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1 Contest

1.1 C++

```
#include <bits/stdc++.h>
 using namespace std;
 #ifdef LOCAL
 #include "cp/debug.h"
 #else
 #define debug(...)
 #endif
 mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
 int main() {
     cin.tie(nullptr)->sync_with_stdio(false);
     // freopen("input.txt", "r", stdin);
     // freopen("output.txt", "w", stdout);
     return 0:
1.2 Debug
 #define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s),
     ',', ''); stringstream _ss(_s); istream_iterator<string> _it(_ss);
     out_error(_it, __VA_ARGS__);}
 void out_error(istream_iterator<string> it) { cerr << '\n'; }</pre>
 template < typename T, typename ... Args >
 void out_error(istream_iterator<string> it, T a, Args... args) {
     cerr << " [" << *it << " = " << a << "] ";
     out_error(++it, args...);
 }
 template<typename T, typename G> ostream& operator<<(ostream &os, const
     pair<T. G> &p) {
     return os << "(" << p.first << ", " << p.second << ")";</pre>
 }
 template < class Con, class = decltype(begin(declval < Con > ()))>
 typename enable_if<!is_same<Con, string>::value, ostream&>::type
 operator << (ostream& os, const Con& container) {</pre>
     os << "{";
     for (auto it = container.begin(); it != container.end(); ++it)
         os << (it == container.begin() ? "" : ", ") << *it;
     return os << "}";</pre>
1.3 Java
 import java.io.BufferedReader;
 import java.util.StringTokenizer;
 import java.io.IOException;
```

```
import java.io.InputStreamReader;
import java.io.PrintWriter;
import java.util.ArrayList;
import java.util.Arrays;
import java.util.Collections;
import java.util.Random;
public class Main {
    public static void main(String[] args) {
        FastScanner fs = new FastScanner();
        PrintWriter out = new PrintWriter(System.out);
        int n = fs.nextInt();
        out.println(n);
        out.close(); // don't forget this line.
    static class FastScanner {
        BufferedReader br:
        StringTokenizer st;
        public FastScanner() {
            br = new BufferedReader(new InputStreamReader(System.in));
            st = null;
        public String next() {
            while (st == null || st.hasMoreTokens() == false) {
                trv {
                    st = new StringTokenizer(br.readLine());
                }
                catch (IOException e) {
                    throw new RuntimeException(e);
            return st.nextToken();
        public int nextInt() {
            return Integer.parseInt(next());
        public long nextLong() {
            return Long.parseLong(next());
        public double nextDouble() {
            return Double.parseDouble(next());
}
    sublime-build
    "cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall",
    "-Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],
    "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$",
```

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2 Data structures

2.1 Sparse table

```
int st[MAXN][K + 1];
for (int i = 0; i < N; i++) {</pre>
    st[i][0] = f(array[i]);
for (int j = 1; j <= K; j++) {
    for (int i = 0; i + (1 << j) <= N; i++) {
        st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
    }
// Range Minimum Queries.
int lg[MAXN + 1];
lq[1] = 0;
for (int i = 2; i \le MAXN; i++) {
    lg[i] = lg[i / 2] + 1;
int j = lg[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
// Range Sum Queries.
long long sum = 0;
for (int j = K; j >= 0; j--) {
    if ((1 << j) <= R - L + 1) {
        sum += st[L][j];
        L += 1 << j;
    }
}
```

2.2 Ordered set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename key_type>
```

```
using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
     tree_order_statistics_node_update>;
 const int INF = 0x3f3f3f3f;
 void example() {
     vector < int > nums = \{1, 2, 3, 5, 10\};
     set_t<int> st(nums.begin(), nums.end());
     cout << *st.find_by_order(0) << '\n'; // 1</pre>
     assert(st.find_by_order(-INF) == st.end());
     assert(st.find_by_order(INF) == st.end());
     cout << st.order_of_key(2) << '\n'; // 1</pre>
     cout << st.order_of_key(4) << '\n'; // 3</pre>
     cout << st.order_of_key(9) << '\n'; // 4</pre>
     cout << st.order_of_key(-INF) << '\n'; // 0</pre>
     cout << st.order_of_key(INF) << '\n'; // 5</pre>
}
2.3
     Dsu
 struct Dsu {
     int n;
     vector<int> par, sz;
     Dsu(int _n) : n(_n) {
         sz.resize(n, 1);
         par.resize(n);
         iota(par.begin(), par.end(), 0);
     int find(int v) {
         // finding leader/parrent of set that contains the element v.
         // with {path compression optimization}.
         return (v == par[v] ? v : par[v] = find(par[v]));
     bool same(int u, int v) {
         return find(u) == find(v);
     bool unite(int u, int v) {
         u = find(u); v = find(v);
         if (u == v) return false;
         if (sz[u] < sz[v]) swap(u, v);
         par[v] = u;
         sz[u] += sz[v];
         return true;
     vector<vector<int>> groups() {
         // returns the list of the "list of the vertices in a connected
     component".
         vector<int> leader(n);
         for (int i = 0; i < n; ++i) {</pre>
             leader[i] = find(i);
         vector < int > id(n, -1);
         int count = 0;
```

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```
for (int i = 0; i < n; ++i) {
             if (id[leader[i]] == -1) {
                 id[leader[i]] = count++;
         }
         vector < int >> result(count);
         for (int i = 0; i < n; ++i) {
             result[id[leader[i]]].push_back(i);
         return result;
     }
 };
2.4
    MinQueue
 /**
  * Description: acts like normal std::queue except it supports get minimum
     value in current queue.
  */
 template <typename T>
 struct MinQueue {
     vector<T> vals;
    int ptr = 0;
     vector<int> st;
     int ptr_idx = 0;
     void push(T val) {
         while ((int) st.size() > ptr_idx && vals[st.back()] >= val) {
             st.pop_back();
         st.push_back((int) vals.size());
         vals.push_back(val);
    }
     void pop() {
         assert(ptr < (int) vals.size());</pre>
         if (ptr_idx < (int) st.size() && st[ptr_idx] == ptr) ptr_idx++;</pre>
         ptr++;
    T get() {
         assert(ptr_idx < (int) st.size());</pre>
         return vals[st[ptr_idx]];
     int front() {
         assert(!empty()); return vals[ptr];
    }
    int back() {
         assert(!empty()); return vals.back();
     bool empty() {
         return (ptr == (int) vals.size());
     int size() {
         return ((int) vals.size() - ptr);
    }
```

}; 2.5 Segment tree * Description: A segment tree with range updates and sum queries that supports three types of operations: * + Increase each value in range [l, r] by x (i.e. a[i] += x). * + Set each value in range [l, r] to x (i.e. a[i] = x). * + Determine the sum of values in range [1, r]. struct SegmentTree { vector<long long> tree, lazy_add, lazy_set; SegmentTree(int _n) : n(_n) { int p = 1;while (p < n) p *= 2;tree.resize(p * 2); lazy_add.resize(p * 2); lazy_set.resize(p * 2); long long merge(const long long &left, const long long &right) { return left + right; void build(int id, int 1, int r, const vector<int> &arr) { **if** (1 == r) { tree[id] += arr[l]; return; int mid = (1 + r) >> 1: build(id * 2, 1, mid, arr); build(id * 2 + 1, mid + 1, r, arr); tree[id] = merge(tree[id * 2], tree[id * 2 + 1]); void push(int id, int l, int r) { if (lazy_set[id] == 0 && lazy_add[id] == 0) return; int mid = (1 + r) >> 1;for (int child : {id * 2, id * 2 + 1}) { int range = (child == id * 2 ? mid - l + 1 : r - mid); if (lazy_set[id] != 0) { lazy_add[child] = 0; lazy_set[child] = lazy_set[id]; tree[child] = range * lazy_set[id]; lazy_add[child] += lazy_add[id]; tree[child] += range * lazy_add[id]; lazy_add[id] = lazy_set[id] = 0; void update(int id, int 1, int r, int u, int v, int amount, bool set_value = false) { if $(r < u \mid | 1 > v)$ return; **if** (u <= 1 && r <= v) {

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```
if (set value) {
                 tree[id] = 1LL * amount * (r - l + 1);
                 lazy_set[id] = amount;
                 lazy_add[id] = 0; // clear all previous updates.
             }
             else {
                 tree[id] += 1LL * amount * (r - l + 1);
                 lazy_add[id] += amount;
             }
             return:
         push(id, 1, r);
         int mid = (1 + r) >> 1;
         update(id * 2, 1, mid, u, v, amount, set_value);
         update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
         tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
     long long get(int id, int l, int r, int u, int v) {
         if (r < u \mid \mid 1 > v) return 0;
         if (u \le 1 \&\& r \le v) {
             return tree[id];
         }
         push(id, 1, r);
         int mid = (1 + r) >> 1;
         long long left = get(id * 2, 1, mid, u, v);
         long long right = get(id * 2 + 1, mid + 1, r, u, v);
         return merge(left, right);
 };
2.6 Efficient segment tree
 template < typename T> struct SegmentTree {
     int n;
     vector<T> tree;
     SegmentTree(int _n) : n(_n), tree(2 * n) {}
     T merge(const T &left, const T &right) {
         return left + right;
     template<typename G>
     void build(const vector<G> &initial) {
         assert((int) initial.size() == n);
         for (int i = 0; i < n; ++i) {
             tree[i + n] = initial[i];
         for (int i = n - 1; i > 0; --i) {
             tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
     void modify(int i, int v) {
         tree[i += n] = v;
         for (i /= 2; i > 0; i /= 2) {
             tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
         }
```

