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## 1 Contest

#### 1.1 C++

```
#include <bits/stdc++.h>
 using namespace std;
 #ifdef LOCAL
 #include "cp/debug.h"
 #else
 #define debug(...)
 #endif
 mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
 int main() {
     cin.tie(nullptr)->sync_with_stdio(false);
     // freopen("input.txt", "r", stdin);
     // freopen("output.txt", "w", stdout);
     return 0:
}
1.2 Debug
 #define debug(...) { string _s = #__VA_ARGS__; replace(begin(_s), end(_s), ',',
     ''); stringstream _ss(_s); istream_iterator<string> _it(_ss);
     out_error(_it, __VA_ARGS__);}
 void out_error(istream_iterator<string> it) { cerr << '\n'; }</pre>
 template < typename T, typename ... Args >
 void out_error(istream_iterator<string> it, T a, Args... args) {
     cerr << " [" << *it << " = " << a << "] ";
     out_error(++it, args...);
 }
 template < typename T, typename G> ostream& operator << (ostream &os, const pair < T,
     return os << "(" << p.first << ", " << p.second << ")";</pre>
 }
 template < class Con, class = decltype(begin(declval < Con > ()))>
 typename enable_if<!is_same<Con, string>::value, ostream&>::type
 operator<<((ostream& os, const Con& container) {</pre>
     os << "{";
     for (auto it = container.begin(); it != container.end(); ++it)
         os << (it == container.begin() ? "" : ", ") << *it;
     return os << "}";</pre>
1.3 Java
 import java.io.BufferedReader;
 import java.util.StringTokenizer;
 import java.io.IOException;
```

```
import java.io.InputStreamReader;
 import java.io.PrintWriter;
 import java.util.ArrayList;
 import java.util.Arrays;
 import java.util.Collections;
 import java.util.Random;
 public class Main {
     public static void main(String[] args) {
         FastScanner fs = new FastScanner();
         PrintWriter out = new PrintWriter(System.out);
         int n = fs.nextInt();
         out.println(n);
         out.close(): // don't forget this line.
     static class FastScanner {
         BufferedReader br;
         StringTokenizer st;
         public FastScanner() {
             br = new BufferedReader(new InputStreamReader(System.in)):
             st = null;
         public String next() {
             while (st == null || st.hasMoreTokens() == false) {
                 trv {
                     st = new StringTokenizer(br.readLine());
                 catch (IOException e) {
                     throw new RuntimeException(e);
             }
             return st.nextToken();
         public int nextInt() {
             return Integer.parseInt(next());
         public long nextLong() {
             return Long.parseLong(next());
         public double nextDouble() {
             return Double.parseDouble(next());
1.4 sublime-build
     "cmd": ["g++", "-std=c++17", "-fmax-errors=5", "-DLOCAL", "-Wall",
     "-Wextra", "-o", "${file_path}/${file_base_name}.out", "${file}"],
     "file_regex": "^(..[^:]*):([0-9]+):?([0-9]+)?:? (.*)$",
```

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```
"working_dir": "${file_path}",
     "selector": "source.cpp, source.c++"
}
1.5 .bashrc
 alias c++='g++ -std=c++2a -fmax-errors=5 -DLOCAL -Wall -Wextra -02 -s'
 #Stress-testing
 function test {
  SOL = 1
  CHECKER=$2
   for i in {1..100};
    do
       ./gen.out > in && $CHECKER < in > ans && $SOL < in > out && diff -Zb out
     ans && echo "Test $i passed!!" || break;
     done
 }
    Data structures
2.1 Sparse table
 int st[MAXN][K + 1];
 for (int i = 0; i < N; i++) {
     st[i][0] = f(array[i]);
 for (int j = 1; j <= K; j++) {
     for (int i = 0; i + (1 << j) <= N; i++) {
         st[i][j] = f(st[i][j-1], st[i+(1 << (j-1))][j-1]);
```

```
}
}
// Range Minimum Queries.
int lg[MAXN + 1];
lq[1] = 0;
for (int i = 2; i <= MAXN; i++) {</pre>
   lg[i] = lg[i / 2] + 1;
}
int j = lg[R - L + 1];
int minimum = min(st[L][j], st[R - (1 << j) + 1][j]);</pre>
// Range Sum Queries.
long long sum = 0;
for (int j = K; j >= 0; j--) {
   if ((1 << j) <= R - L + 1) {
        sum += st[L][j];
        L += 1 << j;
```

## 2.2 Ordered set

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
template<typename key_type>
```

```
using set_t = tree<key_type, null_type, less<key_type>, rb_tree_tag,
     tree_order_statistics_node_update>;
 const int INF = 0x3f3f3f3f;
 void example() {
     vector < int > nums = \{1, 2, 3, 5, 10\};
     set_t<int> st(nums.begin(), nums.end());
     cout << *st.find_by_order(0) << '\n'; // 1</pre>
     assert(st.find_by_order(-INF) == st.end());
     assert(st.find_by_order(INF) == st.end());
     cout << st.order_of_key(2) << '\n'; // 1</pre>
     cout << st.order_of_key(4) << '\n'; // 3</pre>
     cout << st.order_of_key(9) << '\n'; // 4</pre>
     cout << st.order_of_key(-INF) << '\n'; // 0</pre>
     cout << st.order_of_key(INF) << '\n'; // 5</pre>
2.3 Dsu
 struct Dsu {
     int n;
     vector<int> par, sz;
     Dsu(int _n) : n(_n) {
         sz.resize(n, 1);
         par.resize(n);
         iota(par.begin(), par.end(), 0);
     int find(int v) {
         // finding leader/parrent of set that contains the element v.
         // with {path compression optimization}.
         return (v == par[v] ? v : par[v] = find(par[v]));
     bool same(int u, int v) {
         return find(u) == find(v);
     bool unite(int u, int v) {
         u = find(u); v = find(v);
         if (u == v) return false;
         if (sz[u] < sz[v]) swap(u, v);
         par[v] = u;
         sz[u] += sz[v];
         return true;
     vector<vector<int>> groups() {
         // returns the list of the "list of the vertices in a connected
     component".
         vector<int> leader(n);
         for (int i = 0; i < n; ++i) {</pre>
              leader[i] = find(i);
         vector<int> id(n, -1);
         int count = 0;
```

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};

```
for (int i = 0; i < n; ++i) {
             if (id[leader[i]] == -1) {
                 id[leader[i]] = count++;
         }
         vector<vector<int>> result(count);
         for (int i = 0; i < n; ++i) {
             result[id[leader[i]]].push_back(i);
         return result;
     }
 };
2.4 MinOueue
 /**
  * Description: acts like normal std::queue except it supports get minimum
     value in current queue.
 template <typename T>
 struct MinQueue {
     vector<T> vals;
     int ptr = 0;
     vector<int> st:
     int ptr_idx = 0;
     void push(T val) {
         while ((int) st.size() > ptr_idx && vals[st.back()] >= val) {
             st.pop_back();
         st.push_back((int) vals.size());
         vals.push_back(val);
     }
     void pop() {
         assert(ptr < (int) vals.size());</pre>
         if (ptr_idx < (int) st.size() && st[ptr_idx] == ptr) ptr_idx++;</pre>
         ptr++;
    T get() {
         assert(ptr_idx < (int) st.size());</pre>
         return vals[st[ptr_idx]];
     }
     int front() {
         assert(!empty()); return vals[ptr];
     }
     int back() {
         assert(!empty()); return vals.back();
     bool empty() {
         return (ptr == (int) vals.size());
     int size() {
         return ((int) vals.size() - ptr);
     }
```

# 2.5 Segment tree

```
* Description: A segment tree with range updates and sum queries that supports
    three types of operations:
* + Increase each value in range [1, r] by x (i.e. a[i] += x).
* + Set each value in range [1, r] to x (i.e. a[i] = x).
* + Determine the sum of values in range [1, r].
struct SegmentTree {
   vector<long long> tree, lazy_add, lazy_set;
   SegmentTree(int _n) : n(_n) {
       int p = 1;
       while (p < n) p *= 2;
       tree.resize(p * 2);
       lazy_add.resize(p * 2);
       lazy_set.resize(p * 2);
   long long merge(const long long &left, const long long &right) {
       return left + right;
   void build(int id, int 1, int r, const vector<int> &arr) {
       if (1 == r) {
           tree[id] += arr[l];
            return;
       int mid = (1 + r) >> 1:
       build(id * 2, 1, mid, arr);
       build(id * 2 + 1, mid + 1, r, arr);
       tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
   void push(int id, int 1, int r) {
       if (lazy_set[id] == 0 && lazy_add[id] == 0) return;
       int mid = (1 + r) >> 1;
       for (int child : {id * 2, id * 2 + 1}) {
            int range = (child == id * 2 ? mid - l + 1 : r - mid);
           if (lazy_set[id] != 0) {
                lazy_add[child] = 0;
                lazy_set[child] = lazy_set[id];
                tree[child] = range * lazy_set[id];
           lazy_add[child] += lazy_add[id];
            tree[child] += range * lazy_add[id];
       lazy_add[id] = lazy_set[id] = 0;
   void update(int id, int 1, int r, int u, int v, int amount, bool set_value
   = false) {
       if (r < u \mid | 1 > v) return;
       if (u <= 1 && r <= v) {
```

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```
if (set value) {
                 tree[id] = 1LL * amount * (r - l + 1);
                 lazy_set[id] = amount;
                 lazy_add[id] = 0; // clear all previous updates.
             else {
                 tree[id] += 1LL * amount * (r - 1 + 1);
                 lazy_add[id] += amount;
             }
             return:
         push(id, 1, r);
         int mid = (1 + r) >> 1;
         update(id * 2, 1, mid, u, v, amount, set_value);
         update(id * 2 + 1, mid + 1, r, u, v, amount, set_value);
         tree[id] = merge(tree[id * 2], tree[id * 2 + 1]);
     long long get(int id, int l, int r, int u, int v) {
         if (r < u \mid \mid 1 > v) return 0;
         if (u <= 1 && r <= v) {
             return tree[id];
         push(id, 1, r);
         int mid = (1 + r) >> 1;
         long long left = get(id * 2, 1, mid, u, v);
         long long right = get(id * 2 + 1, mid + 1, r, u, v);
         return merge(left, right);
 };
2.6 Efficient segment tree
 template < typename T> struct SegmentTree {
     int n;
     vector<T> tree;
     SegmentTree(int _n) : n(_n), tree(2 * n) {}
    T merge(const T &left, const T &right) {
         return left + right;
     template<typename G>
     void build(const vector<G> &initial) {
         assert((int) initial.size() == n);
         for (int i = 0; i < n; ++i) {
             tree[i + n] = initial[i];
         for (int i = n - 1; i > 0; --i) {
             tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
     void modify(int i, int v) {
         tree[i += n] = v;
         for (i /= 2; i > 0; i /= 2) {
             tree[i] = merge(tree[i * 2], tree[i * 2 + 1]);
         }
```

