



#### Parallel Patterns 4<sup>th</sup> part

058165 – Parallel Computing

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- What is the stencil pattern?
  - Update alternatives
  - 2D Jacobi iteration
- ☐ Implementing stencil with shift
- Stencil and cache optimizations
- Stencil and communication optimizations
- Recurrence

- A stencil pattern is a map where each output depends on a "neighborhood" of inputs
- These inputs are a set of fixed offsets relative to the output position
- A stencil output is a function of a "neighborhood" of elements in an input collection
  - Applies the stencil to select the inputs
- Data access patterns of stencils are regular
  - Stencil is the "shape" of "neighborhood"
  - Stencil remains the same

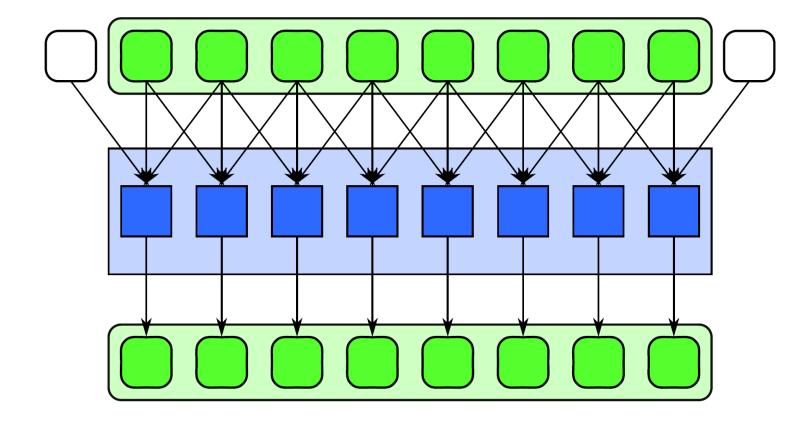
### Serial Stencil Example (part 1)

```
template<
       int NumOff, // number of offsets
       typename In, // type of input locations
3
       typename Out, // type of output locations
4
       typename F // type of function/functor
5
    void stencil(
       int n, // number of elements in data collection
8
       const In a[], // input data collection (n elements)
9
       Out r[], // output data collection (n elements)
10
       In b, // boundary value
11
       F func, // function/functor from neighborhood inputs to output
12
       const int offsets[] // offsets (NumOffsets elements)
13
14
```

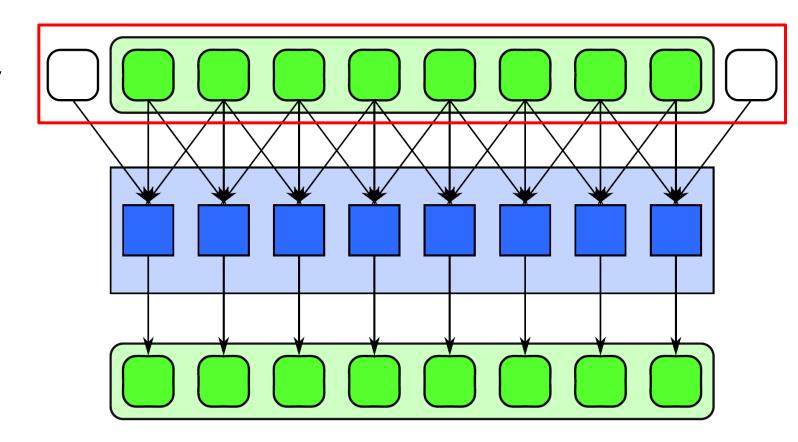
### Serial Stencil Example (part 2)

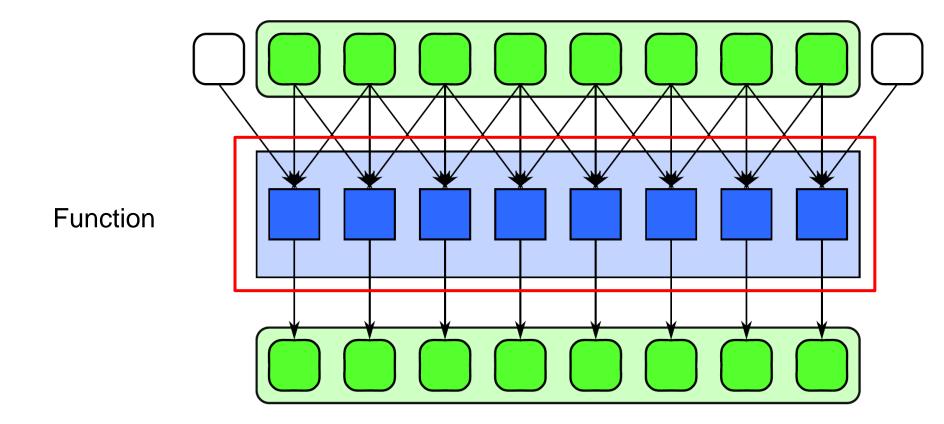
```
// array to hold neighbors
=15
         In neighborhood[NumOff];
 16
         // loop over all output locations
 17
         for (int i = 0; i < n; ++i) {
 18
            // loop over all offsets and gather neighborhood
 19
            for (int j = 0; j < NumOff; ++j) {
 20
                // get index of jth input location
 21
                int k = i + offsets[j];
 22
                if (0 \le k \&\& k \le n) {
 23
                    // read input location
 24
                    neighborhood[j] = a[k];
 25
                } else {
 26
                    // handle boundary case
 27
                    neighborhood[j] = b;
 28
 29
 30
            // compute output value from input neighborhood
 31
            r[i] = func(neighborhood);
 32
 33
 34
```

