

Semantics and Verification

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1 Toward Stone Duality.

Question 1. Show that every Stone space (X, Ω) is Hausdorff (if $x, y \in X$ are distinct, there there are disjoint $U, V \in \Omega$ such that $x \in U$ and $y \in V$).

Let $x, y \in X$ be two distinct points of a Stone space (X, Ω) . As, (X, Ω) is T_0 and without loss of generality, there exists $W \in \Omega$ such that $x \in W$ and $y \notin W$. As (X, Ω) is zero-dimensional, we can write $W =: \bigcup_{i \in I} W_i$ where $W_i \in \mathbf{K}\Omega$ for every $i \in I$. Thus, there exists a clopen set $U := W_i \in \Omega$ such that $x \in W_i \subseteq U$. Define $V := X \setminus U \in \Omega$, and we have that $x \in U, y \in V$ (as $y \notin W \supseteq U$) and the open sets U and V are disjoint. We can conclude that every Stone space is Hausdorff.

Question 2. Show that \leq is a partial order on $\mathcal{L}(\text{LML})$.

We start by showing the following lemma.

Lemma 1. We have $\phi \leq \psi$ if and only if $\llbracket \phi \rrbracket \subseteq \llbracket \psi \rrbracket$.

Proof. We have that $\phi \leq \psi$ iff $\phi \equiv \phi \wedge \psi$ iff $\llbracket \phi \rrbracket = \llbracket \phi \wedge \psi \rrbracket = \llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket$ (that last equality is by definition of $\llbracket - \rrbracket$) iff $\llbracket \phi \rrbracket \subseteq \llbracket \psi \rrbracket$. \square

We can thus easily show that \leq is a partial order.

- ▷ *Reflexivity.* As $\llbracket \phi \rrbracket \subseteq \llbracket \phi \rrbracket$, we have that $\phi \leq \phi$ for every $\phi \in \mathcal{L}(\text{LML})$.
- ▷ *Transitivity.* For any $\phi, \psi, \vartheta \in \mathcal{L}(\text{LML})$, if $\phi \leq \psi$ and $\psi \leq \vartheta$ then, by the lemma, $\llbracket \phi \rrbracket \subseteq \llbracket \psi \rrbracket \subseteq \llbracket \vartheta \rrbracket$, thus we have $\llbracket \phi \rrbracket \subseteq \llbracket \vartheta \rrbracket$, i.e. $\phi \leq \vartheta$.
- ▷ *Antisymmetry.* For any $\phi, \psi \in \mathcal{L}(\text{LML})$, if $\phi \leq \psi$ and $\psi \leq \phi$ then, by double inclusion with the above lemma, $\llbracket \phi \rrbracket = \llbracket \psi \rrbracket$ thus $\phi = \psi$ as we consider LML-formulae quotiented by \equiv .

2 Lattices and Boolean Algebras.

2.1 Semilattices.

Question 3. Let (L, \leq) be a partial order.

1. Show that (L, \leq) is a meet semilattice if, and only if, L has binary meets $\wedge : L \times L \rightarrow L$ and greatest element $\top \in L$.
2. Show that (L, \leq) is a join semilattice if, and only if, L has binary joins $\vee : L \times L \rightarrow L$ and least element $\perp \in L$.
1. If (L, \leq) is a meet semilattice, then L has binary meets and a greatest element $\top = \bigwedge \emptyset$ (any element is a lower bound of \emptyset , thus the greatest lower bound of \emptyset is the greatest element).

Now, suppose (L, \leq) has a binary meet \wedge and a greatest element \top . Consider $\{a_i \mid i \in I\}$ a finite subset of elements of L . By induction on $\#I \in \mathbb{N}$, we define $\bigwedge_{i \in I} a_i \in I$ and show that $\bigwedge_{i \in I} a_i$ is a meet of the finite set $\{a_i \mid i \in I\}$ (like the notation suggests).

