FUZZIFICATION: Anti-Fuzzing Techniques

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Fuzzing Discovers Many Vulnerabilities

50 CVEs in 50 Days: Fuzzing Adobe Reader

December 12, 2018

Research By: Yoav Alon, Netanel Ben-Simon

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Reader

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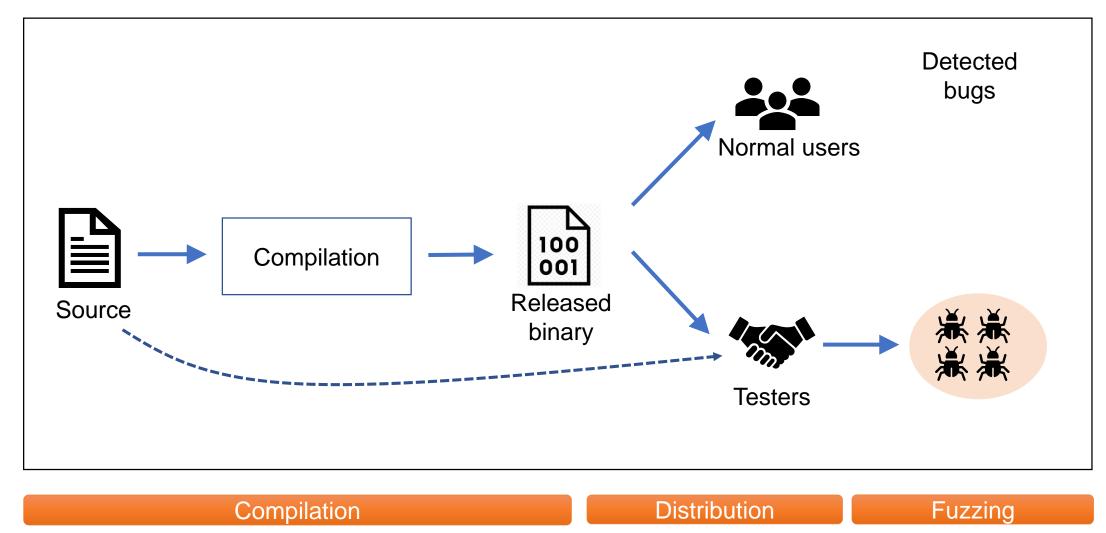
Research By: Yoav Al

Google's automated fuzz bot has found over 9,000 bugs in the past two years

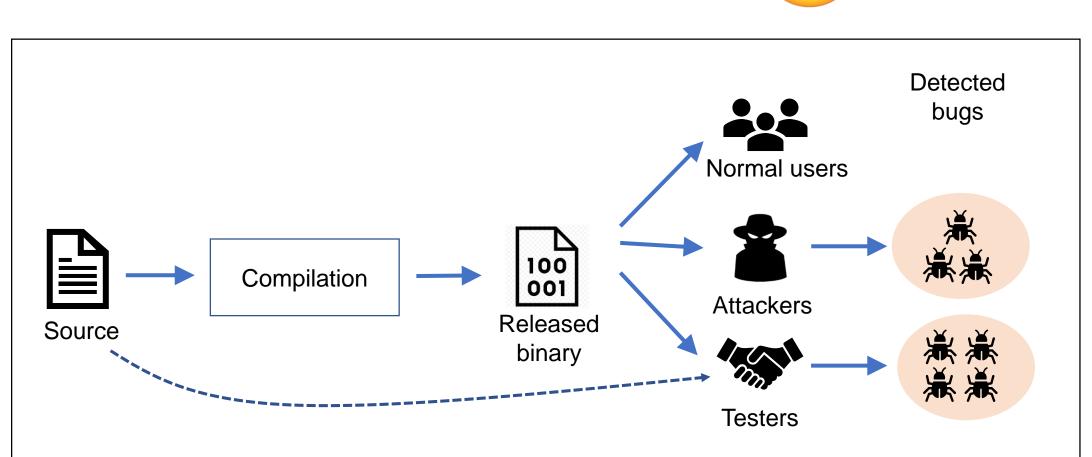
Google improves OSS-Fuzz service, plans to invite new op source projects to join.