```
T get_sum(int 1, int r) {
        // sum of elements from 1 to r - 1.
        T ret{};
        for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
            if (1 & 1) ret = merge(ret, tree[1++]);
            if (r & 1) ret = merge(ret, tree[--r]);
        }
        return ret;
    }
};
   Persistent lazy segment tree
struct Vertex {
    int 1. r:
    long long val, lazy;
    bool has_changed = false;
    Vertex() {}
    Vertex(int _1, int _r, long long _val, int _lazy = 0) : l(_1), r(_r),
    val(_val), lazy(_lazy) {}
};
struct PerSegmentTree {
    vector<Vertex> tree;
    vector<int> root:
    int build(const vector<int> &arr, int 1, int r) {
        if (1 == r) {
            tree.emplace_back(-1, -1, arr[l]);
            return tree.size() - 1;
        int mid = (1 + r) / 2;
        int left = build(arr, 1, mid);
        int right = build(arr, mid + 1, r);
        tree.emplace_back(left, right, tree[left].val + tree[right].val);
        return tree.size() - 1;
    int add(int x, int 1, int r, int u, int v, int amt) {
        if (1 > v \mid | r < u) return x;
        if (u <= 1 && r <= v) {
            tree.emplace_back(tree[x].1, tree[x].r, tree[x].val + 1LL * amt *
    (r - l + 1), tree[x].lazy + amt);
            tree.back().has_changed = true;
            return tree.size() - 1;
        int mid = (1 + r) >> 1;
        push(x, 1, mid, r);
        int left = add(tree[x].1, 1, mid, u, v, amt);
        int right = add(tree[x].r, mid + 1, r, u, v, amt);
        tree.emplace_back(left, right, tree[left].val + tree[right].val, 0);
        return tree.size() - 1;
    long long get_sum(int x, int l, int r, int u, int v) {
        if (r < u \mid | 1 > v) return 0;
        if (u <= 1 && r <= v) return tree[x].val;</pre>
```

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if (1 == r) {

```
int mid = (1 + r) / 2:
         push(x, 1, mid, r);
         return get_sum(tree[x].1, 1, mid, u, v) + get_sum(tree[x].r, mid + 1,
     r, u, v);
     void push(int x, int 1, int mid, int r) {
         if (!tree[x].has_changed) return;
         Vertex left = tree[tree[x].1];
         Vertex right = tree[tree[x].r];
         tree.emplace_back(left);
         tree[x].l = tree.size() - 1;
         tree.emplace_back(right);
         tree[x].r = tree.size() - 1;
         tree[tree[x].l].val += tree[x].lazy * (mid - l + 1);
         tree[tree[x].1].lazy += tree[x].lazy;
         tree[tree[x].r].val += tree[x].lazy * (r - mid);
         tree[tree[x].r].lazy += tree[x].lazy;
         tree[tree[x].1].has_changed = true;
         tree[tree[x].r].has_changed = true;
         tree[x].lazy = 0;
         tree[x].has_changed = false;
};
2.8
   Lichao tree
  * Description: A segment tree that allows insert a new line and query for
     minimum value over all lines at point x.
  * Usage: useful in convex hull trick.
 const long long INF_LL = (long long) 4e18;
 struct Line {
     long long a, b;
    Line(long long _a = 0, long long _b = INF_LL): a(_a), b(_b) {}
    long long operator()(long long x) {
         return a * x + b;
    }
 };
 struct SegmentTree { // min query
     int n;
     vector<Line> tree;
     SeamentTree() {}
     SegmentTree(int _n): n(1) {
         while (n < _n) n *= 2;
         tree.resize(n * 2);
     void insert(int x, int l, int r, Line line) {
```

```
if (line(l) < tree[x](l)) tree[x] = line;</pre>
        int mid = (1 + r) >> 1;
        bool b_left = line(l) < tree[x](l);</pre>
        bool b_mid = line(mid) < tree[x](mid);</pre>
        if (b_mid) swap(tree[x], line);
        if (b_left != b_mid) insert(x * 2, 1, mid, line);
        else insert(x * 2 + 1, mid + 1, r, line);
    long long query(int x, int 1, int r, int at) {
        if (l == r) return tree[x](at);
        int mid = (1 + r) >> 1;
        if (at <= mid) return min(tree[x](at), query(x * 2, 1, mid, at));</pre>
        else return min(tree[x](at), query(x * 2 + 1, mid + 1, r, at));
};
    Old driver tree (Chtholly tree)
 * Description: An optimized brute-force approach to deal with problem that
    has operation of setting an interval to the same number.
 * Note: caution TLE, only works when input is random
struct ODT {
    map<int, long long> tree;
    using It = map<int, long long>::iterator;
    It split(int x) {
        It it = tree.upper_bound(x);
        assert(it != tree.begin());
        --it;
        if (it->first == x) return it;
        return tree.emplace(x, it->second).first;
    }
    void add(int 1, int r, int amt) {
        It it_l = split(l);
        It it_r = split(r + 1);
        while (it_l != it_r) {
            it_l->second += amt;
            ++it_l;
    }
    void set(int 1, int r, int v) {
        It it_l = split(l);
        It it_r = split(r + 1);
        while (it_l != it_r) {
            tree.erase(it_l++);
        tree[1] = v;
```

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```
}
     long long kth_smallest(int 1, int r, int k) {
         // return the k-th smallest value in range [l..r]
         vector<pair<long long, int>> values; // pair(value, count)
         It it_l = split(l);
         It it_r = split(r + 1);
         while (it_l != it_r) {
             It prev = it_l++;
             values.emplace_back(prev->second, it_l->first - prev->first);
         sort(values.begin(), values.end());
         for (auto [value, cnt] : values) {
             if (k <= cnt) return value;</pre>
             k -= cnt;
         }
         return -1;
    int powmod(long long a, long long n, int mod);
     int sum_of_xth_power(int 1, int r, int x, int mod) {
         It it_l = split(l);
         It it_r = split(r + 1);
         int res = 0;
         while (it_l != it_r) {
             It prev = it_l++;
             res = (res + 1LL * (it_l->first - prev->first) *
     powmod(prev->second, x, mod)) % mod;
         }
         return res;
};
2.10
      Disjoint sparse table
  * Description: range query on a static array.
  * Time: O(1) per query.
  * Tested: stress-test.
 const int MOD = (int) 1e9 + 7;
 struct DisjointSparseTable { // product queries.
     int n, h;
     vector<vector<int>> dst:
     vector<int> lg;
     DisjointSparseTable(int _n) : n(_n) {
         h = 1; // in case n = 1: h = 0 !!.
         int p = 1;
         while (p < n) p *= 2, h++;
         lg.resize(p); lg[1] = 0;
         for (int i = 2; i < p; ++i) {</pre>
             lg[i] = 1 + lg[i / 2];
         dst.resize(h, vector<int>(n));
```

```
void build(const vector<int> &A) {
         for (int lv = 0; lv < h; ++lv) {</pre>
             int len = (1 << lv);</pre>
             for (int k = 0; k < n; k += len * 2) {
                 int mid = min(k + len, n);
                 dst[lv][mid - 1] = A[mid - 1] \% MOD;
                 for (int i = mid - 2; i >= k; --i) {
                      dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
                 if (mid == n) break;
                 dst[lv][mid] = A[mid] % MOD;
                 for (int i = mid + 1; i < min(mid + len, n); ++i) {</pre>
                      dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
             }
         }
     int get(int 1, int r) {
         if (1 == r) {
             return dst[0][1];
         int i = lg[l ^ r];
         return 1LL * dst[i][l] * dst[i][r] % MOD;
};
2.11 Fenwick tree
  * Description: range update and range sum query.
 using tree_type = long long;
 struct FenwickTree {
     int n;
     vector<tree_type> fenw_coeff, fenw;
     FenwickTree() {}
     FenwickTree(int _n) : n(_n) {
         fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
         fenw.assign(n, 0); // normal fenwick tree.
     template < typename G>
     void build(const vector<G> &A) {
         assert((int) A.size() == n);
         vector<int> diff(n);
         diff[0] = A[0];
         for (int i = 1; i < n; ++i) {
             diff[i] = A[i] - A[i - 1];
         fenw_coeff[0] = (long long) diff[0] * n;
         fenw[0] = diff[0];
         for (int i = 1; i < n; ++i) {
             fenw_coeff[i] = fenw_coeff[i - 1] + (long long) diff[i] * (n - i);
             fenw[i] = fenw[i - 1] + diff[i];
```

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```
for (int i = n - 1; i >= 0; --i) {
             int j = (i \& (i + 1)) - 1;
             if (j >= 0) {
                 fenw_coeff[i] -= fenw_coeff[j];
                 fenw[i] -= fenw[j];
         }
     }
     void add(vector<tree_type> &fenw, int i, tree_type val) {
         while (i < n) {
             fenw[i] += val;
             i = (i + 1);
         }
     }
     tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
         tree_type res{};
         while (i >= 0) {
             res += fenw[i];
             i = (i \& (i + 1)) - 1:
         }
         return res;
     tree_type prefix_sum(int i) {
         return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i -
     1);
     void range_add(int 1, int r, tree_type val) {
         add(fenw_coeff, 1, (n - 1) * val);
         add(fenw_coeff, r + 1, (n - r - 1) * (-val));
         add(fenw, 1, val);
         add(fenw, r + 1, -val);
     }
     tree_type range_sum(int 1, int r) {
         return prefix_sum(r) - prefix_sum(l - 1);
     }
 };
2.12 Fenwick tree 2D
 /**
  * Description: range update and range sum query on a 2D array.
 using tree_type = long long;
 struct FenwickTree2D {
     int n, m;
     vector<vector<tree_type> > fenw[4];
     FenwickTree2D(int _n, int _m) : n(_n), m(_m) {
         for (int i = 0; i < 4; i++) {
             fenw[i].resize(n, vector<tree_type>(m));
         }
     }
     void add(int u, int v, tree_type val) {
         for (int i = u; i < n; i |= (i + 1)) {
```