- ▷ Define $\bigwedge_{i \in \emptyset} a_i := \top \in L$; as any element is a lower bound of \emptyset , the greatest lower bound of \emptyset is the greatest element.
- ▷ Consider $I := J \sqcup \{i\}$. By induction hypothesis, we have that $\bigwedge_{j \in J} a_j$ exists in L and is a meet of $\{a_j \mid j \in J\}$ in (L, \leq) . Define

$$\bigwedge_{k \in I} a_k := \left(\bigwedge_{j \in J} a_j \right) \wedge a_i \in I.$$

We have that $\bigwedge_{k \in I} a_k$ is a lower bound of $\{a_k \mid k \in I\}$. Consider an element a_k with $k \in I$. If $k \in J$ then $a_k \leq \bigwedge_{j \in J} a_j \leq \bigwedge_{k' \in I} a_{k'}$. Otherwise $k = i$ and we immediacy have that $a_i \leq \bigwedge_{k' \in I} a_{k'}$.

Consider a lower bound $b \in L$ of $\{a_k \mid k \in I\}$, then b is a lower bound of $\{a_j \mid j \in J\}$ and $b \leq a_i$. We have $b \leq \bigwedge_{j \in J} a_j$ and $b \leq a_i$, therefore $b \leq \bigwedge_{k \in I} a_k$.

We can conclude that $\bigwedge_{k \in I} a_k$ is a meet of $\{a_k \mid k \in K\}$.

Finally, we have that (L, \leq) has finite meets.

2. This results follows from 1 when considering the partial order (L, \geq) , by duality. Meets in (L, \geq) are exactly joins in (L, \leq) , and the greatest element of (L, \geq) is the least element of (L, \leq) , and *vice versa*.

Note. In the following, when I will be dealing with multiple partial orders on the same set (e.g. \leq and \geq), I will write \bigwedge_{\leq} for the meet operator in poset (I, \leq) , \bigvee_{\leq} for the join operator in poset (I, \leq) , \top_{\leq} for the greatest element in poset (I, \leq) and \perp_{\leq} for the least element in poset (I, \leq) .

Question 4. Prove the following.

1. Let (L, \leq) be a meet semilattice with binary meets $\wedge : L \times L \rightarrow L$ and greatest element $\top \in L$. Then (L, \wedge, \top) is a commutative monoid in which every element is idempotent. Moreover, we have $a \leq b$ iff $a = a \wedge b$.
2. Let (L, \leq) be a join semilattice with binary joins $\vee : L \times L \rightarrow L$ and least element \perp . Then (L, \vee, \perp) is a commutative monoid in which every element is idempotent. Moreover, we have $a \leq b$ iff $b = a \vee b$.
1. Let $a, b, c \in L$. First, we have that $a \wedge b = \bigwedge\{a, b\} = \bigwedge\{b, a\} = b \wedge a$ thus the binary meet operation \wedge is commutative. Then, as a special case of the previous question, we have that a and $\top \wedge a$ are both meets of $\{a\}$. And, by unicity of meets (i.e. antisymmetry of \leq , mainly), they are equals. Also as a special case of the previous question, we have that elements

$$a \wedge (b \wedge c) = \top \wedge (a \wedge (b \wedge c))$$

and

$$(a \wedge b) \wedge c = c \wedge (a \wedge b) = \top \wedge (c \wedge (a \wedge b))$$

are both meets of the set $\{a, b, c\}$, thus are equal. Next, we have that

$$a \wedge a = \bigwedge\{a, a\} = \bigwedge\{a\} = \top \wedge a = a$$

(penultimate equality is from last question), thus $a \wedge a = a$. Finally, we have that:

▷ if $a = a \wedge b$ then a is a lower bound of $\{a, b\}$ thus $a \leq b$;

- ▷ if $a \leq b$ then $a = a \wedge b$ as a is a lower bound of $\{a, b\}$ and any lower bound c of $\{a, b\}$ must satisfy $c \leq a$.
- 2. Consider the meet semilattice (L, \geq) and apply the results above. Meets in (L, \geq) are exactly joins in (L, \leq) , and the greatest element of (L, \geq) is the least element of (L, \leq) , and *vice versa*. The last statement follows from the equivalence:

$$a \leq b \quad \text{iff} \quad b \geq a \quad \text{iff} \quad b = a \wedge_{\geq} b \quad \text{iff} \quad b = a \vee_{\leq} b,$$

where the second “iff” follows from the result above for (L, \geq) , and the last one follows from the equality $a \wedge_{\geq} b = a \vee_{\leq} b$.

