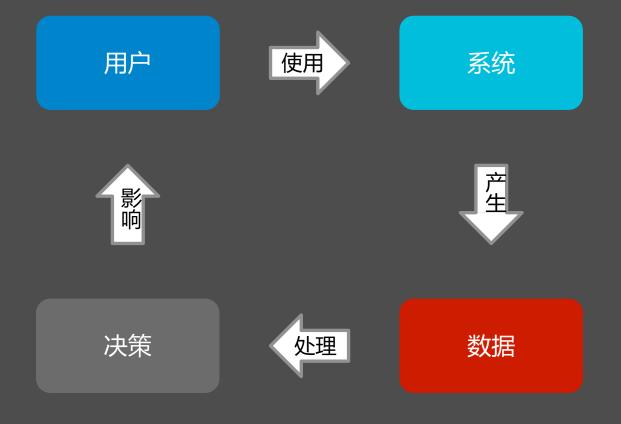
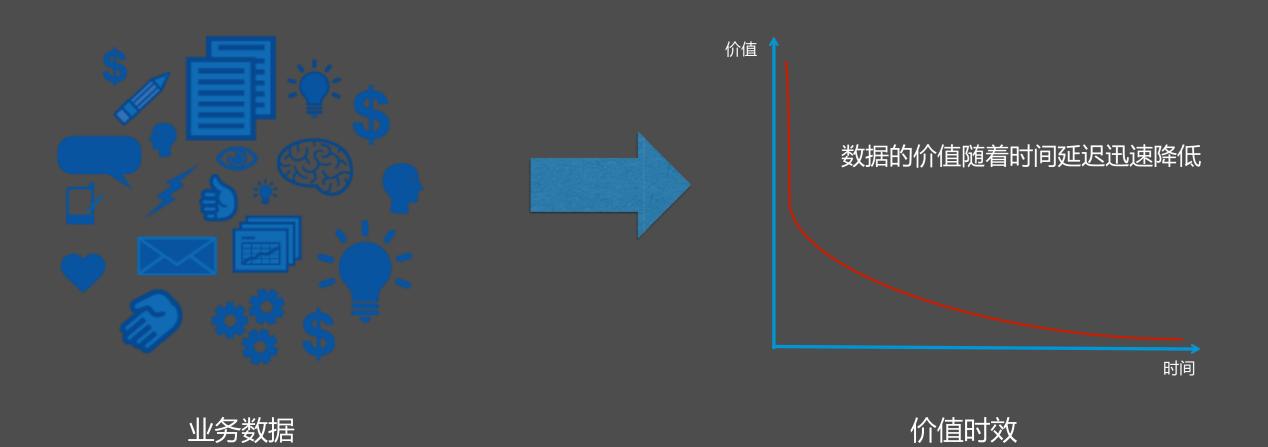