```
for (int j = v; j < m; j |= (j + 1)) {
             fenw[0][i][j] += val;
             fenw[1][i][j] += (u + 1) * val;
             fenw[2][i][j] += (v + 1) * val;
             fenw[3][i][j] += (u + 1) * (v + 1) * val;
         }
     }
     void range_add(int r, int c, int rr, int cc, tree_type val) { // [r, rr]
     x [c, cc].
         add(r, c, val);
         add(r, cc + 1, -val);
         add(rr + 1, c, -val);
         add(rr + 1, cc + 1, val);
     tree_type prefix_sum(int u, int v) {
         tree_type res{};
         for (int i = u; i >= 0; i = (i & (i + 1)) - 1) {
           for (int j = v; j >= 0; j = (j & (j + 1)) - 1) {
             res += (u + 2) * (v + 2) * fenw[0][i][j];
             res -= (v + 2) * fenw[1][i][j];
             res -= (u + 2) * fenw[2][i][j];
             res += fenw[3][i][j];
         }
         return res;
     tree_type range_sum(int r, int c, int rr, int cc) { // [r, rr] x [c, cc].
         return prefix_sum(rr, cc) - prefix_sum(r - 1, cc) - prefix_sum(rr, c
     -1) + prefix_sum(r - 1, c - 1);
 };
2.13 Implicit treap
 struct Node {
     int val, prior, cnt;
     bool rev;
     Node *left, *right;
     Node() {}
     Node(int _val) : val(_val), prior(rng()), cnt(1), rev(false),
     left(nullptr), right(nullptr) {}
 };
 // Binary search tree + min-heap.
 struct Treap {
     Node *root;
     Treap() : root(nullptr) {}
     int get_cnt(Node *n) { return n ? n->cnt : 0; }
     void upd_cnt(Node *&n) {
         if (n) n->cnt = get_cnt(n->left) + get_cnt(n->right) + 1;
     void push_rev(Node *treap) {
         if (!treap || !treap->rev) return;
         treap -> rev = false;
```

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```
swap(treap->left, treap->right);
   if (treap->left) treap->left->rev ^= true;
   if (treap->right) treap->right->rev ^= true;
pair<Node*, Node*> split(Node *treap, int x, int smaller = 0) {
    if (!treap) return {};
    push_rev(treap);
    int idx = smaller + get_cnt(treap->left); // implicit val.
   if (idx <= x) {
        auto pr = split(treap->right, x, idx + 1);
        treap->right = pr.first;
        upd_cnt(treap);
        return {treap, pr.second};
   }
    else {
        auto pl = split(treap->left, x, smaller);
        treap->left = pl.second;
        upd_cnt(treap);
        return {pl.first, treap};
}
Node* merge(Node *1, Node *r) {
    push_rev(l); push_rev(r);
   if (!1 || !r) return (1 ? 1 : r);
   if (l->prior < r->prior) {
        1->right = merge(1->right, r);
        upd_cnt(1);
        return 1;
   }
    else {
        r->left = merge(1, r->left);
        upd_cnt(r);
        return r;
   }
}
void insert(int pos, int val) {
   if (!root) {
        root = new Node(val);
        return:
    Node *1, *m, *r;
   m = new Node(val);
   tie(l, r) = split(root, pos - 1);
   root = merge(l, merge(m, r));
void erase(int pos_l, int pos_r) {
   Node *1, *m, *r;
    tie(l, r) = split(root, pos_l - 1);
    tie(m, r) = split(r, pos_r - pos_l);
    root = merge(1, r);
void reverse(int pos_l, int pos_r) {
    Node *1, *m, *r;
```

```
tie(l, r) = split(root, pos_l - 1);
         tie(m, r) = split(r, pos_r - pos_l);
         m->rev ^= true;
         root = merge(1, merge(m, r));
     int query(int pos_l, int pos_r);
          // returns answer for corresponding types of guery.
     void inorder(Node *n) {
         if (!n) return;
         push_rev(n);
         inorder(n->left);
          cout << n->val << ' ';
         inorder(n->right);
     }
     void print() {
         inorder(root);
         cout << '\n';</pre>
};
2.14 Line container
  * Source: kactl
    Description: container that allow you can add lines in form 'ax + b' and
     query maximum value at 'x'.
 using num_t = int;
 struct Line {
     num_t a, b; // ax + b
     mutable num_t x; // x-intersect with the next line in the hull
     bool operator<(const Line &other) const {</pre>
         return a < other.a;</pre>
     bool operator < (num_t other_x) const {</pre>
         return x < other_x;</pre>
     }
};
 struct LineContainer : multiset<Line, less<>>> { // max-query
     // for doubles, use INF = 1 / 0.0
     static const num_t INF = numeric_limits<num_t>::max();
     num_t floor_div(num_t a, num_t b) {
         return a / b - ((a ^ b) < 0 && a % b != 0);
     bool isect(iterator u, iterator v) {
         if (v == end()) {
              u->x = INF:
              return false;
         if (u->a == v->a) u->x = (u->b > v->b ? INF : -INF);
         else u \rightarrow x = floor_div(v \rightarrow b - u \rightarrow b, u \rightarrow a - v \rightarrow a);
         return u->x >= v->x:
```

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```
void add(num_t a, num_t b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) {
        y = erase(y);
        isect(x, y);
    }
    while ((y = x) != begin() && (--x)->x >= y->x) {
        isect(x, erase(y));
    }
}
num_t query(num_t x) {
    assert(!empty());
    auto it = *lower_bound(x);
    return it.a * x + it.b;
}
```

3 Mathematics

3.1 Trigonometry

};

3.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u)\cos(v) \pm \cos(u)\sin(v)$$

$$\cos(u \pm v) = \cos(u)\cos(v) \mp \sin(u)\sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u)\tan(v)}$$

3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2\cos(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\cos(u) - \cos(v) = -2\sin(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

$$\sin(u) + \sin(v) = 2\sin(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\sin(u) - \sin(v) = 2\cos(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

3.1.3 Product identities

$$\cos(u)\cos(v) = \frac{1}{2}[\cos(u+v) + \cos(u-v)]$$

$$\sin(u)\sin(v) = -\frac{1}{2}[\cos(u+v) - \cos(u-v)]$$

$$\sin(u)\cos(v) = \frac{1}{2}[\sin(u+v) + \sin(u-v)]$$

3.1.4 Double - triple angle identities

$$\sin(2u) = 2\sin(u)\cos(u)$$

$$\cos(2u) = 2\cos^{2}(u) - 1 = 1 - 2\sin^{2}(u)$$

$$\tan(2u) = \frac{2\tan(u)}{1 - \tan^{2}(u)}$$

$$\sin(3u) = 3\sin(u) - 4\sin^{3}(u)$$

$$\cos(3u) = 4\cos^{3}(u) - 3\cos(u)$$

$$\tan(3u) = \frac{3\tan(u) - \tan^{3}(u)}{1 - 3\tan^{2}(u)}$$

3.2 Sums

$$c^{a} + c^{a+1} + \dots + c^{b} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$c + 2c^{2} + \dots + nc^{n} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c - 1)^{2}}, c \neq 1$$

$$1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$$

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(n+1)(2n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$1^{5} + 2^{5} + 3^{5} + \dots + n^{5} = \frac{n^{2}(n+1)^{2}(2n^{2} + 2n - 1)}{12}$$

$$1^{6} + 2^{6} + 3^{6} + \dots + n^{6} = \frac{n(n+1)(2n+1)(3n^{4} + 6n^{3} - 3n + 1)}{42}$$

$$1^{7} + 2^{7} + 3^{7} + \dots + n^{7} = \frac{n^{2}(n+1)^{2}(3n^{4} + 6n^{3} - n^{2} - 4n + 2)}{24}$$

4 String

4.1 Prefix function