Question 5. Prove the following.

1. Given a commutative monoid (L, \wedge, \top) in which every element is idempotent, let $a \leq_{\wedge} b$ iff $a = a \wedge b$. Then (L, \leq_{\wedge}) is a meet semilattice with binary meets given by \wedge and greatest element \top .
2. Given a commutative monoid (L, \vee, \perp) in which every element is idempotent, let $a \leq_{\vee} b$ iff $b = a \vee b$. Then (L, \leq_{\vee}) is a join semilattice with binary joins given by \vee and least element \perp .

1. Let us start by showing that (L, \leq_{\wedge}) is a partial order.

- ▷ *Reflexivity.* As $a \wedge a = a$ by idempotence, we have $a \leq_{\wedge} a$.
- ▷ *Antisymmetry.* If $a \leq_{\wedge} b$ and $b \leq_{\wedge} a$ then, by commutativity, we have $a \wedge b = a = b$.
- ▷ *Transitivity.* If $a \leq_{\wedge} b$ and $b \leq_{\wedge} c$ then, by associativity,

$$a = a \wedge b = a \wedge (b \wedge c) = (a \wedge b) \wedge c = a \wedge c,$$

thus $a \leq_{\wedge} c$.

By question 3, it suffices to show that (L, \leq_{\wedge}) that binary meets for poset (L, \leq_{\wedge}) are \wedge and that \top is the greatest element of poset (L, \leq_{\wedge}) . Consider a, b, c three arbitrary elements of L .

- ▷ For any $b \in L$, we have $b \wedge \top = b$ (as \top is a neutral element) and thus $b \leq_{\wedge} \top$ for all $b \in L$, so \top is the greatest element of (L, \leq_{\wedge}) .

▷ Firstly, element $a \wedge b$ is a lower bound of $\{a, b\}$ as

$$\begin{aligned} a \wedge b \leq_{\wedge} a & \quad \text{iff} \quad a \wedge b = (a \wedge b) \wedge a \\ a \wedge b \leq_{\wedge} b & \quad \text{iff} \quad a \wedge b = (a \wedge b) \wedge b \end{aligned}$$

and the latter equalities are true by idempotence, associativity, and finally commutativity. Secondly, consider $c \in L$ such that we have $c \leq_{\wedge} a$ and $c \leq_{\wedge} b$, then $c \wedge a = c = c \wedge b$. We therefore have that $c \leq_{\wedge} a \wedge b$, as

$$c \wedge (a \wedge b) = (c \wedge a) \wedge b = c \wedge b = c.$$

We can conclude that \wedge is the binary meet operator in (L, \leq_{\wedge}) .

2. Applying the previous result with the commutative monoid (L, \vee, \perp) , we obtain that $(L, \geq_{\vee})^a$ is a meet semilattice where binary meets for \geq_{\vee} are given by \vee and the greatest element for \geq_{\vee} is \perp . We can thus conclude that (L, \leq_{\vee}) is a join semilattice where binary joins for \leq_{\vee} are given by \vee and the least element for \leq_{\vee} is \perp .