Testers Find Bugs with Fuzzing





But Attackers Also Find Bugs

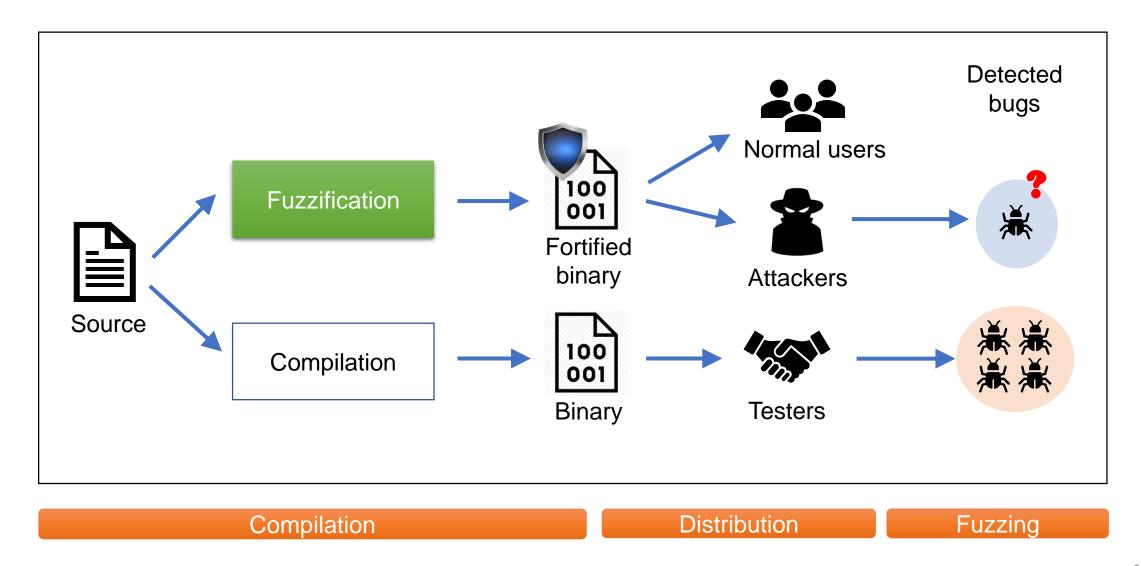


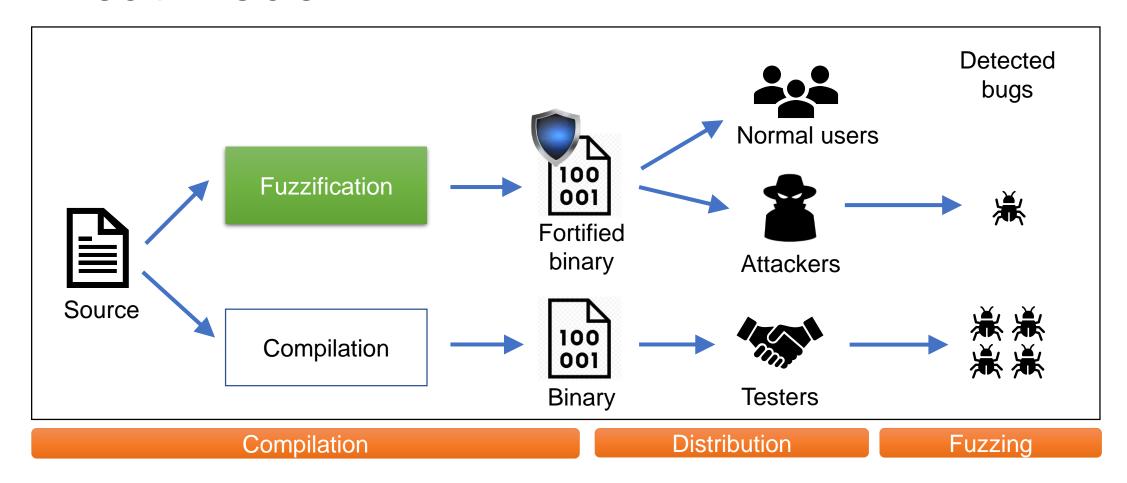
Compilation

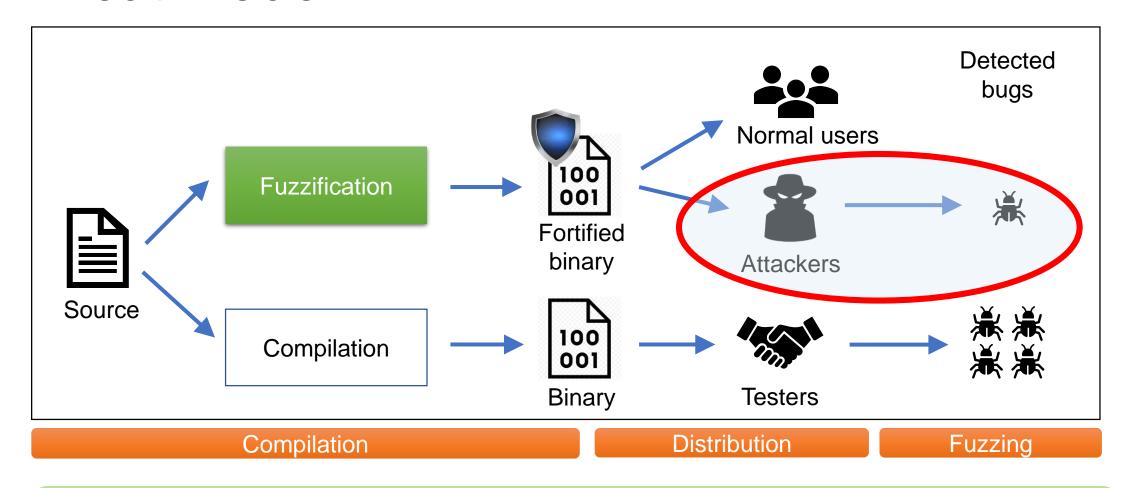
Distribution

Fuzzing

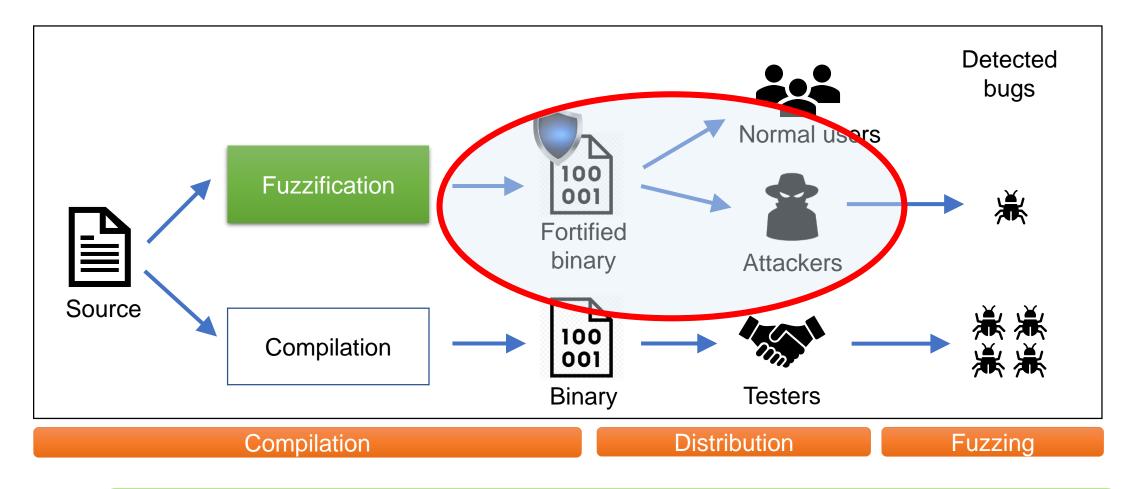
Our work: Make the Fuzzing Only Effective to the Testers



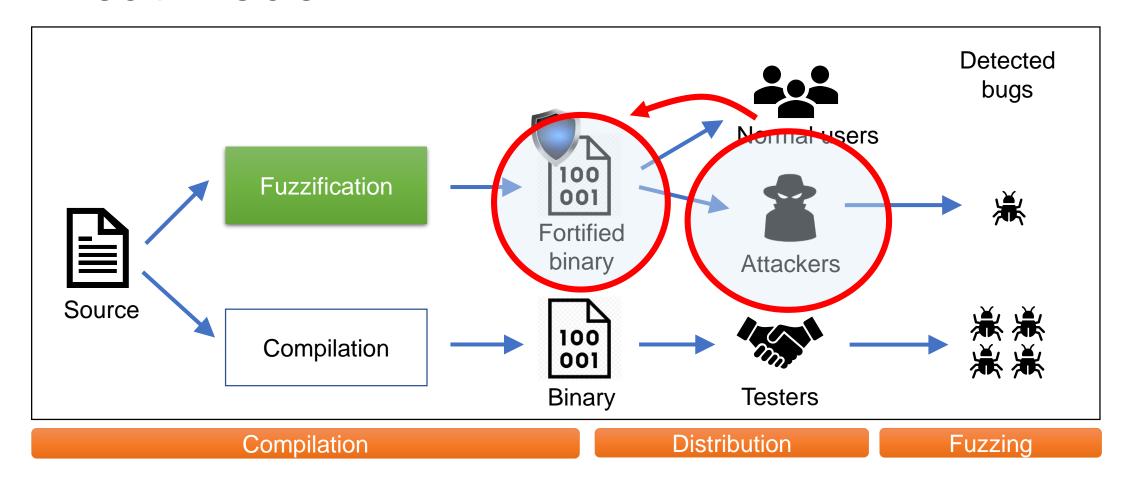




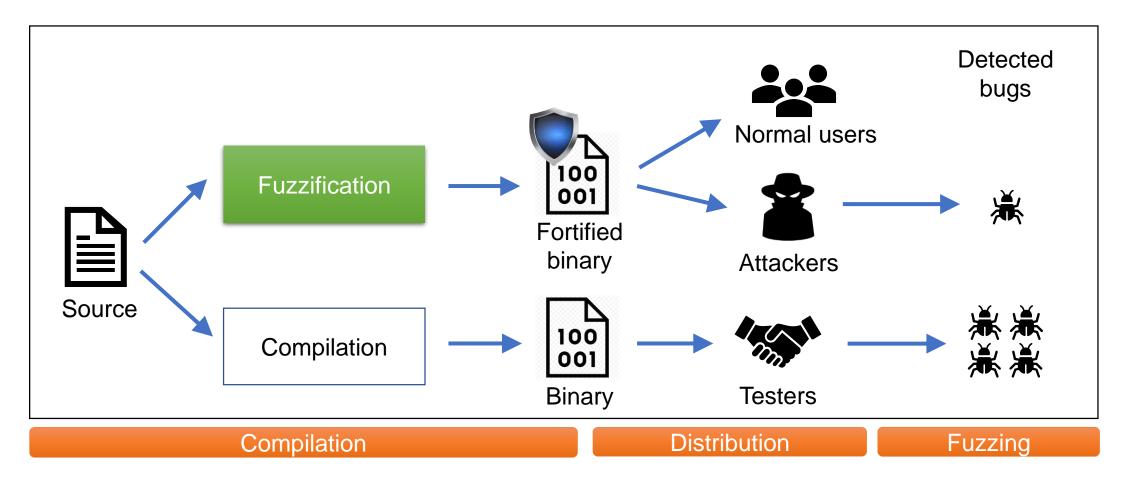
Adversaries try to find vulnerabilities from fuzzing

