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1 Contest

1.1 C++

```
1 #include <bits/stdc++.h>
2 using namespace std;
4 #ifdef LOCAL
5 #include "cp/debug.h"
6 #else
7 #define debug(...)
8 #endif
nt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
12 int main() {
      ios::sync_with_stdio(false); cin.tie(nullptr);
      // freopen("input.txt", "r", stdin);
14
      // freopen("output.txt", "w", stdout);
15
      return 0:
17
18 }
1.2 Debug
1 #define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
       ',', ''); stringstream _ss(_s); istream_iterator<string> _it(_ss);
      out_error(_it, __VA_ARGS__);}
3 void out_error(istream_iterator<string> it) { cerr << '\n'; }</pre>
5 template<typename T, typename ...Args>
6 void out_error(istream_iterator<string> it, T a, Args... args) {
      cerr << " [" << *it << " = " << a << "] ";
      out_error(++it, args...);
8
9 }
11 template<typename T, typename G> ostream& operator<<(ostream &os, const
      pair < T. G > &p) {
      return os << "(" << p.first << ", " << p.second << ")";</pre>
12
13 }
15 template < class Con, class = decltype(begin(declval < Con > ())) >
16 typename enable_if<!is_same<Con, string>::value, ostream&>::type
operator << (ostream& os, const Con& container) {</pre>
      os << "{";
18
      for (auto it = container.begin(); it != container.end(); ++it)
          os << (it == container.begin() ? "" : ", ") << *it;
      return os << "}";</pre>
21
22 }
1.3
      Java
import java.io.BufferedReader;
2 import java.util.StringTokenizer;
3 import java.io.IOException;
```

```
4 import java.io.InputStreamReader;
5 import java.io.PrintWriter;
6 import java.util.ArrayList;
7 import java.util.Arrays;
8 import java.util.Collections;
9 import java.util.Random;
11 public class Main {
      public static void main(String[] args) {
          FastScanner fs = new FastScanner();
          PrintWriter out = new PrintWriter(System.out);
          int n = fs.nextInt();
          out.println(n);
          out.close(); // don't forget this line.
      static class FastScanner {
          BufferedReader br;
          StringTokenizer st:
          public FastScanner() {
              br = new BufferedReader(new InputStreamReader(System.in));
              st = null;
          public String next() {
              while (st == null || st.hasMoreTokens() == false) {
                  try {
                      st = new StringTokenizer(br.readLine());
                  }
                  catch (IOException e) {
                      throw new RuntimeException(e);
              return st.nextToken();
          public int nextInt() {
              return Integer.parseInt(next());
          public long nextLong() {
              return Long.parseLong(next());
          public double nextDouble() {
              return Double.parseDouble(next());
      sublime-build
1 {
      "cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall",
      "-Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],
      "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$",
```

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```
"working_dir": "${file_path}",
      "selector": "source.cpp, source.c++"
5
6 }
1.5
     bashrc
1 alias c++='g++ -std=c++2a -fmax-errors=5 -DLOCAL -Wall -Wextra -O2 -s'
3 #Stress-testing
4 function test {
    SOL = 1
    CHECKER=$2
    for i in {1..100};
        ./gen.out > in && $CHECKER < in > ans && $SOL < in > out && diff -Zb
      out ans && echo "Test $i passed!!" || break;
      done
10
11 }
```

2 Data structures

2.1 Sparse table

```
1 int st[MAXN][K + 1];
2 for (int i = 0; i < N; i++) {
      st[i][0] = f(array[i]);
3
4 }
5 for (int j = 1; j <= K; j++) {</pre>
      for (int i = 0; i + (1 << j) <= N; i++) {
          st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
8
9 }
10 // Range Minimum Queries.
int lg[MAXN + 1];
12 lg[1] = 0;
13 for (int i = 2; i \le MAXN; i++) {
      lg[i] = lg[i / 2] + 1;
15 }
int j = lg[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
18 // Range Sum Queries.
19 long long sum = 0;
20 for (int j = K; j >= 0; j--) {
      if ((1 << j) <= R - L + 1) {
          sum += st[L][j];
22
          L += 1 << j;
23
      }
24
25 }
```

2.2 Ordered set

```
1 #include <ext/pb_ds/assoc_container.hpp>
2 #include <ext/pb_ds/tree_policy.hpp>
3 using namespace __gnu_pbds;
4
5 template<typename key_type>
```

```
6 using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
       tree_order_statistics_node_update>;
9 const int INF = 0x3f3f3f3f;
10 void example() {
       vector < int > nums = \{1, 2, 3, 5, 10\};
       set_t<int> st(nums.begin(), nums.end());
       cout << *st.find_by_order(0) << '\n'; // 1</pre>
       assert(st.find_by_order(-INF) == st.end());
       assert(st.find_by_order(INF) == st.end());
       cout << st.order_of_key(2) << '\n'; // 1</pre>
       cout << st.order_of_key(4) << '\n'; // 3</pre>
       cout << st.order_of_key(9) << '\n'; // 4</pre>
       cout << st.order_of_key(-INF) << '\n'; // 0</pre>
       cout << st.order_of_key(INF) << '\n'; // 5</pre>
 2.3
       Dsu
1 struct Dsu {
       int n;
       vector<int> par, sz;
       Dsu(int _n) : n(_n) {
           sz.resize(n, 1);
           par.resize(n);
           iota(par.begin(), par.end(), 0);
       int find(int v) {
           // finding leader/parrent of set that contains the element v.
           // with {path compression optimization}.
           return (v == par[v] ? v : par[v] = find(par[v]));
12
       bool same(int u, int v) {
           return find(u) == find(v);
       bool unite(int u, int v) {
           u = find(u); v = find(v);
           if (u == v) return false;
           if (sz[u] < sz[v]) swap(u, v);
           par[v] = u;
           sz[u] += sz[v];
22
           return true;
       vector<vector<int>> groups() {
           // returns the list of the "list of the vertices in a connected
       component".
           vector<int> leader(n);
           for (int i = 0; i < n; ++i) {</pre>
               leader[i] = find(i);
           vector < int > id(n, -1);
           int count = 0;
```

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```
for (int i = 0; i < n; ++i) {
33
              if (id[leader[i]] == -1) {
34
                  id[leader[i]] = count++;
          }
          vector<vector<int>> result(count);
          for (int i = 0; i < n; ++i) {
              result[id[leader[i]]].push_back(i);
41
          return result;
42
44 };
      Segment tree
1 /**
2 * Description: A segment tree with range updates and sum queries that
       supports three types of operations:
* + Increase each value in range [1, r] by x (i.e. a[i] += x).
     + Set each value in range [l, r] to x (i.e. a[i] = x).
     + Determine the sum of values in range [1, r].
7 struct SegmentTree {
      int n:
      vector<long long> tree, lazy_add, lazy_set;
      SegmentTree(int _n) : n(_n) {
10
11
          int p = 1;
          while (p < n) p *= 2;
12
          tree.resize(p * 2);
13
          lazy_add.resize(p * 2);
14
          lazy_set.resize(p * 2);
15
16
      long long merge(const long long &left, const long long &right) {
17
          return left + right;
18
      void build(int id, int 1, int r, const vector<int> &arr) {
20
          if (1 == r) {
21
              tree[id] += arr[l];
22
              return:
23
24
          int mid = (1 + r) >> 1;
25
          build(id * 2, 1, mid, arr);
26
          build(id * 2 + 1, mid + 1, r, arr);
27
          tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
28
      }
29
      void push(int id, int l, int r) {
30
          if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
31
          int mid = (1 + r) >> 1;
32
          for (int child : {id * 2, id * 2 + 1}) {
33
              int range = (child == id * 2 ? mid - 1 + 1 : r - mid);
34
              if (lazy_set[id] != 0) {
35
                  lazy_add[child] = 0;
                  lazy_set[child] = lazy_set[id];
                   tree[child] = range * lazy_set[id];
38
```

```
lazy_add[child] += lazy_add[id];
              tree[child] += range * lazy_add[id];
          lazy_add[id] = lazy_set[id] = 0;
      void update(int id, int 1, int r, int u, int v, int amount, bool
      set_value = false) {
          if (r < u || 1 > v) return;
          if (u <= 1 && r <= v) {
              if (set_value) {
                  tree[id] = 1LL * amount * (r - l + 1);
                  lazy_set[id] = amount;
                  lazy_add[id] = 0; // clear all previous updates.
              }
              else {
                  tree[id] += 1LL * amount * (r - 1 + 1);
                  lazy_add[id] += amount;
              return;
          push(id, 1, r);
          int mid = (1 + r) >> 1;
          update(id * 2, 1, mid, u, v, amount, set_value);
          update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
          tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
      long long get(int id, int 1, int r, int u, int v) {
          if (r < u \mid \mid 1 > v) return 0;
          if (u <= 1 && r <= v) {
              return tree[id];
          push(id, 1, r);
          int mid = (1 + r) >> 1;
          long long left = get(id * 2, 1, mid, u, v);
          long long right = get(id * 2 + 1, mid + 1, r, u, v);
          return merge(left, right);
77 };
      Efficient segment tree
1 template < typename T> struct SegmentTree {
      int n;
      vector<T> tree:
      SegmentTree(int _n) : n(_n), tree(2 * n) {}
      T merge(const T &left, const T &right) {
          return left + right;
      template < typename G>
      void build(const vector<G> &initial) {
          assert((int) initial.size() == n);
          for (int i = 0; i < n; ++i) {
```