```
/**
 * Description: The prefix function of a string 's' is defined as an array pi
   of length n,
 * where pi[i] is the length of the longest proper prefix of the substring
 * s[0..i] which is also a suffix of this substring.
 * Time complexity: O(|S|).
 */
```

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```
vector<int> prefix_function(const string &s) {
     int n = (int) s.length();
     vector<int> pi(n);
     pi[0] = 0;
     for (int i = 1; i < n; ++i) {
         int j = pi[i - 1]; // try length pi[i - 1] + 1.
         while (j > 0 \&\& s[j] != s[i]) {
             j = pi[j - 1];
         if (s[j] == s[i]) {
             pi[i] = j + 1;
    }
     return pi;
4.2 Z function
 /**
  * Description: for a given string 's', z[i] = longest common prefix of 's'
     and suffix starting at 'i'.
     z[0] is generally not well defined (this implementation below assume
     z[0] = 0).
  */
 vector<int> z_function(const string &s) {
     int n = (int) s.size();
     vector<int> z(n);
    z[0] = 0;
     // [1, r)
     for (int i = 1, l = 0, r = 0; i < n; ++i) {
         if (i < r) z[i] = min(r - i, z[i - 1]);
         while (i + z[i] < n \& s[z[i]] == s[i + z[i]]) ++z[i];
         if (i + z[i] > r) {
             1 = i;
             r = i + z[i];
        }
    }
     return z;
     Counting occurrences of each prefix
 #include "prefix_function.h"
 vector<int> count_occurrences(const string &s) {
     vector<int> pi = prefix_function(s);
     int n = (int) s.size();
     vector<int> ans(n + 1);
     for (int i = 0; i < n; ++i) {
         ans[pi[i]]++;
    }
     for (int i = n - 1; i > 0; --i) {
         ans[pi[i - 1]] += ans[i];
     for (int i = 0; i <= n; ++i) {</pre>
```

```
ans[i]++;
    return ans;
   // Input: ABACABA
    // Output: 4 2 2 1 1 1 1
   Knuth-Morris-Pratt algorithm
* Searching for a substring in a string.
 * Time complexity: O(N + M).
#include "prefix_function.h"
vector<int> KMP(const string &text, const string &pattern) {
    int n = (int) text.length();
    int m = (int) pattern.length();
    string s = pattern + '$' + text;
    vector<int> pi = prefix_function(s);
    vector<int> indices;
    for (int i = 0; i < (int) s.length(); ++i) {</pre>
       if (pi[i] == m) {
            indices.push_back(i - 2 * m);
    return indices;
    Suffix array
* Description: suffix array is a sorted array of all the suffixes of a given
    string.
    sa[i] = starting index of the i-th smallest suffix.
    rank[i] = rank of the suffix starting at 'i'.
    lcp[i] = longest common prefix between 'sa[i - 1]' and 'sa[i]'
    for arbitrary 'u v', let i = rank[u] - 1, j = rank[v] - 1 (assume i <
       longest\_common\_prefix(u, v) = min(lcp[i + 1], lcp[i + 2], ..., lcp[j])
 * Time: O(NlogN).
struct SuffixArray {
    string s;
   int n, lim;
    vector<int> sa, lcp, rank;
    SuffixArray(const string &_s, int _{lim} = 256) : s(_s), n(s.length() + 1),
        \lim(_{\lim}, sa(n), lcp(n), rank(n) {
       s += '$';
       build(); kasai();
        sa.erase(sa.begin()); lcp.erase(lcp.begin());
        rank.pop_back(); s.pop_back();
```

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```
vector<int> nrank(n), norder(n), cnt(max(n, lim));
        for (int i = 0; i < n; ++i) {
            sa[i] = i; rank[i] = s[i];
        for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = max(1, k * 2),
    lim = rank_cnt + 1) {
            for (int i = 0; i < n; ++i) {</pre>
                norder[i] = (sa[i] - k + n) \% n;
                cnt[rank[i]]++;
            for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];</pre>
            for (int i = n - 1; i >= 0; --i) sa[--cnt[rank[norder[i]]]] =
    norder[i]:
            rank[sa[0]] = rank\_cnt = 0;
            for (int i = 1; i < n; ++i) {
                int u = sa[i], v = sa[i - 1];
                int nu = (u + k) \% n, nv = (v + k) \% n;
                if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
                nrank[sa[i]] = rank cnt:
            for (int i = 0; i < rank_{cnt} + 1; ++i) cnt[i] = 0;
            rank.swap(nrank);
        }
    }
    void kasai() {
        for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
            int j = sa[rank[i] - 1];
            while (s[i + k] == s[j + k]) k++;
            lcp[rank[i]] = k;
    }
};
    Suffix array slow
/**
* Description: an easier way to implement suffix array but run slower
 * Time: O(N * logN^2)
struct SuffixArraySlow {
    string s;
    int n:
    vector<int> sa, lcp, rank;
    SuffixArraySlow(const string &_s): s(_s), n((int) s.size() + 1), sa(n),
    lcp(n), rank(n) {
        s += '$';
        build(); kasai();
        sa.erase(sa.begin()); lcp.erase(lcp.begin());
        rank.pop_back(); s.pop_back();
    bool comp(int i, int j, int k) {
        return make_pair(rank[i], rank[(i + k) % n]) < make_pair(rank[j],</pre>
    rank[(j + k) % n]);
```

void build() {

```
void build() {
        vector<int> nrank(n);
        for (int i = 0; i < n; ++i) {
            sa[i] = i; rank[i] = s[i];
        for (int k = 0; k < n; k = max(1, k * 2)) {
            stable_sort(sa.begin(), sa.end(), [&](int i, int j) {
                return comp(i, j, k);
            });
            for (int i = 0, cnt = 0; i < n; ++i) {
                if (i > 0 \&\& comp(sa[i - 1], sa[i], k)) ++cnt;
                nrank[sa[i]] = cnt;
            }
            rank.swap(nrank);
       }
    void kasai() {
        for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
            int j = sa[rank[i] - 1];
            while (s[i + k] == s[j + k]) ++k;
            lcp[rank[i]] = k;
    }
};
    Manacher's algorithm
 * Description: for each position, computes d[0][i] = half length of
longest palindrome centered on i (rounded up), d[1][i] = half length of
 longest palindrome centered on i and i - 1.
 * Time complexity: O(N).
 * Tested: https://judge.yosupo.jp/problem/enumerate_palindromes,
    stress-tested.
array<vector<int>, 2> manacher(const string &s) {
    int n = (int) s.size();
    array<vector<int>, 2> d;
    for (int z = 0; z < 2; ++z) {
        d[z].resize(n);
        int 1 = 0, r = 0;
        for (int i = 0; i < n; ++i) {
            int mirror = l + r - i + z;
            d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
            int L = i - d[z][i] - z, R = i + d[z][i];
            while (L >= 0 \&\& R < n \&\& s[L] == s[R]) {
                d[z][i]++; L--; R++;
            if (R > r) {
                1 = L; r = R;
       }
```

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```
return d:
4.8 Trie
 struct Trie {
     const static int ALPHABET = 26;
     const static char minChar = 'a';
     struct Vertex {
         int next[ALPHABET];
         bool leaf;
         Vertex() {
             leaf = false;
             fill(next, next + ALPHABET, -1);
         }
     };
     vector<Vertex> trie:
     Trie() { trie.emplace_back(); }
     void insert(const string &s) {
         int i = 0;
         for (const char &ch : s) {
             int j = ch - minChar;
             if (trie[i].next[j] == -1) {
                 trie[i].next[j] = trie.size();
                 trie.emplace_back();
             i = trie[i].next[j];
         trie[i].leaf = true;
     bool find(const string &s) {
         int i = 0:
         for (const char &ch : s) {
             int j = ch - minChar;
             if (trie[i].next[j] == -1) {
                 return false;
             i = trie[i].next[j];
         return (trie[i].leaf ? true : false);
 };
4.9 Hashing
 struct Hash61 {
     static const uint64_t MOD = (1LL << 61) - 1;</pre>
     static uint64_t BASE;
     static vector<uint64_t> pw;
     uint64_t addmod(uint64_t a, uint64_t b) const {
         a += b;
         if (a >= MOD) a -= MOD;
         return a:
     }
```

```
uint64_t submod(uint64_t a, uint64_t b) const {
         a += MOD - b;
         if (a >= MOD) a -= MOD;
         return a;
     uint64_t mulmod(uint64_t a, uint64_t b) const {
         uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
         uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
         uint64_t low = low1 * low2;
         uint64_t mid = low1 * high2 + low2 * high1;
         uint64_t high = high1 * high2;
         uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >> 29)
     + (mid << 35 >> 3) + 1;
         // ret %= MOD:
         ret = (ret >> 61) + (ret & MOD);
         ret = (ret >> 61) + (ret & MOD);
         return ret - 1;
     void ensure_pw(int m) {
         int sz = (int) pw.size();
         if (sz >= m) return;
         pw.resize(m);
         for (int i = sz; i < m; ++i) {</pre>
             pw[i] = mulmod(pw[i - 1], BASE);
         }
    }
     vector<uint64_t> pref;
     int n;
     template < typename T > Hash61(const T &s) { // strings or arrays.
         n = (int) s.size();
         ensure_pw(n);
         pref.resize(n + 1);
         pref[0] = 0:
         for (int i = 0; i < n; ++i) {
             pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
     inline uint64_t operator()(const int from, const int to) const {
         assert(0 \le from \&\& from \le to \&\& to < n);
         // pref[to + 1] - pref[from] * pw[to - from + 1]
         return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
    }
 };
 mt19937 rnd((unsigned int)
     chrono::steady_clock::now().time_since_epoch().count());
 uint64_t Hash61::BASE = (MOD >> 2) + rnd() % (MOD >> 1);
 vector<uint64_t> Hash61::pw = vector<uint64_t>(1, 1);
4.10
     Minimum rotation
 /**
```

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```
* Author: Stjepan Glavina
 * License: Unlicense
 * Source: https://github.com/stjepang/snippets/blob/master/min_rotation.cpp
 * Description: Finds the lexicographically smallest rotation of a string.
 * Time: O(N)
 * Usage:
 * rotate(v.begin(), v.begin()+minRotation(v), v.end());
* Status: Stress-tested
*/
#pragma once
int minRotation(string s) {
 int a = 0, N = (int) s.size(); s += s;
 rep(b, 0, N) rep(k, 0, N) {
   if (a + k == b \mid | s[a + k] < s[b + k]) \{b += max(0, k - 1); break;\}
   if (s[a + k] > s[b + k]) \{ a = b; break; \}
 }
 return a;
```

5 Numerical

5.1 Fast Fourier transform

