

WCET Analysis Lab: Assignment 1

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SS 2012

Problem 1

Recommended (3): As a warm-up exercise, follow the instruction in Timing Analysis Lab: First Steps

Q: How long does it take to execute simple once, according to measurements, and according to the static analysis?

Answer

- Measurements: 763 incl. function call - 70 overhead = 693 cycles
- Static analysis: 705 cycles

We used the compiler flag -O1.

The difference is 12 cycles so the measurement is not far away from the static analysis. After one run we measured 1018 cycles but after reflashing the target the result was 763 cycles again.

Problem 2

Recommended (3): Also extract the instruction trace as outlined in Timing Analysis Lab: First Steps. Then compare the number of cycles needed in one iteration of the loop, with the number of cycles aiT calculated.

Q: Do they coincide? What is the total number of cycles needed to execute simple according to the instruction trace buffer?

Answer

To number of cycles for one iteration we measured 77 cycles. The static analysis results in 77 cycles.

```
10162210  40001210  call  0x400011a0          [40001210]
10162233  40001214  st    %g1, [%l1]         [40000000 07735935]
10162243  400011a0   mov   42, %g3            [0000002a]
10162253  400011a4   mov   0, %g2            [00000000]
1+2+1+1 = 5
```

10162263	400011a8	add	%g3, 1, %g1	[0000002b]
10162273	400011ac	add	%g1, 1, %g1	[0000002c]
10162283	400011b0	add	%g1, 1, %g1	[0000002d]
10162293	400011b4	add	%g1, 1, %g1	[0000002e]
10162303	400011b8	add	%g1, 1, %g1	[0000002f]
10162313	400011bc	add	%g1, 1, %g1	[00000030]
10162323	400011c0	add	%g1, 1, %g1	[00000031]
10162333	400011c4	add	%g1, 1, %g3	[00000032]
10162343	400011c8	add	%g2, 1, %g2	[00000001]
10162353	400011cc	cmp	%g2, 8	[ffffff9]
10162363	400011d0	bne	0x400011a8	[4000125c]
10162373	400011d4	nop		[00000000]

14

10162383	400011a8	add	%g3, 1, %g1	[00000033]
10162393	400011ac	add	%g1, 1, %g1	[00000034]
10162404	400011b0	add	%g1, 1, %g1	[00000035]
10162414	400011b4	add	%g1, 1, %g1	[00000036]
10162424	400011b8	add	%g1, 1, %g1	[00000037]
10162434	400011bc	add	%g1, 1, %g1	[00000038]
10162444	400011c0	add	%g1, 1, %g1	[00000039]
10162454	400011c4	add	%g1, 1, %g3	[0000003a]
10162464	400011c8	add	%g2, 1, %g2	[00000002]
10162474	400011cc	cmp	%g2, 8	[fffffffa]
10162484	400011d0	bne	0x400011a8	[4000125c]
10162494	400011d4	nop		[00000000]
10162504	400011a8	add	%g3, 1, %g1	[0000003b]
10162514	400011ac	add	%g1, 1, %g1	[0000003c]
10162524	400011b0	add	%g1, 1, %g1	[0000003d]
10162534	400011b4	add	%g1, 1, %g1	[0000003e]
10162544	400011b8	add	%g1, 1, %g1	[0000003f]
10162554	400011bc	add	%g1, 1, %g1	[00000040]
10162564	400011c0	add	%g1, 1, %g1	[00000041]
10162574	400011c4	add	%g1, 1, %g3	[00000042]
10162584	400011c8	add	%g2, 1, %g2	[00000003]
10162594	400011cc	cmp	%g2, 8	[fffffffb]
10162604	400011d0	bne	0x400011a8	[4000125c]
10162614	400011d4	nop		[00000000]
10162624	400011a8	add	%g3, 1, %g1	[00000043]
10162634	400011ac	add	%g1, 1, %g1	[00000044]
10162644	400011b0	add	%g1, 1, %g1	[00000045]
10162654	400011b4	add	%g1, 1, %g1	[00000046]
10162664	400011b8	add	%g1, 1, %g1	[00000047]
10162674	400011bc	add	%g1, 1, %g1	[00000048]
10162684	400011c0	add	%g1, 1, %g1	[00000049]