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```
tree[i + n] = initial[i];
12
                                                                                     28
                                                                                     29
                                                                                                int mid = (1 + r) >> 1;
13
14
          for (int i = n - 1; i > 0; --i) {
                                                                                                push(x, 1, mid, r);
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
                                                                                                int left = add(tree[x].1, 1, mid, u, v, amt);
15
                                                                                                int right = add(tree[x].r, mid + 1, r, u, v, amt);
          }
16
      }
                                                                                                tree.emplace_back(left, right, tree[left].val + tree[right].val, 0);
17
      void modify(int i, int v) {
                                                                                                return tree.size() - 1;
18
          tree[i += n] = v;
          for (i /= 2; i > 0; i /= 2) {
                                                                                           long long get_sum(int x, int l, int r, int u, int v) {
20
              tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
                                                                                               if (r < u \mid | 1 > v) return 0;
21
22
          }
                                                                                                if (u <= 1 && r <= v) return tree[x].val;</pre>
      }
                                                                                                int mid = (1 + r) / 2;
23
      T get_sum(int 1, int r) {
                                                                                                push(x, 1, mid, r);
24
          // sum of elements from 1 to r - 1.
                                                                                               return get_sum(tree[x].1, 1, mid, u, v) + get_sum(tree[x].r, mid + 1,
25
          T ret{};
26
          for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2)
              if (1 & 1) ret = merge(ret, tree[1++]);
                                                                                            void push(int x, int 1, int mid, int r) {
28
                                                                                     43
              if (r & 1) ret = merge(ret, tree[--r]);
                                                                                                if (!tree[x].has_changed) return;
29
          }
                                                                                                Vertex left = tree[tree[x].1];
30
                                                                                                Vertex right = tree[tree[x].r];
31
          return ret;
                                                                                                tree.emplace_back(left);
                                                                                                tree[x].l = tree.size() - 1;
33 };
                                                                                                tree.emplace_back(right);
     Persistent lazy segment tree
                                                                                                tree[x].r = tree.size() - 1;
1 struct Vertex {
                                                                                                tree[tree[x].1].val += tree[x].lazy * (mid - 1 + 1);
      int 1, r;
                                                                                                tree[tree[x].1].lazy += tree[x].lazy;
      long long val, lazy;
      bool has_changed = false;
                                                                                                tree[tree[x].r].val += tree[x].lazy * (r - mid);
      Vertex(int _1, int _r, long long _val, int _lazy = 0) : l(_1), r(_r),
                                                                                                tree[tree[x].r].lazy += tree[x].lazy;
      val(_val), lazy(_lazy) {}
                                                                                                tree[tree[x].1].has_changed = true;
7 };
                                                                                                tree[tree[x].r].has_changed = true;
8 struct PerSegmentTree {
                                                                                                tree[x].lazy = 0;
      vector<Vertex> tree;
                                                                                                tree[x].has_changed = false;
      vector<int> root;
10
      int build(const vector<int> &arr, int 1, int r) {
11
          if (1 == r) {
                                                                                     63 };
12
              tree.emplace_back(-1, -1, arr[1]);
13
                                                                                      2.7 Lichao tree
              return tree.size() - 1;
14
                                                                                     1 /**
15
          int mid = (1 + r) / 2;
                                                                                     2 * Description: A segment tree that allows insert a new line and query for
          int left = build(arr, 1, mid);
                                                                                            minimum value over all lines at point x.
17
          int right = build(arr, mid + 1, r);
                                                                                      * Usage: useful in convex hull trick.
          tree.emplace_back(left, right, tree[left].val + tree[right].val);
                                                                                      4 */
          return tree.size() - 1;
20
                                                                                     6 const long long INF_LL = (long long) 4e18;
21
      int add(int x, int 1, int r, int u, int v, int amt) {
22
          if (1 > v \mid | r < u) return x:
                                                                                     8 struct Line {
23
          if (u <= 1 && r <= v) {
                                                                                           long long a, b;
24
              tree.emplace_back(tree[x].1, tree[x].r, tree[x].val + 1LL * amt *
                                                                                           Line(long long a = 0, long long b = INF_LL): a(a), b(b) {}
25
      (r - l + 1), tree[x].lazy + amt);
                                                                                           long long operator()(long long x) {
```

12

13

}

return a * x + b;

tree.back().has_changed = true;

return tree.size() - 1;

27

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```
14 };
                                                                                       23
15
                                                                                       24
16 struct SegmentTree { // min query
      int n;
      vector<Line> tree;
18
      SegmentTree() {}
20
      SegmentTree(int _n): n(1) {
           while (n < _n) n *= 2;</pre>
           tree.resize(n * 2);
22
23
24
      void insert(int x, int l, int r, Line line) {
          if (1 == r) {
25
               if (line(l) < tree[x](l)) tree[x] = line;</pre>
26
               return:
27
28
           int mid = (1 + r) >> 1;
          bool b_left = line(l) < tree[x](l);</pre>
30
          bool b_mid = line(mid) < tree[x](mid);</pre>
31
          if (b_mid) swap(tree[x], line);
          if (b_left != b_mid) insert(x * 2, 1, mid, line);
33
           else insert(x * 2 + 1, mid + 1, r, line);
35
      long long query(int x, int 1, int r, int at) {
36
          if (l == r) return tree[x](at);
37
          int mid = (1 + r) >> 1;
38
          if (at <= mid) return min(tree[x](at), query(x * 2, 1, mid, at));</pre>
           else return min(tree[x](at), query(x * 2 + 1, mid + 1, r, at));
42 };
      Old driver tree (Chtholly tree)
1 /**
2 * Description: An optimized brute-force approach to deal with problem that
      has operation of setting an interval to the same number.
     Note: caution TLE, only works when input is random
4 */
5 struct ODT {
      map<int, long long> tree;
      using It = map<int, long long>::iterator;
      It split(int x) {
          It it = tree.upper_bound(x);
10
           assert(it != tree.begin());
11
          if (it->first == x) return it;
13
          return tree.emplace(x, it->second).first;
14
      }
15
      void add(int 1, int r, int amt) {
17
          It it_l = split(l);
18
          It it_r = split(r + 1);
19
          while (it_l != it_r) {
               it_l->second += amt;
21
```

```
++it_1;
     }
      void set(int 1, int r, int v) {
          It it_l = split(l);
          It it_r = split(r + 1);
          while (it_l != it_r) {
              tree.erase(it_l++);
          tree[1] = v;
      long long kth_smallest(int 1, int r, int k) {
          // return the k-th smallest value in range [l..r]
          vector<pair<long long, int>> values; // pair(value, count)
          It it_l = split(l);
          It it_r = split(r + 1);
          while (it_l != it_r) {
              It prev = it l++:
              values.emplace_back(prev->second, it_l->first - prev->first);
          sort(values.begin(), values.end());
          for (auto [value, cnt] : values) {
              if (k <= cnt) return value;</pre>
              k -= cnt;
          }
          return -1;
      int powmod(long long a, long long n, int mod) {}
      int sum_of_xth_power(int 1, int r, int x, int mod) {
          It it_l = split(l);
          It it_r = split(r + 1);
          int res = 0;
          while (it_l != it_r) {
              It prev = it_l++;
              res = (res + 1LL * (it_l->first - prev->first) *
      powmod(prev->second, x, mod)) % mod;
         }
          return res;
      Disjoint sparse table
2 * Description: range query on a static array.
  * Time: O(1) per query.
4 * Tested: stress-test.
6 const int MOD = (int) 1e9 + 7;
7 struct DisjointSparseTable { // product queries.
      int n, h;
      vector<vector<int>> dst;
```

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```
vector<int> lg;
10
      DisjointSparseTable(int _n) : n(_n) {
11
12
          h = 1; // in case n = 1: h = 0 !!.
          int p = 1;
13
          while (p < n) p *= 2, h++;
14
          lg.resize(p); lg[1] = 0;
15
          for (int i = 2; i < p; ++i) {
               lg[i] = 1 + lg[i / 2];
18
          dst.resize(h, vector<int>(n));
19
20
      void build(const vector<int> &A) {
21
          for (int lv = 0; lv < h; ++lv) {
22
               int len = (1 << lv);</pre>
23
               for (int k = 0; k < n; k += len * 2) {
24
                   int mid = min(k + len, n);
                   dst[lv][mid - 1] = A[mid - 1] % MOD;
                   for (int i = mid - 2; i >= k; --i) {
                       dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
                   if (mid == n) break;
                   dst[lv][mid] = A[mid] % MOD;
                   for (int i = mid + 1; i < min(mid + len, n); ++i) {</pre>
                       dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
              }
          }
37
      }
      int get(int 1, int r) {
38
          if (1 == r) {
39
               return dst[0][1];
40
41
          int i = lg[l ^ r];
          return 1LL * dst[i][1] * dst[i][r] % MOD;
      }
45 };
      Fenwick tree
 2.10
2 * Description: range update and range sum query.
3 */
5 using tree_type = long long;
6 struct FenwickTree {
      int n;
      vector<tree_type> fenw_coeff, fenw;
      FenwickTree() {}
      FenwickTree(int _n) : n(_n) {
          fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
11
          fenw.assign(n, 0); // normal fenwick tree.
12
      }
13
      template<typename G>
14
      void build(const vector<G> &A) {
15
```