```
const double PI = acos(-1);
using Comp = complex < double >;
int reverse_bit(int n, int lg) {
    int res = 0;
   for (int i = 0; i < lg; ++i) {
        if (n & (1 << i)) {
            res = (1 << (lg - i - 1));
       }
   }
    return res;
void fft(vector<Comp> &a, bool invert = false) {
   int n = (int) a.size();
   int lg = 0;
    while (1 << (lg) < n) ++lg;</pre>
    for (int i = 0; i < n; ++i) {
        int rev_i = reverse_bit(i, lg);
        if (i < rev_i) swap(a[i], a[rev_i]);</pre>
    for (int len = 2; len <= n; len *= 2) {</pre>
        double angle = 2 * PI / len * (invert ? -1 : 1);
        Comp w_base(cos(angle), sin(angle));
        for (int i = 0; i < n; i += len) {</pre>
            Comp w(1);
            for (int j = i; j < i + len / 2; ++j) {
                Comp u = a[j], v = a[j + len / 2];
                a[j] = u + w * v;
                a[j + len / 2] = u - w * v;
                w *= w_base;
```

```
if (invert) for (int i = 0; i < n; ++i) a[i] /= n;</pre>
vector<int> mult(vector<int> &a, vector<int> &b) {
    vector<Comp> A(a.begin(), a.end()), B(b.begin(), b.end());
    int n = (int) a.size(), m = (int) b.size(), p = 1;
    while (p < n + m) p *= 2;
    A.resize(p), B.resize(p);
    fft(A, false);
    fft(B, false);
    for (int i = 0; i < p; ++i) {
        A[i] *= B[i];
    }
    fft(A, true);
    vector < int > res(n + m - 1);
    for (int i = 0; i < n + m - 1; ++i) {
        res[i] = (int) round(A[i].real());
    return res:
}
```

6 Number Theory

6.1 Euler's totient function

- Euler's totient function, also known as **phi-function** $\phi(n)$ counts the number of integers between 1 and n inclusive, that are **coprime to** n.
- Properties:
 - Divisor sum property: $\sum_{d|n} \phi(d) = n$.
 - $\phi(n)$ is a **prime number** when n = 3, 4, 6.
 - If *p* is a prime number, then $\phi(p) = p 1$.
 - If *p* is a prime number and *k* ≥ 1, then $\phi(p^k) = p^k p^{k-1}$.
 - If *a* and *b* are **coprime**, then $\phi(ab) = \phi(a) \cdot \phi(b)$.
 - In general, for **not coprime** *a* and *b*, with d = gcd(a, b) this equation holds: $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$.
 - With $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\phi(n) = \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m})$$
$$= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)$$

• Application in Euler's theorem:

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- If gcd(a, M) = 1, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod{\phi(M)}} \pmod{M}$$

- In general, for arbitrary a, M and n ≥ $log_2 M$:

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

Mobius function

• For a positive integer $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$:

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1\\ 0, & \text{if } \exists k_i > 1\\ (-1)^m & \text{otherwise} \end{cases}$$

- Properties:
 - $-\sum_{d|n} \mu(d) = [n=1].$
 - If *a* and *b* are **coprime**, then $\mu(ab) = \mu(a) \cdot \mu(b)$.
 - Mobius inversion: let *f* and *g* be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right)g(d)$$

Primes 6.3

Approximating the number of primes up to *n*:

n	$\pi(n)$	$\frac{n}{\ln n - 1}$
$100 (1e^2)$	25	28
$500 (5e^2)$	95	96
$1000 (1e^3)$	168	169
$5000 (5e^3)$	669	665
$10000 (1e^4)$	1229	1218
$50000 (5e^4)$	5133	5092
$100000 (1e^5)$	9592	9512
$500000 (5e^5)$	41538	41246
$1000000 (1e^6)$	78498	78030
$5000000 (5e^6)$	348513	346622

 $(\pi(n))$ = the number of primes less than or equal to n, $\frac{n}{\ln n - 1}$ is used to approximate $\pi(n)$).

6.4 Wilson's theorem

A positive integer *n* is a prime if and only if:

$$(n-1)! \equiv n-1 \pmod{n}$$

6.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example: 64 = 55 + 8 + 1

```
85 = 55 + 21 + 8 + 1
vector<int> zeckendoft_theorem(int n) {
    vector < int > fibs = \{1, 1\};
    int sz = 2;
    while (fibs.back() <= n) {</pre>
         fibs.push_back(fibs[sz - 1] + fibs[sz - 2]);
    fibs.pop_back();
    vector<int> nums;
    int p = sz - 1;
    while (n > 0) {
        if (n >= fibs[p]) {
             nums.push_back(fibs[p]);
             n -= fibs[p];
        p--;
    }
    return nums;
}
```

Bitwise operation

- $a + b = (a \oplus b) + 2(a \& b)$ $a \mid b = (a \oplus b) + (a \& b)$ $a \& (b \oplus c) = (a \& b) \oplus (a \& c)$ $a \mid (b \& c) = (a \mid b) \& (a \mid c)$ $a \mid (b \& c) = (a \mid b) \& (a \mid c)$ $a \mid (a \otimes v) u$ $a \& (a \mid b) = a$ $a \& (a \mid b) = a$ $a = 2^k \Leftrightarrow !(n \& (n-1)) = 1$ $-a = \sim a + 1$ $4i \oplus (4i + 1) \oplus (4i + 2) \oplus (4i + 2) \oplus (4i + 2)$

- $4i \oplus (4i + 1) \oplus (4i + 2) \oplus (4i + 3) = 0$
- Iterating over all subsets of a set and iterating over all submasks of a mask:

```
int n:
void mask_example() {
    for (int mask = 0; mask < (1 << n); ++mask) {
        for (int i = 0; i < n; ++i) {
            if (mask & (1 << i)) {
                // do something...
        }
```

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```
// Time complexity: 0(n * 2^n).
     for (int mask = 0; mask < (1 << n); ++mask) {
         for (int submask = mask; ; submask = (submask - 1) & mask) {
             // do something...
             if (submask == 0) break;
         // Time complexity: 0(3^n).
    }
 }
6.7 Pollard's rho algorithm
 const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.</pre>
 const int LIMIT = (int) 1e9;
 vector<int> primes;
 int small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 73, 113,
     193, 311, 313, 407521, 299210837};
 void linear sieve(int n):
 num_t mulmod(num_t a, num_t b, num_t mod);
 num_t powmod(num_t a, num_t n, num_t mod);
 bool miller_rabin(num_t a, num_t d, int s, num_t mod) {
     num_t x = powmod(a, d, mod);
    if (x == mod - 1 || x == 1) 
         return true;
    }
    for (int i = 0; i < s - 1; ++i) {
         x = mulmod(x, x, mod);
         if (x == mod - 1) return true;
    }
     return false;
 bool is_prime(num_t n, int tests = 10) {
    if (n < 4) return (n > 1);
     num_t d = n - 1;
     int s = 0;
     while (d % 2 == 0) { d >>= 1; s++; }
     for (int i = 0; i < tests; ++i) {</pre>
         int a = small_primes[i];
        if (n == a) return true;
         if (n % a == 0 || !miller_rabin(a, d, s, n)) return false;
    }
     return true;
 num_t f(num_t x, int c, num_t mod) { // f(x) = (x^2 + c) % mod.}
    x = mulmod(x, x, mod);
     x += c;
     if (x >= mod) x -= mod;
     return x;
 num_t pollard_rho(num_t n, int c) {
     // algorithm to find a random divisor of 'n'.
    // using random function: f(x) = (x^2 + c) \% n.
```

```
num_t x = 2, y = x, d;
    long long p = 1;
    int dist = 0;
    while (true) {
        y = f(y, c, n);
        dist++:
        d = \_\_gcd(llabs(x - y), n);
        if (d > 1) break;
        if (dist == p) { dist = 0; p *= 2; x = y; }
    return d;
void factorize(int n, vector<num_t> &factors);
void llfactorize(num_t n, vector<num_t> &factors) {
    if (n < 2) return;
    if (is_prime(n)) {
        factors.emplace_back(n);
        return:
    if (n < LIMIT) {</pre>
        factorize(n, factors);
        return;
    }
    num_t d = n;
    for (int c = 2; d == n; c++) {
        d = pollard_rho(n, c);
    llfactorize(d, factors);
   llfactorize(n / d, factors);
vector<num_t> gen_divisors(vector<pair<num_t, int>> &factors) {
    vector<num_t> divisors = {1};
    for (auto &x : factors) {
        int sz = (int) divisors.size();
        for (int i = 0; i < sz; ++i) {
            num_t cur = divisors[i];
            for (int j = 0; j < x.second; ++j) {
                cur *= x.first:
                divisors.push_back(cur);
       }
    return divisors; // this array is NOT sorted yet.
    Segment divisor sieve
const int MAXN = (int) 1e6; // R - L + 1 <= N.
int divisor_count[MAXN + 3];
void segment_divisor_sieve(long long L, long long R) {
    for (long long i = 1; i \le (long long) sqrt(R); ++i) {
        long long start1 = ((L + i - 1) / i) * i;
        long long start2 = i * i;
        long long j = max(start1, start2);
```

