10/11

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Theory of Computation

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Outline

1. Regular Expressions (RE)
2. Connection Between REs and Regular Languages
3. Regular Grammars

Specifying Language

How do we specify languages?

- If language is finite, you can list all of its strings.
  - $L = \{a, aa, aba, aca\}$
- Descriptive:
  - $L = \{x \mid n_a(x) = n_b(x)\}$
- Using basic Language operations
  - $L = \{aa, ab\}^* \cup \{b\}\{bb\}^*$
- Regular languages are described using the last method

Regular Expressions

Regular expressions describe regular languages and the notation involves a combination of:

- Strings of symbols from some alphabet Σ
- Parentheses ()
- Operators +, ·, * 0次以上

or concate

Regular Expressions

Important thing to remember

- A regular expression is not a language
  ↳ 是一個 modeling 的工具
- A regular expression is used to describe a language.
- X It is incorrect to say that for a language L, L = (a + b + c)*
- O But it's okay to say that L is described by (a + b + c)*

Regular Expressions

All finite languages can be described by regular expressions

0次以上

or concate

Example: $(a \overset{\text{or}}{+} b \overset{\text{concate}}{\cdot} c)^*$ $\Longleftrightarrow$ {{a} U {bc}}*

describes the language

$$\{a, bc\}^* = \{\lambda, a, bc, aa, abc, bca, \ldots\}$$

Definition 3.1

Let Σ be a given alphabet. Then

1. ϕ, λ, and $a \in \Sigma$ are all regular expressions. These are called primitive regular expressions.
2. If r_1 and r_2 are regular expressions, so are r_1+r_2, $r_1 \cdot r_2$, r_1^* and (r_1). 以做出来的都是 regular expression
3. A string is a regular expression ***iff*** it can be derived from the primitive regular expressions by a finite number of applications of the rules in (2).

Example 3.1

A regular expression: $(a+b\cdot c)^{*}\cdot(c+\emptyset)$

Not a regular expression: $(a + b +)$

a^{n}

a^{+}

Languages of Regular Expressions

$L(r)$: language of regular expression r

Example

$$L((a+b\cdot c)^*)=\{\lambda,a,bc,aa,abc,bca,...\}$$

Definition 3.2

- For primitive regular expressions:

$$L(\varnothing) = \varnothing \qquad (1)$$

$$L(\lambda) = \{\lambda\} \qquad (2)$$

$$L(a) = \{a\} \qquad (3)$$

Definition (continued)

For regular expressions r_1 and r_2

$$L(r_1 + r_2) = L(r_1) \cup L(r_2) \qquad (4)$$

$$L(r_1 \cdot r_2) = L(r_1)\, L(r_2) \qquad (5)$$

$$L(r_1 *) = (L(r_1))* \qquad (6)$$

$$L((r_1)) = L(r_1) \qquad (7)$$

Example 3.2

Regular expression: $(a+b)\cdot a^*$

$$
\begin{aligned}
L((a+b)\cdot a^*) &= L((a+b))\,L(a^*) \\
&= L(a+b)\,L(a^*) \\
&= (L(a)\cup L(b))\,(L(a))^* \\
&= (\{a\}\cup\{b\})\,(\{a\})^* \\
&= \{a,b\}\,\{\lambda,a,aa,aaa,\ldots\} \\
&= \{a,aa,aaa,\ldots,b,ba,baa,\ldots\}
\end{aligned}
$$

Priority of Operators

- Regular expression: $r = a \cdot b + c$

  ✓ $r_1 = a \cdot b \quad r_2 = c$ or ✗ $r_1 = a \quad r_2 = b + c$

  $L(r) = \{ab, c\} \neq \{ab, ac\}$

- Star closure (*) precedes concatenation (·) precedes union (+)

優先權

Example 3.3

$\Sigma=\{a,b\}$

- Regular expression $r=(a+b)^*(a+bb)$

↳ 結尾為 a or bb

Stands for any string of a's and b's

↳ 所有 string

$$L(r)=\{a,bb,aa,abb,ba,bbb,\ldots\}$$

L(r) is the set of all strings on {a,b}, terminated by either an a or a bb

Example 3.4

- Regular expression $r = (aa)^*(bb)^*b$

$$L(r) = \{a^{2n}b^{2m+1} : \quad n, m \geq 0\}$$

偶a奇b

L(r) is the set of all strings with an even number of a's followed by an odd number of b's

Example 3.5

- For $\Sigma = \{0, 1\}$, give a regular expression r such that

$L(r) = \{ w \in \Sigma^* : w$ has at least one pair of consecutive 0 $\}$

連續

$r = (0+1)^* 00 (0+1)^*$

Example 3.6

L(r) = { all strings with no pairs of consecutive 0s }

↳ 0的前後都要是1

- Regular expression (I) $\boxed{r = (1+01)^*(0+\lambda)}$

(II) $r = (1^* 0 1 1^*)^* \, (0+\lambda) + 1^*(0+\lambda)$

Add 1 immediately after a 0 — String ending in 0 — String with all 1's

There are an unlimited number of REs for any given language!

Equivalent Regular Expressions

Definition:

Regular expressions r_1 and r_2

are **equivalent** if $L(r_1) = L(r_2)$

Example

L = { all strings without two consecutive 0 }

$$r_1 = (1 + 01)^*(0 + \lambda)$$

$$r_2 = (1^*011^*)^*(0 + \lambda) + 1^*(0 + \lambda)$$

$L(r_1) = L(r_2) = L$ $\Rightarrow$ r_1 and r_2 are equivalent regular expressions

More Examples

- $L_1 = \{a, aa, aba, aca\}$
- $L_1 = \{a\} \cup \{aa\} \cup \{aba\} \cup \{aca\}$
- Regular expression describing L_1:

  $(a + aa + aba + aca)$

More Examples

(I) • $L_2 = \{x \in \{0,1\}^* \mid |x| \text{ is even}\}$

(II) • $L_2 = \{00, 01, 10, 11\}^*$

• Regular expressions describing L_2:

(III) $(00 + 01 + 10 + 11)^*$

(IV) $((0 + 1)(0 + 1))^*$

More Examples

- $L_3 = \{x \in \{0,1\}^* \mid x \text{ does not end in } 01\}$ 以非0.1結尾

  If x does not end in 01, then either

  $|x| < 2$ or x ends in 00, 10, or 11

- A regular expression that describes L_3 is:

  $(\lambda + 0 + 1) + (0 + 1)^*(00 + 10 + 11)$

More Examples

- $L_4 = \{x \in \{0,1\}^* \mid x$ contains an odd number of 0s $\}$

  Express $x = yz$

  y is a string of the form $y=1^i01^j$

  In z, there must be an even number of 0's

  $z = (01^k01^m)^*$

- A regular expression that describes L_4 is:

  $(1^*01^*)(01^*01^*)^*$

Short Quiz

- Give regular expressions for the following language on Σ= {a, b, c}.
  - All strings containing exactly one a

$r = (b+c)^*a(b+c)^*$

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1 Regular Expressions (RE)

2 Connection Between REs and Regular Languages

3 Regular Grammars

Theorem

$$\left\{\begin{array}{l}\text{Languages}\\\text{Described by}\\\text{Regular Expressions}\end{array}\right\} = \left\{\begin{array}{l}\text{Regular}\\\text{Languages}\end{array}\right\}$$

For every regular language there is a regular expression
For every regular expression there is a regular language

Kleene Theorem:
Regular expressions and Finite Automata are equivalent (w.r.t. the languages they describe/accept)

↳ Regular expression ≡ DFA ≡ NFA

Theorem - Part 1

$$\left\{\begin{array}{l}\text{Languages} \\ \text{Described by} \\ \text{Regular Expressions}\end{array}\right\} \subseteq \left\{\begin{array}{l}\text{Regular} \\ \text{Languages}\end{array}\right\}$$

1. For any regular expression r the language $L(r)$ is regular

■Theorem 3.1

Theorem - Part 2

$$\left\{\begin{array}{l}\text{Languages} \\ \text{Described by} \\ \text{Regular Expressions}\end{array}\right\} \supseteq \left\{\begin{array}{l}\text{Regular} \\ \text{Languages}\end{array}\right\}$$

2. For any regular language L there is a regular expression r with $L(r) = L$

■ Theorem 3.2

Proof - Part 1 $\left\{ \begin{array}{l} \text{Languages} \\ \text{Described by} \\ \text{Regular Expressions} \end{array} \right\} \subseteq \left\{ \begin{array}{l} \text{Regular} \\ \text{Languages} \end{array} \right\}$

1. For any regular expression r
   the language $L(r)$ is regular

If we have any regular expression r,
we can construct an NFA(DFA) that accepts L(r)

与有NFA or DFA来描述 L(r)

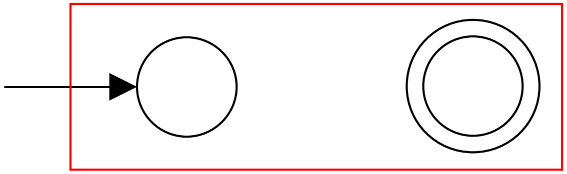
Proof by induction on the size of r

Induction Basis

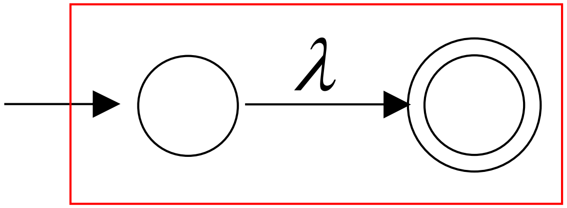
- Primitive Regular Expressions: $\emptyset, \quad \lambda, \quad a$

NFAs

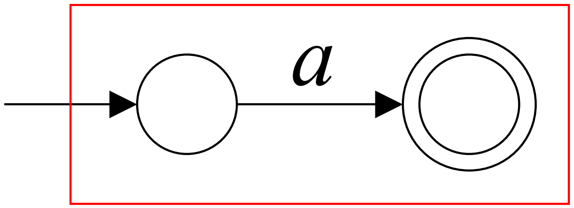
從Primitive Regular Expression 畫NFA



$$L(M_1) = \emptyset = L(\emptyset)$$



$$L(M_2) = \{\lambda\} = L(\lambda)$$



$$L(M_3) = \{a\} = L(a)$$

regular languages

Inductive Hypothesis

Assume

for regular expressions r_1 and r_2

that

$L(r_1)$ and $L(r_2)$ are regular languages

Inductive Step

∵ REs are derived from these four rules:

We will prove:

$$\left.\begin{array}{l} L(r_1 + r_2) \\ \\ L(r_1 \cdot r_2) \\ \\ L(r_1 *) \\ \\ L((r_1)) \end{array}\right\}$$

Are regular Languages

- By definition of regular expressions:

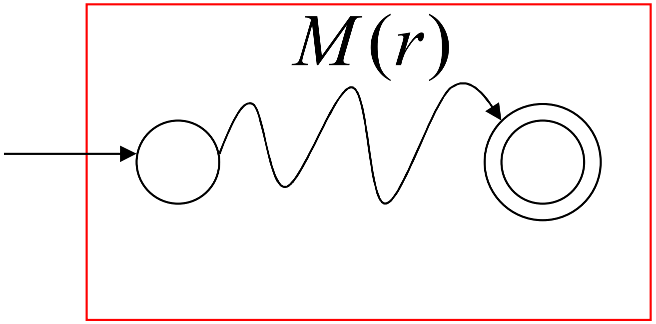
$$L(r_1 + r_2) = L(r_1) \cup L(r_2)$$

$$L(r_1 \cdot r_2) = L(r_1)\, L(r_2)$$

$$L(r_1 *) = (L(r_1))*$$

$$L((r_1)) = L(r_1)$$

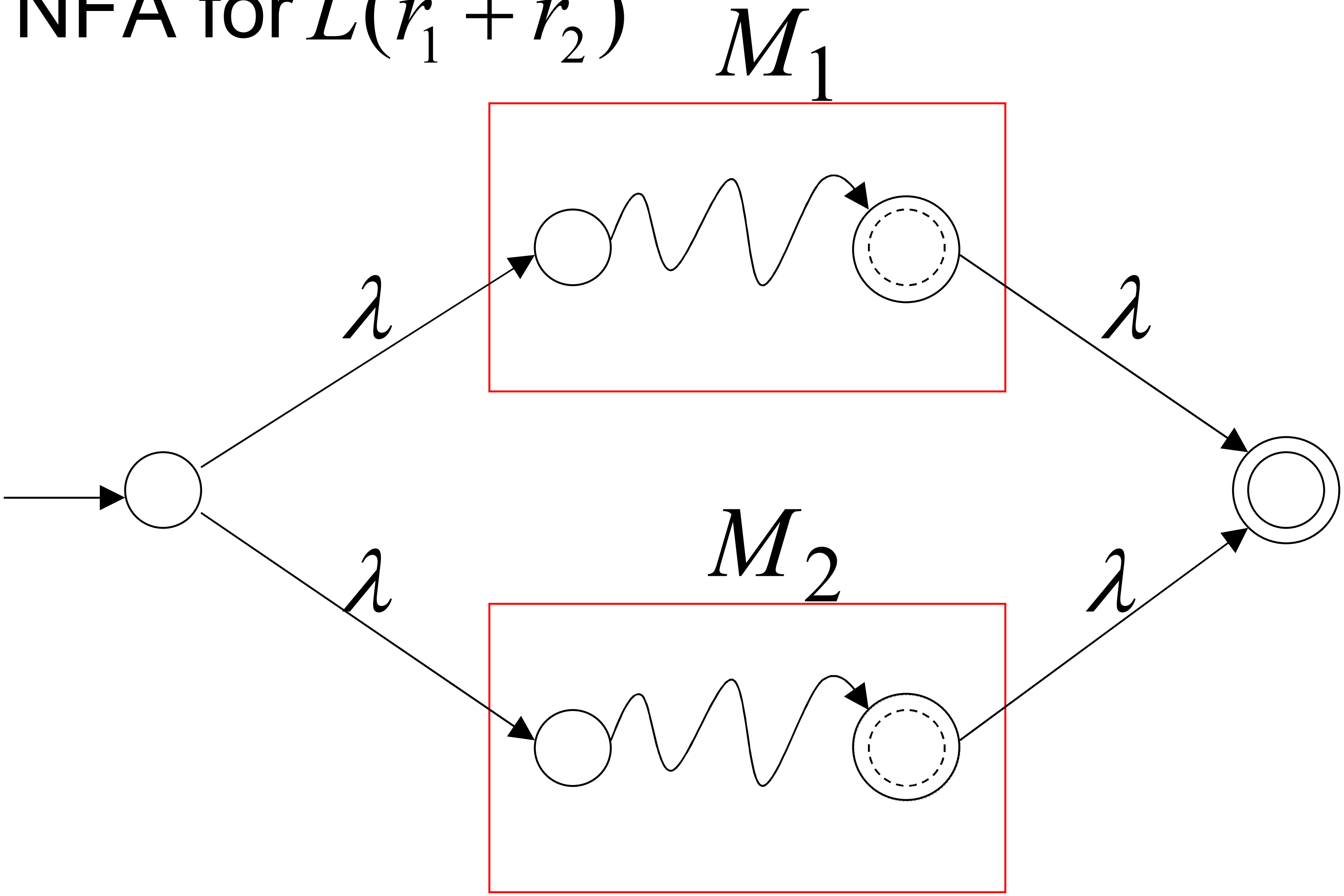
Schematic representation of an NFA (M(r)) accepting L(r)