```
assert((int) A.size() == n);
          vector<int> diff(n);
17
          diff[0] = A[0];
          for (int i = 1; i < n; ++i) {</pre>
              diff[i] = A[i] - A[i - 1];
          fenw_coeff[0] = (long long) diff[0] * n;
          fenw[0] = diff[0];
          for (int i = 1; i < n; ++i) {
              fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n - i);
              fenw[i] = fenw[i - 1] + diff[i];
          for (int i = n - 1; i >= 0; --i) {
              int j = (i \& (i + 1)) - 1;
              if (j >= 0) {
                  fenw_coeff[i] -= fenw_coeff[j];
                  fenw[i] -= fenw[j];
          }
      void add(vector<tree_type> &fenw, int i, tree_type val) {
          while (i < n) {
              fenw[i] += val;
              i = (i + 1);
          }
      tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
          tree_type res{};
          while (i >= 0) {
              res += fenw[i];
              i = (i \& (i + 1)) - 1;
          }
          return res;
      tree_type prefix_sum(int i) {
          return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i -
      1);
      void range_add(int 1, int r, tree_type val) {
          add(fenw_coeff, 1, (n - 1) * val);
          add(fenw_coeff, r + 1, (n - r - 1) * (-val));
          add(fenw, l, val);
          add(fenw, r + 1, -val);
      tree_type range_sum(int 1, int r) {
          return prefix_sum(r) - prefix_sum(l - 1);
62 };
 2.11 Fenwick tree 2D
1 /**
* Description: range update and range sum query on a 2D array.
3 */
```

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```
8 // Binary search tree + min-heap.
4 using tree_type = long long;
5 struct FenwickTree2D {
                                                                                      9 struct Treap {
                                                                                            Node *root:
      int n, m;
      vector<vector<tree_type> > fenw[4];
                                                                                            Treap() : root(nullptr) {}
      FenwickTree2D(int _n, int _m) : n(_n), m(_m) {
                                                                                            int get_cnt(Node *n) { return n ? n->cnt : 0; }
          for (int i = 0; i < 4; i++) {
                                                                                            void upd_cnt(Node *&n) {
                                                                                                if (n) n->cnt = get_cnt(n->left) + get_cnt(n->right) + 1;
              fenw[i].resize(n, vector<tree_type>(m));
          }
11
12
      }
                                                                                            void push_rev(Node *treap) {
      void add(int u, int v, tree_type val) {
                                                                                                if (!treap || !treap->rev) return;
13
14
          for (int i = u; i < n; i |= (i + 1)) {
                                                                                                treap->rev = false;
            for (int j = v; j < m; j |= (j + 1)) {
                                                                                                swap(treap->left, treap->right);
15
              fenw[0][i][i] += val;
                                                                                                if (treap->left) treap->left->rev ^= true;
                                                                                     21
              fenw[1][i][j] += (u + 1) * val;
                                                                                                if (treap->right) treap->right->rev ^= true;
17
              fenw[2][i][j] += (v + 1) * val;
18
                                                                                            pair<Node*, Node*> split(Node *treap, int x, int smaller = 0) {
              fenw[3][i][j] += (u + 1) * (v + 1) * val;
                                                                                                if (!treap) return {};
20
          }
                                                                                                push_rev(treap);
21
      }
                                                                                                int idx = smaller + get_cnt(treap->left); // implicit val.
22
      void range_add(int r, int c, int rr, int cc, tree_type val) { // [r, rr]
                                                                                                if (idx <= x) {
23
      x [c, cc].
                                                                                                    auto pr = split(treap->right, x, idx + 1);
                                                                                                    treap->right = pr.first;
          add(r, c, val);
24
          add(r, cc + 1, -val);
                                                                                                    upd_cnt(treap);
25
                                                                                                    return {treap, pr.second};
26
          add(rr + 1, c, -val);
          add(rr + 1, cc + 1, val);
27
      }
                                                                                                else {
28
      tree_type prefix_sum(int u, int v) {
                                                                                                    auto pl = split(treap->left, x, smaller);
29
                                                                                                    treap -> left = pl.second;
30
          tree_type res{};
          for (int i = u; i >= 0; i = (i & (i + 1)) - 1) {
                                                                                                    upd_cnt(treap);
31
            for (int j = v; j >= 0; j = (j & (j + 1)) - 1) {
                                                                                                    return {pl.first, treap};
32
              res += (u + 2) * (v + 2) * fenw[0][i][j];
                                                                                               }
33
              res -= (v + 2) * fenw[1][i][j];
                                                                                           }
              res -= (u + 2) * fenw[2][i][j];
                                                                                            Node* merge(Node *1, Node *r) {
              res += fenw[3][i][j];
                                                                                                push_rev(1); push_rev(r);
                                                                                                if (!1 || !r) return (1 ? 1 : r);
            }
                                                                                     42
37
                                                                                                if (1->prior < r->prior) {
38
                                                                                                    1->right = merge(1->right, r);
          return res;
                                                                                                    upd_cnt(1);
      tree_type range_sum(int r, int c, int rr, int cc) { // [r, rr] x [c, cc].
                                                                                                    return 1;
41
          return prefix_sum(rr, cc) - prefix_sum(r - 1, cc) - prefix_sum(rr, c
                                                                                                }
42
      -1) + prefix_sum(r - 1, c - 1);
                                                                                                else {
                                                                                                    r->left = merge(1, r->left);
44 };
                                                                                                    upd_cnt(r);
                                                                                                    return r;
2.12 Implicit treap
                                                                                                }
1 struct Node {
                                                                                            void insert(int pos, int val) {
      int val, prior, cnt;
                                                                                                if (!root) {
      bool rev:
      Node *left, *right;
                                                                                                    root = new Node(val);
                                                                                                    return;
      Node() {}
      Node(int _val) : val(_val), prior(rng()), cnt(1), rev(false),
                                                                                                Node *1, *m, *r;
      left(nullptr), right(nullptr) {}
7 };
                                                                                                m = new Node(val);
```

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```
tie(l, r) = split(root, pos - 1);
61
           root = merge(l, merge(m, r));
62
63
      void erase(int pos_l, int pos_r) {
64
           Node *1, *m, *r;
           tie(l, r) = split(root, pos_l - 1);
           tie(m, r) = split(r, pos_r - pos_l);
           root = merge(1, r);
      }
      void reverse(int pos_l, int pos_r) {
70
71
           Node *1, *m, *r;
           tie(l, r) = split(root, pos_l - 1);
72
           tie(m, r) = split(r, pos_r - pos_l);
          m->rev ^= true:
74
           root = merge(l, merge(m, r));
75
      int query(int pos_l, int pos_r);
            // returns answer for corresponding types of query.
78
       void inorder(Node *n) {
79
           if (!n) return:
80
           push_rev(n);
82
           inorder(n->left);
           cout << n->val << ' ';
           inorder(n->right);
      void print() {
87
           inorder(root);
88
           cout << '\n';
90 };
        Line container
       Source: kactl
      Description: container that allow you can add lines in form 'ax + b' and
       query maximum value at 'x'.
4 **/
5 using num_t = int;
6 struct Line {
       num_t a, b; // ax + b
       mutable num_t x; // x-intersect with the next line in the hull
       bool operator < (const Line &other) const {</pre>
           return a < other.a;</pre>
10
11
12
      bool operator<(num_t other_x) const {</pre>
           return x < other_x;</pre>
13
14
15 };
17 struct LineContainer : multiset<Line, less<>> { // max-query
       // for doubles, use INF = 1 / 0.0
       static const num_t INF = numeric_limits<num_t>::max();
```

20

```
21
        num_t floor_div(num_t a, num_t b) {
            return a / b - ((a ^ b) < 0 && a % b != 0);
22
23
       bool isect(iterator u, iterator v) {
            if (v == end()) {
                u \rightarrow x = INF;
                return false:
            if (u->a == v->a) u->x = (u->b > v->b? INF : -INF);
            else u \rightarrow x = floor_div(v \rightarrow b - u \rightarrow b, u \rightarrow a - v \rightarrow a);
            return u \rightarrow x >= v \rightarrow x;
       void add(num_t a, num_t b) {
            auto z = insert({a, b, 0}), y = z++, x = y;
            while (isect(y, z)) z = erase(z);
            if (x != begin() && isect(--x, y)) {
                y = erase(y);
                isect(x, y);
            while ((y = x) != begin() && (--x)->x >= y->x) {
                isect(x, erase(y));
       }
        num_t query(num_t x) {
            assert(!empty());
            auto it = *lower_bound(x);
            return it.a * x + it.b:
49 };
```

3 Mathematics

3.1 Trigonometry

3.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u)\cos(v) \pm \cos(u)\sin(v)$$

$$\cos(u \pm v) = \cos(u)\cos(v) \mp \sin(u)\sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u)\tan(v)}$$

3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2\cos(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\cos(u) - \cos(v) = -2\sin(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

$$\sin(u) + \sin(v) = 2\sin(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\sin(u) - \sin(v) = 2\cos(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

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3.1.3 Product identities

$$\cos(u)\cos(v) = \frac{1}{2}[\cos(u+v) + \cos(u-v)]$$

$$\sin(u)\sin(v) = -\frac{1}{2}[\cos(u+v) - \cos(u-v)]$$

$$\sin(u)\cos(v) = \frac{1}{2}[\sin(u+v) + \sin(u-v)]$$

3.1.4 Double - triple angle identities

$$\sin(2u) = 2\sin(u)\cos(u)$$

$$\cos(2u) = 2\cos^{2}(u) - 1 = 1 - 2\sin^{2}(u)$$

$$\tan(2u) = \frac{2\tan(u)}{1 - \tan^{2}(u)}$$

$$\sin(3u) = 3\sin(u) - 4\sin^{3}(u)$$

$$\cos(3u) = 4\cos^{3}(u) - 3\cos(u)$$

$$\tan(3u) = \frac{3\tan(u) - \tan^{3}(u)}{1 - 3\tan^{2}(u)}$$

3.2 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$c + 2c^{2} + \dots + nc^{n} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c - 1)^{2}}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$1^{5} + 2^{5} + 3^{5} + \dots + n^{5} = \frac{n^{2}(n+1)^{2}(2n^{2} + 2n - 1)}{12}$$

$$1^{6} + 2^{6} + 3^{6} + \dots + n^{6} = \frac{n(n+1)(2n+1)(3n^{4} + 6n^{3} - 3n + 1)}{42}$$

$$1^{7} + 2^{7} + 3^{7} + \dots + n^{7} = \frac{n^{2}(n+1)^{2}(3n^{4} + 6n^{3} - n^{2} - 4n + 2)}{24}$$

4 String

4.1 Prefix function