```
T get_sum(int 1, int r) {
         // sum of elements from 1 to r - 1.
         T ret{};
         for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
             if (1 & 1) ret = merge(ret, tree[1++]);
             if (r \& 1) ret = merge(ret, tree[--r]);
         }
         return ret;
     }
 };
2.7 Persistent lazy segment tree
 struct Vertex {
     int 1. r:
     long long val, lazy;
     bool has_changed = false;
     Vertex() {}
     Vertex(int _l, int _r, long long _val, int _lazy = 0) : l(_l), r(_r),
      val(_val), lazy(_lazy) {}
 };
 struct PerSegmentTree {
     vector<Vertex> tree:
     vector<int> root:
     int build(const vector<int> &arr, int 1, int r) {
         if (1 == r) {
             tree.emplace_back(-1, -1, arr[1]);
             return tree.size() - 1;
         int mid = (1 + r) / 2;
         int left = build(arr, 1, mid);
         int right = build(arr, mid + 1, r);
         tree.emplace_back(left, right, tree[left].val + tree[right].val);
         return tree.size() - 1;
     int add(int x, int 1, int r, int u, int v, int amt) {
         if (1 > v \mid | r < u) return x;
         if (u <= 1 && r <= v) {
             tree.emplace_back(tree[x].1, tree[x].r, tree[x].val + 1LL * amt *
      (r - l + 1), tree[x].lazy + amt);
             tree.back().has_changed = true;
             return tree.size() - 1;
         int mid = (1 + r) >> 1;
         push(x, 1, mid, r);
         int left = add(tree[x].1, 1, mid, u, v, amt);
         int right = add(tree[x].r, mid + 1, r, u, v, amt);
         tree.emplace_back(left, right, tree[left].val + tree[right].val, 0);
         return tree.size() - 1;
     long long get_sum(int x, int l, int r, int u, int v) {
         if (r < u \mid | 1 > v) return 0;
         if (u <= 1 && r <= v) return tree[x].val;</pre>
```

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```
int mid = (1 + r) / 2;
         push(x, 1, mid, r);
         return get_sum(tree[x].1, 1, mid, u, v) + get_sum(tree[x].r, mid + 1,
     r, u, v);
     void push(int x, int 1, int mid, int r) {
         if (!tree[x].has_changed) return;
         Vertex left = tree[tree[x].1];
         Vertex right = tree[tree[x].r];
         tree.emplace_back(left);
         tree[x].l = tree.size() - 1;
         tree.emplace_back(right);
         tree[x].r = tree.size() - 1;
         tree[tree[x].1].val += tree[x].lazy * (mid - l + 1);
         tree[tree[x].1].lazy += tree[x].lazy;
         tree[tree[x].r].val += tree[x].lazy * (r - mid);
         tree[tree[x].r].lazy += tree[x].lazy;
         tree[tree[x].1].has_changed = true;
         tree[tree[x].r].has_changed = true;
         tree[x].lazy = 0;
         tree[x].has_changed = false;
};
2.8 Lichao tree
 /**
 * Description: A segment tree that allows insert a new line and query for
     minimum value over all lines at point x.
  * Usage: useful in convex hull trick.
 const long long INF_LL = (long long) 4e18;
 struct Line {
    long long a, b;
    Line(long long _a = 0, long long _b = INF_LL): a(_a), b(_b) {}
    long long operator()(long long x) {
         return a * x + b;
    }
};
 struct SegmentTree { // min query
     int n;
     vector<Line> tree;
     SeamentTree() {}
     SegmentTree(int _n): n(1) {
         while (n < _n) n *= 2;
         tree.resize(n * 2);
     void insert(int x, int l, int r, Line line) {
```

```
if (1 == r) {
            if (line(l) < tree[x](l)) tree[x] = line;</pre>
        int mid = (1 + r) >> 1;
        bool b_left = line(l) < tree[x](l);</pre>
        bool b_mid = line(mid) < tree[x](mid);</pre>
        if (b_mid) swap(tree[x], line);
        if (b_left != b_mid) insert(x * 2, 1, mid, line);
        else insert(x * 2 + 1, mid + 1, r, line);
    long long query(int x, int 1, int r, int at) {
        if (l == r) return tree[x](at);
        int mid = (1 + r) >> 1:
        if (at <= mid) return min(tree[x](at), query(x * 2, 1, mid, at));</pre>
        else return min(tree[x](at), query(x * 2 + 1, mid + 1, r, at));
};
    Old driver tree (Chtholly tree)
 * Description: An optimized brute-force approach to deal with problem that has
    operation of setting an interval to the same number.
 * Note: caution TLE. only works when input is random
struct ODT {
    map<int, long long> tree;
    using It = map<int, long long>::iterator;
    It split(int x) {
        It it = tree.upper_bound(x);
        assert(it != tree.begin());
        --it;
        if (it->first == x) return it;
        return tree.emplace(x, it->second).first;
    }
    void add(int 1, int r, int amt) {
        It it_l = split(l);
        It it_r = split(r + 1);
        while (it_l != it_r) {
            it 1->second += amt:
            ++it_l;
    }
    void set(int 1, int r, int v) {
        It it_l = split(l);
        It it_r = split(r + 1);
        while (it_l != it_r) {
            tree.erase(it_l++);
        tree[1] = v;
```

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```
}
    long long kth_smallest(int 1, int r, int k) {
         // return the k-th smallest value in range [l..r]
         vector<pair<long long, int>> values; // pair(value, count)
         It it_l = split(l);
         It it_r = split(r + 1);
         while (it_l != it_r) {
             It prev = it_l++;
             values.emplace_back(prev->second, it_l->first - prev->first);
         sort(values.begin(), values.end());
         for (auto [value, cnt] : values) {
             if (k <= cnt) return value;</pre>
             k -= cnt;
         }
         return -1;
    int powmod(long long a, long long n, int mod);
    int sum_of_xth_power(int 1, int r, int x, int mod) {
         It it_l = split(l);
        It it_r = split(r + 1);
         int res = 0;
         while (it_l != it_r) {
             It prev = it_l++;
             res = (res + 1LL * (it_l->first - prev->first) *
     powmod(prev->second, x, mod)) % mod;
         }
         return res;
};
2.10
      Disjoint sparse table
  * Description: range query on a static array.
  * Time: O(1) per query.
  * Tested: stress-test.
 const int MOD = (int) 1e9 + 7;
 struct DisjointSparseTable { // product queries.
    int n, h;
     vector<vector<int>> dst:
     vector<int> lg;
     DisjointSparseTable(int _n) : n(_n) {
         h = 1; // in case n = 1: h = 0 !!.
         int p = 1;
         while (p < n) p *= 2, h++;
         lg.resize(p); lg[1] = 0;
         for (int i = 2; i < p; ++i) {
             lg[i] = 1 + lg[i / 2];
         dst.resize(h, vector<int>(n));
    }
```

```
void build(const vector<int> &A) {
         for (int lv = 0; lv < h; ++lv) {
             int len = (1 << lv);
             for (int k = 0; k < n; k += len * 2) {
                  int mid = min(k + len, n);
                  dst[lv][mid - 1] = A[mid - 1] % MOD;
                  for (int i = mid - 2; i >= k; --i) {
                      dst[lv][i] = 1LL * A[i] * dst[lv][i + 1] % MOD;
                  if (mid == n) break;
                  dst[lv][mid] = A[mid] % MOD;
                  for (int i = mid + 1; i < min(mid + len, n); ++i) {</pre>
                      dst[lv][i] = 1LL * A[i] * dst[lv][i - 1] % MOD;
             }
     int get(int 1, int r) {
         if (1 == r) {
             return dst[0][1];
         int i = lg[l ^ r];
         return 1LL * dst[i][1] * dst[i][r] % MOD;
 };
2.11 Fenwick tree
 /**
  * Description: range update and range sum query.
 using tree_type = long long;
 struct FenwickTree {
     vector<tree_type> fenw_coeff, fenw;
     FenwickTree() {}
     FenwickTree(int _n) : n(_n) {
         fenw_coeff.assign(n, 0); // fenwick tree with coefficient (n - i).
         fenw.assign(n, 0); // normal fenwick tree.
     template<typename G>
     void build(const vector<G> &A) {
         assert((int) A.size() == n);
         vector<int> diff(n);
         diff[0] = A[0];
         for (int i = 1; i < n; ++i) {
             diff[i] = A[i] - A[i - 1];
         fenw_coeff[0] = (long long) diff[0] * n;
         fenw[0] = diff[0];
         for (int i = 1; i < n; ++i) {
             fenw_coeff[i] = fenw_coeff[i - 1] + (long_long) diff[i] * (n - i);
             fenw[i] = fenw[i - 1] + diff[i];
```

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```
for (int i = n - 1; i >= 0; --i) {
             int j = (i \& (i + 1)) - 1;
             if (j >= 0) {
                 fenw_coeff[i] -= fenw_coeff[j];
                 fenw[i] -= fenw[j];
         }
     }
     void add(vector<tree_type> &fenw, int i, tree_type val) {
         while (i < n) {
             fenw[i] += val;
             i = (i + 1);
         }
     }
     tree_type __prefix_sum(vector<tree_type> &fenw, int i) {
         tree_type res{};
         while (i >= 0) {
             res += fenw[i];
             i = (i \& (i + 1)) - 1:
         return res;
     tree_type prefix_sum(int i) {
         return __prefix_sum(fenw_coeff, i) - __prefix_sum(fenw, i) * (n - i -
     1);
     }
     void range_add(int 1, int r, tree_type val) {
         add(fenw_coeff, 1, (n - 1) * val);
         add(fenw_coeff, r + 1, (n - r - 1) * (-val));
         add(fenw, l, val);
         add(fenw, r + 1, -val);
     tree_type range_sum(int 1, int r) {
         return prefix_sum(r) - prefix_sum(l - 1);
    }
 };
2.12 Fenwick tree 2D
 /**
  * Description: range update and range sum query on a 2D array.
 using tree_type = long long;
 struct FenwickTree2D {
     int n, m;
     vector<vector<tree_type> > fenw[4];
     FenwickTree2D(int _n, int _m) : n(_n), m(_m) {
         for (int i = 0; i < 4; i++) {
             fenw[i].resize(n, vector<tree_type>(m));
         }
     }
     void add(int u, int v, tree_type val) {
```