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```
if (j == start2) {
            divisor_count[j - L] += 1;
           j += i:
       for ( ; j <= R; j += i) {</pre>
            divisor_count[j - L] += 2;
   }
   Linear sieve
* Description: Finding primes and computing value for multiplicative
    function in O(N)
const int N = (int) 1e6;
bool is_prime[N + 1];
int spf[N + 1]; // smallest prime factor
int phi[N + 1]; // euler's totient function
int mu[N + 1]; // mobius function
int func[N + 1]; // a multiplicative function, f(p^k) = k
int cnt[N + 1]; // cnt[i] = the power of the smallest prime factor of i
int pw[N + 1]; // pw[i] = p^cnt[i] where p is the smallest prime factor of i
vector<int> primes;
void sieve(int n = N) {
   spf[0] = spf[1] = -1;
   phi[1] = mu[1] = func[1] = 1;
   for (int x = 2; x <= n; ++x) {
       if (spf[x] == 0) {
           primes.push_back(spf[x] = x);
           is_prime[x] = true;
            phi[x] = x - 1;
           mu[x] = -1;
            func[x] = 1;
           cnt[x] = 1;
           pw[x] = x;
       for (int p : primes) {
           if (p > spf[x] || x * p > n) break;
            spf[x * p] = p;
           if (p == spf[x]) {
                phi[x * p] = phi[x] * p;
                mu[x * p] = 0;
                func[x * p] = func[x / pw[x]] * (cnt[x] + 1);
                cnt[x * p] = cnt[x] + 1;
                pw[x * p] = pw[x] * p;
           }
            else {
                phi[x * p] = phi[x] * phi[p];
                mu[x * p] = mu[x] * mu[p]; // or -mu[x]
                func[x * p] = func[x] * func[p];
```

```
cnt[x * p] = 1;
                 pw[x * p] = p;
     }
 }
6.10 Bitset sieve
  * Description: sieve of eratosthenes for large n (up to 1e9).
  * Time and space (tested on codeforces):
  * + For n = 1e8: ~200 ms, 6 MB.
  * + For n = 1e9: ~4000 ms. 60 MB.
 const int N = (int) 1e8;
 bitset<N / 2 + 1> isPrime:
 void sieve(int n = N) {
     isPrime.flip();
     isPrime[0] = false;
     for (int i = 3; i \le (int) sqrt(n); i += 2) {
         if (isPrime[i >> 1]) {
             for (int j = i * i; j <= n; j += 2 * i) {
                 isPrime[j >> 1] = false;
         }
     }
 }
 void example(int n) {
     sieve(n):
     int primeCnt = (n >= 2);
     for (int i = 3; i \le n; i += 2) {
         if (isPrime[i >> 1]) {
             primeCnt++;
         }
     }
     cout << primeCnt << '\n';</pre>
 }
6.11 Block sieve
  * Description: very fast sieve of eratosthenes for large n (up to 1e9).
  * Source: kactl.
  * Time and space (tested on codeforces):
  * + For n = 1e8: ~160 ms, 60 MB.
  * + For n = 1e9: ~1600 ms, 505 MB.
  * Need to check memory limit.
 const int N = (int) 1e8;
 bitset<N + 1> is_prime;
 vector<int> fast_sieve() {
     const int S = (int) sqrt(N), R = N / 2;
     vector<int> primes = {2}:
     vector<bool> sieve(S + 1, true);
```

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```
vector<array<int, 2>> cp;
for (int i = 3; i <= S; i += 2) {
    if (sieve[i]) {
        cp.push_back({i, i * i / 2});
        for (int j = i * i; j <= S; j += 2 * i) {
            sieve[j] = false;
        }
    }
}
for (int L = 1; L <= R; L += S) {
    array<bool, S> block{};
    for (auto &[p, idx] : cp) {
        for (; idx < S + L; idx += p) block[idx - L] = true;
    }
    for (int i = 0; i < min(S, R - L); ++i) {
        if (!block[i]) primes.push_back((L + i) * 2 + 1);
    }
}
for (int p : primes) is_prime[p] = true;
return primes;</pre>
```

6.12 Combinatorics

6.12.1 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{n!(n+1)!}$$

$$C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}, \ C_0 = 1, \ C_n = \frac{4n-2}{n+1} C_{n-1}$$

• The first 12 Catalan numbers (n = 0, 1, 2, ..., 11):

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$$

- Applications of Catalan numbers:
 - difference binary search trees with n vertices from 1 to n.
 - rooted binary trees with n + 1 leaves (vertices are not numbered).
 - correct bracket sequence of length 2 * n.
 - permutation [n] with no 3-term increasing subsequence (i.e. doesn't exist i < j < k for which a[i] < a[j] < a[k]).
 - ways a convex polygon of n + 2 sides can split into triangles by connecting vertices.

6.12.2 Fibonacci numbers

$$F_n = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F_{n-1} + F_{n-2}, & \text{otherwise} \end{cases}$$

• The first 20 Fibonacci numbers $(n = 0, 1, 2, \dots, 19)$:

$$F_n = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181$$

• Properties:

$$F_{2n+1} = F_n^2 + F_{n+1}^2$$

$$F_{2n} = F_{n-1} \cdot F_n + F_n \cdot F_{n+1}$$

$$F_{n+1} \cdot F_{n-1} - F_n^2 = (-1)^n$$

$$n \mid m \iff F_n \mid F_m$$

$$(F_n, F_m) = F_{(n,m)}$$

6.12.3 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* non-empty groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=1}^{k} (-1)^{k-i} \binom{k}{i} i^n$$

6.12.4 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixied point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

7 Geometry

7.1 Fundamentals

7.1.1 **Point**

```
#pragma once

const double PI = acos(-1);
const double EPS = 1e-9;
typedef double ftype;
struct Point {
   ftype x, y;
   Point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
```

```
Point& operator+=(const Point& other) {
        x += other.x; y += other.y; return *this;
   Point& operator -= (const Point& other) {
        x -= other.x; y -= other.y; return *this;
    Point& operator*=(ftype t) {
        x *= t; y *= t; return *this;
    Point& operator/=(ftype t) {
        x /= t; y /= t; return *this;
    Point operator+(const Point& other) const {
        return Point(*this) += other;
    Point operator-(const Point& other) const {
        return Point(*this) -= other;
   }
    Point operator*(ftype t) const {
        return Point(*this) *= t:
    }
    Point operator/(ftype t) const {
        return Point(*this) /= t;
    Point rotate(double angle) const {
        return Point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
    cos(angle));
    }
    friend istream& operator>>(istream &in, Point &t);
    friend ostream& operator<<(ostream &out, const Point& t);</pre>
    bool operator < (const Point& other) const {</pre>
        if (fabs(x - other.x) < EPS)</pre>
            return y < other.y;</pre>
        return x < other.x;</pre>
    }
};
istream& operator>>(istream &in, Point &t) {
    in >> t.x >> t.y;
    return in:
ostream& operator << (ostream &out, const Point& t) {
    out << t.x << ' ' << t.y;
    return out;
ftype dot(Point a, Point b) {return a.x * b.x + a.y * b.y;}
ftype norm(Point a) {return dot(a, a);}
ftype abs(Point a) {return sqrt(norm(a));}
ftype angle(Point a, Point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
ftype proj(Point a, Point b) {return dot(a, b) / abs(b);}
ftype cross(Point a, Point b) {return a.x * b.y - a.y * b.x;}
bool ccw(Point a, Point b, Point c) {return cross(b - a, c - a) > EPS;}
```

```
bool collinear(Point a, Point b, Point c) {return fabs(cross(b - a, c - a)) <</pre>
 Point intersect(Point a1, Point d1, Point a2, Point d2) {
     double t = cross(a2 - a1, d2) / cross(d1, d2);
     return a1 + d1 * t;
}
7.1.2 Line
 #include "point.h"
 struct Line {
     double a, b, c;
     Line (double a = 0, double b = 0, double c = 0): a(a), b(b), c(c) {}
     friend ostream & operator<<(ostream& out, const Line& 1);</pre>
 }:
 ostream & operator<<(ostream& out, const Line& 1) {</pre>
     out << 1.a << ' ' << 1.b << ' ' << 1.c;
     return out;
 void PointsToLine(const Point& p1, const Point& p2, Line& 1) {
     if (fabs(p1.x - p2.x) < EPS)
         1 = \{1.0, 0.0, -p1.x\};
     else {
         l.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
         1.b = 1.0;
         1.c = -1.a * p1.x - 1.b * p1.y;
 void PointsSlopeToLine(const Point& p, double m, Line& l) {
     1.a = -m:
     1.b = 1:
     1.c = -1.a * p.x - 1.b * p.y;
 bool areParallel(const Line& 11, const Line& 12) {
     return fabs(11.a - 12.a) < EPS && fabs(11.b - 12.b) < EPS;</pre>
 bool areSame(const Line& 11, const Line& 12) {
     return areParallel(11, 12) && fabs(11.c - 12.c) < EPS;</pre>
 bool areIntersect(Line 11, Line 12, Point& p) {
     if (areParallel(l1, l2)) return false;
     p.x = -(11.c * 12.b - 11.b * 12.c) / (11.a * 12.b - 11.b * 12.a);
     if (fabs(11.b) > EPS) p.y = -(11.c + 11.a * p.x);
     else p.y = -(12.c + 12.a * p.x);
     return 1;
 double distToLine(Point p, Point a, Point b, Point& c) {
     double t = dot(p - a, b - a) / norm(b - a);
     c = a + (b - a) * t;
     return abs(c - p);
 double distToSegment(Point p, Point a, Point b, Point& c) {
     double t = dot(p - a, b - a) / norm(b - a);
     if (t > 1.0)
```

