

CS 161

HW 1

Xia Hua

86211 8335

1.

ISA	IPC	clock
A	10	500 MHz
B	2	600 MHz

$$A. \text{ MIPS} = \frac{\text{Clock Rate}}{\text{CPI} \times 10^6}$$

$$\text{CPI} = \frac{\text{cycles}}{\text{instruction}} = \frac{1}{\text{IPC}}$$

$$\text{IPC} = \frac{1}{\text{CPI}}$$

$$(A) \text{ MIPS} = \frac{\text{clock}}{\frac{1}{\text{IPC}} \times 10^6}$$

$$= \frac{500 \text{ MHz}}{\frac{1}{10} \times 10^6} = \frac{500 \times 10^6}{1 \times 10^5} = 5000$$

$$b. (B) \text{ MIPS} = \frac{600 \text{ MHz}}{\frac{1}{2} \times 10^6} = \frac{600 \times 10^6}{0.5 \times 10^6} = 1200$$

C. Unable to define, Because A and B are using different ISA instruction set, Even though we can have A and B's MIPS.

2.

Processor	CPI	clock Rate
P1	1.5	3.0 GHz
P2	1.0	2.5 GHz
P3	2.2	4.0 GHz

a. Calculate IPS.

$$IPS = \frac{\text{clock} \times 10^9}{CPI}$$

$$P1: \frac{3.0 \times 10^9}{1.5} = 2 \times 10^9$$

$$P2: \frac{2.5 \times 10^9}{1.0} = 2.5 \times 10^9$$

$$P3: \frac{4.0 \times 10^9}{2.2} = 1.8 \times 10^9$$

b. number of cycles | number of instructions (in 10 seconds)

$$P1: 3 \times 10^9 \times 10 = 3 \times 10^{10}$$

$$P2: 2.5 \times 10^9 \times 10 = 2.5 \times 10^{10}$$

$$P3: 4.0 \times 10^9 \times 10 = 4 \times 10^{10}$$

$$P1: 2 \times 10^9 \times 10 = 2 \times 10^{10}$$

$$P2: 2.5 \times 10^9 \times 10 = 2.5 \times 10^{10}$$

$$P3: 1.8 \times 10^9 \times 10 = 1.8 \times 10^{10}$$

c. execution time = $\frac{\text{Instruction count} \times CPI}{\text{clock Rate}}$

$$\text{clock Rate} = \frac{\text{Instruction count} \times CPI (120\%)}{\text{execution time} (70\%)}$$

$$\frac{120\%}{70\%} = 1.714$$

increase of 71.4% on clock Rate

3.

Processor	CPI				clock Rate
	A	B	C	D	
P1	1	2	3	3	2.5 GHz
P2	2	2	2	2	3.0 GHz

case 1

case 2.

Weighted average CPI

$$CPI = \frac{\text{clock Cycle}}{\text{Instruction Count}} = \sum_{i=1}^n \left(CPI_i \times \frac{\text{Instruction Count}_i}{\text{Instruction Count}} \right)$$

$$\begin{aligned} \text{a. P1: } & 1 \times \frac{1}{10} + 2 \times \frac{2}{10} + 3 \times \frac{5}{10} + 3 \times \frac{2}{10} \\ & = 0.1 + 0.4 + 1.5 + 0.6 = 2.6 \end{aligned}$$

CPI (global)

$$\begin{aligned} \text{P2: } & 2 \times \frac{1}{10} + 2 \times \frac{2}{10} + 2 \times \frac{5}{10} + 2 \times \frac{2}{10} \\ & = 0.2 + 0.4 + 1 + 0.4 = 2 \end{aligned}$$

CPI (global)

$$\text{b. clock Cycle} = \text{Instruction Count} \cdot CPI$$

$$\text{P1: } 1 \times 10^6 \cdot 2.6 = 2.6 \times 10^6$$

$$\text{P2: } 1 \times 10^6 \cdot 2 = 2 \times 10^6$$

4.

Compiler	instruction count	execution time
A	1×10^9	1.1s
B	1.2×10^9	1.3s

a. $1ns = 1 \times 10^{-9} \cdot s$

$$CPU \text{ time} = \text{Execution time} = \frac{\text{Instruction Count} \times CPI}{\text{clock rate}}$$

$$\text{Execution time} = \text{Instruction count} \times CPI \times \text{clock cycle time}$$

$$(AVG) CPI = \frac{\text{Execution time}}{\text{Instruction Count} \times \text{clock Cycle Time}}$$

A: $\frac{1.1s}{1 \times 10^9 \cdot 1 \times 10^{-9}s} = 1.1$ B: $\frac{1.3s}{1.2 \times 10^9 \cdot 1 \times 10^{-9}s} = 1.25$

b. Unknown variable: execution time.