初识流计算Streaming

A Whirlwind Tour of Streaming Data Processing



商业和数据形成闭环







系统



越快 越好



决策



数据

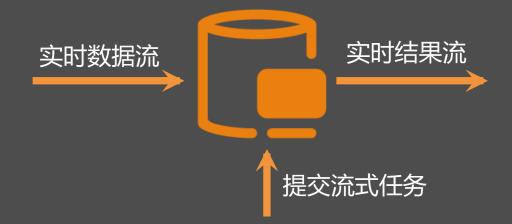
越快越有竞争优势 大数据实时化——流式处 理

离线 (批量) 计算



批量计算是一种批量、高时延、主动发起的计算任务

流计算



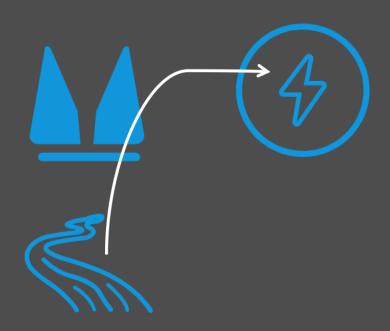
流计算是一种持续、低时延、事件触发的计算任务

离线 (批量) 计算



开船去湖里打鱼

流计算

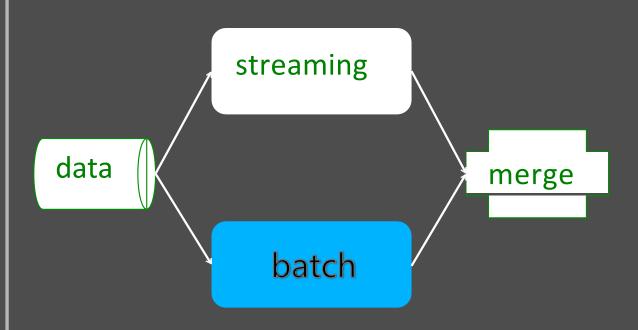


拦河建坝发电

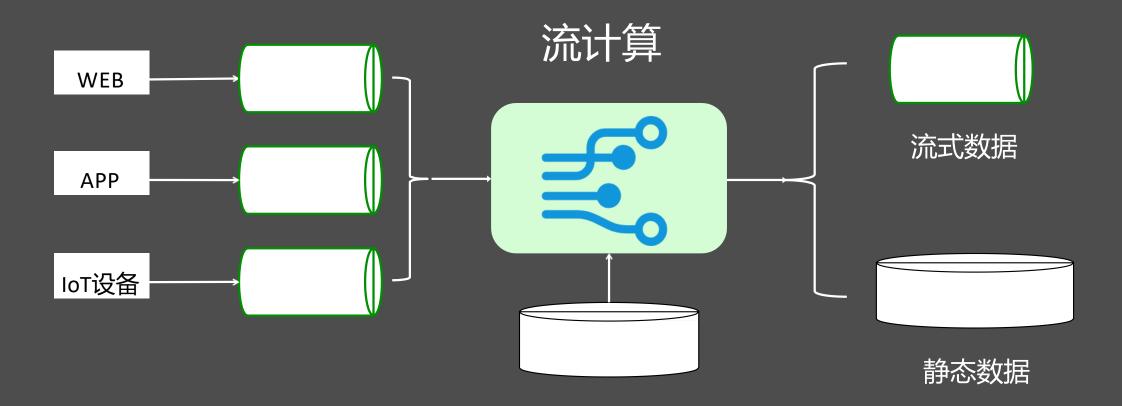
理论



实践



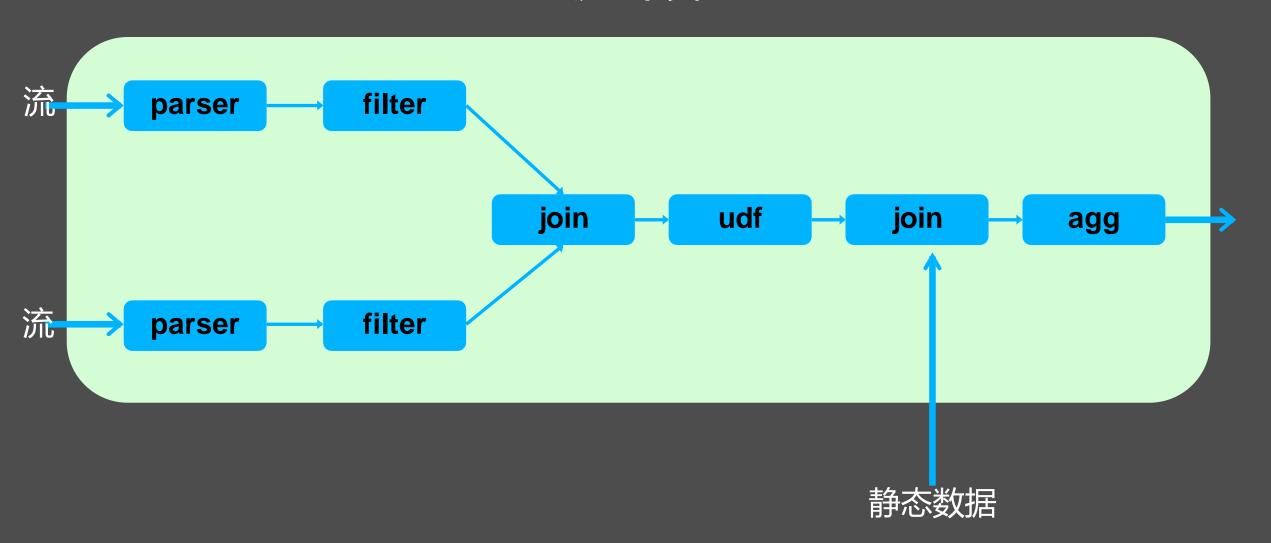
典型架构



流式数据

静态数据

流计算



典型场景





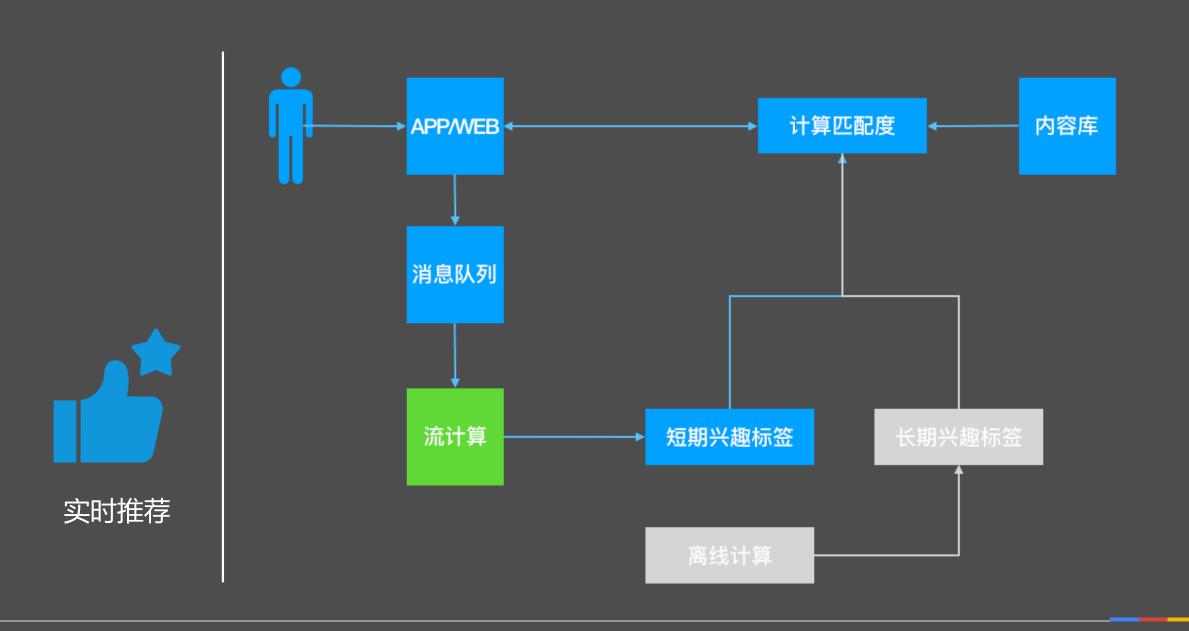


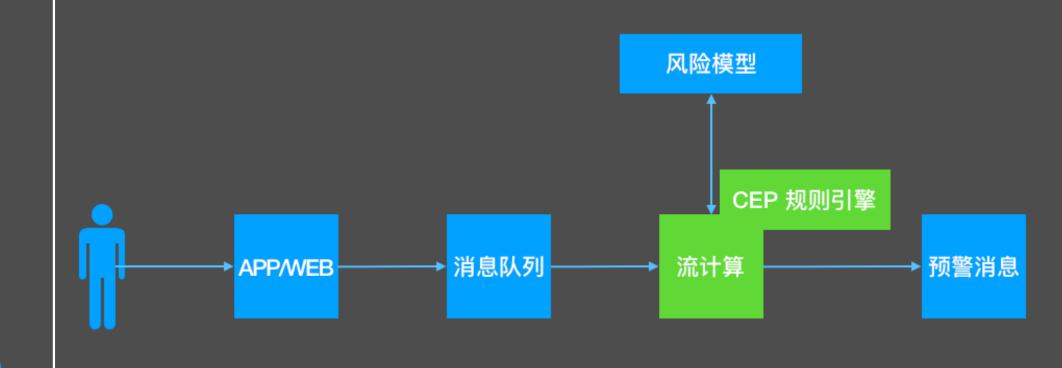


欺诈检测



实时数仓/报表







欺诈检测

相关术语

□ Q1: 什么是streaming流式计算

