

Optimal Transport Gromov-Wasserstein Problem for Dimensionality Reduction and Graph Analysis

Hugues Van Assel



Cédric Vincent-Cuaz



Rémi Flamary



Nicolas Courty



Pascal Frossard



Titouan Vayer

My talk

Overview of **dimensionality reduction**

Optimal Transport : from linear OT
to **Gromov Wasserstein**

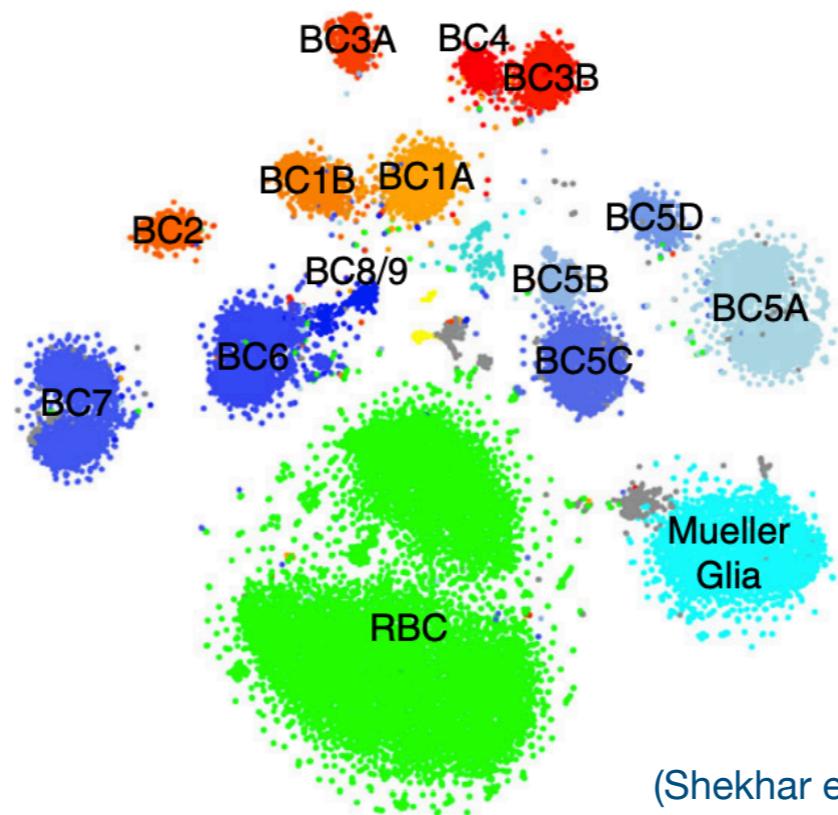
vectorial
data

Distributional Reduction : a framework
to embed distributions

graph
data

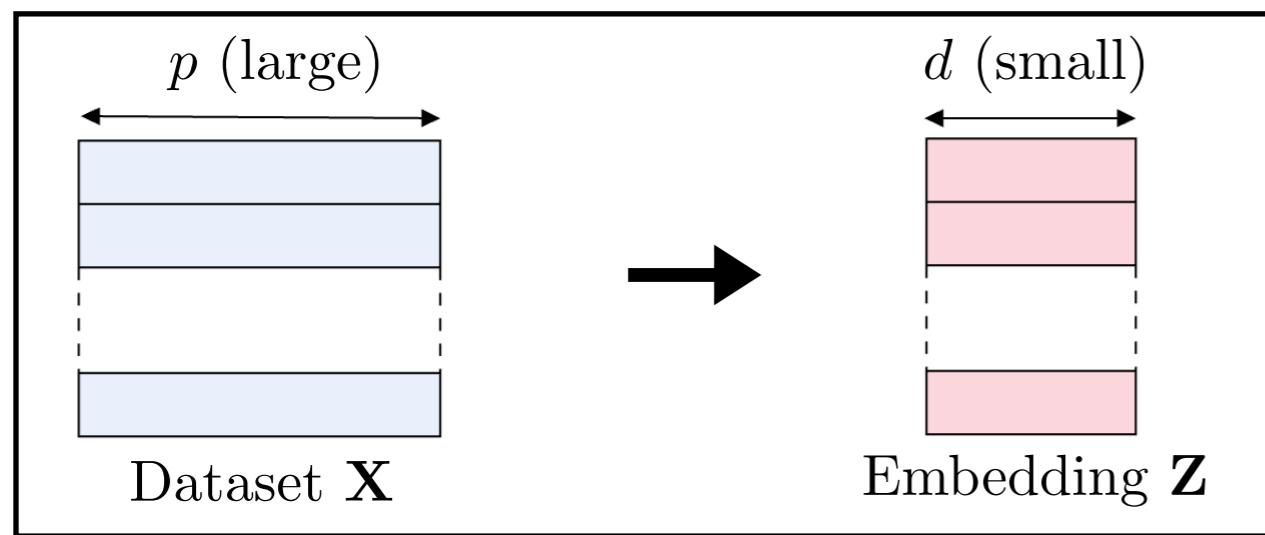
Gromov Wasserstein for graph analysis
Application to graph **generative modeling**

Dimension Reduction

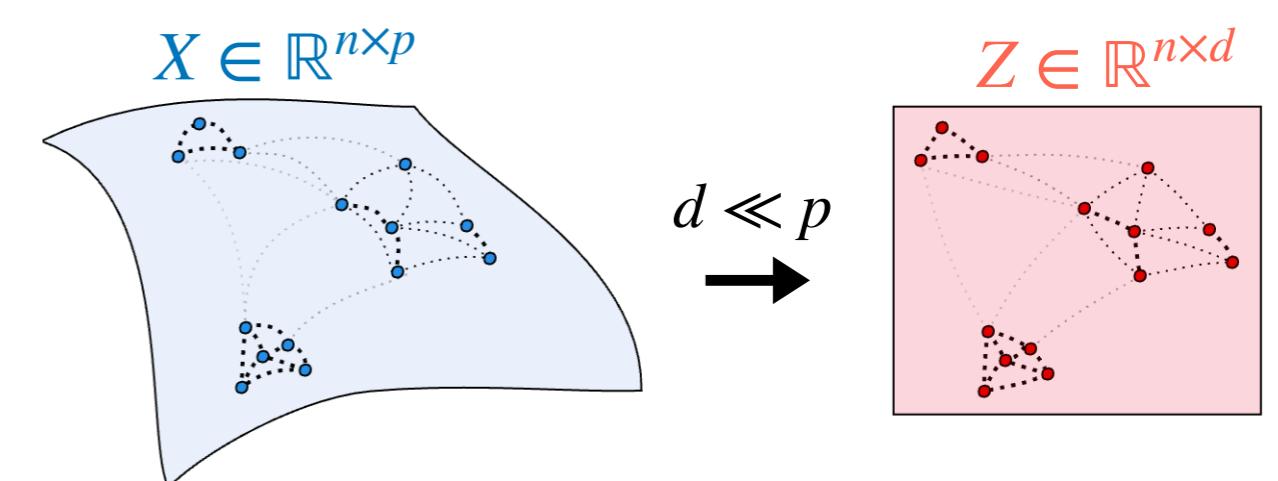


(Shekhar et al., 2016)

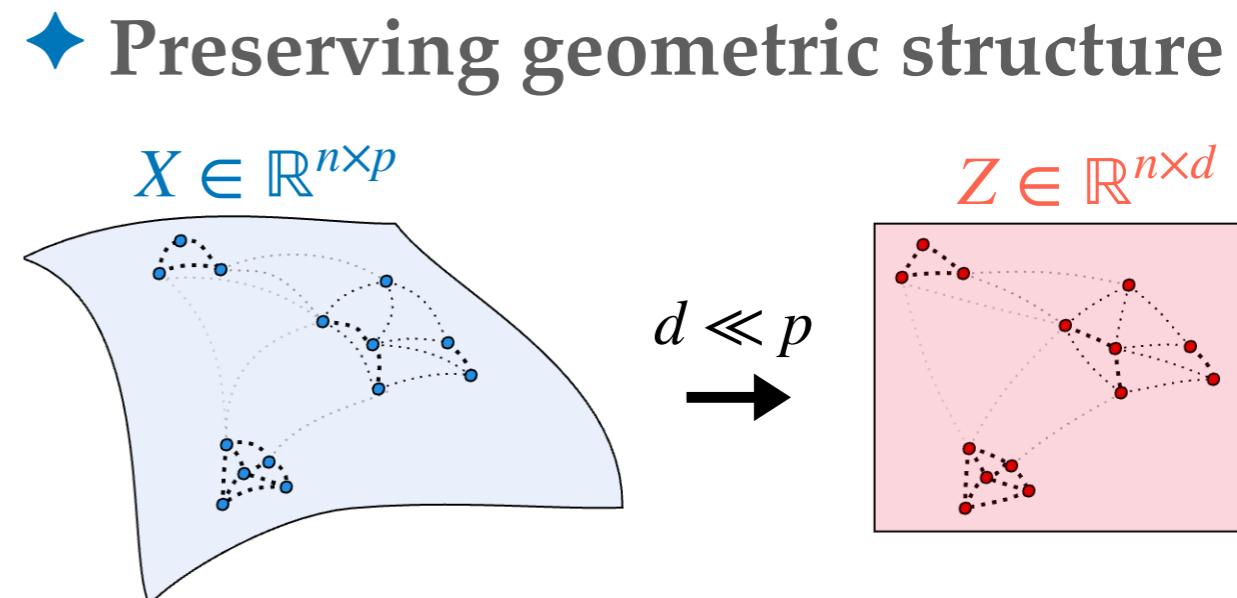
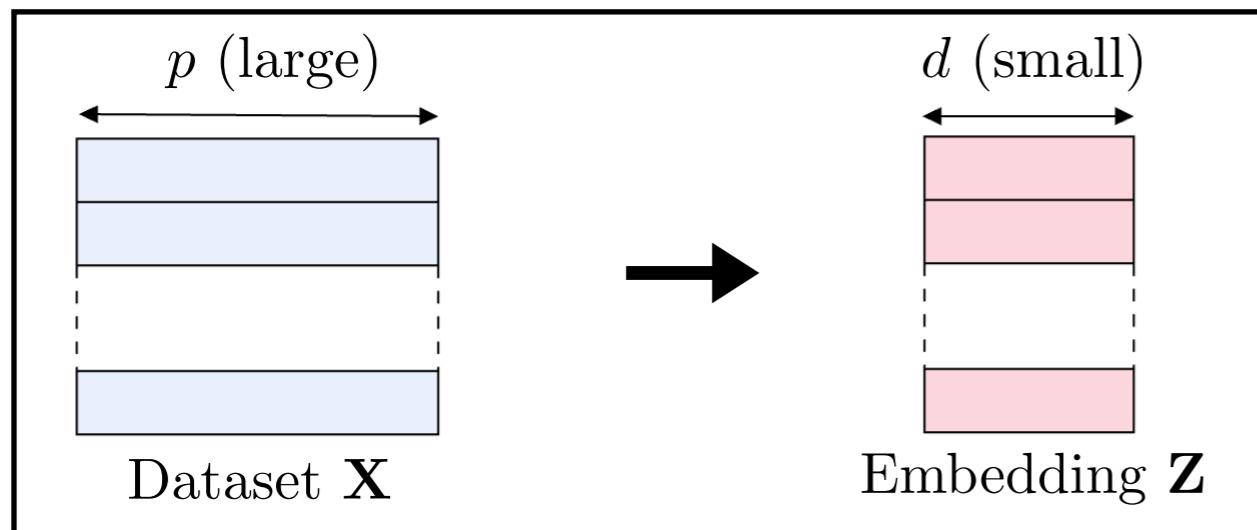
Dimension reduction



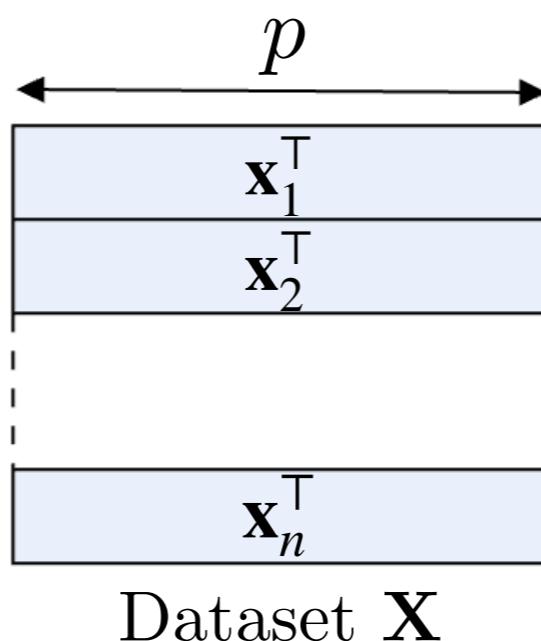
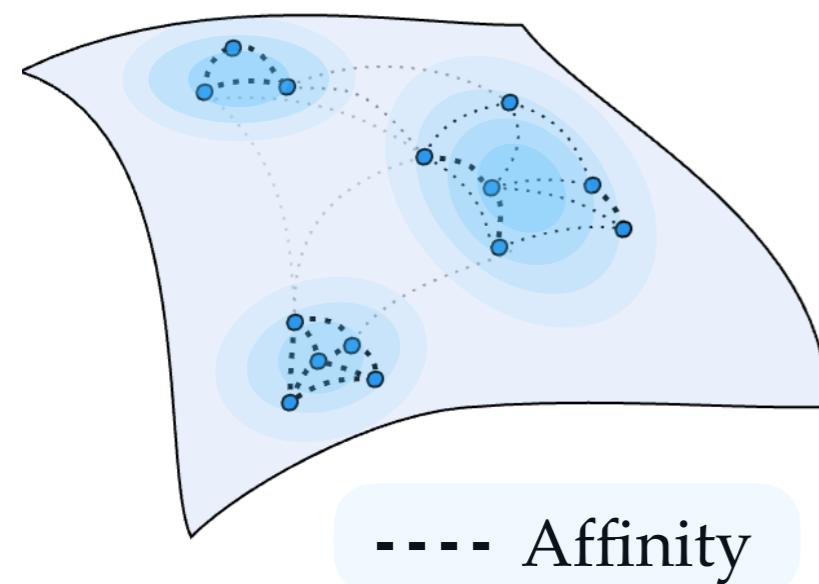
◆ Preserving geometric properties



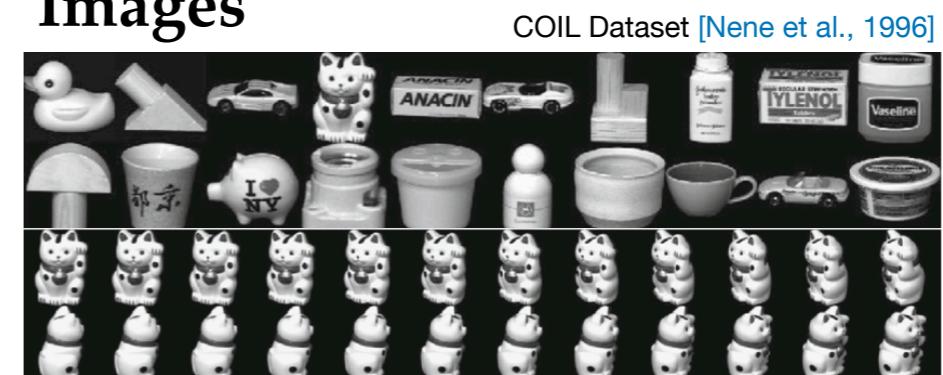
Dimension reduction



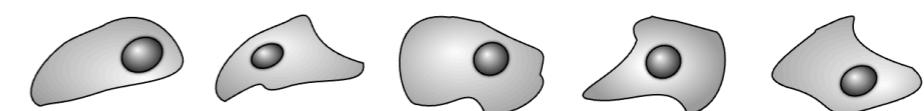
◆ Affinity Matrices



Images



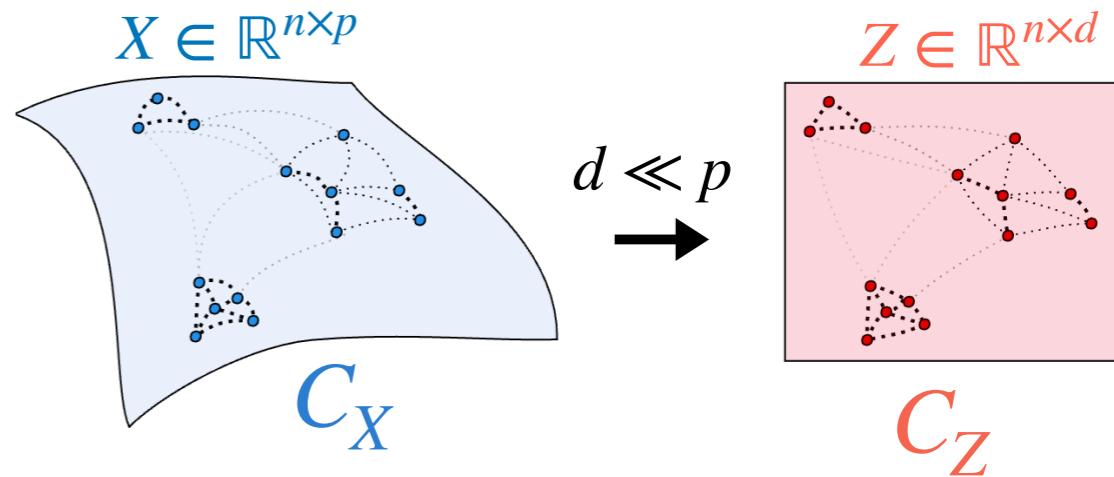
Cells



Symmetric matrix with non-negative coefficients.

Coefficient (i, j) = similarity between \mathbf{x}_i and \mathbf{x}_j .

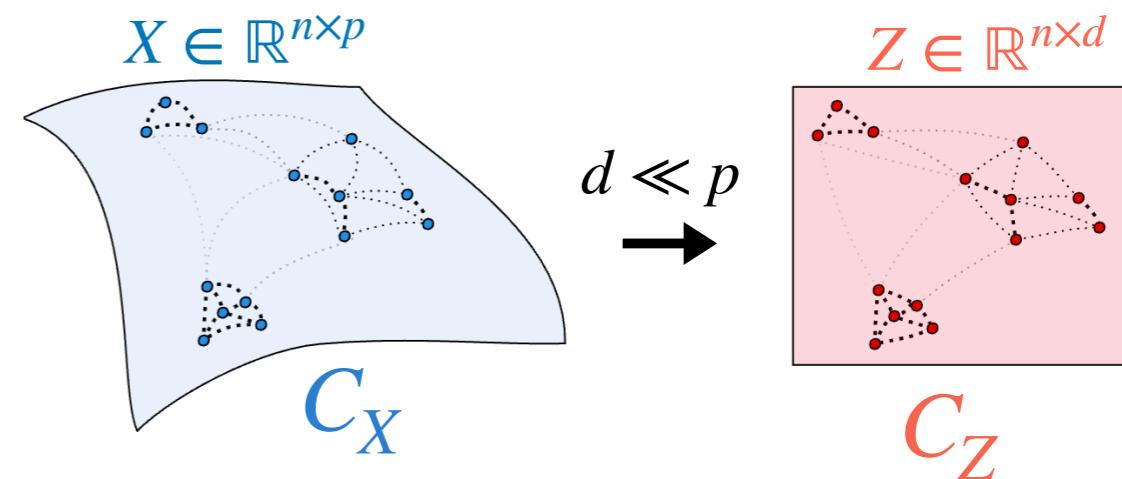
Dimension reduction



◆ A general optimization problem

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right) \text{ for some loss } L : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$$

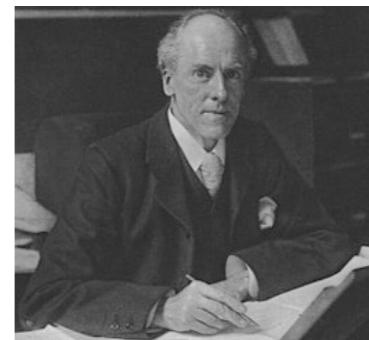
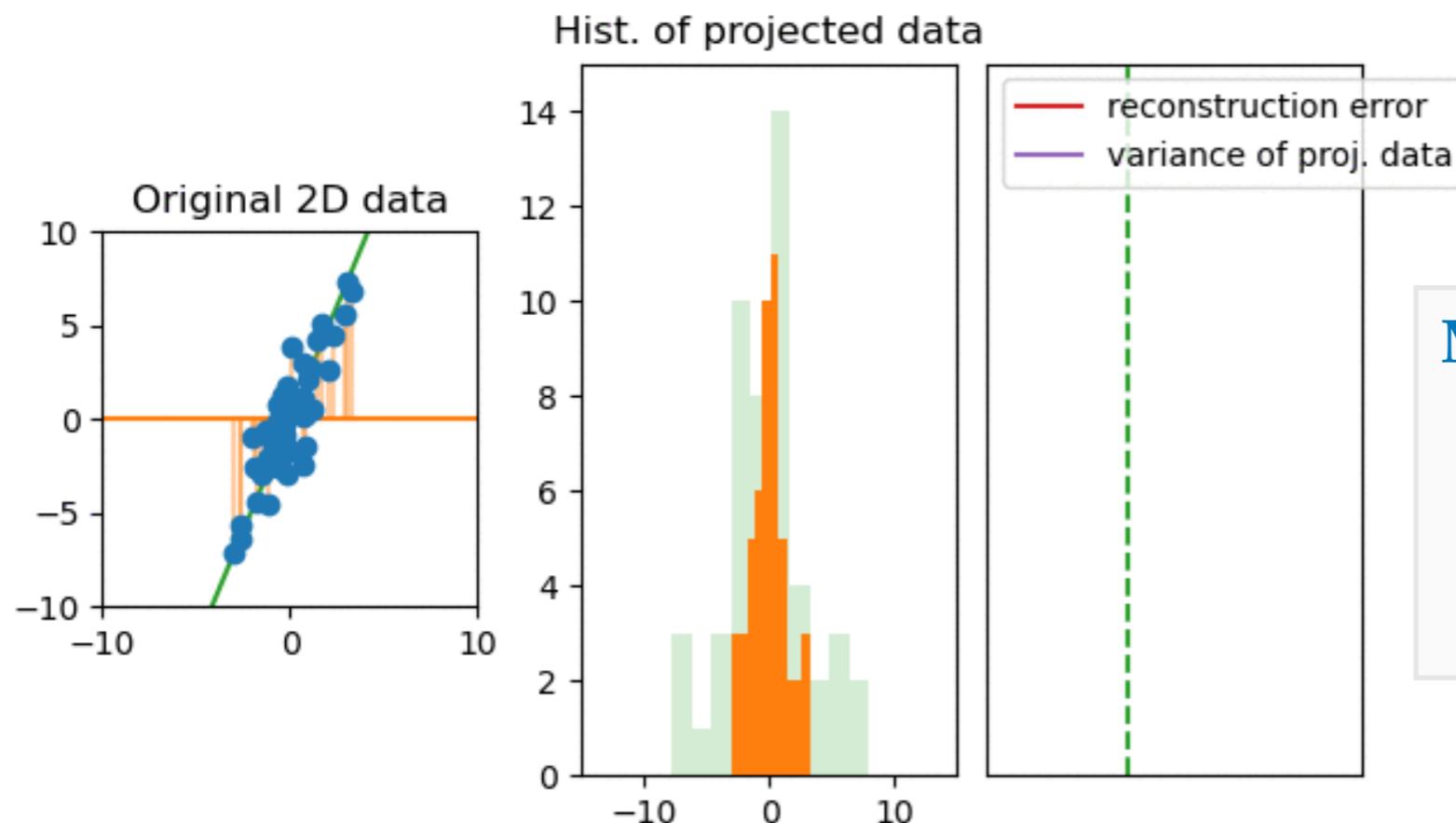
Dimension reduction



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♦ Principal components analysis

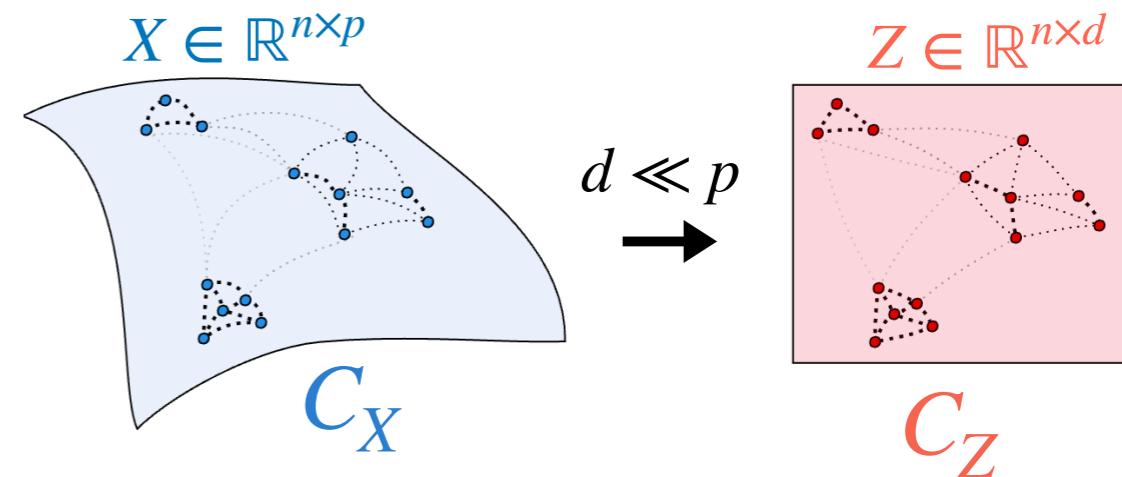


(Pearson, 1901)

Minimizing the reconstruction error

$$\min_{H: \dim(H)=d} \frac{1}{n} \sum_{i=1}^n \|\mathbf{x}_i - P_H(\mathbf{x}_i)\|_2^2$$