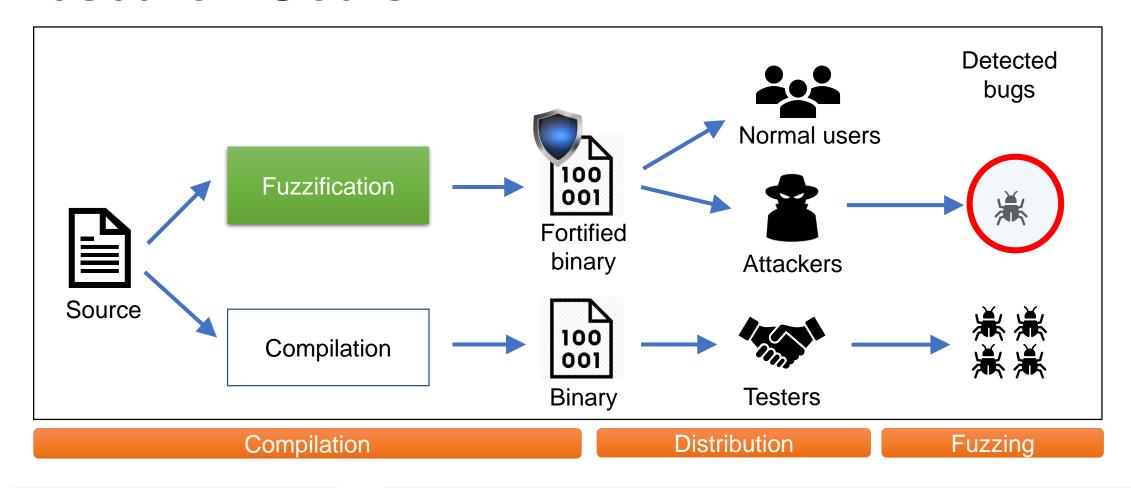


Adversaries only have a copy of fortified binary



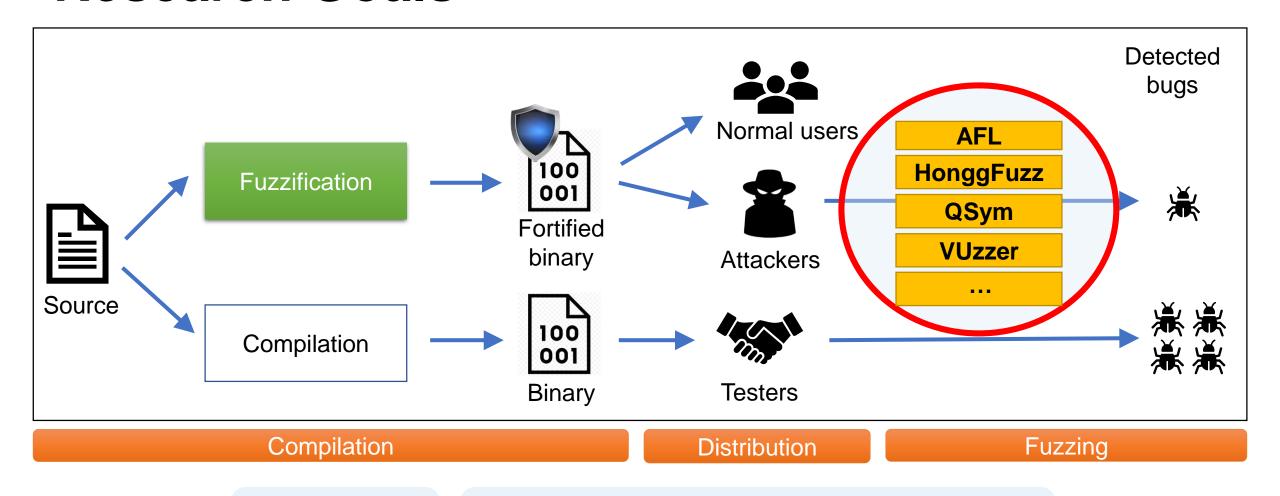
Adversaries know Fuzzification and try to nullify





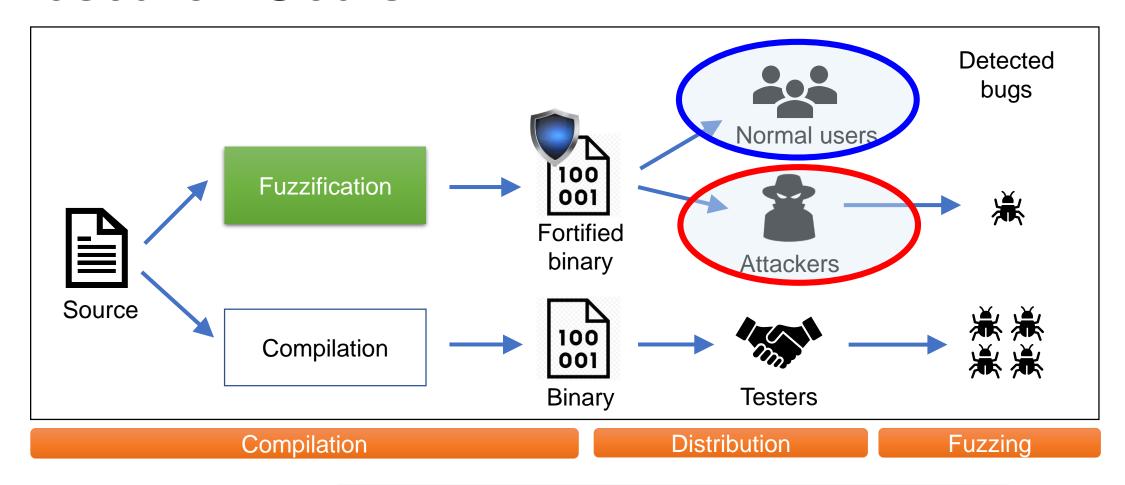
Hinder Fuzzing

Reduce the number of detected bugs



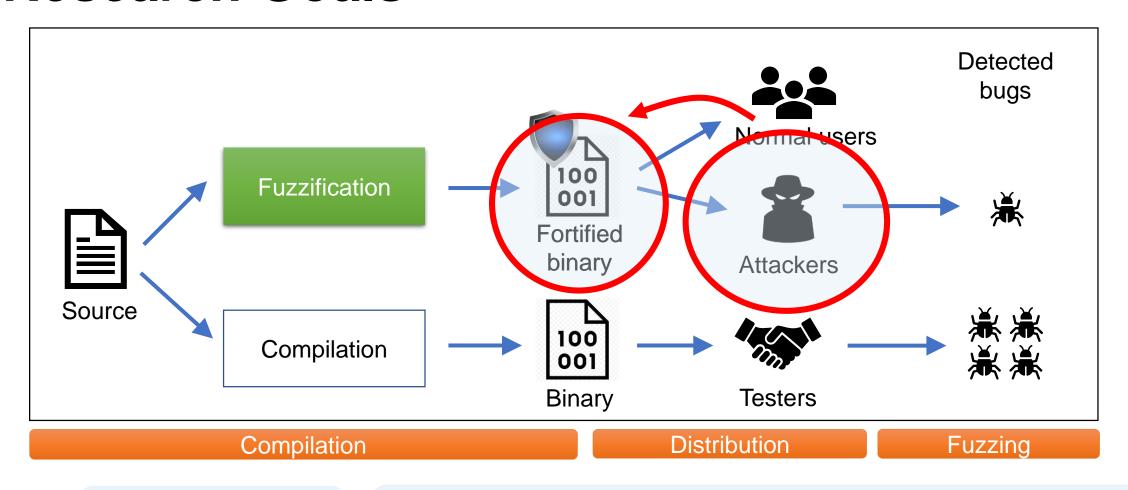
Generic

Affect most of the fuzzers



Overhead

Low overhead to normal user High overhead to attackers



Resiliency

Resilient to the adversarial analysis

Method	Generic to most fuzzers	Low overhead	Resilient to adversary
Packing or obfuscation	0	X	Ο

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Bug injection	Ο	Ο	X

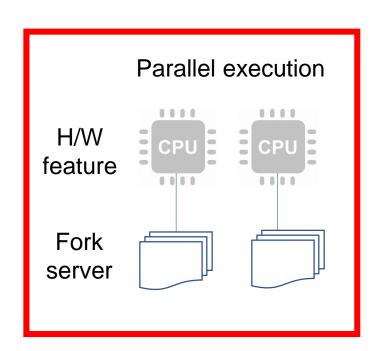
Method	Generic to most fuzzers	_	Resilient to adversary
Packing or obfuscation	Ο	X	0
Bug injection	Ο	O	X
Fuzzer detection	X	Ο	X

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Packing or obfuscation	O	X	O
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Fuzzer detection	X	0	X
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Fuzzification	Ο	0	Ο