a type of data processing engine that is designed with infinite data sets in mind (一种被设计用于处理无穷数据集的数据处理引擎)

□ Q2: 无穷数据集

a type of ever-growing, essentially infinite data set (一种持续生成,本质上是无穷尽的数据集)

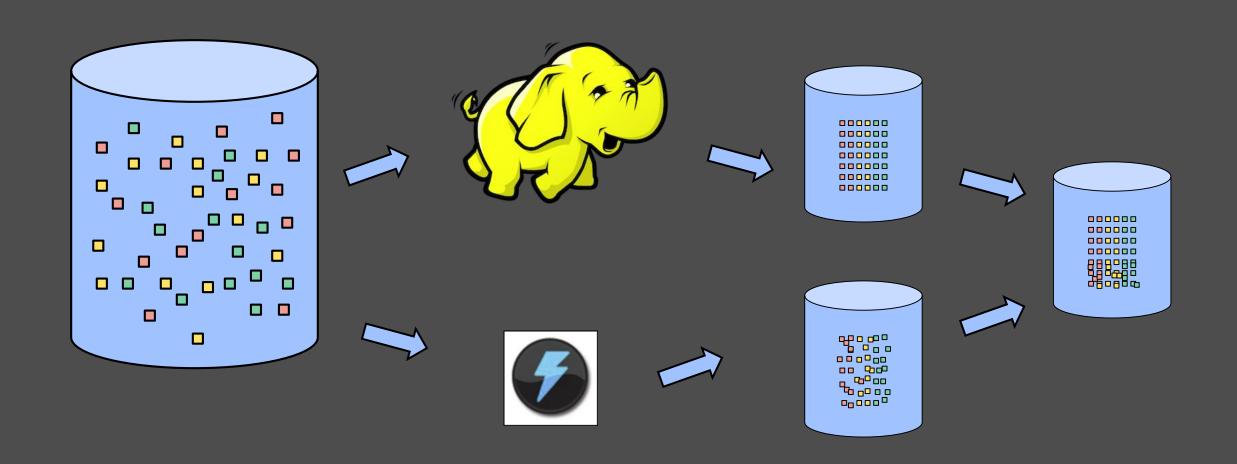
Agenda

1 State of the Streaming World

2 Streaming 101

State of the Streaming World

The Lambda Architecture



流计算最夸张的限制

□ 长久以来,流计算系统被认为是专为提供低延迟、不精确结果的某些特定场景而设计, 并配合一个更强大的批处理系统来提供最终准确的结果,如Lambda架构(Lambda Architecture)。

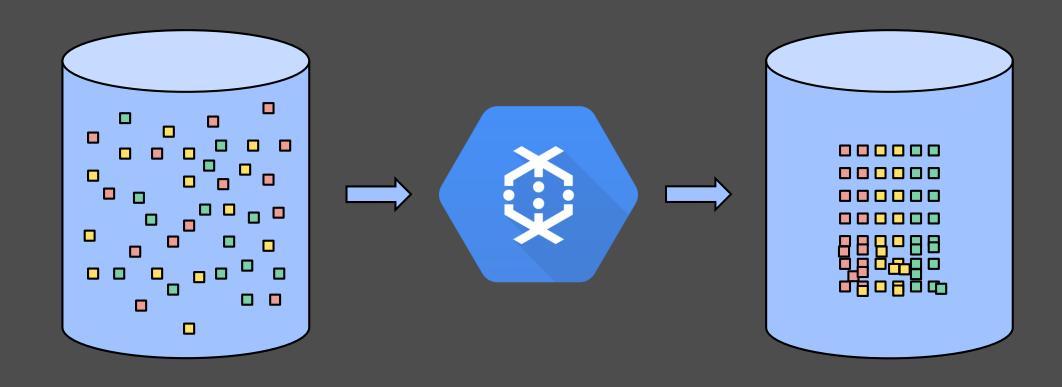
□ Lambda架构的基本思想

与批处理系统一起运行流计算系统,同时进行几乎一样的计算。流计算系统提供低延迟、不准确的结果(或是因为使用了近似算法,或是因为流计算系统本身没能提供足够准确的结果),而一段时间之后当批处理计算完成,再给出正确的结果。

设计良好的流计算系统的能力是批处理系统的功能的超集。



The Evolution of Streaming



What does it take?