```
2 * Description: The prefix function of a string 's' is defined as an array pi
      where pi[i] is the length of the longest proper prefix of the substring
      s[0..i] which is also a suffix of this substring.
  * Time complexity: O(|S|).
7 vector<int> prefix_function(const string &s) {
      int n = (int) s.length();
      vector<int> pi(n);
      pi[0] = 0;
      for (int i = 1; i < n; ++i) {
          int j = pi[i - 1]; // try length pi[i - 1] + 1.
          while (j > 0 \&\& s[j] != s[i]) {
              i = pi[i - 1];
          if (s[j] == s[i]) {
              pi[i] = i + 1;
      return pi;
      Counting occurrences of each prefix
#include "prefix_function.h"
vector<int> count_occurrences(const string &s) {
      vector<int> pi = prefix_function(s);
      int n = (int) s.size();
      vector<int> ans(n + 1);
      for (int i = 0; i < n; ++i) {
          ans[pi[i]]++;
      for (int i = n - 1; i > 0; --i) {
          ans[pi[i - 1]] += ans[i];
      for (int i = 0; i <= n; ++i) {</pre>
          ans[i]++;
      return ans;
      // Input: ABACABA
      // Output: 4 2 2 1 1 1 1
      Knuth-Morris-Pratt algorithm
    Searching for a substring in a string.
  * Time complexity: O(N + M).
6 #include "prefix_function.h"
8 vector<int> KMP(const string &text, const string &pattern) {
      int n = (int) text.length();
```

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```
int m = (int) pattern.length();
10
       string s = pattern + '$' + text;
11
                                                                                        39
12
      vector<int> pi = prefix_function(s);
      vector<int> indices;
13
      for (int i = 0; i < (int) s.length(); ++i) {</pre>
14
           if (pi[i] == m) {
15
               indices.push_back(i - 2 * m);
          }
      }
                                                                                        46 };
18
       return indices;
19
20 }
      Suffix array
struct SuffixArray {
       string s;
      int n, lim;
3
      vector<int> sa, lcp, rank;
      SuffixArray(const string &_s, int _{lim} = 256): s(_s), n(s.length() + 1),
                                                                                               }
       lim(_lim), sa(n), lcp(n), rank(n) {
          s += '$';
          build();
           kasai();
           sa.erase(sa.begin());
          lcp.erase(lcp.begin());
10
           s.pop_back();
11
12
      }
      void build() {
13
           vector<int> nrank(n), norder(n), cnt(max(n, lim));
14
           for (int i = 0; i < n; ++i) {
15
               sa[i] = i; rank[i] = s[i];
17
           for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = \max(1, k * 2),
                                                                                               }
18
       lim = rank_cnt + 1) {
               // counting sort.
               for (int i = 0; i < n; ++i) norder[i] = (sa[i] - k + n) % n;
20
               for (int i = 0; i < n; ++i) cnt[rank[i]]++;</pre>
                                                                                        23 }
               for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];</pre>
22
               for (int i = n - 1; i >= 0; --i) sa[--cnt[rank[norder[i]]]] =
23
                                                                                         1 /**
       norder[i];
               rank[sa[0]] = rank\_cnt = 0;
24
               for (int i = 1; i < n; ++i) {
25
                   int u = sa[i], v = sa[i - 1];
26
                   int nu = u + k, nv = v + k;
                   if (nu >= n) nu -= n;
                   if (nv >= n) nv -= n;
29
                   if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
                                                                                         7 */
                   nrank[sa[i]] = rank_cnt;
32
               for (int i = 0; i < rank_cnt + 1; ++i) cnt[i] = 0;</pre>
               rank.swap(nrank);
34
          }
                                                                                        12
35
      }
                                                                                        13
      void kasai() {
37
```

```
for (int i = 0; i < n; ++i) rank[sa[i]] = i;
          for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
              int j = sa[rank[i] - 1];
              while (s[i + k] == s[j + k]) k++;
              lcp[rank[i]] = k;
          // Note: lcp[i] = longest common prefix(sa[i - 1], sa[i]).
      Suffix array slow
vector<int> suffix_array_slow(string &s) {
      s += '$':
      int n = (int) s.size();
      vector<int> order(n), rank(n);
      for (int i = 0; i < n; ++i) {
          order[i] = i; rank[i] = s[i];
      for (int k = 0; k < n; k = max(1, k * 2)) {
          stable_sort(order.begin(), order.end(), [&](int i, int j) {
              return make_pair(rank[i], rank[(i + k) % n]) < make_pair(rank[j],</pre>
      rank[(j + k) % n]);
         });
          vector<int> nrank(n):
          for (int i = 0, cnt = 0; i < n; ++i) {
              if (i > 0 && rank[order[i]] != rank[order[i - 1]]) ++cnt;
              else if (i > 0 && rank[(order[i] + k) % n] != rank[(order[i - 1]
      + k) % n]) ++cnt;
              nrank[order[i]] = cnt;
          rank.swap(nrank);
      s.pop_back(); order.erase(order.begin());
      return order:
      // Time complexity: O(N * log(N)^2).
      Manacher's algorithm
2 * Description: for each position, computes d[0][i] = half length of
3 longest palindrome centered on i (rounded up), d[1][i] = half length of
  longest palindrome centered on i and i - 1.
  * Time complexity: O(N).
  * Tested: https://judge.yosupo.jp/problem/enumerate_palindromes,
      stress-tested.
8 array<vector<int>, 2> manacher(const string &s) {
      int n = (int) s.size();
      array<vector<int>, 2> d;
      for (int z = 0; z < 2; ++z) {
          d[z].resize(n);
          int 1 = 0. r = 0:
          for (int i = 0; i < n; ++i) {
```

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```
int mirror = 1 + r - i + z;
15
               d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
16
               int L = i - d[z][i] - z, R = i + d[z][i];
               while (L >= 0 \&\& R < n \&\& s[L] == s[R]) {
                   d[z][i]++; L--; R++;
               }
20
               if (R > r) 
21
                   1 = L; r = R;
23
24
25
      }
      return d;
27 }
 4.7
      Trie
1 struct Trie {
      const static int ALPHABET = 26;
      const static char minChar = 'a';
3
      struct Vertex {
4
          int next[ALPHABET];
          bool leaf;
          Vertex() {
               leaf = false;
               fill(next, next + ALPHABET, -1);
          }
10
11
      };
      vector<Vertex> trie;
12
      Trie() { trie.emplace_back(); }
13
14
      void insert(const string &s) {
15
          int i = 0;
16
           for (const char &ch : s) {
17
               int j = ch - minChar;
18
               if (trie[i].next[j] == -1) {
                   trie[i].next[j] = trie.size();
                   trie.emplace_back();
22
               i = trie[i].next[j];
23
24
           trie[i].leaf = true;
25
26
      bool find(const string &s) {
27
          int i = 0;
28
           for (const char &ch : s) {
29
               int j = ch - minChar;
30
               if (trie[i].next[j] == -1) {
31
                   return false;
               i = trie[i].next[j];
34
          return (trie[i].leaf ? true : false);
38 };
```

4.8 Hashing

```
1 struct Hash61 {
      static const uint64_t MOD = (1LL << 61) - 1;</pre>
      static uint64_t BASE;
     static vector<uint64_t> pw;
      uint64_t addmod(uint64_t a, uint64_t b) const {
         a += b;
         if (a >= MOD) a -= MOD;
         return a;
      uint64_t submod(uint64_t a, uint64_t b) const {
         a += MOD - b;
         if (a >= MOD) a -= MOD;
         return a;
      uint64_t mulmod(uint64_t a, uint64_t b) const {
         uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
         uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
         uint64_t low = low1 * low2;
         uint64_t mid = low1 * high2 + low2 * high1;
         uint64_t high = high1 * high2;
         uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >> 29)
      + (mid << 35 >> 3) + 1;
         // ret %= MOD:
         ret = (ret >> 61) + (ret & MOD);
         ret = (ret >> 61) + (ret & MOD);
         return ret - 1;
     void ensure_pw(int m) {
         int sz = (int) pw.size();
         if (sz >= m) return;
         pw.resize(m);
         for (int i = sz; i < m; ++i) {</pre>
              pw[i] = mulmod(pw[i - 1], BASE);
         }
     }
     vector<uint64_t> pref;
      template < typename T > Hash61(const T &s) { // strings or arrays.
         n = (int) s.size();
         ensure_pw(n);
         pref.resize(n + 1);
         pref[0] = 0;
          for (int i = 0; i < n; ++i) {
              pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
     inline uint64_t operator()(const int from, const int to) const {
         assert(0 <= from && from <= to && to < n);
         // pref[to + 1] - pref[from] * pw[to - from + 1]
```

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```
return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
53
54 };
55 rng((unsigned int) chrono::steady_clock::now().time_since_epoch().count());
56 uint64_t Hash61::BASE = (MOD >> 2) + rng() % (MOD >> 1);
57 vector<uint64_t> Hash61::pw = vector<uint64_t>(1, 1);
      Minimum rotation
     Author: Stjepan Glavina
     License: Unlicense
     Source: https://github.com/stjepang/snippets/blob/master/min_rotation.cpp
     Description: Finds the lexicographically smallest rotation of a string.
     Time: O(N)
     Usage:
      rotate(v.begin(), v.begin()+minRotation(v), v.end());
     Status: Stress-tested
  */
11 #pragma once
12
int minRotation(string s) {
    int a = 0, N = (int) s.size(); s += s;
    rep(b, 0, N) rep(k, 0, N) {
      if (a + k == b \mid | s[a + k] < s[b + k]) \{b += max(0, k - 1); break;\}
      if (s[a + k] > s[b + k]) \{ a = b; break; \}
17
    return a;
19
```