Fast execution

Coverage-guidance

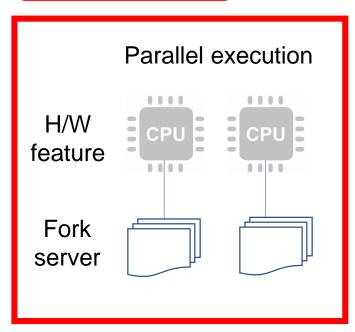


Fast execution

Coverage-guidance

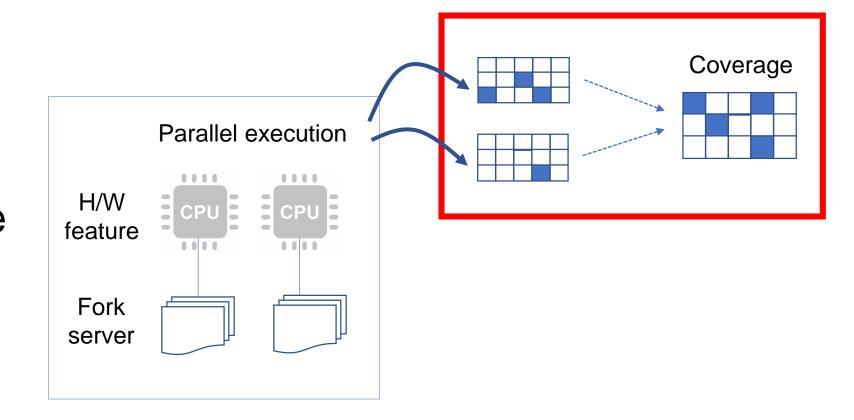
Hybrid approach

SpeedBump



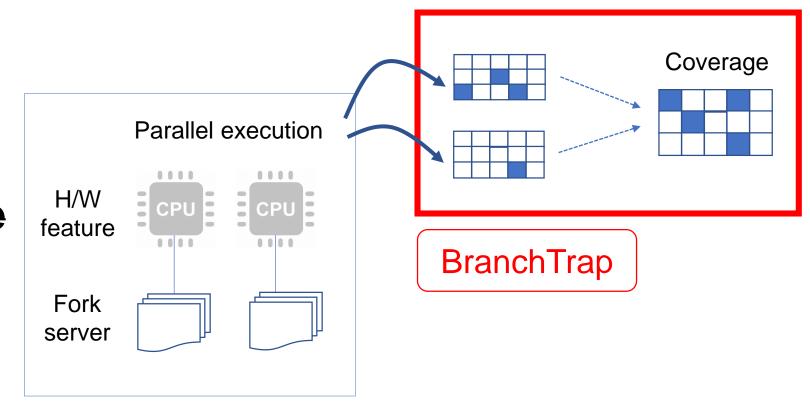
Fast execution

Coverage-guidance



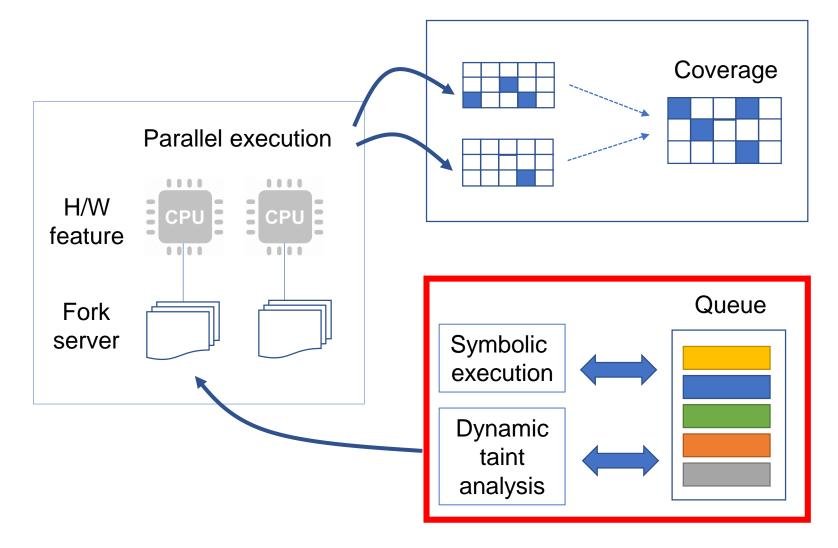
Fast execution

Coverage-guidance



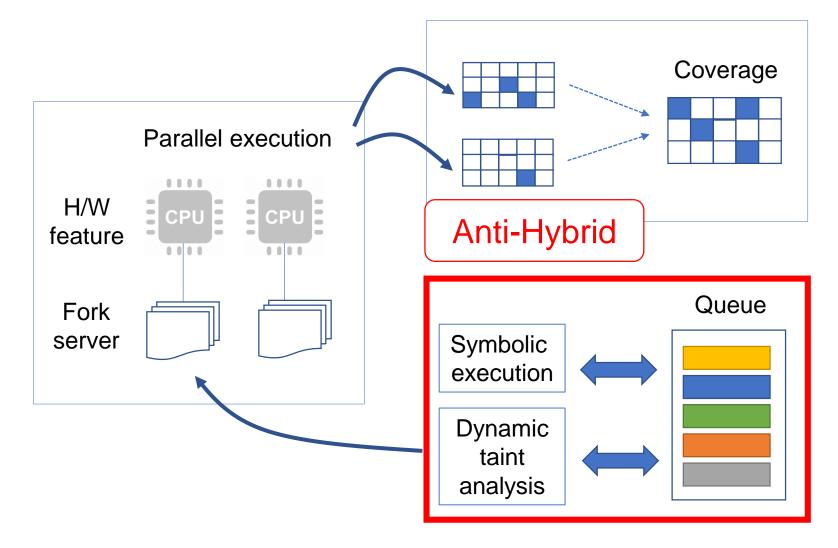
Fast execution

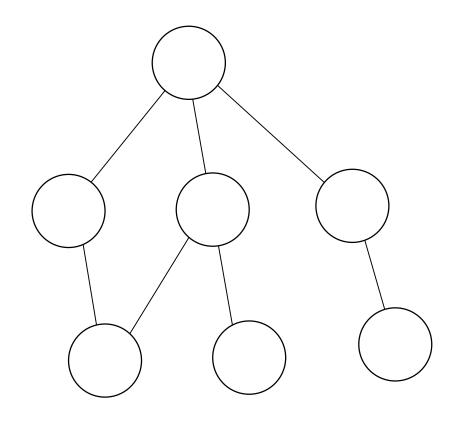
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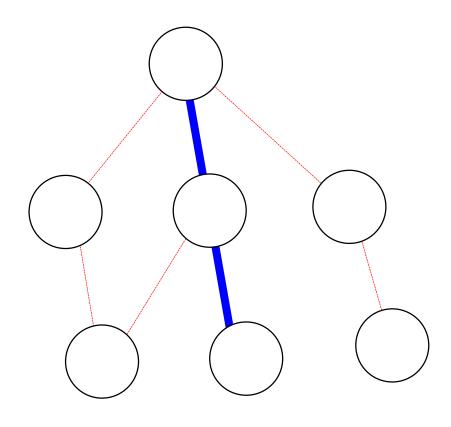
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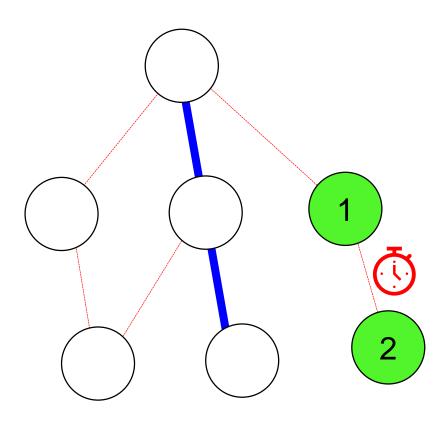


Basic block



 Identify frequently and rarely visited paths

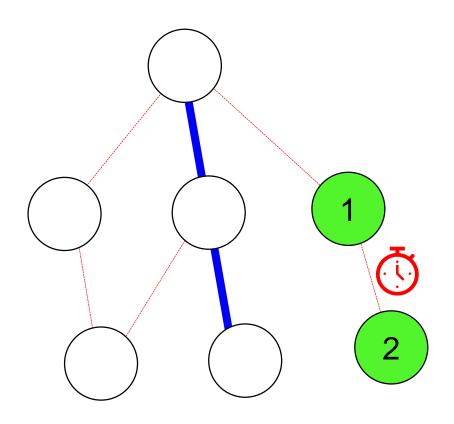
Basic blockRarely visited pathFrequently visited path



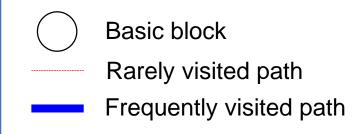
 Identify frequently and rarely visited paths

 Inject delays from the most rarely visited edges

Basic block
Rarely visited path
Frequently visited path



- Why this is effective?
 - User: follows common paths
 - Attacker: searches for new paths
 - → Impact of delay is more significant to attackers



SpeedBump: How to delay?