Strong Consistency

Tools for Reasoning About Time

Why consistency is important

- Mostly correct is not good enough
- Required for exactly-once processing
- Required for repeatable results
- Cannot replace batch without it

Fault tolerance in streaming



☐ How do we ensure the results are always correct?

☐ Failures should not lead to data loss or incorrect results

Fault tolerance in streaming

at least once: ensure all operators see all.

Storm: Replay stream in failure case (acking of individual records)

Exactly once: ensure that operators do not perform duplicate updates to their state

Flink: Distributed Snapshots

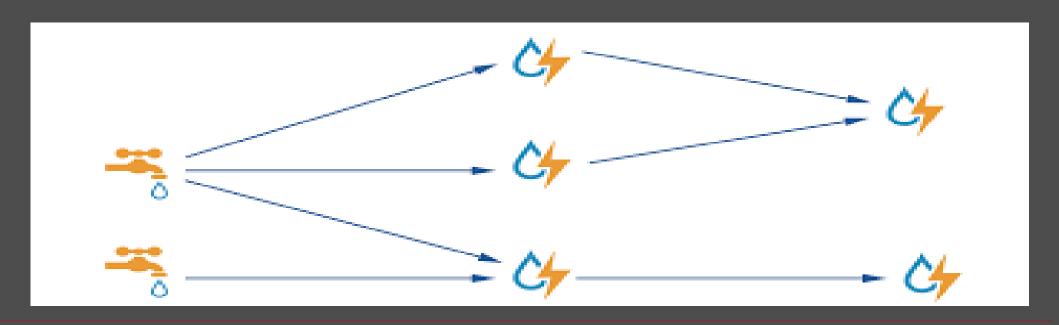
Spark: Micro-batches on batch runtime

Google Cloud Dataflow: Transactional updates

Storm Record acknowledgement

□ Record acknowledgement模式

Topology的Source会保留其产生的所有记录备份用来处理Fail情况。当源头一条记录的所有派生记录都被整个Topology处理完成,Source节点就可以删除其备份;当系统出现部分Fail情况,例如一条记录并没有收到其下游的派生记录的确认,Source就会重新发送该记录到下游的Topology以便重新进行计算。



Storm Record acknowledgement

□缺点

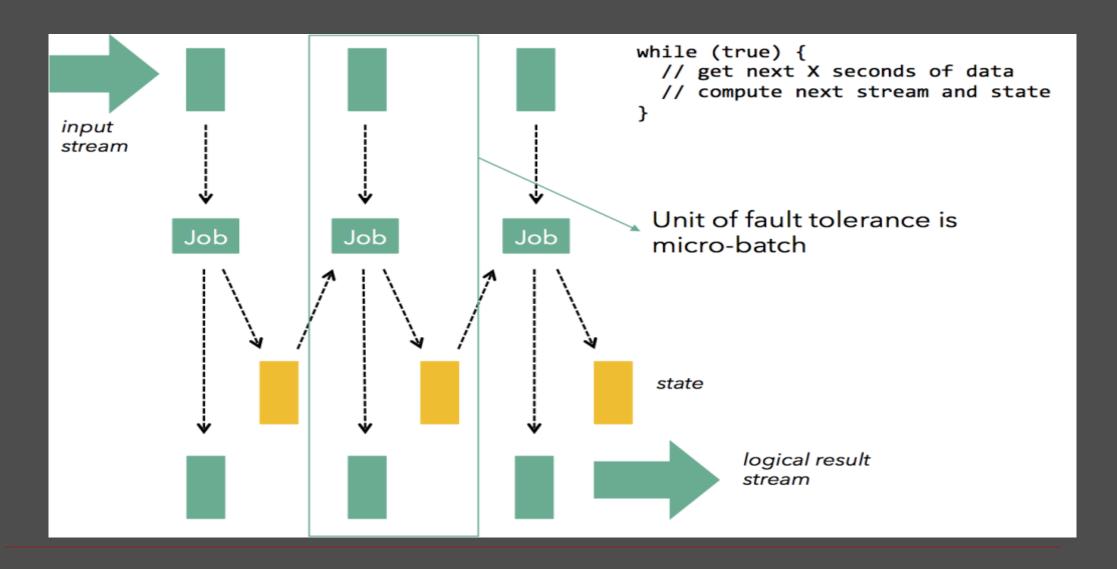
- 1、性能低:每个都会产生一条对应的系统消息到acker,同时有额外的计算消耗,并且acker消息会消耗大量的网络带宽。Single ack注定与高吞吐量无缘。
- 2、不支持Exactly once语义。(Trident支持Exactly once语义,可以归类为Micro batches,性能下降严重。)
- 3、Back Pressure:记录确认的容错方式会导致上游节点错误地认为数据处理出现了Fail(实际上仅仅是由于back pressure导致记录处理不及时,而无法ack)

Spark: Micro-batches(stream discretization)