Question 6. Show the following, for the partial order $(\mathfrak{L}(\text{LML}), \leq)$:

1. $(\mathfrak{L}(\text{LML}), \leq)$ is a meet semilattice with greatest element \top and binary joins given by

$$\begin{aligned} - \wedge - : \mathfrak{L}(\text{LML}) \times \mathfrak{L}(\text{LML}) & \longrightarrow \mathfrak{L}(\text{LML}) \\ (\phi, \psi) & \longmapsto \phi \wedge \psi; \end{aligned}$$

2. $(\mathfrak{L}(\text{LML}), \leq)$ is a join semilattice with least element \perp and binary joins given by

$$\begin{aligned} - \vee - : \mathfrak{L}(\text{LML}) \times \mathfrak{L}(\text{LML}) & \longrightarrow \mathfrak{L}(\text{LML}) \\ (\phi, \psi) & \longmapsto \phi \vee \psi. \end{aligned}$$

We will use the lemma proven in question 2 (lemma 1, page 1).

^aThe notation is, in a way, “context-sensitive,” as for an arbitrary monoid (M, \otimes, \mathbf{I}) , we can either define $a \leq_{\otimes} b$ as $a \otimes b = a$ or $a \leq_{\otimes} b$ as $a \otimes b = b$.

1. We only need to show that $-\wedge-$ defines a binary meet for $(\mathfrak{L}(\text{LML}), \leq)$ and that \top is a greatest element.

For any $\phi \in \mathfrak{L}(\text{LML})$, we have $\phi \leq \top$ as $\llbracket \phi \rrbracket \subseteq \llbracket \top \rrbracket = (\mathbf{2}^{\text{AP}})^\omega$, thus \top is the greatest element.

For any formulae $\phi, \psi \in \mathfrak{L}(\text{LML})$, we have that $\phi \wedge \psi \leq \phi$ and $\phi \wedge \psi \leq \psi$ as both $\llbracket \phi \rrbracket$ and $\llbracket \psi \rrbracket$ are supersets of $\llbracket \phi \wedge \psi \rrbracket = \llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket$ (by definition of interpretation $\llbracket - \rrbracket$). Then, if $\vartheta \leq \phi$ and $\vartheta \leq \psi$, we have that $\llbracket \vartheta \rrbracket \subseteq \llbracket \phi \rrbracket$ and $\llbracket \vartheta \rrbracket \subseteq \llbracket \psi \rrbracket$ thus $\llbracket \vartheta \rrbracket \subseteq \llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket = \llbracket \phi \wedge \psi \rrbracket$, therefore $\vartheta \leq \phi \wedge \psi$.

We can conclude that $(\mathfrak{L}(\text{LML}), \leq)$ is a meet semilattice with greatest element \top and binary meets given by $-\wedge-$.

2. We only need to show that $-\vee-$ defines a binary join for $(\mathfrak{L}(\text{LML}), \leq)$ and that \perp is a least element.

For any $\phi \in \mathfrak{L}(\text{LML})$, we have $\perp \leq \phi$ as $\emptyset \subseteq \llbracket \perp \rrbracket \subseteq \llbracket \phi \rrbracket$, thus \perp is the least element.

For any formulae $\phi, \psi \in \mathfrak{L}(\text{LML})$, we have that $\phi \leq \phi \vee \psi$ and $\psi \leq \phi \vee \psi$ as both $\llbracket \phi \rrbracket$ and $\llbracket \psi \rrbracket$ are subsets of $\llbracket \phi \vee \psi \rrbracket = \llbracket \phi \rrbracket \cup \llbracket \psi \rrbracket$ (by definition of interpretation $\llbracket - \rrbracket$). Then, if $\phi \leq \vartheta$ and $\psi \leq \vartheta$, we have that $\llbracket \phi \rrbracket \subseteq \llbracket \vartheta \rrbracket$ and $\llbracket \psi \rrbracket \subseteq \llbracket \vartheta \rrbracket$ thus $\llbracket \phi \vee \psi \rrbracket = \llbracket \phi \rrbracket \cup \llbracket \psi \rrbracket \subseteq \llbracket \vartheta \rrbracket$, therefore $\phi \vee \psi \leq \vartheta$.

We can conclude that $(\mathfrak{L}(\text{LML}), \leq)$ is a join semilattice with least element \perp and binary joins given by $-\vee-$.

