$n \cdot Q(k, n) \leq \binom{k-1}{n-1}$ Claim # dishibuts h I'l Heigr to a sid reseptents, # distiluting eath gets at least 1. lethings to n isentral recepients s. t., each gots at last me { no repet r permete rose. then paramete recepients. Obvious since purshing he into a district part, and bleau parmute is "included" in distributing be rit in parts blackstack s.t. and part gets at least one. (I really den't hern lun alse to emplois this...) Pthyn) & (K+M, N) & lest w/ my Q(k,n)=7 > So Q(k,n)=Q(k-n,n). If all justs bare size ?? Her remove lagest whom, and = set a(h-n,h) n-1 ports -> get & (k-n,n-1)

Supp 5	P(n,z) = ?
7.	z U
	If n is every then can split my n
	or have n-1 aptions
	-) get " choices
	If n is odd, then have [n] side clivius
	So $P(n,z) = \begin{bmatrix} n \\ z \end{bmatrix}$ $n \ge 0$.
Supp 6	If (k = 3) then P(h, h-2) = P(3, 1) = 1
	IF (k ? 3) Kun P(k, h-2) = EP(k-(k-2), 5)
	Σ P (2, 2°)
	Con also
	lut this is timind

() Shm: P(h,n) (gtits fine at next me) - P (mn-k... sinto no more them n perto of size of most m-1). For ruch justifien of kints

n ports sit maxe size sis m,

(unique)

Ne have a justifien of

(mn-k) sits no more flean

n ports of size of most

m-/ (ley Yong's diaprom) So, # perhitions of kints in ports s.t. near is equal to the # portions of (mn-h) into no neare them n perts of fire at most m-1.