We can claim that for every NFA there is only one final state (by exercise 7, section 2.3) ↳ NFA可畫成只有一個final

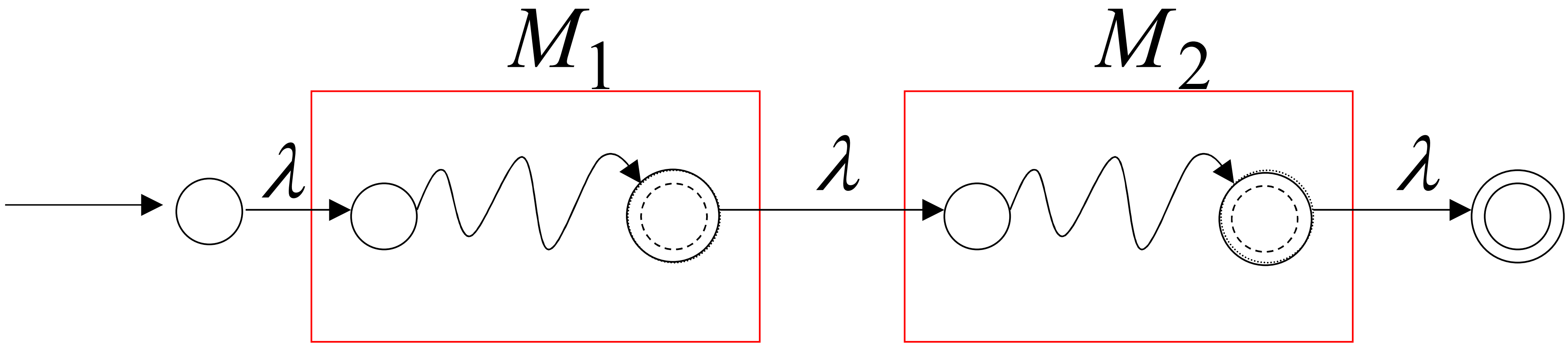
Union

- NFA for $L(r_1 + r_2)$



Concatenation

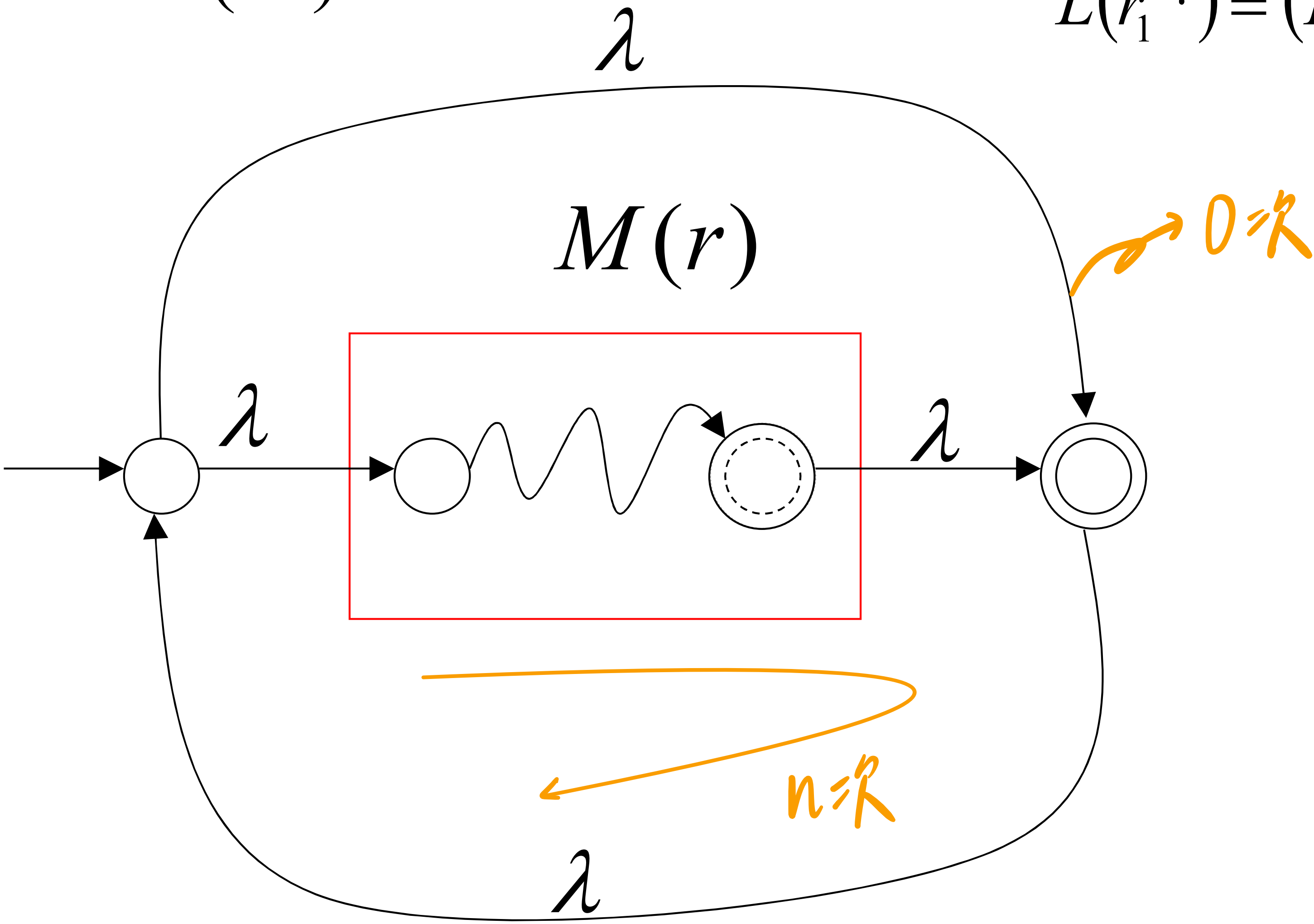
- NFA for $L(r_1 r_2)$



Star Operation

- NFA for $L(r^*)$

$$L(r_1{}^*) = (L(r_1))^*$$



By inductive hypothesis we know:

$L(r_1)$ and $L(r_2)$ are regular languages

We also know:

Regular languages are closed under:

Union	$L(r_1) \cup L(r_2)$
Concatenation	$L(r_1)\,L(r_2)$
Star	$(L(r_1))^*$

Therefore:

$$\left.\begin{aligned} L(r_1 + r_2) &= L(r_1) \cup L(r_2) \\[2em] L(r_1 \cdot r_2) &= L(r_1)\, L(r_2) \\[2em] L(r_1 *) &= (L(r_1))^* \end{aligned}\right\} \text{Are regular languages}$$

And trivially:

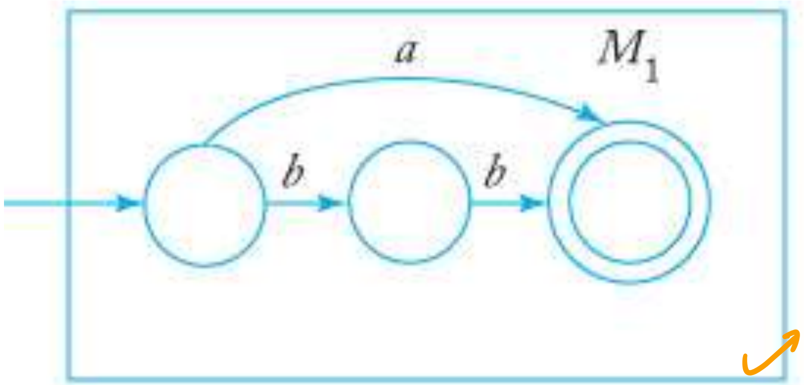
$L((r_1))$ is a regular language

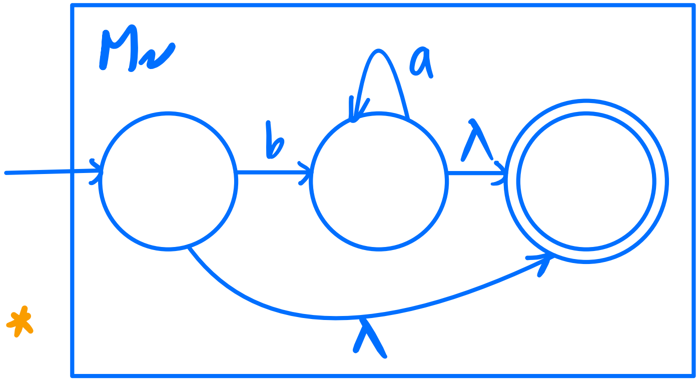
$\therefore$ For any regular expression r the language $L(r)$ is regular

Example 3.7

- Find an NFA that accepts L(r), where

$$r = (a + bb)^*(ba^* + \lambda)$$





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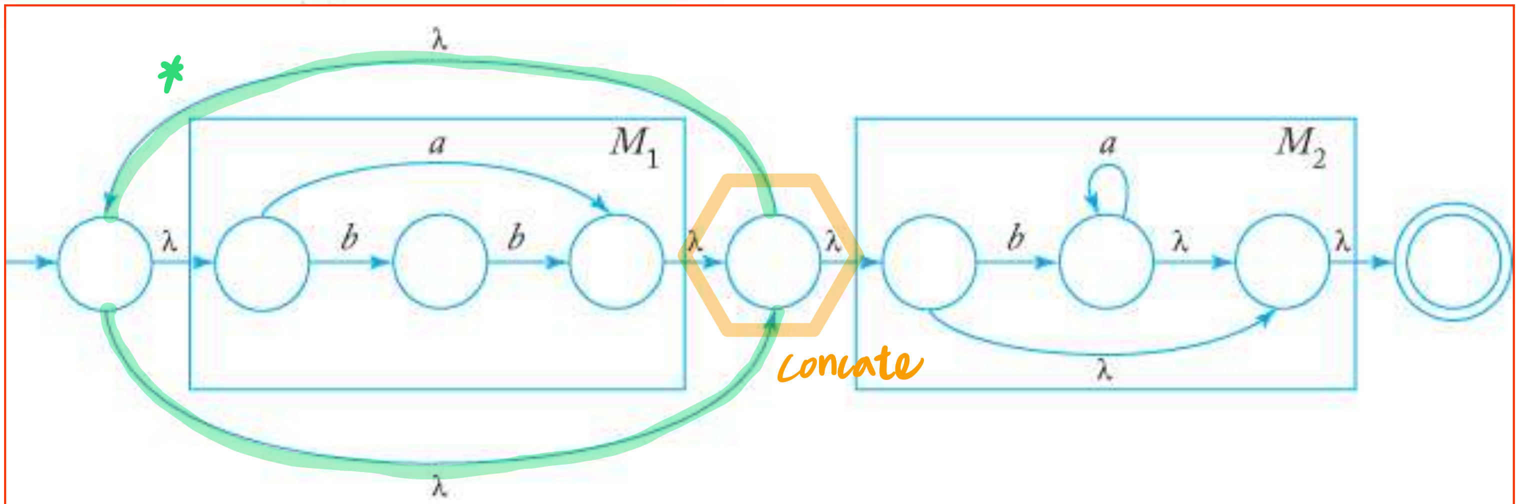


Proof – Part 2 $\left\{ \begin{array}{l}\text{Languages} \\ \text{Described by} \\ \text{Regular Expressions}\end{array} \right\} \supseteq \left\{ \begin{array}{l}\text{Regular} \\ \text{Languages}\end{array} \right\}$

2. For any regular language L there is a regular expression r with $L(r) = L$

Since any regular language has an associated NFA and hence a transition graph, all we need to do is to find a regular expression capable of generating the labels of all the walks from q_0 to any final state.

Proof by construction of regular expression

Generalized Transition Graphs (GTG)