- Strawman: using sleep()
 - trivially removed by adversary

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 - Use randomly generated code
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SpeedBump: How to delay?

- Strawman: using sleep()
 - trivially removed by adversary

- Counter to advanced adversary
 - Use randomly generated code
 - avoid static-pattern
 - Impose control-flow and data-flow dependency
 - avoid automated analysis

```
int rarely_executed_code ()
{
    return 0;
}
```

```
int rarely_executed_code ()
    return 0;
//define global variables
int global1 = 1;
int global2 = 2;
int rarely_executed_code ()
    //inject delay function
    int pass = 20;
    global2 = func(pass);
    return 0;
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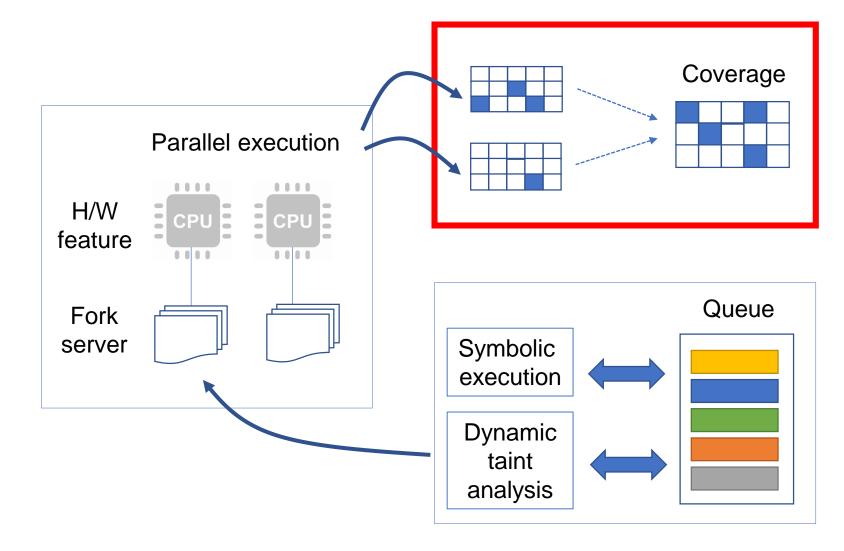
```
int func(int p6) {
  int local1[10];
 // affect global1 variable
 global1 = 45;
  int local2 = global1;
 for (int i = 0; i < 1000; i++)
  // affect local1 variable
   local1[i] = p6 + local2 + i;
 // affect global2 variable
  return local1[5];
```

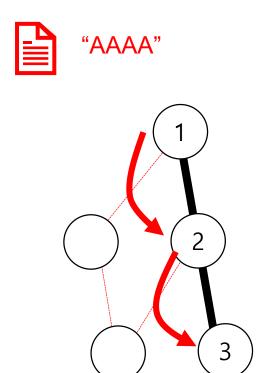
BranchTrap Hinders Coverage Management

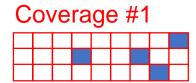
Fast execution