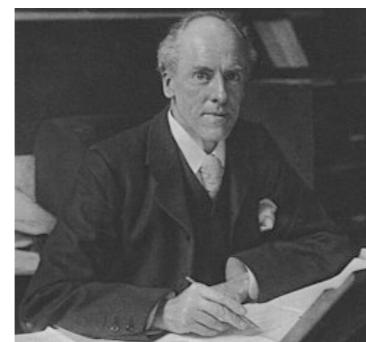
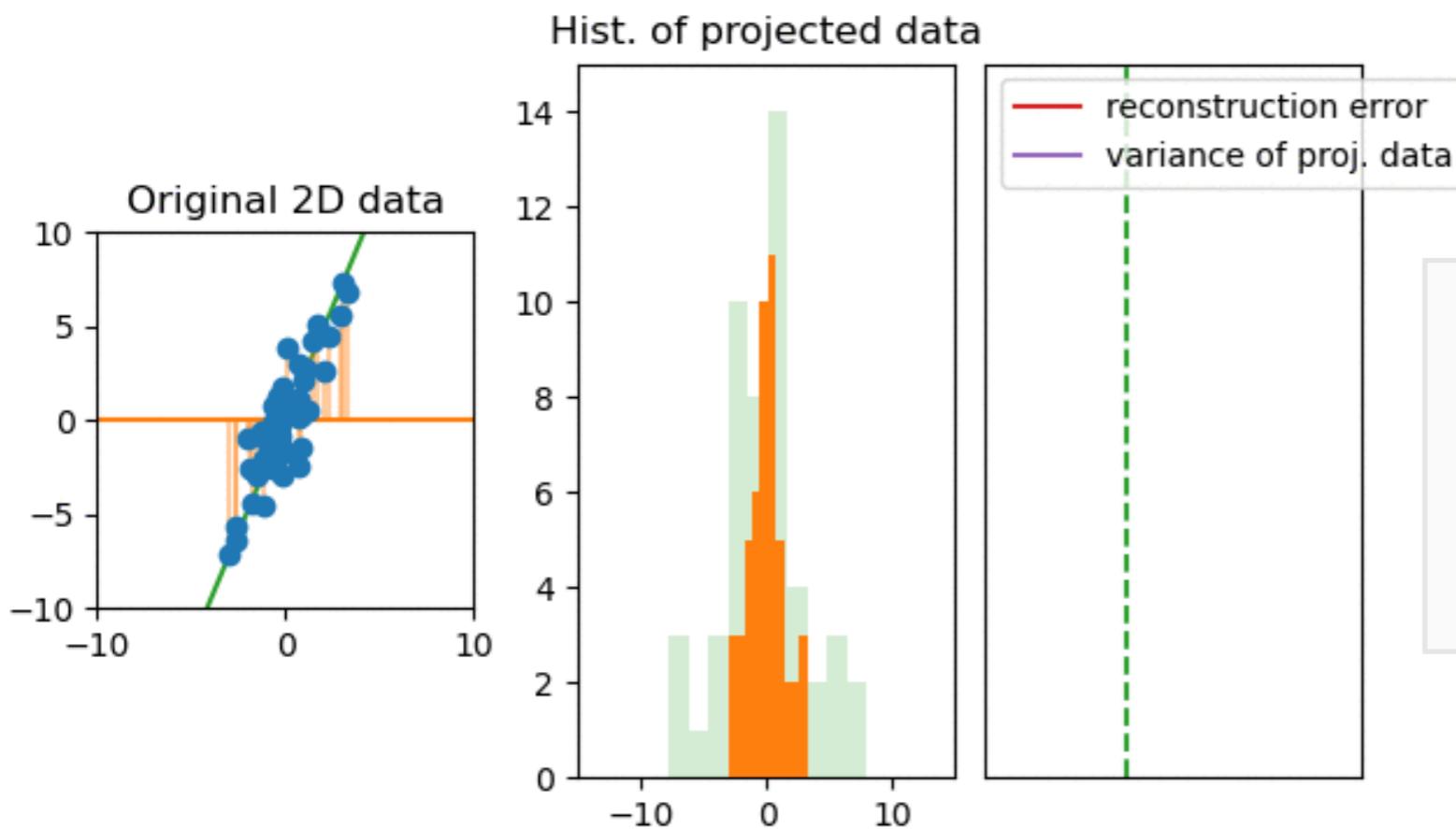
Dimension reduction



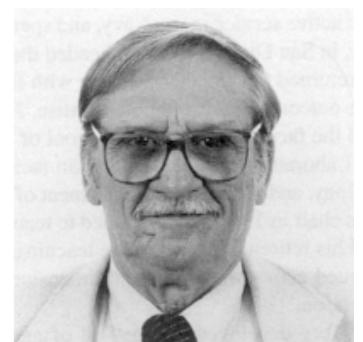
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◆ Principal components analysis



(Pearson, 1901)



(Torgerson, 1958)

Preserving the inner products

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i=1}^n \left(\langle \mathbf{x}_i, \mathbf{x}_j \rangle - \langle \mathbf{z}_i, \mathbf{z}_j \rangle \right)^2$$

$$Z \leftarrow \text{EVD}\left(\frac{1}{n} \mathbf{X} \mathbf{X}^\top\right)$$

Dimension reduction

♦ Spectral methods

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i=1}^n \left([C_X]_{ij} - \langle z_i, z_j \rangle \right)^2$$

Dimension reduction

♦ Spectral methods

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i=1}^n \left([C_X]_{ij} - \langle \mathbf{z}_i, \mathbf{z}_j \rangle \right)^2 \xrightarrow[\substack{C_X \succeq 0 \\ \text{solution} \\ (\text{Eckart \& Young, 1936})}]{} Z^\star = (\sqrt{\lambda_1} \mathbf{v}_1, \dots, \sqrt{\lambda_d} \mathbf{v}_d)^\top$$

λ_i i-th largest eigenvalue of C_X
with eigenvector \mathbf{v}_i

$$[C_X]_{ij} = \langle \phi(X_i), \phi(X_j) \rangle_{\mathcal{H}}$$

Dimension reduction

♦ Spectral methods

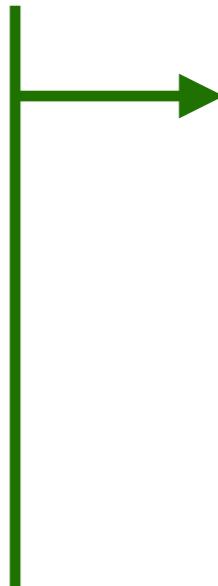
$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i=1}^n \left([C_X]_{ij} - \langle \mathbf{z}_i, \mathbf{z}_j \rangle \right)^2 \xrightarrow[\text{(Eckart & Young, 1936)}]{\begin{matrix} C_X \succeq 0 \\ \text{solution} \end{matrix}}$$

$Z^\star = (\sqrt{\lambda_1} \mathbf{v}_1, \dots, \sqrt{\lambda_d} \mathbf{v}_d)^\top$
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with eigenvector \mathbf{v}_i

♦ Kernel PCA $C_X \succeq 0$ $[C_X]_{ij} = \langle \phi(X_i), \phi(X_j) \rangle_{\mathcal{H}}$



(Schölkopf, 1997)



$$\text{PCA: } C_X = XX^\top \quad (Z \leftarrow \text{SVD}(X))$$

Dimension reduction

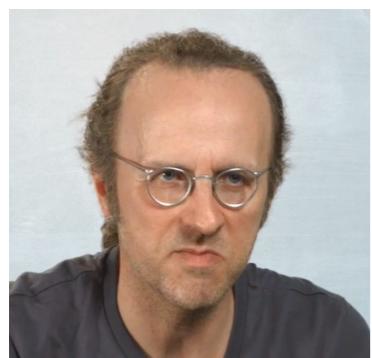
♦ Spectral methods

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i=1}^n \left([C_X]_{ij} - \langle \mathbf{z}_i, \mathbf{z}_j \rangle \right)^2 \quad \begin{matrix} C_X \succeq 0 \\ \text{solution} \\ (\text{Eckart \& Young, 1936}) \end{matrix}$$

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(Schölkopf, 1997)

→ PCA: $C_X = XX^\top$ ($Z \leftarrow \text{SVD}(X)$)
→ (classical) Multidimensional scaling: $C_X = -\frac{1}{2}HD_XH$

Dimension reduction

◆ Spectral methods

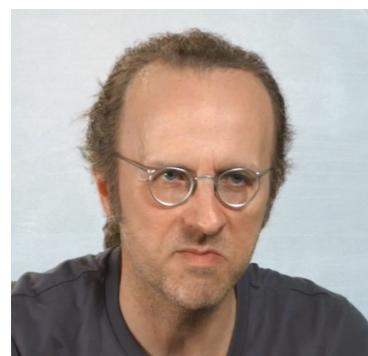
$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i=1}^n \left([C_X]_{ij} - \langle \mathbf{z}_i, \mathbf{z}_j \rangle \right)^2$$

C_X ≥ 0
solution →
(Eckart & Young, 1936)

$$Z^\star = (\sqrt{\lambda_1} \mathbf{v}_1, \dots, \sqrt{\lambda_d} \mathbf{v}_d)^\top$$

λ_i i-th largest eigenvalue of C_X
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◆ Kernel PCA $C_X \succeq 0$ $[C_X]_{ij} = \langle \phi(X_i), \phi(X_j) \rangle_{\mathcal{H}}$



(Schölkopf, 1997)

► PCA: $C_X = XX^\top$ ($Z \leftarrow \text{SVD}(X)$)

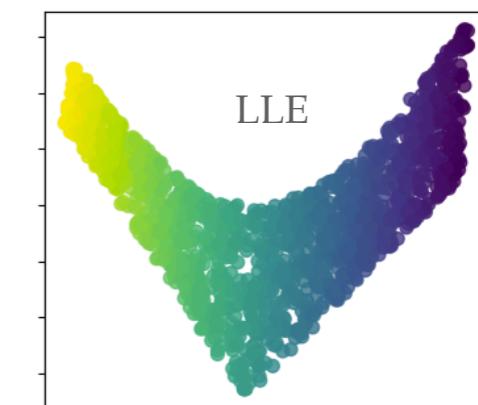
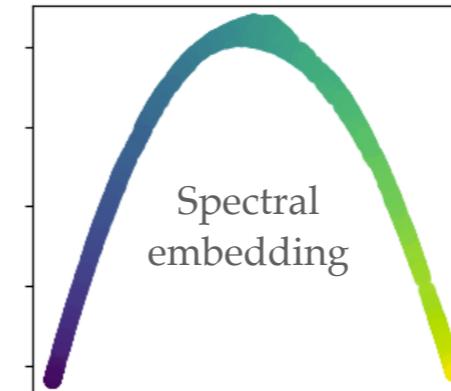
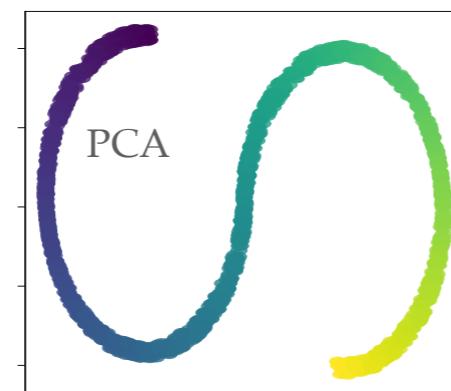
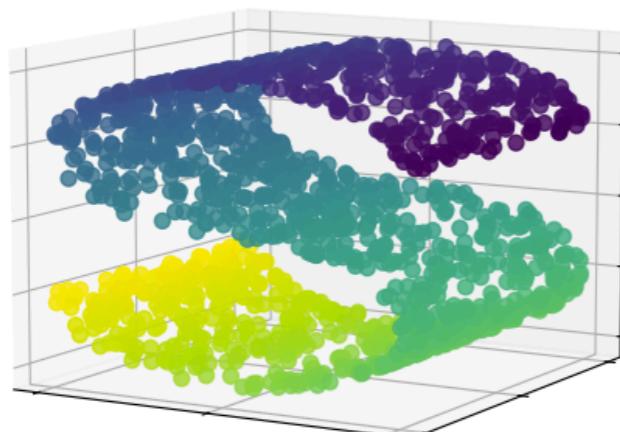
→ (classical) Multidimensional scaling: $C_X = -\frac{1}{\sigma} HD_X H$

- Laplacian Eigenmap (spectral embedding): $C_X = L_X^\dagger$
(Belkin & Niyogi, 2003)

► Locally Linear Embedding, Diffusion Map ...

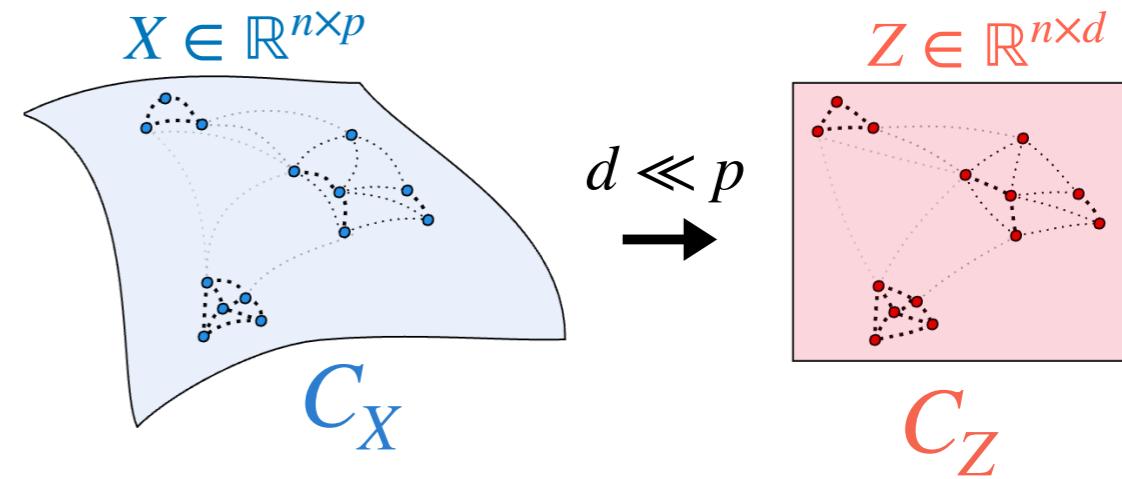
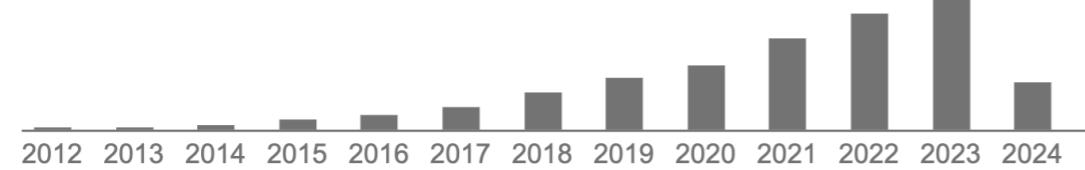
(Roweis & Saul, 2000)

(Coifman & Lafon, 2006)



Dimension reduction

Total citations Cited by 36223

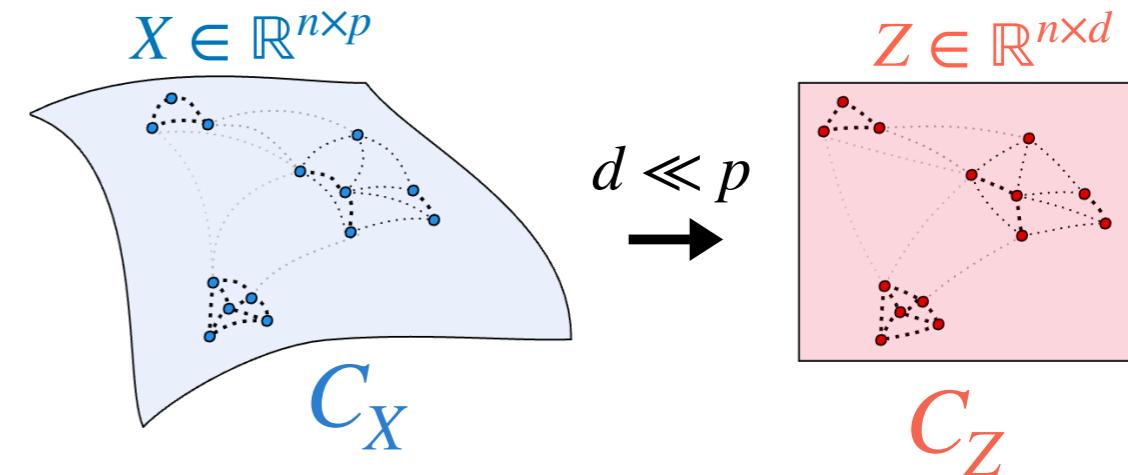
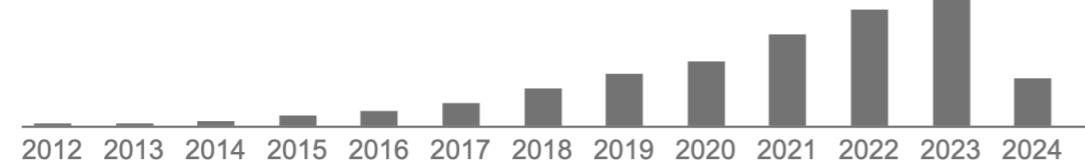


♦ Neighbor embedding methods

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n \text{KL} \left([C_X]_{ij}, [C_Z]_{ij} \right)$$

Dimension reduction

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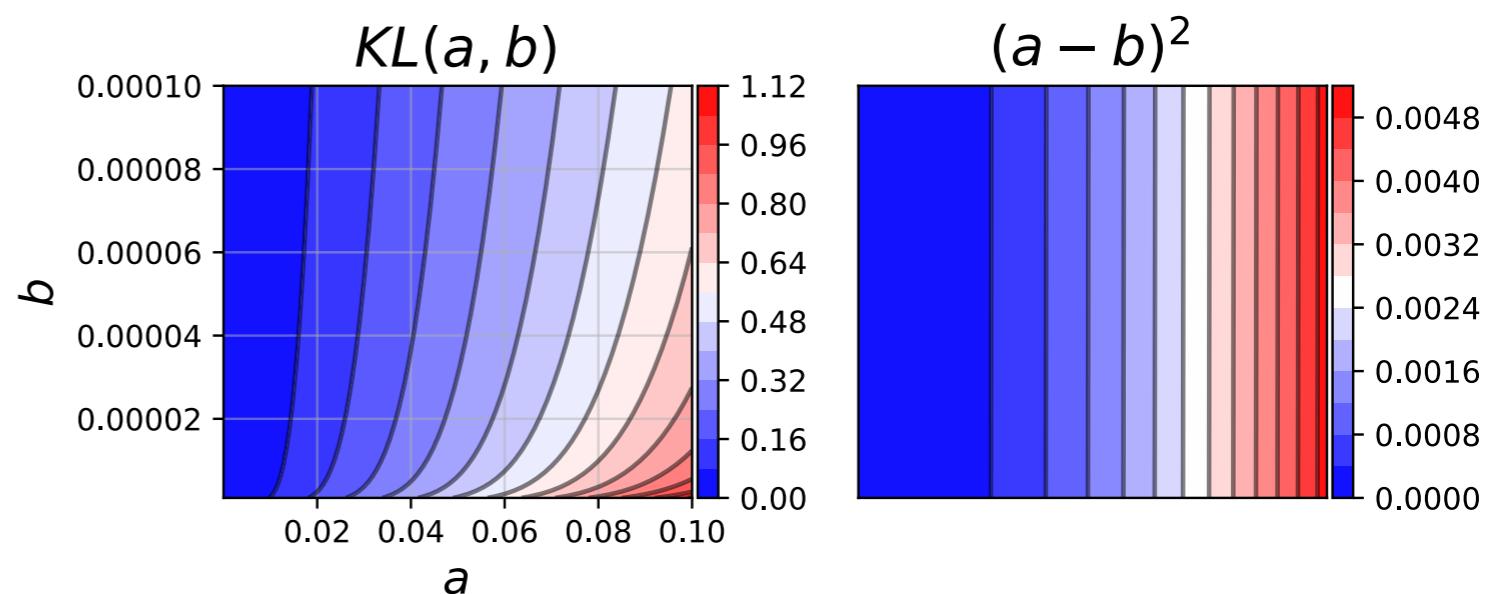
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$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n \text{KL} \left([C_X]_{ij}, [C_Z]_{ij} \right)$$

♦ Kullback-Leiber divergence

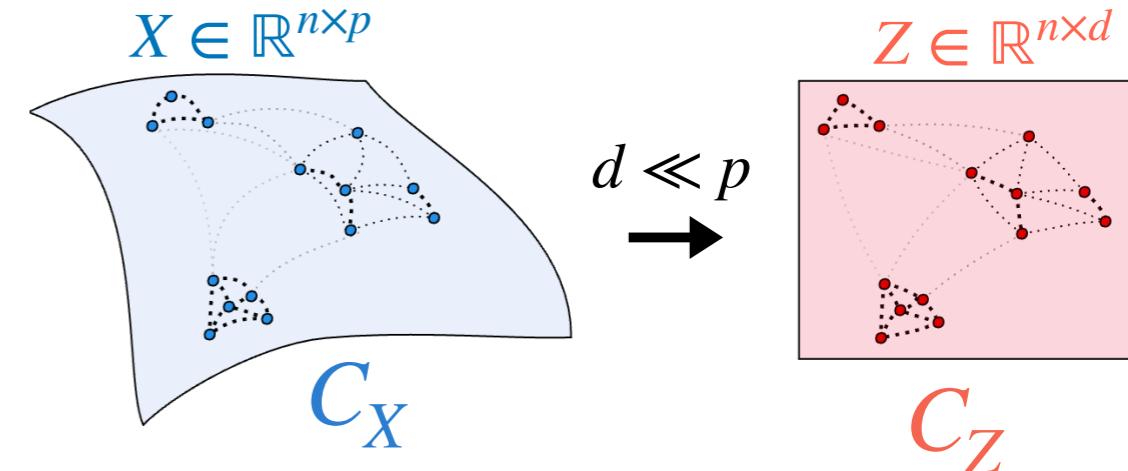
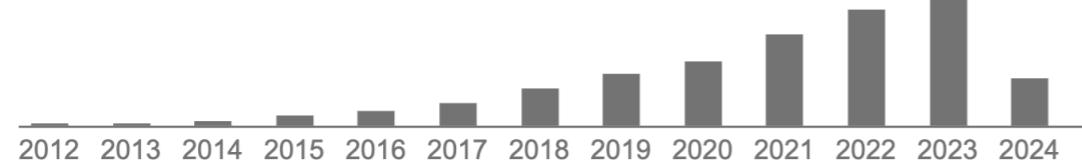
$$\text{KL}(a, b) = a \log(a/b) - a + b = D_\phi(a, b)$$

$$\text{Shannon–Boltzman entropy } \phi(x) = x \log(x) - x + 1$$



Dimension reduction

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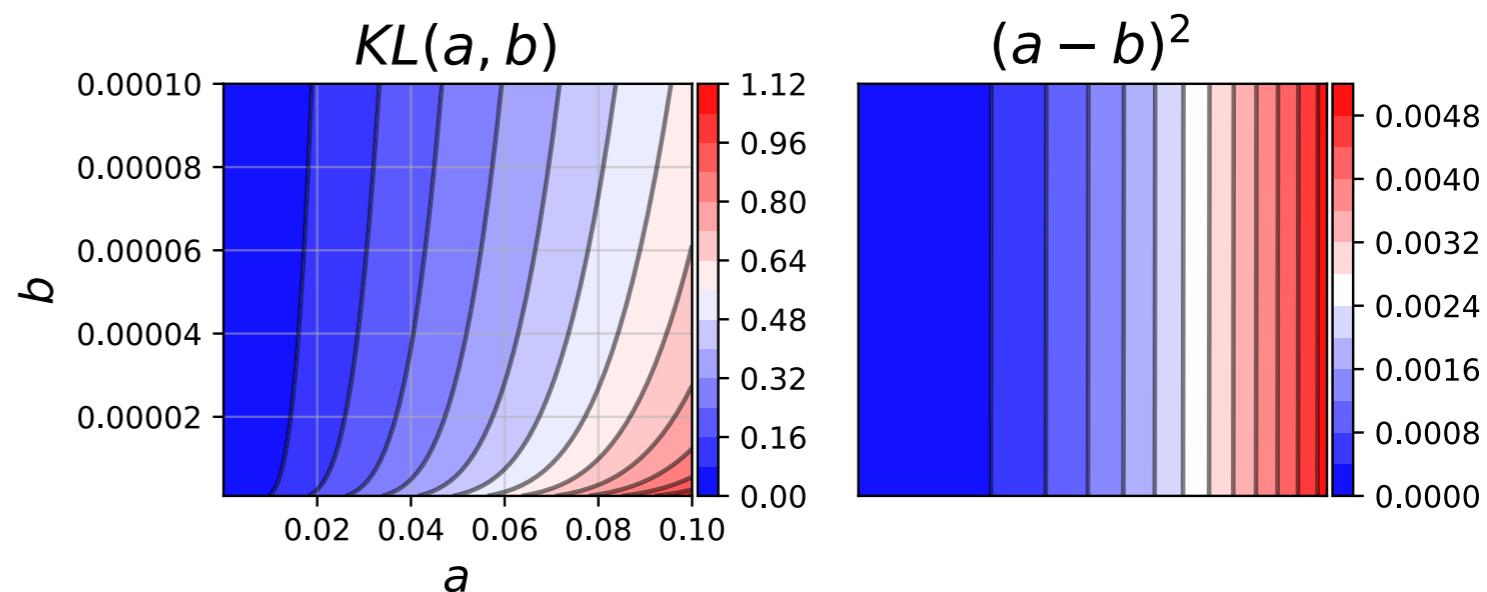
When $\sum_{i,j} [C_X]_{ij} = \sum_{i,j} [C_Z]_{ij}$ (same mass)

$$\sim \min_{Z \in \mathbb{R}^{n \times d}} \frac{\sum_{i,j=1}^n [C_X]_{ij} \log \left(\frac{[C_X]_{ij}}{[C_Z]_{ij}} \right)}{\text{KL}}$$