Question 7. Show that a map of meet (resp. join) semilattices is monotone.

Let $f : L \rightarrow L'$ be an arbitrary function where (L, \leq) and (L', \leq') are partial orders.

1. Suppose $f : (L, \leq) \rightarrow (L', \leq')$ is a map of meet semilattices. Let $a, b \in L$. If $a \leq b$, then $a \wedge b = a$ and, as f preserves finite meets,

$$f(a) \wedge' f(b) = f(a \wedge b) = f(a),$$

and thus $f(a) \leq' f(b)$. Therefore, f is monotone.

2. Suppose $f : (L, \leq) \rightarrow (L', \leq')$ is a map of join semilattices. Let $a, b \in L$. If $a \leq b$, then $a \vee b = b$ and, as f preserves finite joins,

$$f(a) \vee' f(b) = f(a \vee b) = f(b),$$

and thus $f(a) \leq' f(b)$. Therefore, f is monotone.

2.2 Lattices.

Question 8. Consider the partial order (L, \sqsubseteq) where

$$L := \mathbb{N} \cup \{\alpha, \beta, \top\},$$

where \sqsubseteq is the reflexive-transitive closure of \sqsubset , where

$$a \sqsubset b \quad \text{iff} \quad \begin{cases} a < b \text{ in } \mathbb{N} \\ \text{or} \\ a \in \mathbb{N} \text{ and } b \in \{\alpha, \beta\} \\ \text{or} \\ a \in \{\alpha, \beta\} \text{ and } b = \top. \end{cases}$$

Show that (L, \sqsubseteq) is a join semilattice but is not a lattice.

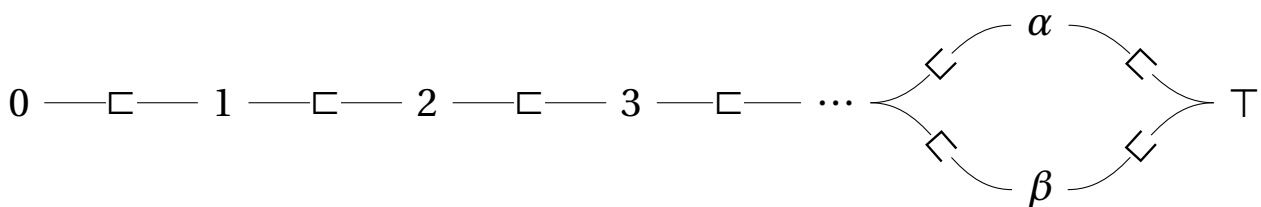


Figure 1 | Hasse diagram of (L, \sqsubseteq) from question 8

Note: Hasse diagrams are usually read bottom-to-top, but this one is drawn left-to-right for convenience.

The relation \sqsubseteq is a partial order. Reflexivity and transitivity is true by definition of \sqsubseteq as the reflexive and transitive closure of \sqsubset . For antisymmetry, we have that:

- ▷ for $n, m \in \mathbb{N}$, $n \sqsubseteq m$ iff $n \leq m$;

- ▷ for any $n \in \mathbb{N}$ and $m \in L \setminus \mathbb{N}$, we have $n \sqsubseteq m$ and $m \not\sqsubseteq n$;
- ▷ $\alpha \sqsubseteq \top$, $\beta \sqsubseteq \top$, $\top \not\sqsubseteq \alpha$, $\top \not\sqsubseteq \beta$, $\alpha \not\sqsubseteq \beta$ and $\beta \not\sqsubseteq \alpha$;

(this can be shown by induction on the relation \sqsubseteq).

We have that 0 is the least element in (L, \sqsubseteq) : we have that $0 \sqsubseteq a$ for all $a \in L$. For $a, b \in L$, we can define $a \vee b$ as:

- ▷ if $a, b \in \mathbb{N}$, let $a \vee b := \min_{\leq \mathbb{N}}(a, b)$;
- ▷ if $a \in \mathbb{N}$ and $b \in L \setminus \mathbb{N}$, let $a \vee b, b \vee a := b$;
- ▷ otherwise let $\alpha \vee \beta := \top$, $a \wedge a := a$, $a \wedge \top, \top \wedge a := \top$ for $a \in \{\alpha, \beta, \top\}$.

