

CS181 HW3

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1.

a) Codes in 1a.cpp

b) Codes in 1b.cpp

ciphertext in binary: 1 0 0 0 0 1 0 1 0 0 1 0 0 1 1 1 0 1 1 0 0

c) Codes in 1c&d.cpp

keystream:

181 210 86 5 165 245 93 50 124 85 172 216 184 91 155 127 193 164 7 131 207 15
52 13

Keystream in binary:

1 0 1 1 0 1 0 1 1 1 0 1 0 0 1 0 0 1 0 1 0 1 1 0 0 0 0 0 0 0 1 0 1 1 0 1 0 0 1
0 1 1 1 1 1 0 1 0 1 0 1 0 1 1 1 0 1 0 0 1 1 0 0 1 0 0 1 1 1 1 1 0 0 0 1 0 1
0 1 0 1 1 0 1 0 1 1 0 0 1 1 0 1 1 0 0 0 1 0 1 1 1 0 0 0 0 0 1 0 1 1 0 1 1 1 0
0 1 1 0 1 1 0 1 1 1 1 1 1 1 1 1 0 0 0 0 0 1 1 0 1 0 0 1 0 0 0 0 0 0 0 0 1 1 1
1 0 0 0 0 0 1 1 1 1 0 0 1 1 1 1 0 0 0 0 1 1 1 1 0 0 1 1 0 1 0 0 0 0 0 0 1 1
0 1

d) Codes in 1c&d.cpp

plaintext in binary after ciphertext XORing with keystream:

```
0 1 0 1 1 0 0 1 0 1 1 0 1 1 1 1 0 1 1 1 0 1 0 0 1 0 0 1 1 1 0 1 1 1 0 1
1 0 0 1 1 0 0 1 0 1 0 0 1 0 0 0 0 0 1 1 0 0 1 1 0 0 1 1 0 1 1 1 1 0 1 1 1
0 1 0 1 0 1 1 0 1 1 1 0 0 1 1 0 0 1 0 0 0 0 1 0 0 0 0 0 0 1 1 1 0 1 0 0 0 1
1 0 1 0 0 0 0 1 1 0 0 1 0 1 0 0 1 0 0 0 0 0 0 1 1 0 0 0 0 1 0 1 1 0 1 1 1 0
0 1 1 1 0 0 1 1 0 1 1 1 0 1 1 1 0 1 1 0 0 1 0 1 0 1 1 1 0 0 1 0 0 0 1 0 0 0
0 1
```

Convert to ASCII using online translator:

plaintext:

You've found the answer!

From

To

Binary

Text

Open File

Open Bin File

Q

Paste binary numbers or drop file:

0 1 0 1 1 0 0 1 0 1 1 0 1 1 1 1 0 1 1 1 0 1 0 1 0 0 1 0 0
1 1 1 0 1 1 1 0 1 1 0 0 1 1 0 0 1 0 1 0 0 1 0 0 0 0 0 1
1 0 0 1 1 0 0 1 1 0 1 1 1 1 0 1 1 1 0 1 0 1 0 1 1 0 1 1 1
0 0 1 1 0 0 1 0 0 0 0 1 0 0 0 0 0 0 1 1 1 0 1 0 0 0 1 1 0
1 0 0 0 0 1 1 0 0 1 0 1 0 0 1 0 0 0 0 0 0 1 1 0 0 0 0 1 0
1 1 0 1 1 1 0 0 1 1 1 0 0 1 1 0 1 1 1 0 1 1 1 0 1 1 0 0 1

Character encoding (optional)

ASCII

Convert

Reset

Swap

You've found the answer!

2.

a)

- As decryption only requires $M_i = C_i + E_k(C_{i-1})$, so if C_2 is wrongly transmission, then, it will have effect on M_2 , and M_3 which use $E_k(C_2)$;
- So Bob can still decrypt M_1, M_4, \dots, M_{10} correctly, since they do not use C_2 to decrypt.
- M_2 will have 2 bits off, since it's XORing the C_2 with other stuff, as C_2 has two bits off, M_2 is gonna also has 2 bits off.
- M_3 seems to be completely off, since some slight change in the input of the $E_k()$ can have great effect on the output, so $E_k(C_2)$ is gonna be huge off even if C_2 is just 2 bits off, and then the decrypted M_3 will also be huge different.

b)

If the CBC uses the same IV, it will be **deterministic**, which means that same plaintext will always be mapped to the same ciphertext. It is not save against CPA attack because the attacker can test on whether some specific plaintext map to some ciphertext he knows or not.

For Eve, as she know M and M' , and their corresbonding C and C' . She can basically ask Bob to encrypt one of the M 's. And since the encryption is deterministic, it will definitely output one of the C 's. And based on the result, she can see the ciphertext of M and M' .

Reference:

<https://www.rapidtables.com/convert/number/binary-to-ascii.html>