♦ Kullback-Leiber divergence

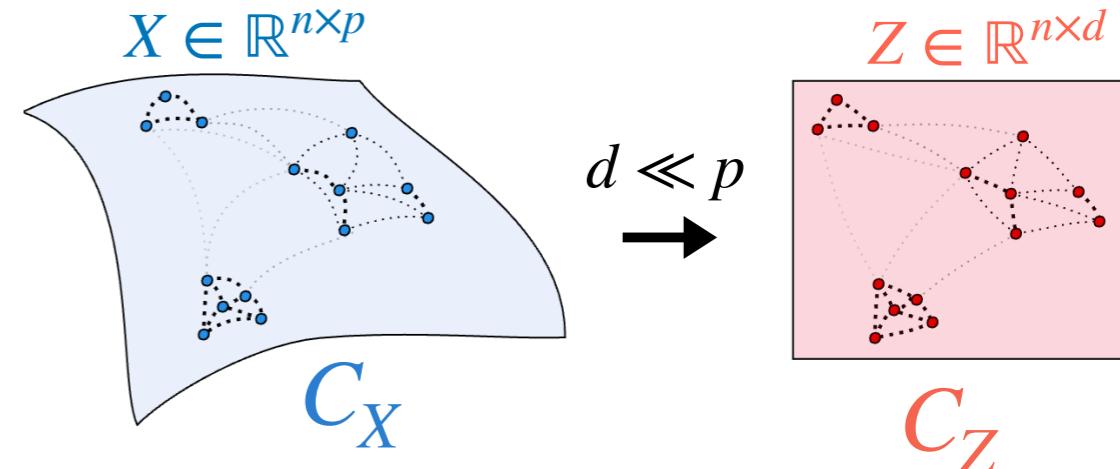
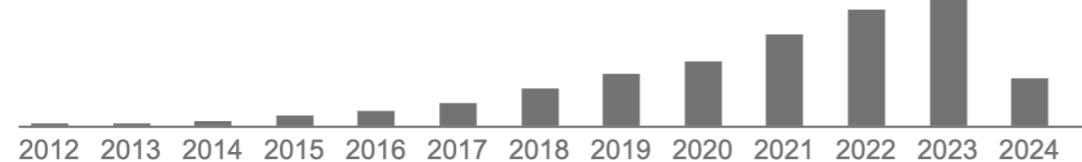
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Dimension reduction

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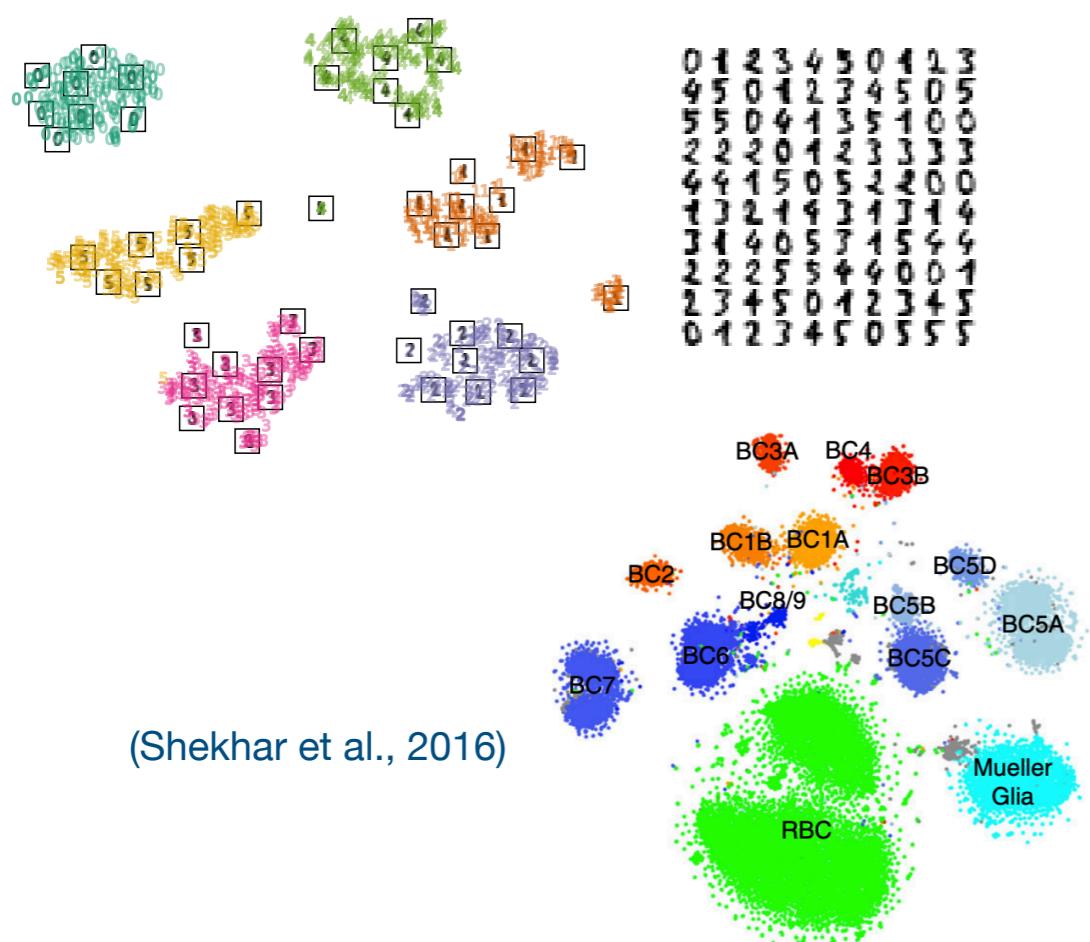


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$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n \text{KL}\left([\mathbf{C}_X]_{ij}, [\mathbf{C}_Z]_{ij}\right)$$

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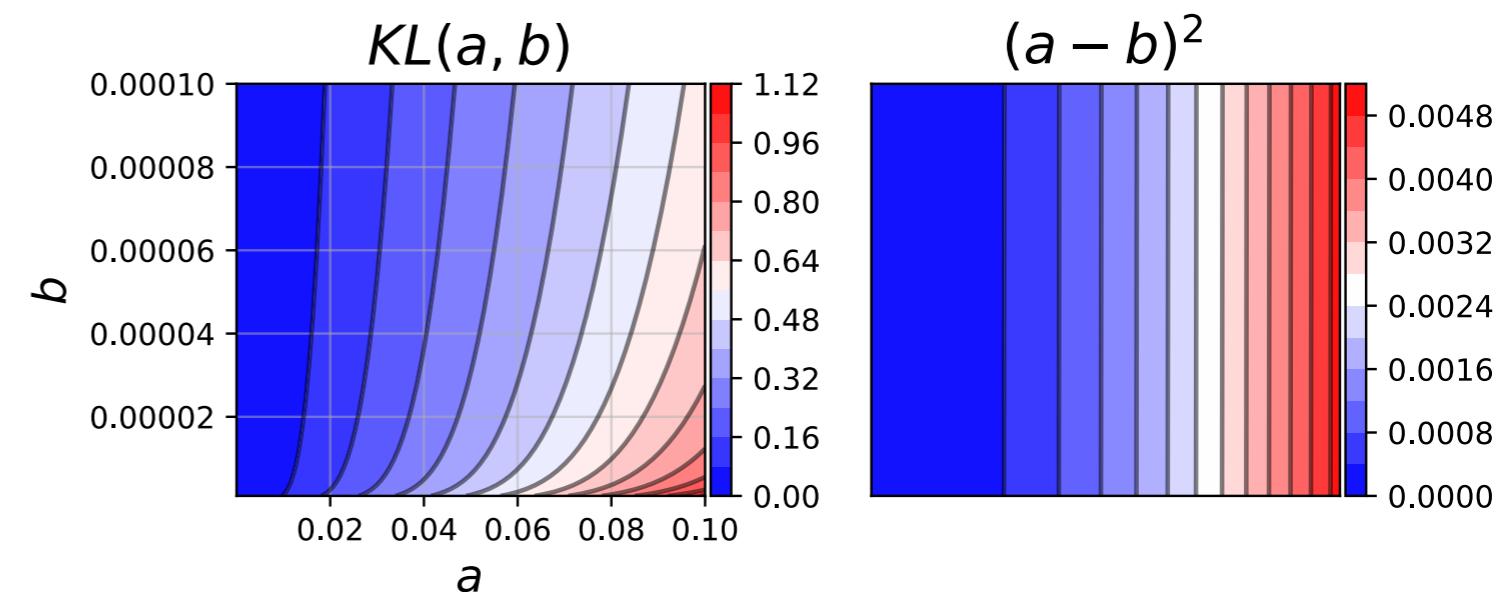
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♦ Kullback-Leiber divergence

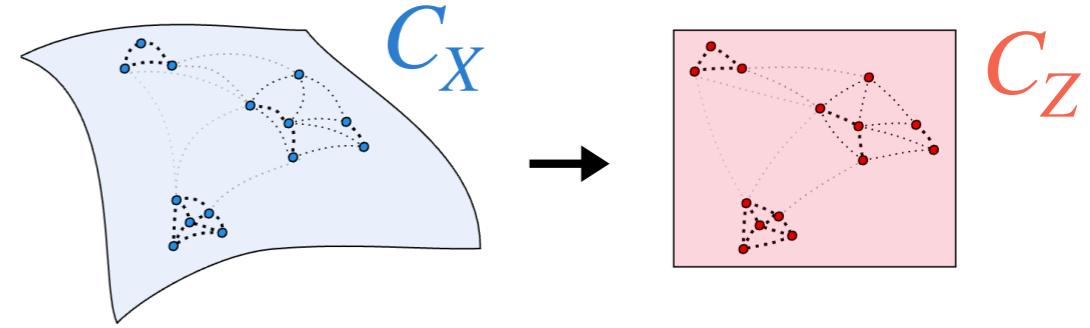
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Dimension reduction

◆ SNE (Hinton & Roweis, 2002)

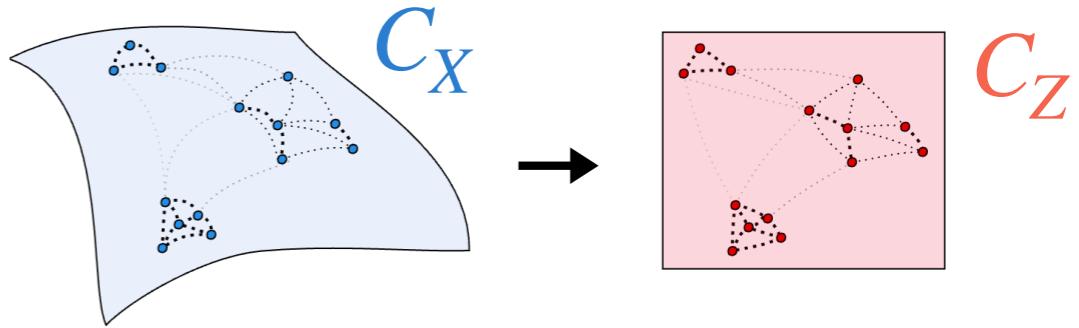


Embedding space

$$[C_Z]_{ij} = \frac{\exp(-\|\mathbf{z}_i - \mathbf{z}_j\|_2^2)}{\sum_k \exp(-\|\mathbf{z}_i - \mathbf{z}_k\|_2^2)} (= \mathbb{P}(j | i))$$

Dimension reduction

◆ SNE (Hinton & Roweis, 2002)



Input space

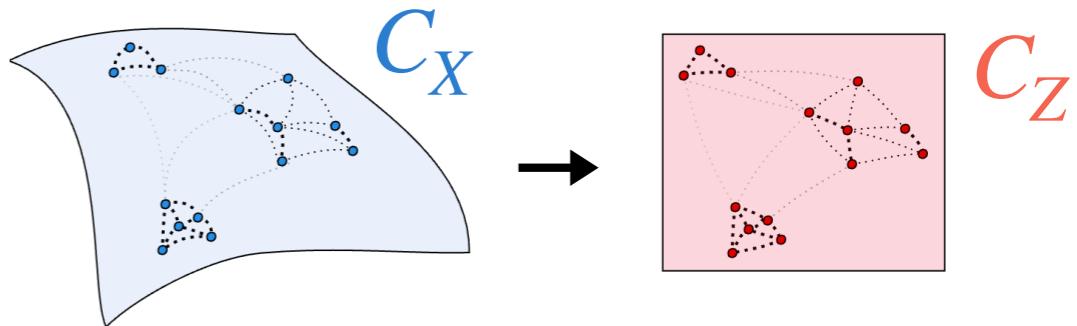
$$[C_X]_{ij} = \frac{\exp(-\|\mathbf{x}_i - \mathbf{x}_j\|_2^2 / 2\sigma_i^2)}{\sum_k \exp(-\|\mathbf{x}_i - \mathbf{x}_k\|_2^2 / 2\sigma_i^2)}$$

Embedding space

$$[C_Z]_{ij} = \frac{\exp(-\|\mathbf{z}_i - \mathbf{z}_j\|_2^2)}{\sum_k \exp(-\|\mathbf{z}_i - \mathbf{z}_k\|_2^2)} (= \mathbb{P}(j | i))$$

Dimension reduction

◆ SNE (Hinton & Roweis, 2002)



Input space

$$[C_X]_{ij} = \frac{\exp(-\|\mathbf{x}_i - \mathbf{x}_j\|_2^2 / 2\sigma_i^2)}{\sum_k \exp(-\|\mathbf{x}_i - \mathbf{x}_k\|_2^2 / 2\sigma_i^2)}$$

◆ Local bandwidths **optimized** s.t.

$$\forall i, \text{entropy}([C_X]_{i,:}) = \log(\text{perplexity})$$

◆ Perplexity = effective number of **neighbors**

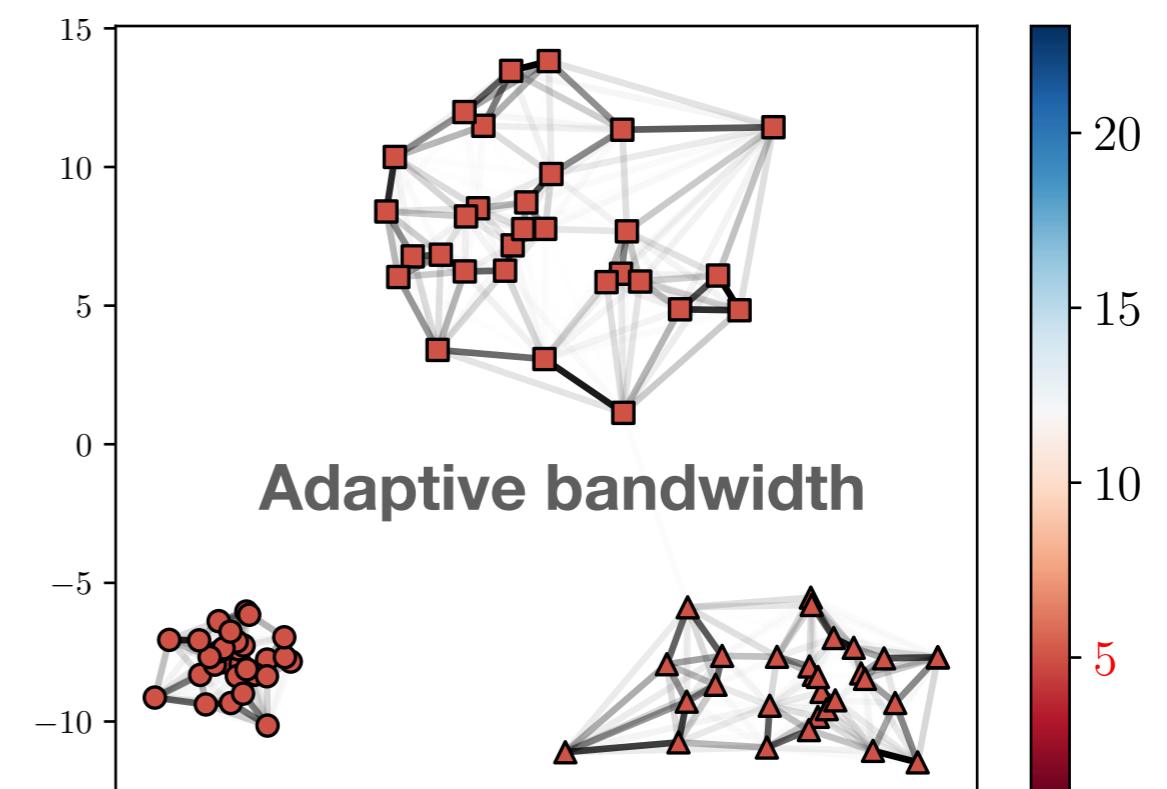
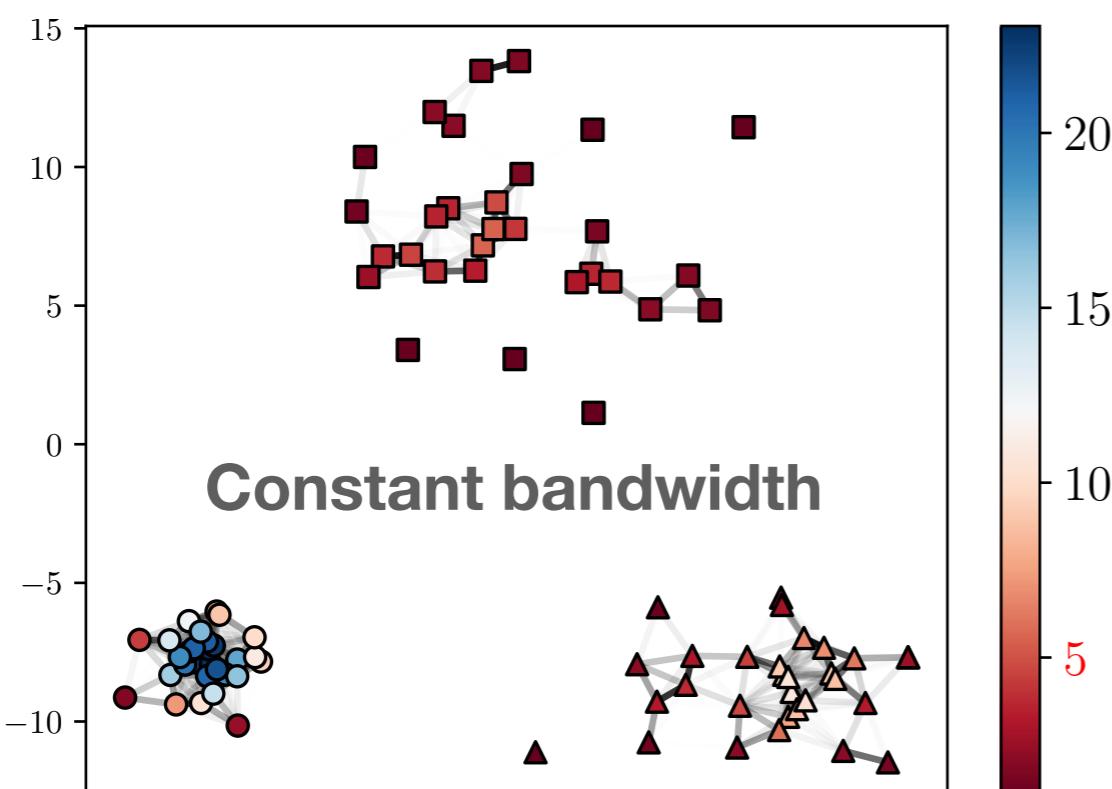
◆ Account for **varying density**

Embedding space

$$[C_Z]_{ij} = \frac{\exp(-\|\mathbf{z}_i - \mathbf{z}_j\|_2^2)}{\sum_k \exp(-\|\mathbf{z}_i - \mathbf{z}_k\|_2^2)} (= \mathbb{P}(j | i))$$

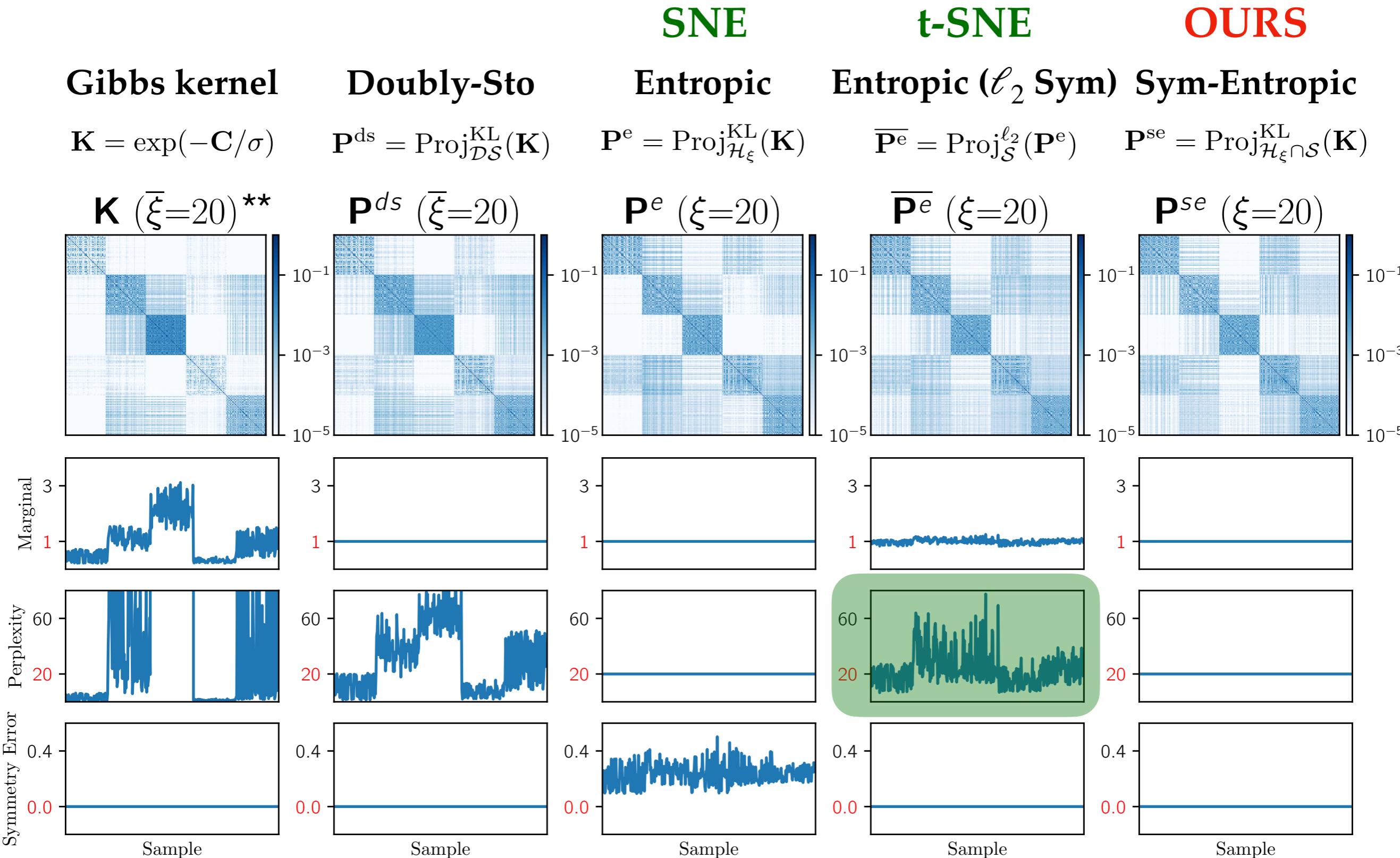
◆ (t)-SNE (Van der Maaten & Hinton, 2008)

◆ Crowding effect: Student t-distribution instead of Gaussian in Z



Dimension reduction

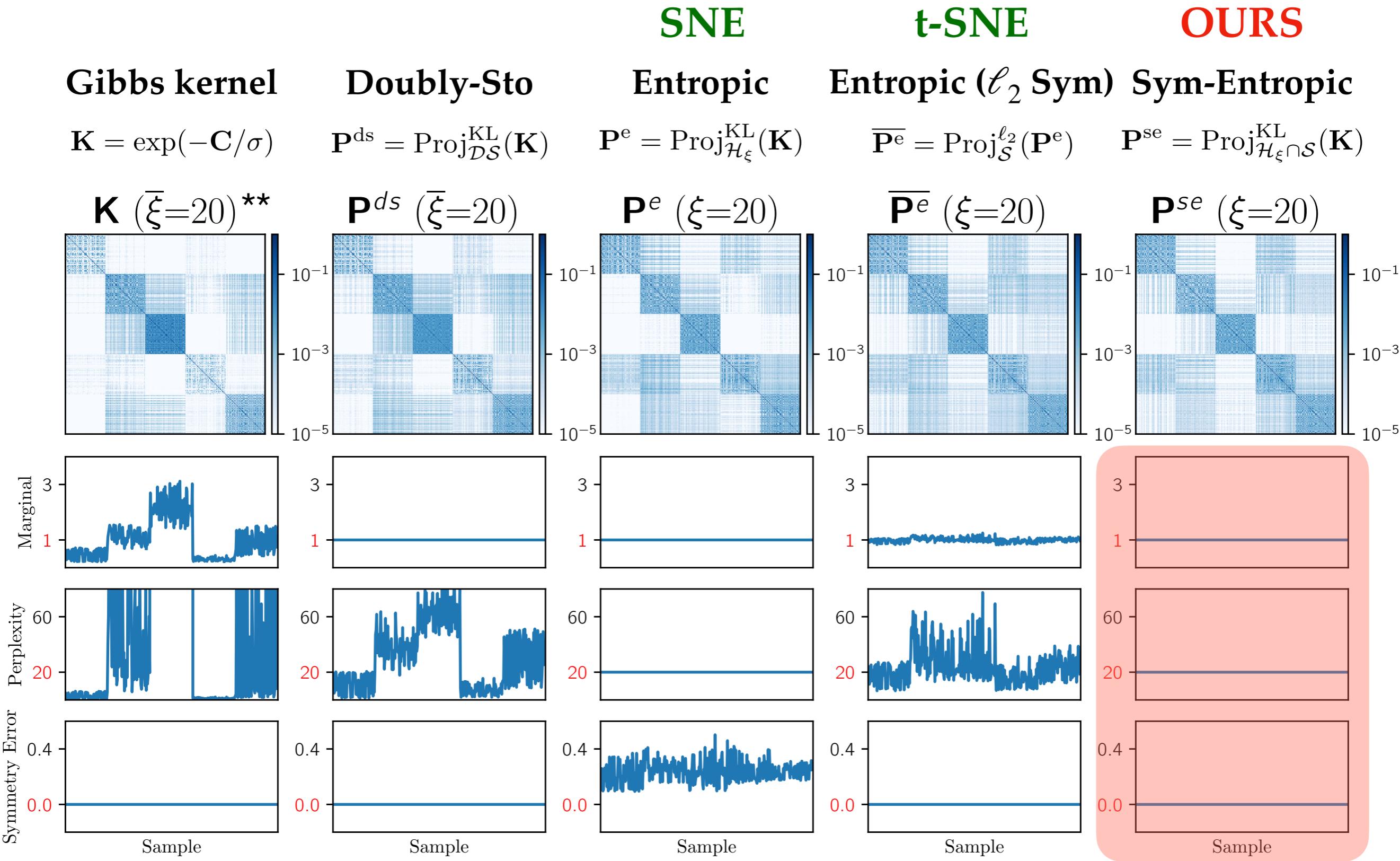
(Van Assel et al., 2023)



t-SNE fails at controlling the entropy when symmetrizing

Dimension reduction

(Van Assel et al., 2023)



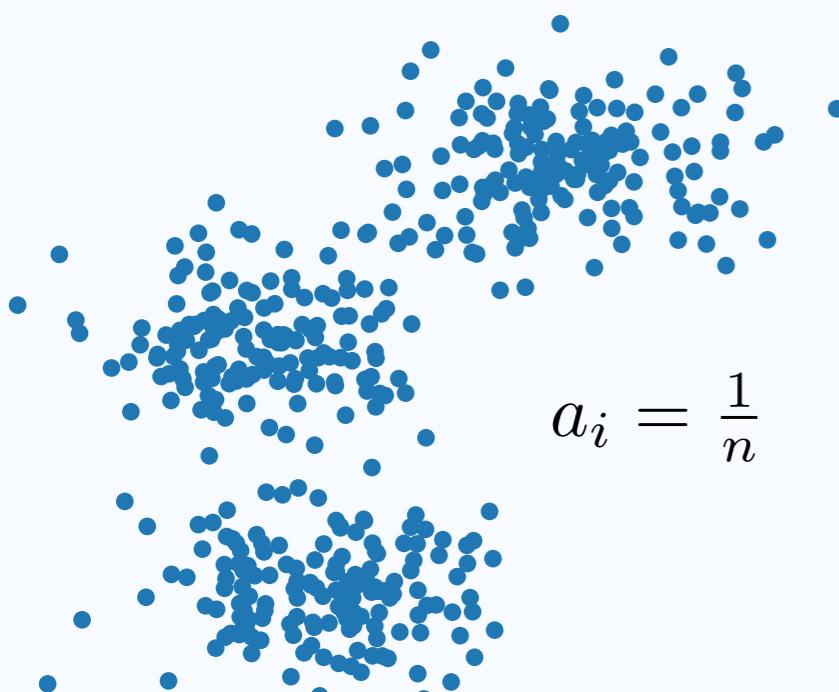
→ Controls ℓ_1 norm, entropy and symmetry at the same time.