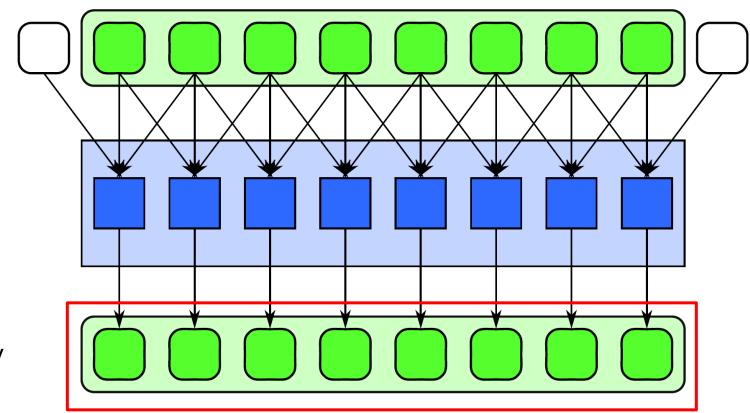
How would we parallelize this?



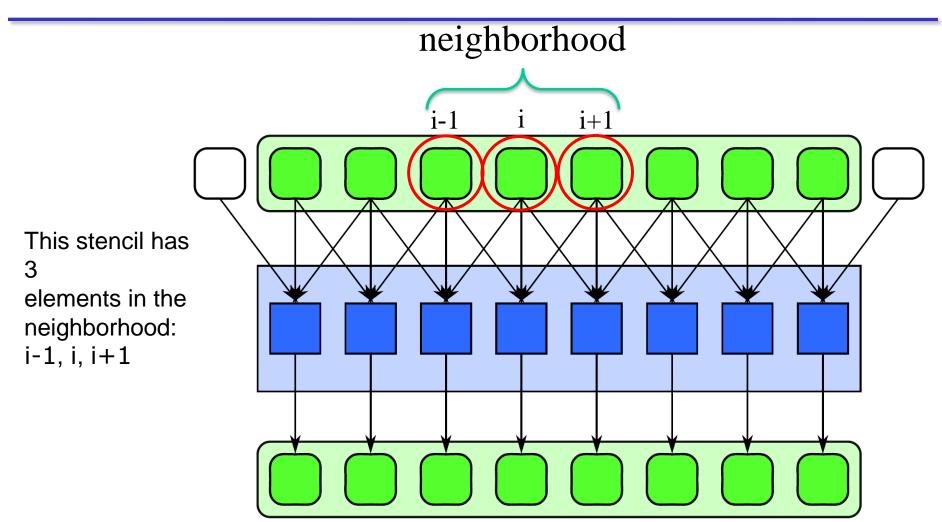
Input array



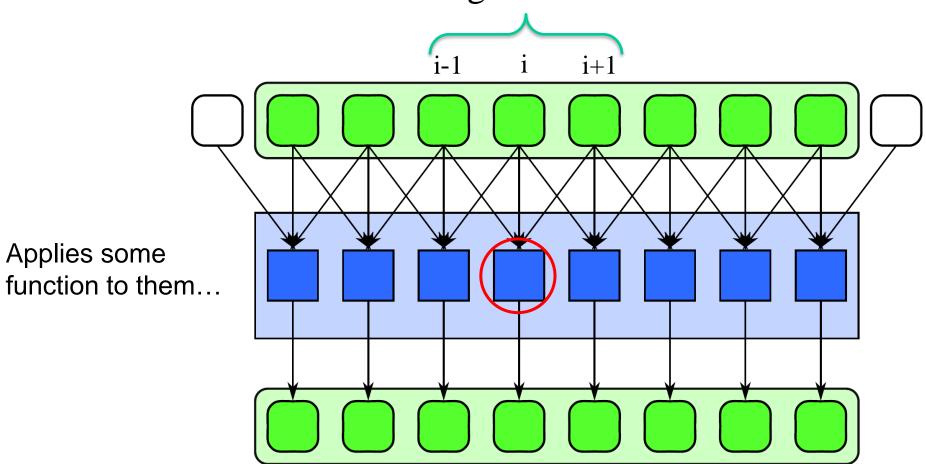


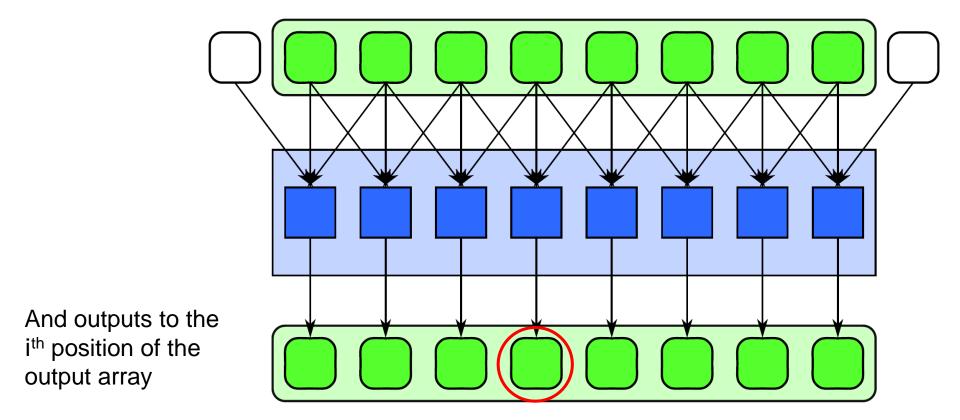


**Output Array** 



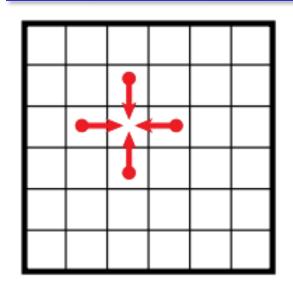




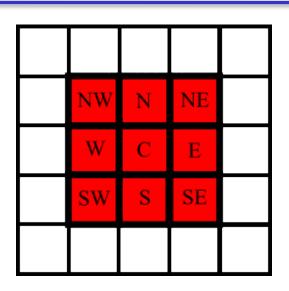


- Stencils can operate on one dimensional and multidimensional data
- Stencil neighborhoods can range from compact to sparse, square to cube, and anything else!
- It is the pattern of the stencil that determines how the stencil operates in an application

### 2-Dimensional Stencils



	<i>D</i> ( <i>x</i> , <i>y</i> -1)	
$D_{(x-1,y)}$	$P_{(x,y)}$	<i>D</i> ( <i>x</i> +1, <i>y</i> )
	D $(x,y+1)$	



4-point stencil

Center cell (P) is not used

