## DDA4230 Reinforcement learning

Iterative methods

## Lecture 11

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## 1 Goal of this lecture

In this lecture we will extend the iterative approaches for discrete MDPs, including policy evaluation, policy iteration, value iteration, and the their connections to Q-learning. **Suggested reading**: Chapter 3, 4 and 5 of *Reinforcement learning*: An introduction; Chapter 1 and 2 of *Reinforcement learning*: Theory and algorithms.

## 2 Recap: Solving discrete MDPs

For policy evaluation, by the Bellman equation we can obtain the value function as

$$V = (I - \gamma P)^{-1} r.$$

When P and r are unknown, a model-based methods is to estimate them before computing V.

For value optimization, by the Bellman optimality equation it amounts to solve

minimize 
$$\mathbf{e}^T V$$
  
subject to  $(I - \gamma P_j)V - r_j \ge 0$ ,  $j = 1, \dots, m$ ,

where e is the all-one vector. The dual of the linear program describes policy optimization

maximize 
$$\sum_{\lambda_1,\dots,\lambda_m} \lambda_j^T r_j$$
 subject to 
$$\sum_{j} (I - \gamma P_j^T) \lambda_j = \mathbf{e},$$
 
$$\lambda_i > 0, \quad j = 1,\dots,m.$$

When P and r are unknown, a model-based methods is to estimate them before computing V and  $\lambda_i$ .

## 3 Iterative policy evaluation

Recall that the iterative policy evaluation algorithm construct a contraction when  $\gamma < 1$ , which gives an arbitrarily close value function estimation of a given policy.

## **Algorithm 1:** Iterative policy evaluation

```
Input: Policy \pi, threshold \epsilon > 0
Output: Value function estimation V \approx V^*
Initialize \Delta > \epsilon and V arbitrarily
while \Delta > \epsilon do
\begin{array}{c|c} \Delta = 0 \\ \text{for } s \in \mathcal{S} \text{ do} \\ v = V(s) \\ V(s) = \sum_a \pi(a|s) \sum_{s',r} \mathbb{P}(s',r|s,a) \left[r + \gamma V(s')\right] \\ \Delta = \max(\Delta, |v - V(s)|) \end{array}
```

## 3.1 Dynamic programming

It worth mention that for a finite horizon MDP, the iterative policy evaluation algorithm requires the iteration to go through the index with a non-stationary value function. This process is also known as dynamic programming. Both the algorithms look surprisingly similar, to the point that it is hard to tell the difference. For dynamic programming, by the Bellman equation,

$$V_t(s) = R(s) + \gamma \sum_{s' \in S} P(s' \mid s, \pi) V_{t+1}(s') , \forall t = 0, \dots, T - 1,$$
  
$$V_T(s) = 0 .$$

For episodic MDPs, R and P can be stochastic and we run this process for many episodes (usually denoted as T/H episodes with horizon H).

We leave it to the reader to prove equations (3.1) for a finite horizon Markov process.

These equations immediately lend themselves to a dynamic programming solution whose pseudo-code is outlined in Algorithm 2. The algorithm takes as input a finite horizon Markov reward process  $\mathcal{S}, \mathbb{P}, \mathcal{R}$ , and computes the value function for all states and at all times.

#### **Algorithm 2:** Iterative policy evaluation with finite horizon

```
Input: S, \mathbb{P}, \mathcal{R}, T

For all states s \in S, V_T(s) \leftarrow 0

t \leftarrow T - 1

while t \ge 0 do

For all states s \in S, V_t(s) = \sum_a \pi(a|s) \sum_{s',r} \mathbb{P}(s',r|s,a) [r + \gamma V_{t+1}(s')]

t \leftarrow t - 1

return V_t(s) for all s \in S and t = 0, \dots, T
```

# 4 Iterative policy search

The policy evaluation algorithm immediately renders itself to a brute force algorithm called policy search to find the optimal value function  $V^*$  and an optimal policy  $\pi^*$ , as described

in pseudo-code in Algorithm 3. The algorithm takes as input an infinite horizon MDP  $\mathcal{M} = (\mathcal{S}, \mathcal{A}, \mathbb{P}, \mathcal{R}, \gamma)$  with arbitrary initial state distribution  $\rho_0$  and a tolerance  $\epsilon$  for accuracy of policy evaluation, and returns the optimal value function and an optimal policy.

```
Algorithm 3: Policy search

Input: \mathcal{M}, \epsilon
\Pi \leftarrow \text{All stationary deterministic policies of M}
\pi^* \leftarrow \text{Randomly choose a policy } \pi \in \Pi
V^* \leftarrow \text{POLICY EVALUATION } (\mathcal{M}, \pi^*, \epsilon)
for \pi \in \Pi do
V^{\pi} \leftarrow \text{POLICY EVALUATION } (\mathcal{M}, \pi, \epsilon)
if V^{\pi}(s) \geq V^*(s) \ \forall \ s \in S then
V^* \leftarrow V^{\pi}
```