Framing DR in terms of distributions

Measure and probability distributions are at the core of Machine learning.

Data: $(\mathbf{x}_i)_{i \in \llbracket n \rrbracket}$; $\mathbf{x}_i \in \mathbb{R}^d \longrightarrow$ A probability distribution describing the data

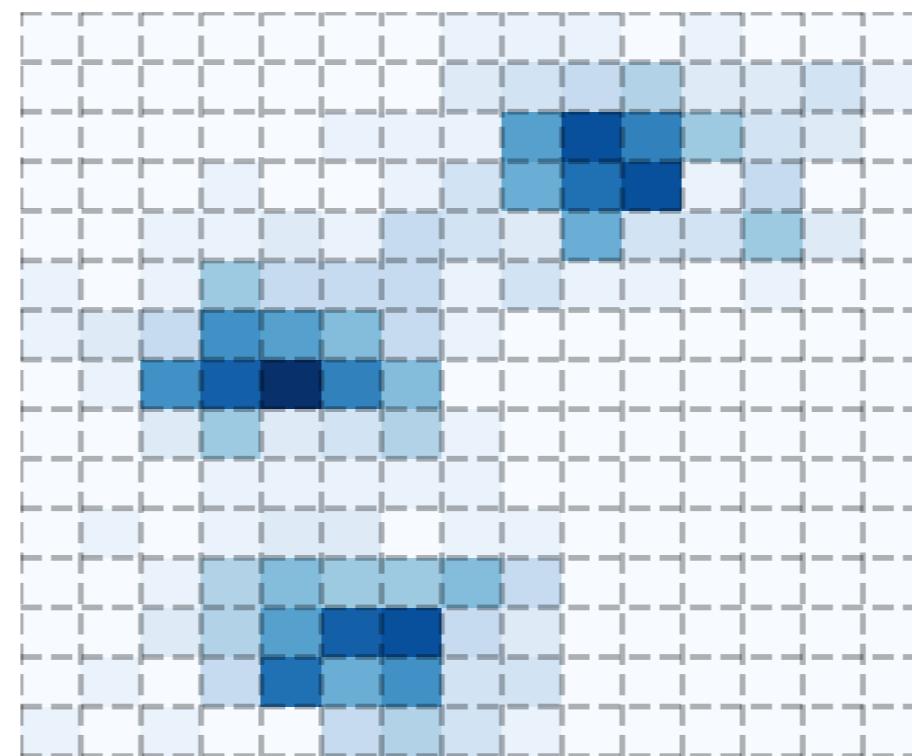
Lagrangian: $\sum_{i=1}^n a_i \delta_{x_i}$



(point clouds)

$$\delta_{\mathbf{x}_i}(\mathbf{x}) = 1 \text{ if } \mathbf{x} = \mathbf{x}_i \text{ else } 0$$

Eulerian: $\sum_{i=1}^N a_i \delta_{\hat{x}_i}$



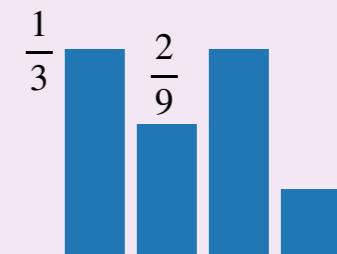
(histograms)

$$\hat{x}_i \text{ fixed position (grid)}$$

Probability simplex

$$\mathbf{a} = (a_i)_{i \in \llbracket n \rrbracket} \in \Sigma_n$$

$$a_i \geq 0, \sum_{i=1}^n a_i = 1$$



Framing DR in terms of distributions

Measure and probability distributions are at the core of Machine learning.

A point of view on the data

Data: $(\mathbf{x}_i)_{i \in [\![n]\!]} ; \mathbf{x}_i \in \mathbb{R}^d \longrightarrow$ A probability distribution describing the data

A formalism for many machine learning paradigms

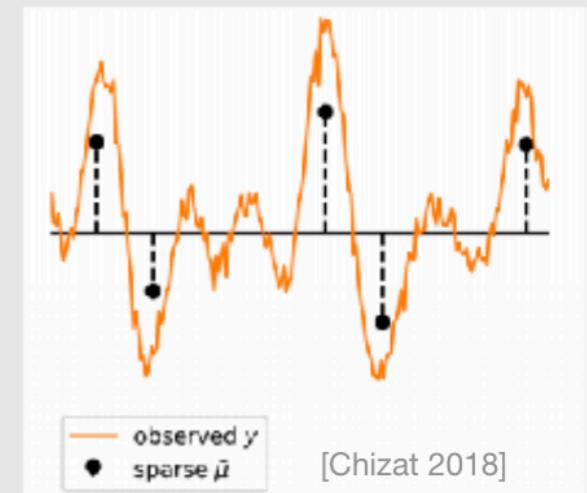
$$\begin{array}{|c} \text{(ERM)} \\ \hline \end{array} \quad \min_f \mathbb{E}_{(x,y) \sim \mu} [L(f(x), y)] \quad \mu = \frac{1}{n} \sum_{i=1}^n \delta_{(\mathbf{x}_i, y_i)}$$

$$\begin{array}{|c} \text{(Likelihood)} \\ \hline \end{array} \quad \max_{\theta \in \Theta} \mathbb{E}_{\mathbf{x} \sim \mu} [\log(\mathbb{P}_{\theta}(\mathbf{x}))] \quad \mu = \frac{1}{n} \sum_{i=1}^n \delta_{\mathbf{x}_i}$$

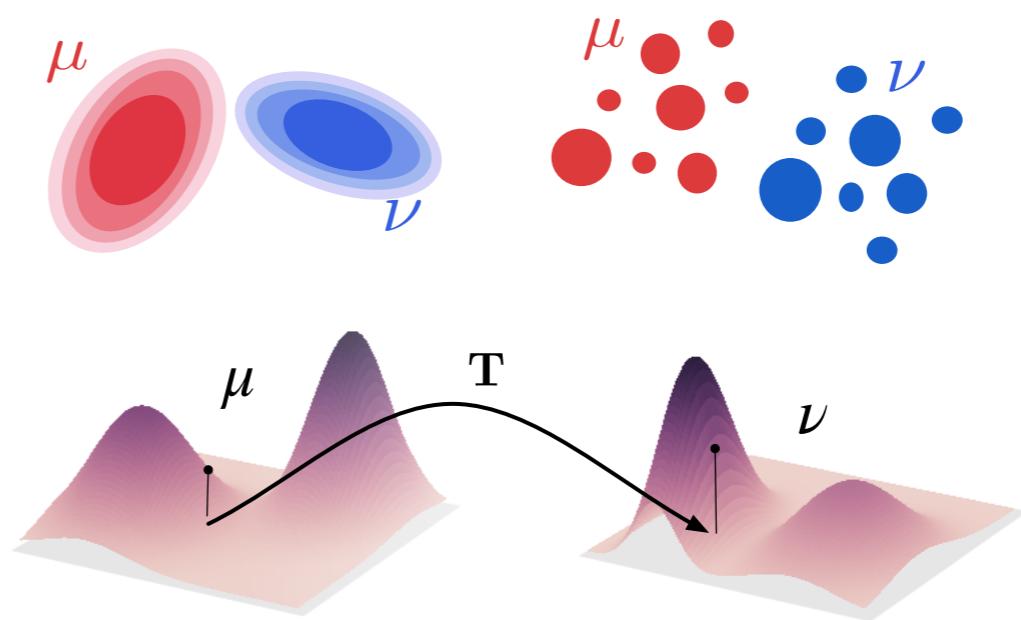
$$\begin{array}{|c} \text{(GAN)} \\ \hline \end{array} \quad \min_{\theta \in \Theta} D(\mu_{\theta}, \nu)$$

(Signal processing) Recover a sparse signal

$$\min_{\mu \in \mathcal{M}(\Theta)} \frac{1}{2} \|\mathbf{y} - \phi * \mu\|_{L^2}^2 + R(\mu) \quad \bar{\mu} = \sum_i w_i \delta_{\theta_i}$$



From linear Optimal Transport to Gromov-Wasserstein



From Wasserstein to Gromov-Wasserstein

♦ Classical optimal transport (in a nutshell)

Kantorovitch Formulation

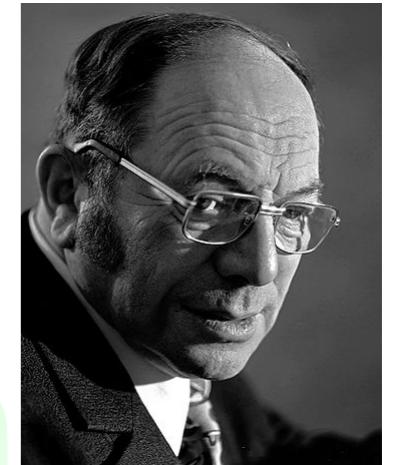
Two probability distributions

$$\mu \in \mathcal{P}(\mathcal{X})$$

$$\nu \in \mathcal{P}(\mathcal{Y})$$

A cost function

$$c(x, y) : \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}$$



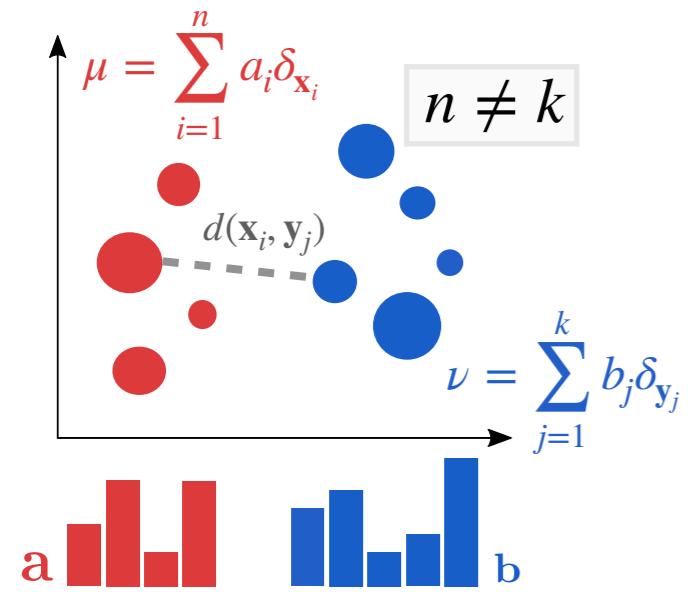
$\mu \in \mathcal{P}(\mathcal{X})$ is transported to $\nu \in \mathcal{P}(\mathcal{Y})$ by a **transport plan** $T \in \mathcal{P}(\mathcal{X} \times \mathcal{Y})$

We want to find the plan that **minimizes the overall cost** of moving all the points.

$$\inf_{T \in \Pi(\mu, \nu)} \int c(x, y) dT(x, y)$$

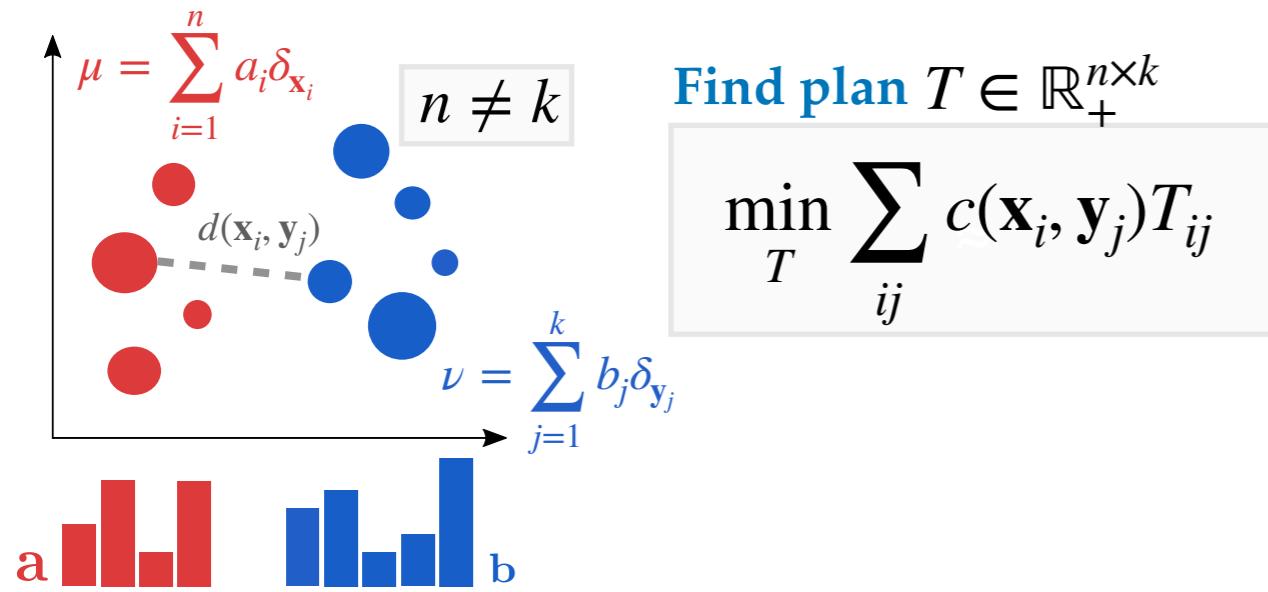
From Wasserstein to Gromov-Wasserstein

♦ Classical optimal transport (in a nutshell)



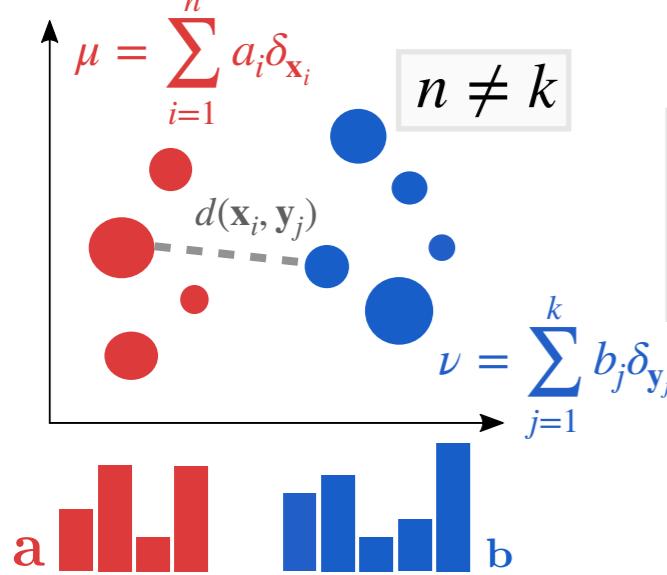
From Wasserstein to Gromov-Wasserstein

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From Wasserstein to Gromov-Wasserstein

♦ Classical optimal transport (in a nutshell)



Find plan $T \in \mathbb{R}_+^{n \times k}$

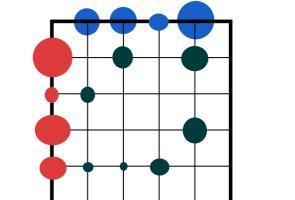
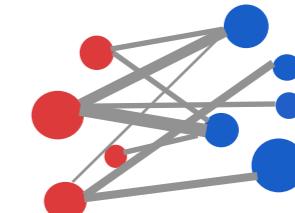
$$\min_T \sum_{ij} c(x_i, y_j) T_{ij}$$

which constraints ?

Coupling $\Pi(a, b)$

$$T^\top 1_n = a$$

$$T 1_k = b$$



Bakeries = quantity of breads

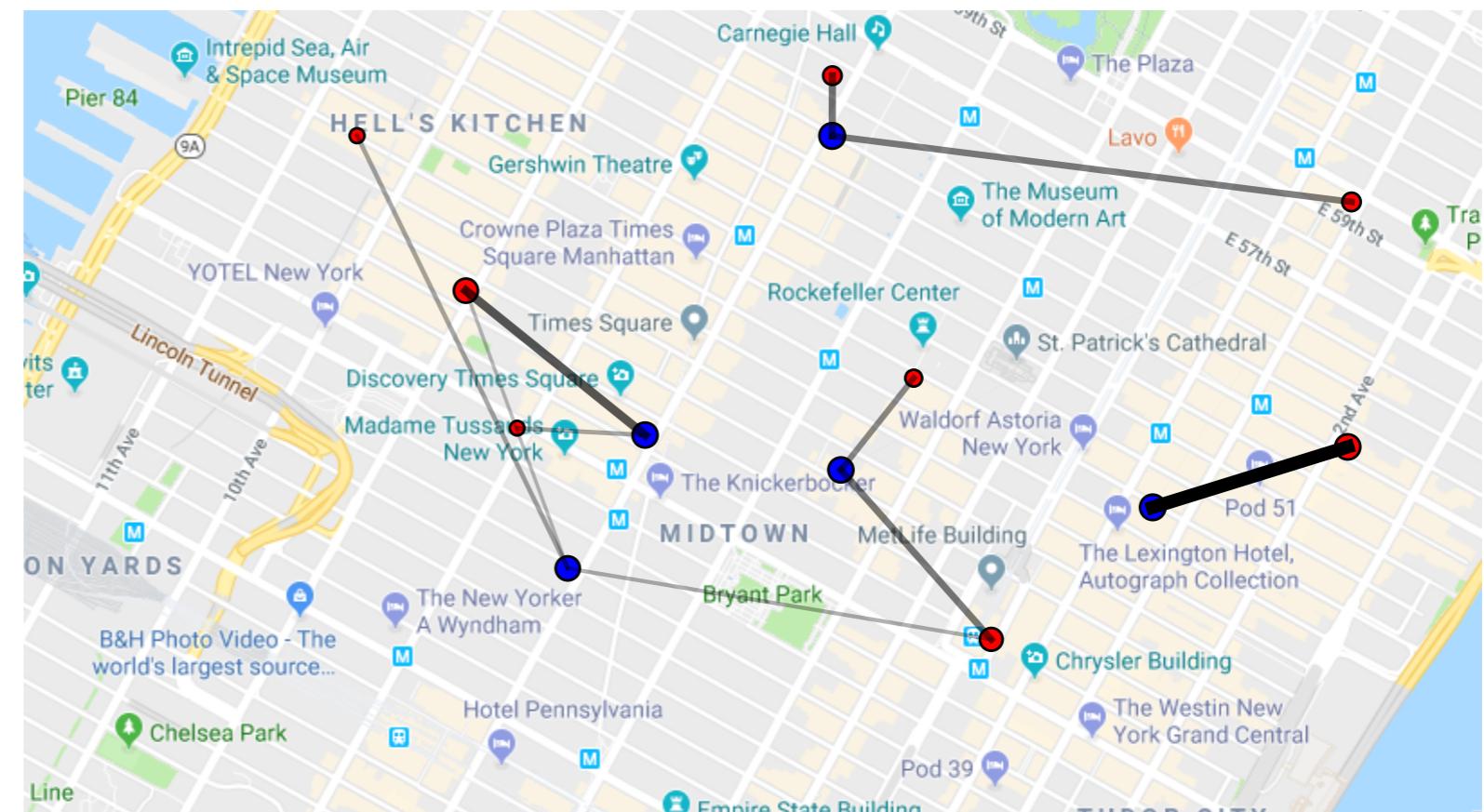
loc: \mathbf{x}_i quantity: a_i

Cafés = demand of breads

loc: \mathbf{y}_j demand: b_j

Distance between bakeries and cafés

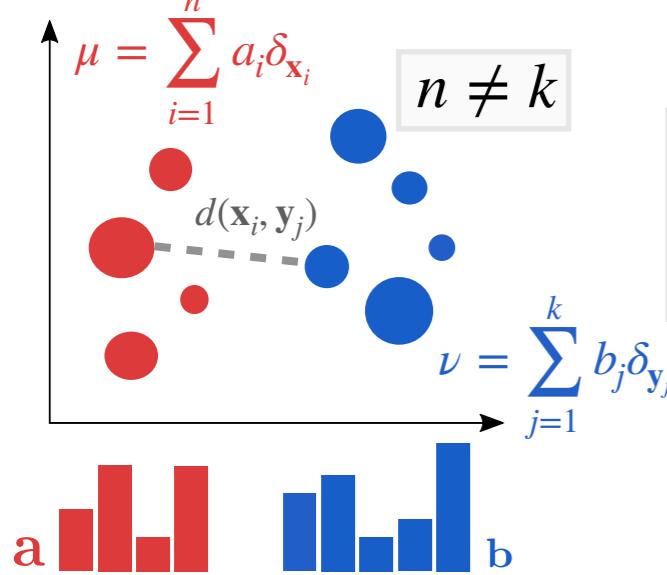
$$c(\mathbf{x}_i, \mathbf{y}_j)$$



We want to route all the breads from bakeries to cafés the cheapest way

From Wasserstein to Gromov-Wasserstein

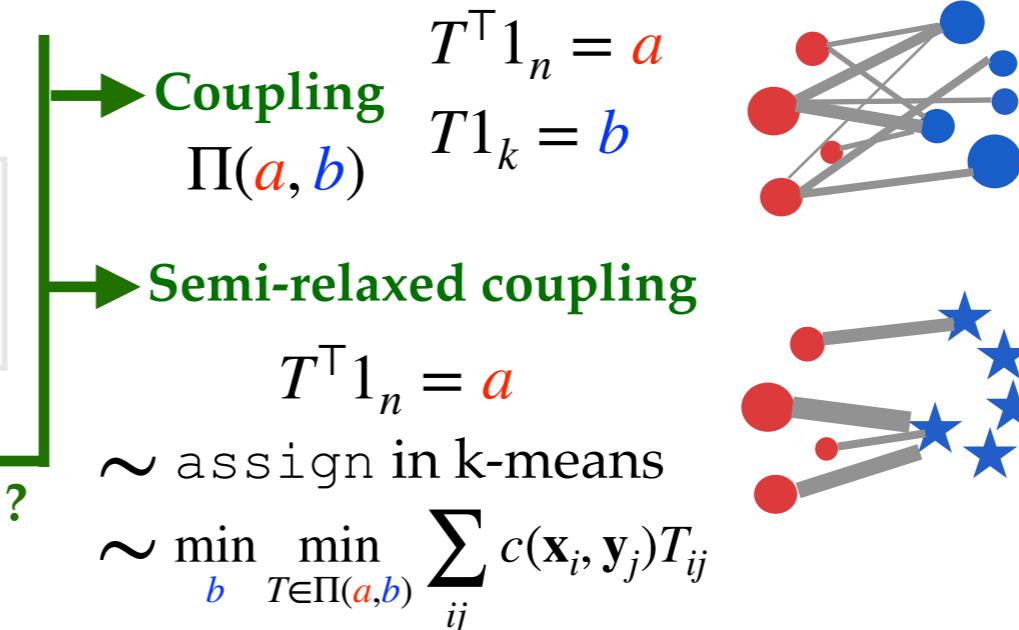
♦ Classical optimal transport (in a nutshell)



Find plan $T \in \mathbb{R}_+^{n \times k}$

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Bakeries = quantity of breads

