

Lecture 9: Scaling Latency

<https://web3.princeton.edu/principles-of-blockchains/>

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This lecture:
Improving latency of Bitcoin

Three Modules

- Bitcoin (lectures 2-7)
- Scaling Bitcoin (lectures 8-14)
- Beyond Bitcoin (lectures 15-20)

Scaling Bitcoin

- Scaling Bitcoin (lecture 8-14)
- L1 Scaling: Improve Bitcoin performance while still retain basic structure of the longest chain protocol
 - Throughput (#8)
 - Latency (#9)
 - Storage & compute (#10) – Sharding
 - Energy (#11) – Proof of Stake
- L2 Scaling: Improve performance via an “overlay” on Bitcoin
 - Payment Channels (#12)
 - Data Availability (#13)
 - Rollups: Optimistic and Cryptographic proofs of compressed ledgers (#14)

Bitcoin latency

Time from when a transaction was broadcast until the transaction is confirmed in the ledger

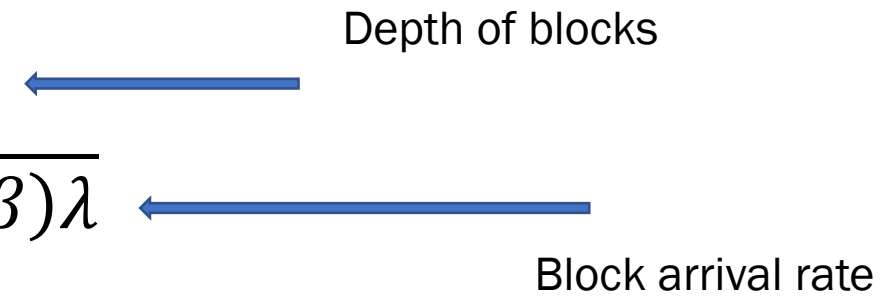
- τ_1 : Time from when a transaction was broadcast until the transaction is put into a mined block B
- τ_2 : Time from when the transaction was put into a mined block B until block B is k -deep in the longest chain

$$\tau = \tau_1 + \tau_2$$

τ_2 is the real bottleneck, depends on how large k is.

Bitcoin latency

Assume low forking ($\lambda\Delta \ll 1$),

$$\tau = \frac{k}{(1 - \beta)\lambda}$$


Depth of blocks

Block arrival rate

From Lecture 6, error probability

$$\epsilon = e^{-ck}$$

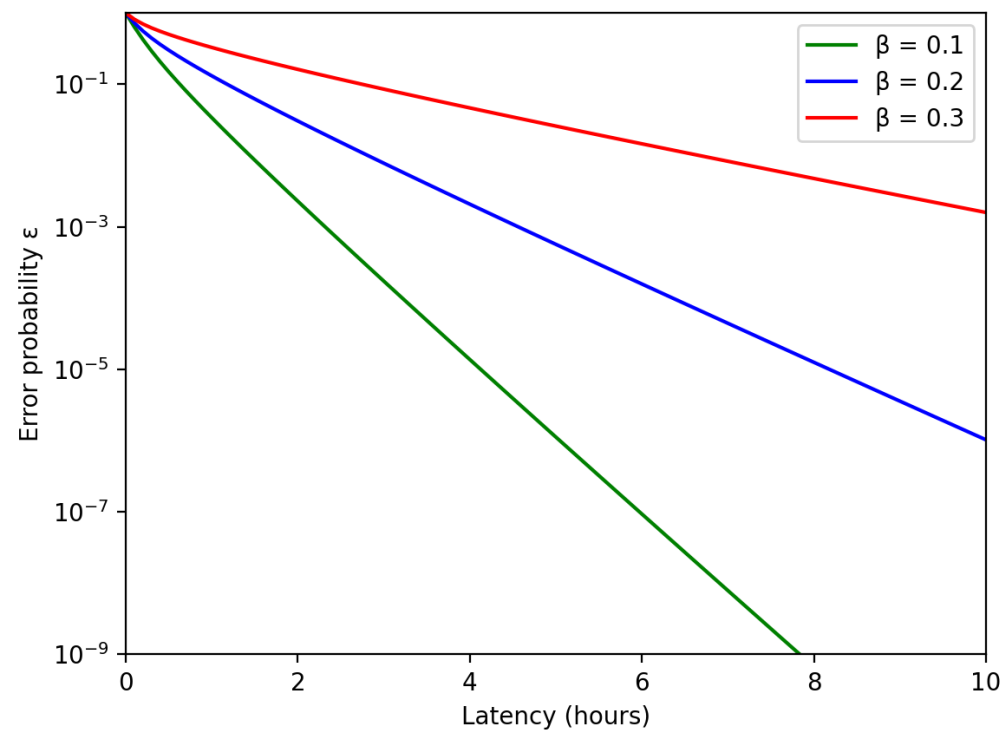
$$\tau = \frac{\frac{1}{c} \log(\frac{1}{\epsilon})}{(1 - \beta)\lambda} = O\left(\frac{1}{\lambda} \log\left(\frac{1}{\epsilon}\right)\right)$$