10162694	400011c4	add	%g1, 1, %g3	[0000004a]
10162704	400011c8	add	%g2, 1, %g2	[00000004]
10162714	400011cc	cmp	%g2, 8	[fffffffc]
10162724	400011d0	bne	0x400011a8	[4000125c]
10162734	400011d4	nop		[00000000]
10162744	400011a8	add	%g3, 1, %g1	[0000004b]
10162754	400011ac	add	%g1, 1, %g1	[0000004c]
10162764	400011b0	add	%g1, 1, %g1	[0000004d]
10162774	400011b4	add	%g1, 1, %g1	[0000004e]
10162784	400011b8	add	%g1, 1, %g1	[0000004f]
10162795	400011bc	add	%g1, 1, %g1	[00000050]
10162805	400011c0	add	%g1, 1, %g1	[00000051]
10162815	400011c4	add	%g1, 1, %g3	[00000052]
10162825	400011c8	add	%g2, 1, %g2	[00000005]
10162835	400011cc	cmp	%g2, 8	[fffffffd]
10162845	400011d0	bne	0x400011a8	[4000125c]
10162855	400011d4	nop		[00000000]
10162865	400011a8	add	%g3, 1, %g1	[00000053]
10162875	400011ac	add	%g1, 1, %g1	[00000054]
10162885	400011b0	add	%g1, 1, %g1	[00000055]
10162895	400011b4	add	%g1, 1, %g1	[00000056]
10162905	400011b8	add	%g1, 1, %g1	[00000057]
10162915	400011bc	add	%g1, 1, %g1	[00000058]
10162925	400011c0	add	%g1, 1, %g1	[00000059]
10162935	400011c4	add	%g1, 1, %g3	[0000005a]
10162945	400011c8	add	%g2, 1, %g2	[00000006]
10162955	400011cc	cmp	%g2, 8	[fffffffe]
10162965	400011d0	bne	0x400011a8	[4000125c]
10162975	400011d4	nop		[00000000]
10162985	400011a8	add	%g3, 1, %g1	[0000005b]
10162995	400011ac	add	%g1, 1, %g1	[0000005c]
10163005	400011b0	add	%g1, 1, %g1	[0000005d]
10163015	400011b4	add	%g1, 1, %g1	[0000005e]
10163025	400011b8	add	%g1, 1, %g1	[0000005f]
10163035	400011bc	add	%g1, 1, %g1	[00000060]
10163045	400011c0	add	%g1, 1, %g1	[00000061]
10163055	400011c4	add	%g1, 1, %g3	[00000062]
10163065	400011c8	add	%g2, 1, %g2	[00000007]
10163075	400011cc	cmp	%g2, 8	[fffffff]
10163085	400011d0	bne	0x400011a8	[4000125c]
10163095	400011d4	nop		[00000000]
10163105	400011a8	add	%g3, 1, %g1	[00000063]
10163115	400011ac	add	%g1, 1, %g1	[00000064]
10163125	400011b0	add	%g1, 1, %g1	[00000065]
10163135	400011b4	add	%g1, 1, %g1	[00000066]
10163145	400011b8	add	%g1, 1, %g1	[00000067]

```

10163155  400011bc  add  %g1, 1, %g1          [00000068]
10163165  400011c0  add  %g1, 1, %g1          [00000069]
10163175  400011c4  add  %g1, 1, %g3          [0000006a]
10163186  400011c8  add  %g2, 1, %g2          [00000008]
10163196  400011cc  cmp  %g2, 8               [00000000]
10163206  400011d0  bne  0x400011a8           [4000125c]
10163216  400011d4  nop                        [00000000]
8*14
10163228  400011d8  retl                       [400011d8]
1

```

5+8*14+1

Just based on the number of instructions we calculated the number of cycles needed = 118 cycles. This is much shorter than the result of the static analysis. The reason for this is that we don't know the exact memory timing and the instruction decode takes more time.

Problem 3

Mandatory (4): First, create a project containing the files contained in the insertion sort folder of the task specification. Now complete the function `main.c:run()`, executing insertion sort a few times, with array size 32 and different input data. Measure the minimum and maximum time needed to execute the sort function.

Q: What were the results of the measurement? How many test sets would do you need to cover all possible execution path?

Answer

We used the compiler flag `-Os`.