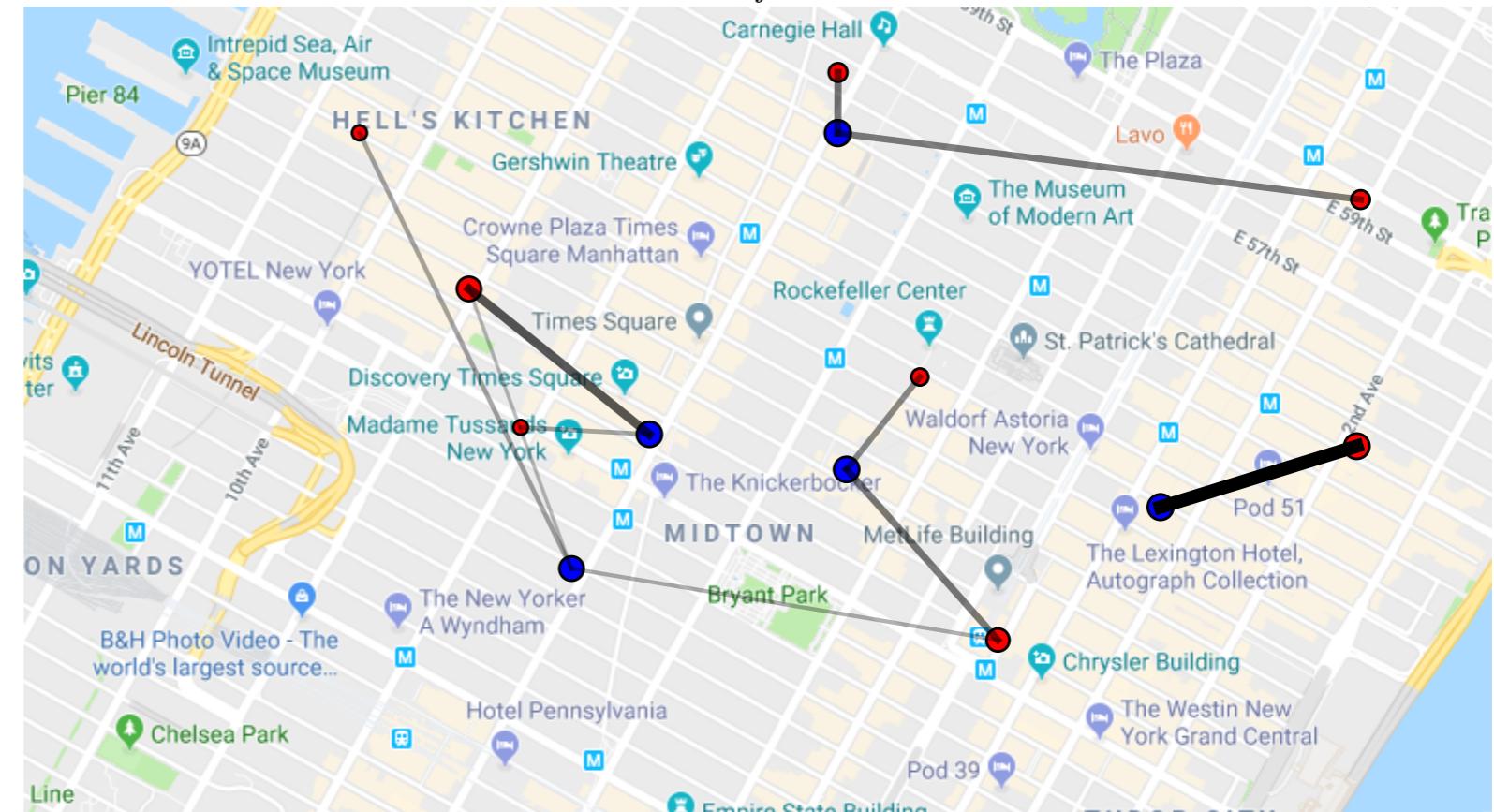
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Distance between bakeries and cafés

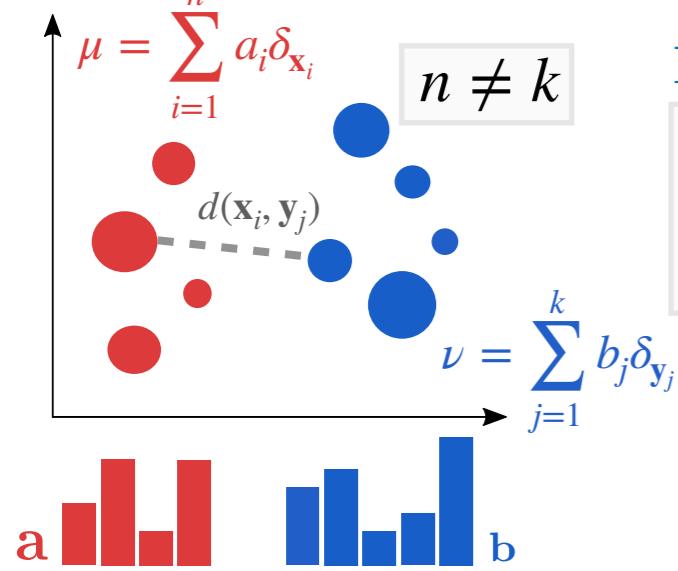
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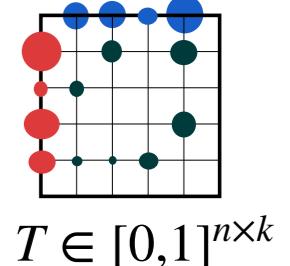
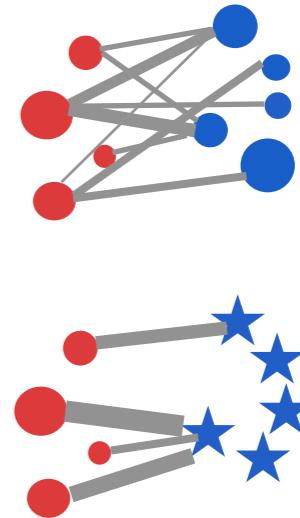
$$T \mathbf{1}_k = b$$

→ **Semi-relaxed coupling**

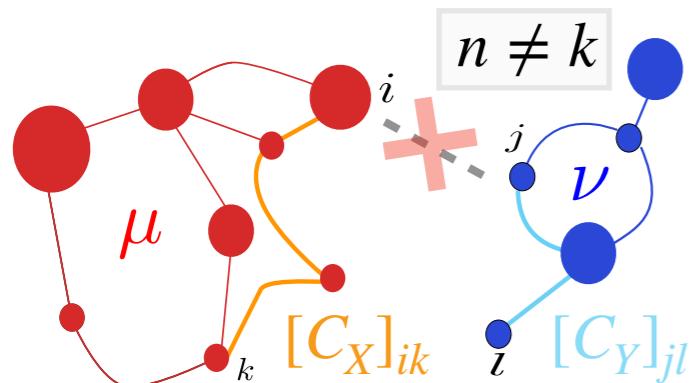
$$T^\top \mathbf{1}_n = a$$

~ assign in k-means

~ $\min_b \min_{T \in \Pi(a, b)} \sum_{ij} c(\mathbf{x}_i, \mathbf{y}_j) T_{ij}$

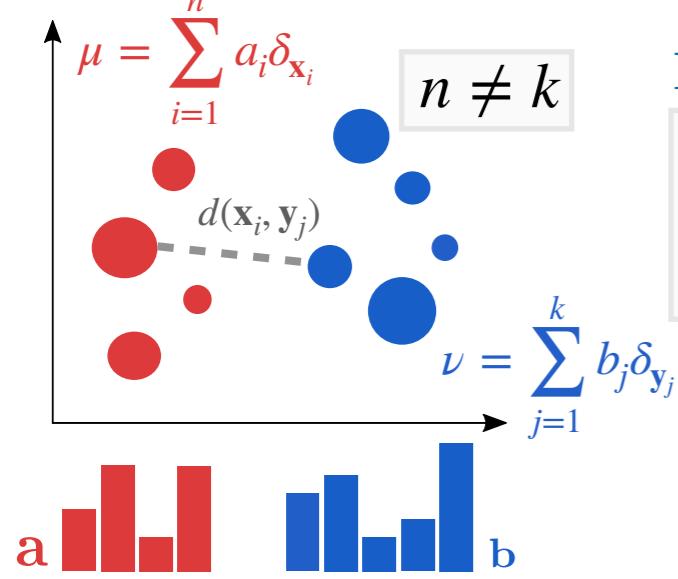


♦ Gromov-Wasserstein



From Wasserstein to Gromov-Wasserstein

♦ Classical optimal transport (in a nutshell)



Find plan $T \in \mathbb{R}_+^{n \times k}$

$$\min_T \sum_{ij} c(\mathbf{x}_i, \mathbf{y}_j) T_{ij}$$

which constraints ?

→ **Coupling**
 $\Pi(a, b)$

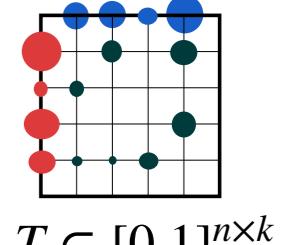
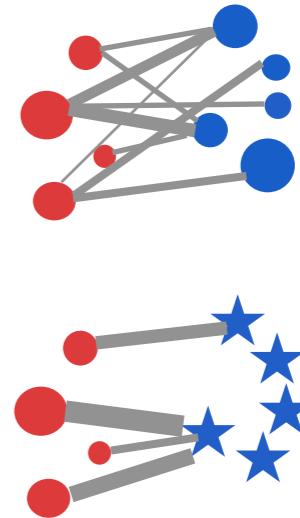
$$T^\top \mathbf{1}_n = a$$

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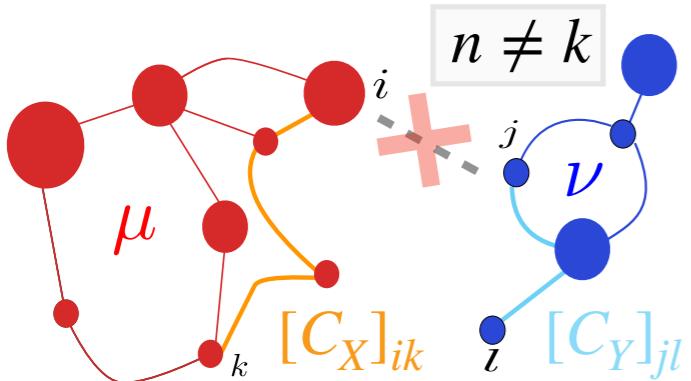
→ **Semi-relaxed coupling**

$$T^\top \mathbf{1}_n = a$$

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♦ Gromov-Wasserstein

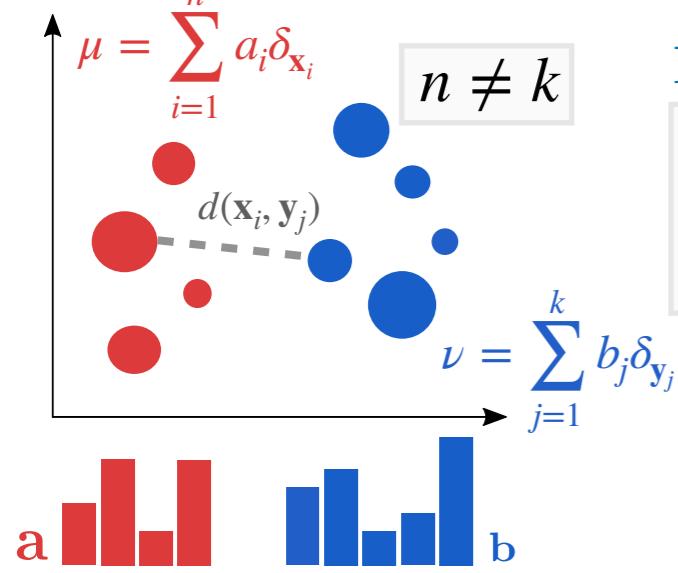


Quadratic OT: find the plan

$$\min_{T \in \Pi(a, b)} \sum_{ijkl} L([C_X]_{ik}, [C_Y]_{jl}) T_{ij} T_{kl}$$

From Wasserstein to Gromov-Wasserstein

♦ Classical optimal transport (in a nutshell)



Find plan $T \in \mathbb{R}_+^{n \times k}$

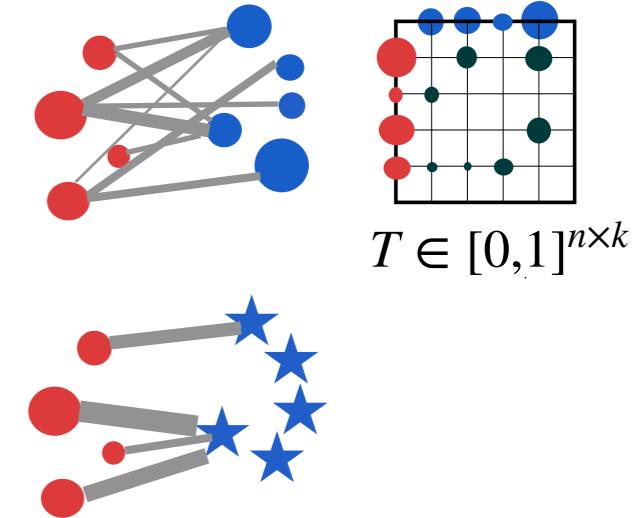
$$\min_T \sum_{ij} c(\mathbf{x}_i, \mathbf{y}_j) T_{ij}$$

which constraints ?

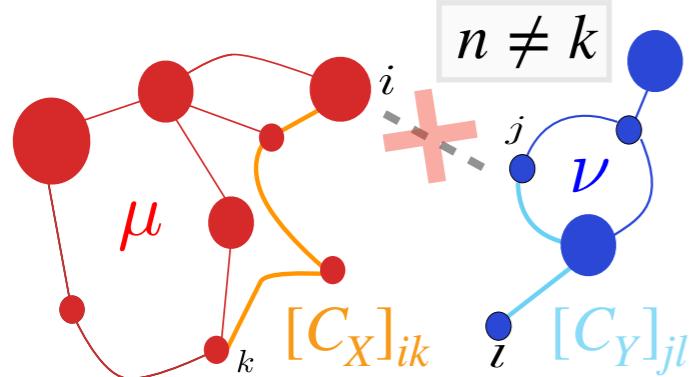
→ **Coupling** $\Pi(a, b)$ $T^\top \mathbf{1}_n = a$
 $T \mathbf{1}_k = b$

→ **Semi-relaxed coupling**

$T^\top \mathbf{1}_n = a$
~ assign in k-means
~ $\min_b \min_{T \in \Pi(a, b)} \sum_{ij} c(\mathbf{x}_i, \mathbf{y}_j) T_{ij}$



♦ Gromov-Wasserstein



♦ L measures distortion

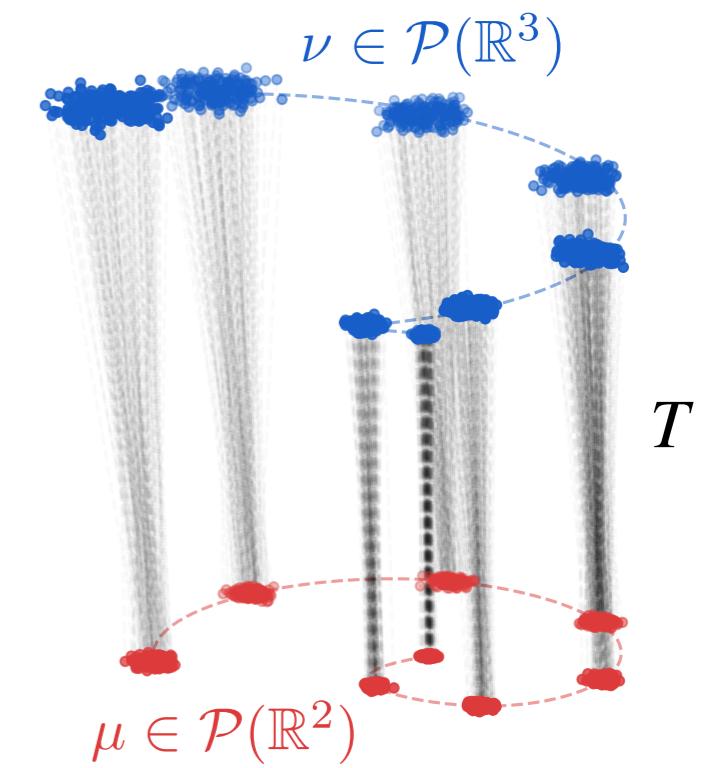
$$\left| [C_X]_{ik} - [C_Y]_{jl} \right|^2$$

♦ Goal : preserving pairwise connectivity

- ♦ Distance w.r.t. isomorphisms
- ♦ Difficult quadratic problem (NP-hard)

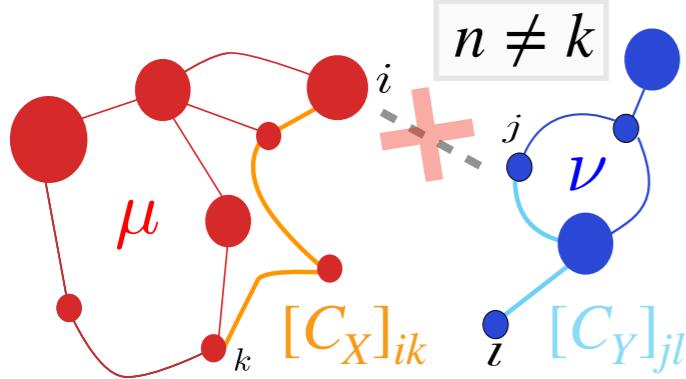
Quadratic OT: find the plan

$$\min_{T \in \Pi(a, b)} \sum_{ijkl} L([C_X]_{ik}, [C_Y]_{jl}) T_{ij} T_{kl}$$



From Wasserstein to Gromov-Wasserstein

♦ Gromov-Wasserstein



(Sturm, 2012) (Mémoli, 2011)

♦ L measures distortion

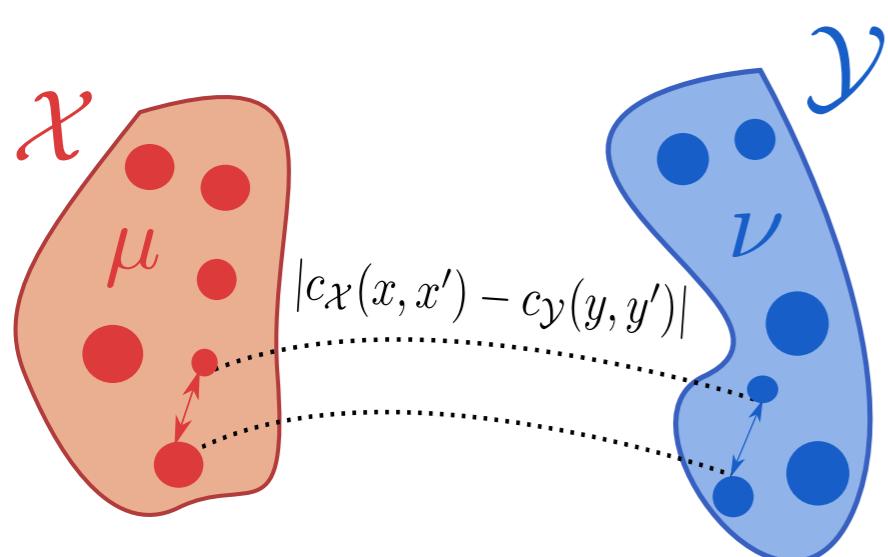
$$\left| [C_X]_{ik} - [C_Y]_{jl} \right|^2$$

♦ Goal : preserving pairwise connectivity

♦ Non-convex quadratic problem (NP-hard)

Quadratic OT: find the plan

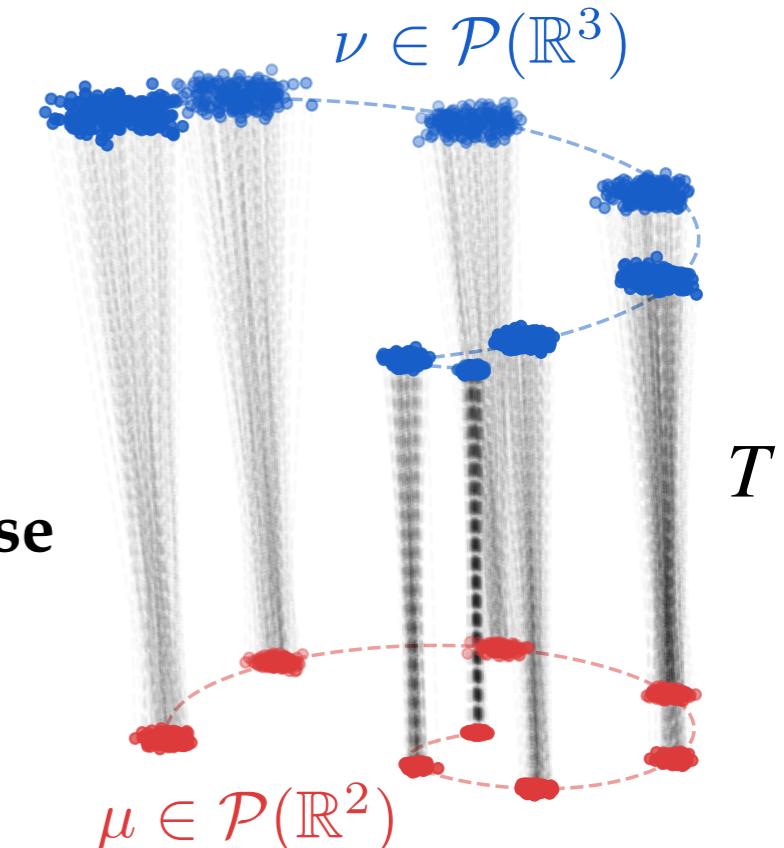
$$\min_{T \in \Pi(\mathbf{a}, \mathbf{b})} \sum_{ijkl} L \left([C_X]_{ik}, [C_Y]_{jl} \right) T_{ij} T_{kl}$$



♦ Distance w.r.t. isomorphisms, on the space of metric measure spaces

$$\mathbb{X} = (\mathcal{X}, c_{\mathcal{X}}, \mu \in \mathcal{P}(\mathcal{X}))$$

$$\mathbb{Y} = (\mathcal{Y}, c_{\mathcal{Y}}, \nu \in \mathcal{P}(\mathcal{Y}))$$



$$GW(\mathbb{X}, \mathbb{Y}) = 0 \text{ iff}$$

$$\exists \phi : \mathcal{X} \rightarrow \mathcal{Y}$$

Isometry

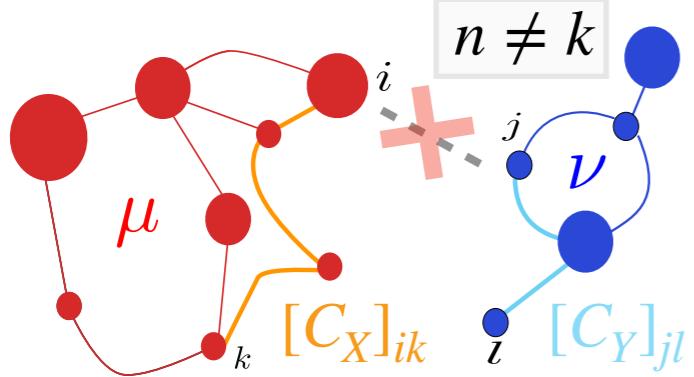
$$c_{\mathcal{X}}(x, x') = c_{\mathcal{Y}}(\phi(x), \phi(x'))$$

Measure preserving

$$\phi \# \mu = \nu$$

From Wasserstein to Gromov-Wasserstein

◆ Gromov-Wasserstein



(Sturm, 2012) (Mémoli, 2011)

◆ L measures distortion

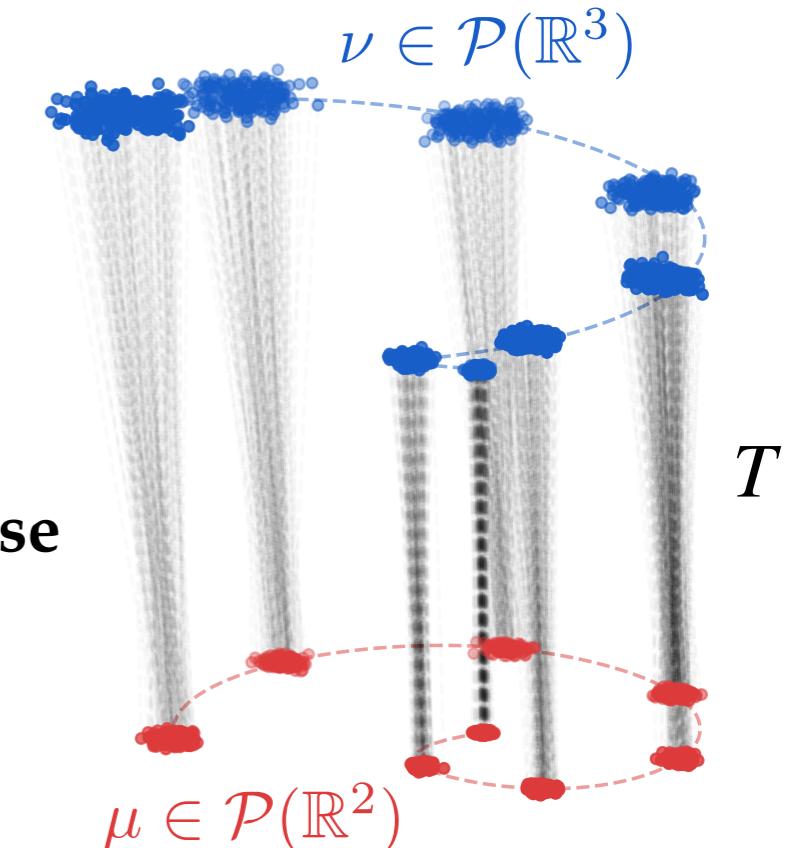
$$\left| [C_X]_{ik} - [C_Y]_{jl} \right|^2$$

◆ Goal : preserving pairwise connectivity

◆ Non-convex quadratic problem (NP-hard)

Quadratic OT: find the plan

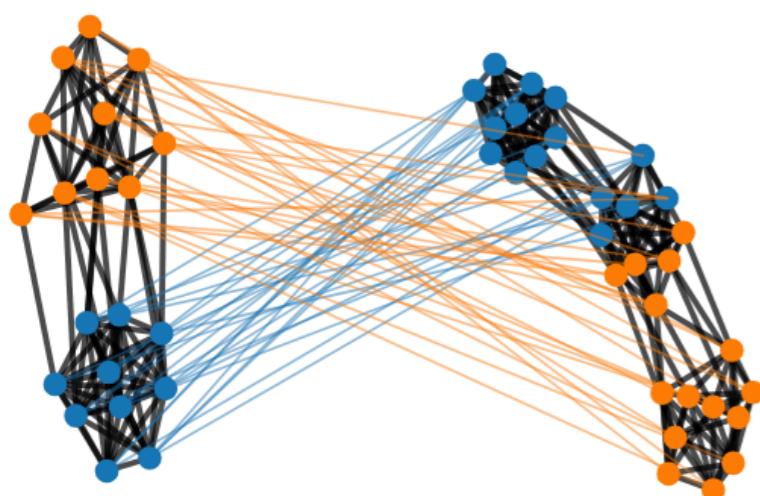
$$\min_{T \in \Pi(\mathbf{a}, \mathbf{b})} \sum_{ijkl} L([C_X]_{ik}, [C_Y]_{jl}) T_{ij} T_{kl}$$