5-point stencil

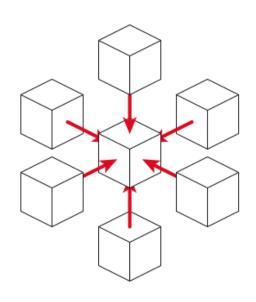
Center cell (P) is used as well

9-point stencil

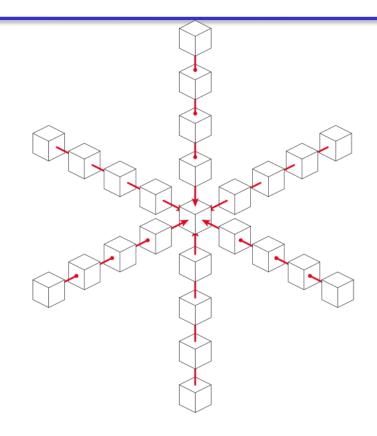
Center cell (C) is used as well

Source: http://en.wikipedia.org/wiki/Stencil\_code

## 3-Dimensional Stencils



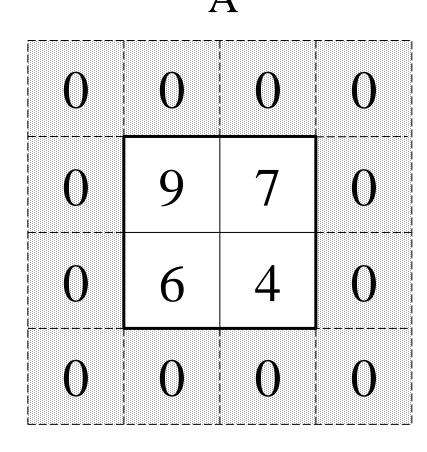
6-point stencil (7-point stencil)



24-point stencil (25-point stencil)

Source: http://en.wikipedia.org/wiki/Stencil\_code

☐ Here is our array, A



### ☐ Here is our array **A**

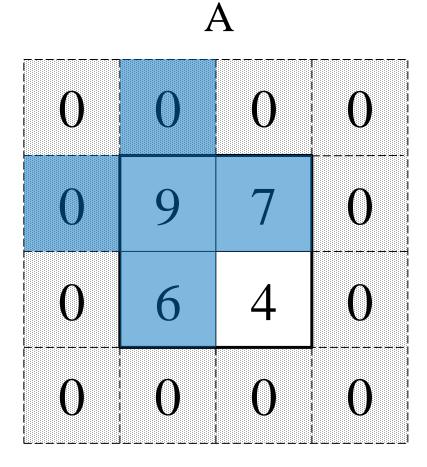
- □ B is the output array
  - ▶ Initialize to all 0
- ■Apply a stencil operation to the inner square of the form:

$$B(i,j) = avg(A(i,j),$$
  
 $A(i-1,j), A(i+1,j),$   
 $A(i,j-1), A(i,j+1)$ 

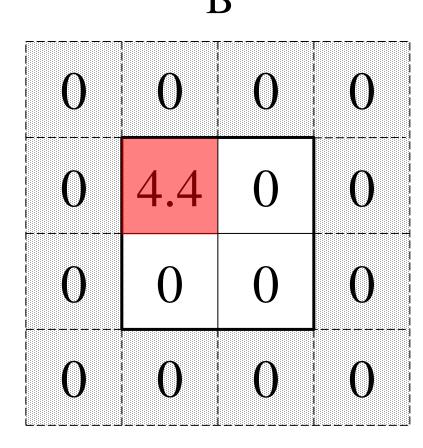
) What is the stencil?