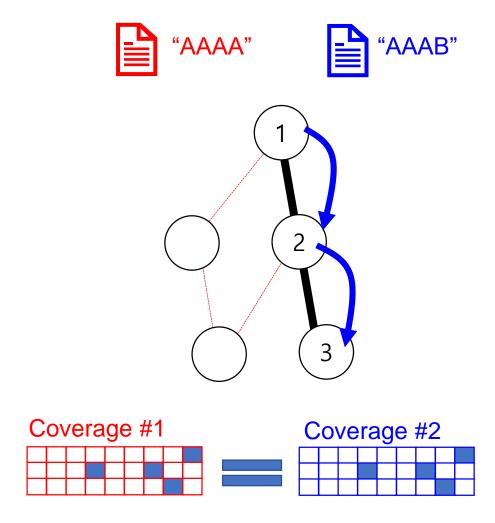
Coverage-guidance

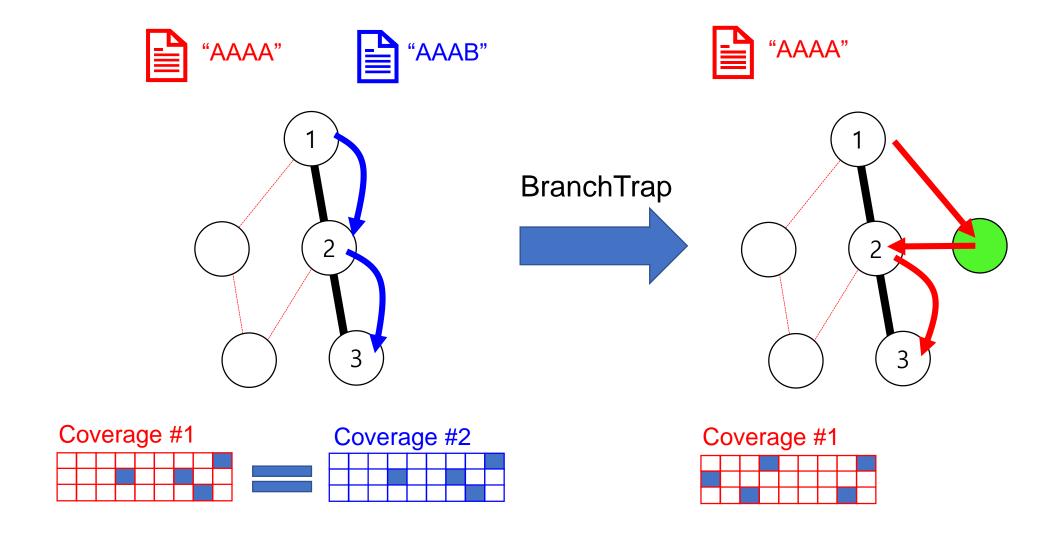
Hybrid approach

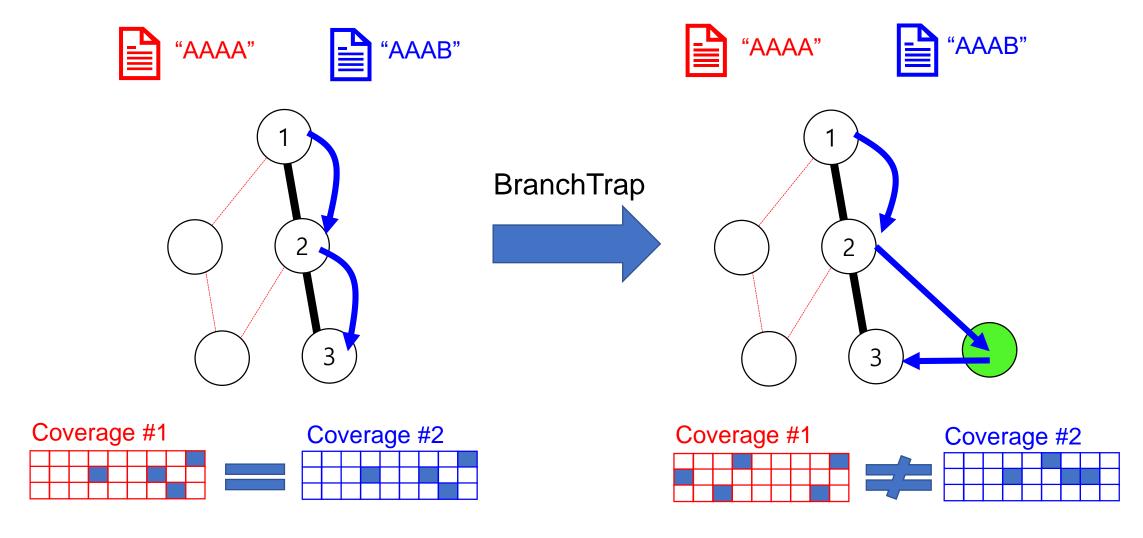






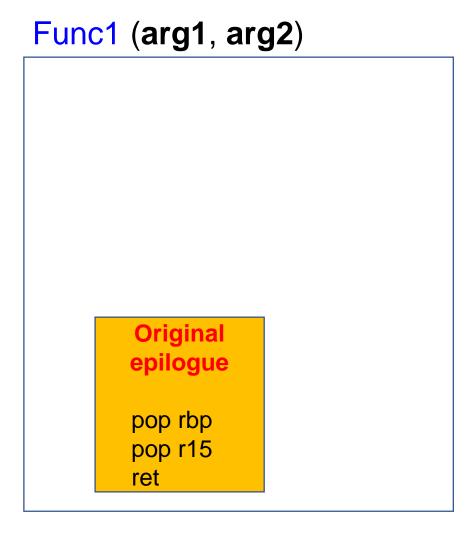






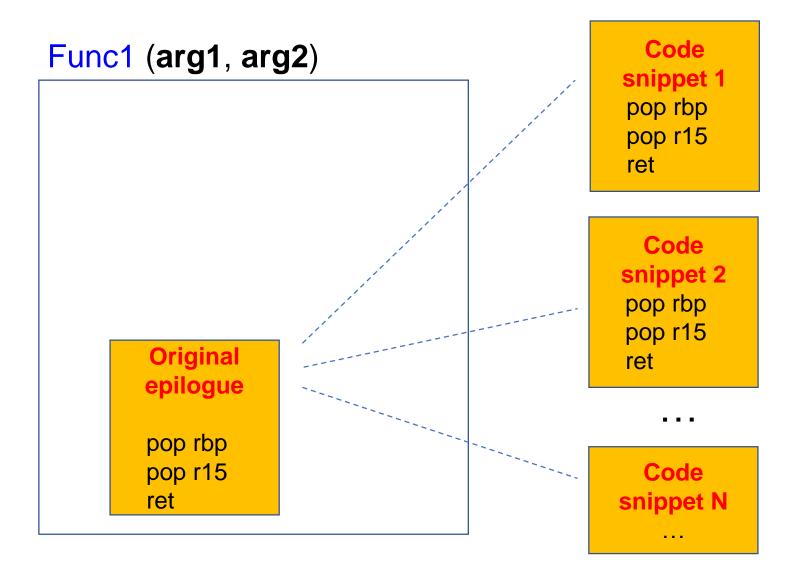


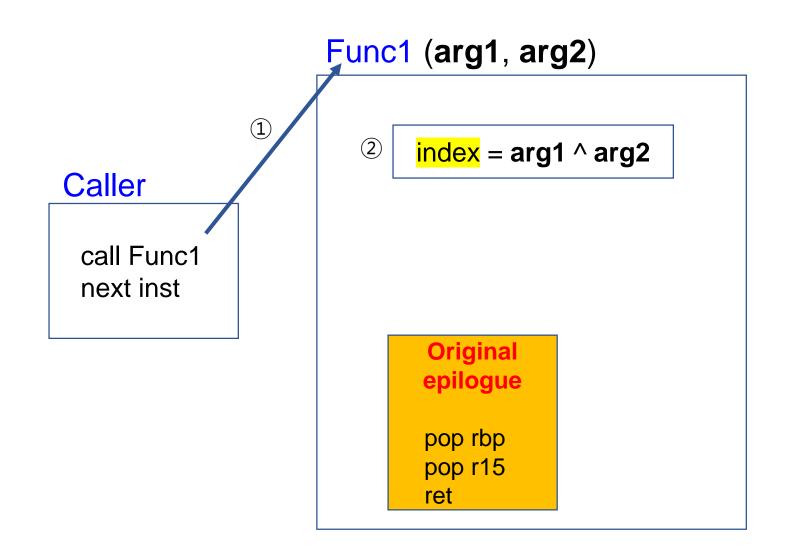
call Func1 next inst



Caller

call Func1 next inst



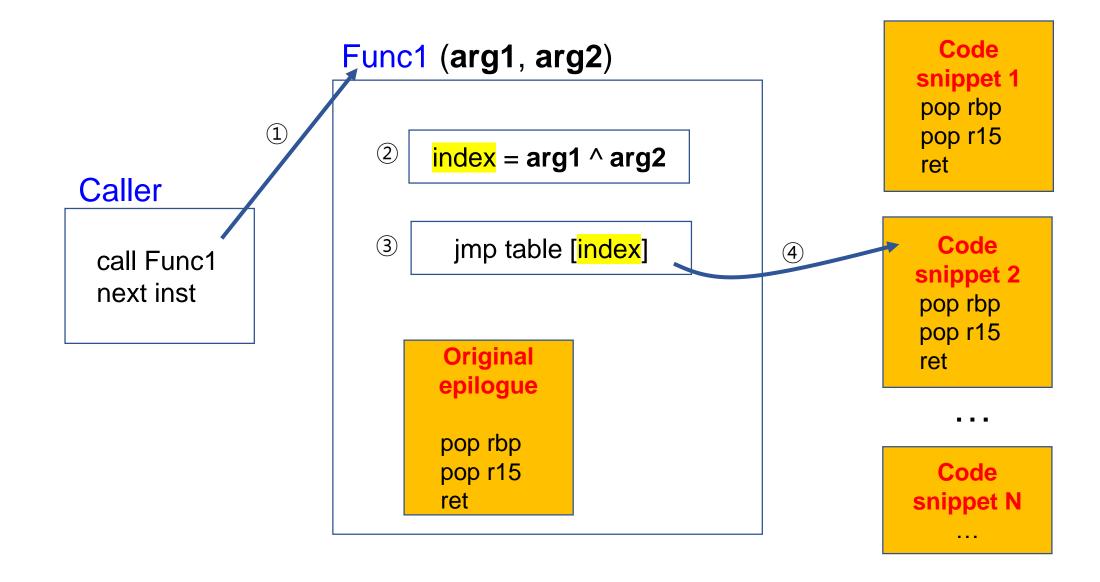


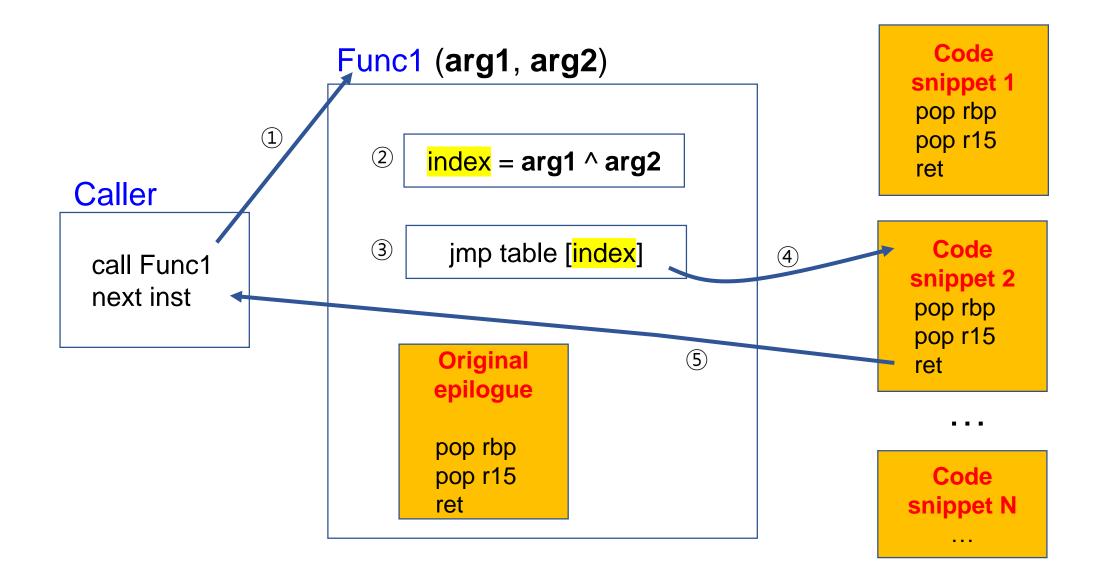
Code
snippet 1
pop rbp
pop r15
ret

Code snippet 2 pop rbp pop r15 ret

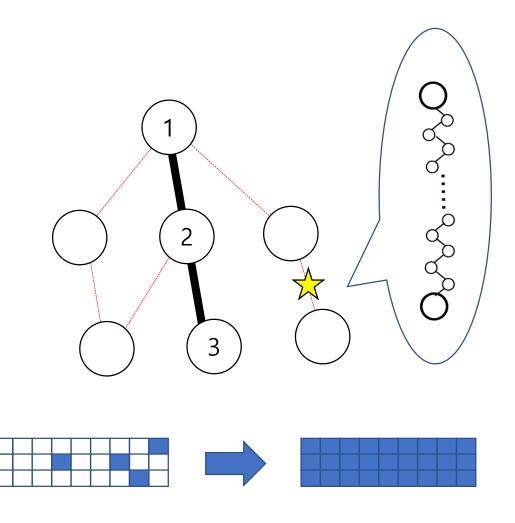
- -

Code snippet N





BranchTrap#2: Saturate Feedback State



One-time visit makes effect

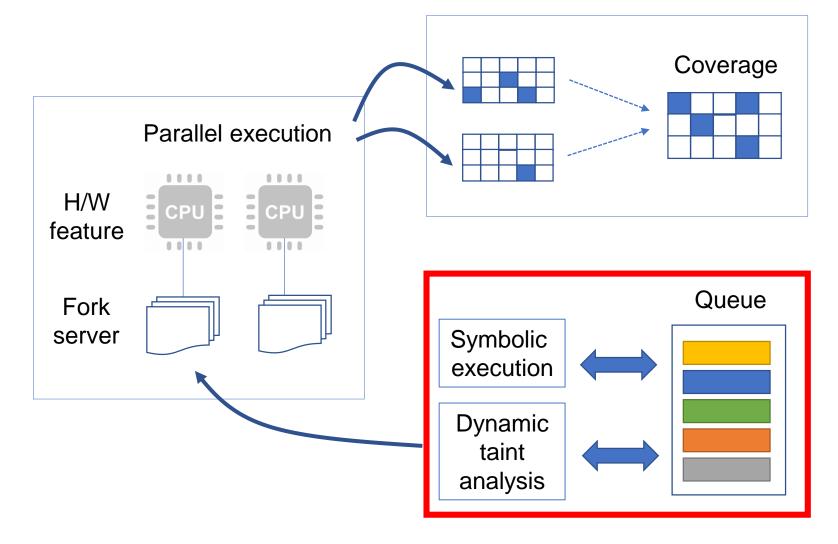
- BranchTrap:
 - Saturates bitmap data
 - Prevents coverage recording

AntiHybrid Hinders Hybrid Fuzzing

Fast execution

Coverage-guidance

Hybrid approach



Challenge of Hybrid Fuzzing

- Dynamic taint analysis
 - Expensive implicit flow

Transform <u>explicit</u> data-flow → <u>implicit</u> data-flow

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- Dynamic taint analysis
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Transform <u>explicit</u> data-flow → <u>implicit</u> data-flow

- Symbolic execution
 - Path explosion