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```
c = Point(b.x, b.y);
     else if (t < 0.0)
         c = Point(a.x, a.y);
     else
         c = a + (b - a) * t;
     return abs(c - p);
 bool intersectTwoSegment(Point a, Point b, Point c, Point d) {
     ftype ABxAC = cross(b - a, c - a);
     ftype ABxAD = cross(b - a, d - a);
     ftype CDxCA = cross(d - c, a - c);
     ftype CDxCB = cross(d - c, b - c);
     if (ABxAC == 0 \mid \mid ABxAD == 0 \mid \mid CDxCA == 0 \mid \mid CDxCB == 0) {
         if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;
         if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;</pre>
         if (CDxCA == 0 && dot(c - a, d - a) <= 0) return true;
         if (CDxCB == 0 && dot(c - b, d - b) <= 0) return true;</pre>
         return false:
     return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0):
 void perpendicular(Line 11, Point p, Line& 12) {
     if (fabs(l1.a) < EPS)</pre>
         12 = \{1.0, 0.0, -p.x\};
     else {
         12.a = -11.b / 11.a;
         12.b = 1.0:
         12.c = -12.a * p.x - 12.b * p.y;
     }
}
7.1.3 Circle
 #include "point.h"
 int insideCircle(const Point& p, const Point& center, ftype r) {
     ftype d = norm(p - center);
     ftype rSq = r * r;
     return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
 bool circle2PointsR(const Point& p1, const Point& p2, ftype r, Point& c) {
     double h = r * r - norm(p1 - p2) / 4.0;
     if (fabs(h) < 0) return false;</pre>
     h = sqrt(h);
     Point perp = (p2 - p1).rotate(PI / 2.0);
     Point m = (p1 + p2) / 2.0;
     c = m + perp * (h / abs(perp));
     return true;
 }
7.1.4 Triangle
 #include "point.h"
 #include "line.h"
 double areaTriangle(double ab, double bc, double ca) {
```

```
double p = (ab + bc + ca) / 2;
     return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
 double rInCircle(double ab, double bc, double ca) {
     double p = (ab + bc + ca) / 2;
     return areaTriangle(ab, bc, ca) / p;
 double rInCircle(Point a, Point b, Point c) {
     return rInCircle(abs(a - b), abs(b - c), abs(c - a));
 bool inCircle(Point p1, Point p2, Point p3, Point &ctr, double &r) {
     r = rInCircle(p1, p2, p3);
     if (fabs(r) < EPS) return false;</pre>
     Line 11. 12:
     double ratio = abs(p2 - p1) / abs(p3 - p1);
     Point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
     PointsToLine(p1, p, l1);
     ratio = abs(p1 - p2) / abs(p2 - p3);
     p = p1 + (p3 - p1) * (ratio / (1 + ratio));
     PointsToLine(p2, p, 12);
     areIntersect(l1, l2, ctr);
     return true;
 double rCircumCircle(double ab, double bc, double ca) {
     return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
 double rCircumCircle(Point a, Point b, Point c) {
     return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
7.1.5 Convex hull
 #include "point.h"
 vector<Point> CH_Andrew(vector<Point> &Pts) { // overall O(n log n)
     int n = Pts.size(), k = 0;
     vector<Point> H(2 * n);
     sort(Pts.begin(), Pts.end());
     for (int i = 0; i < n; ++i) {
         while ((k \ge 2) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
         H[k++] = Pts[i];
     for (int i = n - 2, t = k + 1; i >= 0; --i) {
         while ((k \ge t) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
         H\lceil k++\rceil = Pts\lceil i\rceil:
     H.resize(k);
     return H;
}
7.1.6 Polygon
#include "point.h"
 double perimeter(const vector<Point> &P) {
     double ans = 0.0;
```

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```
for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
         ans += abs(P[i] - P[i + 1]);
     return ans;
 double area(const vector<Point> &P) {
     double ans = 0.0;
     for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
         ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
     return fabs(ans) / 2.0;
 bool isConvex(const vector<Point> &P) {
     int n = (int)P.size();
     if (n <= 3) return false;</pre>
     bool firstTurn = ccw(P[0], P[1], P[2]);
     for (int i = 1; i < n - 1; ++i)
         if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
             return false;
     return true:
 int insidePolygon(Point pt, const vector<Point> &P) {
     int n = (int)P.size();
     if (n <= 3) return -1;
     bool on_polygon = false;
     for (int i = 0; i < n - 1; ++i)
         if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1]))
     < EPS)
             on_polygon = true;
     if (on_polygon) return 0;
     double sum = 0.0;
     for (int i = 0; i < n - 1; ++i) {
         if (ccw(pt, P[i], P[i + 1]))
             sum += angle(P[i] - pt, P[i + 1] - pt);
         else
             sum -= angle(P[i] - pt, P[i + 1] - pt);
     }
     return fabs(sum) > PI ? 1 : -1;
7.2 Minimum enclosing circle
  * Description: computes the minimum Circle that encloses all the given
     Points.
  */
 #include "point.h"
 // #include "circle.h"
 // TODO:
 Point center_from(double bx, double by, double cx, double cy) {
     double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by *
     return Point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
 }
```

```
Circle Circle_from(Point A, Point B, Point C) {
     Point I = center_from(B.x - A.x, B.y - A.y, C.x - A.x, C.y - A.y);
     return Circle(I + A, abs(I));
 }
 const int N = 100005;
 int n, x[N], y[N];
 Point a[N];
 Circle emo_welzl(int n, vector<Point> T) {
     if (T.size() == 3 || n == 0) {
         if (T.size() == 0) return Circle(Point(0, 0), -1);
         if (T.size() == 1) return Circle(T[0], 0);
         if (T.size() == 2) return Circle((T[0] + T[1]) / 2, abs(T[0] - T[1])
         return Circle_from(T[0], T[1], T[2]);
     random\_shuffle(a + 1, a + n + 1);
     Circle Result = emo welzl(0. T):
     for (int i = 1; i <= n; i++)</pre>
         if (abs(Result.x - a[i]) > Result.y + 1e-9) {
             T.push_back(a[i]);
             Result = emo_welzl(i - 1, T);
             T.pop_back();
     return Result;
    Linear algebra
8.1 Gauss elimination
 const double EPS = 1e-9;
 const int INF = 2; // it doesn't actually have to be infinity or a big number
 int gauss (vector < vector < double > > a, vector < double > & ans) {
     int n = (int) a.size();
     int m = (int) a[0].size() - 1;
     vector<int> where (m, -1);
     for (int col=0, row=0; col<m && row<n; ++col) {</pre>
         int sel = row;
         for (int i=row; i<n; ++i) {
             if (abs (a[i][col]) > abs (a[sel][col])) sel = i;
         if (abs (a[sel][col]) < EPS) continue;</pre>
         for (int i=col; i<=m; ++i) {</pre>
             swap (a[sel][i], a[row][i]);
```

where[col] = row;

for (int i=0; i<n; ++i) {
 if (i != row) {</pre>

double c = a[i][col] / a[row][col];

for (**int** j=col; j<=m; ++j) {

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```
a[i][j] -= a[row][j] * c;
         }
    }
    ++row;
ans.assign (m, 0);
for (int i=0; i<m; ++i) {</pre>
    if (where[i] != -1) {
         ans[i] = a[where[i]][m] / a[where[i]][i];
for (int i=0; i<n; ++i) {</pre>
    double sum = 0;
    for (int j=0; j<m; ++j) {</pre>
         sum += ans[j] * a[i][j];
    if (abs (sum - a[i][m]) > EPS) return 0;
for (int i=0; i<m; ++i) {</pre>
    if (where[i] == -1) return INF;
}
return 1;
```

9 Graph

9.1 Bellman-Ford algorithm

```
* Description: single source shortest path in a weighted (negative or
    positive) directed graph.
 * Time: O(N * M).
 * Tested: https://open.kattis.com/problems/shortestpath3
const int64_t INF = (int64_t) 2e18;
struct Edge {
    int u, v; // u -> v
    int64_t w;
    Edge() {}
    Edge(int \underline{u}, int \underline{v}, int64\underline{t} \underline{w}) : \underline{u}(\underline{u}), \underline{v}(\underline{v}), \underline{w}(\underline{w}) {}
};
int n;
vector<Edge> edges;
vector<int64_t> bellmanFord(int s) {
    // dist[stating] = 0.
    // dist[u] = +INF, if u is unreachable.
    // dist[u] = -INF, if there is a negative cycle on the path from s to u.
    // -INF < dist[u] < +INF, otherwise.</pre>
    vector<int64_t> dist(n, INF);
    dist[s] = 0;
    for (int i = 0; i < n - 1; ++i) {
         bool any = false;
         for (auto [u, v, w] : edges) {
```

```
if (dist[u] != INF && dist[v] > w + dist[u]) {
                dist[v] = w + dist[u];
                any = true;
            }
        if (!any) break;
    // handle negative cycles
    for (int i = 0; i < n - 1; ++i) {
        for (auto [u, v, w] : edges) {
            if (dist[u] != INF && dist[v] > w + dist[u]) {
                dist[v] = -INF;
        }
    }
    return dist;
    Articulation point and Bridge
 * Description: finding articulation points and bridges in a simple
    undirected graph.
 * Tested: https://oj.vnoi.info/problem/graph_
 */
const int N = (int) 1e5;
vector<int> g[N];
int num[N], low[N], dfs_timer;
bool joint[N];
vector<pair<int, int>> bridges;
void dfs(int u, int prev) {
    low[u] = num[u] = ++dfs_timer;
    int child = 0;
    for (int v : g[u]) {
        if (v == prev) continue;
        if (num[v]) low[u] = min(low[u], num[v]);
        else {
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            child++;
            if (low[v] >= num[v]) {
                bridges.emplace_back(u, v);
            if (u != prev && low[v] >= num[u]) joint[u] = true;
        }
    }
    if (u == prev && child > 1) joint[u] = true;
}
int solve() {
    int n, m;
    cin >> n >> m;
    for (int i = 0; i < m; ++i) {
        int u, v;
```

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(i.e, topo.size() < n).

