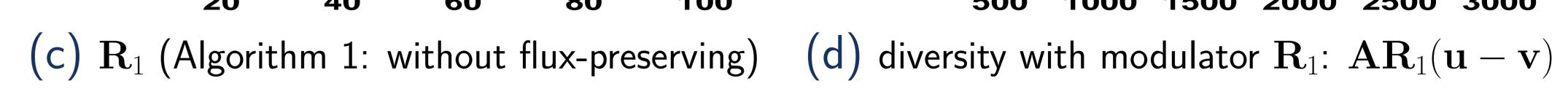
Real data  $\mathbf{u}, \mathbf{v} \in \mathbb{R}^{101}_+$ ,  $\mathbf{A} \in \mathbb{R}^{3040 \times 101}_+$ 

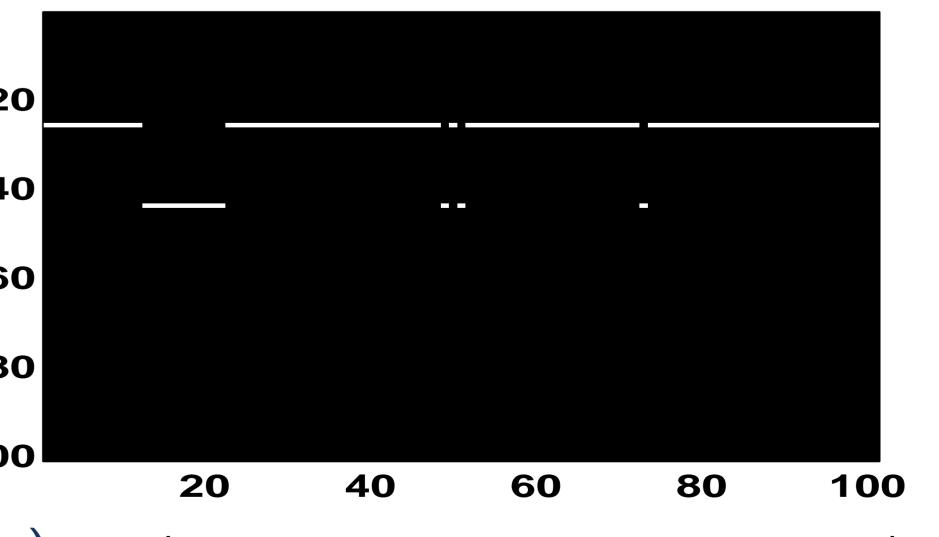


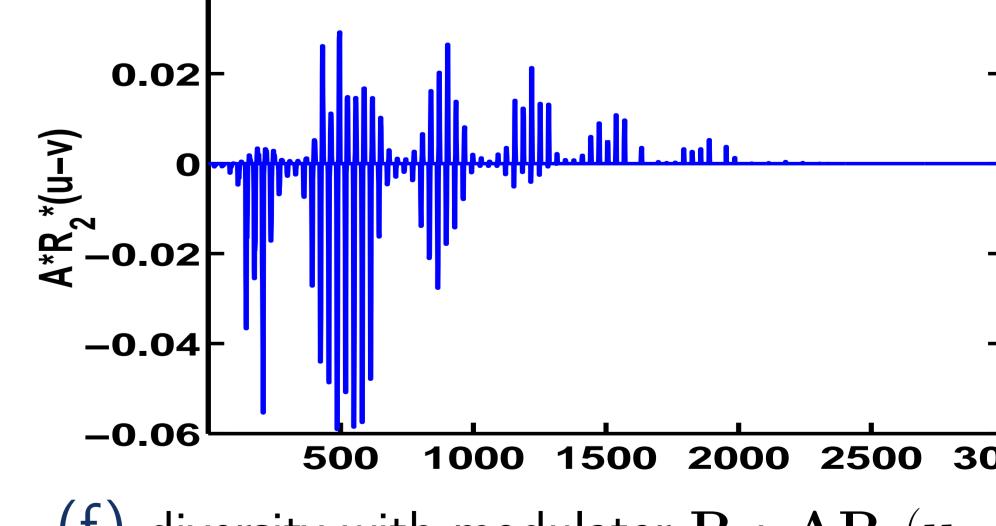












(e)  ${f R}_2$  (Algorithm 2: with flux-preserving) (f) diversity with modulator  ${f R}_2$ :  ${f AR}_2({f u}-{f v})$ 

Figure: Hypotheses(a), modulator(c,e) and diversity of measurements(b,d,f)

Table: Classification Accuracy no modulator modulator  ${f R}_1$  modulator  ${f R}_2$ 76.8%100.0%82.6%