for (int i = u; i < n; i |= (i + 1)) {

```
for (int j = v; j < m; j | = (j + 1)) {
             fenw[0][i][j] += val;
             fenw[1][i][j] += (u + 1) * val;
             fenw[2][i][j] += (v + 1) * val;
             fenw[3][i][j] += (u + 1) * (v + 1) * val;
         }
     void range_add(int r, int c, int rr, int cc, tree_type val) { // [r, rr] x
     [c, cc].
         add(r, c, val);
         add(r, cc + 1, -val);
         add(rr + 1, c, -val);
         add(rr + 1, cc + 1, val);
     tree_type prefix_sum(int u, int v) {
         tree_type res{};
         for (int i = u; i >= 0; i = (i & (i + 1)) - 1) {
           for (int j = v; j >= 0; j = (j & (j + 1)) - 1) {
             res += (u + 2) * (v + 2) * fenw[0][i][j];
             res -= (v + 2) * fenw[1][i][j];
             res -= (u + 2) * fenw[2][i][j];
             res += fenw[3][i][j];
           }
         }
         return res;
     tree_type range_sum(int r, int c, int rr, int cc) { // [r, rr] x [c, cc].
         return prefix_sum(rr, cc) - prefix_sum(r - 1, cc) - prefix_sum(rr, c -
     1) + prefix_sum(r - 1, c - 1);
 };
2.13 Implicit treap
 struct Node {
     int val, prior, cnt;
     bool rev;
     Node *left, *right;
     Node() {}
     Node(int _val) : val(_val), prior(rng()), cnt(1), rev(false),
     left(nullptr), right(nullptr) {}
 // Binary search tree + min-heap.
 struct Treap {
     Node *root;
     Treap() : root(nullptr) {}
     int get_cnt(Node *n) { return n ? n->cnt : 0; }
     void upd_cnt(Node *&n) {
         if (n) n->cnt = get_cnt(n->left) + get_cnt(n->right) + 1;
     void push_rev(Node *treap) {
         if (!treap || !treap->rev) return;
         treap->rev = false;
```

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```
swap(treap->left, treap->right);
    if (treap->left) treap->left->rev ^= true;
    if (treap->right) treap->right->rev ^= true;
pair<Node*, Node*> split(Node *treap, int x, int smaller = 0) {
    if (!treap) return {};
    push_rev(treap);
    int idx = smaller + get_cnt(treap->left); // implicit val.
    if (idx <= x) {
        auto pr = split(treap->right, x, idx + 1);
        treap->right = pr.first;
        upd_cnt(treap);
        return {treap, pr.second};
    }
    else {
        auto pl = split(treap->left, x, smaller);
        treap->left = pl.second;
        upd_cnt(treap);
        return {pl.first, treap};
}
Node* merge(Node *1, Node *r) {
    push_rev(1); push_rev(r);
    if (!1 || !r) return (1 ? 1 : r);
    if (1->prior < r->prior) {
        l->right = merge(l->right, r);
        upd_cnt(1);
        return 1;
    }
    else {
        r->left = merge(1, r->left);
        upd_cnt(r);
        return r;
    }
}
void insert(int pos, int val) {
    if (!root) {
        root = new Node(val);
        return:
    }
    Node *1, *m, *r;
    m = new Node(val);
    tie(1, r) = split(root, pos - 1);
    root = merge(l, merge(m, r));
}
void erase(int pos_l, int pos_r) {
    Node *1, *m, *r;
    tie(l, r) = split(root, pos_l - 1);
    tie(m, r) = split(r, pos_r - pos_l);
    root = merge(1, r);
void reverse(int pos_l, int pos_r) {
    Node *1, *m, *r;
```

```
tie(l, r) = split(root, pos_l - 1);
          tie(m, r) = split(r, pos_r - pos_l);
          m->rev ^= true;
          root = merge(1, merge(m, r));
     int query(int pos_l, int pos_r);
           // returns answer for corresponding types of guery.
     void inorder(Node *n) {
          if (!n) return;
          push_rev(n);
          inorder(n->left);
          cout << n->val << ' ';</pre>
          inorder(n->right);
     }
     void print() {
          inorder(root);
          cout << '\n';
 };
2.14 Line container
 /**
  * Source: kactl
  * Description: container that allow you can add lines in form 'ax + b' and
      query maximum value at 'x'.
 using num_t = int;
 struct Line {
     num t a. b: // ax + b
     mutable num_t x; // x-intersect with the next line in the hull
     bool operator<(const Line &other) const {</pre>
          return a < other.a;</pre>
     bool operator<(num_t other_x) const {</pre>
          return x < other_x;</pre>
     }
 };
 struct LineContainer : multiset<Line, less<>>> { // max-query
     // for doubles, use INF = 1 / 0.0
     static const num_t INF = numeric_limits<num_t>::max();
     num_t floor_div(num_t a, num_t b) {
          return a / b - ((a ^ b) < 0 && a % b != 0);
     bool isect(iterator u, iterator v) {
         if (v == end()) {
              u->x = INF:
              return false:
          if (u->a == v->a) u->x = (u->b > v->b ? INF : -INF);
          else u \rightarrow x = floor_div(v \rightarrow b - u \rightarrow b, u \rightarrow a - v \rightarrow a);
          return u->x>=v->x:
```

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```
void add(num_t a, num_t b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin() && isect(--x, y)) {
        y = erase(y);
        isect(x, y);
    }
    while ((y = x) != begin() && (--x)->x >= y->x) {
        isect(x, erase(y));
    }
}
num_t query(num_t x) {
    assert(!empty());
    auto it = *lower_bound(x);
    return it.a * x + it.b;
}
```

## 3 Mathematics

};

## 3.1 Trigonometry

#### 3.1.1 Sum - difference identities

$$\sin(u \pm v) = \sin(u)\cos(v) \pm \cos(u)\sin(v)$$

$$\cos(u \pm v) = \cos(u)\cos(v) \mp \sin(u)\sin(v)$$

$$\tan(u \pm v) = \frac{\tan(u) \pm \tan(v)}{1 \mp \tan(u)\tan(v)}$$

## 3.1.2 Sum to product identities

$$\cos(u) + \cos(v) = 2\cos(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\cos(u) - \cos(v) = -2\sin(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

$$\sin(u) + \sin(v) = 2\sin(\frac{u+v}{2})\cos(\frac{u-v}{2})$$

$$\sin(u) - \sin(v) = 2\cos(\frac{u+v}{2})\sin(\frac{u-v}{2})$$

#### 3.1.3 Product identities

$$\cos(u)\cos(v) = \frac{1}{2}[\cos(u+v) + \cos(u-v)]$$
  

$$\sin(u)\sin(v) = -\frac{1}{2}[\cos(u+v) - \cos(u-v)]$$
  

$$\sin(u)\cos(v) = \frac{1}{2}[\sin(u+v) + \sin(u-v)]$$

#### 3.1.4 Double - triple angle identities

$$\sin(2u) = 2\sin(u)\cos(u)$$

$$\cos(2u) = 2\cos^{2}(u) - 1 = 1 - 2\sin^{2}(u)$$

$$\tan(2u) = \frac{2\tan(u)}{1 - \tan^{2}(u)}$$

$$\sin(3u) = 3\sin(u) - 4\sin^{3}(u)$$

$$\cos(3u) = 4\cos^{3}(u) - 3\cos(u)$$

$$\tan(3u) = \frac{3\tan(u) - \tan^{3}(u)}{1 - 3\tan^{2}(u)}$$

#### 3.2 Sums

$$\sum_{i=a}^{b} c^{i} = \frac{c^{b+1} - c^{a}}{c - 1}, c \neq 1$$

$$\sum_{i=0}^{n} i c^{i} = \frac{nc^{n+2} - (n+1)c^{n+1} + c}{(c - 1)^{2}}, c \neq 1$$

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i=1}^{n} i^{3} = \left(\frac{n(n+1)}{2}\right)^{2}$$

$$\sum_{i=1}^{n} i^{5} = \frac{n^{2}(n+1)^{2}(2n^{2} + 2n - 1)}{12}$$

$$\sum_{i=1}^{n} i^{5} = \frac{n^{2}(n+1)^{2}(3n^{4} + 6n^{3} - n^{2} - 4n + 2)}{24}$$

$$\sum_{i=0}^{n} i^{7} = \frac{n^{2}(n+1)^{2}(3n^{4} + 6n^{3} - n^{2} - 4n + 2)}{24}$$

$$\sum_{i=0}^{n} i^{n} \binom{n}{i} a^{n-i}b^{i} = (a+b)^{n}$$

$$\sum_{i=0}^{n} \binom{n}{i} = n2^{n-1}$$

$$\sum_{i=0}^{n} \binom{n}{i} = n2^{n-1}$$

$$\sum_{i=0}^{n} \binom{n}{i} = \binom{n+1}{k+1}$$

$$\sum_{i=0}^{n} \binom{n+k}{k} = \binom{n+m+1}{n+1}$$

$$\sum_{i=0}^{n} \binom{n+k}{k} = \binom{n+m+1}{k+1}$$

# 3.3 Pythagorean triple

- A Pythagorean triple is a triple of positive integers a, b, and c such that  $a^2 + b^2 = c^2$ .
- If (a, b, c) is a Pythagorean triple, then so is (ka, kb, kc) for any positive integer k.
- A primitive Pythagorean triple is one in which *a*, *b*, and *c* are coprime.
- Generating Pythagorean triple

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– Euclid's formula: with arbitrary 0 < n < m, then:

$$a = m^2 - n^2$$
,  $b = 2mn$ ,  $c = m^2 + n^2$ 

form a Pythagorean triple.

- To obtain primitive Pythagorean triple, this condition must hold: *m* and *n* are coprime, m and n have opposite parity.