```
cin >> u >> v;
         u--; v--;
         g[u].push_back(v);
        g[v].push_back(u);
     for (int i = 0; i < n; ++i) {
         if (!num[i]) dfs(i, i);
    }
     return 0;
}
     Strongly connected components
 /**
  * Description: Tarjan's algorithm finds strongly connected components
  * in a directed graph. If vertices u and v belong to the same component,
  * then scc_id[u] == scc_id[v].
  * Tested: https://judge.yosupo.jp/problem/scc
 const int N = (int) 5e5;
 vector<int> q[N], st;
 int low[N], num[N], dfs_timer, scc_id[N], scc;
 bool used[N];
 void Tarjan(int u) {
    low[u] = num[u] = ++dfs_timer;
     st.push_back(u);
     for (int v : g[u]) {
        if (used[v]) continue;
        if (num[v] == 0) {
             Tarjan(v);
             low[u] = min(low[u], low[v]);
        }
         else {
             low[u] = min(low[u], num[v]);
        }
    }
     if (low[u] == num[u]) {
        int v;
         do {
             v = st.back(); st.pop_back();
             debug(u, v)
             used[v] = true;
             scc_id[v] = scc;
        } while (v != u);
         scc++;
    }
9.4
   Topo sort
  * Description: A topological sort of a directed acyclic graph
  * is a linear ordering of its vertices such that for every directed edge
  * from vertex u to vertex v, u comes before v in the ordering.
  * Note: If there are cycles, the returned list will have size smaller than n
```

```
* Tested: https://judge.yosupo.jp/problem/scc
vector<int> topo_sort(const vector<vector<int>> &g) {
    int n = (int) g.size();
    vector<int> indeg(n);
    for (int u = 0; u < n; ++u) {
        for (int v : g[u]) indeg[v]++;
    queue < int > q; // Note: use min-heap to get the smallest lexicographical
    for (int u = 0; u < n; ++u) {
        if (indeg[u] == 0) q.emplace(u);
    vector<int> topo;
    while (!q.empty()) {
        int u = q.front(); q.pop();
        topo.emplace_back(u);
        for (int v : g[u]) {
            if (--indeg[v] == 0) q.emplace(v);
        }
    return topo;
}
    K-th smallest shortest path
/** Description: Finding the k-th smallest shortest path from vertex s to
    vertex t.
 * each vertex can be visited more than once.
using adj_list = vector<vector<pair<int, int>>>;
vector<long long> k_smallest(const adj_list &g, int k, int s, int t) {
    int n = (int) q.size();
    vector<long long> ans;
    vector<int> cnt(n);
    using pli = pair<long long, int>;
    priority_queue<pli, vector<pli>, greater<pli>> pq;
    pq.emplace(0, s);
    while (!pq.empty() && cnt[t] < k) {</pre>
        int u = pq.top().second;
        long long d = pq.top().first;
        pq.pop();
        if (cnt[u] == k) continue;
        cnt[u]++;
        if (u == t) {
            ans.push_back(d);
        for (auto [v, cost] : g[u]) {
            pq.emplace(d + cost, v);
    assert(k == (int) ans.size());
    return ans;
```

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```
}
9.6 Eulerian path
9.6.1 Directed graph
 /**
  * Hierholzer's algorithm.
  * Description: An Eulerian path in a directed graph is a path that visits
     all edges exactly once.
  * An Eulerian cycle is a Eulerian path that is a cycle.
  * Time complexity: O(|E|).
  */
  vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
     int n = (int) q.size();
     vector<int> stack, cur_edge(n), vertices;
     stack.push_back(s);
     while (!stack.empty()) {
         int u = stack.back();
         stack.pop_back();
         while (cur_edge[u] < (int) g[u].size()) {</pre>
             stack.push_back(u);
             u = g[u][cur_edge[u]++];
         vertices.push_back(u);
     reverse(vertices.begin(), vertices.end());
     return vertices;
 }
9.6.2 Undirected graph
 /**
  * Hierholzer's algorithm.
  * Description: An Eulerian path in a undirected graph is a path that visits
     all edges exactly once.
  * An Eulerian cycle is a Eulerian path that is a cycle.
  * Time complexity: O(|E|).
  */
 struct Edge {
     int to:
     list<Edge>::iterator reverse_edge;
     Edge(int _to) : to(_to) {}
 };
 vector<int> vertices;
 void find_path(vector<list<Edge>> &g, int u) {
     while (!g[u].empty()) {
         int v = g[u].front().to;
         g[v].erase(g[u].front().reverse_edge);
         g[u].pop_front();
         find_path(g, v);
     vertices.emplace_back(u); // reversion list.
 void add_edge(vector<list<Edge>> &g, int u, int v) {
     g[u].emplace_front(v);
```

```
g[v].emplace_front(u);
     g[u].front().reverse_edge = g[v].begin();
     g[v].front().reverse_edge = g[u].begin();
}
9.7 HLD
 const int INF = 0x3f3f3f3f3f;
 template < class SegmentTree >
 struct HLD { // vertex update and max query on path u -> v
     int n:
     vector<vector<int>> g;
     SegmentTree seg_tree;
     vector<int> par, top, depth, sz, id;
     int timer = 0;
     bool VAL_IN_EDGE = false;
     HLD() {}
     HLD(int _n): n(_n), g(n), seg_tree(n), par(n), top(n), depth(n), sz(n),
     id(n) {}
     void build() {
         dfs_sz(0);
         dfs_hld(0);
     void add_edge(int u, int v) {
         g[u].push_back(v);
         g[v].push_back(u);
     void dfs_sz(int u) {
         sz[u] = 1;
         for (int &v : g[u]) { // MUST BE ref for the swap below
             par[v] = u;
             depth[v] = depth[u] + 1;
             g[v].erase(find(g[v].begin(), g[v].end(), u));
             dfs_sz(v);
             sz[u] += sz[v];
             if (sz[v] > sz[g[u][0]]) swap(v, g[u][0]);
         }
     }
     void dfs_hld(int u) {
         id[u] = timer++;
         for (int v : g[u]) {
             top[v] = (v == g[u][0] ? top[u] : v);
             dfs hld(v):
         }
     int lca(int u, int v) {
         while (top[u] != top[v]) {
             if (depth[top[u]] > depth[top[v]]) swap(u, v);
             v = par[top[v]];
         // now u, v is in the same heavy-chain
         return (depth[u] < depth[v] ? u : v);</pre>
     void set_vertex(int v, int x) {
```

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```
seg_tree.set(id[v], x);
    void set_edge(int u, int v, int x) {
       if (u != par[v]) swap(u, v);
        seg_tree.set(id[v], x);
    void set_subtree(int v, int x) {
        // modify segment_tree so that it supports range update
        seg_tree.set_range(id[v] + VAL_IN_EDGE, id[v] + sz[v] - 1, x);
   int query_path(int u, int v) {
       int res = -INF;
       while (top[u] != top[v]) {
            if (depth[top[u]] > depth[top[v]]) swap(u, v);
            int cur = seg_tree.query(id[top[v]], id[v]);
            res = max(res, cur);
            v = par[top[v]];
       if (depth[u] > depth[v]) swap(u, v);
       int cur = seg_tree.query(id[u] + VAL_IN_EDGE, id[v]);
        res = max(res, cur);
       return res;
};
```

10 Misc.

10.1 Ternary search

```
* Description: given an unimodal function f(x), find the maximum/minimum of
    f(x).
* Unimodal means The function strictly increases/decreases first,
* reaches a maximum/minimum (at a single point or over an interval),
* and then strictly decreases/increases.
const double eps = 1e-9;
template<typename T>
inline T func(T x) { return x * x; }
// these two functions below find min, for find max: change '<' below to '>'.
double ternary_search(double 1, double r) { // min
    while (r - l > eps) {
        double mid1 = 1 + (r - 1) / 3;
        double mid2 = r - (r - 1) / 3;
       if (func(mid1) < func(mid2)) r = mid2;</pre>
       else 1 = mid1;
   }
   return 1;
int ternary_search(int 1, int r) { // min
    while (1 < r) {
        int mid = 1 + (r - 1) / 2;
       if (func(mid) < func(mid + 1)) r = mid;</pre>
```

```
else l = mid + 1:
     return 1;
}
10.2 Matrix
 using matrix_type = int;
 const int MOD = (int) 1e9 + 7;
 struct Matrix {
     static const matrix_type INF = numeric_limits<matrix_type>::max();
     int N, M;
     vector<vector<matrix_type>> mat;
     Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
         mat.assign(N, vector<matrix_type>(M, v));
     static Matrix identity(int n) { // return identity matrix.
         Matrix I(n. n):
         for (int i = 0; i < n; ++i) {</pre>
             I[i][i] = 1;
         return I;
     }
     vector<matrix_type>& operator[](int r) { return mat[r]; }
     const vector<matrix_type>& operator[](int r) const { return mat[r]; }
     Matrix& operator*=(const Matrix &other) {
         assert(M == other.N); // [N x M] [other.N x other.M]
         Matrix res(N, other.M);
         for (int r = 0; r < N; ++r) {
             for (int c = 0; c < other.M; ++c) {</pre>
                 long long square_mod = (long long) MOD * MOD;
                 long long sum = 0;
                 for (int g = 0; g < M; ++g) {
                     sum += (long long) mat[r][g] * other[g][c];
                     if (sum >= square_mod) sum -= square_mod;
                 res[r][c] = sum % MOD;
             }
         mat.swap(res.mat); return *this;
};
```