5 Numerical

20 }

5.1 Fast Fourier transform

```
1 const double PI = acos(-1);
2 using Comp = complex < double >;
3 int reverse_bit(int n, int lg) {
       int res = 0;
       for (int i = 0; i < lg; ++i) {</pre>
           if (n & (1 << i)) {
               res |= (1 << (lg - i - 1));
      }
10
       return res;
11 }
void fft(vector<Comp> &a, bool invert = false) {
      int n = (int) a.size();
13
14
      int lq = 0;
      while (1 << (lg) < n) ++lg;</pre>
15
       for (int i = 0; i < n; ++i) {
           int rev_i = reverse_bit(i, lg);
18
           if (i < rev_i) swap(a[i], a[rev_i]);
       for (int len = 2; len <= n; len *= 2) {</pre>
20
           double angle = 2 * PI / len * (invert ? -1 : 1);
21
```

```
Comp w_base(cos(angle), sin(angle));
22
23
           for (int i = 0; i < n; i += len) {</pre>
24
               Comp w(1);
               for (int j = i; j < i + len / 2; ++j) {
                   Comp u = a[j], v = a[j + len / 2];
                   a[j] = u + w * v;
                   a[i + len / 2] = u - w * v;
                   w *= w_base;
               }
          }
      if (invert) for (int i = 0; i < n; ++i) a[i] /= n;
35 vector<int> mult(vector<int> &a, vector<int> &b) {
       vector<Comp> A(a.begin(), a.end()), B(b.begin(), b.end());
       int n = (int) a.size(), m = (int) b.size(), p = 1;
       while (p < n + m) p *= 2;
       A.resize(p), B.resize(p);
       fft(A, false);
       fft(B. false):
       for (int i = 0; i < p; ++i) {
          A[i] *= B[i];
       fft(A, true);
       vector < int > res(n + m - 1);
       for (int i = 0; i < n + m - 1; ++i) {
           res[i] = (int) round(A[i].real());
      }
       return res;
51 }
```

6 Number Theory

6.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n.
- Properties:
 - Divisor sum property: $\sum_{d|n} \phi(d) = n$.
 - $\phi(n)$ is a **prime number** when n = 3, 4, 6.
 - If *p* is a prime number, then $\phi(p) = p 1$.
 - If *p* is a prime number and $k \ge 1$, then $\phi(p^k) = p^k p^{k-1}$.
 - If *a* and *b* are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.
 - In general, for **not coprime** a and b, with d = gcd(a, b) this equation holds: $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$.

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- With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\phi(n) = \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m})$$
$$= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)$$

- Application in Euler's theorem:
 - If gcd(a, M) = 1, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod{\phi(M)}} \pmod{M}$$

- In general, for arbitrary a, M and n ≥ $\log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

6.2 Mobius function

• For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1\\ 0, & \text{if } \exists k_i > 1\\ (-1)^m & \text{otherwise} \end{cases}$$

- Properties:
 - $-\sum_{d|n}\mu(d)=[n=1].$
 - If *a* and *b* are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.
 - Mobius inversion: let *f* and *g* be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right)g(d)$$

6.3 Primes

Approximating the number of primes up to *n*:

n	$\pi(n)$	$\frac{n}{\ln n - 1}$
$100 (1e^2)$	25	28
$500 (5e^2)$	95	96
$1000 (1e^3)$	168	169
$5000 (5e^3)$	669	665
$10000 (1e^4)$	1229	1218
$50000 (5e^4)$	5133	5092
$100000 (1e^5)$	9592	9512
$500000 (5e^5)$	41538	41246
$1000000 (1e^6)$	78498	78030
$5000000 (5e^6)$	348513	346622

 $(\pi(n))$ = the number of primes less than or equal to n, $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

6.4 Wilson's theorem

A positive integer *n* is a prime if and only if:

$$(n-1)! \equiv n-1 \pmod{n}$$

6.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example:

$$64 = 55 + 8 + 1$$
$$85 = 55 + 21 + 8 + 1$$

```
vector<int> zeckendoft_theorem(int n) {
       vector < int > fibs = \{1, 1\};
       int sz = 2;
       while (fibs.back() <= n) {</pre>
           fibs.push_back(fibs[sz - 1] + fibs[sz - 2]);
           SZ++;
      }
       fibs.pop_back();
       vector<int> nums;
       int p = sz - 1;
       while (n > 0) {
           if (n >= fibs[p]) {
               nums.push_back(fibs[p]);
               n -= fibs[p];
           p--;
18
       return nums;
19 }
```

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6.6 Bitwise operation

```
• a + b = (a \oplus b) + 2(a \& b)

• a \mid b = (a \oplus b) + (a \& b)

• a \& (b \oplus c) = (a \& b) \oplus (a \& c)

• a \mid (b \& c) = (a \mid b) \& (a \mid c)

• a \& (b \mid c) = (a \& b) \mid (a \& c)

• a \& (b \mid c) = (a \& b) \mid (a \& c)

• a \mid (a \& b) = a

• a \& (a \mid b) = a
```

• Iterating over all subsets of a set and iterating over all submasks of a mask:

```
1 int n;
void mask_example() {
      for (int mask = 0; mask < (1 << n); ++mask) {
          for (int i = 0; i < n; ++i) {
              if (mask & (1 << i)) {
                  // do something...
          // Time complexity: 0(n * 2^n).
10
      for (int mask = 0; mask < (1 << n); ++mask) {
11
          for (int submask = mask; ; submask = (submask - 1) & mask) {
12
              // do something...
13
              if (submask == 0) break;
14
15
          // Time complexity: 0(3^n).
17
18 }
```

6.7 Pollard's rho algorithm

```
1 using num_t = long long;
const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.</pre>
3 const int LIMIT = (int) 1e9;
4 vector<int> primes;
5 int small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 73, 113,
      193, 311, 313, 407521, 299210837};
6 void linear_sieve(int n);
7 num_t mulmod(num_t a, num_t b, num_t mod);
8 num_t powmod(num_t a, num_t n, num_t mod);
9 bool miller_rabin(num_t a, num_t d, int s, num_t mod) {
      num_t x = powmod(a, d, mod);
      if (x == mod - 1 || x == 1) {
11
          return true;
12
13
      for (int i = 0; i < s - 1; ++i) {
14
          x = mulmod(x, x, mod);
          if (x == mod - 1) return true;
16
      }
17
      return false:
18
19 }
20 bool is_prime(num_t n, int tests = 10) {
      if (n < 4) return (n > 1);
```

```
num_t d = n - 1;
      int s = 0;
       while (d % 2 == 0) { d >>= 1; s++; }
       for (int i = 0; i < tests; ++i) {</pre>
           int a = small_primes[i];
          if (n == a) return true;
          if (n % a == 0 || !miller_rabin(a, d, s, n)) return false;
      return true;
31 }
|_{32} num_t f(num_t x, int c, num_t mod) { // f(x) = (x^2 + c) % mod.
      x = mulmod(x, x, mod);
      x += c;
      if (x >= mod) x -= mod;
38 num_t pollard_rho(num_t n, int c) {
       // algorithm to find a random divisor of 'n'.
       // using random function: f(x) = (x^2 + c) \% n.
      num_t x = 2, y = x, d;
      long long p = 1;
      int dist = 0;
       while (true) {
          y = f(y, c, n);
          dist++:
          d = \_gcd(llabs(x - y), n);
          if (d > 1) break;
           if (dist == p) { dist = 0; p *= 2; x = y; }
       return d;
52 }
|sa void factorize(int n, vector<num_t> &factors);
54 void llfactorize(num_t n, vector<num_t> &factors) {
      if (n < 2) return;</pre>
       if (is_prime(n)) {
           factors.emplace_back(n);
           return;
      if (n < LIMIT) {</pre>
           factorize(n, factors);
           return;
      num_t d = n;
       for (int c = 2; d == n; c++) {
           d = pollard_rho(n, c);
      llfactorize(d, factors);
      llfactorize(n / d, factors);
|71 vector<num_t> gen_divisors(vector<pair<num_t, int>> &factors) {
       vector<num_t> divisors = {1};
       for (auto &x : factors) {
           int sz = (int) divisors.size();
```