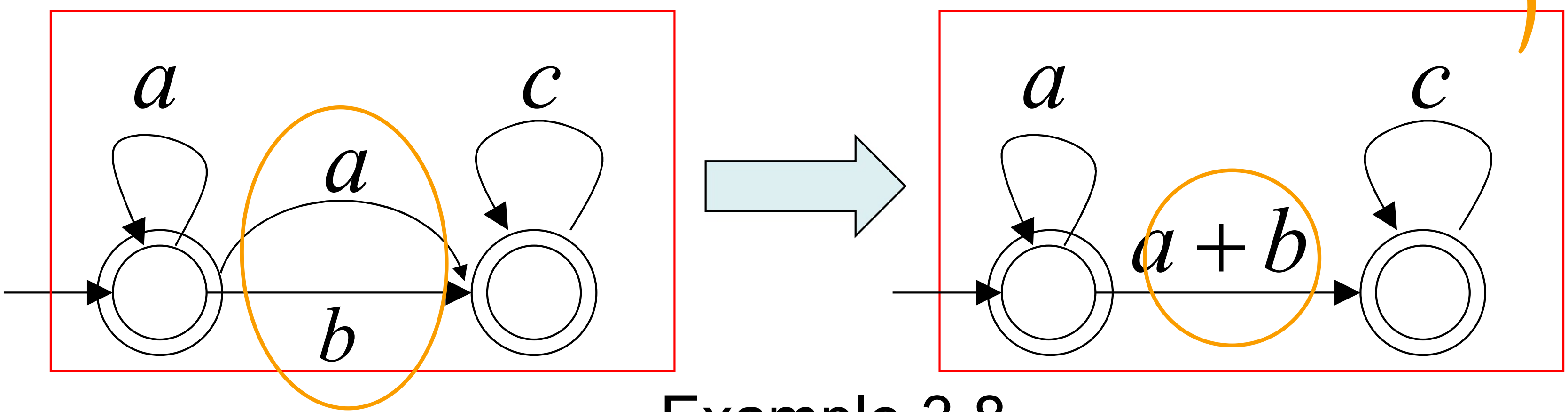
From M construct the equivalent

Generalized Transition Graph ↳ 把edge寫regular expression

in which transition labels are regular expressions

Example: M

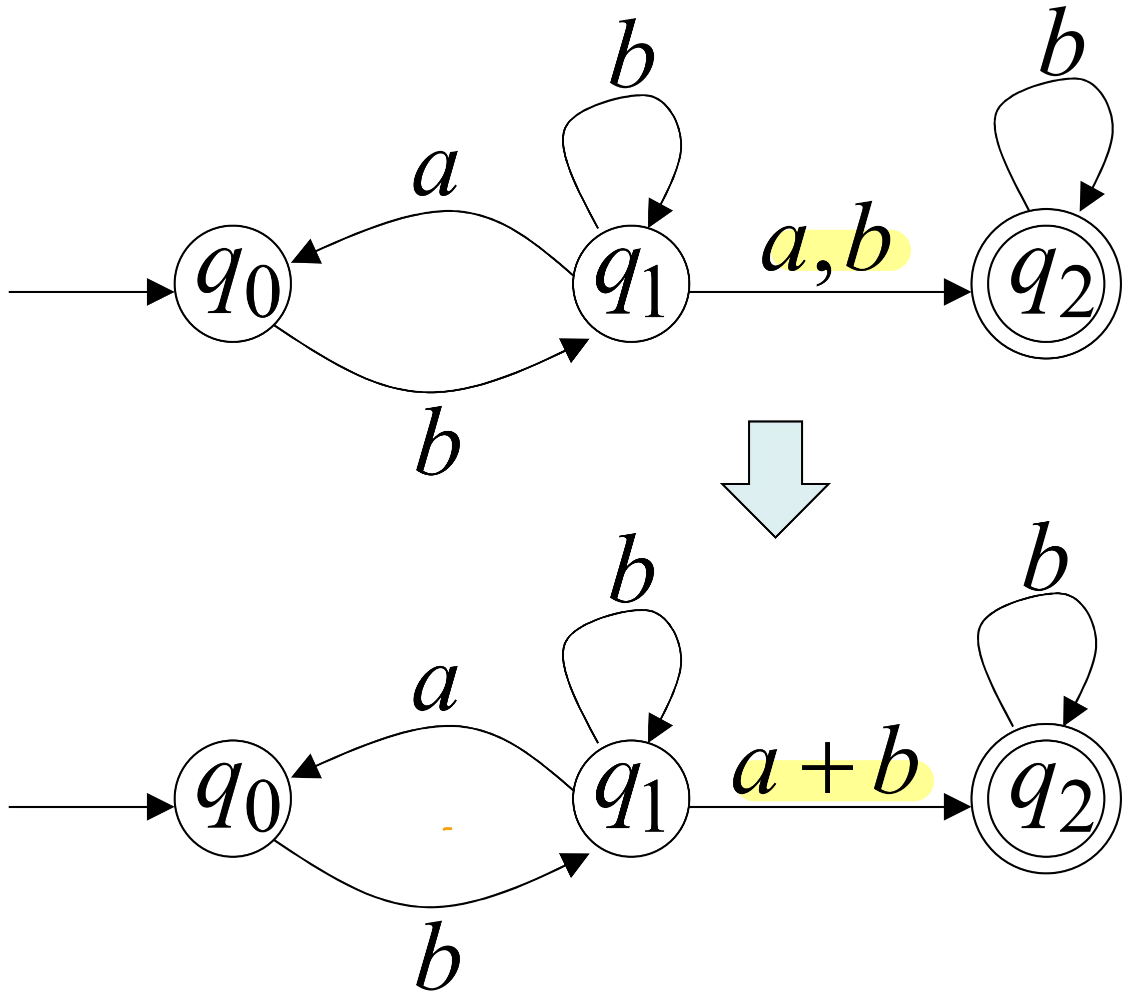
L(a* + a*(a + b)c*)



Example 3.8

GTG may have many states

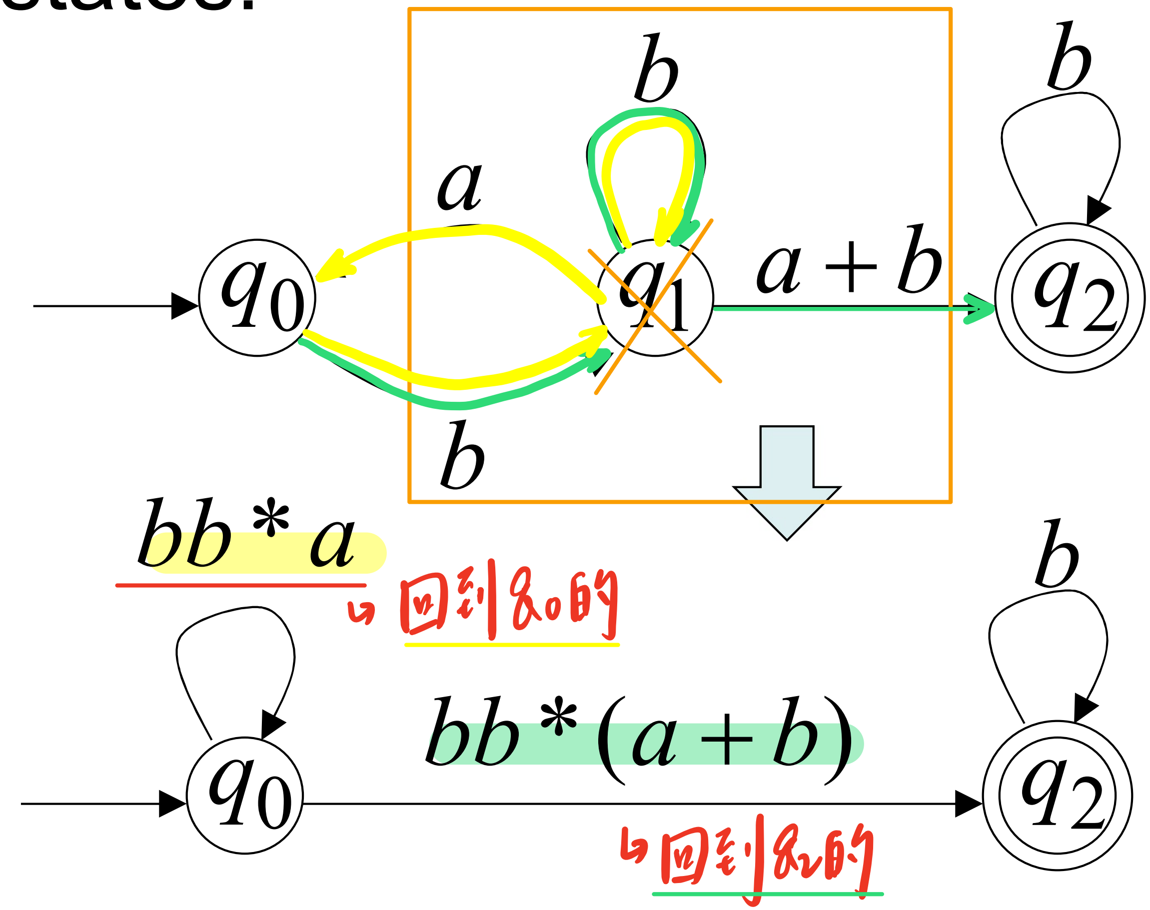
Enumerating all walks is time-consuming



Reducing the states:

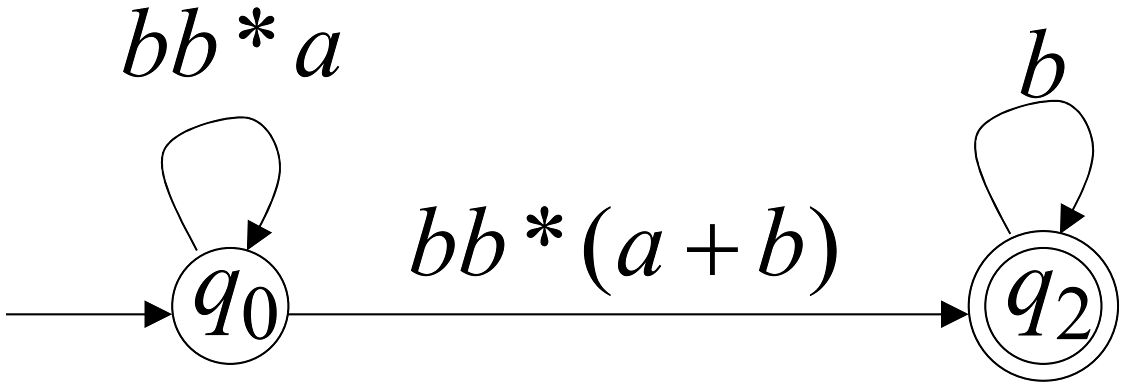
Ex. reduce q_1

分段處理



Simple two-state GTG

Resulting Regular Expression:



$$r=(bb^*a)^*bb^*(a+b)b^*$$

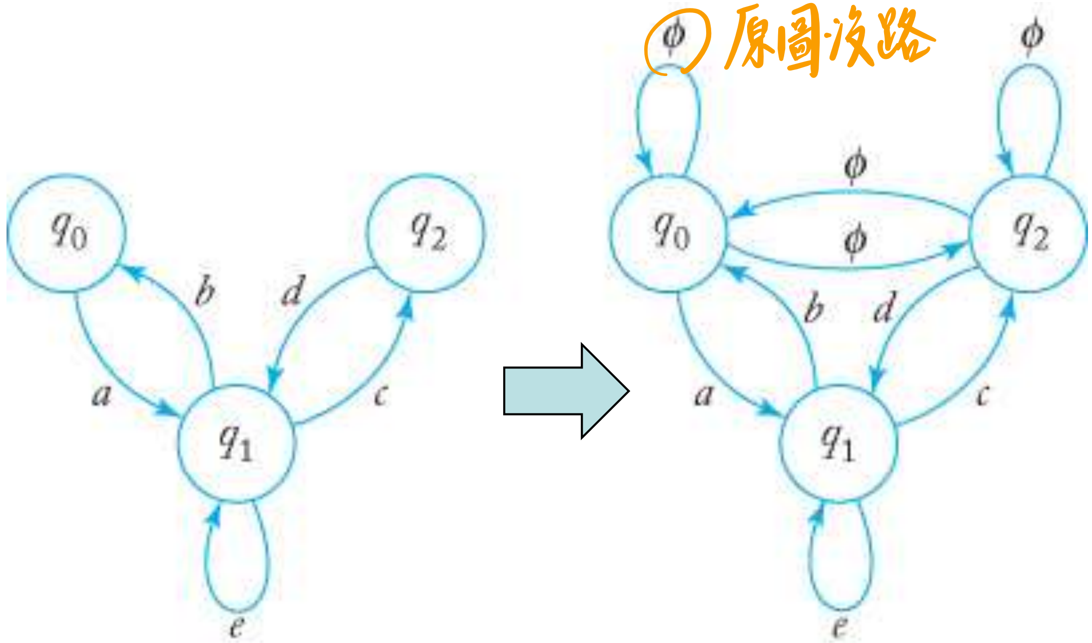
↳ reduce 到只剩 initial 和 final 即可寫出 regular expression

$$L(r)=L(M)=L$$

Complete GTG

GTG

Complete GTG

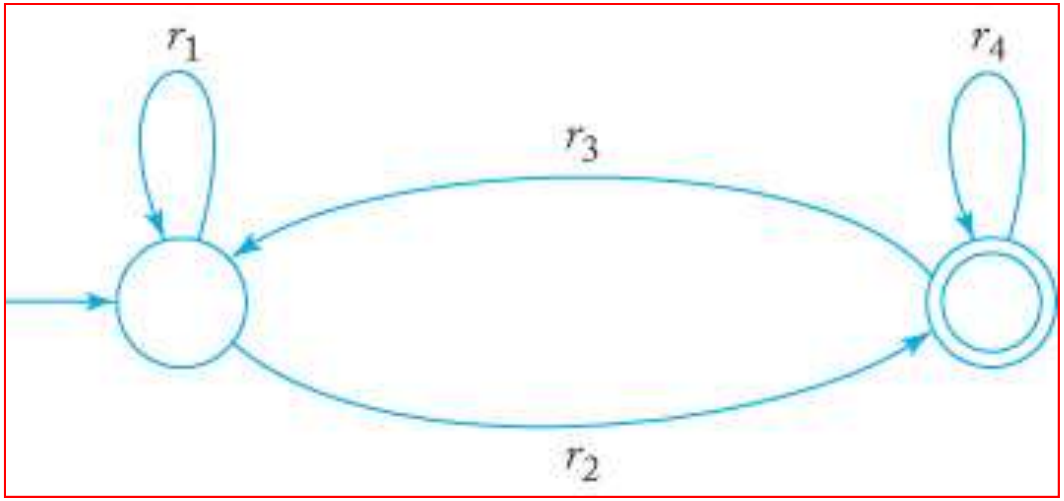


(a) (b)

補足自己走到所有人(含自己)的路徑

- If a GTG, after conversion from an NFA, has some edges missing, we put them in and label them with ϕ
- A complete GTG with |V| vertices has exactly $|V|^2$ edges

Example 3.9



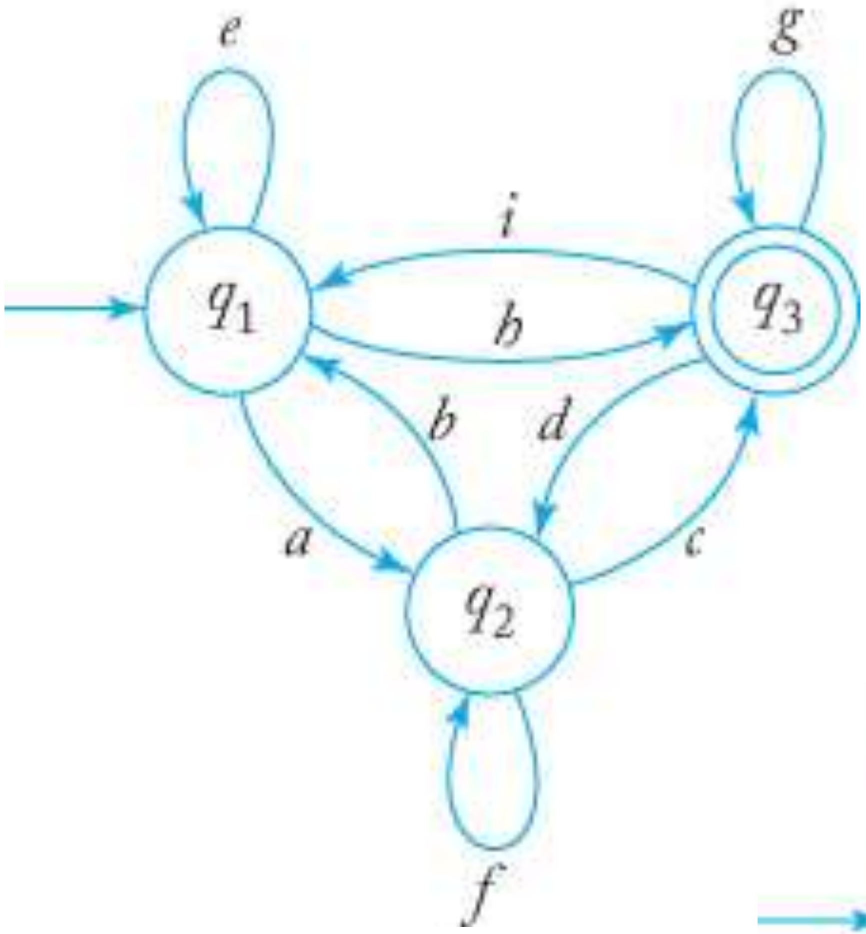
RE?

$r = r_1^*r_2(r_4 + r_3r_1^*r_2)^*$

How about a GTG with more than two states?