It is clear that Algorithm 3 always terminates as it checks all  $|\Pi| = |\mathcal{A}|^{|\mathcal{S}|}$  deterministic stationary policies. Thus the run-time complexity of this algorithm is  $O(|\mathcal{A}|^{|\mathcal{S}|})$ . It is possible to prove correctness of the algorithm when  $\epsilon = 0$ , i.e. when in each iteration the policy evaluation is done exactly. In practice  $\epsilon$  is set to a small number such as  $10^{-9}$  to  $10^{-12}$ .

Claim 1 Algorithm 3 returns the optimal value function and an optimal policy when  $\epsilon = 0$ .

**Proof:** Let  $\pi^*$  be an optimal policy, and thus  $V^{\pi^*}(s) = V^*(s)$  for all states  $s \in S$ . Since the algorithm checks every policy in  $\Pi$ , it means that  $\pi^*$  must get selected at some iteration of the algorithm. Thus for the policies considered in future iterations the value function can no longer strictly increase. Future iterations may select a different policy with the same optimal value function, thus completing the proof.

For the MDP examples shown in LN3, we leave it to the reader to discuss (a) How many deterministic stationary policies does the agent have? (b) If  $\gamma < 1$ , is the optimal policy unique? (c) If  $\gamma = 1$ , is the optimal policy unique?

#### 4.1 Policy iteration

We now discuss a more efficient algorithm than policy search called policy iteration. The algorithm is a straightforward application of the Bellman operator, which states that given any stationary policy  $\pi$ , we can find a deterministic stationary policy that is no worse than the existing policy. In particular the theorem also applies to deterministic policies. This simple step has a special name called policy improvement, whose pseudo-code is presented in Algorithm 4.

The output of Algorithm 4 is always guaranteed to be at least as good as the policy  $\pi$  corresponding to the input value function  $V^{\pi}$ , and represents a greedy attempt to improve the policy. When performed iteratively with the policy evaluation algorithm (Algorithm

## Algorithm 4: Policy improvement

```
Input: V^{\pi}
\hat{\pi}(s) \leftarrow \operatorname*{arg\,max}_{a \in A} \left[ R(s,a) + \gamma \sum_{s' \in S} P(s'|s,a) V^{\pi}(s') \right] , \ \forall \ s \in S
return \hat{\pi}(s) for all s \in S
```

1), this gives rise to the policy iteration algorithm. The pseudo-code of policy iteration is outlined in Algorithm 5.

```
Algorithm 5: Policy iteration
```

The proof of correctness of Algorithm 5 is left to the reader as an exercise. Note that the algorithm will always terminate as there are a finite number of stationary deterministic policies, as a conclusion of the following claim. The reader should also establish the worst case run-time complexity of this algorithm.

Claim 2 Consider an infinite horizon MDP with  $\gamma < 1$ . The following statements hold.

- 1. When Algorithm 5 is run with  $\epsilon = 0$ , it finds the optimal value function and an optimal policy.
- 2. If the policy does not change during a policy improvement step, then the policy cannot improve in future iterations.
- 3. The value functions corresponding to the policies in each iteration of the algorithm form a non-decreasing sequence for every  $s \in S$ .

## 5 Value iteration

Value iteration is yet another technique that can be used to compute the optimal value function and an optimal policy, given an MDP. To motivate this method we will need the following theorem.

Suppose we are now given an MDP  $\mathcal{M} = (\mathcal{S}, \mathcal{A}, \mathbb{P}, \mathcal{R}, \gamma)$  and consider the finite dimensional Banach space  $\mathbb{R}^{|S|}$  equipped with the infinity norm  $||\cdot||_{\infty}$ . Then for every element  $U \in \mathbb{R}^{|S|}$  the Bellman optimality backup operator  $B^*$  is defined as

$$(B^*U)(s) = \max_{a \in A} \left[ R(s, a) + \gamma \sum_{s' \in S} P(s'|s, a)U(s') \right] , \quad \forall s \in S.$$
 (1)

This operator converts the LFS of the Bellman optimality equation to its RHS, while the argument is not necessarily the optimal value function.