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```
mu[x] = -1;
          for (int i = 0; i < sz; ++i) {
              num_t cur = divisors[i];
                                                                                     24
                                                                                                    func[x] = 1;
76
              for (int j = 0; j < x.second; ++j) {
                                                                                                    cnt[x] = 1;
                  cur *= x.first;
                                                                                                    pw[x] = x;
                  divisors.push_back(cur);
                                                                                                for (int p : primes) {
              }
                                                                                                    if (p > spf[x] \mid \mid x * p > n) break;
                                                                                                    spf[x * p] = p;
      return divisors; // this array is NOT sorted yet.
                                                                                                    if (p == spf[x]) {
83
84 }
                                                                                                        phi[x * p] = phi[x] * p;
                                                                                                        mu[x * p] = 0;
      Segment divisor sieve
                                                                                                        func[x * p] = func[x / pw[x]] * (cnt[x] + 1);
1 const int MAXN = (int) 1e6; // R - L + 1 <= N.
                                                                                                        cnt[x * p] = cnt[x] + 1;
2 int divisor_count[MAXN + 3];
                                                                                                        pw[x * p] = pw[x] * p;
3 void segment_divisor_sieve(long long L, long long R) {
                                                                                                   }
      for (long long i = 1; i \le (long long) sqrt(R); ++i) {
                                                                                                    else {
          long long start1 = ((L + i - 1) / i) * i;
                                                                                                        phi[x * p] = phi[x] * phi[p];
          long long start2 = i * i;
                                                                                                        mu[x * p] = mu[x] * mu[p]; // or -mu[x]
          long long j = max(start1, start2);
                                                                                                        func[x * p] = func[x] * func[p];
                                                                                                        cnt[x * p] = 1;
          if (j == start2) {
                                                                                                        pw[x * p] = p;
              divisor_count[j - L] += 1;
              j += i;
11
12
          for (; j <= R; j += i) {
              divisor_count[j - L] += 2;
                                                                                     47 }
13
          }
                                                                                            Bitset sieve
15
                                                                                     1 /**
16 }
                                                                                     2 * Description: sieve of eratosthenes for large n (up to 1e9).
      Linear sieve
                                                                                        * Time and space (tested on codeforces):
                                                                                        * + For n = 1e8: ~200 ms, 6 MB.
     Description: Finding primes and computing value for multiplicative
                                                                                     * + For n = 1e9: ~4000 \text{ ms}, 60 \text{ MB}.
      function in O(N)
                                                                                     7 const int N = (int) 1e8;
                                                                                     8 bitset<N / 2 + 1> isPrime;
                                                                                     9 void sieve(int n = N) {
5 const int N = (int) 1e6;
                                                                                           isPrime.flip();
6 bool is_prime[N + 1];
7 int spf[N + 1]; // smallest prime factor
                                                                                           isPrime[0] = false;
                                                                                            for (int i = 3; i <= (int) sqrt(n); i += 2) {
8 int phi[N + 1]; // euler's totient function
9 int mu[N + 1]; // mobius function
                                                                                                if (isPrime[i >> 1]) {
                                                                                                    for (int j = i * i; j <= n; j += 2 * i) {
10 int func[N + 1]; // a multiplicative function, f(p^k) = k
int cnt[N + 1]; // cnt[i] = the power of the smallest prime factor of i
                                                                                                        isPrime[j >> 1] = false;
12 int pw[N + 1]; // pw[i] = p^cnt[i] where p is the smallest prime factor of i
vector<int> primes;
                                                                                           }
15 void sieve(int n = N) {
      spf[0] = spf[1] = -1;
                                                                                     20 void example(int n) {
                                                                                            sieve(n):
      phi[1] = mu[1] = func[1] = 1;
17
      for (int x = 2; x <= n; ++x) {
                                                                                           int primeCnt = (n >= 2);
18
                                                                                            for (int i = 3; i \le n; i += 2) {
          if (spf[x] == 0) {
19
              primes.push_back(spf[x] = x);
                                                                                               if (isPrime[i >> 1]) {
20
              is_prime[x] = true;
                                                                                     25
                                                                                                    primeCnt++;
              phi[x] = x - 1;
                                                                                               }
```

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```
cout << primeCnt << '\n';</pre>
29 }
      Block sieve
     Description: very fast sieve of eratosthenes for large n (up to 1e9).
     Source: kactl.
     Time and space (tested on codeforces):
      + For n = 1e8: ~160 ms, 60 MB.
      + For n = 1e9: ~1600 ms, 505 MB.
     Need to check memory limit.
9 const int N = (int) 1e8;
10 bitset<N + 1> is_prime;
vector<int> fast_sieve() {
      const int S = (int) sqrt(N), R = N / 2;
      vector<int> primes = {2};
13
      vector<bool> sieve(S + 1, true);
14
      vector<array<int, 2>> cp;
15
      for (int i = 3; i \le S; i += 2) {
          if (sieve[i]) {
               cp.push_back({i, i * i / 2});
18
               for (int j = i * i; j \le S; j += 2 * i) {
                   sieve[j] = false;
20
22
      }
23
      for (int L = 1; L <= R; L += S) {
24
          array<bool, S> block{};
25
          for (auto &[p, idx] : cp) {
               for (; idx < S + L; idx += p) block[idx - L] = true;</pre>
27
28
          for (int i = 0; i < min(S, R - L); ++i) {
              if (!block[i]) primes.push_back((L + i) * 2 + 1);
32
      for (int p : primes) is_prime[p] = true;
      return primes;
35 }
       Combinatorics
 6.12.1 Catalan numbers
```

$$C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}, C_0 = 1, C_n = \frac{4n-2}{n+1} C_{n-1}$$

• The first 12 Catalan numbers (n = 0, 1, 2, ..., 11):

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:
 - difference binary search trees with n vertices from 1 to n.
 - rooted binary trees with n + 1 leaves (vertices are not numbered).
 - correct bracket sequence of length 2 * n.
 - permutation [n] with no 3-term increasing subsequence (i.e. doesn't exist i < j < k for which a[i] < a[j] < a[k]).
 - ways a convex polygon of n + 2 sides can split into triangles by connecting vertices.

6.12.2 Fibonacci numbers

$$F_n = \begin{cases} 0, & \text{if } n = 0\\ 1, & \text{if } n = 1\\ F_{n-1} + F_{n-2}, & \text{otherwise} \end{cases}$$

• The first 20 Fibonacci numbers (n = 0, 1, 2, ..., 19):

$$F_n = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181$$

• Properties:

$$F_{2n+1} = F_n^2 + F_{n+1}^2$$

$$F_{2n} = F_{n-1} \cdot F_n + F_n \cdot F_{n+1}$$

$$F_{n+1} \cdot F_{n-1} - F_n^2 = (-1)^n$$

$$n \mid m \iff F_n \mid F_m$$

$$(F_n, F_m) = F_{(n,m)}$$

6.12.3 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* non-empty groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=1}^{k} (-1)^{k-i} {k \choose i} i^n$$

6.12.4 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixied point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

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7 Geometry

7.1 Fundamentals

7.1.1 **Point**

```
1 #pragma once
3 const double PI = acos(-1);
4 const double EPS = 1e-9;
5 typedef double ftype;
6 struct Point {
       ftype x, y;
      Point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
      Point& operator+=(const Point& other) {
          x += other.x; y += other.y; return *this;
10
11
      Point& operator -= (const Point& other) {
12
          x -= other.x; y -= other.y; return *this;
13
14
      Point& operator*=(ftype t) {
15
          x *= t; y *= t; return *this;
      Point& operator/=(ftype t) {
18
          x /= t; y /= t; return *this;
19
20
      Point operator+(const Point& other) const {
21
           return Point(*this) += other;
      }
23
      Point operator-(const Point& other) const {
24
           return Point(*this) -= other;
25
      Point operator*(ftype t) const {
27
           return Point(*this) *= t;
28
29
      Point operator/(ftype t) const {
          return Point(*this) /= t;
31
32
      Point rotate(double angle) const {
           return Point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
       cos(angle));
35
      friend istream& operator>>(istream &in, Point &t);
       friend ostream& operator<<(ostream &out, const Point& t);</pre>
       bool operator<(const Point& other) const {</pre>
          if (fabs(x - other.x) < EPS)</pre>
               return y < other.y;</pre>
          return x < other.x:</pre>
41
43 };
45 istream& operator>>(istream &in, Point &t) {
      in >> t.x >> t.y;
       return in;
48 }
```

```
49 ostream& operator << (ostream &out, const Point& t) {
      out << t.x << ' ' << t.y;
      return out;
52 }
54 ftype dot(Point a, Point b) {return a.x * b.x + a.y * b.y;}
55 ftype norm(Point a) {return dot(a, a);}
56 ftype abs(Point a) {return sqrt(norm(a));}
ftype angle(Point a, Point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
58 ftype proj(Point a, Point b) {return dot(a, b) / abs(b);}
ftype cross(Point a, Point b) {return a.x * b.y - a.y * b.x;}
60 bool ccw(Point a, Point b, Point c) {return cross(b - a, c - a) > EPS;}
61 bool collinear(Point a, Point b, Point c) {return fabs(cross(b - a, c - a)) <
62 Point intersect(Point a1, Point d1, Point a2, Point d2) {
       double t = cross(a2 - a1, d2) / cross(d1, d2);
      return a1 + d1 * t;
65 }
 7.1.2 Line
1 #include "point.h"
3 struct Line {
       double a, b, c;
      Line (double _a = 0, double _b = 0, double _c = 0): a(_a), b(_b), c(_c) {}
      friend ostream & operator<<(ostream& out, const Line& 1);</pre>
s ostream & operator<<(ostream& out. const Line& 1) {</pre>
       out << 1.a << ' ' << 1.b << ' ' << 1.c;
      return out;
12 void PointsToLine(const Point& p1, const Point& p2, Line& 1) {
      if (fabs(p1.x - p2.x) < EPS)
          1 = \{1.0, 0.0, -p1.x\};
      else {
          1.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
          1.b = 1.0;
          1.c = -1.a * p1.x - 1.b * p1.y;
20 }
void PointsSlopeToLine(const Point& p, double m, Line& 1) {
      1.a = -m:
      1.b = 1:
24
      1.c = -1.a * p.x - 1.b * p.y;
25 }
26 bool areParallel(const Line& 11, const Line& 12) {
      return fabs(l1.a - l2.a) < EPS && fabs(l1.b - l2.b) < EPS;</pre>
28 }
29 bool areSame(const Line& 11, const Line& 12) {
       return areParallel(11, 12) && fabs(11.c - 12.c) < EPS;</pre>
32 bool areIntersect(Line 11, Line 12, Point& p) {
      if (areParallel(l1, l2)) return false;
      p.x = -(11.c * 12.b - 11.b * 12.c) / (11.a * 12.b - 11.b * 12.a);
```

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```
if (fabs(11.b) > EPS) p.y = -(11.c + 11.a * p.x);
      else p.y = -(12.c + 12.a * p.x);
36
      return 1;
38 }
39 double distToLine(Point p, Point a, Point b, Point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
      c = a + (b - a) * t;
      return abs(c - p);
43 }
44 double distToSegment(Point p, Point a, Point b, Point& c) {
      double t = dot(p - a, b - a) / norm(b - a);
      if (t > 1.0)
          c = Point(b.x, b.y);
      else if (t < 0.0)
          c = Point(a.x, a.y);
      else
          c = a + (b - a) * t;
51
      return abs(c - p);
52
53 }
54 bool intersectTwoSegment(Point a. Point b. Point c. Point d) {
      ftype ABxAC = cross(b - a, c - a);
      ftype ABxAD = cross(b - a, d - a);
      ftype CDxCA = cross(d - c, a - c);
      ftype CDxCB = cross(d - c, b - c);
      if (ABxAC == 0 \mid | ABxAD == 0 \mid | CDxCA == 0 \mid | CDxCB == 0) {
          if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;</pre>
          if (ABxAD == 0 &\& dot(a - d, b - d) <= 0) return true;
          if (CDxCA == 0 \&\& dot(c - a, d - a) <= 0) return true;
          if (CDxCB == 0 \&\& dot(c - b, d - b) <= 0) return true;
          return false;
      return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);</pre>
67 }
68 void perpendicular(Line 11, Point p, Line& 12) {
      if (fabs(l1.a) < EPS)</pre>
69
          12 = \{1.0, 0.0, -p.x\};
70
71
          12.a = -11.b / 11.a;
72
          12.b = 1.0;
73
          12.c = -12.a * p.x - 12.b * p.y;
74
75
76 }
7.1.3 Circle
1 #include "point.h"
3 int insideCircle(const Point& p, const Point& center, ftype r) {
      ftype d = norm(p - center);
      ftype rSq = r * r;
      return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
7 }
8 bool circle2PointsR(const Point& p1, const Point& p2, ftype r, Point& c) {
      double h = r * r - norm(p1 - p2) / 4.0;
      if (fabs(h) < 0) return false;</pre>
```