Latency and security are coupled

Bitcoin latency

$$\tau = O\left(\frac{1}{\lambda} \log\left(\frac{1}{\epsilon}\right)\right)$$

Bitcoin: $\frac{1}{\lambda} = 10$ minutes



Improve Bitcoin latency

Only way to improve latency is to

- reduce k ; but this reduces security
- Increase λ ; but this also reduces security

Ethereum: $\frac{1}{\lambda} = 15\text{s}; k = 100$

- latency = 25 minutes
- Way better than Bitcoin performance; improvement simply by picking better parameters.

Improve Bitcoin latency

Question: can we make relatively small changes to the longest chain protocol and PoW mining while scaling latency?

Key Requirement:

- Do not want latency to depend on security level
- **Decouple security from latency**

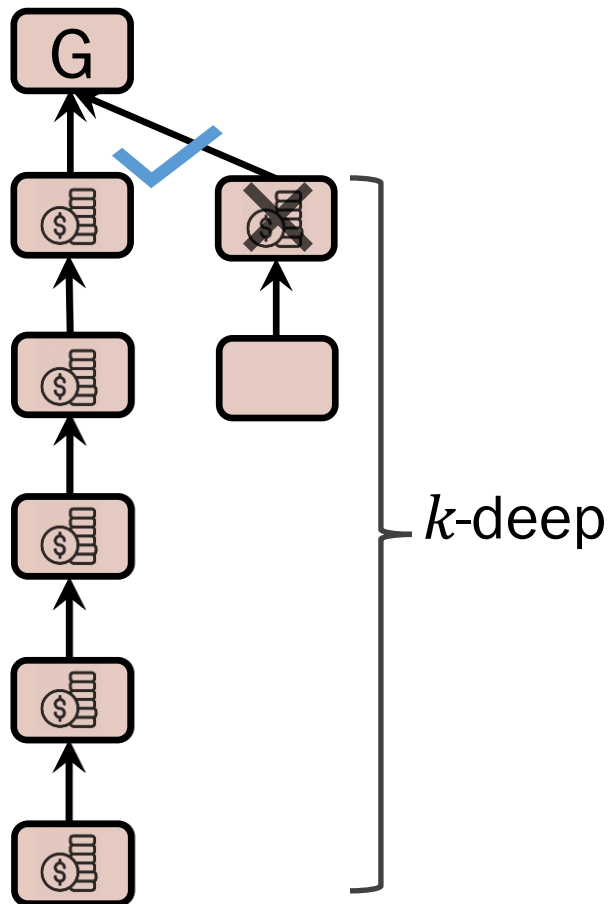
Prism

Prism achieves optimal latency

- **Decoupling principle**: separate performance from security
- Prism 1.0 achieves optimal throughput; last lecture

Decoupling voting

k-deep confirmation rule is a form of voting



Satoshi's Table

$q=0.3$	
$z=0$	$P=1.0000000$
$z=5$	$P=0.1773523$
$z=10$	$P=0.0416605$
$z=15$	$P=0.0101008$
$z=20$	$P=0.0024804$
$z=25$	$P=0.0006132$