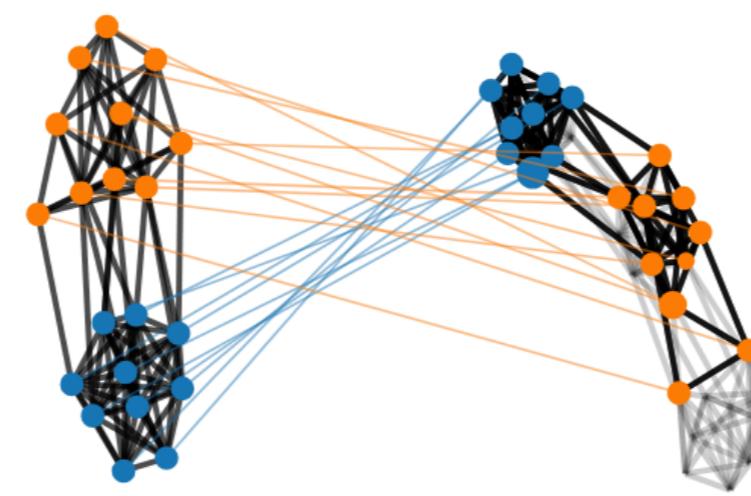
◆ Semi relaxed Gromov-Wasserstein

(Vincent-Cuaz, 2022)

$$GW(\mathbf{C}, \mathbf{h}, \bar{\mathbf{C}}, \bar{\mathbf{h}}) = 0.219$$



$$srGW(\mathbf{C}, \mathbf{h}, \bar{\mathbf{C}}) = 0.05$$

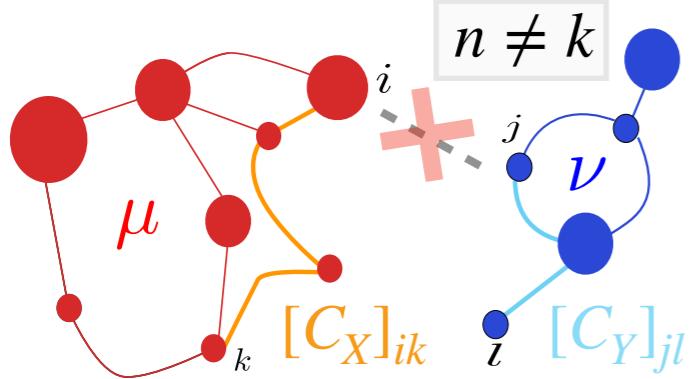


$$\min_T \sum_{ijkl} L([C_X]_{ik}, [C_Y]_{jl}) T_{ij} T_{kl}$$

$$| T^\top \mathbf{1}_n = \mathbf{a}$$

From Wasserstein to Gromov-Wasserstein

◆ Gromov-Wasserstein



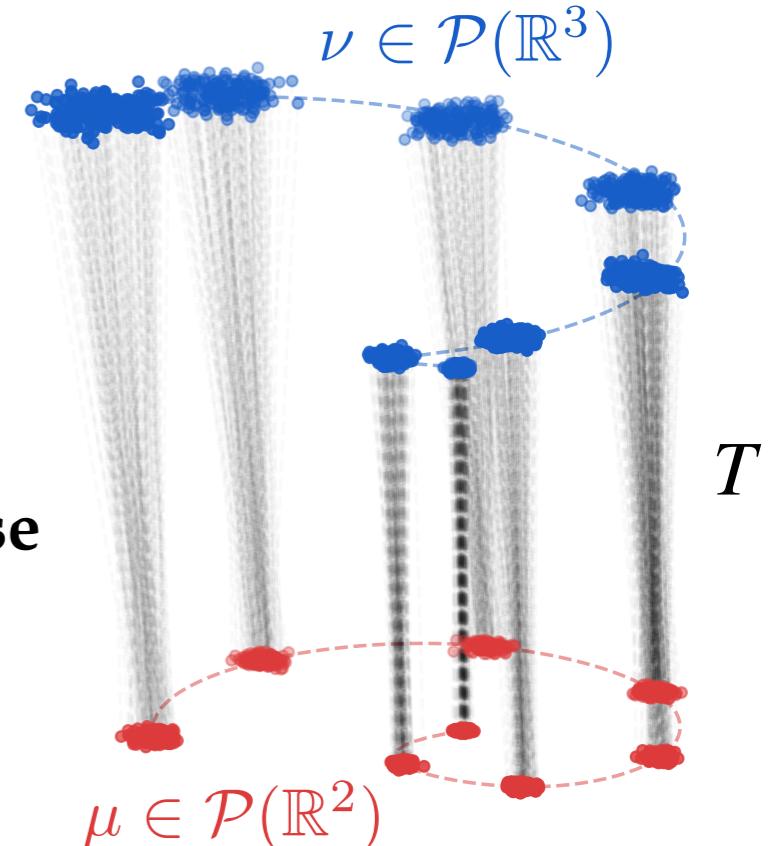
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◆ L measures distortion

$$\left| [C_X]_{ik} - [C_Y]_{jl} \right|^2$$

◆ Goal : preserving pairwise connectivity

◆ Non-convex quadratic problem (NP-hard)

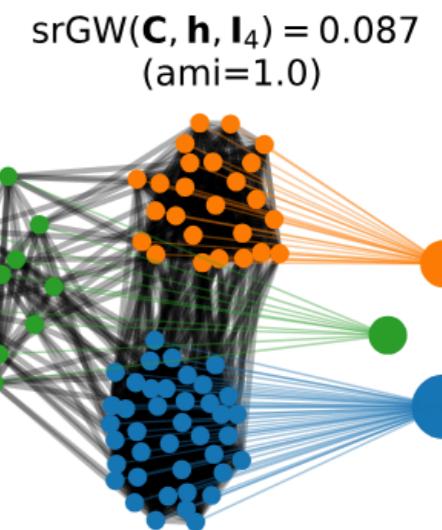
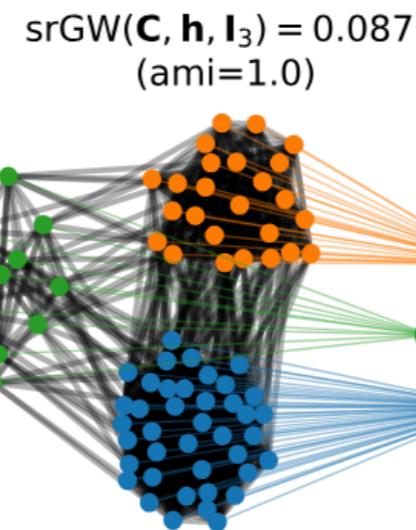
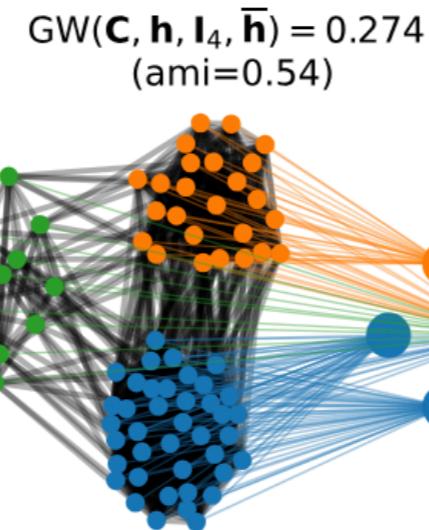
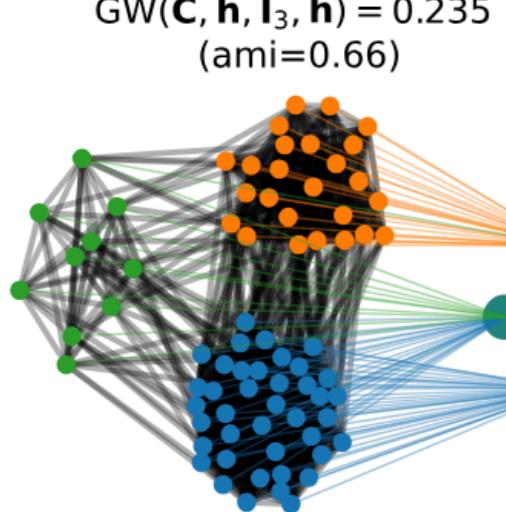


Quadratic OT: find the plan

$$\min_{T \in \Pi(\mathbf{a}, \mathbf{b})} \sum_{ijkl} L \left([C_X]_{ik}, [C_Y]_{jl} \right) T_{ij} T_{kl}$$

◆ Semi relaxed Gromov-Wasserstein

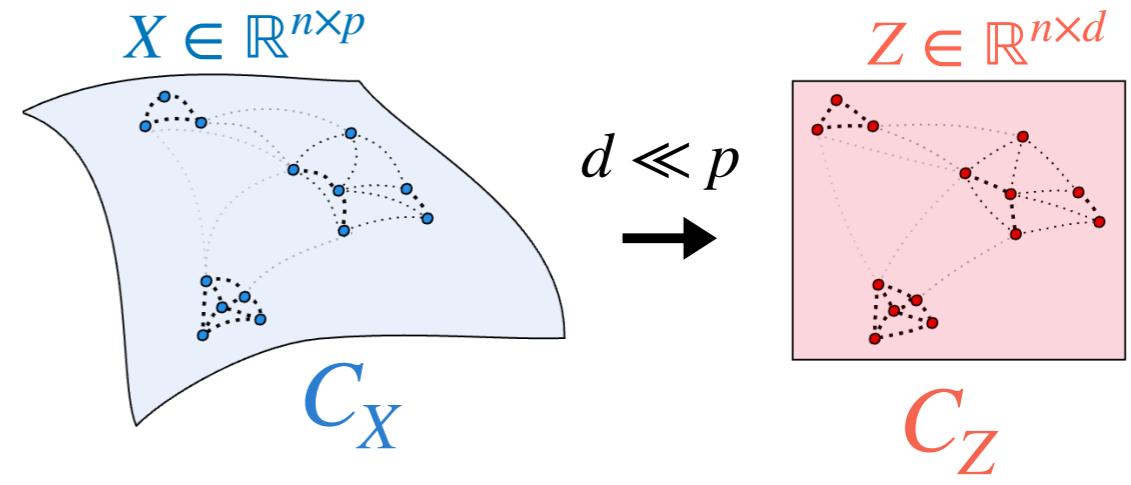
◆ Clustering of nodes



DR as OT in disguise

♦ Dimension reduction

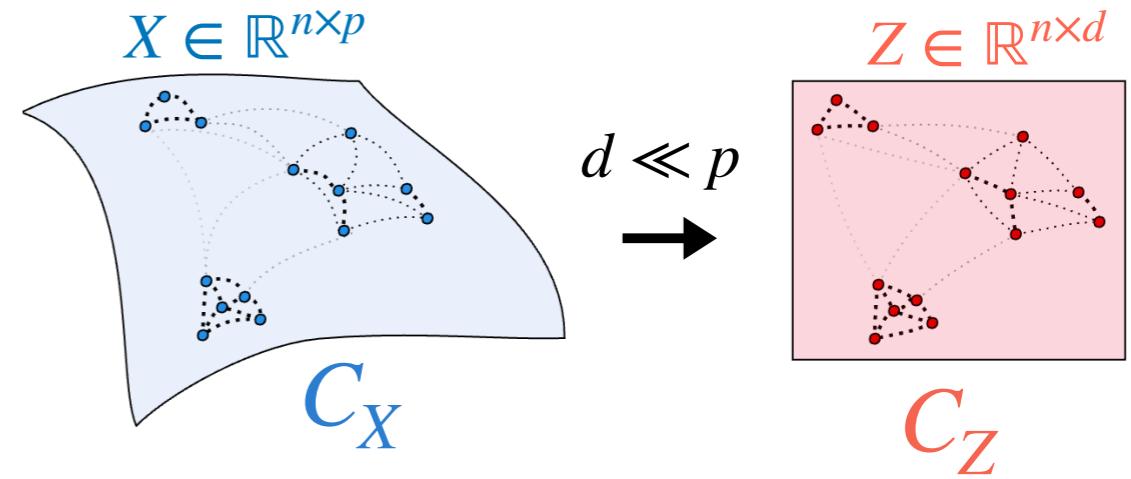
$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$



DR as OT in disguise

♦ Dimension reduction

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$



↑
equiv
↓

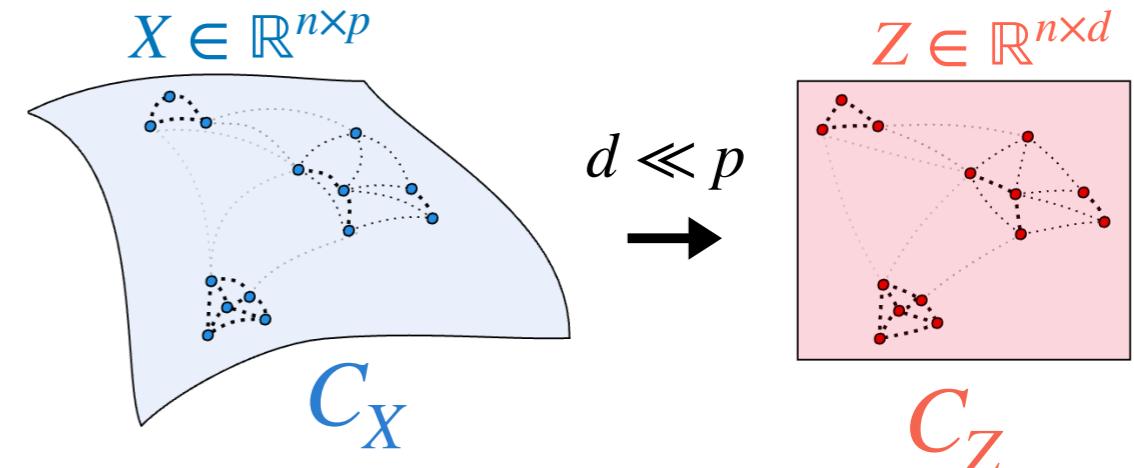
Permutation equivariance
 $\forall P, C_{PZ} = PC_ZP^\top$

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{\sigma \in S_n} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{\sigma(i)\sigma(j)}\right)$$

DR as OT in disguise

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$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$



↑
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↓

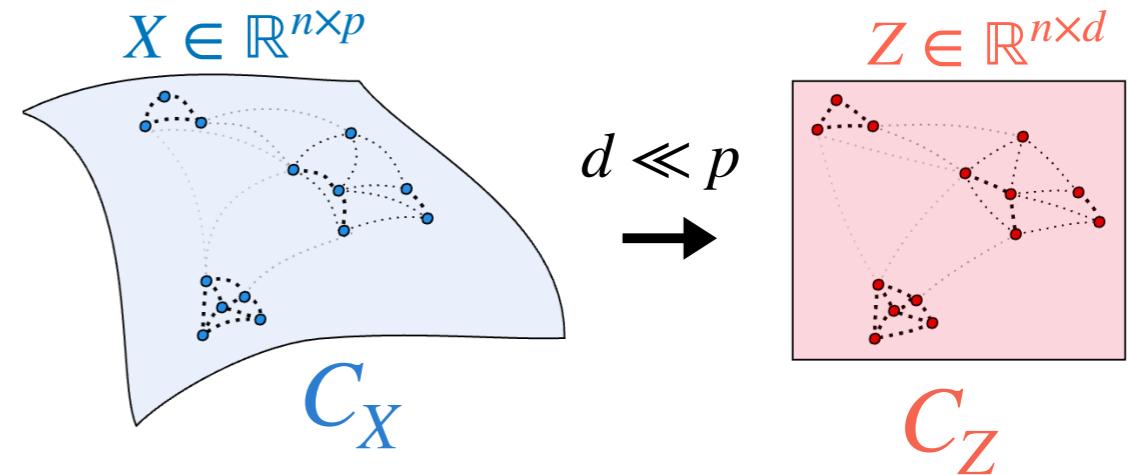
$$\min_{Z \in \mathbb{R}^{n \times d}} \min_P \sum_{i,j,k,l=1}^n L\left([C_X]_{ik}, [C_Z]_{jl}\right) P_{ij}P_{kl}$$

Gromov-Monge

DR as OT in disguise

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$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$



Permutation equivariance

$$\forall P, C_{PZ} = PC_ZP^\top$$

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$$\min_{Z \in \mathbb{R}^{n \times d}} \min_P \sum_{i,j,k,l=1}^n L\left([C_X]_{ik}, [C_Z]_{jl}\right) P_{ij}P_{kl}$$

Gromov-Monge



♦ Gromov-Wasserstein projection

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T \in \Pi(\frac{1}{n}, \frac{1}{n})} \sum_{ijkl} L\left([C_X]_{ik}, [C_Z]_{jl}\right) T_{ij}T_{kl}$$

Gromov-Wasserstein

DR as OT in disguise

♦ Dimension reduction

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$

equiv
↔

Permutation equivariance

$$\forall P, C_{PZ} = PC_Z P^\top$$

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{\sigma \in S_n} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{\sigma(i)\sigma(j)}\right)$$

equiv
↔

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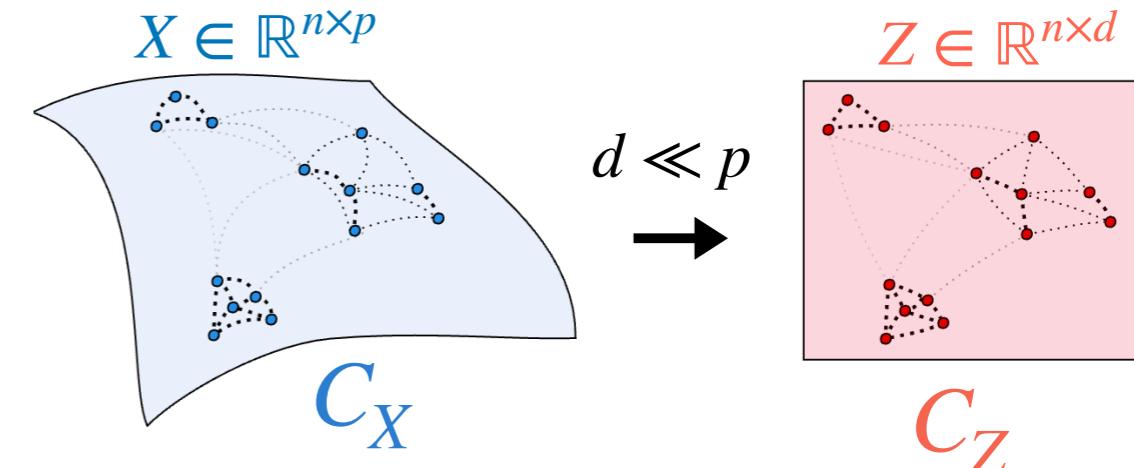
Gromov-Monge

↑?
↓

♦ Gromov-Wasserstein projection

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T \in \Pi(\frac{\mathbf{1}_n}{n}, \frac{\mathbf{1}_n}{n})} \sum_{ijkl} L\left([C_X]_{ik}, [C_Z]_{jl}\right) T_{ij} T_{kl}$$

Gromov-Wasserstein



♦ Equivalence holds for

Spectral methods

♦ C_X any matrix, $L = |\cdot|^2$, $C_Z = ZZ^\top$

DR as OT in disguise

♦ Dimension reduction

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$

equiv
↔

Permutation equivariance

$$\forall P, C_{PZ} = PC_Z P^\top$$

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{\sigma \in S_n} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{\sigma(i)\sigma(j)}\right)$$

equiv
↔

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_P \sum_{i,j,k,l=1}^n L\left([C_X]_{ik}, [C_Z]_{jl}\right) P_{ij} P_{kl}$$

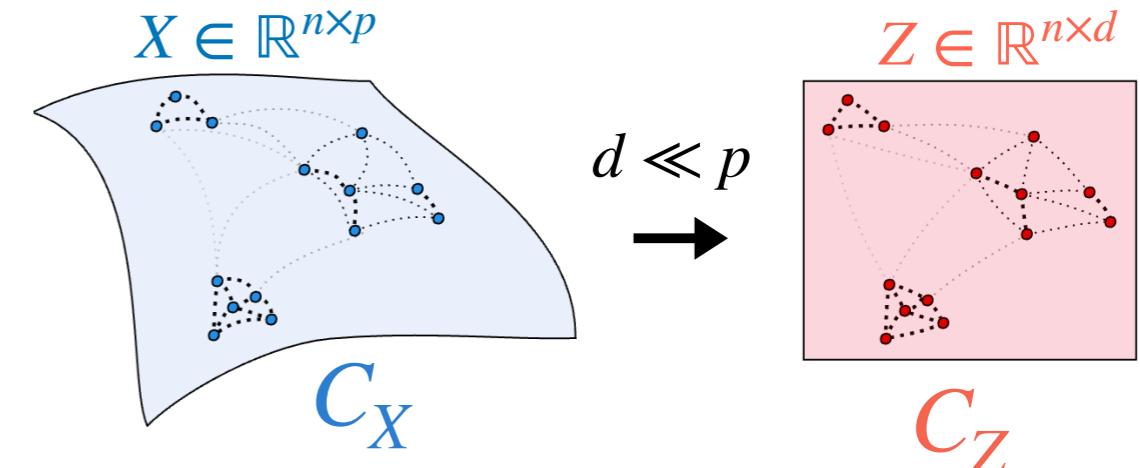
Gromov-Monge

↑?
?

♦ Gromov-Wasserstein projection

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T \in \Pi(\frac{\mathbf{1}_n}{n}, \frac{\mathbf{1}_n}{n})} \sum_{ijkl} L\left([C_X]_{ik}, [C_Z]_{jl}\right) T_{ij} T_{kl}$$

Gromov-Wasserstein



♦ Equivalence holds for

Spectral methods

♦ C_X any matrix, $L = |\cdot|^2$, $C_Z = ZZ^\top$

A is CPD : $\forall x$ **s.t.** $x^\top 1 = 0$, $x^\top Ax \geq 0$

DR as OT in disguise

♦ Dimension reduction

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$

equiv
↔

Permutation equivariance

$$\forall P, C_{PZ} = PC_ZP^\top$$

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{\sigma \in S_n} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{\sigma(i)\sigma(j)}\right)$$

equiv
↔

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_P \sum_{i,j,k,l=1}^n L\left([C_X]_{ik}, [C_Z]_{jl}\right) P_{ij}P_{kl}$$

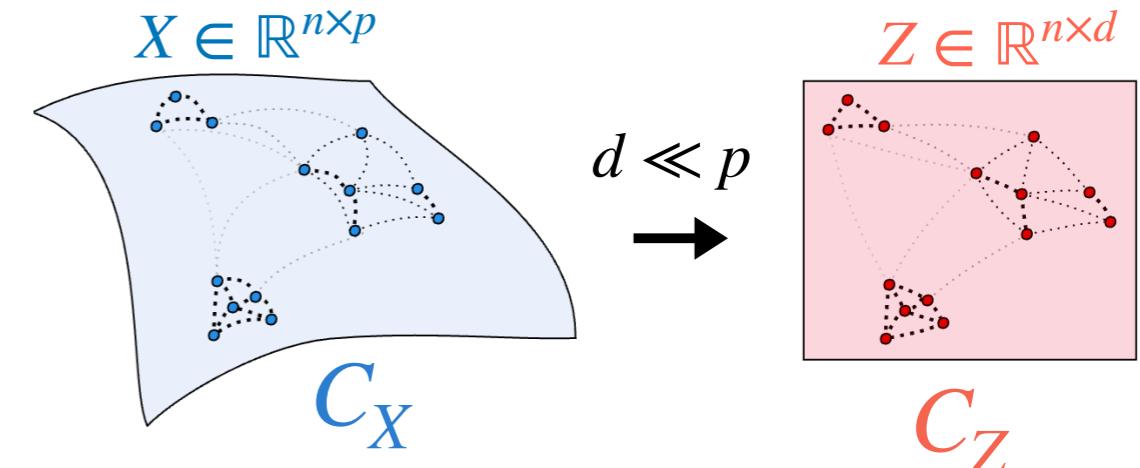
Gromov-Monge

↑?
?

♦ Gromov-Wasserstein projection

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T \in \Pi(\frac{\mathbf{1}_n}{n}, \frac{\mathbf{1}_n}{n})} \sum_{ijkl} L\left([C_X]_{ik}, [C_Z]_{jl}\right) T_{ij}T_{kl}$$

Gromov-Wasserstein



♦ Equivalence holds for

Spectral methods

♦ C_X any matrix, $L = |\cdot|^2$, $C_Z = ZZ^\top$

A is CPD : $\forall x$ **s.t.** $x^\top 1 = 0$, $x^\top Ax \geq 0$

Neighbor embedding methods

♦ C_X is CPD, $L = KL$

$$C_Z = \text{diag}(\alpha_Z) K_Z \text{ diag}(\beta_Z)$$

where $\log(K_Z)$ is CPD

DR as OT in disguise

♦ Dimension reduction

$$\min_{Z \in \mathbb{R}^{n \times d}} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{ij}\right)$$

equiv
↔

Permutation equivariance

$$\forall P, C_{PZ} = PC_ZP^\top$$

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{\sigma \in S_n} \sum_{i,j=1}^n L\left([C_X]_{ij}, [C_Z]_{\sigma(i)\sigma(j)}\right)$$

equiv
↔

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_P \sum_{i,j,k,l=1}^n L\left([C_X]_{ik}, [C_Z]_{jl}\right) P_{ij}P_{kl}$$

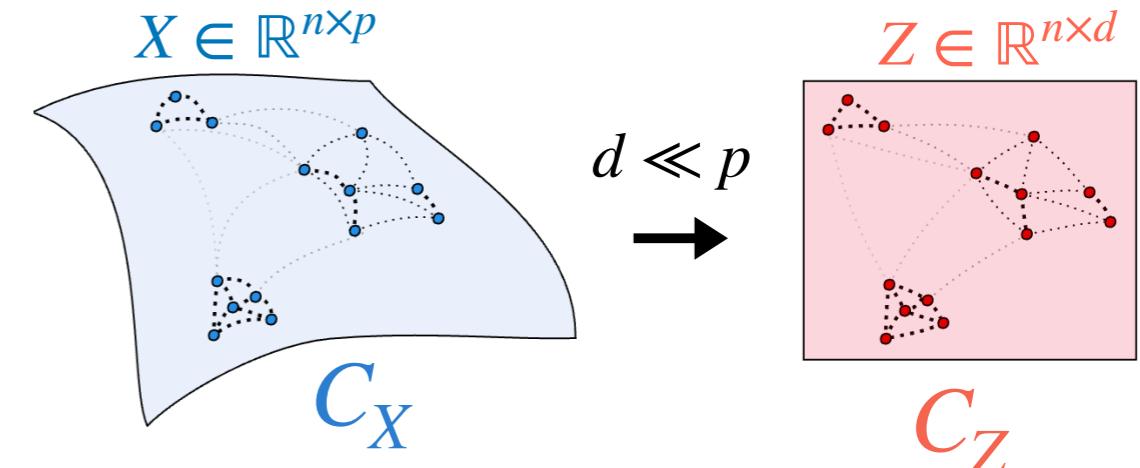
Gromov-Monge

↑?
?