```
h = sqrt(h);
      Point perp = (p2 - p1).rotate(PI / 2.0);
      Point m = (p1 + p2) / 2.0;
      c = m + perp * (h / abs(perp));
      return true:
16 }
 7.1.4 Triangle
1 #include "point.h"
2 #include "line.h"
4 double areaTriangle(double ab, double bc, double ca) {
       double p = (ab + bc + ca) / 2;
       return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
8 double rInCircle(double ab, double bc, double ca) {
       double p = (ab + bc + ca) / 2;
      return areaTriangle(ab, bc, ca) / p;
11 }
12 double rInCircle(Point a, Point b, Point c) {
      return rInCircle(abs(a - b), abs(b - c), abs(c - a));
14 }
15 bool inCircle(Point p1, Point p2, Point p3, Point &ctr, double &r) {
      r = rInCircle(p1, p2, p3);
      if (fabs(r) < EPS) return false;</pre>
      Line 11, 12;
       double ratio = abs(p2 - p1) / abs(p3 - p1);
      Point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
      PointsToLine(p1, p, l1);
      ratio = abs(p1 - p2) / abs(p2 - p3);
      p = p1 + (p3 - p1) * (ratio / (1 + ratio));
      PointsToLine(p2, p, 12);
      areIntersect(l1, l2, ctr);
      return true;
28 double rCircumCircle(double ab, double bc, double ca) {
       return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
30 }
31 double rCircumCircle(Point a, Point b, Point c) {
       return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
 7.1.5 Convex hull
1 #include "point.h"
3 vector<Point> CH_Andrew(vector<Point> &Pts) { // overall O(n log n)
      int n = Pts.size(), k = 0;
       vector<Point> H(2 * n);
      sort(Pts.begin(), Pts.end());
       for (int i = 0; i < n; ++i) {
          while ((k \ge 2) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
          H[k++] = Pts[i];
      for (int i = n - 2, t = k + 1; i >= 0; --i) {
```

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```
while ((k \ge t) \& ! ccw(H[k - 2], H[k - 1], Pts[i])) --k;
12
          H[k++] = Pts[i];
13
14
      }
15
      H.resize(k);
      return H;
16
17 }
7.1.6 Polygon
1 #include "point.h"
2
3 double perimeter(const vector<Point> &P) {
      double ans = 0.0;
      for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
           ans += abs(P[i] - P[i + 1]);
      return ans:
8 }
9 double area(const vector<Point> &P) {
      double ans = 0.0;
10
11
      for (int i = 0; i < (int)P.size() - 1; ++i)
           ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
      return fabs(ans) / 2.0;
13
14 }
15 bool isConvex(const vector<Point> &P) {
      int n = (int)P.size();
      if (n <= 3) return false;</pre>
      bool firstTurn = ccw(P[0], P[1], P[2]);
18
      for (int i = 1; i < n - 1; ++i)
           if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
21
               return false;
      return true;
22
23 }
24 int insidePolygon(Point pt, const vector<Point> &P) {
      int n = (int)P.size();
      if (n <= 3) return -1;
      bool on_polygon = false;
27
      for (int i = 0; i < n - 1; ++i)
28
          if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1]))
      < EPS)
               on_polygon = true;
      if (on_polygon) return 0;
31
32
      double sum = 0.0;
      for (int i = 0; i < n - 1; ++i) {
          if (ccw(pt, P[i], P[i + 1]))
               sum += angle(P[i] - pt, P[i + 1] - pt);
           else
               sum -= angle(P[i] - pt, P[i + 1] - pt);
      }
38
      return fabs(sum) > PI ? 1 : -1;
39
40 }
       Minimum enclosing circle
```

1 /^^
2 * Description: computes the minimum Circle that encloses all the given
Points.

```
3 */
5 #include "point.h"
6 // #include "circle.h"
7 // TODO:
9 Point center_from(double bx, double by, double cx, double cy) {
       double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *
       return Point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
12 }
14 Circle Circle_from(Point A, Point B, Point C) {
      Point I = center_from(B.x - A.x, B.y - A.y, C.x - A.x, C.y - A.y);
       return Circle(I + A, abs(I));
17 }
19 const int N = 100005;
20 int n, x[N], y[N];
21 Point a[N];
23 Circle emo_welzl(int n, vector<Point> T) {
      if (T.size() == 3 || n == 0) {
          if (T.size() == 0) return Circle(Point(0, 0), -1);
          if (T.size() == 1) return Circle(T[0], 0);
          if (T.size() == 2) return Circle((T[0] + T[1]) / 2, abs(T[0] - T[1])
       / 2);
          return Circle_from(T[0], T[1], T[2]);
      random\_shuffle(a + 1, a + n + 1);
       Circle Result = emo_welzl(0, T);
       for (int i = 1; i <= n; i++)</pre>
          if (abs(Result.x - a[i]) > Result.y + 1e-9) {
              T.push_back(a[i]);
              Result = emo_welzl(i - 1, T);
              T.pop_back();
       return Result;
```

8 Linear algebra

8.1 Gauss elimination

```
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be infinity or a big number
int gauss (vector < vector<double> > a, vector<double> & ans) {
   int n = (int) a.size();
   int m = (int) a[0].size() - 1;
   vector<int> where (m, -1);
   for (int col=0, row=0; col<m && row<n; ++col) {
      int sel = row;
      for (int i=row; i<n; ++i) {
        if (abs (a[i][col]) > abs (a[sel][col])) sel = i;
   }
}
```

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```
11
           if (abs (a[sel][col]) < EPS) continue;</pre>
12
            for (int i=col; i<=m; ++i) {</pre>
13
                swap (a[sel][i], a[row][i]);
14
15
            where[col] = row;
            for (int i=0; i<n; ++i) {</pre>
                if (i != row) {
19
                     double c = a[i][col] / a[row][col];
20
21
                     for (int j=col; j<=m; ++j) {</pre>
                         a[i][j] -= a[row][j] * c;
                }
24
           }
25
            ++row;
       }
27
       ans.assign (m, 0);
28
       for (int i=0; i<m; ++i) {</pre>
29
           if (where[i] != -1) {
30
                ans[i] = a[where[i]][m] / a[where[i]][i];
           }
32
       }
33
       for (int i=0; i<n; ++i) {</pre>
34
            double sum = 0;
35
           for (int j=0; j<m; ++j) {</pre>
                sum += ans[j] * a[i][j];
           if (abs (sum - a[i][m]) > EPS) return 0;
      }
40
       for (int i=0; i<m; ++i) {</pre>
41
           if (where[i] == -1) return INF;
42
       }
43
       return 1;
45 }
```

9 Graph

9.1 Bellman-Ford algorithm

```
1 /**
2  * Description: single source shortest path in a weighted (negative or positive) directed graph.
3  * Time: O(N * M).
4  * Tested: https://open.kattis.com/problems/shortestpath3
5  */
6  const int64_t INF = (int64_t) 2e18;
7  struct Edge {
8    int u, v; // u -> v
9    int64_t w;
10  Edge() {}
11  Edge() {}
12  };
13  int n;
```

```
14 vector<Edge> edges;
15 vector<int64_t> bellmanFord(int s) {
      // dist[stating] = 0.
      // dist[u] = +INF, if u is unreachable.
      // dist[u] = -INF, if there is a negative cycle on the path from s to u.
      // -INF < dist[u] < +INF, otherwise.</pre>
      vector<int64_t> dist(n, INF);
       dist[s] = 0;
       for (int i = 0; i < n - 1; ++i) {
          bool any = false;
          for (auto [u, v, w] : edges) {
              if (dist[u] != INF && dist[v] > w + dist[u]) {
                   dist[v] = w + dist[u];
                   any = true;
          if (!any) break;
      // handle negative cycles
      for (int i = 0; i < n - 1; ++i) {
           for (auto [u, v, w] : edges) {
              if (dist[u] != INF && dist[v] > w + dist[u]) {
                   dist[v] = -INF;
              }
          }
      }
      return dist;
       Articulation point and Bridge
1 /**
2 * Description: finding articulation points and bridges in a simple
       undirected graph.
3 * Tested: https://oj.vnoi.info/problem/graph_
5 const int N = (int) 1e5;
6 vector<int> g[N];
7 int num[N], low[N], dfs_timer;
8 bool joint[N];
9 vector<pair<int, int>> bridges;
10 void dfs(int u, int prev) {
      low[u] = num[u] = ++dfs_timer;
      int child = 0;
       for (int v : g[u]) {
          if (v == prev) continue;
          if (num[v]) low[u] = min(low[u], num[v]);
              dfs(v, u);
              low[u] = min(low[u], low[v]);
              child++;
              if (low[v] >= num[v]) {
                   bridges.emplace_back(u, v);
              }
```

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```
if (u != prev && low[v] >= num[u]) joint[u] = true;
23
24
25
       if (u == prev && child > 1) joint[u] = true;
26
27 }
28
29 int solve() {
       int n, m;
       cin >> n >> m;
31
       for (int i = 0; i < m; ++i) {
32
33
           int u, v;
           cin >> u >> v;
34
           u--; v--;
           g[u].push_back(v);
           g[v].push_back(u);
37
38
       for (int i = 0; i < n; ++i) {
           if (!num[i]) dfs(i, i);
      }
41
42
       return 0;
```