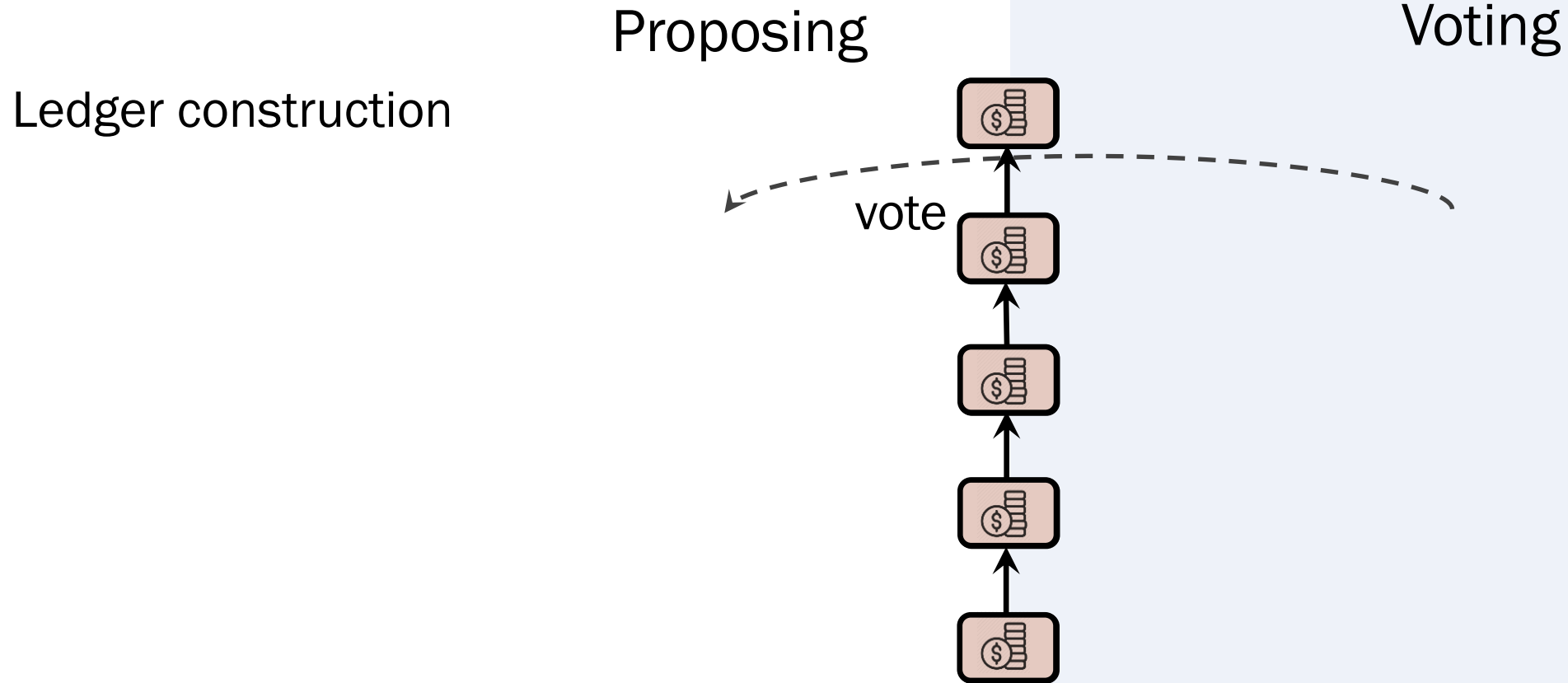
~~1 deep~~ => .45
25 deep => 0.0006

Can think of one block =
one vote underneath B

k-deep = k votes in
sequence

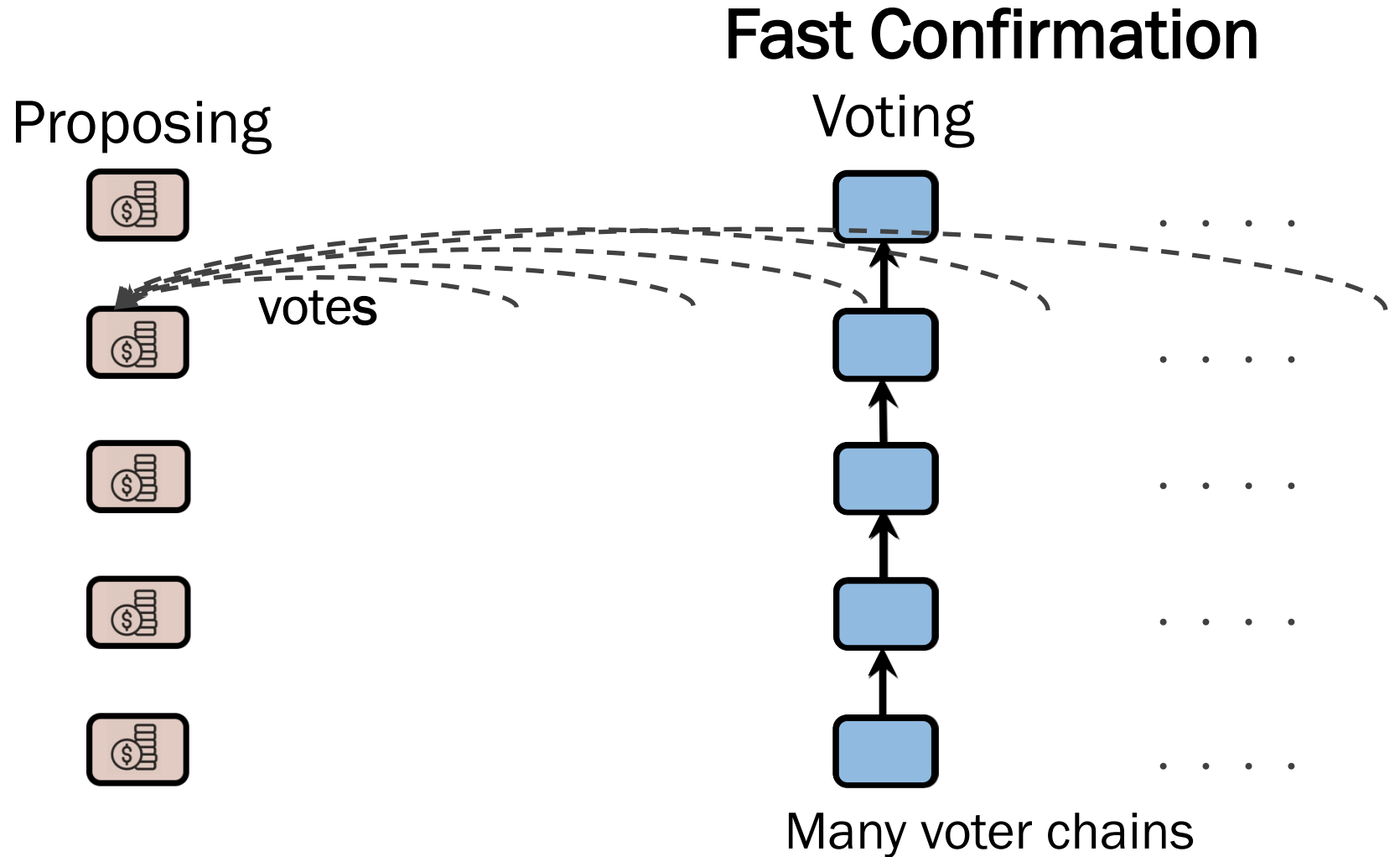
Really need k large to
sample the miners