We can find an equivalent graph by removing one state at a time

Example 3.10



To remove q_2, we create edges as follows:

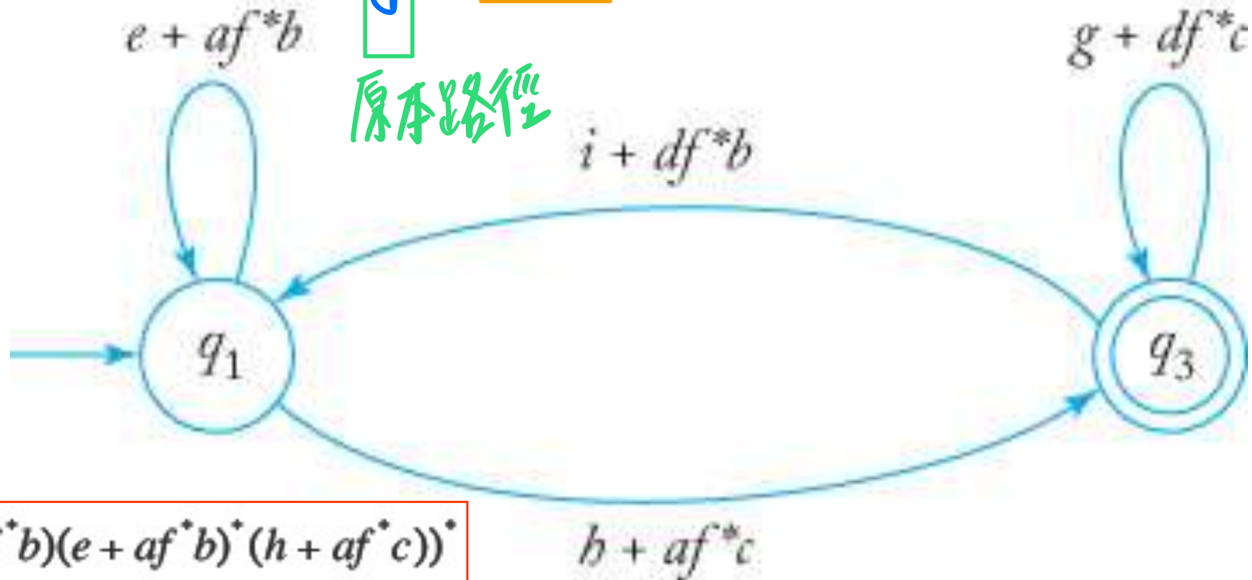
$\overrightarrow{q_1q_1} \rightarrow$ e or $+ \ af^*b$

$\overrightarrow{q_1q_3} \rightarrow$ $h + af^*c$

$\overrightarrow{q_3q_1} \rightarrow$ $i + df^*b$ 經q_2

$\overrightarrow{q_3q_3} \rightarrow$ $g + df^*c$

原本路徑



$$r = r_1^*r_2(r_4 + r_3r_1^*r_2)^*$$

$$\boxed{(e+af^*b)^*(h+af^*c)((g+df^*c)+(i+df^*b)(e+af^*b)^*(h+af^*c))^*}$$

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NFA → RE

1. Convert the NFA (with single final state) into a complete GTG. Let r_{ij} stand for the label of the edge from q_i to q_j.
2. If the GTG has only two states with $q_i \in q_0$ and $q_j \in F$, as its associated RE is:

$$r = r_{ii}^* r_{ij} (r_{jj} + r_{ji} r_{ii}^* r_{ij})^*$$

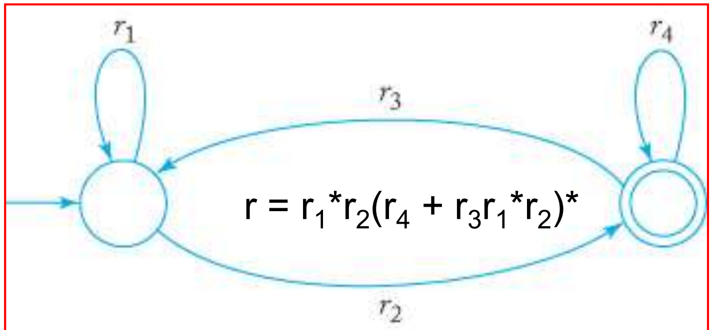
↳ 只剩2 state時使用

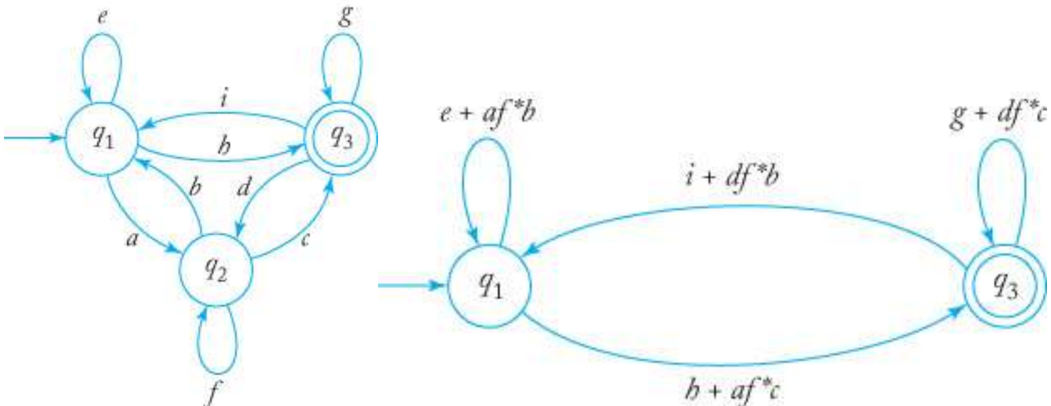
3. If the GTG has three states with $q_i \in q_0$, $q_j \in F$, and $q_k \in Q$, introduce new edges, labeled:

$$r_{pq} + r_{pk} r_{kk}^* r_{kq}$$

↳ 中間state消除

for p = i, j, q = i, j. When this is done, remove vertex q_k and its associated edges.





NFA → RE

4. If the GTG has four or more states, pick a state q_k to be removed. Apply rule 3 for all pairs of states (q_i, q_j), $i \neq k$, $j \neq k$. At each step apply the simplifying rules

$$r \overset{\text{or}}{+} \phi = r, \quad r \overset{\text{concate}}{\cdot} \phi = \phi, \quad \phi^* = \lambda$$

wherever possible. When this is done, remove state q_k.

5. Repeat step 2 to 4 until the correct RE is obtained.

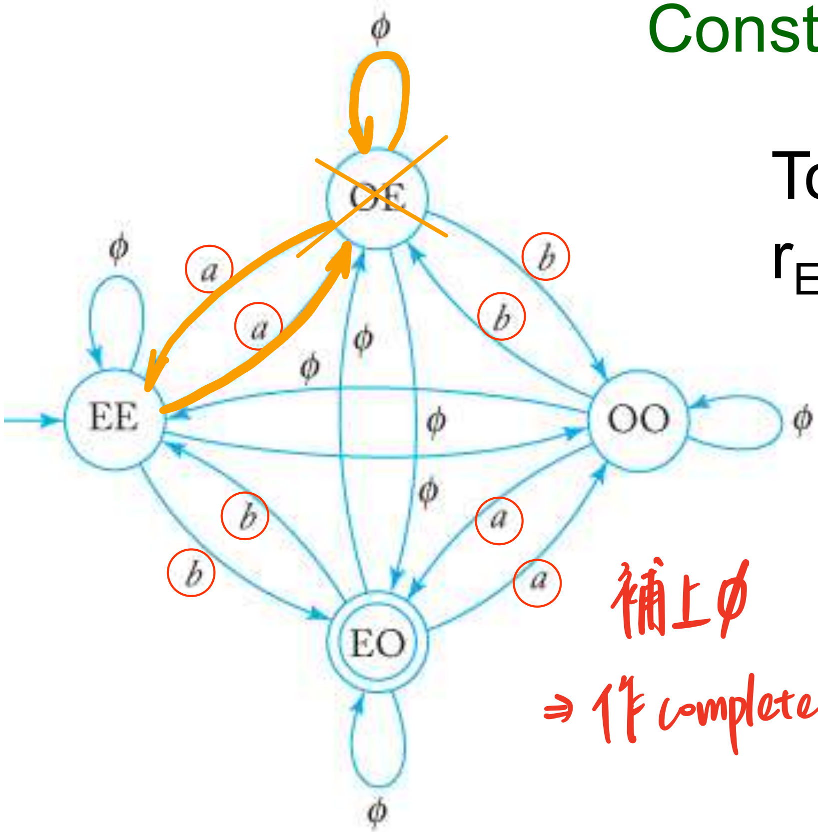
Example 3.11

- Find a RE for the language

a偶 b奇

$$L=\{w\in\{a,b\}^{*}: n_a(w) \text{ is even and } n_b(w) \text{ is odd}\}.$$

Construct NFA first!



補上ϕ

⇒ 作 complete GTG

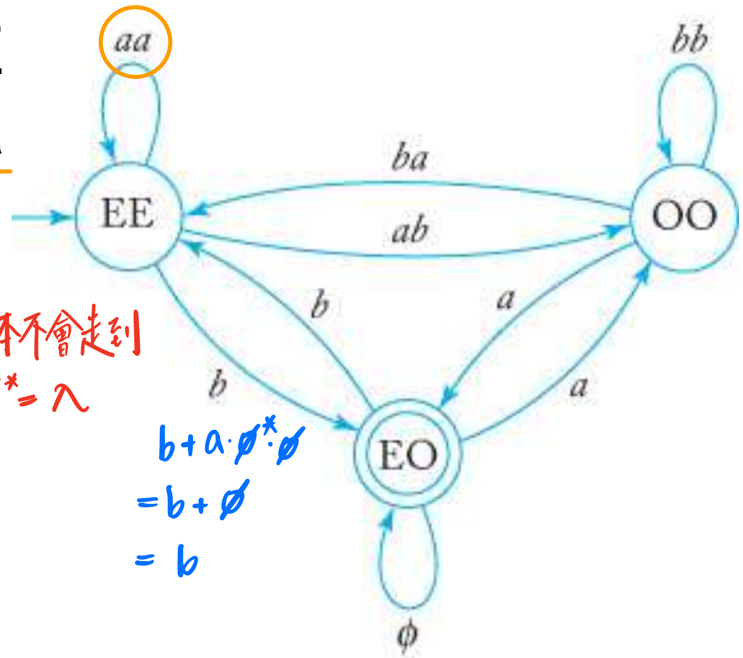
To remove OE

$r_{EE} = \phi + a\,\phi^{*}a$

$= aa$

∵根本不會走到

∴ $\phi^{*} = \lambda$



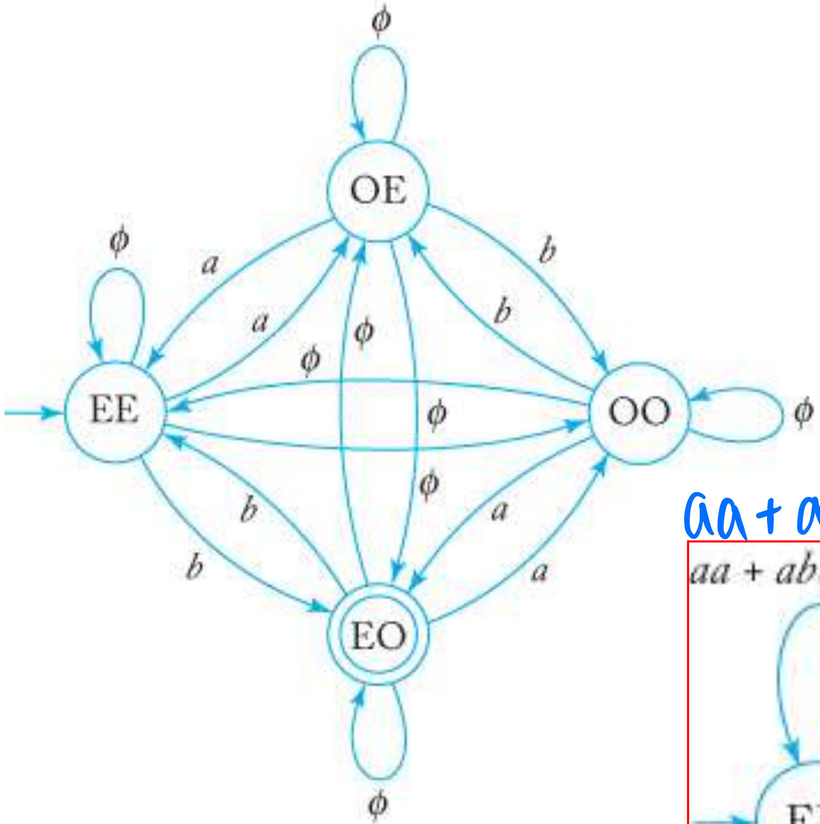
$b + a\cdot\phi^{*}\cdot\phi$

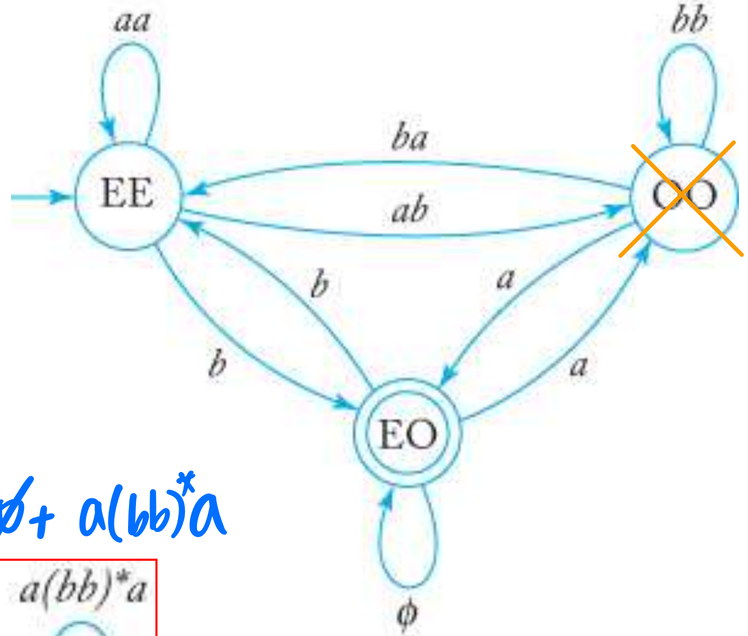
$= b + \phi$

$= b$

Example 3.11

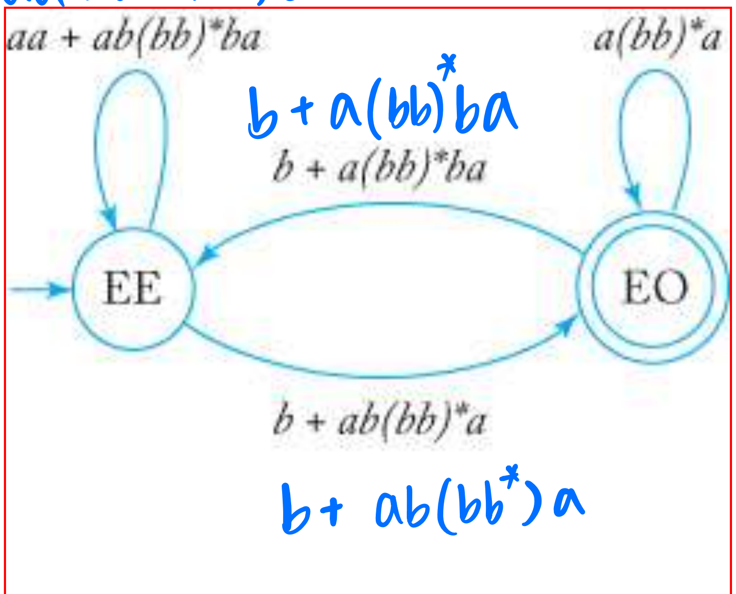
$$L = \{w \in \{a,b\}^* : n_a(w) \text{ is even and } n_b(w) \text{ is odd}\}.$$





aa + ab(bb)*ba

∅ + a(bb)*a



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Grammar Recap

- A grammar G is defined as a 4-tuple:

$$G = (V, T, S, P)$$

where

- V is a finite set of variables
- T is a finite set of terminals
- $S \in V$, called start variable
- P is a finite set of production rules

Grammar Recap

- Let $G = (V, T, S, P)$ be a grammar. Then the set

$$L(G) = \{w \in T^*: S \stackrel{*}{\Rightarrow} w\}$$

is the language generated by G

- If $w \in L(G)$, then the sequence

$$S \Rightarrow w_1 \Rightarrow w_2 \Rightarrow \ldots \Rightarrow w_n \Rightarrow w$$

is a derivation of the sentence w.