□ 流式处理系统中的算子都是在record级别进行计算同步和容错,由此带来 了在record如此低层次上进行处理的复杂和开销。

□ 把连续的数据流不要切分到record级别,而是收敛切分为一批一批微批的、原子的数据进行类似Batch的计算。这样,每个batch数据可能会成功或者 失败处理,我们就对当前失败的这一小批数据进行处理即可。

Spark: Micro-batches(stream discretization)



Spark: Micro-batches(stream discretization)

- □缺点:
 - 1、延迟高
 - 2、窗口受限:不支持count或session窗口
- 3、Back Pressure:如果某个下游的算子处理较慢,此时如果负责数据流切分的 Operator速度快于下游的阻塞节点,就会导致数据切分比原有的配置时间更长。导致越来越 多的批次在内存排队等待被处理,最终内存OOM,Runtime不够稳定。
- 4、微批处理模型的最大限制可能是它连接了两个不应连接的概念:应用程序定义的窗口 大小和系统的批处理间隔。

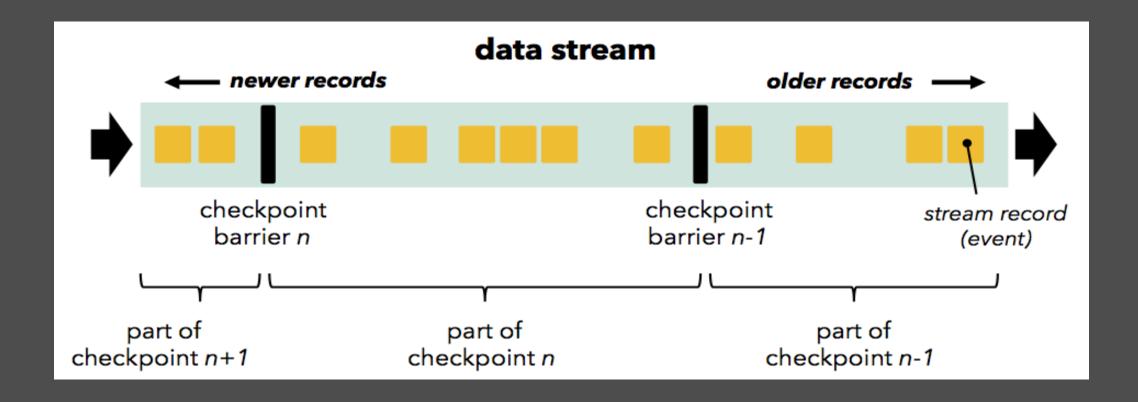
Flink: Distributed Snapshots

□ Flink使用的是Chandy Lamport算法的一个变种,定期对正在运行的流拓扑的状态做快照, 并将这些快照存储到持久存储。

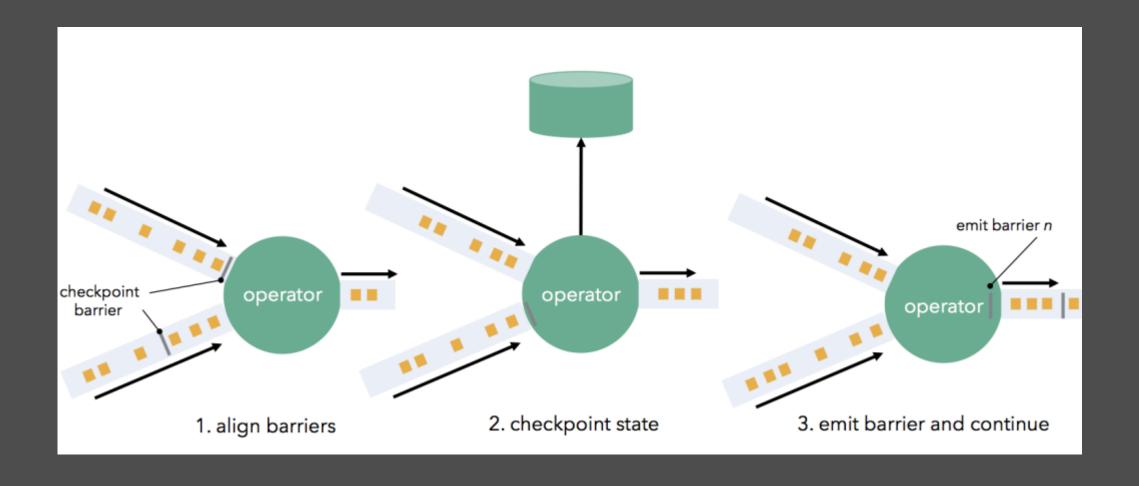
」 遵循真正的持续计算模型(低延迟,流量控制和真正的流编程模型)和高吞吐量的优点, 并且也是在Chandy-Lamport算法中被严格证明的Exactly Once保证。

□ 除了持久化有状态计算的状态(其他容错机制也需要这样做)之外,这种容错机制几乎 没有开销。对于小状态(例如,计数或其他统计摘要),这种持久化开销通常可忽略不 计。