♦ Gromov-Wasserstein projection

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T \in \Pi(\frac{1_n}{n}, \frac{1_n}{n})} \sum_{ijkl} L\left([C_X]_{ik}, [C_Z]_{jl}\right) T_{ij}T_{kl}$$

Gromov-Wasserstein



♦ Equivalence holds for

Spectral methods

♦ C_X any matrix, $L = |\cdot|^2$, $C_Z = ZZ^\top$

A is CPD : $\forall x$ **s.t.** $x^\top 1 = 0$, $x^\top Ax \geq 0$

Neighbor embedding methods

♦ C_X is CPD, $L = KL$

$$C_Z = \text{diag}(\alpha_Z) K_Z \text{ diag}(\beta_Z)$$

where $\log(K_Z)$ is CPD

| e.g. $K_Z = \exp(-\|\mathbf{z}_i - \mathbf{z}_j\|_2^2)$
and its usual normalizations

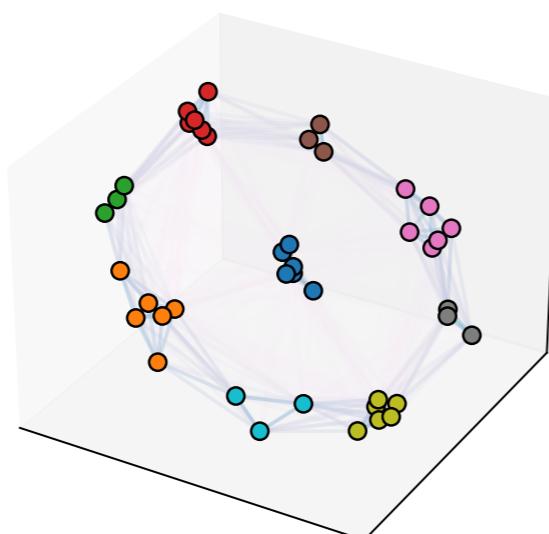
$$1_n^\top K_Z 1_n = 1, K_Z 1_n = 1_n, + K_Z^\top 1_n = 1_n \\ K_Z 1_n = 1_n$$

(Sinkhorn & Knopp, 1967)

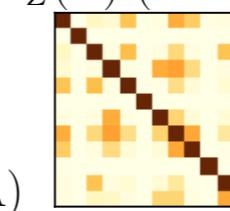
| Beware that C_X is not always CPD.

Distributional reduction

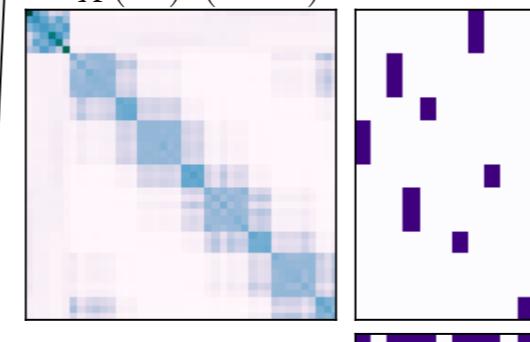
Input \mathbf{X}



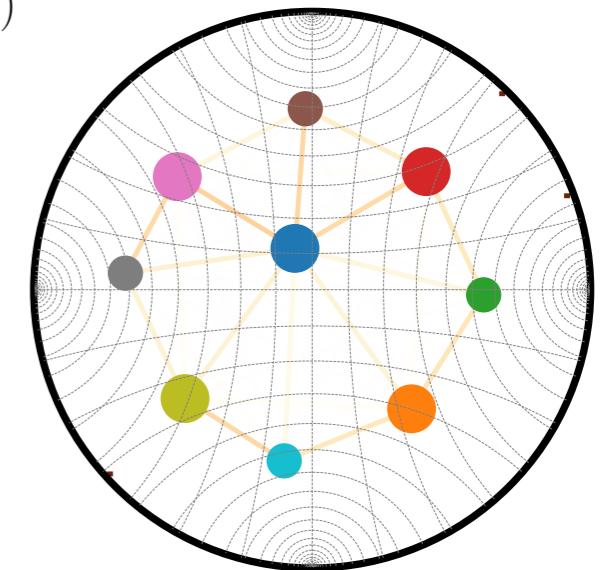
$C_Z(\mathbf{Z})$ (Lorentz)



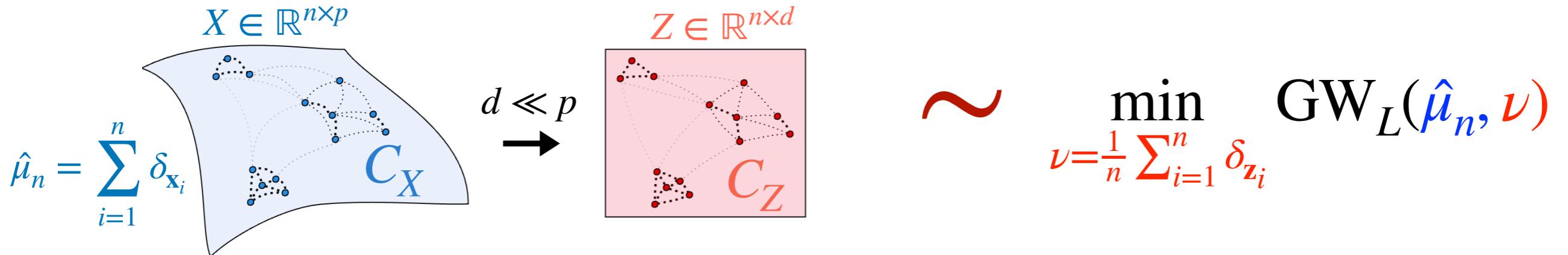
$C_X(\mathbf{X})$ (SEA)



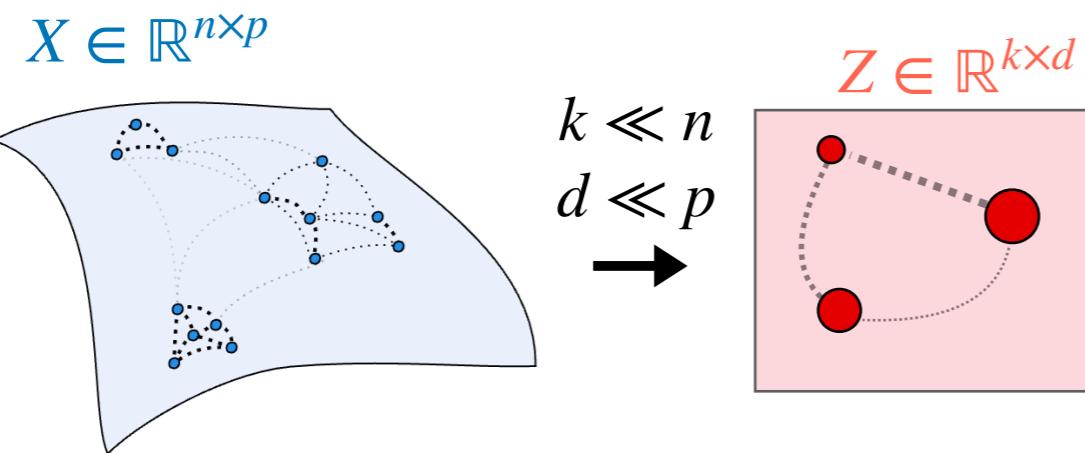
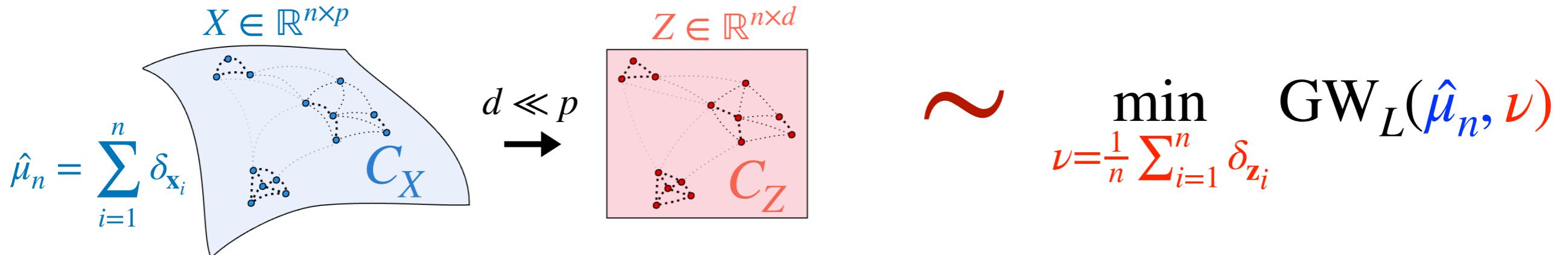
Embedding \mathbf{Z}



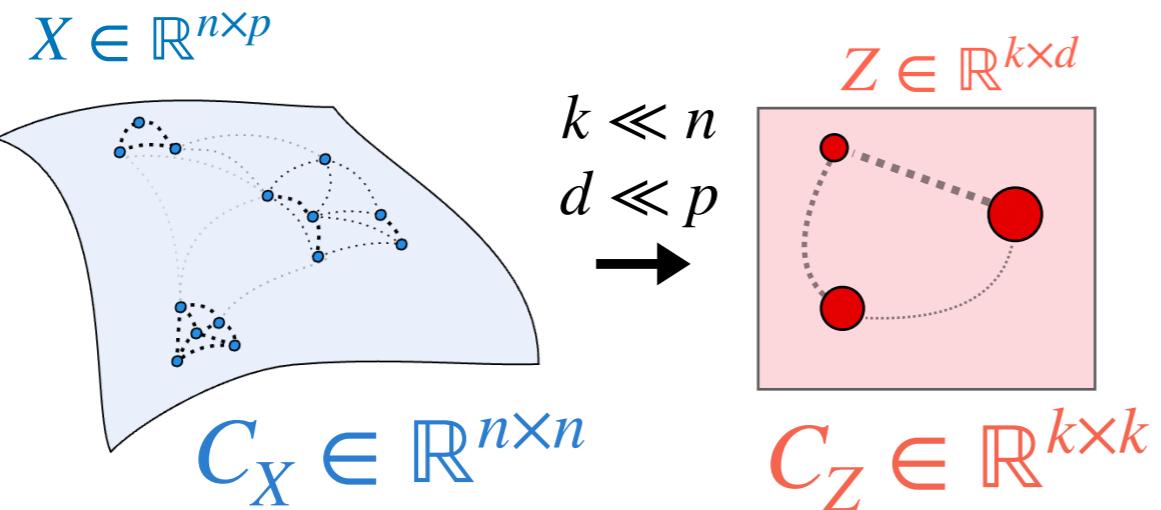
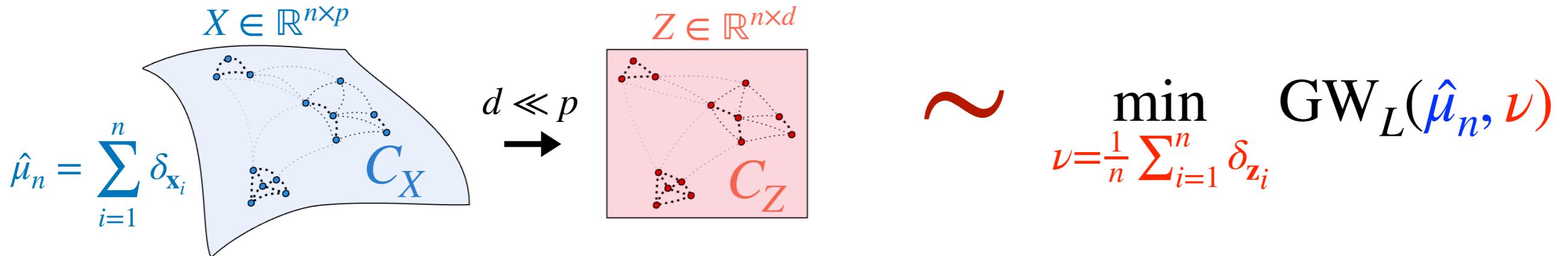
Distributional Reduction



Distributional Reduction



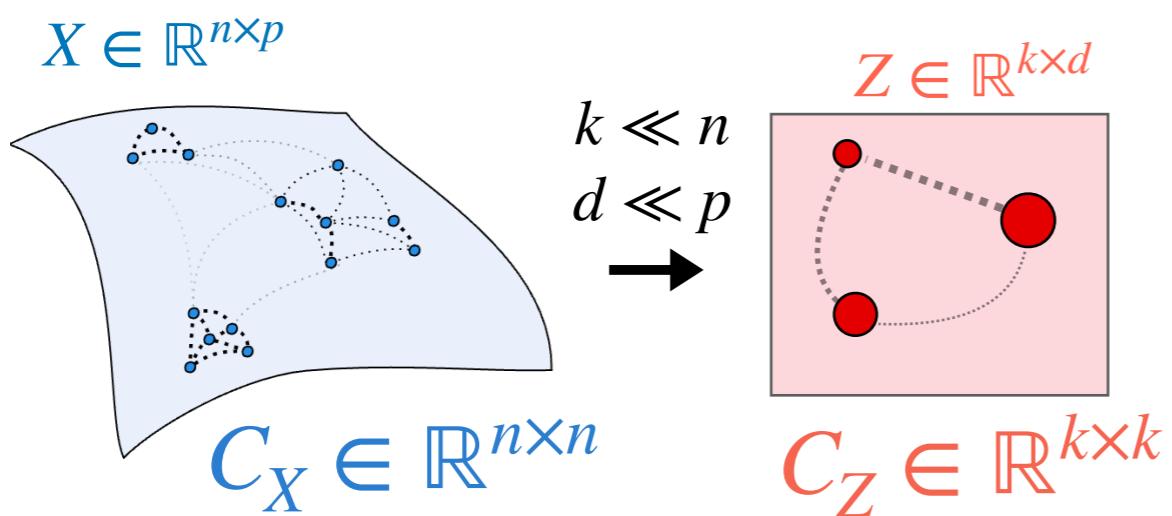
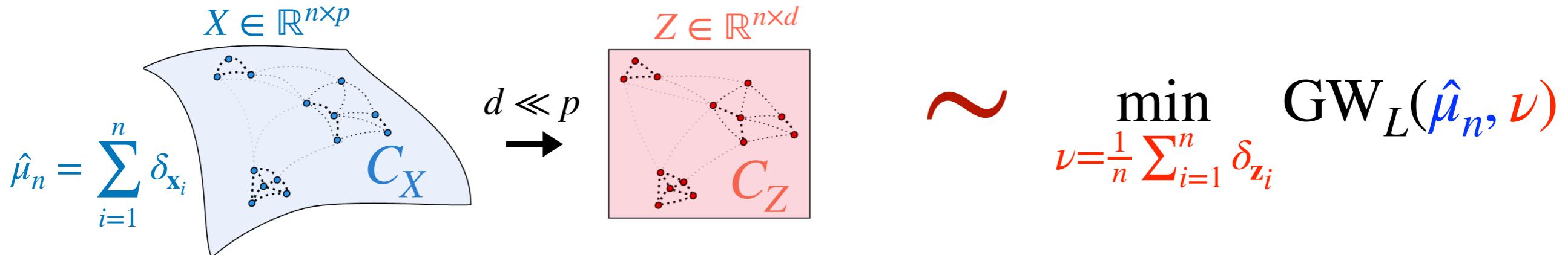
Distributional Reduction



◆ **GW projection**

$$\min_{\nu \in \mathcal{P}_k(\mathbb{R}^d)} \text{GW}(\hat{\mu}_n, \nu) \quad \nu = \sum_{j=1}^k b_j \delta_{\mathbf{z}_j}$$

Distributional Reduction



◆ GW projection

$$\min_{\nu \in \mathcal{P}_k(\mathbb{R}^d)} \text{GW}(\hat{\mu}_n, \nu)$$

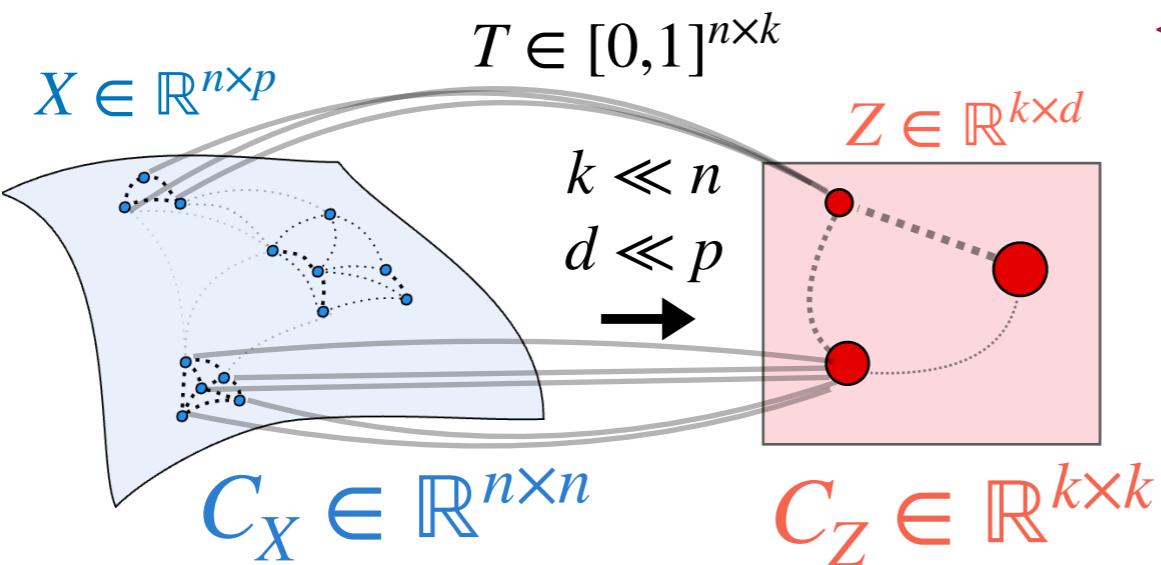
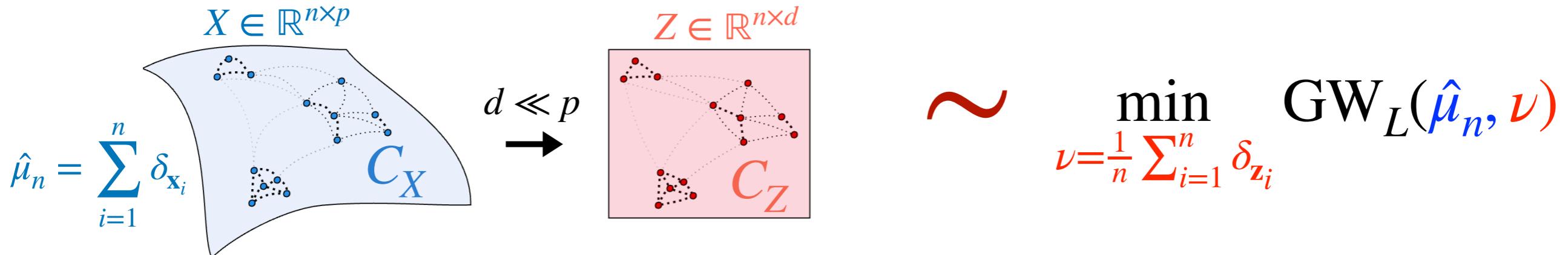
$$\nu = \sum_{j=1}^k b_j \delta_{\mathbf{z}_j}$$

◆ Optimization problem

$$\min_{Z \in \mathbb{R}^{k \times d}} \min_{\mathbf{b} \in \Sigma_k} \text{GW}_L(C_X, C_Z, \frac{1}{n}, \mathbf{b})$$

- ◆ Find few prototypes in low dim.
- ◆ Find the weights/cluster size

Distributional Reduction



◆ **GW projection**

$$\min_{\nu \in \mathcal{P}_k(\mathbb{R}^d)} \text{GW}(\hat{\mu}_n, \nu)$$

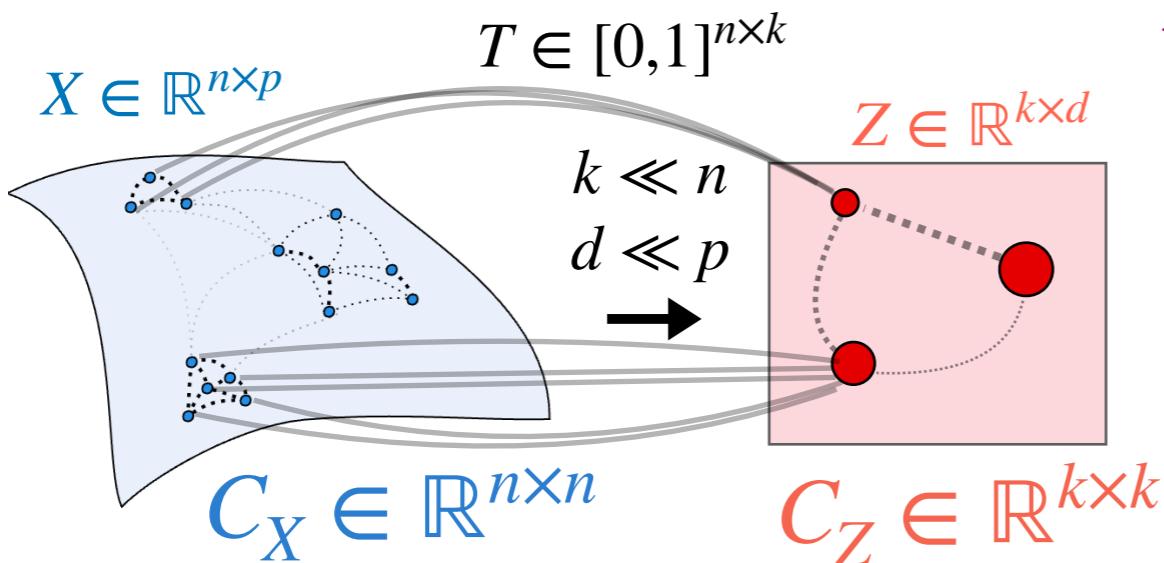
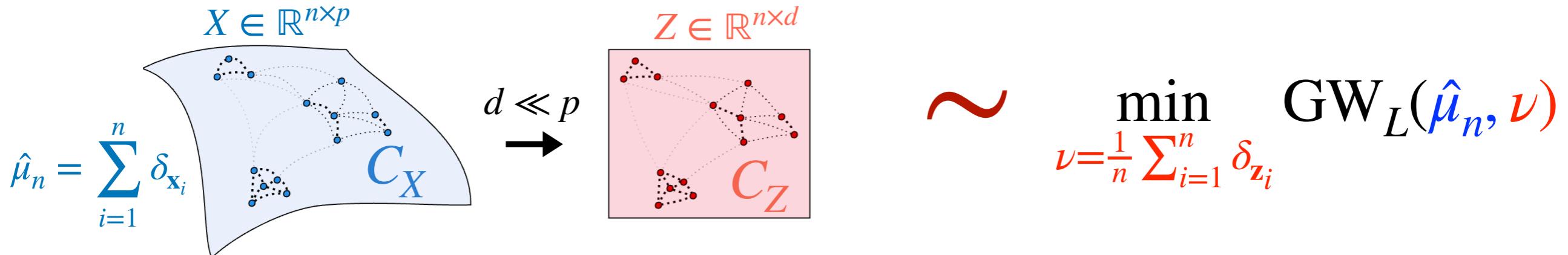
$$\nu = \sum_{j=1}^k b_j \delta_{\mathbf{z}_j}$$

◆ **Optimization problem**

$$\min_{Z \in \mathbb{R}^{k \times d}} \min_{b \in \Sigma_k} \text{GW}_L(C_X, C_Z, \frac{1}{n}, b)$$

- ◆ Find few prototypes in low dim.
- ◆ Find the weights/cluster size
- ◆ Clustering via the coupling T (soft-assignment)
- ◆ Sufficient conditions for hard assignment (see paper)

Distributional Reduction



◆ GW projection

$$\min_{\nu \in \mathcal{P}_k(\mathbb{R}^d)} \text{GW}(\hat{\mu}_n, \nu)$$

$$\nu = \sum_{j=1}^k b_j \delta_{\mathbf{z}_j}$$

◆ Optimization problem

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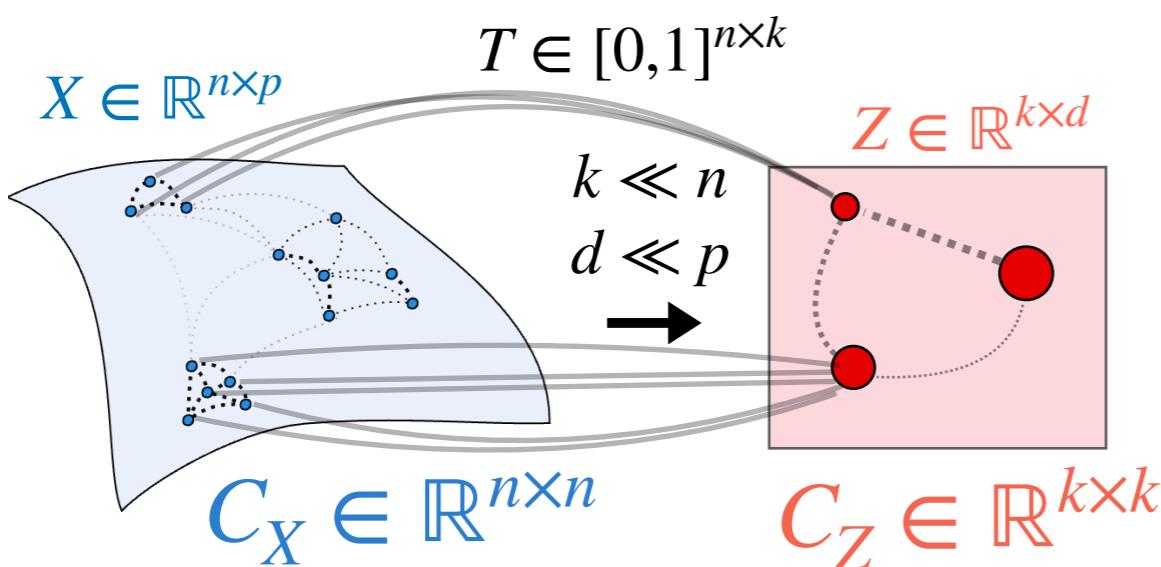
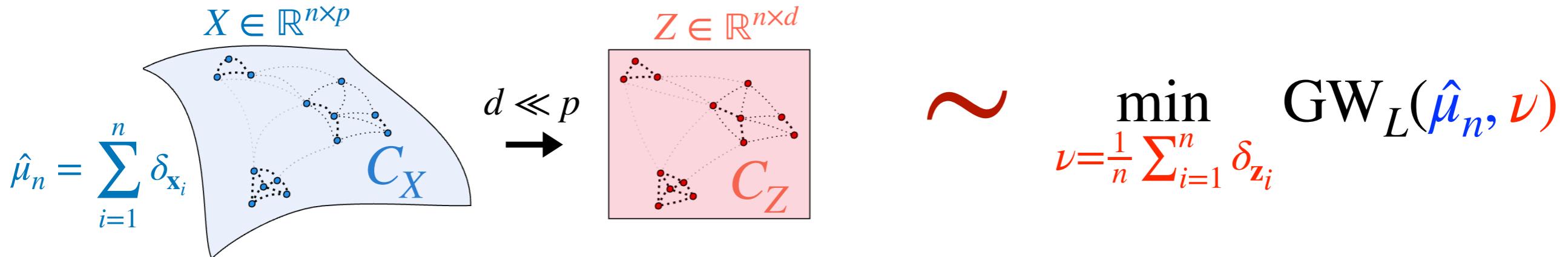
◆ A semi-relaxed objective