# String

#### 4.1 Prefix function

```
/**
* Description: The prefix function of a string 's' is defined as an array pi
    of length n,
 * where pi[i] is the length of the longest proper prefix of the substring
 * s[0..i] which is also a suffix of this substring.
* Time complexity: O(|S|).
vector<int> prefix_function(const string &s) {
   int n = (int) s.length();
   vector<int> pi(n);
   pi[0] = 0;
   for (int i = 1; i < n; ++i) {
        int j = pi[i - 1]; // try length pi[i - 1] + 1.
        while (j > 0 \&\& s[j] != s[i]) {
           j = pi[j - 1];
       if (s[j] == s[i]) {
           pi[i] = j + 1;
   return pi;
   Z function
    suffix starting at 'i'.
```

```
* Description: for a given string 's', z[i] = longest common prefix of 's' and
    z[0] is generally not well defined (this implementation below assume z[0]
    = 0).
vector<int> z_function(const string &s) {
   int n = (int) s.size();
   vector<int> z(n);
   z[0] = 0;
   // [1, r)
   for (int i = 1, l = 0, r = 0; i < n; ++i) {
       if (i < r) z[i] = min(r - i, z[i - l]);
        while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]]) ++z[i];
       if (i + z[i] > r) {
```

```
1 = i:
            r = i + z[i];
    return z;
}
#include "prefix_function.h"
```

## Counting occurrences of each prefix

```
vector<int> count_occurrences(const string &s) {
    vector<int> pi = prefix_function(s);
   int n = (int) s.size();
   vector < int > ans(n + 1);
   for (int i = 0; i < n; ++i) {</pre>
        ans[pi[i]]++;
   for (int i = n - 1; i > 0; --i) {
        ans[pi[i - 1]] += ans[i];
   for (int i = 0; i <= n; ++i) {
        ans[i]++;
   return ans;
   // Input: ABACABA
    // Output: 4 2 2 1 1 1 1
```

#### Knuth-Morris-Pratt algorithm 4.4

```
* Searching for a substring in a string.
* Time complexity: O(N + M).
#include "prefix_function.h"
vector<int> KMP(const string &text, const string &pattern) {
    int n = (int) text.length();
    int m = (int) pattern.length();
    string s = pattern + '$' + text;
    vector<int> pi = prefix_function(s);
    vector<int> indices;
    for (int i = 0; i < (int) s.length(); ++i) {</pre>
        if (pi[i] == m) {
            indices.push_back(i - 2 * m);
        }
    return indices;
```

#### 4.5 Suffix array

\* Description: suffix array is a sorted array of all the suffixes of a given string.

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```
* Usage:
     sa[i] = starting index of the i-th smallest suffix.
    rank[i] = rank of the suffix starting at 'i'.
   lcp[i] = longest common prefix between 'sa[i - 1]' and 'sa[i]'
     for arbitrary 'u v', let i = rank[u] - 1, j = rank[v] - 1 (assume i < j),
       longest\_common\_prefix(u, v) = min(lcp[i + 1], lcp[i + 2], ..., lcp[j])
 * Time: O(NlogN).
struct SuffixArray {
    string s;
    int n, lim;
    vector<int> sa, lcp, rank;
    SuffixArray(const string &_s, int _lim = 256) : s(_s), n(s.length() + 1),
        \lim(_{\lim}, \sin(n), \log(n), rank(n))
        s += '$';
        build(); kasai();
        sa.erase(sa.begin()); lcp.erase(lcp.begin());
        rank.pop_back(); s.pop_back();
    void build() {
        vector<int> nrank(n), norder(n), cnt(max(n, lim));
        for (int i = 0; i < n; ++i) {
            sa[i] = i; rank[i] = s[i];
        for (int k = 0, rank_cnt = 0; rank_cnt < n - 1; k = max(1, k * 2), lim
    = rank_cnt + 1) {
            for (int i = 0; i < n; ++i) {
                norder[i] = (sa[i] - k + n) \% n;
                cnt[rank[i]]++;
            for (int i = 1; i < lim; ++i) cnt[i] += cnt[i - 1];</pre>
            for (int i = n - 1; i >= 0; --i) sa[--cnt[rank[norder[i]]]] =
    norder[i];
            rank[sa[0]] = rank\_cnt = 0;
            for (int i = 1; i < n; ++i) {
                int u = sa[i], v = sa[i - 1];
                int nu = (u + k) \% n, nv = (v + k) \% n;
                if (rank[u] != rank[v] || rank[nu] != rank[nv]) ++rank_cnt;
                nrank[sa[i]] = rank_cnt;
            for (int i = 0; i < rank_cnt + 1; ++i) cnt[i] = 0;</pre>
            rank.swap(nrank);
        }
    }
    void kasai() {
        for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
            int j = sa[rank[i] - 1];
            while (s[i + k] == s[j + k]) k++;
            lcp[rank[i]] = k;
        }
   }
};
```

## 4.6 Suffix array slow

```
/**
 * Description: an easier way to implement suffix array but run slower
 * Time: O(N * logN^2)
struct SuffixArraySlow {
    string s;
    int n;
    vector<int> sa, lcp, rank;
    SuffixArraySlow(const string &_s): s(_s), n((int) s.size() + 1), sa(n),
    lcp(n), rank(n) {
        s += '$';
        build(); kasai();
        sa.erase(sa.begin()); lcp.erase(lcp.begin());
        rank.pop_back(); s.pop_back();
    bool comp(int i, int j, int k) {
        return make_pair(rank[i], rank[(i + k) % n]) < make_pair(rank[j],</pre>
    rank[(i + k) % n]);
    void build() {
        vector<int> nrank(n);
        for (int i = 0; i < n; ++i) {
            sa[i] = i; rank[i] = s[i];
        for (int k = 0; k < n; k = max(1, k * 2)) {
            stable_sort(sa.begin(), sa.end(), [&](int i, int j) {
                return comp(i, j, k);
            });
            for (int i = 0, cnt = 0; i < n; ++i) {
                if (i > 0 \&\& comp(sa[i - 1], sa[i], k)) ++cnt;
                nrank[sa[i]] = cnt;
            rank.swap(nrank);
        }
    void kasai() {
        for (int i = 0, k = 0; i < n - 1; ++i, k = max(0, k - 1)) {
            int j = sa[rank[i] - 1];
            while (s[i + k] == s[j + k]) ++k;
            lcp[rank[i]] = k;
        }
};
    Manacher's algorithm
/**
 * Description: for each position, computes d[0][i] = half length of
 longest palindrome centered on i (rounded up), d[1][i] = half length of
 longest palindrome centered on i and i - 1.
 * Time complexity: O(N).
 * Tested: https://judge.yosupo.jp/problem/enumerate_palindromes, stress-tested.
```

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```
array<vector<int>, 2> manacher(const string &s) {
    int n = (int) s.size();
    array<vector<int>, 2> d;
    for (int z = 0; z < 2; ++z) {
        d[z].resize(n);
        int 1 = 0, r = 0;
        for (int i = 0; i < n; ++i) {
            int mirror = 1 + r - i + z;
            d[z][i] = (i > r ? 0 : min(d[z][mirror], r - i));
            int L = i - d[z][i] - z, R = i + d[z][i];
            while (L >= 0 \&\& R < n \&\& s[L] == s[R]) {
                d[z][i]++; L--; R++;
            if (R > r) {
                l = L; r = R;
            }
        }
    return d;
    Trie
struct Trie {
    const static int ALPHABET = 26;
    const static char minChar = 'a';
    struct Vertex {
        int next[ALPHABET];
        bool leaf;
        Vertex() {
            leaf = false;
            fill(next, next + ALPHABET, -1);
    };
    vector<Vertex> trie;
   Trie() { trie.emplace_back(); }
    void insert(const string &s) {
        int i = 0;
        for (const char &ch : s) {
            int j = ch - minChar;
            if (trie[i].next[j] == -1) {
                trie[i].next[j] = trie.size();
                trie.emplace_back();
            i = trie[i].next[j];
        trie[i].leaf = true;
    bool find(const string &s) {
        int i = 0;
        for (const char &ch : s) {
            int j = ch - minChar;
            if (trie[i].next[j] == -1) {
```

```
return false:
             i = trie[i].next[j];
         return (trie[i].leaf ? true : false);
 };
4.9 Hashing
 struct Hash61 {
     static const uint64_t MOD = (1LL << 61) - 1;</pre>
     static uint64_t BASE;
     static vector<uint64_t> pw;
     uint64_t addmod(uint64_t a, uint64_t b) const {
         a += b:
         if (a >= MOD) a -= MOD;
         return a:
     }
     uint64_t submod(uint64_t a, uint64_t b) const {
         a += MOD - b;
         if (a >= MOD) a -= MOD;
         return a;
     uint64_t mulmod(uint64_t a, uint64_t b) const {
         uint64_t low1 = (uint32_t) a, high1 = (a >> 32);
         uint64_t low2 = (uint32_t) b, high2 = (b >> 32);
         uint64_t low = low1 * low2;
         uint64_t mid = low1 * high2 + low2 * high1;
         uint64_t high = high1 * high2;
         uint64_t ret = (low & MOD) + (low >> 61) + (high << 3) + (mid >> 29) +
      (mid << 35 >> 3) + 1;
         // ret %= MOD:
         ret = (ret >> 61) + (ret & MOD);
         ret = (ret >> 61) + (ret & MOD);
         return ret - 1;
     void ensure_pw(int m) {
         int sz = (int) pw.size();
         if (sz >= m) return;
         pw.resize(m):
         for (int i = sz; i < m; ++i) {</pre>
              pw[i] = mulmod(pw[i - 1], BASE);
         }
     }
     vector<uint64_t> pref;
     template < typename T > Hash61(const T &s) { // strings or arrays.
         n = (int) s.size();
         ensure_pw(n);
         pref.resize(n + 1);
```