0	0	0	0
0	9	7	0
0	6	4	0
0	0	0	0

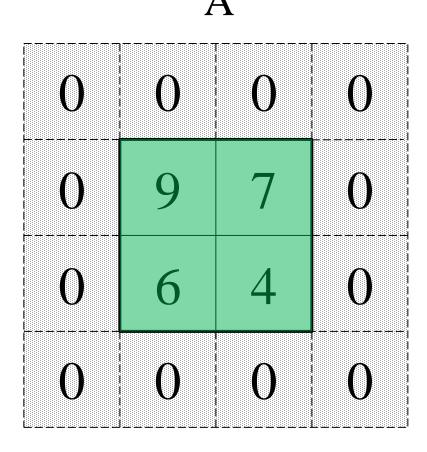
1) Average all blue squares



- 1) Average all blue squares
- 2) Store result in B

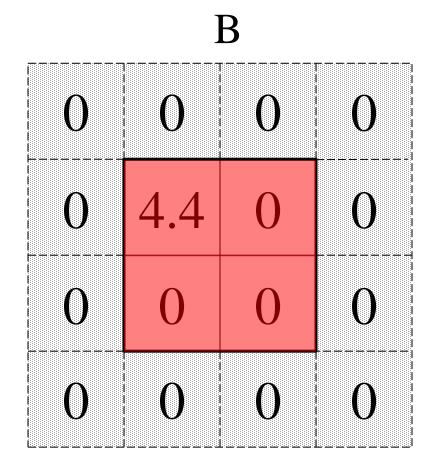


- 1) Average all blue squares
- 2) Store result in B
- 3) Repeat 1 and 2 for all green squares

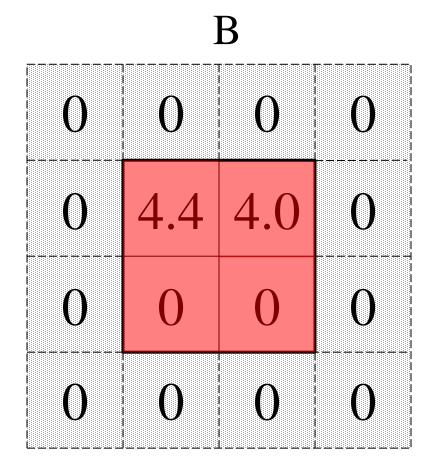


# Practice!

A				
0	0	0	0	
0	9	7	0	
0	6	4	0	
0	0	0	0	



A				
0	0	0	0	
0	9	7	0	
0	6	4	0	
0	0	0	0	



A				
0	0	0	0	
0	9	7	0	
0	6	4	0	
0	0	0	0	

	В				
0	0	0	0		
0	4.4	4.0	0		
0	3.8	0	0		
0	0	0	0		

A				
0	0	0	0	
0	9	7	0	
0	6	4	0	
0	0		0	

	]	В	
		0	0
0	4.4	4.0	0
0	3.8	3.4	0
0	0	0	0

#### **Outline**

- Partitioning
- What is the stencil pattern?
  - Update alternatives
  - ▶ 2D Jacobi iteration
- ☐ Implementing stencil with shift
- Stencil and cache optimizations
- Stencil and communication optimizations
- Recurrence

### Serial Stencil Example (part 1)

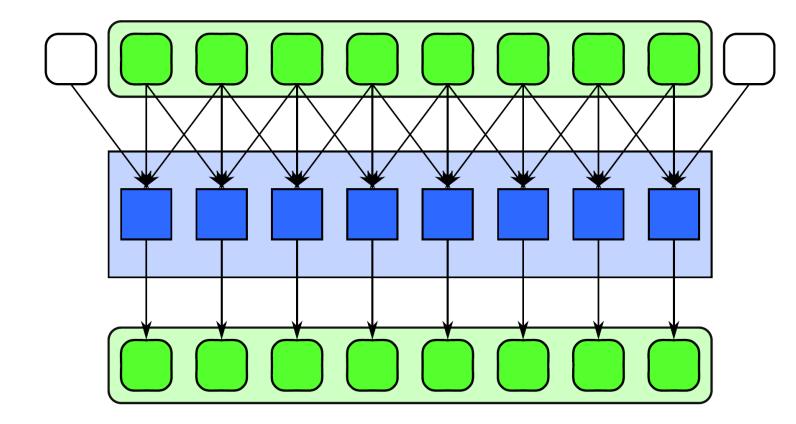
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10
              // boundary value
       In b.
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       F func, // function/functor from neighborhood inputs to output
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       const int offsets[] // offsets (NumOffsets elements)
13
14
```

### Serial Stencil Example (part 2)

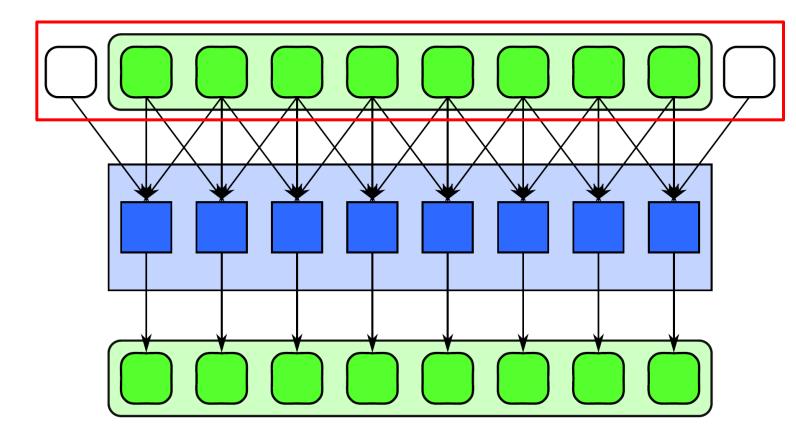
```
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=15
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         // loop over all output locations
 17
         for (int i = 0; i < n; ++i) {
 18
             // loop over all offsets and gather neighborhood
 19
            for (int j = 0; j < NumOff; ++j) {
 20
                // get index of jth input location
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                int k = i + offsets[j];
 22
                if (0 \le k \&\& k \le n) {
 23
                    // read input location
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                    neighborhood[j] = a[k];
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                } else {
 26
                    // handle boundary case
 27
                    neighborhood[j] = b;
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 29
 30
             // compute output value from input neighborhood
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     a[i] <del>r[i]</del> = func(neighborhood);
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 33
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```