(Vincent-Cuaz et al., 2022)

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T: T1_k = \frac{1}{n}} \sum_{ijkl} L([C_X]_{ik}, [C_Z]_{jl}) T_{ij} T_{kl}$$

easier than GW

Distributional Reduction



◆ GW projection

$$\min_{\nu \in \mathcal{P}_k(\mathbb{R}^d)} \text{GW}(\hat{\mu}_n, \nu)$$

$$\nu = \sum_{j=1}^k b_j \delta_{\mathbf{z}_j}$$

◆ Optimization problem

$$\min_{Z \in \mathbb{R}^{k \times d}} \min_{b \in \Sigma_k} \text{GW}_L(C_X, C_Z, \frac{1}{n}, b)$$

- ◆ Find few prototypes in low dim.
- ◆ Find the weights/cluster size
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◆ A semi-relaxed objective

(Vincent-Cuaz et al., 2022)

$$\min_{Z \in \mathbb{R}^{n \times d}} \min_{T: T1_k = \frac{1}{n}} \sum_{ijkl} L \left([C_X]_{ik}, [C_Z]_{jl} \right) T_{ij} T_{kl}$$

easier than GW

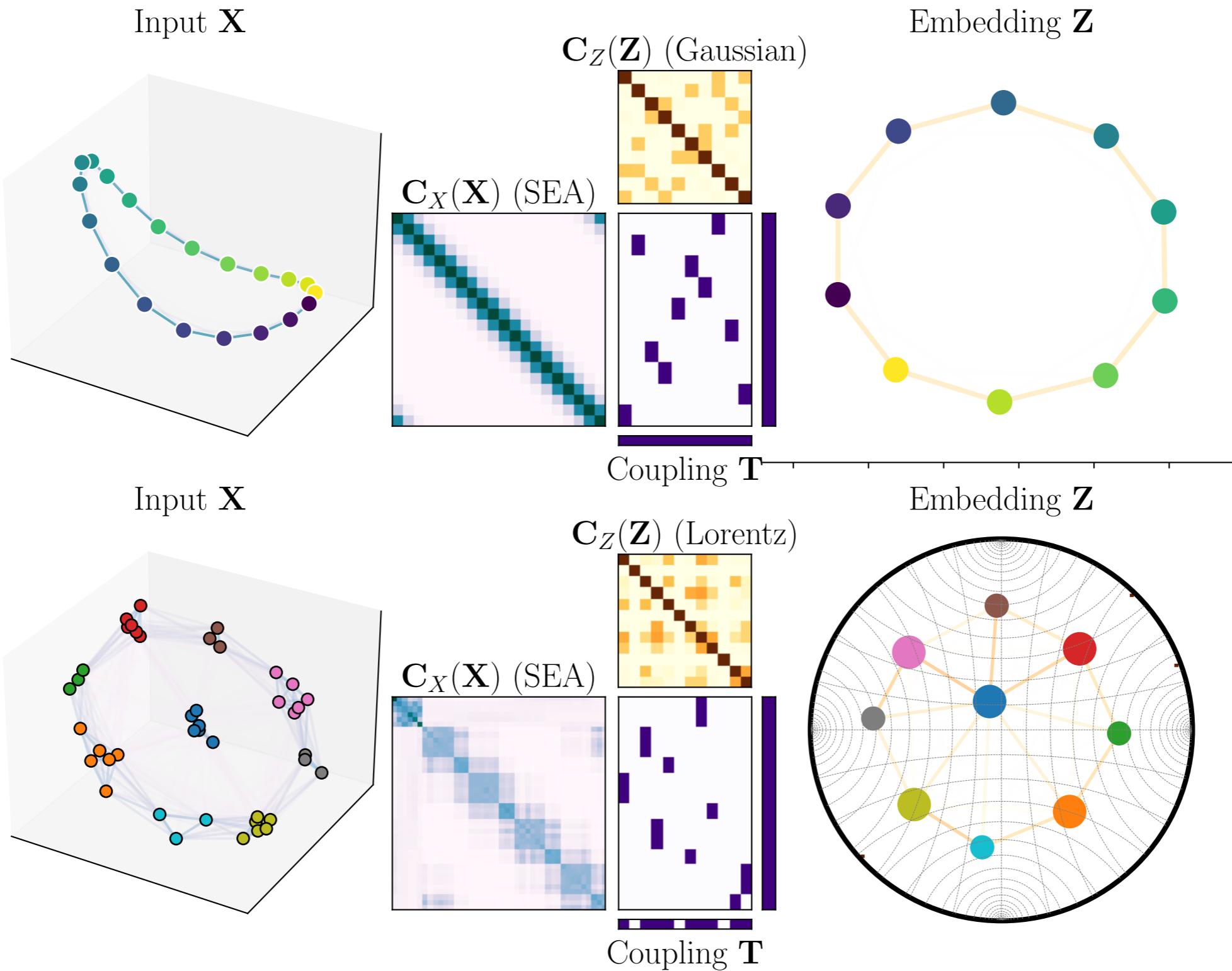
- ◆ Non-convex problem

- ◆ BCD: alternates optim in Z, in T

- ◆ Optim in T: CG solver in $O(n^2 k)$ for $L \in \{\text{KL}, |\cdot|^2\}$

- ◆ With low-rank structures $O(nkr + n^2)$

Distributional Reduction



Motivation

◆ Single-cell RNA-seq

Technical noise due to partial sampling of RNA molecules within cells.

METHOD

Open Access

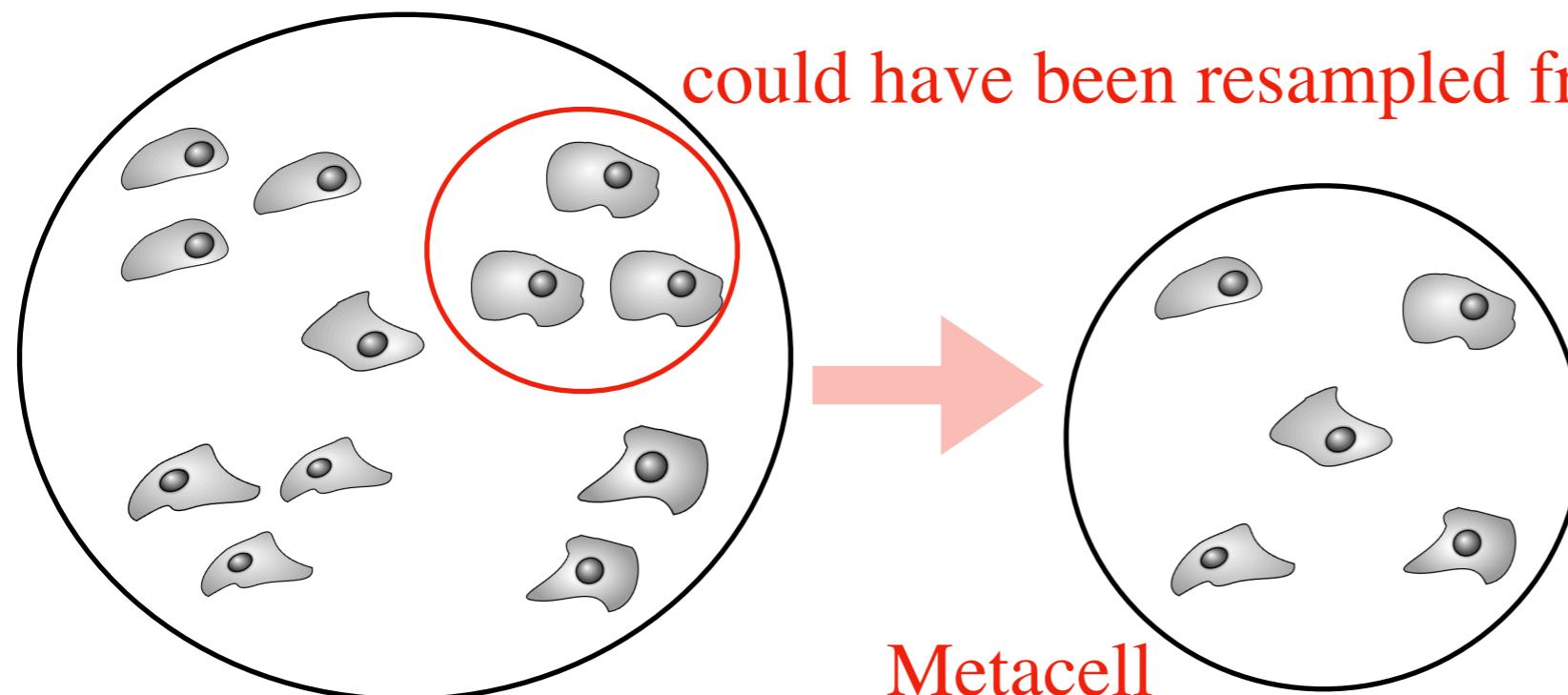


MetaCell: analysis of single-cell RNA-seq data using *K*-nn graph partitions

Yael Baran¹, Akhiad Bercovich¹, Arnau Sebe-Pedros¹, Yaniv Lubling¹, Amir Giladi², Elad Chomsky¹, Zohar Meir¹, Michael Hoichman¹, Aviezer Lifshitz¹ and Amos Tanay^{1*}

Problem : impossible to resample a cell

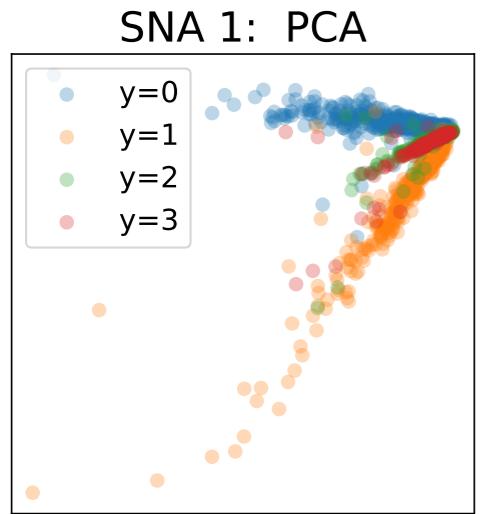
- integration of data from different cells
- need to **separate the sampling effect from biological variance**



- ◆ We would like to choose the granularity of the output data

Distributional Reduction

- ◆ **Single-cell dataset** (Chen et al., 2019) Only solve $\min_{\mathbf{b} \in \Sigma_k} \text{GW}_L(\mathbf{C}_X, \mathbf{C}_Z, \frac{\mathbf{1}_n}{n}, \mathbf{b})$



Distributional Reduction

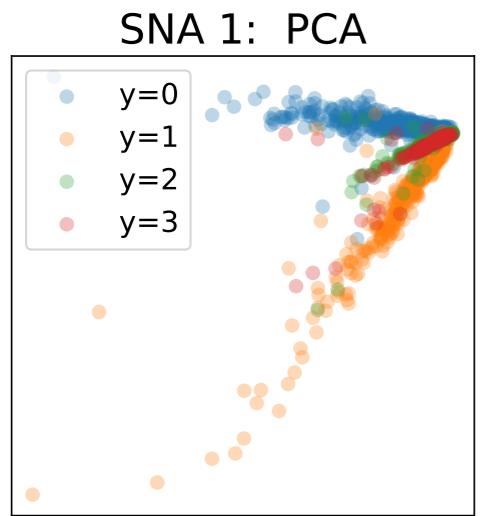
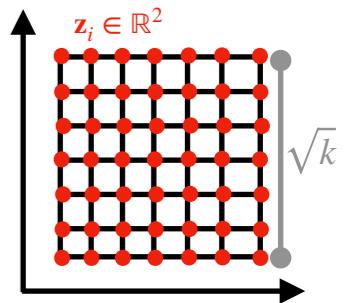
- ◆ **Single-cell dataset** (Chen et al., 2019) Only solve $\min_{\mathbf{b} \in \Sigma_k} \text{GW}_L(\mathbf{C}_X, \mathbf{C}_Z, \frac{\mathbf{1}_n}{n}, \mathbf{b})$ with $\mathbf{C}_X = \mathbf{X}\mathbf{X}^\top$

and

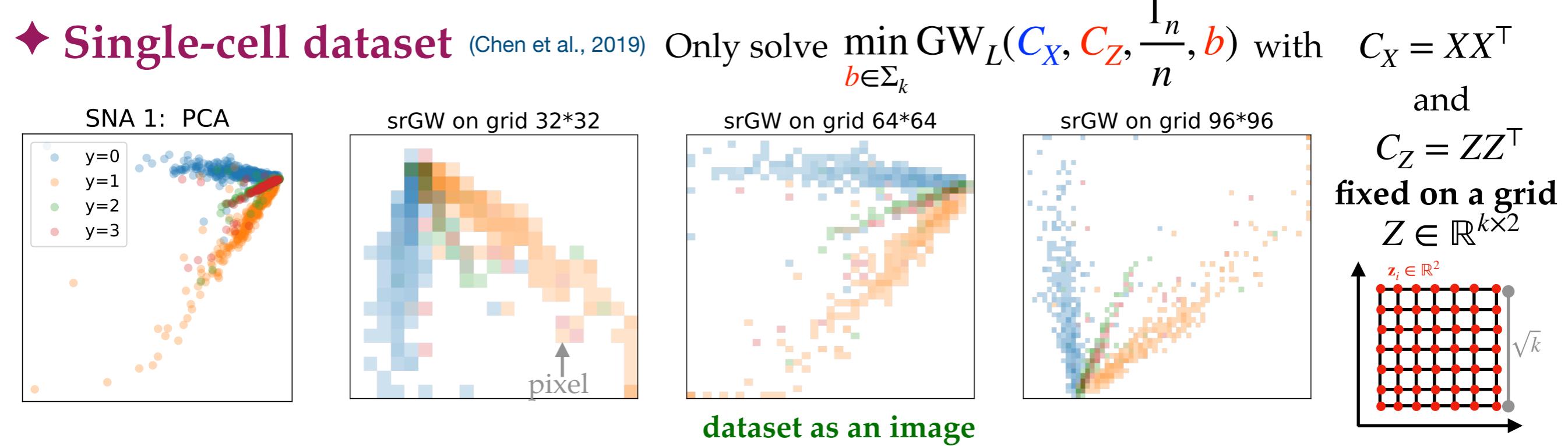
$$\mathbf{C}_Z = \mathbf{Z}\mathbf{Z}^\top$$

fixed on a grid

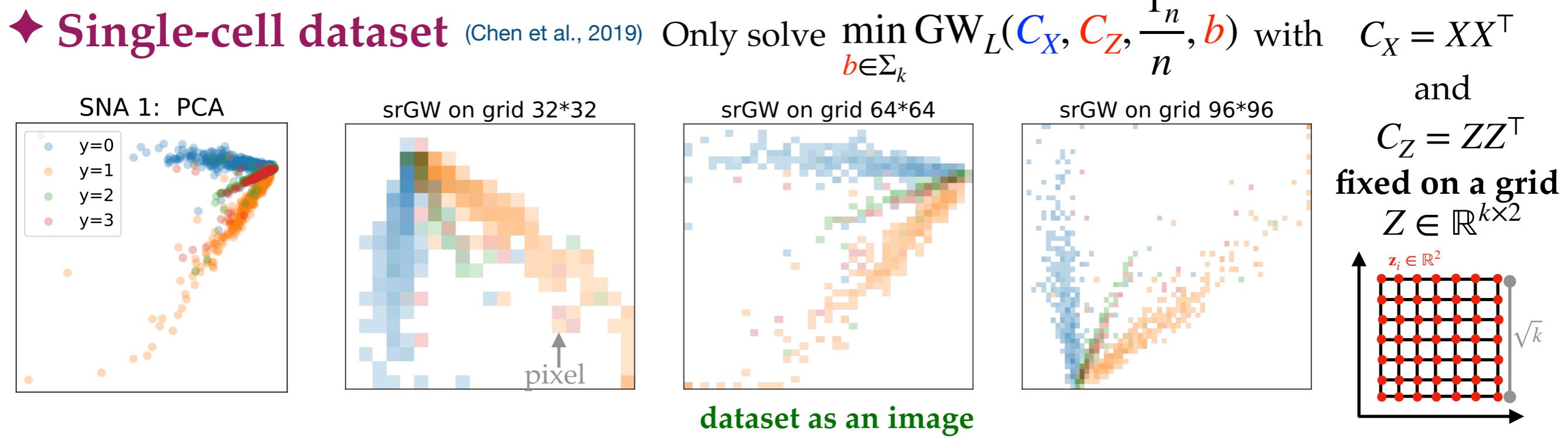
$$\mathbf{Z} \in \mathbb{R}^{k \times 2}$$



Distributional Reduction



Distributional Reduction

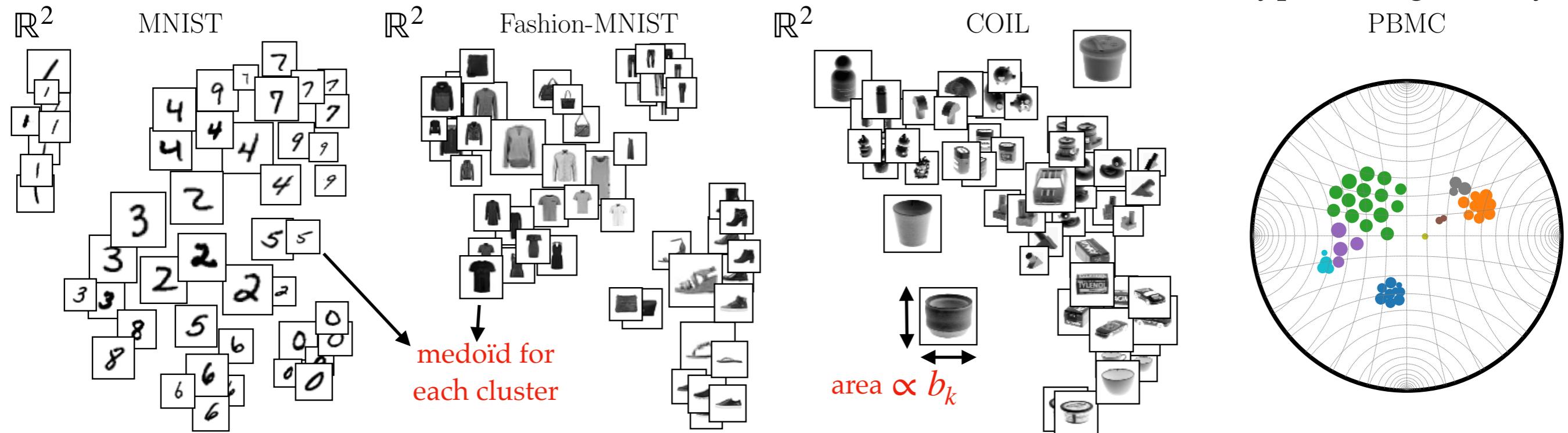


◆ Image datasets

\mathbf{C}_X symmetric entropic aff. (Van Assel et al., 2023) \mathbf{C}_Z Student t-kernel

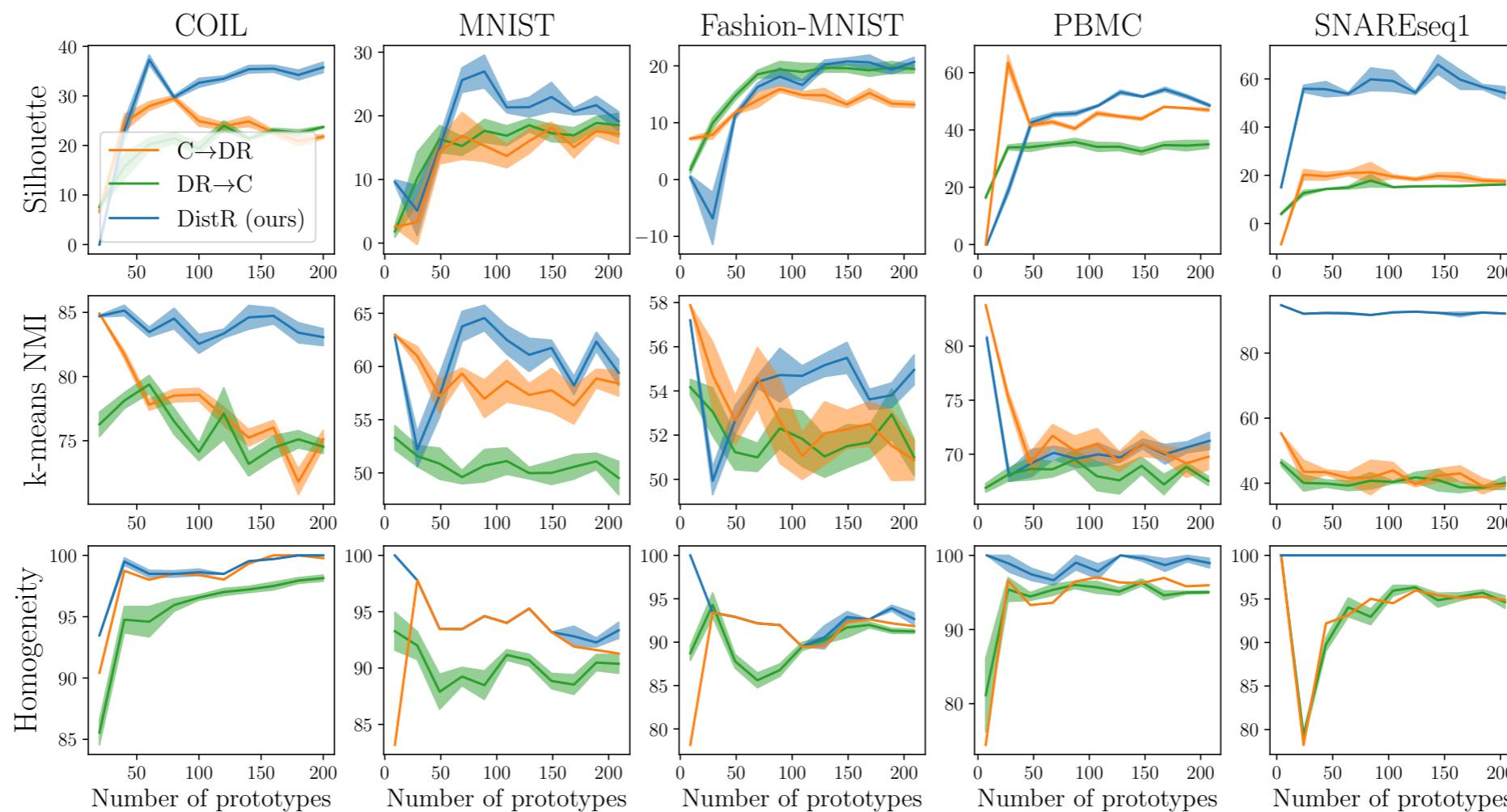
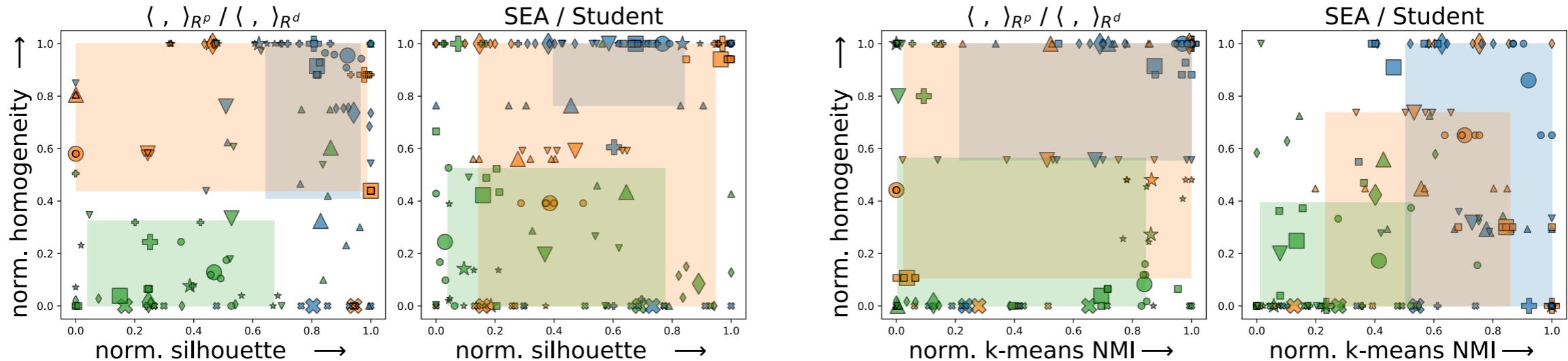
Hyperbolic geometry

PBMC



Distributional Reduction

♦ Comparison with DR then clustering or clustering then DR



Thank you!

Open-source implementations are available here:



- | DR library in pytorch
 - Implements popular DR algos
 - Modular
 - Efficiency (GPU, KeOps)

| Python Optimal Transport

- OT LP solver, Sinkhorn
- Barycenters
- Gromov, graphs OT...

| **Url:** <https://github.com/PythonOT/POT>