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```
pref[0] = 0;
         for (int i = 0; i < n; ++i) {
             pref[i + 1] = addmod(mulmod(pref[i], BASE), s[i]);
     inline uint64_t operator()(const int from, const int to) const {
         assert(0 \le from \&\& from \le to \&\& to < n);
         // pref[to + 1] - pref[from] * pw[to - from + 1]
         return submod(pref[to + 1], mulmod(pref[from], pw[to - from + 1]));
     }
 };
 mt19937 rnd((unsigned int)
     chrono::steady_clock::now().time_since_epoch().count());
 uint64_t Hash61::BASE = (MOD >> 2) + rnd() % (MOD >> 1);
 vector<uint64_t> Hash61::pw = vector<uint64_t>(1, 1);
      Minimum rotation
 /**
  * Author: Stjepan Glavina
  * License: Unlicense
  * Source: https://github.com/stjepang/snippets/blob/master/min_rotation.cpp
  * Description: Finds the lexicographically smallest rotation of a string.
  * Time: O(N)
  * Usage:
  * rotate(v.begin(), v.begin()+minRotation(v), v.end());
  * Status: Stress-tested
 #pragma once
 int minRotation(string s) {
   int a = 0, N = (int) s.size(); s += s;
   rep(b, 0, N) rep(k, 0, N) {
    if (a + k == b \mid | s[a + k] < s[b + k]) \{b += max(0, k - 1); break;\}
     if (s[a + k] > s[b + k]) \{ a = b; break; \}
   return a;
    Numerical
5.1 Fast Fourier transform
 const double PI = acos(-1):
```

```
const double PI = acos(-1);
using Comp = complex<double>;
int reverse_bit(int n, int lg) {
    int res = 0;
    for (int i = 0; i < lg; ++i) {
        if (n & (1 << i)) {
            res |= (1 << (lg - i - 1));
        }
    }
    return res;
}
void fft(vector<Comp> &a, bool invert = false) {
```

```
int n = (int) a.size();
    int lg = 0;
    while (1 << (lg) < n) ++lg;</pre>
    for (int i = 0; i < n; ++i) {
        int rev_i = reverse_bit(i, lg);
        if (i < rev_i) swap(a[i], a[rev_i]);</pre>
    for (int len = 2; len <= n; len *= 2) {</pre>
        double angle = 2 * PI / len * (invert ? -1 : 1);
        Comp w_base(cos(angle), sin(angle));
        for (int i = 0; i < n; i += len) {</pre>
            Comp w(1);
            for (int j = i; j < i + len / 2; ++j) {
                 Comp u = a[j], v = a[j + len / 2];
                a[j] = u + w * v;
                a[i + len / 2] = u - w * v;
                w *= w_base;
            }
        }
    if (invert) for (int i = 0; i < n; ++i) a[i] /= n;</pre>
vector<int> mult(vector<int> &a, vector<int> &b) {
    vector<Comp> A(a.begin(), a.end()), B(b.begin(), b.end());
    int n = (int) a.size(), m = (int) b.size(), p = 1;
    while (p < n + m) p *= 2;
    A.resize(p), B.resize(p);
    fft(A, false);
    fft(B, false);
    for (int i = 0; i < p; ++i) {
        A[i] *= B[i];
    fft(A, true);
    vector < int > res(n + m - 1);
    for (int i = 0; i < n + m - 1; ++i) {
        res[i] = (int) round(A[i].real());
    return res;
```

# 6 Number Theory

#### 6.1 Euler's totient function

- Euler's totient function, also known as **phi-function**  $\phi(n)$  counts the number of integers between 1 and n inclusive, that are **coprime to** n.
- Properties:
  - Divisor sum property:  $\sum_{d|n} \phi(d) = n$ .
  - $\phi(n)$  is a **prime number** when n = 3, 4, 6.
  - If p is a prime number, then  $\phi(p) = p 1$ .

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- If *p* is a prime number and *k* ≥ 1, then  $φ(p^k) = p^k p^{k-1}$ .
- If *a* and *b* are **coprime**, then  $\phi(ab) = \phi(a) \cdot \phi(b)$ .
- In general, for **not coprime** a and b, with d = gcd(a,b) this equation holds:  $\phi(ab) = \phi(a) \cdot \phi(b) \cdot \frac{d}{\phi(d)}$ .
- With  $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$ :

$$\phi(n) = \phi(p_1^{k_1}) \cdot \phi(p_2^{k_2}) \cdots \phi(p_m^{k_m})$$
$$= n \cdot \left(1 - \frac{1}{p_1}\right) \cdot \left(1 - \frac{1}{p_2}\right) \cdots \left(1 - \frac{1}{p_m}\right)$$

- Application in Euler's theorem:
  - If gcd(a, M) = 1, then:

$$a^{\phi(M)} \equiv 1 \pmod{M} \Rightarrow a^n \equiv a^{n \bmod{\phi(M)}} \pmod{M}$$

- In general, for arbitrary a, M and n ≥  $\log_2 M$ :

$$a^n \equiv a^{\phi(M) + [n \bmod \phi(M)]} \pmod{M}$$

#### 6.2 Mobius function

• For a positive integer  $n = p_1^{k_1} \cdot p_2^{k_2} \cdots p_m^{k_m}$ :

$$\mu(n) = \begin{cases} 1, & \text{if } n = 1 \\ 0, & \text{if } \exists k_i > 1 \\ (-1)^m & \text{otherwise} \end{cases}$$

- Properties:
  - $-\sum_{d|n}\mu(d)=[n=1].$
  - If *a* and *b* are **coprime**, then  $\mu(ab) = \mu(a) \cdot \mu(b)$ .
  - Mobius inversion: let *f* and *g* be arithmetic functions:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right)g(d)$$

6.3 Primes

Approximating the number of primes up to *n*:

п	$\pi(n)$	$\frac{n}{\ln n - 1}$
$100 (1e^2)$	25	28
$500 (5e^2)$	95	96
$1000 (1e^3)$	168	169
$5000 (5e^3)$	669	665
$10000 (1e^4)$	1229	1218
$50000 (5e^4)$	5133	5092
$100000 (1e^5)$	9592	9512
$500000 (5e^5)$	41538	41246
$1000000 (1e^6)$	78498	78030
$5000000 (5e^6)$	348513	346622

 $(\pi(n))$  = the number of primes less than or equal to n,  $\frac{n}{\ln n - 1}$  is used to approximate  $\pi(n)$ ).

## 6.4 Wilson's theorem

A positive integer *n* is a prime if and only if:

$$(n-1)! \equiv n-1 \pmod{n}$$

#### 6.5 Zeckendorf's theorem

The Zeckendorf's theorem states that every positive integer n can be represented uniquely as a sum of one or more distinct non-consecutive Fibonacci numbers. For example: 64 = 55 + 8 + 1

$$85 = 55 + 21 + 8 + 1$$

```
vector<int> zeckendoft_theorem(int n) {
    vector<int> fibs = {1, 1};
    int sz = 2;
    while (fibs.back() <= n) {
        fibs.push_back(fibs[sz - 1] + fibs[sz - 2]);
        sz++;
    }
    fibs.pop_back();
    vector<int> nums;
    int p = sz - 1;
    while (n > 0) {
        if (n >= fibs[p]) {
            nums.push_back(fibs[p]);
            n -= fibs[p];
        }
        p--;
    }
    return nums;
}
```

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## 6.6 Bitwise operation

```
• a + b = (a \oplus b) + 2(a \& b)

• a \mid b = (a \oplus b) + (a \& b)

• a \& (b \oplus c) = (a \& b) \oplus (a \& c)

• a \mid (b \& c) = (a \mid b) \& (a \mid c)

• a \& (b \mid c) = (a \& b) \mid (a \& c)

• a \& (b \mid c) = (a \& b) \mid (a \& c)

• a \& (b \mid c) = (a \& b) \mid (a \& c)

• a \mid (a \& b) = a

• a \& (a \mid b) = a
```

• Iterating over all subsets of a set and iterating over all submasks of a mask:

## 6.7 Pollard's rho algorithm

```
using num_t = long long;
const int PRIME_MAX = (int) 4e4; // for handle numbers <= 1e9.</pre>
const int LIMIT = (int) 1e9;
vector<int> primes;
int small_primes[] = {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 73, 113, 193,
    311, 313, 407521, 299210837};
void linear_sieve(int n);
num_t mulmod(num_t a, num_t b, num_t mod);
num_t powmod(num_t a, num_t n, num_t mod);
bool miller_rabin(num_t a, num_t d, int s, num_t mod) {
    num_t x = powmod(a, d, mod);
   if (x == mod - 1 || x == 1) {
        return true:
    for (int i = 0; i < s - 1; ++i) {
        x = mulmod(x, x, mod);
        if (x == mod - 1) return true;
    return false:
bool is_prime(num_t n, int tests = 10) {
   if (n < 4) return (n > 1);
```