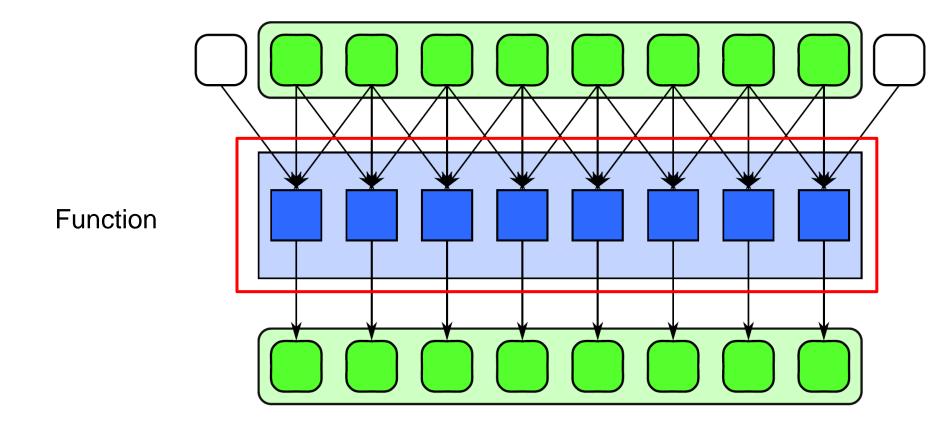
How would we parallelize this?

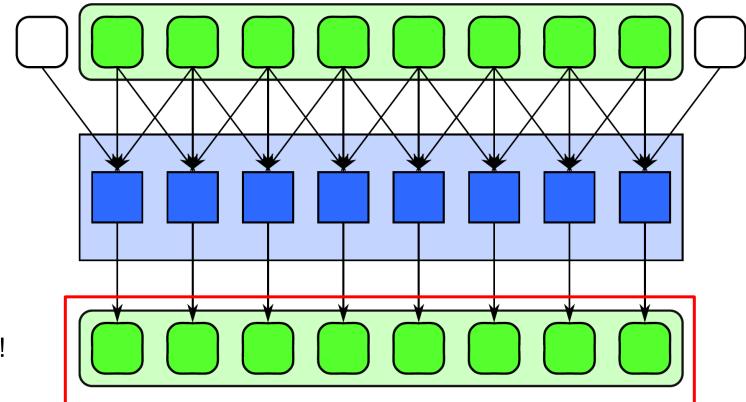
Updates occur in place!!!



Input array







Input Array !!!

☐ Here is our array, A

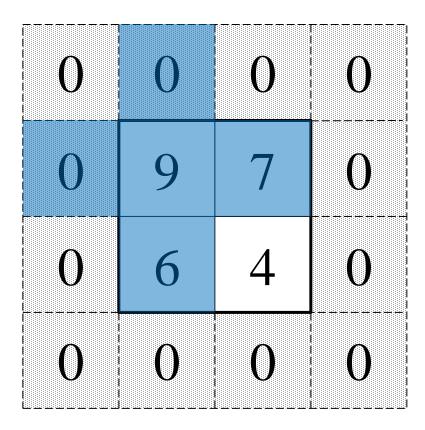
- ☐ Here is our array A
- ■Update A in place
- ■Apply a stencil operation to the inner square of the form:

$$A(i,j) = avg(A(i,j),$$
 $A(i-1,j), A(i+1,j),$ 
 $A(i,j-1), A(i,j+1)$ 

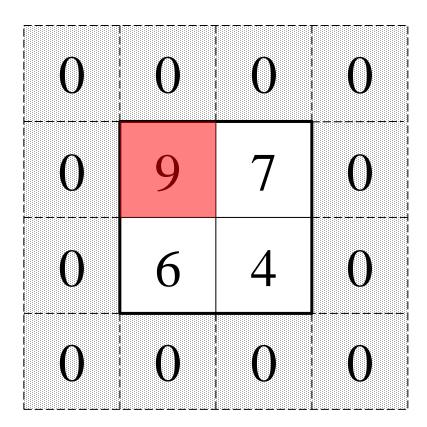
0	0	0	0
0	9	7	0
0	6	4	0
0	0	0	0

What is the stencil?