Introduce an arbitrary path explosions

AntiHybrid Avoids Dynamic Taint Analysis

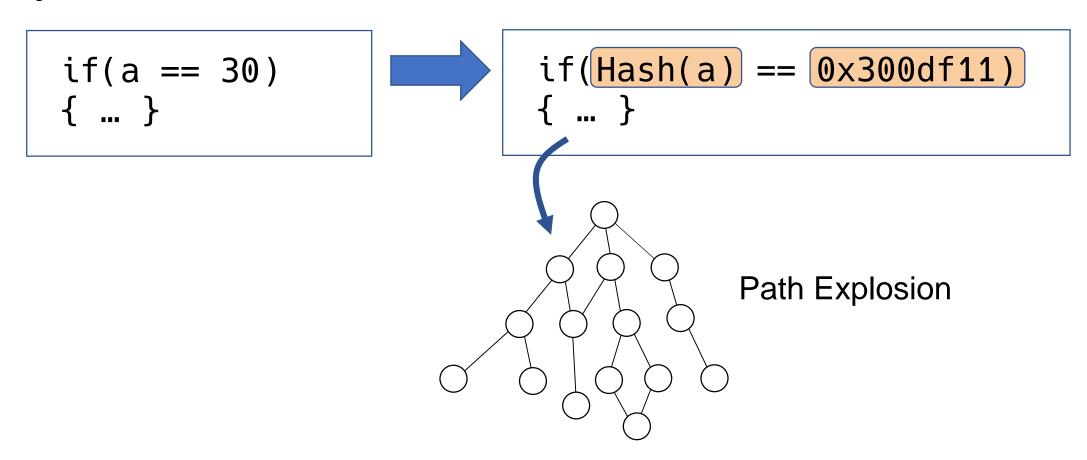
Transform explicit data-flow to implicit data-flow

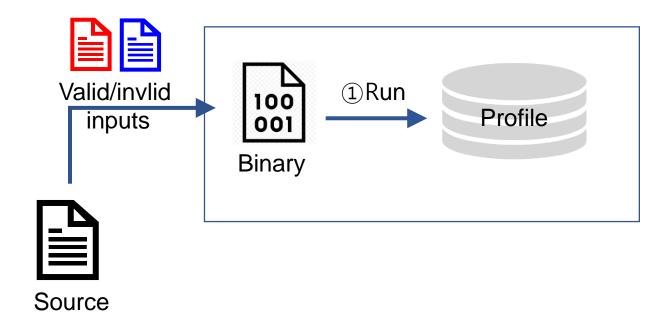
```
char input = 'a';
if (!strcmp(input, 'a'))
{ ... }
                       input
              Unable to
                taint
```

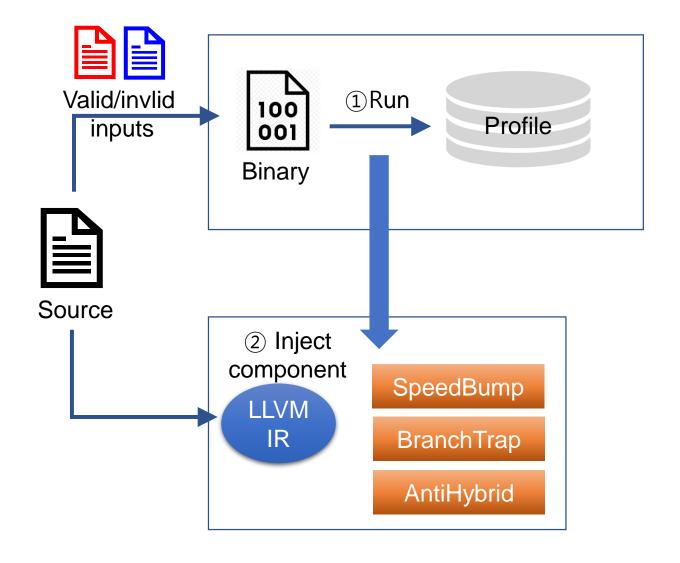
```
char input = 'a';
char anti_dta;
if (input == 97)
  anti dta = 'a';
if (!strcmp(anti_dta, 'a'))
```

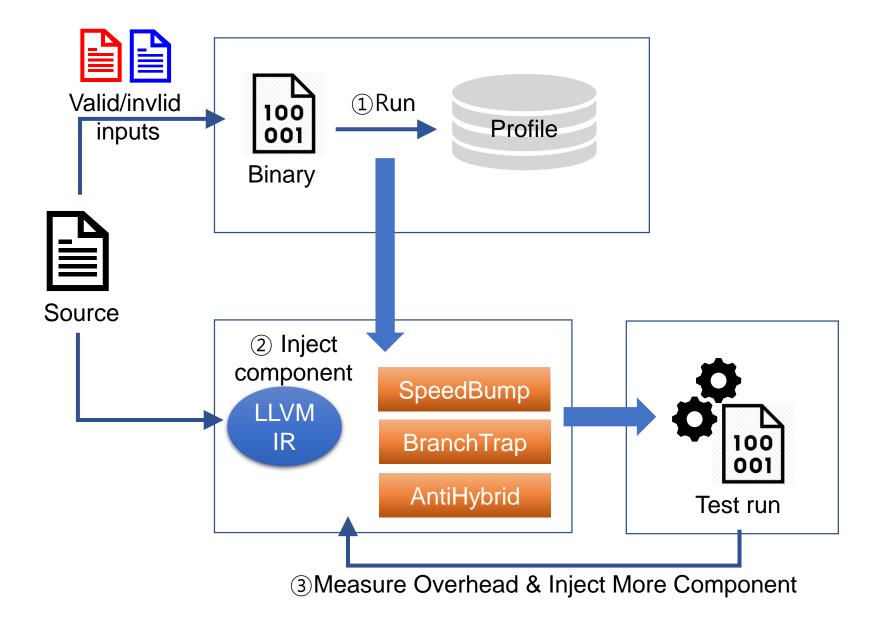
AntiHybrid Incurs Path Explosions

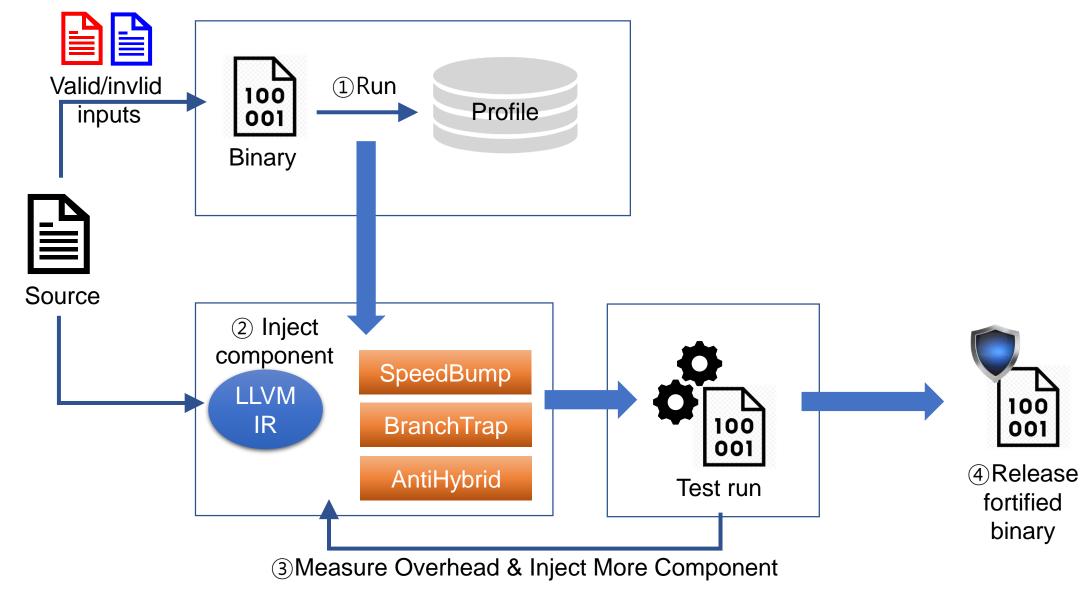
Inject hash calculations into branches









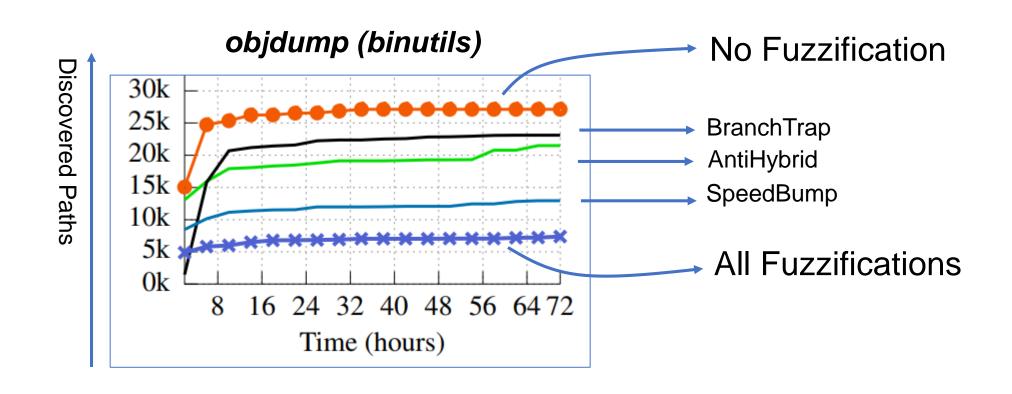


Evaluation Summary

- Implementation
 - □ 6,599 lines of Python and 758 lines of C++

- Evaluation questions:
 - Effective in "Reducing discovered paths and bugs?"
 - Effective on "Various fuzzers?
 - Impose "Low overhead" to the normal user?

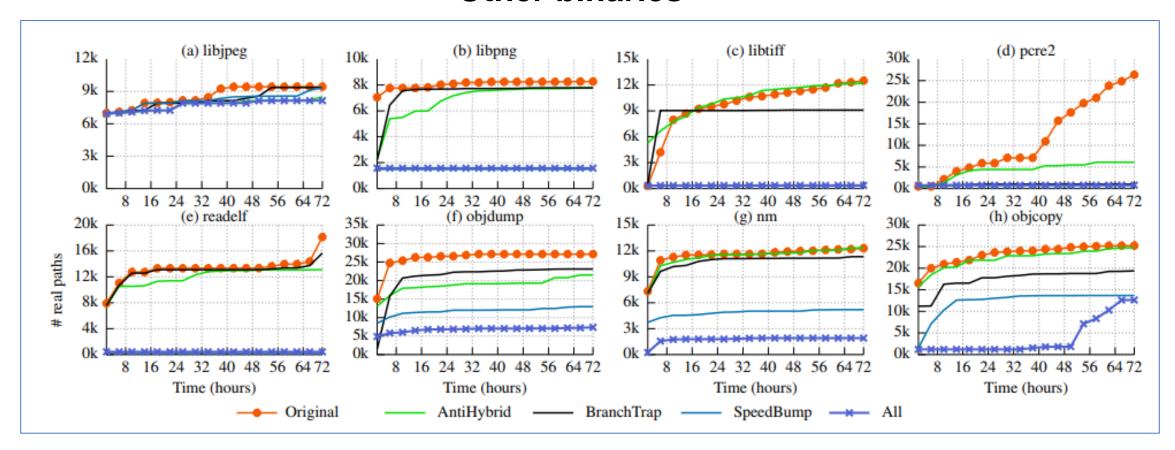
Reduced the Discovered Coverage By 71%



^{*} Fuzzing result on AFL-QEMU

Reduced the Discovered Coverage By 71%

Other binaries



^{*} Fuzzing result on AFL-QEMU

Fuzzification is Effective on Various Fuzzers

Fuzzer	Result