Using the previous results on \sqsubseteq , we have that $- \vee -$ really is a join.

This concludes the proof that (L, \sqsubseteq) is a join semilattice.

We also have that (L, \sqsubseteq) is not a lattice. Suppose it is a lattice, and consider the element $a := \alpha \wedge \beta$. Necessarily, we have that $a \in \mathbb{N}$ (if $a = \alpha$ then we would have $\alpha \sqsubseteq \beta$, which is false). As $a = \alpha \wedge \beta$ and $a + 1 \sqsubseteq \alpha, \beta$ we have, by definition of meet, that $a + 1 \sqsubseteq a$, thus $a + 1 \leq a$ (since $a, a + 1 \in \mathbb{N}$) which is **absurd**. We can conclude that (L, \sqsubseteq) is not a lattice.

Question 9. Consider a set L equipped with two binary operations $\wedge, \vee : L \times L \rightarrow L$ and two constants $\top, \perp \in L$. Assume that (L, \wedge, \top) and (L, \vee, \perp) are commutative monoids in which every element is idempotent. Show that the following are equivalent.

1. The partial order \leq_\vee induced by (L, \vee, \perp) coincides with the partial order \leq_\wedge induced by (L, \wedge, \top) .
2. $(L, \vee, \wedge, \perp, \top)$ satisfies the two following **absorptive laws**:

$$\forall a, b \in L, \quad a \vee (a \wedge b) = a \quad (\text{abs}_1)$$

$$\forall a, b \in L, \quad a \wedge (a \vee b) = a \quad (\text{abs}_2)$$

- ▷ Let us show that 1 implies 2. Let $a, b \in L$. We have that $a \wedge b \leq_\wedge a$ and, assuming \leq_\wedge and \leq_\vee coincide, $a \wedge b \leq_\vee a$, thus $a \vee (a \wedge b) = a$, i.e. (abs_1) holds. Similarly, $a \leq_\vee a \vee b$ thus $a \leq_\wedge a \vee b$, so $(a \wedge b) \vee a = a$ holds, and we can recover (abs_2) by using commutativity.
- ▷ Let us show that 2 implies 1.

- Suppose $b \leq_{\wedge} a$, then $b \wedge a = b$. By (abs₁) and commutativity, we have $b \vee a = (b \wedge a) \vee a = a$, thus $b \leq_{\vee} a$.
- Suppose $b \leq_{\vee} a$, then $b \vee a = a$. By (abs₂), we have

$$b \wedge a = b \wedge (b \vee a) = b,$$

thus $b \leq_{\wedge} a$.

Thus the two order coincide.

Question 10. Show that the partial order $(\mathfrak{L}(\text{LML}), \leq)$ is a lattice.

We have shown that $(\mathfrak{L}(\text{LML}), \leq)$ has a greatest element \top , a least element \perp , binary meets given by $-\wedge-$ and binary joins given by $-\vee-$ (question 6). Thus it has all finite meets and finite joins (as seen in question 3), i.e. $(\mathfrak{L}(\text{LML}), \leq)$ is a lattice.

Question 11. Show that the function

$$\begin{aligned} \circ : \mathfrak{L}(\text{LML}) &\longrightarrow \mathfrak{L}(\text{LML}) \\ \phi &\longmapsto \circ\phi \end{aligned}$$

is a morphism of lattices.