Flink: Distributed Snapshots



Flink: Distributed Snapshots



Streaming 101

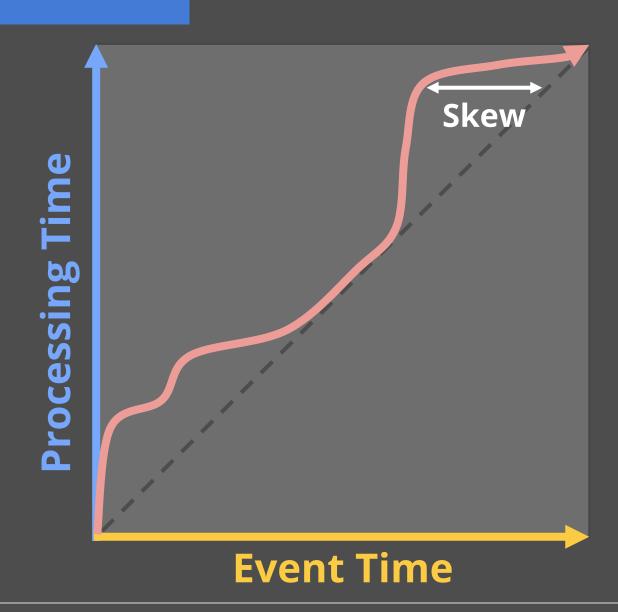
Event Time vs Processing Time Batch vs Streaming Approaches

Event Time vs Processing Time

Event Time - When Events Happened

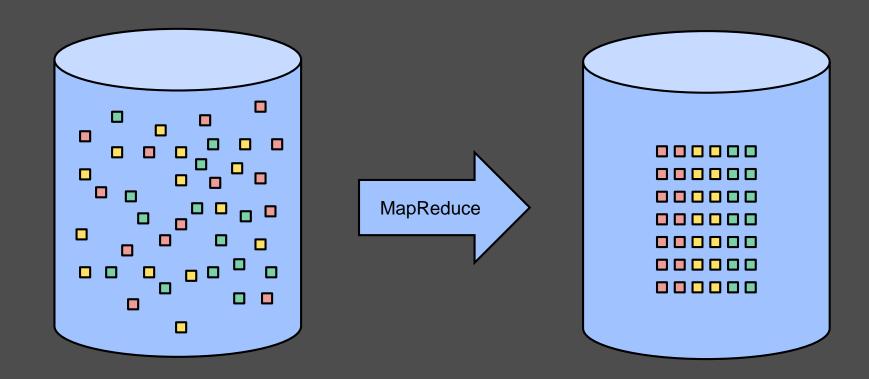
Processing Time - When Events Are Processed