**Theorem 3** For a MDP with  $\gamma < 1$ , let the fixed point of the Bellman optimality backup operator  $B^*$  be denoted by  $V^* \in \mathbb{R}^{|S|}$ . Then the policy given by

$$\pi^*(s) = \underset{a \in A}{\arg \max} \left[ R(s, a) + \gamma \sum_{s' \in S} P(s'|s, a) V^*(s') \right] , \ \forall \ s \in S ,$$
 (2)

is a stationary deterministic policy. The value function of this policy  $V^{\pi^*}$  satisfies the identity  $V^{\pi^*} = V^*$ , and thus  $V^*$  is also the fixed point of the operator  $B^{\pi^*}$ . In particular this implies that there exists a stationary deterministic policy  $\pi^*$  whose value function is the fixed point of  $B^*$ . Moreover,  $\pi^*$  is an optimal policy.

**Proof:** We start by noting that  $\pi^*$  as defined in (2) is a stationary deterministic policy, and so we can conclude that  $R^{\pi^*}(s) = R(s, \pi^*(s))$  and  $P^{\pi^*}(s'|s) = P(s'|s, \pi^*(s))$  for all  $s \in S$  and  $a \in A$ .

As  $V^*$  is the fixed point of  $B^*$ , we have  $B^*V^* = V^*$ . So using definition (1) of  $B^*$ , and (2) we can write

$$V^*(s) = \max_{a \in A} \left[ R(s, a) + \gamma \sum_{s' \in S} P(s'|s, a) V^*(s') \right]$$

$$= R(s, \pi^*(s)) + \gamma \sum_{s' \in S} P(s'|s, \pi^*(s)) V^*(s')$$

$$= R^{\pi^*}(s) + \gamma \sum_{s' \in S} P^{\pi^*}(s'|s) V^*(s')$$

$$= V^{\pi^*}(s)$$

for all  $s \in S$ , completing the proof of the first part of the theorem.

To prove that  $\pi^*$  is an optimal policy, we show that if an optimal policy exists then its value function must be a fixed point of the operator  $B^*$ . Assume that an optimal policy exists, then we can specify a stationary deterministic policy to be optimal, and let us denote it as  $\mu$  and the corresponding optimal value function as  $V^{\mu}$ . Suppose that  $V^{\mu}$  is not a fixed point of  $B^*$ . Then there exists  $s \in S$  such that  $V^{\mu}(s) \neq (B^*V^{\mu})(s)$ , which implies that  $V^{\mu}(s) > (B^*V^{\mu})(s)$ . Then the greedy policy  $\hat{\pi}$  induced by  $V^{\mu}(s)$  implies that there exists a policy  $\hat{\pi}$  which is strictly better than  $\mu$ , which leads to a contradiction. This proves that  $V^{\mu}$  must be the unique fixed point of  $B^*$ . Combining this fact with the first part implies

that  $V^*$  must be the optimal value function and  $\pi^*$  is an optimal policy. This completes the proof.

Theorem 3 suggests a straightforward way to calculate the optimal value function  $V^*$  and an optimal policy  $\pi^*$ . The idea is to run fixed point iterations to find the fixed point of  $B^*$ . Once we have  $V^*$ , an optimal policy  $\pi^*$  can be extracted using the arg max operator in the Bellman optimality equation. The pseudo-code of this algorithm is given in Algorithm 6, which takes as input an infinite horizon MDP  $\mathcal{M} = (\mathcal{S}, \mathcal{A}, \mathbb{P}, \mathcal{R}, \gamma)$  and a tolerance  $\epsilon$ , and returns the optimal value function and an optimal policy.

### **Algorithm 6:** Value iteration

```
Input: \epsilon
For all states s \in S, V'(s) \leftarrow 0, V(s) \leftarrow \infty
while ||V - V'||_{\infty} > \epsilon do
 \begin{array}{c} V \leftarrow V' \\ \text{For all states } s \in S, \ V'(s) = \max_{a \in A} \left[ R(s,a) + \gamma \ \sum_{s' \in S} P(s'|s,a) V(s') \right] \\ V^* \leftarrow V \text{ for all } s \in S \\ \pi^* \leftarrow \arg\max_{a \in A} \left[ R(s,a) + \gamma \ \sum_{s' \in S} P(s'|s,a) V^*(s') \right] \ , \ \forall \ s \in S \end{array}
return V^*(s), \pi^*(s) for all s \in S
```

If Algorithm 6 is run with  $\epsilon = 0$ , we can recover the optimal value function and an optimal policy exactly. However in practice,  $\epsilon$  is set to be a small number such as  $10^{-9}$ - $10^{-12}$ .

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