Bitcoin → Deconstruct



1. Select **votes** along **longest** voter chain
2. Order the proposer blocks by votes

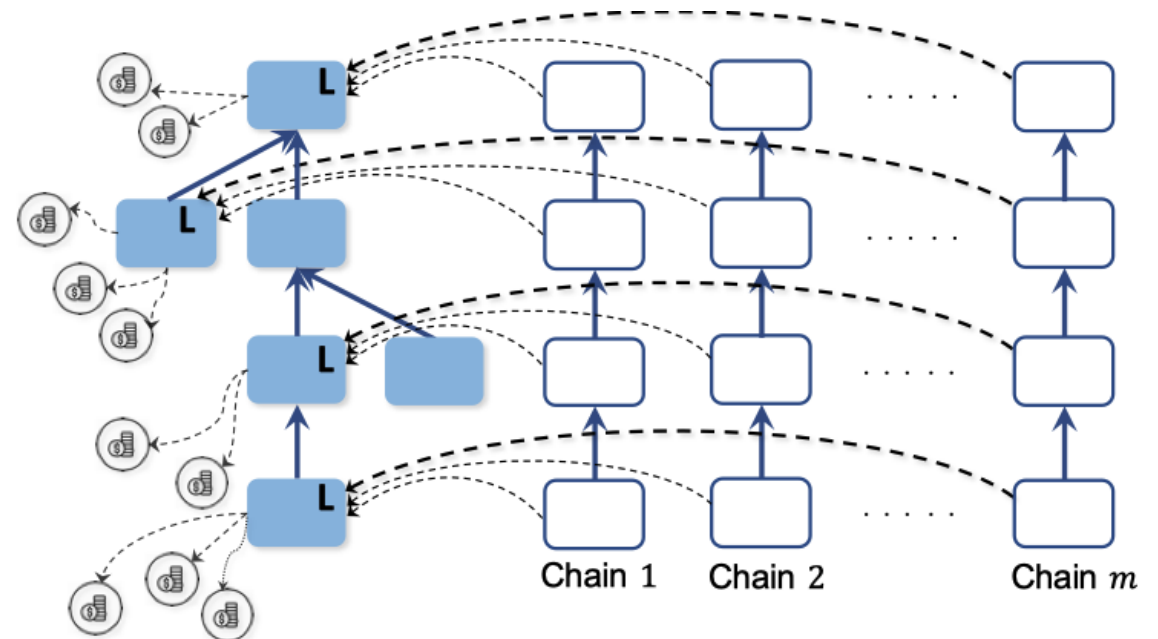
Bitcoin → Deconstruct → Prism



Ledger Construction: For each level choose the proposer block with **maximum** votes