1) Average all blue squares



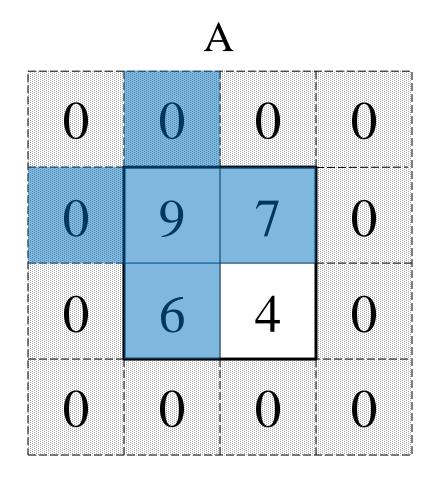
- 1) Average all blue squares
- 2) Store result in red square

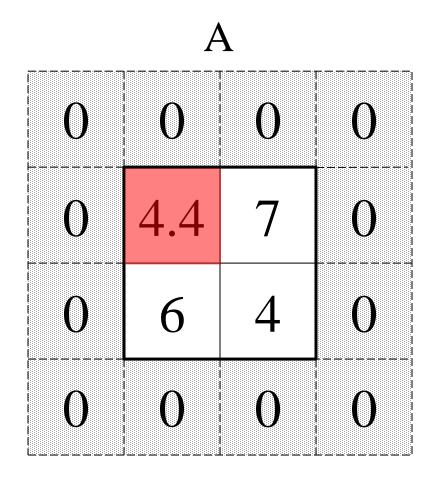


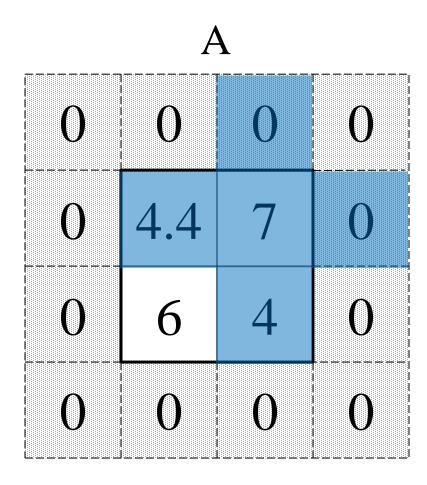
- 1) Average all blue squares
- 2) Store result in red square
- 3) Repeat 1 and 2 for all green squares

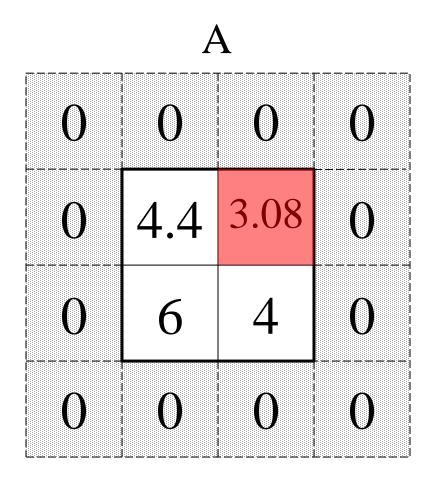
0	0	0	0
0	9	7	0
0	6	4	0
0	0	0	0

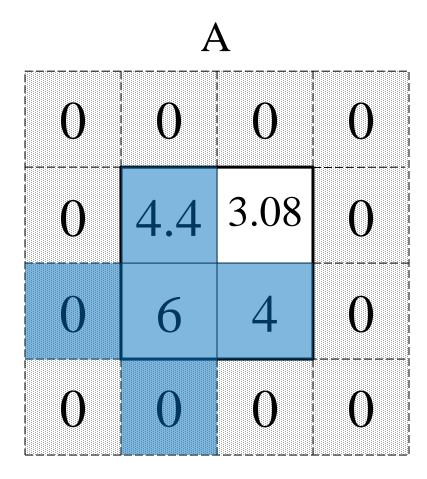
## Practice!