- $S, w_1, w_2, \ldots, w_n$ are called sentential forms

  ↳至少含有一個 variable

Linear Grammars

Grammars with

at most one variable at the right side

of a production

variable : 大寫

terminal : 小寫

Examples:

$$S \to aSb$$

$$S \to \lambda$$

$$S \to Ab$$

$$A \to aAb$$

$$A \to \lambda$$

Another Linear Grammar

↳右邊只有一個 variable

Grammar G :

$$S \to A$$

or

$$A \to aB \mid \lambda$$

$$B \to Ab$$

A
aB
aAb
aaBb
aaAbb
aaλbb
aabb

$$L(G) = \{a^n b^n : n \geq 0\}$$

↳非 regular：畫不出 NFA、DFA

A Non-Linear Grammar

↳ 右邊有2個variable

Grammar G:

$$S \to SS$$

$$S \to \lambda$$

$$S \to aSb$$

$$S \to bSa$$

$$L(G) = \{w: \quad n_a(w) = n_b(w)\}$$

Number of a in string w

Right-Linear Grammars

↳ variable 只能在最右邊 or 無 variable

- All productions have form: $A \to xB$

  or

  $A \to x$

  string of terminals

- Example: $S \to abS$

  $S \to a$

Left-Linear Grammars

↳ variable 只能在最左邊 or 無 variable

- All productions have form: $A \to Bx$

  or

  $A \to x$

  string of terminals

- Example: $S \to Aab$

  $A \to Aab \mid B$

  $B \to a$

Regular Grammars

A regular grammar is either right-linear or left-linear grammar

Examples:

G_1

$$S \rightarrow abS$$

$$S \rightarrow a$$

G_2

$$S \rightarrow Aab$$

$$A \rightarrow Aab \mid B$$

$$B \rightarrow a$$

G_3

$$S \rightarrow A$$

$$A \rightarrow aB \mid \lambda$$

$$B \rightarrow Ab$$

Observation

Regular grammars generate regular languages

A regular grammar is always linear, but not all linear grammars are regular.

G_3 is linear grammar but not regular grammar

$$G_3$$

$$S \to A$$

$$A \to aB \mid \lambda$$

$$B \to Ab$$

Example 3.13

Regular grammars generate regular languages

G_1

$$S \rightarrow abS$$

$$S \rightarrow a$$

G_2

$$S \rightarrow Aab$$

$$A \rightarrow Aab \mid B$$

$$B \rightarrow a$$

$$L(G_1) = (ab)^* a$$

$$L(G_2) = aab(ab)^*$$

Theorem <proof>

$$\left\{ \begin{array}{l} \text{Languages} \\ \text{Generated by} \\ \text{Regular Grammars} \end{array} \right\} = \left\{ \begin{array}{l} \text{Regular} \\ \text{Languages} \end{array} \right\}$$

Theorem - Part 1

$$\left\{ \begin{array}{l} \text{Languages} \\ \text{Generated by} \\ \text{Regular Grammars} \end{array} \right\} \subseteq \left\{ \begin{array}{l} \text{Regular} \\ \text{Languages} \end{array} \right\}$$

Any regular grammar generates a regular language

- Theorem 3.3

Theorem - Part 2

$$\left\{ \begin{array}{l} \text{Languages} \\ \text{Generated by} \\ \text{Regular Grammars} \end{array} \right\} \supseteq \left\{ \begin{array}{l} \text{Regular} \\ \text{Languages} \end{array} \right\}$$

Any regular language is generated by a regular grammar

■ Theorem 3.4

Proof – Part 1

$$\left\{\begin{array}{l}\text{Languages}\\ \text{Generated by}\\ \text{Regular Grammars}\end{array}\right\} \subseteq \left\{\begin{array}{l}\text{Regular}\\ \text{Languages}\end{array}\right\}$$

The language $L(G)$ generated by any regular grammar G is regular

The case of Right-Linear Grammars

Let G be a right-linear grammar

We will prove: $L(G)$ is regular ⇒ 找NFA、DFA

Proof idea: We will construct NFA M with $L(M) = L(G)$

* Grammar G is right-linear

Example:

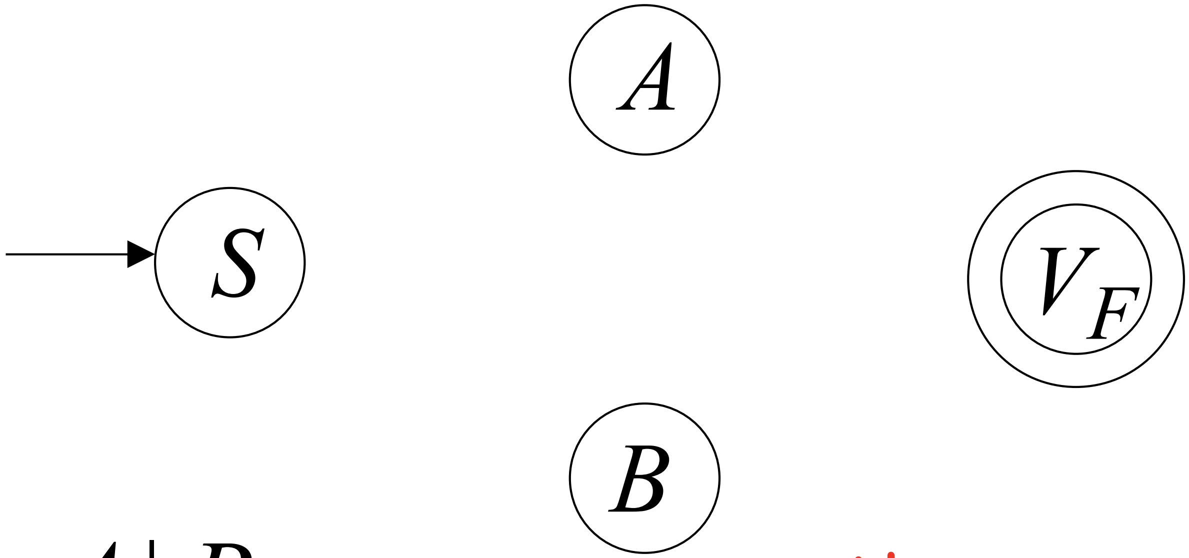
$$S \to aA \mid B$$

$$A \to aa\, B$$

$$B \to b\, B \mid a$$

right-linear

Construct NFA M such that every state is a grammar variable:



special final state

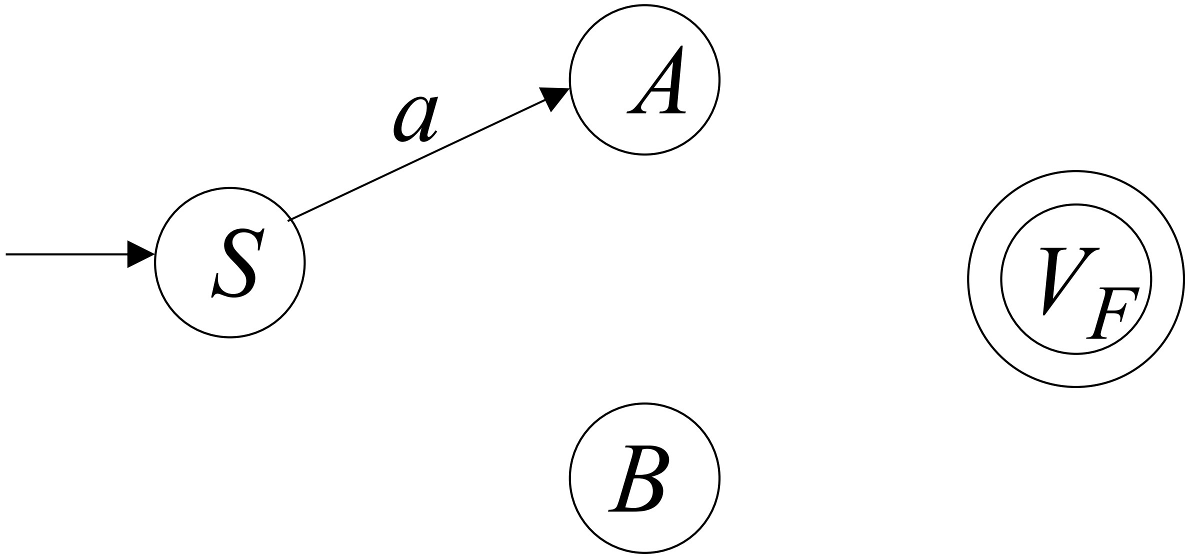
↳ variable

$$S \to aA \mid B$$

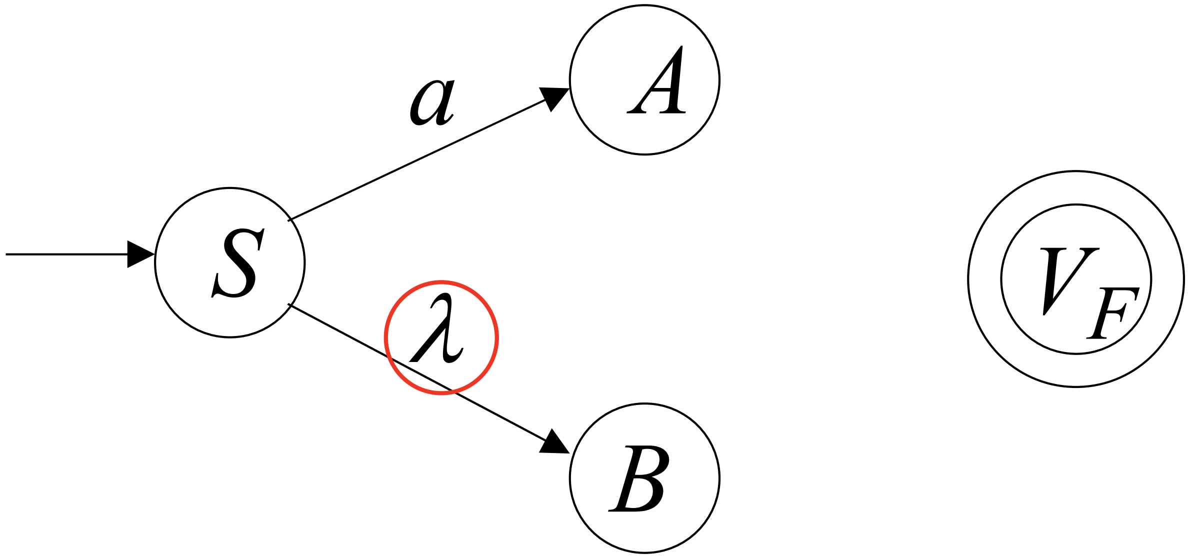
$$A \to aa\,B$$

$$B \to b\,B \mid a$$

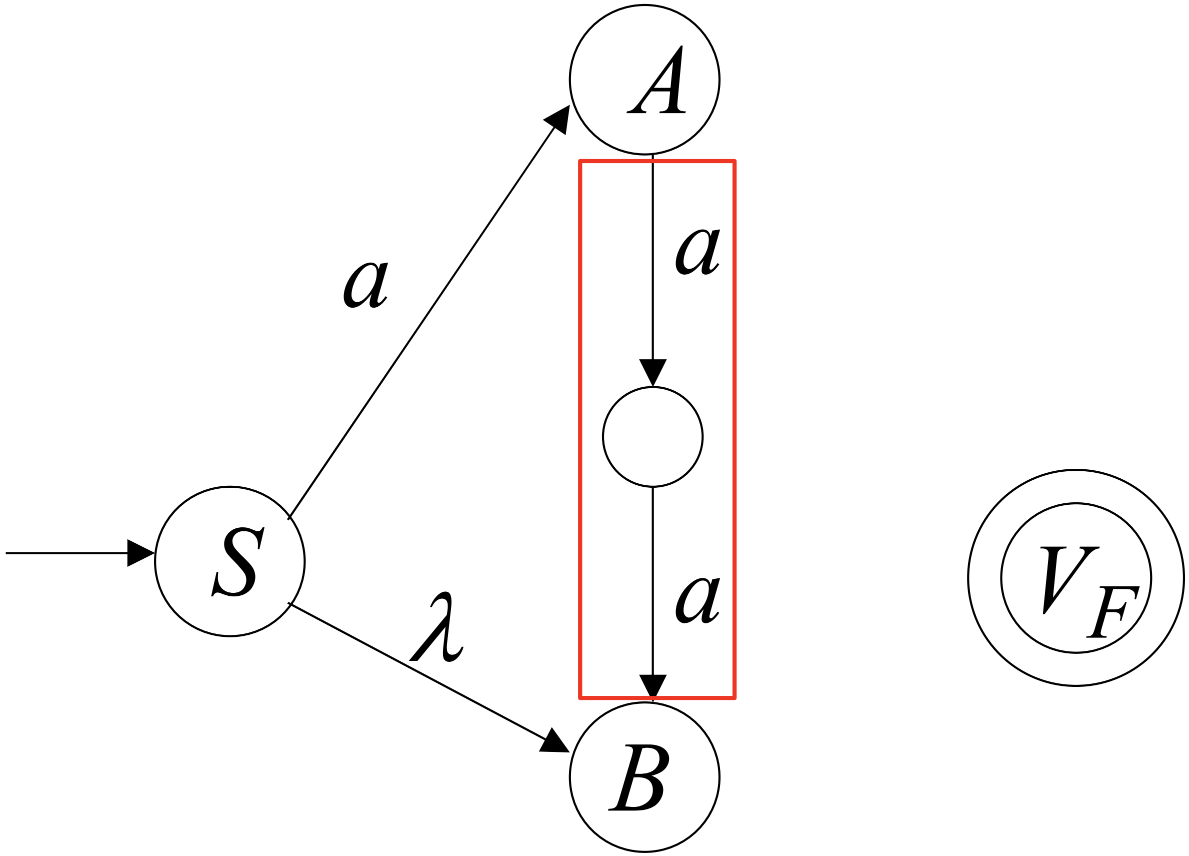
- Add edges for each production:



$$S \to aA$$

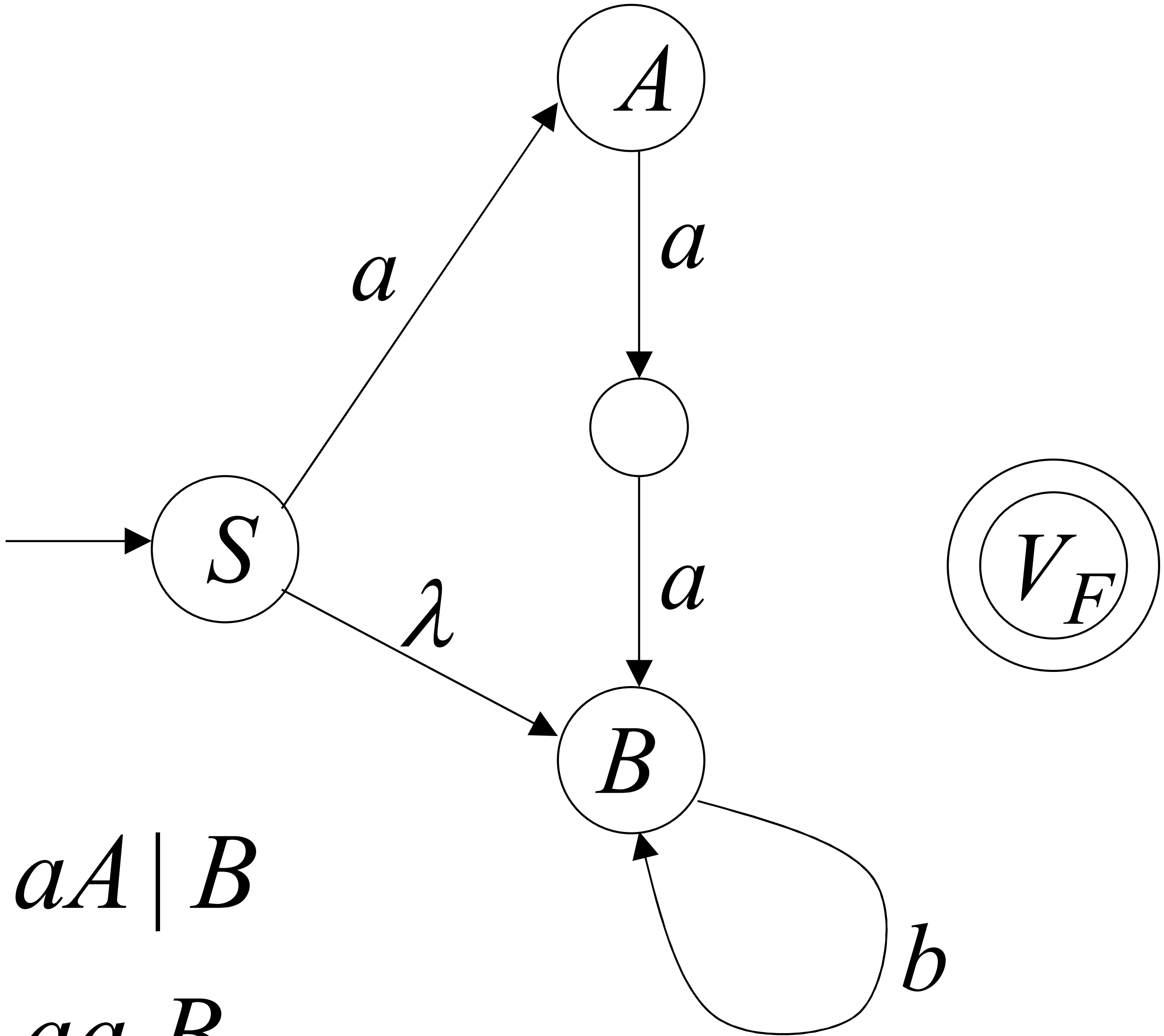


$$S \rightarrow aA \mid B$$



$$S \rightarrow aA \mid B$$

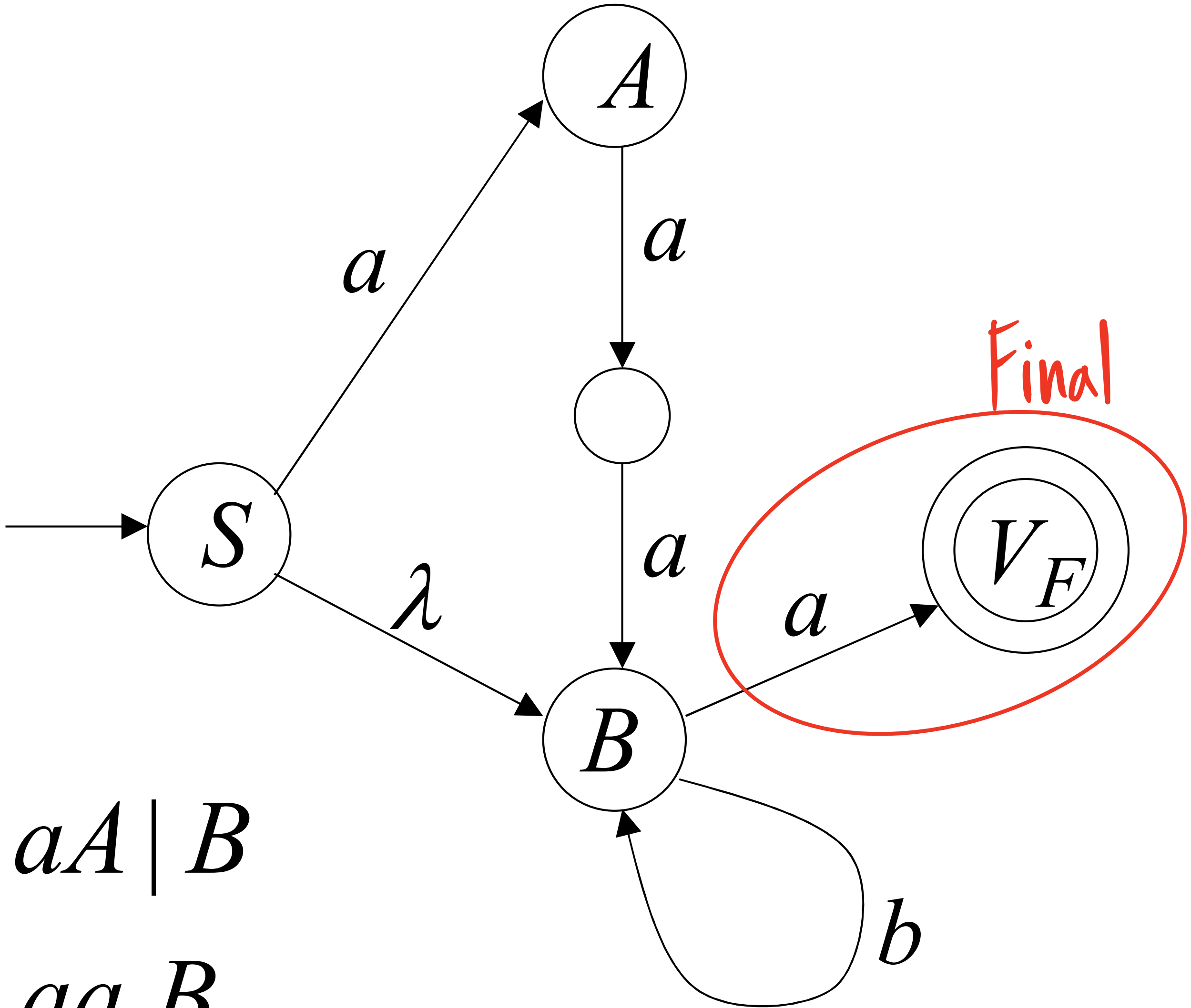
$$A \rightarrow \boxed{aa}\, B$$



$$S \to aA \mid B$$

$$A \to aa\ B$$

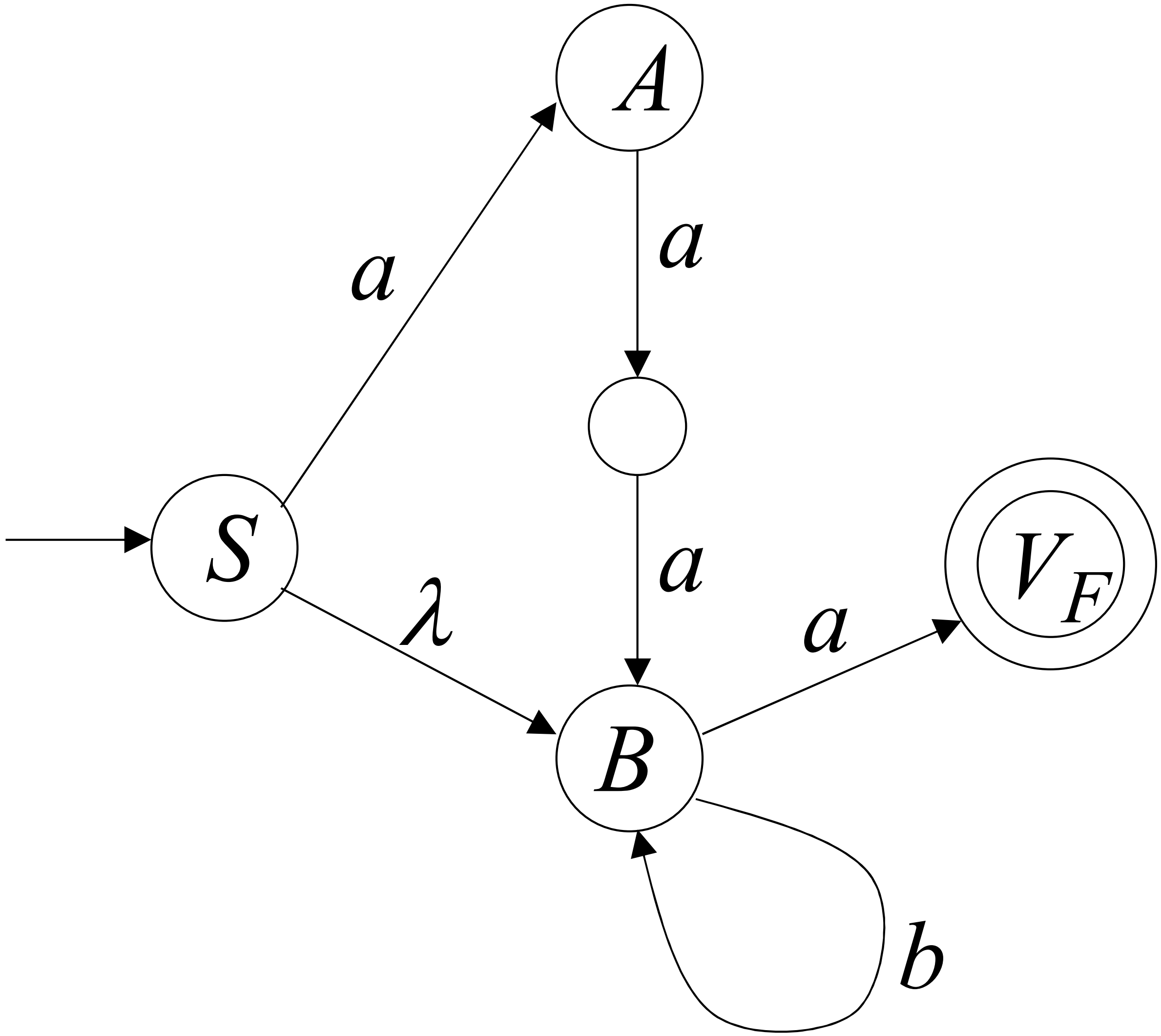
$$B \to bB$$



$$S \rightarrow aA \mid B$$

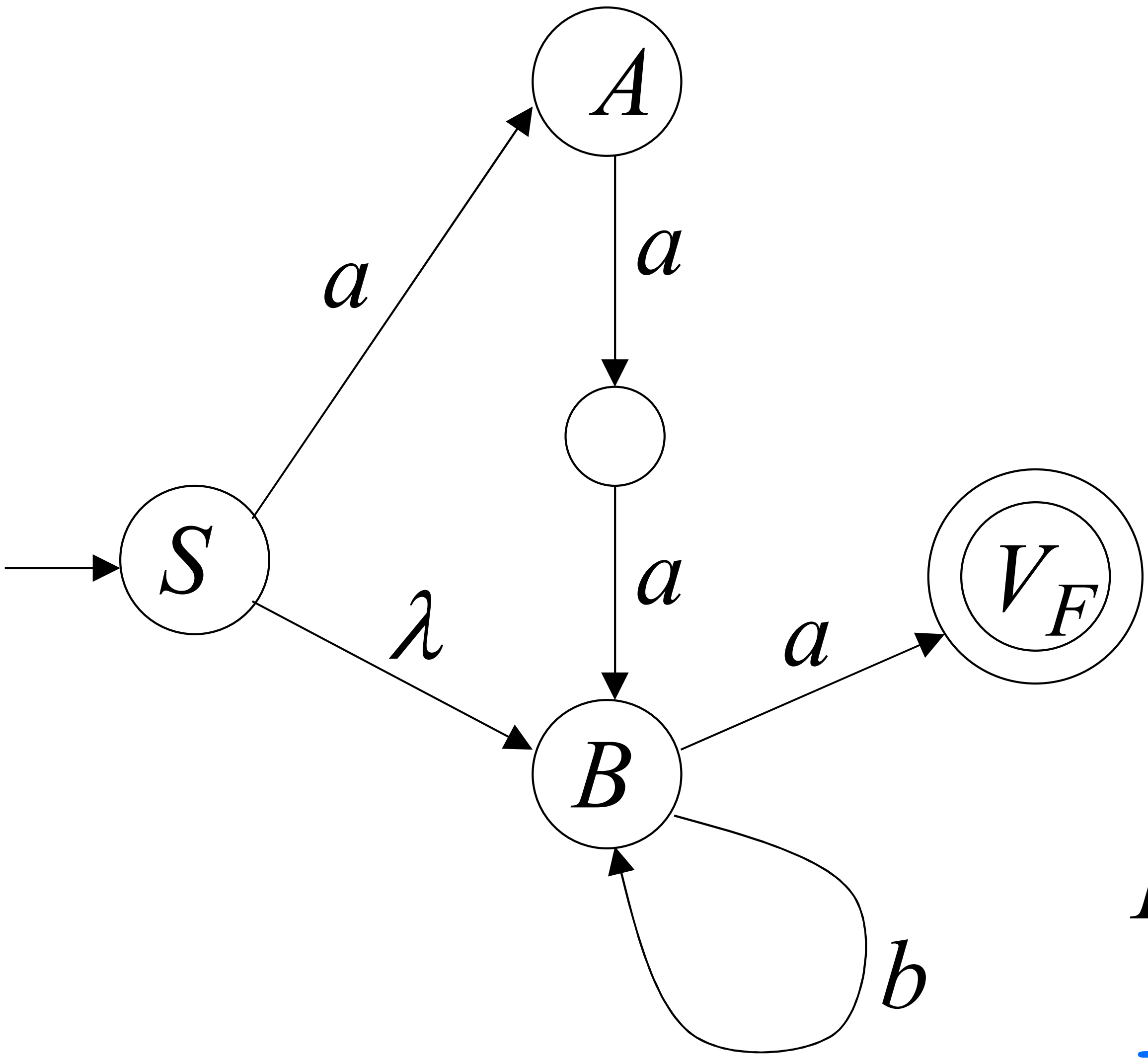
$$A \rightarrow aa\, B$$

$$B \rightarrow bB \mid a$$



$$S \Rightarrow aA \Rightarrow aaaB \Rightarrow aaabB \Rightarrow aaaba$$

NFA M



Grammar

G

$$S \to aA \mid B$$

$$A \to aa\,B$$

$$B \to bB \mid a$$

$$L(M) = L(G)$$

$= aaab^*a + b^*a$

In General

A right-linear grammar G

has variables: $V_0, V_1, V_2, \ldots$

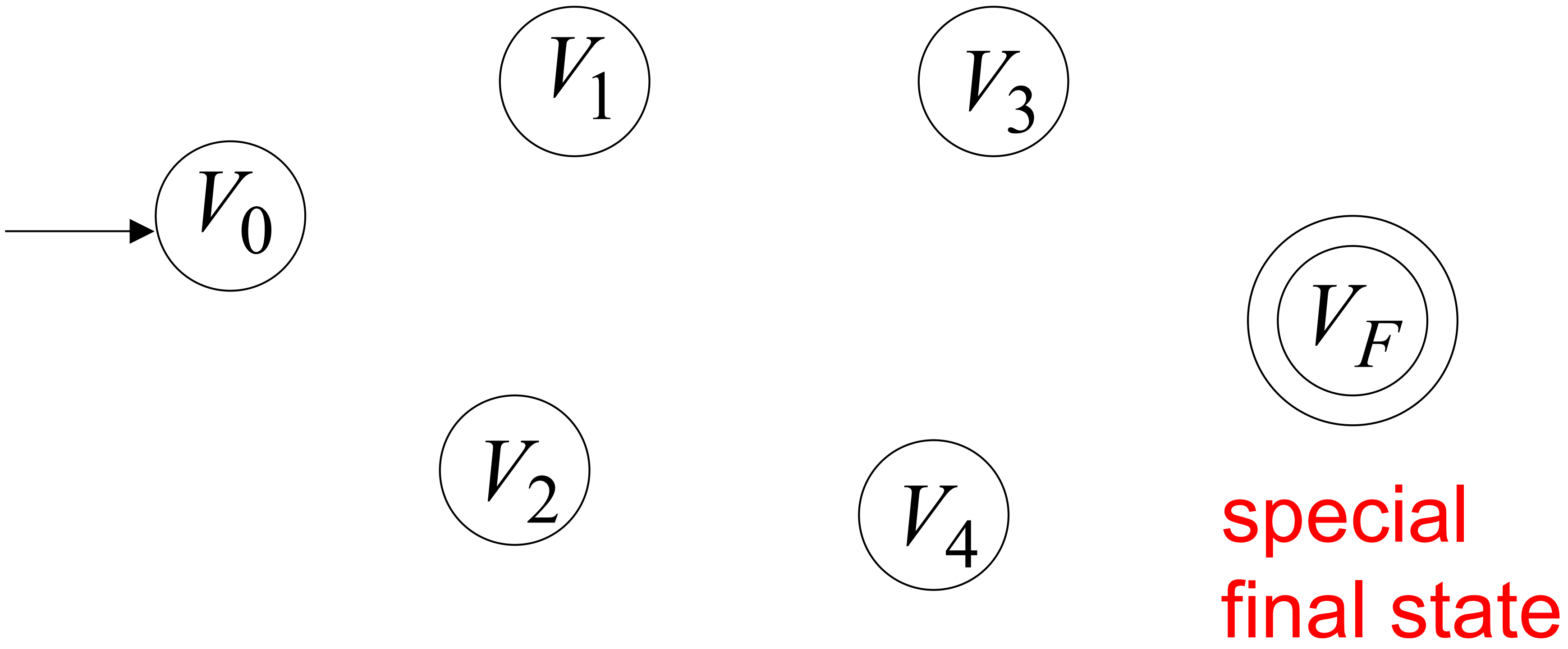
and productions: $V_i \rightarrow a_1 a_2 \cdots a_m V_j$

or

$$V_i \rightarrow a_1 a_2 \cdots a_m$$

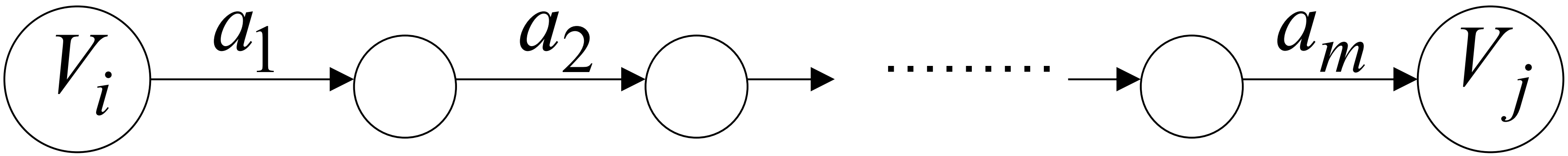
We construct the NFA M such that:

each variable V_i corresponds to a node:



For each production: $V_i \to \overbrace{a_1 a_2 \cdots a_m}^{\text{多個terminal}} V_j$

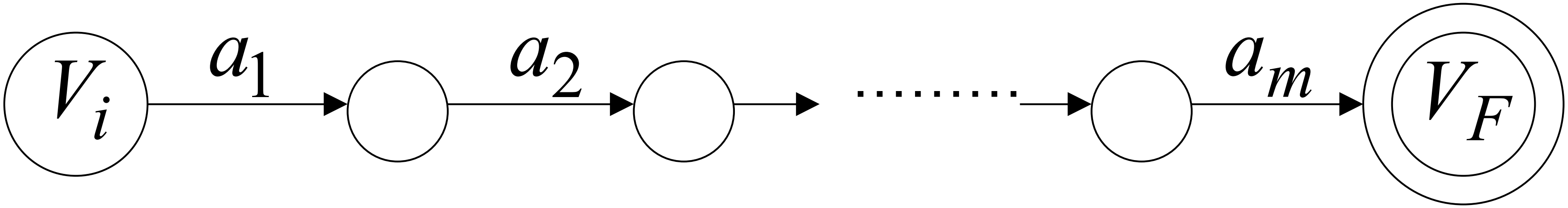
we add transitions and intermediate nodes



不具名的state

For each production: $V_i \to a_1 a_2 \cdots a_m$

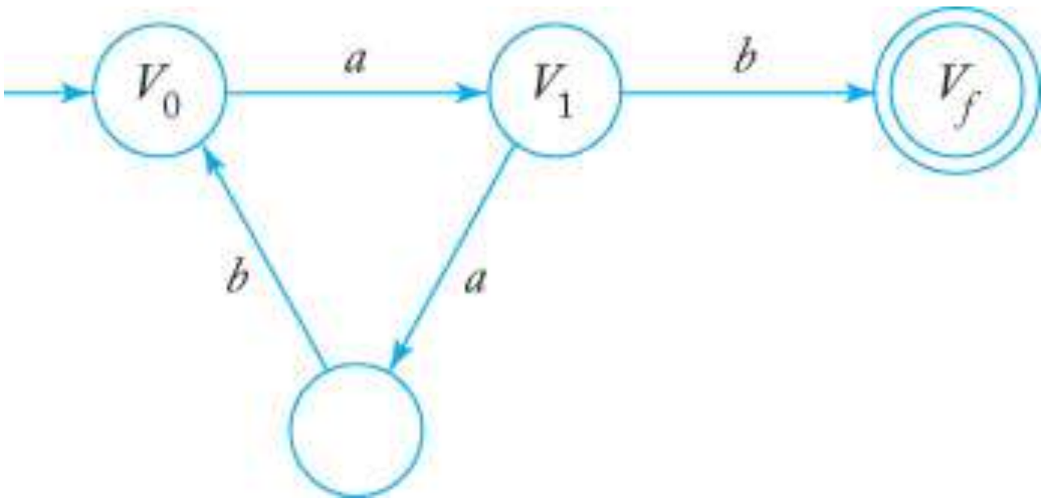
we add transitions and intermediate nodes



Example 3.15

- Construct a FA that accepts the language generated by the grammar

$$V_0 \rightarrow aV_1,$$
$$V_1 \rightarrow abV_0 | b$$



L(G)= $((aab)^* ab)$

The case of Left-Linear Grammars