Event Time Skew

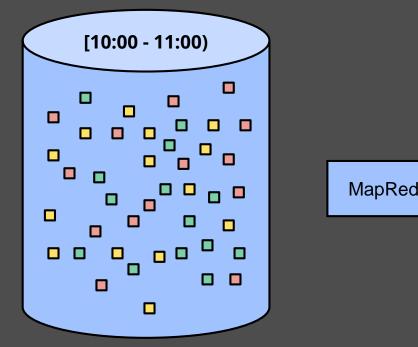


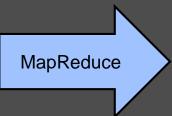
Batch vs Streaming

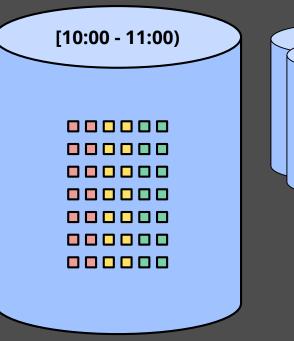
Batch

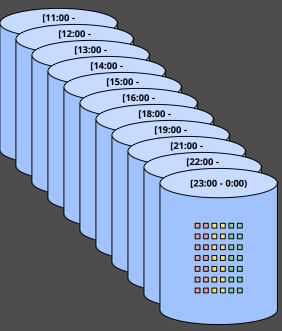


Batch: Fixed Windows

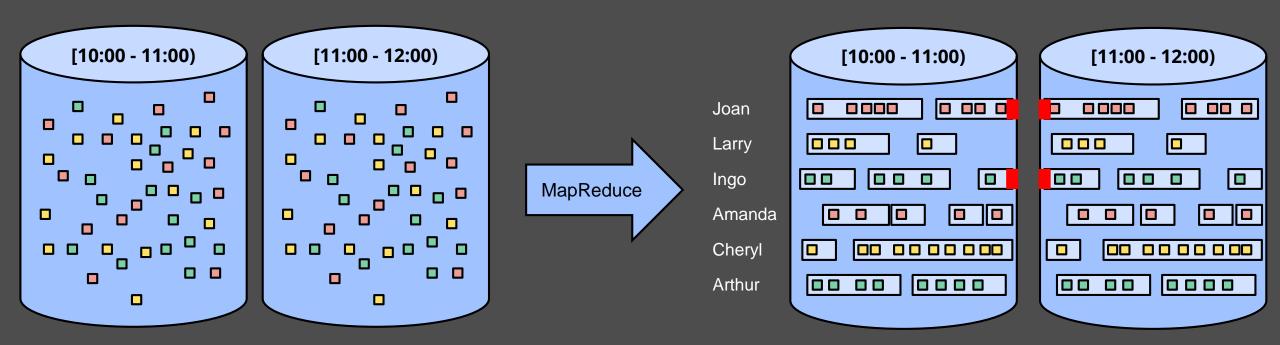




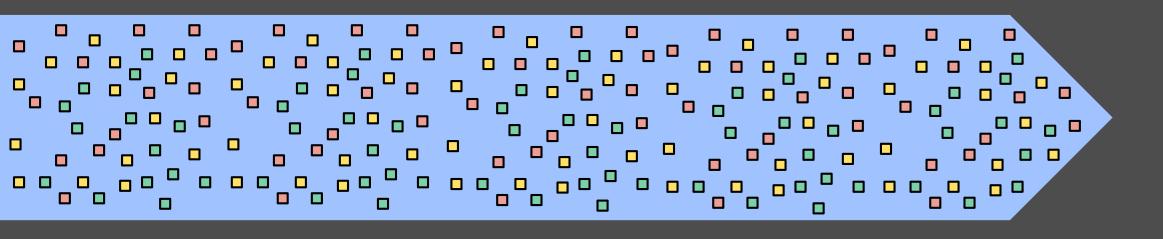




Batch: User Sessions



Streaming



Confounding characteristics of data streams

Unordered

Unbounded

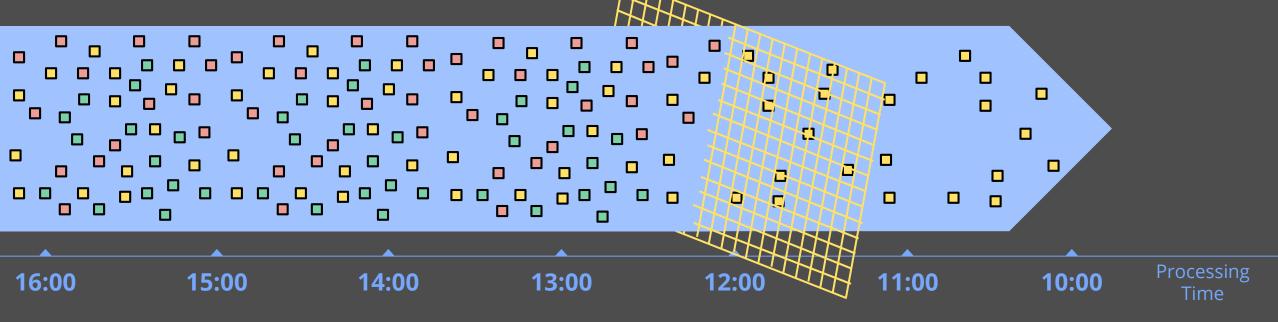
Of Varying Event Time Skew

Approaches

Approaches to reasoning about time

- 1. Time-Agnostic Processing
- 2. Approximation
- 3. Processing Time Windowing
- 4. Event Time Windowing

1. Time-Agnostic Processing - Filters



Example Input: Web server traffic logs

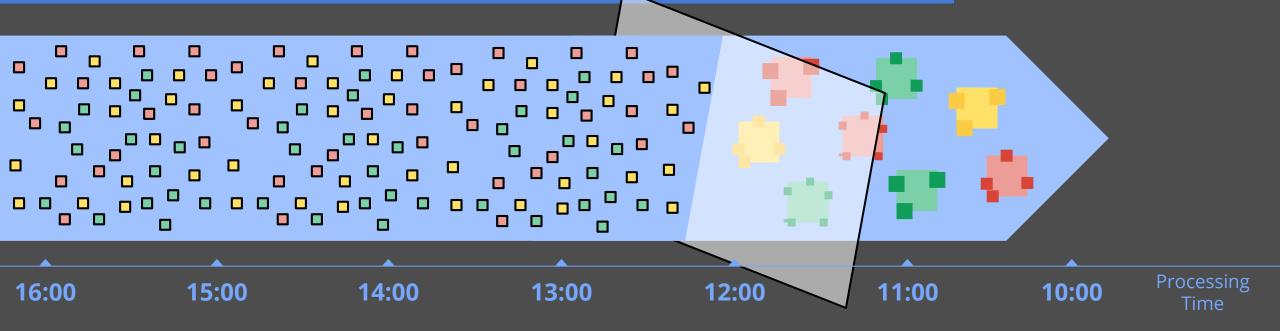
Example Output: All traffic from specific domains

Pros: Straightforward

Efficient

Cons: Limited utility

2. Approximation via Online Algorithms



Example Input: Twitter hashtags

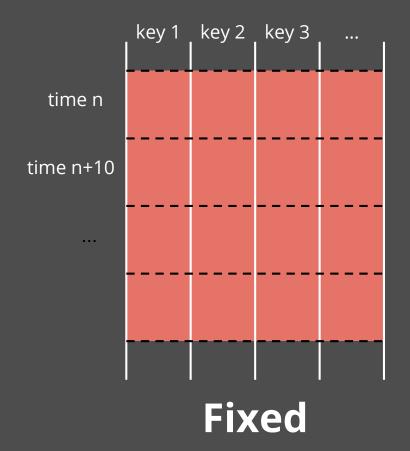
Example Output: Approximate top N hashtags per prefix