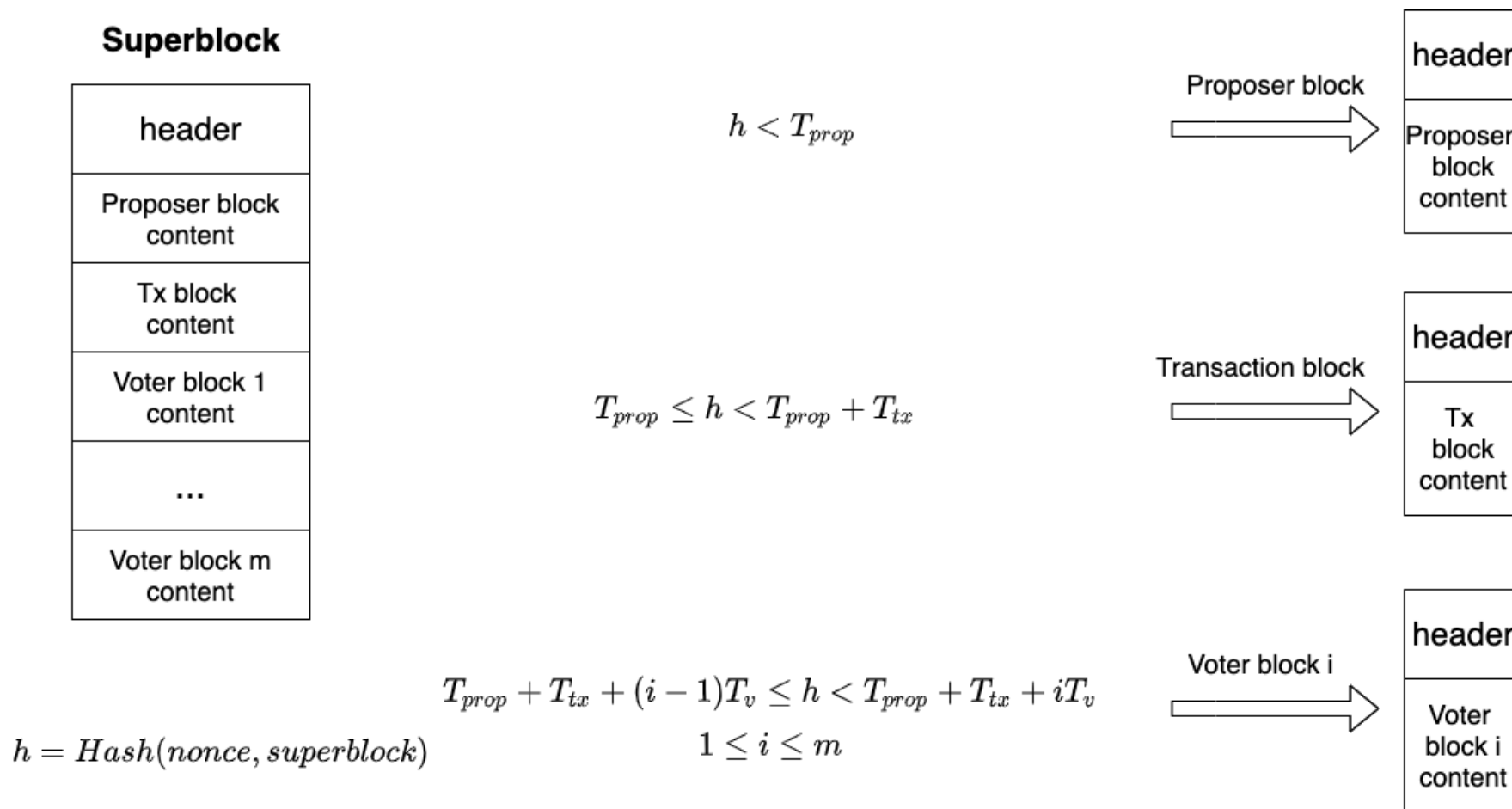
Prism

- Proposal rule: longest chain
- Voting rule:
 - a) each voter chain votes for one and only one proposer block at each level
 - b) each voter block votes for all the proposer levels that have not been voted by its parent.
- Mining rule: honest miner picks to be proposer/voter/transaction block at random