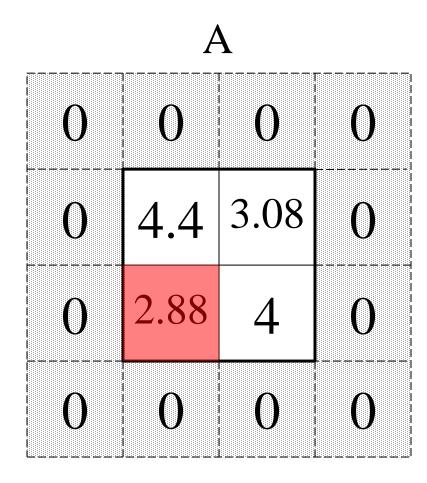


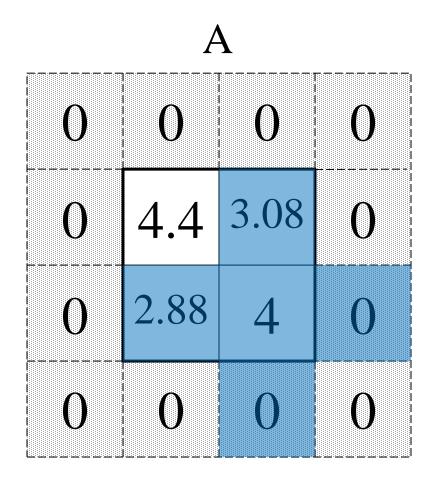


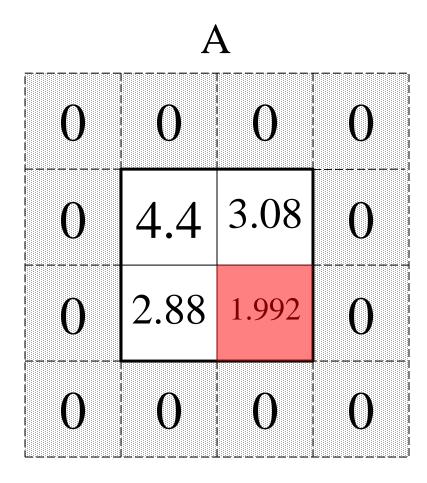






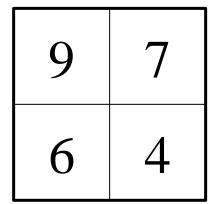




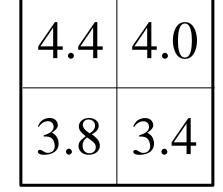


### Input

Separate output array



### Output



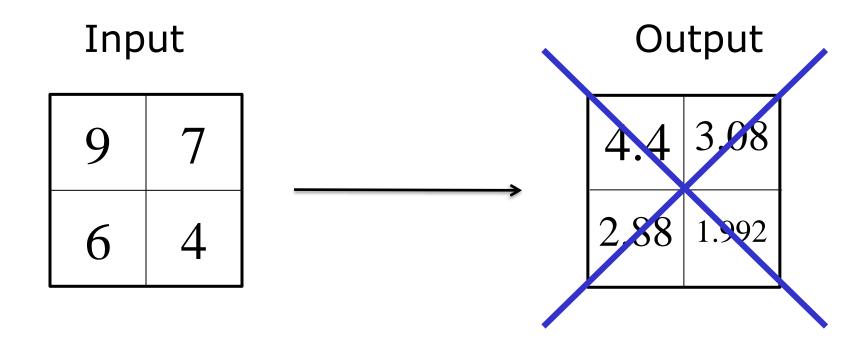
## Input

Updates occur in place

9	7
6	4

### Output

4.4	3.08
2.88	1.992



Is this output incorrect?

#### **Outline**

- What is the stencil pattern?
  - Update alternatives
  - 2D Jacobi iteration
- ☐ Implementing stencil with shift
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#### **Iterative Codes**

- ☐ Iterative codes are ones that update their data in steps
  - At each step, a new value of an element is computed using a formula based on other elements
  - Once all elements are updated, the computation proceeds to the next step or completes
- Iterative codes are most commonly found in computer simulations of physical systems for scientific and engineering applications
  - Computational fluid dynamics
  - ► Electromagnetics modeling
- They are often applied to solve partial differential equations
  - Jacobi iteration
  - Gauss-Seidel iteration
  - Successive over relaxation

- Stencils essentially define which elements are used in the update formula
- Because the data is organized in a regular manner, stencils can be applied across the data uniformly

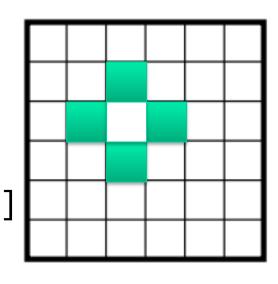
Consider the following code

```
for k=1, 1000

for i=1, N-2

for j = 1, N-2

a[i][j] = 0.25 * (a[i][j] + a[i-1][j] + a[i+1][j] + a[i][j-1] + a[i][j-1]
```



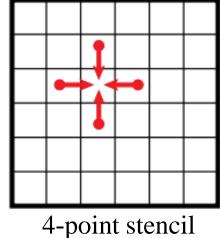
5-point stencil

} } }

Do you see anything interesting?

How would you parallelize?

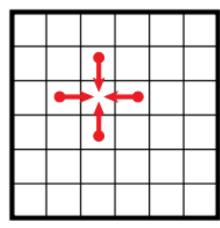
- ☐ Consider a 2D array of elements
- ☐ Initialize each array element to some value
- □ At each step, update each array element to the arithmetic mean of its N, S, E, W neighbors
- ☐ Iterate until array values converge
- ☐ Here we are using a 4-point stencil

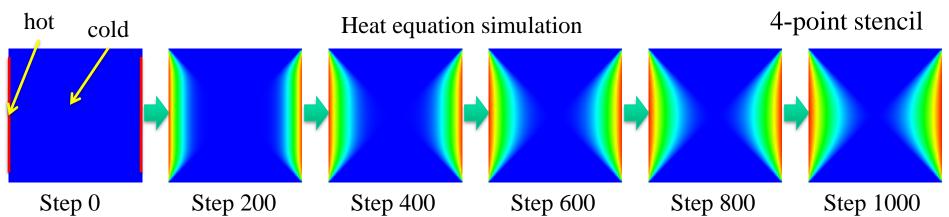


. .

☐ It is different from before because we want to update all array elements simultaneously ... How?