$$\text{clock cycle time} = \frac{\text{Execution time}}{\text{Instruction Count} \cdot CPI} =$$

$$\frac{A}{B} = \frac{\frac{\text{Execution time A}}{\text{Instruction count A} \cdot CPI_A}}{\frac{\text{Execution time B}}{\text{Instruction count B} \cdot CPI_B}} = \frac{\text{Execution time A}}{\text{Execution time B}} \cdot \frac{\text{Instruction count B} \cdot CPI_B}{\text{Instruction count A} \cdot CPI_A}$$

result is unknown

due to unknown execution time

$$= \frac{\text{Execution time A}}{\text{Execution time B}} \cdot \frac{1.2 \times 10^9 \cdot 1.25}{1 \times 10^9 \cdot 1.1}$$

$$= 1.36 \frac{\text{Execution time A}}{\text{Execution time A}}$$

c. instruction count $\cdot CPI = 6 \times 10^8 \cdot 1.1 = 6.6 \times 10^8$

$$\frac{\text{new compiler}}{A} = \frac{6.6 \times 10^8}{1.1 \times 10^8} = 0.6$$

$$\frac{1}{0.6} = 1.67 \text{ performance increase}$$

$$\frac{\text{new compiler}}{B} = \frac{6.6 \times 10^8}{1.2 \times 10^8 \cdot 1.25} = 0.44$$

$$\frac{1}{0.44} = 2.27 \text{ increase}$$

5.

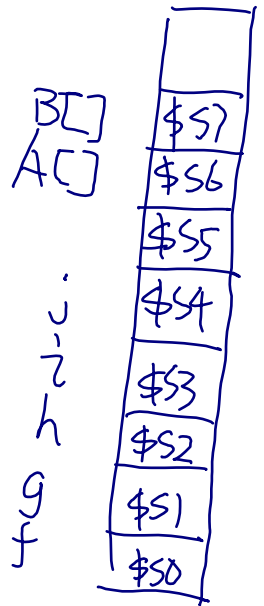
$$\text{Speed up}_{\text{total}} = \frac{1}{\left(1 - \text{Fraction}_{\text{enhanced1}} - \text{Fraction}_{\text{enhanced2}}\right) + \frac{\text{Fraction}_{\text{enhanced1}}}{\text{Speedup}_{\text{enhanced1}}} + \frac{\text{Fraction}_{\text{enhanced2}}}{\text{Speedup}_{\text{enhanced2}}}}$$

$$= \frac{1}{(1 - 0.3 - 0.8 - 0.3 - 0.2 - 0.5) + \frac{0.3 \cdot 0.8}{4} + \frac{0.3 \cdot 0.2 \cdot 0.5}{2}}$$

$$= \frac{1}{0.805} = 1.24$$

6.

sll \$t0, \$s0, 2	# \$t0 = f * 4
= add \$t0, \$s6, \$t0	# \$t0 = &A[f]
sll \$t1, \$s1, 2	# \$t1 = g * 4
= add \$t1, \$s7, \$t1	# \$t1 = &B[g] ☆
lw \$s0, 0(\$t0)	# f = A[f]
= addi \$t2, \$t0, 4	# \$t2 = &A[f+1]
lw \$t0, 0(\$t2)	# \$t0 = A[f+1]
= add \$t0, \$t0, \$s0	# \$t0 = A[f+1] + A[f]
sw \$t0, 0(\$t1)	# B[g] = A[f+1] + A[f]



f = A[f];

f = A[f+1] + A[f]

B[g] = f

7. $B[z] = A[i] + A[j]$;

```

sll $t0, $s3, 2      # $t0 = i * 4
add $t0, $t0, $s6     # $t0 = &A[i]
lw  $t0, 0($t0)       # $t0 = A[i]

sll $t1, $s4, 2      # $t1 = j * 4
add $t1, $t1, $s6     # $t1 = &A[j]
lw  $t1, 0($t1)       # $t1 = A[j]

add $t2, $t0, $t1     # $t2 = A[i] + A[j]
sw  $t2, 32($s7)      # B[z] = A[i] + A[j]
      ↑
    (4 * 8)
  
```

B[]	\$s7
A[]	\$s6
j	\$s4
i	\$s3
h	\$s2
g	\$s1
f	\$s0

8.

addi	\$t0, \$s6, 4	# \$t0 = &A[1]
add	\$t1, \$s6, \$0	# \$t1 = &A[0]
sw	\$t1, 0(\$t0)	# A[0] = A[1]
lw	\$t0, 0(\$t0)	# \$t0 = A[0]
add	\$s0, \$t1, \$t0	# f = A[0] + A[1]

$A[1] = A[0];$

$f = A[0] + A[1]$

9, t0 has 0x00101000

```
slt    $t2, $0, $t0      #  
bne    $t2, $0, ELSE     #  
j      DONE
```

```
ELSE:  
    addi $t2, $t2, 2      #
```

DONE:

① Compare \$0 and \$t0:

if $\$0 < \$t0$

$\$t2 = 1$; ✓ \Rightarrow next ②

else

$\$t2 = 0$;

② Not equal

if ($\$t2 \neq \0)

{ ELSE; } ✓ \Rightarrow next ③

③ Adding 2 to \$t2 .

$\$t2 + 2 = 3$.

$\$t2 = 3$.