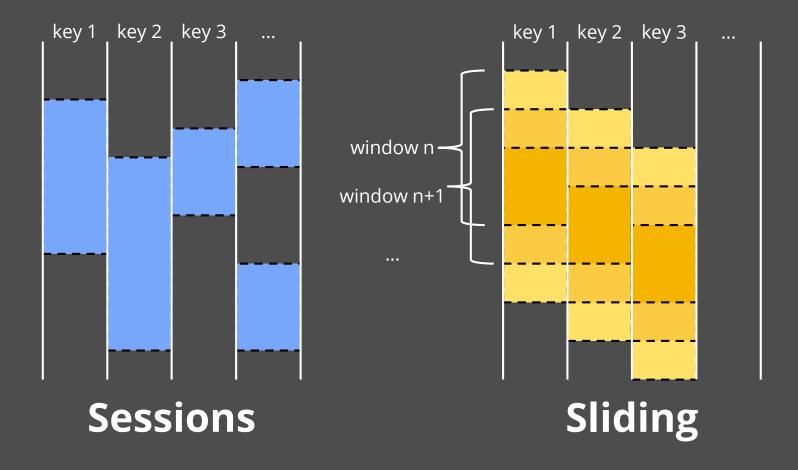
Pros: Efficient

Cons: Inexact

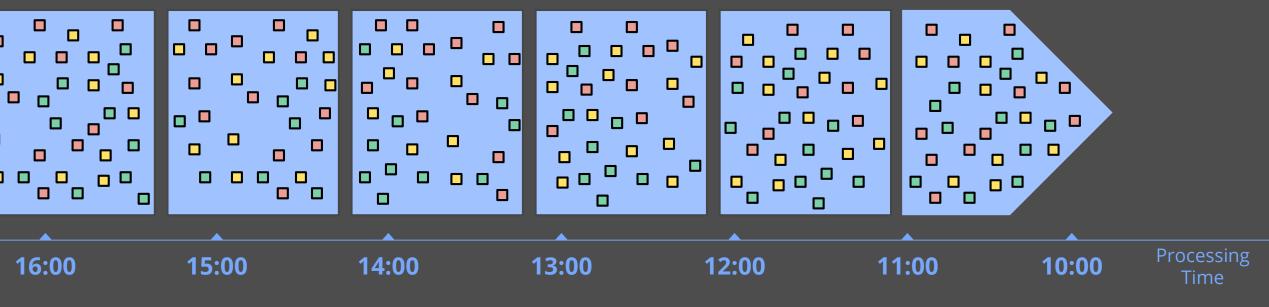
Complicated Algorithms

Windowing





3. Windowing by Processing Time



Example Input: Web server request traffic

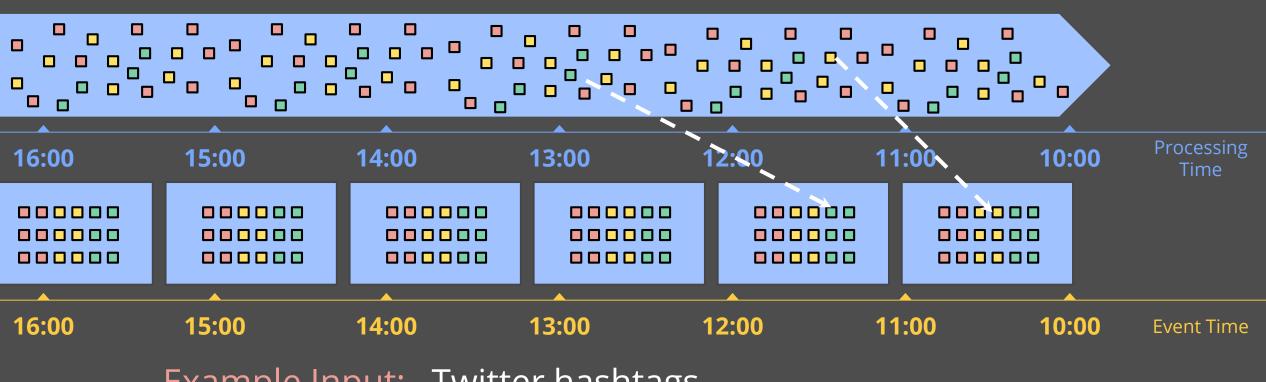
Example Output: Per-minute rate of received requests

Pros: Straightforward

Results reflect contents of stream

Cons: Results don't reflect events as they happened If approximating event time, usefulness varies

4. Windowing by Event Time - Fixed Windows



Example Input: Twitter hashtags

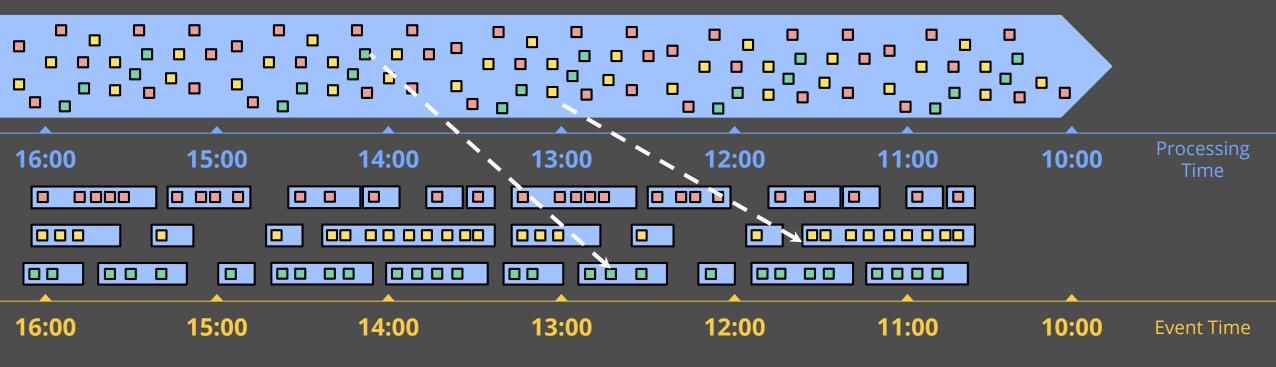
Example Output: Top N hashtags by prefix per hour.

Pros: Reflects events as they occurred

Cons: More complicated buffering

Completeness issues

4. Windowing by Event Time - Sessions



Example Input: User activity stream

Example Output: Per-session group of activities

Pros: Reflects events as they occurred

Cons: More complicated buffering

Completeness issues

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