```
num_t d = n - 1;
    int s = 0;
    while (d % 2 == 0) { d >>= 1; s++; }
    for (int i = 0; i < tests; ++i) {
        int a = small_primes[i];
        if (n == a) return true;
        if (n % a == 0 || !miller_rabin(a, d, s, n)) return false;
    return true;
num_t f(num_t x, int c, num_t mod) { // f(x) = (x^2 + c) % mod.
    x = mulmod(x, x, mod);
    x += c;
    if (x >= mod) x -= mod;
    return x;
num_t pollard_rho(num_t n, int c) {
    // algorithm to find a random divisor of 'n'.
    // using random function: f(x) = (x^2 + c) \% n.
    num_t x = 2, y = x, d;
    long long p = 1;
    int dist = 0;
    while (true) {
        y = f(y, c, n);
        dist++:
        d = \_gcd(llabs(x - y), n);
        if (d > 1) break;
        if (dist == p) { dist = 0; p *= 2; x = y; }
    return d;
void factorize(int n, vector<num_t> &factors);
void llfactorize(num_t n, vector<num_t> &factors) {
    if (n < 2) return;</pre>
    if (is_prime(n)) {
        factors.emplace_back(n);
        return;
    if (n < LIMIT) {</pre>
        factorize(n, factors);
        return;
    num_t d = n;
    for (int c = 2; d == n; c++) {
        d = pollard_rho(n, c);
    llfactorize(d, factors);
    llfactorize(n / d, factors);
vector<num_t> gen_divisors(vector<pair<num_t, int>> &factors) {
    vector<num_t> divisors = {1};
    for (auto &x : factors) {
        int sz = (int) divisors.size();
```

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```
for (int i = 0; i < sz; ++i) {
            num_t cur = divisors[i];
            for (int j = 0; j < x.second; ++j) {
                cur *= x.first;
                divisors.push_back(cur);
            }
        }
    }
    return divisors; // this array is NOT sorted yet.
}
    Segment divisor sieve
const int MAXN = (int) 1e6; // R - L + 1 <= N.
int divisor_count[MAXN + 3];
void segment_divisor_sieve(long long L, long long R) {
    for (long long i = 1; i \leftarrow (long long) sqrt(R); ++i) {
        long long start1 = ((L + i - 1) / i) * i;
        long long start2 = i * i;
        long long j = max(start1, start2);
        if (i == start2) {
            divisor_count[j - L] += 1;
            j += i;
        }
        for ( ; j <= R; j += i) {
            divisor_count[j - L] += 2;
    Linear sieve
 * Description: Finding primes and computing value for multiplicative function
    in O(N)
const int N = (int) 1e6;
bool is_prime[N + 1];
int spf[N + 1]; // smallest prime factor
int phi[N + 1]; // euler's totient function
int mu[N + 1]; // mobius function
int func[N + 1]; // a multiplicative function, f(p^k) = k
int cnt[N + 1]; // cnt[i] = the power of the smallest prime factor of i
int pw[N + 1]; // pw[i] = p^cnt[i] where p is the smallest prime factor of i
vector<int> primes;
void sieve(int n = N) {
    spf[0] = spf[1] = -1;
    phi[1] = mu[1] = func[1] = 1;
    for (int x = 2; x <= n; ++x) {
        if (spf[x] == 0) {
            primes.push_back(spf[x] = x);
            is_prime[x] = true;
            phi[x] = x - 1;
```

```
mu[x] = -1;
             func[x] = 1;
             cnt[x] = 1;
             pw[x] = x;
         for (int p : primes) {
             if (p > spf[x] \mid | x * p > n) break;
             spf[x * p] = p;
             if (p == spf[x]) {
                 phi[x * p] = phi[x] * p;
                 mu[x * p] = 0;
                 func[x * p] = func[x / pw[x]] * (cnt[x] + 1);
                 cnt[x * p] = cnt[x] + 1;
                 pw[x * p] = pw[x] * p;
             else {
                 phi[x * p] = phi[x] * phi[p];
                 mu[x * p] = mu[x] * mu[p]; // or -mu[x]
                 func[x * p] = func[x] * func[p];
                 cnt[x * p] = 1;
                 pw[x * p] = p;
 }
6.10 Bitset sieve
  * Description: sieve of eratosthenes for large n (up to 1e9).
  * Time and space (tested on codeforces):
  * + For n = 1e8: ~200 ms, 6 MB.
  * + For n = 1e9: ~4000 ms, 60 MB.
 const int N = (int) 1e8;
 bitset<N / 2 + 1> isPrime;
 void sieve(int n = N) {
     isPrime.flip();
     isPrime[0] = false;
     for (int i = 3; i \le (int) sqrt(n); i += 2) {
         if (isPrime[i >> 1]) {
             for (int j = i * i; j <= n; j += 2 * i) {
                 isPrime[j >> 1] = false;
     }
 void example(int n) {
     sieve(n):
     int primeCnt = (n >= 2);
     for (int i = 3; i \le n; i += 2) {
         if (isPrime[i >> 1]) {
             primeCnt++;
         }
```

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```
cout << primeCnt << '\n';</pre>
6.11 Block sieve
  * Description: very fast sieve of eratosthenes for large n (up to 1e9).
  * Source: kactl.
  * Time and space (tested on codeforces):
  * + For n = 1e8: ~160 ms, 60 MB.
  * + For n = 1e9: ~1600 ms, 505 MB.
  * Need to check memory limit.
  */
 const int N = (int) 1e8;
 bitset<N + 1> is_prime;
 vector<int> fast_sieve() {
     const int S = (int)  sqrt(N), R = N / 2;
     vector<int> primes = {2};
     vector<bool> sieve(S + 1, true);
     vector<array<int, 2>> cp;
     for (int i = 3; i \le S; i += 2) {
         if (sieve[i]) {
             cp.push_back({i, i * i / 2});
             for (int j = i * i; j \le S; j += 2 * i) {
                 sieve[i] = false;
     for (int L = 1; L <= R; L += S) {
         array<bool, S> block{};
         for (auto &[p, idx] : cp) {
             for (; idx < S + L; idx += p) block[idx - L] = true;</pre>
         for (int i = 0; i < min(S, R - L); ++i) {
             if (!block[i]) primes.push_back((L + i) * 2 + 1);
     for (int p : primes) is_prime[p] = true;
     return primes;
      Combinatorics
6.12.1 Catalan numbers
                             C_n = \frac{1}{n+1} {2n \choose n} = \frac{(2n)!}{n!(n+1)!}
```

- $C_{n+1} = \sum_{i=0}^{n} C_i C_{n-i}, \ C_0 = 1, \ C_n = \frac{4n-2}{n+1} C_{n-1}$ Catalan numbers  $(n = 0, 1, 2, \dots, 11)$ :
- The first 12 Catalan numbers (n = 0, 1, 2, ..., 11):
  - $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786$

- Applications of Catalan numbers:
  - difference binary search trees with n vertices from 1 to n.
  - rooted binary trees with n + 1 leaves (vertices are not numbered).
  - correct bracket sequence of length 2 \* n.
  - permutation [n] with no 3-term increasing subsequence (i.e. doesn't exist i < j < k for which a[i] < a[j] < a[k]).
  - ways a convex polygon of n + 2 sides can split into triangles by connecting vertices.

#### 6.12.2 Fibonacci numbers

$$F_n = \begin{cases} 0, & \text{if } n = 0\\ 1, & \text{if } n = 1\\ F_{n-1} + F_{n-2}, & \text{otherwise} \end{cases}$$

• The first 20 Fibonacci numbers (n = 0, 1, 2, ..., 19):

$$F_n = 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181$$

• Properties:

$$F_{2n+1} = F_n^2 + F_{n+1}^2$$

$$F_{2n} = F_{n-1} \cdot F_n + F_n \cdot F_{n+1}$$

$$F_{n+1} \cdot F_{n-1} - F_n^2 = (-1)^n$$

$$n \mid m \iff F_n \mid F_m$$

$$(F_n, F_m) = F_{(n,m)}$$

#### 6.12.3 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* non-empty groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{i=1}^{k} (-1)^{k-i} {k \choose i} i^n$$

## 6.12.4 Derangements

Permutation of the elements of a set, such that no element appears in its original position (no fixied point). Recursive formulas:

$$D(n) = (n-1)[D(n-1) + D(n-2)] = nD(n-1) + (-1)^n$$

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# 7 Geometry

#### 7.1 Fundamentals

#### 7.1.1 **Point**

```
#pragma once
const double PI = acos(-1);
const double EPS = 1e-9;
typedef double ftype;
struct Point {
    ftype x, y;
    Point(ftype _x = 0, ftype _y = 0): x(_x), y(_y) {}
    Point& operator+=(const Point& other) {
        x += other.x; y += other.y; return *this;
    Point& operator -= (const Point& other) {
        x -= other.x; y -= other.y; return *this;
    Point& operator*=(ftype t) {
        x *= t; y *= t; return *this;
    Point& operator/=(ftype t) {
        x /= t; y /= t; return *this;
    Point operator+(const Point& other) const {
        return Point(*this) += other;
   }
    Point operator-(const Point& other) const {
        return Point(*this) -= other;
    Point operator*(ftype t) const {
        return Point(*this) *= t;
    Point operator/(ftype t) const {
        return Point(*this) /= t;
    Point rotate(double angle) const {
        return Point(x * cos(angle) - y * sin(angle), x * sin(angle) + y *
    cos(angle));
    friend istream& operator>>(istream &in, Point &t);
    friend ostream& operator<<(ostream &out, const Point& t);</pre>
    bool operator<(const Point& other) const {</pre>
        if (fabs(x - other.x) < EPS)</pre>
            return y < other.y;</pre>
        return x < other.x:</pre>
   }
};
istream& operator>>(istream &in, Point &t) {
    in >> t.x >> t.y;
    return in:
}
```