- Consider a 2D array of elements
- ☐ Initialize each array element to some value
- □ At each step, update each array element to the arithmetic mean of its N, S, E, W neighbors
- ☐ Iterate until array values converge





#### **Outline**

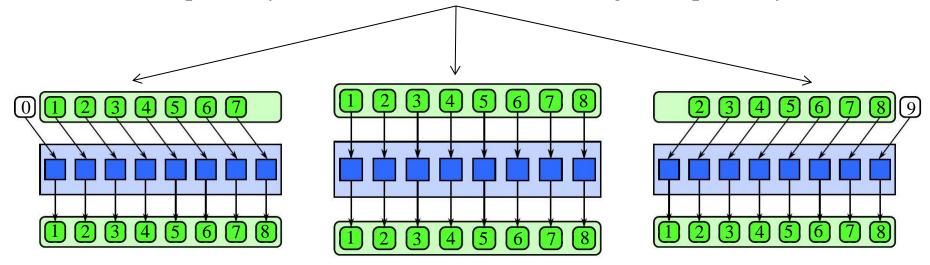
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- One possible implementation of the stencil pattern includes shifting the input data
- □ For each offset in the stencil, we gather a new input vector by shifting the original input by the offset amount

#### **Implementing Stencil with Shift**

All input arrays are derived from the same original input array



#### **Implementing Stencil with Shift**

- □ This implementation is only beneficial for one dimensional stencils or the memory-contiguous dimension of a multidimensional stencil
- Memory traffic to external memory is not reduced with shifts
- But, shifts allow vectorization of the data reads, which may reduce the total number of instructions

- Partitioning
- What is the stencil pattern?
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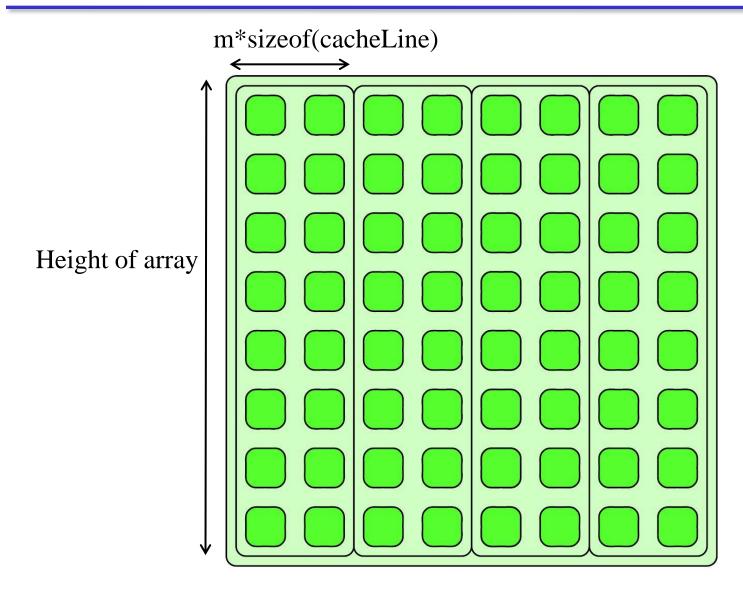
- Assuming 2D array where rows are contiguous in memory...
  - ► Horizontally related data will tend to belong to the same cache line
  - Vertical offset accesses will most likely result in cache misses

- Assigning rows to cores:
  - Maximizes horizontal data locality
  - Assuming vertical offsets in stencil, this will create redundant reads of adjacent rows from each core
- Assigning columns to cores:
  - Redundantly read data from same cache line
  - Create false sharing as cores write to same cache line

- Assigning "strips" to each core can be a better solution
- Strip-mining: an optimization in a stencil computation that groups elements in a way that avoids redundant memory accesses and aligns memory accesses with cache lines

- A strip's size is a multiple of a cache line in width, and the height of the 2D array
- Strip widths are in increments of the cache line size so as to avoid false sharing and redundant reads
- Each strip is processed serially from top to bottom within each core

#### **Stencil and Cache Optimizations**



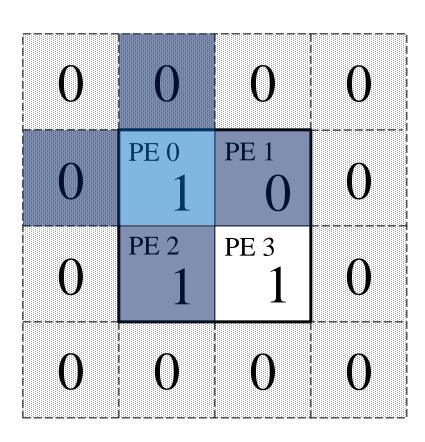
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# Stencil and Communication Optimizations

■ When data is distributed, ghost cells must be explicitly communicated among nodes between loop iterations

- ☐ Darker cells are PE 0's ghost cells
- ☐ After first iteration of stencil computation
  - O PE 0 must request PE 1 & PE 2's stencil results
  - PE 0 can perform another iteration of stencil



## Stencil and Communication Optimizations

- Generally better to replicate ghost cells in each local memory and swap after each iteration than to share memory
  - Fine-grained sharing can lead to increased communication cost

# Stencil and Communication Optimizations

- □ Halo: set of all ghost cells
- □ Halo must contain all neighbors needed for one iteration
- ☐ Larger halo (deep halo)
  - ▶ Trade off
    - less communications and more independence, but...
    - more redundant computation and more memory used

■ Latency Hiding: Compute interior of stencil while waiting for ghost cell updates