Let G be a left-linear grammar

We will prove: $L(G)$ is regular

Proof idea:

We will construct a right-linear grammar G' with $L(G) = L(G')^R$

Since G is left-linear grammar the productions look like:

$$A \rightarrow Ba_1a_2\cdots a_k$$

$$A \rightarrow a_1a_2\cdots a_k$$

- Construct right-linear grammar G'

Left linear G

$$A \to B a_1 a_2 \cdots a_k$$

$$A \to Bv$$

$$\Big\Downarrow \text{Reverse}$$

Right linear G'

$$A \to a_k \cdots a_2 a_1 B$$

$$A \to v^R B$$

- Construct right-linear grammar G'

Left linear G

$$A \rightarrow a_1 a_2 \cdots a_k$$

$$A \rightarrow v$$

Right linear G'

$$A \rightarrow a_k \cdots a_2 a_1$$

$$A \rightarrow v^R$$

It is easy to see that: $L(G) = L(G')^R$

Since G' is right-linear, we have:

$L(G') \Rightarrow L(G')^R \Rightarrow L(G)$

Regular Language | Regular Language | Regular Language

Proof - Part 2

$$\left\{\begin{array}{l}\text{Languages} \\ \text{Generated by} \\ \text{Regular Grammars}\end{array}\right\} \supseteq \left\{\begin{array}{l}\text{Regular} \\ \text{Languages}\end{array}\right\}$$

Any regular language L is generated by some regular grammar G

Any regular language L is generated by some regular grammar G

Proof idea:

Let M be the NFA with $L = L(M)$.

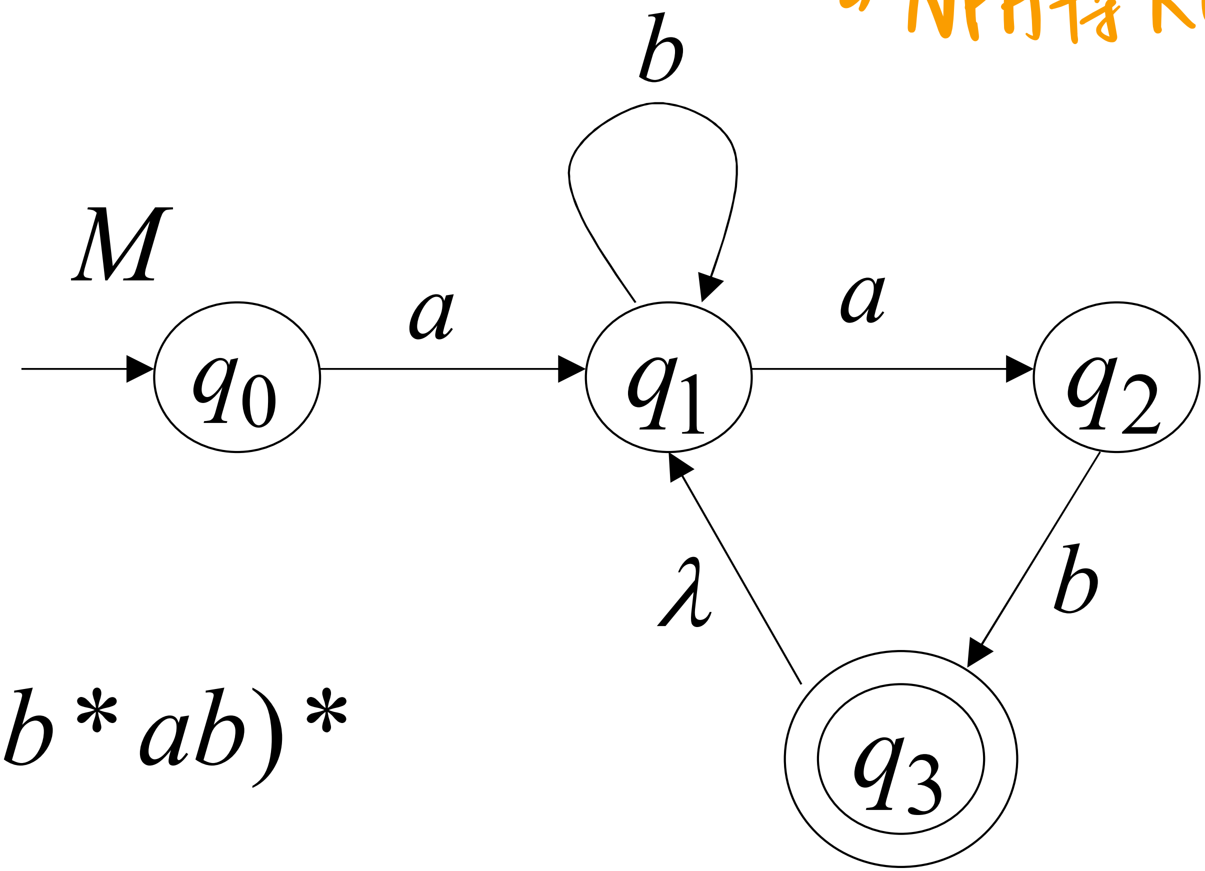
Construct from M to a regular grammar G such that $L(M) = L(G)$

Since L is regular

there is an NFA M such that $L = L(M)$

↳ NFA轉 RG

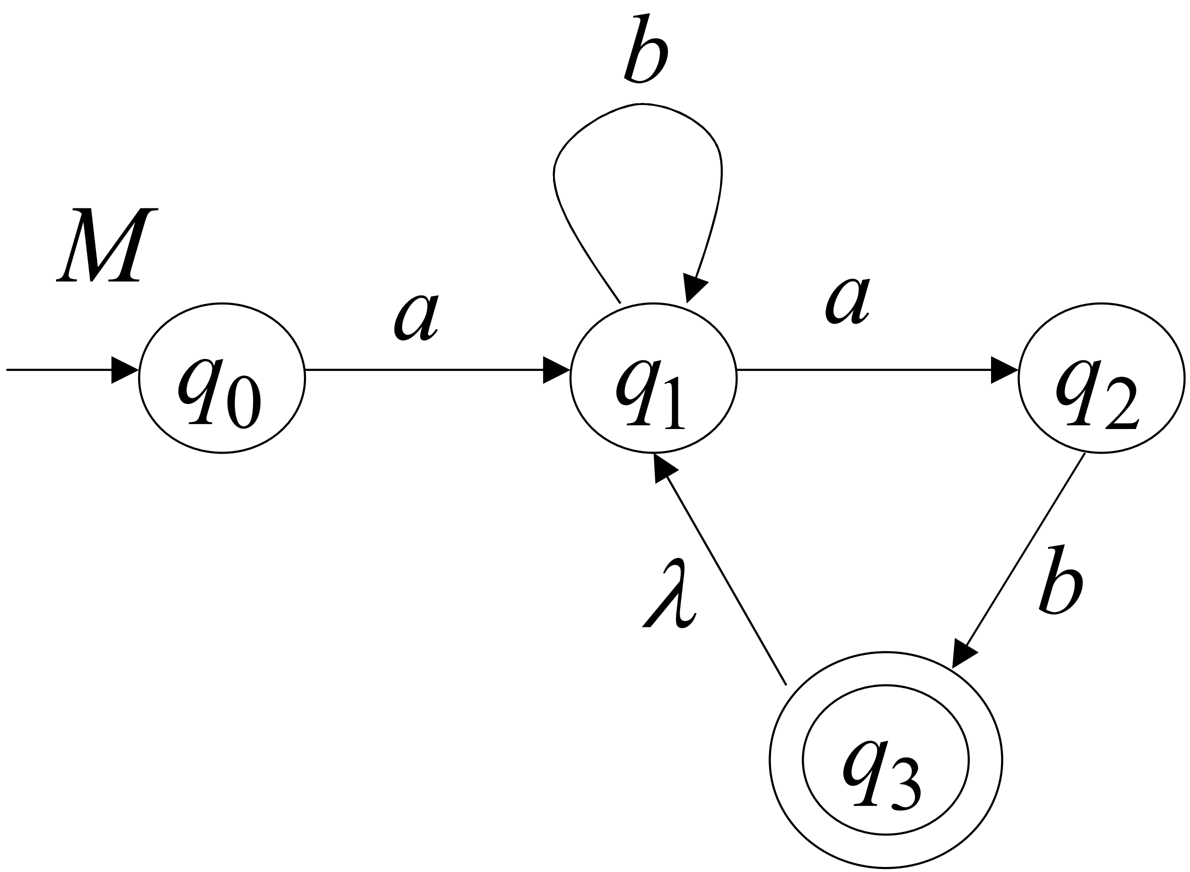
Example:



$L = ab^*ab(b^*ab)^*$

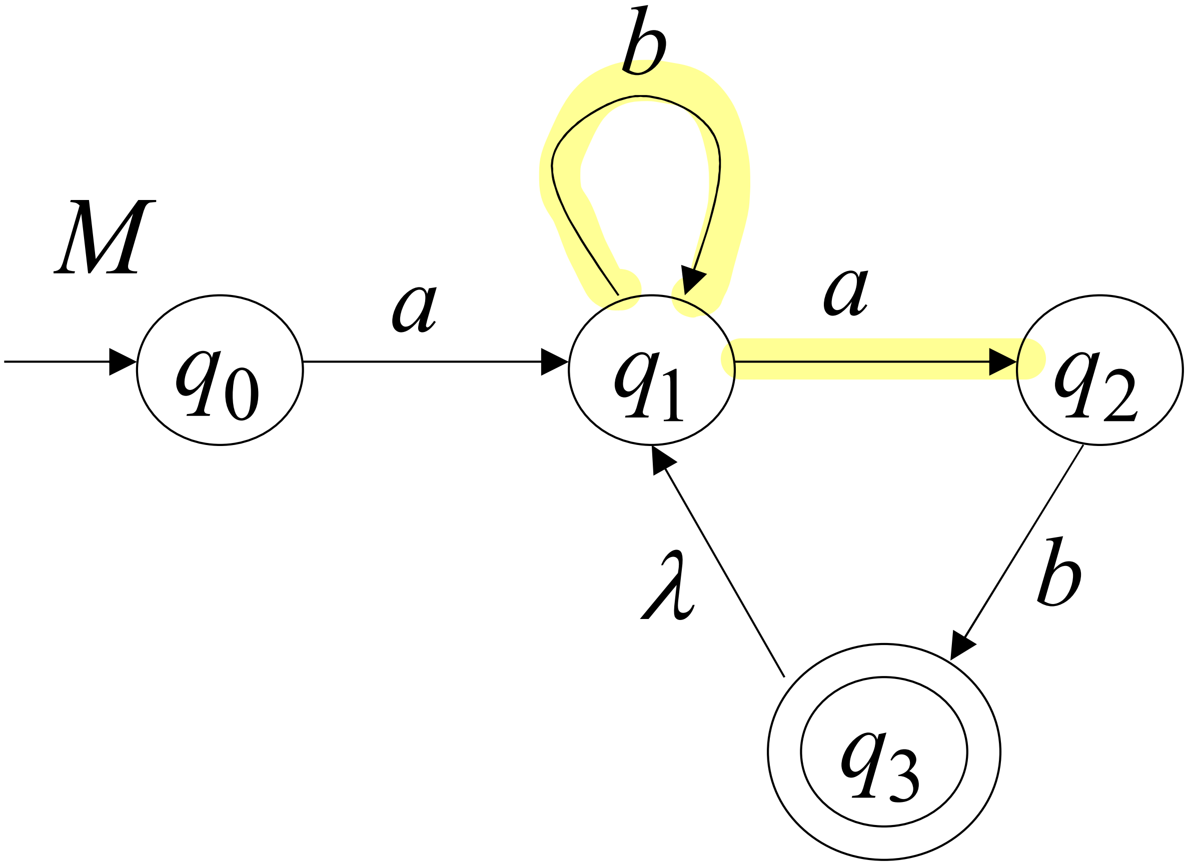
$L = L(M)$

Convert M to a right-linear grammar (BFS-like)



$$q_0 \rightarrow aq_1$$

M



$$q_0 \rightarrow aq_1$$

$$q_1 \rightarrow bq_1$$

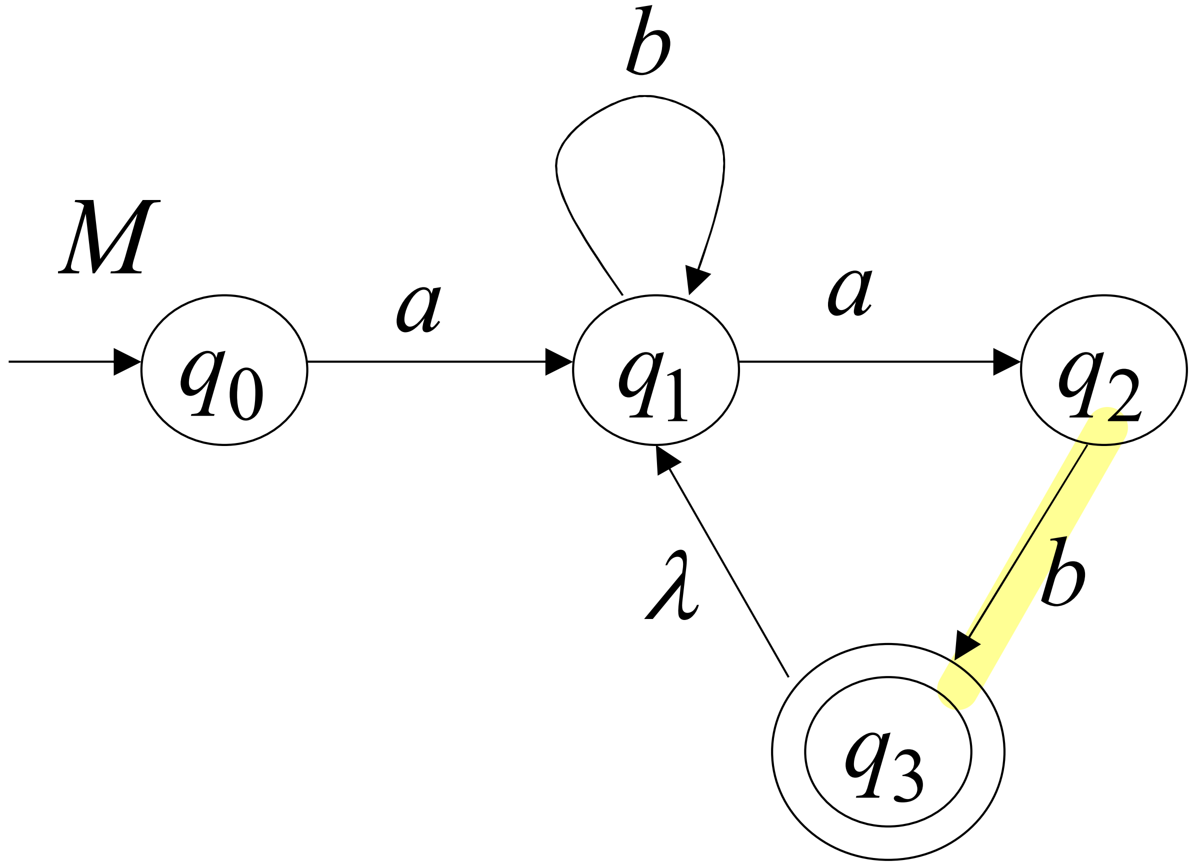
$$q_1 \rightarrow aq_2$$

$$q_0 \rightarrow aq_1$$

$$q_1 \rightarrow bq_1$$

$$q_1 \rightarrow aq_2$$

$$\boxed{q_2 \rightarrow bq_3}$$



$$L(G) = L(M) = L$$

G

$q_0 \to aq_1$

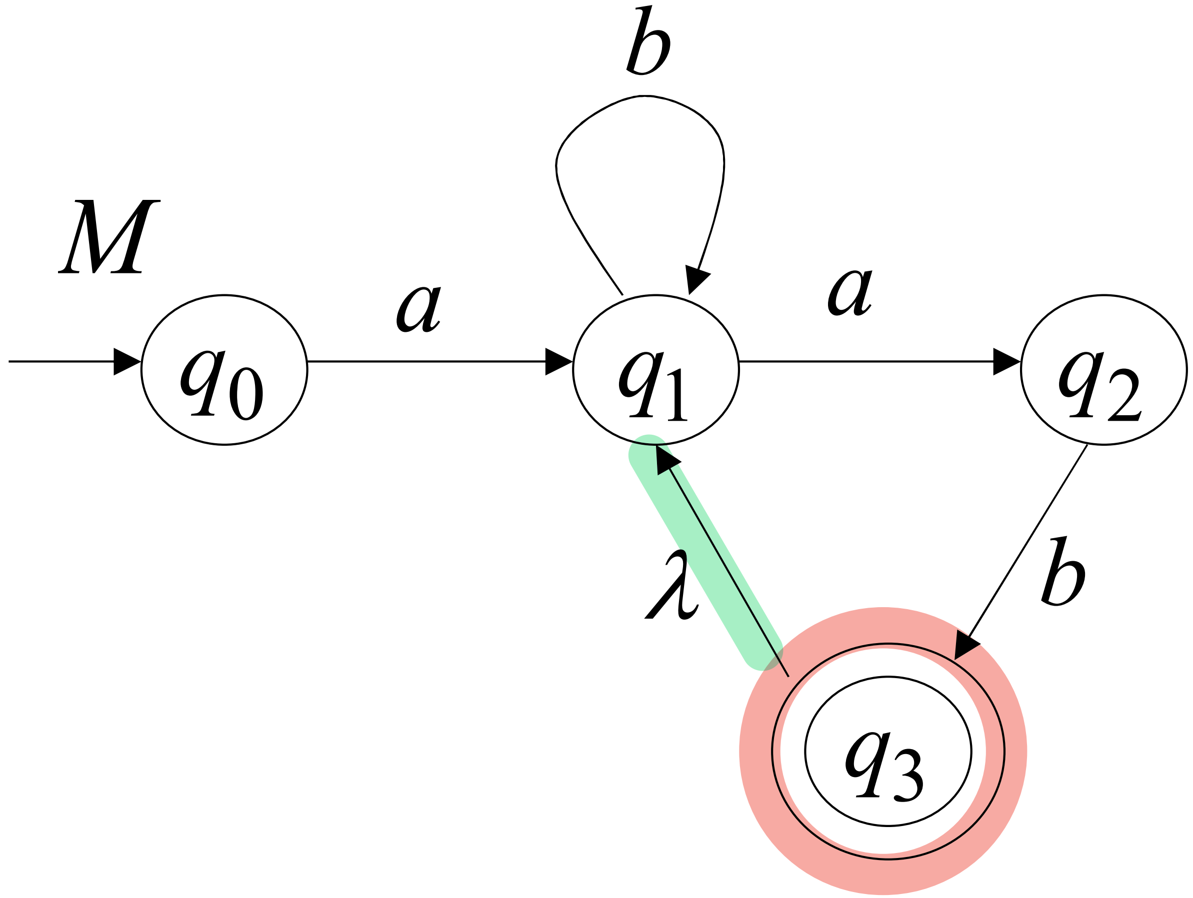
$q_1 \to bq_1$

$q_1 \to aq_2$

$q_2 \to bq_3$

$q_3 \to q_1$

$q_3 \to \lambda$ ↳ final state 需通向 λ



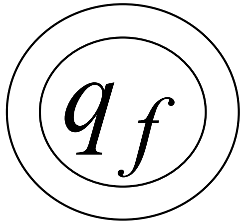
In General

For any transition: $q \xrightarrow{a} p$

Add production: $q \rightarrow ap$

variable terminal variable

For any final state:



Add production:

$$q_f \to \lambda$$

Since G is right-linear grammar

G is also a regular grammar

with $L(G) = L(M) = L$

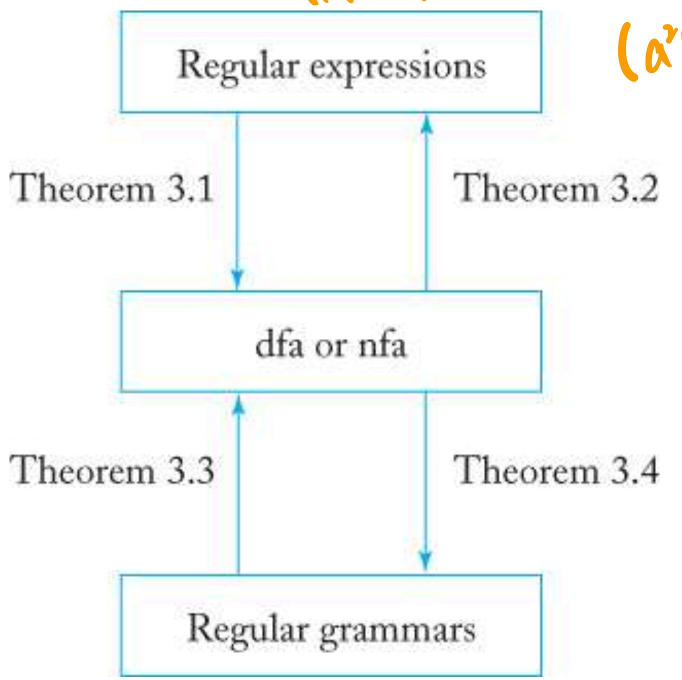
Summary

- We now have several ways of describing regular languages:

↳無記憶性，∴ $a^n b^n$ 找不到
($a^2 b^2$、$a^3 b^3$ … 可以)

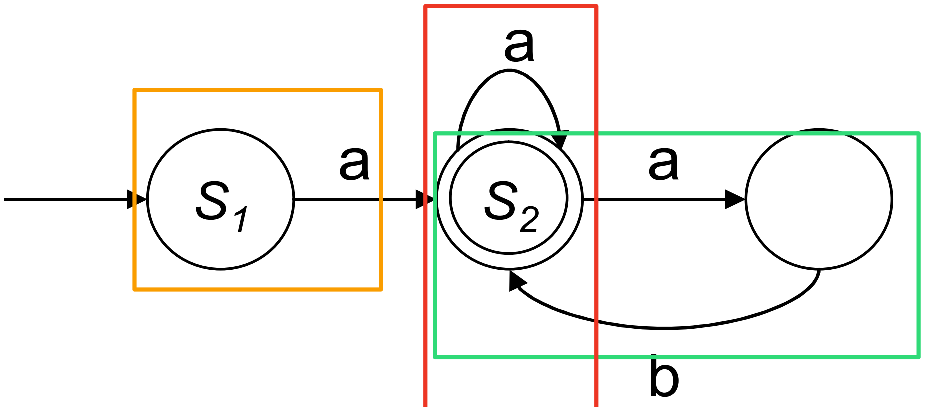
model
工具

  - DFA
  - NFA
  - RE
  - RG



Short Quiz

- Find a regular grammar that generates the language L(aa*(ab+a)*).



$$S_1 \to aS_2$$
$$S_2 \to aS_2 \mid abS_2 \mid \lambda$$

Questions?