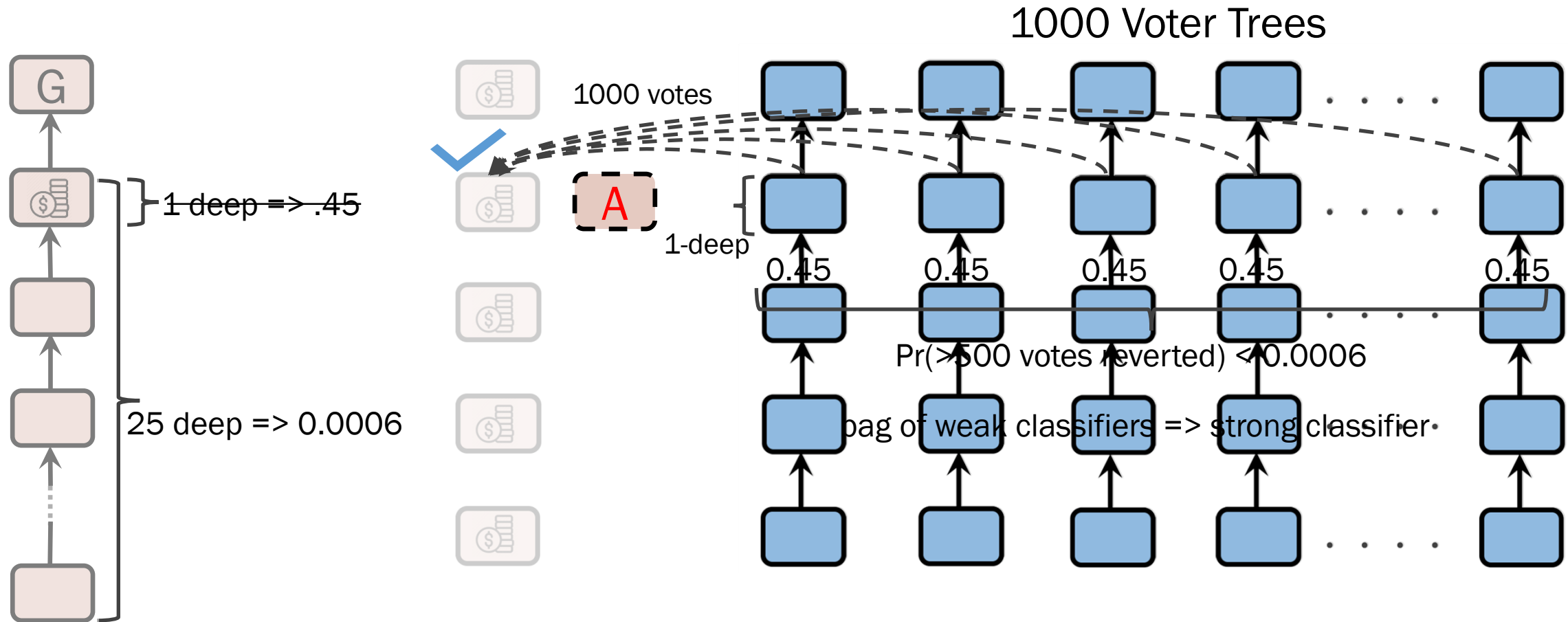
Cryptographic sortition

How do you prevent adversary from focusing its mining power on a specific type of blocks or on a specific voter chain?



Fast confirmation

Bitcoin



Ledger Construction: For each level choose the proposer block with **maximum** votes