Case	measurement[cycles]
Best-case (pre sorted)	3659
Worst-case (upside down sorted)	44267
Average-case (unsorted)	24875

Problem 4

Mandatory (5): Add loop bounds and additional flow facts for `insertion sort.c:insertion sort()`, using the symbolic name `@size` for the size of the array to be sorted. Next, analyze the WCET of insertion sort, assuming an array size of 32. Keep the array size as a symbolic name (user register `@size`). Finally, write a test function which calls insertion sort more than once, with different array sizes (e.g., 16, 32 and 64). Also repeat the static analysis with different array sizes.

Q: How many cycles do you need to execute insertion sort according to the static analysis?

Answer

The static analysis resulted in 74694 cycles.

Q: What results do you get for an array size of 8,16 or 64, using measurements and static analysis?

Answer

method	size 16[cycles]	size 32[cycles]	size 64[cycles]
measurement	11527	34579	132899
static	18232	77224	318088

Q: In addition to the size of the array, what other aspects of the input data might influence the WCET?

Answer

The structure of the data. For instance: The worst-case occurs when the input data is sorted upside down. The best-case occurs if the input data is already sorted.

Problem 5

Recommended (8 Points): Assume that your goal was to find out the WCET of `task.c:task()`. Before analyzing the execution time, you should answer a few basic questions about the input data for the monitoring task, and analyze the control flow on the source code level.

Q: What is the set of input data which might influence the execution time of the task at the software side?

Answer

If many samples are missing the interpolation of the data has big influence on the runtime. Most of the other loops just depend on the amount of samples used for calculation but this does not depend on the input data and is determined at compile time.

Q: Is it tractable to enumerate every possible input?

Answer

No it's not traceable. The result of the interpolation is dependent on the history of the last samples so enumerating the history and the inputs would require many different testcases.

Q: Which loops need to be bounded?

Answer

The loops we found are in the merge samples function.

Q: Add all loop bounds and flow facts you can find to the file `task.c` (as source code annotations).

Answer

```
if (! IS_VALUE_MISSING(x))
{
```

```

/* Only interpolate if we interpolate at most MAX_CONSECUTIVE_MISSING samples
int missing_samples = i - valid - 1;
if(missing_samples > 0 && missing_samples <= MAX_CONSECUTIVE_MISSING)
{
    /* TODO: loop bound */
    for(j = i-1; j > valid; --j)
    {
        /* At most once for each invalid input sample */
        /* TODO: flow fact */
        sample_value_t y = sample_buffer_get(sbuf, j);
        if(! IS_VALUE_MISSING(y)) break;
        y = iinterpolate16(valid, sample_buffer_get(sbuf, valid), i, x, j);
        sample_buffer_set(sbuf, j, y);
    }
}
valid = i;
}

```

Problem 6

Recommended (8 Points): Analyze the `fft()` function called in `task.c`. Try to find loop bounds for the Fast Fourier Transform implementation (`fixedpoint.c:fp_radix2fft` withscaling) first. If you have difficulties finding them, add a debug statement and run the transform with different input data sizes. Add flow constraints relating the execution frequency of the inner loops with the functions execution frequency. Finally, try to analyze the execution time using `aiT`. There is already a timing measurement for the `fft` in the executable, so it is easy to compare the number of cycles estimated to execute the function.

Q: Compare the worst-case number of iterations for the inner loop with and without using these flow constraints. Finally, think about the complexity of calculating loop bounds for FFT.

Answer

Comparison static analysis with measurement:

type	cycles
measurement	178807
static with flow constraints	160344

Comparison of inner loop in `fp_radix2fft` withscaling with/without flow constraints:

type
static without flow
static with flow co

Q: Does the FFT loop bound depend on the input data?

Answer

It just depends on the ammount of the input data and not the data itself.

Problem 7

Optional Challenge (5 Bonus Points): Try to analyze the WCET of `task.c:task()` using

aiT. If you attempt to solve this challenge, use the control flow graph and disassembling capabilities of aiT, and be sure to understand the source code you are analyzing.

Problem 8

Mandatory (4): Answer the following questions

Q: How much time did you spend writing annotations and analyzing the code? Was it less or more than you expected? How much time did you spend on this first assignment?

Answer

- Mandatory Part: 6 hours
- Recommended Part:

Q: What is the ratio between observed and actual execution time? Discuss the causes of the overestimation.

Answer

Q: As you learned, sometimes it is necessary to annotate the assembler code. Why? What problems can you see because of this?

Answer