We know \circ is a map of meet iff $\circ\top = \top$ and $\circ(\phi \wedge \psi) = \circ\phi \wedge \circ\psi$. Both are true as,

$$\begin{aligned} \llbracket \circ\top \rrbracket &= \{ \sigma \in (\mathbf{2}^{\text{AP}})^{\omega} \mid \sigma \upharpoonright 1 \in \llbracket \top \rrbracket = (\mathbf{2}^{\text{AP}})^{\omega} \} = (\mathbf{2}^{\text{AP}})^{\omega} = \llbracket \top \rrbracket \\ \llbracket \circ(\phi \wedge \psi) \rrbracket &= \{ \sigma \in (\mathbf{2}^{\text{AP}})^{\omega} \mid \sigma \upharpoonright 1 \in \llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket \} = \llbracket \circ\phi \rrbracket \cap \llbracket \circ\psi \rrbracket = \llbracket \circ\phi \wedge \circ\psi \rrbracket. \end{aligned}$$

Very similarly, \circ is a map of joins iff $\circ\perp = \perp$ and $\circ(\phi \vee \psi) = \circ\phi \vee \circ\psi$. One can show that both equalities hold by applying $\llbracket - \rrbracket$ and showing the equality of the sets like above.

Thus $\circ : (\mathfrak{L}(\text{LML}), \leq) \rightarrow (\mathfrak{L}(\text{LML}), \leq)$ is a morphism of lattices.

2.3 Distributive Lattices.

Question 12. Show that the following two **distributive laws** are equivalent in a lattice $(L, \vee, \wedge, \perp, \top)$:

$$\forall a, b, c \in L, \quad a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c) \quad (\text{dist}_1)$$

$$\forall a, b, c \in L, \quad a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c) \quad (\text{dist}_2)$$

Suppose (dist_1) holds and let us show (dist_2) is true for $a, b, c \in L$:

$$\begin{aligned} (a \vee b) \wedge (a \vee c) &= ((a \vee b) \wedge a) \vee ((a \vee b) \wedge c) && \text{by } (\text{dist}_1) \\ &= a \vee ((a \vee b) \wedge c) && \text{by } (\text{abs}_2) \\ &= a \vee (a \wedge b) \vee (b \wedge c) && \text{by } (\text{dist}_1) \\ &= a \vee (b \wedge c) && \text{by } (\text{abs}_1). \end{aligned}$$

To prove that (dist_1) holds when (dist_2) is true, we can apply the previous result to the lattice $(L, \leq)^{\text{op}} = (L, \geq)$. This gives exactly the implication “ (dist_2) implies (dist_1) ,” as wanted.

Thus, the two distributive laws (dist_1) and (dist_2) are equivalent.

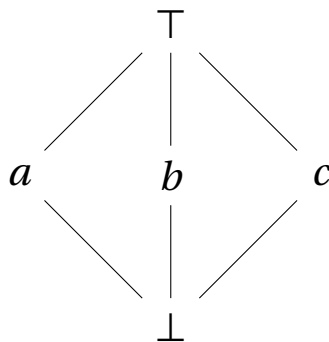
Question 13. Show that the lattice $(\mathcal{L}(\text{LML}), \leq)$ is distributive.

Let $\phi, \psi, \vartheta \in \mathcal{L}(\text{LML})$. We have that

$$\llbracket \phi \wedge (\psi \vee \vartheta) \rrbracket = \llbracket \phi \rrbracket \cap (\llbracket \psi \rrbracket \cup \llbracket \vartheta \rrbracket) = (\llbracket \phi \rrbracket \cap \llbracket \psi \rrbracket) \cup (\llbracket \phi \rrbracket \cap \llbracket \vartheta \rrbracket) = \llbracket (\phi \wedge \psi) \vee (\phi \wedge \vartheta) \rrbracket,$$

thus $\phi \wedge (\psi \vee \vartheta) = (\phi \wedge \psi) \vee (\phi \wedge \vartheta)$.

Question 14. Consider the following lattice M_3 :



(i.e. $\perp \leq a, b, c \leq \top$ with a, b, c incomparable). Show that M_3 is not distributive.