AFL (QEMU)	74%
HonggFuzz (PT)	61%
QSym (AFL-QEMU)	80%
Average	71%

Reduced code coverage

Reduced the Discovered Bugs

Result

88%

98%

94%

93%

binutils v2.3.0

Fuzzer

Average

AFL (QEMU)

HonggFuzz (PT)

QSym (AFL-QEMU)

LAVA-M dataset

Fuzzer	Result
Vuzzer	56%
QSym (AFL-QEMU)	78%
Average	67%

File size & CPU Overheads

binutils v2.3.0

OverheadResultFile Size1.4MB (62.1%)CPU Overhead3.7%

Real-world applications (e.g., GUI)

Overhead	Result
File Size	1.3MB (5.4%)
CPU Overhead	0.73%

^{*} Both overheads are configurable

Discussion

Best-effort protections against adversarial analysis

- Complementary to other defense techniques
 - Not hiding all vulnerabilities
 - But introducing significant cost on attacker

Comparison: Fuzzification vs. AntiFuzz

Component	Fuzzification	AntiFuzz
Delay execution	• (+ cold path)	
Fake coverage	(randomized return)	• (fake code)
Saturate coverage		
Prevent crash		
Anti-hybrid	• (+ anti-DTA)	
Countermeasures		64

Conclusion

Make the fuzzing only effective to the testers

- SpeedBump: Inject delays and only affects attackers
- BranchTrap: Insert input-sensitive branches
- AntiHybrid: Hinder hybrid fuzzing techniques

https://github.com/sslab-gatech/fuzzification