9.3 Strongly connected components

```
1 /**
     Description: Tarjan's algorithm finds strongly connected components
2
      in a directed graph. If vertices u and v belong to the same component,
      then scc_id[u] == scc_id[v].
     Tested: https://judge.yosupo.jp/problem/scc
6 */
7 const int N = (int) 5e5;
8 vector<int> g[N], st;
9 int low[N], num[N], dfs_timer, scc_id[N], scc;
10 bool used[N];
void Tarjan(int u) {
      low[u] = num[u] = ++dfs_timer;
12
      st.push_back(u);
13
14
      for (int v : g[u]) {
           if (used[v]) continue;
15
          if (num[v] == 0) {
               Tarjan(v);
17
               low[u] = min(low[u], low[v]);
18
20
           else {
               low[u] = min(low[u], num[v]);
22
          }
23
      if (low[u] == num[u]) {
24
          int v:
25
           do {
26
               v = st.back(); st.pop_back();
27
               debug(u, v)
28
               used[v] = true;
               scc_id[v] = scc;
30
```

```
} while (v != u);
32
           scc++;
33
      }
34 }
      Topo sort
2 * Description: A topological sort of a directed acyclic graph
     is a linear ordering of its vertices such that for every directed edge
      from vertex u to vertex v, u comes before v in the ordering.
   * Note: If there are cycles, the returned list will have size smaller than n
       (i.e, topo.size() < n).
   * Tested: https://judge.yosupo.jp/problem/scc
8 vector<int> topo_sort(const vector<vector<int>> &g) {
       int n = (int) g.size();
       vector<int> indeq(n);
       for (int u = 0; u < n; ++u) {
           for (int v : g[u]) indeg[v]++;
12
       queue < int > q; // Note: use min-heap to get the smallest lexicographical
       for (int u = 0; u < n; ++u) {
          if (indeg[u] == 0) q.emplace(u);
       vector<int> topo;
       while (!q.empty()) {
          int u = q.front(); q.pop();
21
          topo.emplace_back(u);
           for (int v : g[u]) {
              if (--indeg[v] == 0) q.emplace(v);
24
25
      }
       return topo;
      K-th smallest shortest path
      Description: Finding the k-th smallest shortest path from vertex s to
       vertex t.
     each vertex can be visited more than once.
4 using adj_list = vector<vector<pair<int, int>>>;
5 vector<long long> k_smallest(const adj_list &g, int k, int s, int t) {
      int n = (int) g.size();
       vector<long long> ans;
       vector<int> cnt(n);
       using pli = pair<long long, int>;
       priority_queue<pli, vector<pli>, greater<pli>> pq;
      pq.emplace(0, s);
       while (!pq.empty() && cnt[t] < k) {</pre>
           int u = pq.top().second;
14
          long long d = pq.top().first;
          pq.pop();
```

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```
if (cnt[u] == k) continue;
16
           cnt[u]++;
17
          if (u == t) {
               ans.push_back(d);
20
          for (auto [v, cost] : g[u]) {
21
22
               pq.emplace(d + cost, v);
         }
      }
24
      assert(k == (int) ans.size());
25
26
      return ans;
27 }
      Eulerian path
 9.6.1
       Directed graph
     Hierholzer's algorithm.
     Description: An Eulerian path in a directed graph is a path that visits
       all edges exactly once.
      An Eulerian cycle is a Eulerian path that is a cycle.
     Time complexity: O(|E|).
6
   vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
7
      int n = (int) q.size();
      vector<int> stack, cur_edge(n), vertices;
      stack.push_back(s);
      while (!stack.empty()) {
11
12
          int u = stack.back();
          stack.pop_back();
          while (cur_edge[u] < (int) g[u].size()) {</pre>
14
               stack.push_back(u);
15
               u = g[u][cur\_edge[u]++];
          vertices.push_back(u);
      reverse(vertices.begin(), vertices.end());
20
      return vertices;
21
22 }
9.6.2
      Undirected graph
     Hierholzer's algorithm.
     Description: An Eulerian path in a undirected graph is a path that visits
      all edges exactly once.
      An Eulerian cycle is a Eulerian path that is a cycle.
     Time complexity: O(|E|).
6 */
7 struct Edge {
      list<Edge>::iterator reverse_edge;
      Edge(int _to) : to(_to) {}
11 };
12 vector<int> vertices;
```

```
13 void find_path(vector<list<Edge>> &g, int u) {
      while (!g[u].empty()) {
          int v = g[u].front().to;
          g[v].erase(g[u].front().reverse_edge);
          g[u].pop_front();
           find_path(g, v);
      vertices.emplace_back(u); // reversion list.
21 }
22 void add_edge(vector<list<Edge>> &g, int u, int v) {
      g[u].emplace_front(v);
      g[v].emplace_front(u);
      g[u].front().reverse_edge = g[v].begin();
      g[v].front().reverse_edge = g[u].begin();
27 }
      Misc.
 10
 10.1 Ternary search
 const double eps = 1e-9;
2 double f(double x) { return 42; }
3 double ternary_search_max(double 1, double r) {
      // find x0 such that: f(x0) > f(x), \all x: 1 \le x \le r.
      while (r - 1 > eps) {
          double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) < f(mid2)) 1 = mid1;
          else r = mid2;
      }
      return 1;
12 }
13 double ternary_search_min(double 1, double r) {
      // find x0 such that: f(x0) < f(x), \all x: 1 <= x <= r.
      while (r - 1 > eps) {
          double mid1 = 1 + (r - 1) / 3;
          double mid2 = r - (r - 1) / 3;
          if (f(mid1) > f(mid2)) 1 = mid1;
          else r = mid2;
      }
      return 1;
22 }
 10.2
      Ternary search 2
2 * Source: https://github.com/icpcvn/icpcvn.github.io/blob/master/2017/
   * national/solutions/E/RR_ternary.cpp
4 */
6 #include <bits/stdc++.h>
7 using namespace std;
8 #include "../geometry/point.h"
```

10 Point A, B, C;

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```
11 double f(Point P) {
      return abs(P - A) + abs(P - B) + abs(P - C);
12
13 }
14 double f(double x, double& y) {
      double 1 = -2000, r = 2000;
15
      for (int i = 0; i < 500; ++i) {
16
           double y1 = (1 * 2 + r) / 3.0;
17
           double y2 = (1 + r * 2) / 3.0;
19
           double f1 = f(Point(x, y1));
20
21
           double f2 = f(Point(x, y2));
22
          if (f1 < f2) r = y2;
           else l = y1;
24
      }
25
      y = (1 + r) / 2.0;
      return f(Point(x, y));
27
28 }
29 int32_t main() {
30
      ios::sync_with_stdio(0);
      cin.tie(0);
31
      cout << fixed << setprecision(9) << boolalpha;</pre>
32
      cin >> A >> B >> C;
33
34
      double 1 = -2000, r = 2000;
35
      for (int i = 0; i < 500; ++i) {
           double x1 = (1 * 2 + r) / 3.0;
37
38
           double x2 = (1 + r * 2) / 3.0;
           double y;
40
           double f1 = f(x1, y);
41
           double f2 = f(x2, y);
42
          if (f1 < f2) r = x2;
          else 1 = x1;
45
      }
      double y;
      f(1, y);
      cout << 1 << ' ' << y << '\n';
      return 0;
50
51 }
       Dutch flag national problem
void dutch_flag_national(vector<int> &arr) {
      // All elements that are LESS than pivot are moved to the LEFT.
      // All elements that are GREATER than pivot are moved to the RIGHT.
      // E.g. [1, 2, 0, 0, 2, 2, 1], pivot = 1 -> [0, 0, 1, 1, 2, 2, 2].
      int n = (int) arr.size();
      int i = 0, j = 0, k = n - 1;
      int pivot = 1;
      // 0....i....j....k....n
      while (i <= k) {
          if (arr[j] < pivot) {</pre>
10
```

```
swap(arr[i], arr[j]);
11
12
              i++;
              j++;
14
          else if (arr[j] > pivot) {
              swap(arr[j], arr[k]);
              k--;
          }
          else {
              j++;
          }
      // 0 <= index <= i - 1: arr[index] < mid.
      // i <= index <= k: arr[index] = mid.</pre>
      // k + 1 <= index < sz: arr[index] > mid.
 10.4 Matrix
using matrix_type = int;
2 const int MOD = (int) 1e9 + 7;
3 struct Matrix {
       static const matrix_type INF = numeric_limits<matrix_type>::max();
      int N, M;
      vector<vector<matrix_type>> mat;
       Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
          mat.assign(N, vector<matrix_type>(M, v));
       static Matrix identity(int n) { // return identity matrix.
          Matrix I(n, n);
          for (int i = 0; i < n; ++i) {
              I[i][i] = 1;
          return I;
      }
       vector<matrix_type>& operator[](int r) { return mat[r]; }
       const vector<matrix_type>& operator[](int r) const { return mat[r]; }
       Matrix& operator*=(const Matrix &other) {
           assert(M == other.N); // [N x M] [other.N x other.M]
          Matrix res(N, other.M);
          for (int r = 0; r < N; ++r) {
              for (int c = 0; c < other.M; ++c) {</pre>
                   long long square_mod = (long long) MOD * MOD;
                   long long sum = 0;
                   for (int g = 0; g < M; ++g) {
                       sum += (long long) mat[r][g] * other[g][c];
                       if (sum >= square_mod) sum -= square_mod;
32
                   res[r][c] = sum % MOD;
              }
          }
```

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