```
ostream& operator << (ostream &out, const Point& t) {
     out << t.x << ' ' << t.y;
     return out;
 ftype dot(Point a, Point b) {return a.x * b.x + a.y * b.y;}
 ftype norm(Point a) {return dot(a, a);}
 ftype abs(Point a) {return sqrt(norm(a));}
 ftype angle(Point a, Point b) {return acos(dot(a, b) / (abs(a) * abs(b)));}
 ftype proj(Point a, Point b) {return dot(a, b) / abs(b);}
 ftype cross(Point a, Point b) {return a.x * b.y - a.y * b.x;}
 bool ccw(Point a, Point b, Point c) {return cross(b - a, c - a) > EPS;}
 bool collinear(Point a, Point b, Point c) {return fabs(cross(b - a, c - a)) <</pre>
 Point intersect(Point a1, Point d1, Point a2, Point d2) {
     double t = cross(a2 - a1, d2) / cross(d1, d2);
     return a1 + d1 * t;
7.1.2 Line
 #include "point.h"
 struct Line {
     double a, b, c;
     Line (double _a = 0, double _b = 0, double _c = 0): a(_a), b(_b), c(_c) {}
     friend ostream & operator<<(ostream& out, const Line& 1);</pre>
 ostream & operator << (ostream & out. const Line & 1) {
     out << 1.a << ' ' << 1.b << ' ' << 1.c;
     return out;
 void PointsToLine(const Point& p1, const Point& p2, Line& 1) {
     if (fabs(p1.x - p2.x) < EPS)
         1 = \{1.0, 0.0, -p1.x\};
     else {
         1.a = - (double)(p1.y - p2.y) / (p1.x - p2.x);
         1.b = 1.0;
         1.c = -1.a * p1.x - 1.b * p1.y;
 void PointsSlopeToLine(const Point& p, double m, Line& 1) {
     1.a = -m:
     1.b = 1:
     1.c = -1.a * p.x - 1.b * p.y;
 bool areParallel(const Line& 11, const Line& 12) {
     return fabs(l1.a - l2.a) < EPS && fabs(l1.b - l2.b) < EPS;</pre>
 bool areSame(const Line& 11, const Line& 12) {
     return areParallel(11, 12) && fabs(11.c - 12.c) < EPS;</pre>
 bool areIntersect(Line 11, Line 12, Point& p) {
     if (areParallel(l1, l2)) return false;
     p.x = -(11.c * 12.b - 11.b * 12.c) / (11.a * 12.b - 11.b * 12.a);
```

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```
if (fabs(11.b) > EPS) p.y = -(11.c + 11.a * p.x);
     else p.y = -(12.c + 12.a * p.x);
     return 1;
 double distToLine(Point p, Point a, Point b, Point& c) {
     double t = dot(p - a, b - a) / norm(b - a);
     c = a + (b - a) * t;
     return abs(c - p);
 double distToSegment(Point p, Point a, Point b, Point& c) {
     double t = dot(p - a, b - a) / norm(b - a);
     if (t > 1.0)
         c = Point(b.x, b.y);
     else if (t < 0.0)
         c = Point(a.x, a.y);
     else
         c = a + (b - a) * t;
     return abs(c - p);
 bool intersectTwoSegment(Point a. Point b. Point c. Point d) {
     ftype ABxAC = cross(b - a, c - a);
     ftype ABxAD = cross(b - a, d - a);
     ftype CDxCA = cross(d - c, a - c);
     ftype CDxCB = cross(d - c, b - c);
     if (ABxAC == 0 \mid | ABxAD == 0 \mid | CDxCA == 0 \mid | CDxCB == 0) 
         if (ABxAC == 0 && dot(a - c, b - c) <= 0) return true;</pre>
         if (ABxAD == 0 && dot(a - d, b - d) <= 0) return true;
         if (CDxCA == 0 && dot(c - a, d - a) <= 0) return true;</pre>
         if (CDxCB == 0 &\& dot(c - b, d - b) <= 0) return true;
         return false;
     }
     return (ABxAC * ABxAD < 0 && CDxCA * CDxCB < 0);
 void perpendicular(Line 11, Point p, Line& 12) {
     if (fabs(l1.a) < EPS)</pre>
         12 = \{1.0, 0.0, -p.x\};
     else {
         12.a = -11.b / 11.a;
         12.b = 1.0;
         12.c = -12.a * p.x - 12.b * p.y;
 }
7.1.3 Circle
 #include "point.h"
 int insideCircle(const Point& p, const Point& center, ftype r) {
     ftype d = norm(p - center);
     ftype rSq = r * r;
     return fabs(d - rSq) < EPS ? 0 : (d - rSq >= EPS ? 1 : -1);
 bool circle2PointsR(const Point& p1, const Point& p2, ftype r, Point& c) {
     double h = r * r - norm(p1 - p2) / 4.0;
     if (fabs(h) < 0) return false;</pre>
```

```
h = sqrt(h);
     Point perp = (p2 - p1).rotate(PI / 2.0);
     Point m = (p1 + p2) / 2.0;
     c = m + perp * (h / abs(perp));
     return true:
 }
7.1.4 Triangle
 #include "point.h"
 #include "line.h"
 double areaTriangle(double ab, double bc, double ca) {
     double p = (ab + bc + ca) / 2;
     return sqrt(p) * sqrt(p - ab) * sqrt(p - bc) * sqrt(p - ca);
 double rInCircle(double ab, double bc, double ca) {
     double p = (ab + bc + ca) / 2;
     return areaTriangle(ab, bc, ca) / p;
 double rInCircle(Point a, Point b, Point c) {
     return rInCircle(abs(a - b), abs(b - c), abs(c - a));
 bool inCircle(Point p1, Point p2, Point p3, Point &ctr, double &r) {
     r = rInCircle(p1, p2, p3);
     if (fabs(r) < EPS) return false;</pre>
     Line 11, 12;
     double ratio = abs(p2 - p1) / abs(p3 - p1);
     Point p = p2 + (p3 - p2) * (ratio / (1 + ratio));
     PointsToLine(p1, p, l1);
     ratio = abs(p1 - p2) / abs(p2 - p3);
     p = p1 + (p3 - p1) * (ratio / (1 + ratio));
     PointsToLine(p2, p, 12);
     areIntersect(l1, l2, ctr);
     return true;
 double rCircumCircle(double ab, double bc, double ca) {
     return ab * bc * ca / (4.0 * areaTriangle(ab, bc, ca));
 double rCircumCircle(Point a, Point b, Point c) {
     return rCircumCircle(abs(b - a), abs(c - b), abs(a - c));
7.1.5 Convex hull
 #include "point.h"
 vector<Point> CH_Andrew(vector<Point> &Pts) { // overall O(n log n)
     int n = Pts.size(), k = 0;
     vector<Point> H(2 * n);
     sort(Pts.begin(), Pts.end());
     for (int i = 0; i < n; ++i) {
         while ((k \ge 2) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
         H[k++] = Pts[i];
     for (int i = n - 2, t = k + 1; i >= 0; --i) {
```

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```
while ((k \ge t) \&\& !ccw(H[k - 2], H[k - 1], Pts[i])) --k;
         H[k++] = Pts[i];
     H.resize(k);
     return H;
 }
7.1.6 Polygon
 #include "point.h"
 double perimeter(const vector<Point> &P) {
     double ans = 0.0;
     for (int i = 0; i < (int)P.size() - 1; ++i)
         ans += abs(P[i] - P[i + 1]);
     return ans:
 double area(const vector<Point> &P) {
     double ans = 0.0;
     for (int i = 0; i < (int)P.size() - 1; ++i)</pre>
         ans += (P[i].x * P[i + 1].y - P[i + 1].x * P[i].y);
     return fabs(ans) / 2.0;
 bool isConvex(const vector<Point> &P) {
     int n = (int)P.size();
     if (n <= 3) return false;</pre>
     bool firstTurn = ccw(P[0], P[1], P[2]);
     for (int i = 1; i < n - 1; ++i)
         if (ccw(P[i], P[i + 1], P[(i + 2) == n ? 1 : i + 2]) != firstTurn)
             return false;
     return true;
 }
 int insidePolygon(Point pt, const vector<Point> &P) {
     int n = (int)P.size();
     if (n <= 3) return -1;
     bool on_polygon = false;
     for (int i = 0; i < n - 1; ++i)
         if (fabs(abs(P[i] - pt) + abs(pt - P[i + 1]) - abs(P[i] - P[i + 1])) <
     EPS)
             on_polygon = true;
     if (on_polygon) return 0;
     double sum = 0.0;
     for (int i = 0; i < n - 1; ++i) {
         if (ccw(pt, P[i], P[i + 1]))
             sum += angle(P[i] - pt, P[i + 1] - pt);
         else
             sum -= angle(P[i] - pt, P[i + 1] - pt);
     }
     return fabs(sum) > PI ? 1 : -1;
7.2 Minimum enclosing circle
 /**
  * Description: computes the minimum Circle that encloses all the given Points.
```

```
*/
```

```
#include "point.h"
// #include "circle.h"
// TODO:
Point center_from(double bx, double by, double cx, double cy) {
    double B = bx * bx + by * by, C = cx * cx + cy * cy, D = bx * cy - by * cx;
    return Point((cy * B - by * C) / (2 * D), (bx * C - cx * B) / (2 * D));
Circle Circle_from(Point A, Point B, Point C) {
    Point I = center_from(B.x - A.x, B.y - A.y, C.x - A.x, C.y - A.y);
    return Circle(I + A, abs(I));
}
const int N = 100005;
int n, x[N], y[N];
Point a[N];
Circle emo_welzl(int n, vector<Point> T) {
    if (T.size() == 3 || n == 0) {
        if (T.size() == 0) return Circle(Point(0, 0), -1);
        if (T.size() == 1) return Circle(T[0], 0);
        if (T.size() == 2) return Circle((T[0] + T[1]) / 2, abs(T[0] - T[1]) / 2
        return Circle_from(T[0], T[1], T[2]);
    random\_shuffle(a + 1, a + n + 1);
    Circle Result = emo_welzl(0, T);
    for (int i = 1; i <= n; i++)
        if (abs(Result.x - a[i]) > Result.y + 1e-9) {
            T.push_back(a[i]);
            Result = emo_welzl(i - 1, T);
            T.pop_back();
        }
    return Result:
```

# Linear algebra

## 8.1 Gauss elimination

```
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be infinity or a big number
int gauss (vector < vector < double > > a, vector < double > & ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
        int sel = row:
        for (int i=row; i<n; ++i) {</pre>
            if (abs (a[i][col]) > abs (a[sel][col])) sel = i;
        if (abs (a[sel][col]) < EPS) continue;</pre>
```

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```
for (int i=col; i<=m; ++i) {</pre>
        swap (a[sel][i], a[row][i]);
    where[col] = row;
    for (int i=0; i<n; ++i) {</pre>
        if (i != row) {
             double c = a[i][col] / a[row][col];
             for (int j=col; j<=m; ++j) {</pre>
                 a[i][j] -= a[row][j] * c;
             }
    }
    ++row;
ans.assign (m, 0);
for (int i=0; i<m; ++i) {</pre>
    if (where[i] != -1) {
        ans[i] = a[where[i]][m] / a[where[i]][i];
for (int i=0; i<n; ++i) {</pre>
    double sum = 0;
    for (int j=0; j<m; ++j) {
        sum += ans[j] * a[i][j];
    if (abs (sum - a[i][m]) > EPS) return 0;
for (int i=0; i<m; ++i) {
    if (where[i] == -1) return INF;
}
return 1;
```

# 9 Graph

}

# 9.1 Bellman-Ford algorithm

```
/**
 * Description: single source shortest path in a weighted (negative or positive) directed graph.
 * Time: O(N * M).
 * Tested: https://open.kattis.com/problems/shortestpath3
 */
const int64_t INF = (int64_t) 2e18;
struct Edge {
   int u, v; // u -> v
   int64_t w;
   Edge() {}
   Edge(int _u, int _v, int64_t _w) : u(_u), v(_v), w(_w) {}
};
int n;
vector<Edge> edges;
vector<int64_t> bellmanFord(int s) {
```

```
// dist[stating] = 0.
   // dist[u] = +INF, if u is unreachable.
   // dist[u] = -INF, if there is a negative cycle on the path from s to u.
   // -INF < dist[u] < +INF, otherwise.</pre>
   vector<int64_t> dist(n, INF);
   dist[s] = 0;
   for (int i = 0; i < n - 1; ++i) {
       bool any = false;
       for (auto [u, v, w] : edges) {
            if (dist[u] != INF && dist[v] > w + dist[u]) {
                dist[v] = w + dist[u];
                any = true;
           }
       if (!any) break;
   // handle negative cycles
   for (int i = 0; i < n - 1; ++i) {
       for (auto [u, v, w] : edges) {
           if (dist[u] != INF && dist[v] > w + dist[u]) {
                dist[v] = -INF;
           }
   return dist;
   Articulation point and Bridge
* Description: finding articulation points and bridges in a simple undirected
* Tested: https://oj.vnoi.info/problem/graph_
const int N = (int) 1e5;
vector<int> g[N];
int num[N], low[N], dfs_timer;
bool joint[N];
vector<pair<int, int>> bridges;
void dfs(int u, int prev) {
   low[u] = num[u] = ++dfs_timer;
   int child = 0;
   for (int v : g[u]) {
       if (v == prev) continue;
       if (num[v]) low[u] = min(low[u], num[v]);
       else {
            dfs(v, u);
            low[u] = min(low[u], low[v]);
            child++:
            if (low[v] >= num[v]) {
                bridges.emplace_back(u, v);
            if (u != prev && low[v] >= num[u]) joint[u] = true;
       }
```