Suppose M_3 is distributive. As a, b, c are incomparable, we have that

$$a \wedge b = a \wedge c = \perp \quad \text{and} \quad b \vee c = \top,$$

and thus,

$$a = a \wedge \top = a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c) = \perp \vee \perp = \perp,$$

which is **absurd**. Thus M_3 is not distributive.

2.4 Booleans algebras.

Question 15. Show that if (L, \leq) is a distributive lattice then $a \in L$ has at most one complement.

Consider $c, c' \in L$ two complements of $a \in L$. Then, we have that

$$c = c \wedge \top = c \wedge (a \vee c') \stackrel{(\text{dist}_1)}{=} (c \wedge a) \vee (c \wedge c') = \perp \vee (c \wedge c') = c \wedge c',$$

and,

$$c' = c' \wedge \top = c' \wedge (a \vee c) \stackrel{(\text{dist}_1)}{=} (c' \wedge a) \vee (c' \wedge c) = \perp \vee (c' \wedge c) = c' \wedge c.$$

We can conclude that $c = c'$ by commutativity of meets.

Question 16. Show that $(\mathcal{L}(\text{LML}), \leq)$ is a Boolean algebra.

Let us show that $\neg\phi$ is a complement for $\phi \in \mathcal{L}(\text{LML})$. We have to check that $\phi \wedge \neg\phi = \perp$ and $\phi \vee \neg\phi = \top$ hold. Both equalities can be easily checked with interpretations:

$$\llbracket \phi \wedge \neg\phi \rrbracket = \llbracket \phi \rrbracket \cap \llbracket \phi \rrbracket^c = \emptyset = \llbracket \perp \rrbracket,$$

and

$$\llbracket \phi \vee \neg\phi \rrbracket = \llbracket \phi \rrbracket \cup \llbracket \phi \rrbracket^c = (2^{\text{AP}})^\omega = \llbracket \top \rrbracket.$$

Thus, $\neg\phi$ is **the** complement of ϕ in $(\mathcal{L}(\text{LML}), \leq)$, which is, as a consequence, a Boolean algebra.

Question 17. Show that the following **De Morgan Laws** hold in every Boolean algebra $(B, \vee, \wedge, \perp, \top)$:

$$a \wedge b = \neg(\neg a \vee \neg b) \quad a \vee b = \neg(\neg a \wedge \neg b) \quad a = \neg\neg a.$$

We start by showing that $\neg a \vee \neg b$ is a complement to $(a \wedge b)$:

$$\begin{aligned} (a \wedge b) \wedge (\neg a \vee \neg b) &= (a \wedge b \wedge \neg b) \vee (a \wedge b \wedge \neg a) = \perp \vee \perp = \perp \\ (a \wedge b) \vee (\neg a \vee \neg b) &= (a \vee \neg a \vee \neg b) \wedge (b \vee \neg a \vee \neg b) = \top \wedge \top = \top, \end{aligned}$$

by using the Boolean algebra laws. Thus $a \wedge b = \neg(\neg a \vee \neg b)$.

For $a \vee b = \neg(\neg a \wedge \neg b)$, we proceed by duality by applying the previous result to the Boolean algebra $(B, \wedge, \vee, \top, \perp)$, as a complement for a in (B, \geq) is exactly a complement for a in (B, \leq) .

We can easily check that $\neg\top = \perp$, and then

$$a = a \wedge \top = \neg(\neg a \vee \neg\top) = \neg(\neg a \vee \perp) = \neg\neg a.$$

Question 18. Show that if f is a map of Boolean algebras from (B, \leq) to (B', \leq') then f preserves complements.

We have that

$$\perp' = f(\perp) = f(a \wedge \neg a) = f(a) \wedge' f(\neg a),$$

and

$$\top = f(\top) = f(a \vee \neg a) = f(a) \vee' f(\neg a),$$

thus $f(\neg a)$ is a complement of $f(a)$ and, by unicity, $f(\neg a) = \neg f(a)$. We can conclude that a map of Boolean algebras preserves complements.

3 Representation of Boolean Algebras.

The rest will be given in the next part.