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```
if (u == prev && child > 1) joint[u] = true;
}
int solve() {
   int n. m:
    cin >> n >> m;
    for (int i = 0; i < m; ++i) {
        int u, v;
        cin >> u >> v;
        u--; v--;
        g[u].push_back(v);
        g[v].push_back(u);
    for (int i = 0; i < n; ++i) {</pre>
        if (!num[i]) dfs(i, i);
   }
    return 0;
    Strongly connected components
 * Description: Tarjan's algorithm finds strongly connected components
 * in a directed graph. If vertices u and v belong to the same component,
 * then scc_id[u] == scc_id[v].
* Tested: https://judge.yosupo.jp/problem/scc
const int N = (int) 5e5;
vector<int> g[N], st;
int low[N], num[N], dfs_timer, scc_id[N], scc;
bool used[N];
void Tarjan(int u) {
    low[u] = num[u] = ++dfs_timer;
    st.push_back(u);
    for (int v : g[u]) {
        if (used[v]) continue;
        if (num[v] == 0) {
            Tarjan(v);
            low[u] = min(low[u], low[v]);
        else {
            low[u] = min(low[u], num[v]);
   if (low[u] == num[u]) {
        int v;
        do {
            v = st.back(); st.pop_back();
            debug(u, v)
            used[v] = true;
```

scc\_id[v] = scc;

} while (v != u);

scc++;

```
}
9.4 Topo sort
  * Description: A topological sort of a directed acyclic graph
  * is a linear ordering of its vertices such that for every directed edge
  * from vertex u to vertex v, u comes before v in the ordering.
  * Note: If there are cycles, the returned list will have size smaller than n
      (i.e, topo.size() < n).
  * Tested: https://judge.yosupo.jp/problem/scc
 vector<int> topo_sort(const vector<vector<int>> &g) {
     int n = (int) g.size();
     vector<int> indeg(n);
     for (int u = 0; u < n; ++u) {
         for (int v : q[u]) indeq[v]++;
     queue<int> q; // Note: use min-heap to get the smallest lexicographical
     for (int u = 0; u < n; ++u) {
         if (indeg[u] == 0) q.emplace(u);
     vector<int> topo:
     while (!q.empty()) {
         int u = q.front(); q.pop();
         topo.emplace_back(u);
         for (int v : q[u]) {
             if (--indeg[v] == 0) q.emplace(v);
     return topo;
9.5 K-th smallest shortest path
 /** Description: Finding the k-th smallest shortest path from vertex s to
  * each vertex can be visited more than once.
 using adj_list = vector<vector<pair<int, int>>>;
 vector<long long> k_smallest(const adj_list &g, int k, int s, int t) {
     int n = (int) g.size();
     vector<long long> ans;
     vector<int> cnt(n);
     using pli = pair<long long, int>;
     priority_queue<pli, vector<pli>, greater<pli>> pq;
     pq.emplace(0, s);
     while (!pq.empty() && cnt[t] < k) {</pre>
         int u = pq.top().second;
         long long d = pq.top().first;
         pq.pop();
         if (cnt[u] == k) continue;
         cnt[u]++;
```

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```
if (u == t) {
             ans.push_back(d);
         for (auto [v, cost] : g[u]) {
             pq.emplace(d + cost, v);
         }
     assert(k == (int) ans.size());
     return ans;
 }
9.6 Eulerian path
9.6.1 Directed graph
 /**
  * Hierholzer's algorithm.
  * Description: An Eulerian path in a directed graph is a path that visits all
     edges exactly once.
    An Eulerian cycle is a Eulerian path that is a cycle.
  * Time complexity: O(|E|).
  vector<int> find_path_directed(const vector<vector<int>> &g, int s) {
     int n = (int) g.size();
     vector<int> stack, cur_edge(n), vertices;
     stack.push_back(s);
     while (!stack.empty()) {
         int u = stack.back();
         stack.pop_back();
         while (cur_edge[u] < (int) g[u].size()) {</pre>
             stack.push_back(u);
             u = g[u][cur_edge[u]++];
         vertices.push_back(u);
     reverse(vertices.begin(), vertices.end());
     return vertices;
 }
9.6.2 Undirected graph
 /**
  * Hierholzer's algorithm.
  * Description: An Eulerian path in a undirected graph is a path that visits
     all edges exactly once.
  * An Eulerian cycle is a Eulerian path that is a cycle.
  * Time complexity: O(|E|).
 struct Edge {
     int to:
     list<Edge>::iterator reverse_edge;
     Edge(int _to) : to(_to) {}
 vector<int> vertices;
 void find_path(vector<list<Edge>> &g, int u) {
     while (!g[u].empty()) {
```

```
int v = g[u].front().to;
         g[v].erase(g[u].front().reverse_edge);
         g[u].pop_front();
         find_path(g, v);
     vertices.emplace_back(u); // reversion list.
 void add_edge(vector<list<Edge>> &g, int u, int v) {
     g[u].emplace_front(v);
     g[v].emplace_front(u);
     g[u].front().reverse_edge = g[v].begin();
     g[v].front().reverse_edge = g[u].begin();
}
9.7 HLD
 const int INF = 0x3f3f3f3f3f:
 template < class SegmentTree >
 struct HLD { // vertex update and max query on path u -> v
     int n;
     vector<vector<int>> g;
     SegmentTree seg_tree;
     vector<int> par, top, depth, sz, id;
     int timer = 0;
     bool VAL_IN_EDGE = false;
     HLD() {}
     HLD(int_n): n(n), g(n), seg\_tree(n), par(n), top(n), depth(n), sz(n),
     id(n) {}
     void build() {
         dfs_sz(0);
         dfs_hld(0);
     void add_edge(int u, int v) {
         g[u].push_back(v);
         g[v].push_back(u);
     }
     void dfs_sz(int u) {
         sz[u] = 1;
         for (int &v : g[u]) { // MUST BE ref for the swap below
             par[v] = u;
             depth[v] = depth[u] + 1;
             g[v].erase(find(g[v].begin(), g[v].end(), u));
             dfs_sz(v);
             sz[u] += sz[v];
             if (sz[v] > sz[g[u][0]]) swap(v, g[u][0]);
         }
     void dfs_hld(int u) {
         id[u] = timer++;
         for (int v : g[u]) {
             top[v] = (v == g[u][0] ? top[u] : v);
             dfs_hld(v);
         }
     }
```

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```
int lca(int u, int v) {
        while (top[u] != top[v]) {
            if (depth[top[u]] > depth[top[v]]) swap(u, v);
            v = par[top[v]];
        // now u, v is in the same heavy-chain
        return (depth[u] < depth[v] ? u : v);</pre>
    void set_vertex(int v, int x) {
        seg_tree.set(id[v], x);
    void set_edge(int u, int v, int x) {
        if (u != par[v]) swap(u, v);
        seg_tree.set(id[v], x);
   }
    void set_subtree(int v, int x) {
        // modify segment_tree so that it supports range update
        seg_tree.set_range(id[v] + VAL_IN_EDGE, id[v] + sz[v] - 1, x);
    int querv path(int u. int v) {
        int res = -INF;
        while (top[u] != top[v]) {
            if (depth[top[u]] > depth[top[v]]) swap(u, v);
            int cur = seg_tree.query(id[top[v]], id[v]);
            res = max(res, cur);
            v = par[top[v]];
        if (depth[u] > depth[v]) swap(u, v);
        int cur = seg_tree.query(id[u] + VAL_IN_EDGE, id[v]);
        res = max(res, cur);
        return res;
   }
};
```

## 10 Misc.

## 10.1 Ternary search

```
/**
 * Description: given an unimodal function f(x), find the maximum/minimum of
   f(x).
 * Unimodal means The function strictly increases/decreases first,
 * reaches a maximum/minimum (at a single point or over an interval),
 * and then strictly decreases/increases.
 */
const double eps = 1e-9;
template < typename T >
inline T func(T x) { return x * x; }

// these two functions below find min, for find max: change '<' below to '>'.
double ternary_search(double 1, double r) { // min
   while (r - 1 > eps) {
      double mid1 = 1 + (r - 1) / 3;
      double mid2 = r - (r - 1) / 3;
}
```

```
if (func(mid1) < func(mid2)) r = mid2;</pre>
         else 1 = mid1;
     }
     return 1;
 int ternary_search(int 1, int r) { // min
     while (1 < r) {
         int mid = 1 + (r - 1) / 2;
         if (func(mid) < func(mid + 1)) r = mid;</pre>
         else l = mid + 1;
     return 1;
10.2 Matrix
 using matrix_type = int;
 const int MOD = (int) 1e9 + 7;
 struct Matrix {
     static const matrix_type INF = numeric_limits<matrix_type>::max();
     vector<vector<matrix_type>> mat;
     Matrix(int _N, int _M, matrix_type v = 0) : N(_N), M(_M) {
         mat.assign(N. vector<matrix type>(M. v)):
     static Matrix identity(int n) { // return identity matrix.
         Matrix I(n, n);
         for (int i = 0; i < n; ++i) {
             I[i][i] = 1;
         return I;
     }
     vector<matrix_type>& operator[](int r) { return mat[r]; }
     const vector<matrix_type>& operator[](int r) const { return mat[r]; }
     Matrix& operator*=(const Matrix &other) {
         assert(M == other.N); // [N x M] [other.N x other.M]
         Matrix res(N, other.M);
         for (int r = 0; r < N; ++r) {
             for (int c = 0; c < other.M; ++c) {
                  long long square_mod = (long long) MOD * MOD;
                 long long sum = 0;
                 for (int g = 0; g < M; ++g) {
                     sum += (long long) mat[r][g] * other[g][c];
                     if (sum >= square_mod) sum -= square_mod;
                 res[r][c] = sum % MOD;
             }
         }
         mat.swap(res.mat); return *this;
 };
```