

Homework One
Theory of Computation 2023

Important Note:

Please remember that you should return your answer at 10/11 (Wednesday) 15:10 and your HW should be handwritten. We will take your HW during the class. After 10/11 15:10, you must upload your HW to moodle. But remember penalty for late submission: 20% per day.

Q1: For $\Sigma = \{a, b\}$, construct dfa's that accept the sets consisting of

- (a) all strings of even length.
- (b) all strings of length greater than 5.
- (c) all strings with an even number of a's and an odd number of b's.
- (d) all the strings with exactly two a's and more than three b's.
- (e) all strings with no more than two a's.

Q2: Find dfa's for the following languages on $\Sigma = \{a, b\}$.

- (a) $L = \{w : |w| \bmod 5 = 0\}$.
- (b) $L = \{w : (n_a(w) - n_b(w)) \bmod 3 = 0\}$.

Q3: Let L be the language accepted by automaton in Figure 1. Find a dfa that accepts L^2 .



Figure 1: Automaton diagram.

Q4: Find a dfa that accepts the language defined by the nfa in Figure 2.

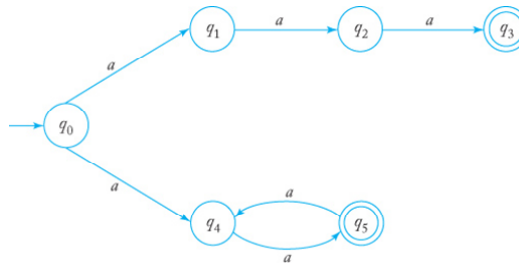


Figure 2: NFA diagram.

Q5: Design an nfa with no more than five states for the set $\{abab^n : n \geq 0\} \cup \{aba^n : n \geq 0\}$

Q6: Find an nfa with four states for $L = \{a^n : n \geq 0\} \cup \{b^n a : n \geq 1\}$.

Q7: Convert the nfa defined by

$$\delta(q_0, a) = \{q_0, q_1\}$$

$$\delta(q_1, b) = \{q_1, q_2\}$$

$$\delta(q_2, a) = \{q_2\}$$

$$\delta(q_0, \lambda) = \{q_2\}$$

with initial state q_0 and final state q_2 into an equivalent dfa.

Q8: Find a minimal dfa for the following language, and prove that the result is minimal.

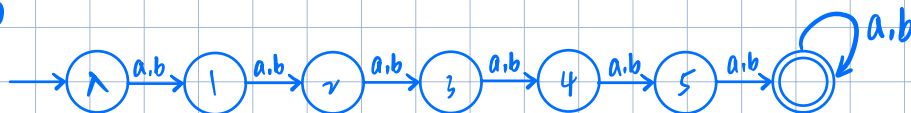
$$L = \{a^n : n \geq 0, n \neq 2\}.$$

- Q1: For $\Sigma = \{a, b\}$, construct dfa's that accept the sets consisting of
- all strings of even length.
 - all strings of length greater than 5.
 - all strings with an even number of a's and an odd number of b's.
 - all the strings with exactly two a's and more than three b's.
 - all strings with no more than two a's.

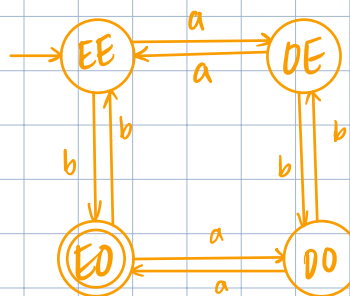
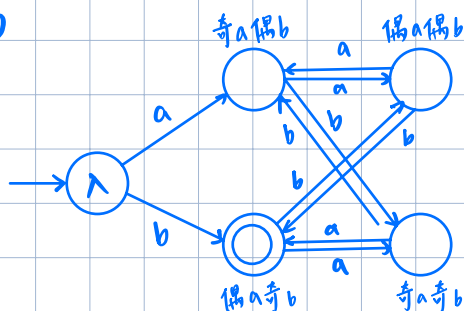
(a)



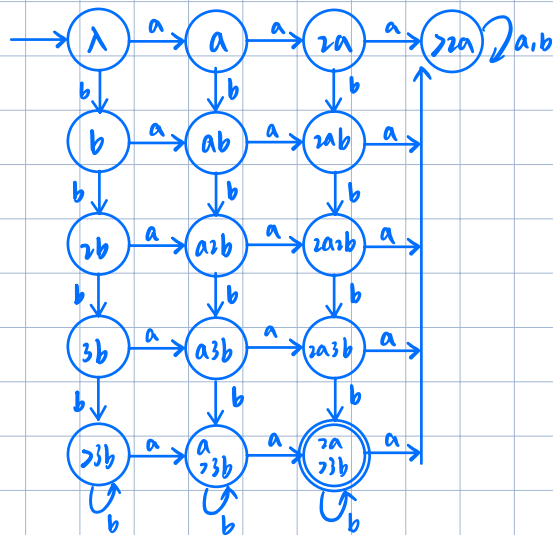
(b)



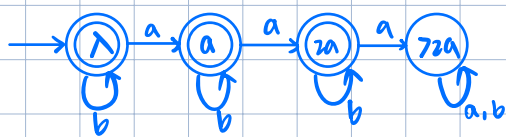
(c)



(d)



(e)

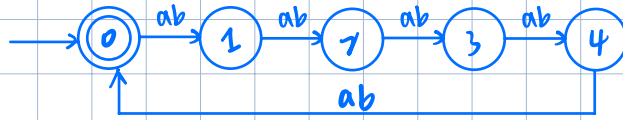


Q2: Find dfa's for the following languages on $\Sigma = \{a, b\}$.

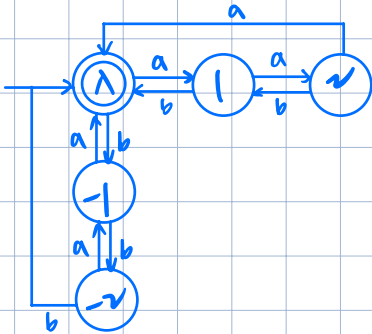
(a) $L = \{w : |w| \bmod 5 = 0\}$.

(b) $L = \{w : (n_a(w) - n_b(w)) \bmod 3 = 0\}$.

(a)



(b) i.e. $(n_a(w) - n_b(w)) \% 3 = 0, 1, 2$

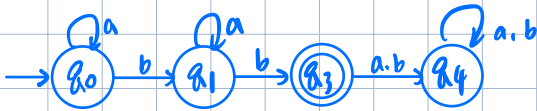


Q3: Let L be the language accepted by automaton in Figure 1. Find a dfa that accepts L^2 .



Figure 1: Automaton diagram.

$L = \{a^n b : n \geq 0\}$, $L^2 = \{a^n b a^m b : n, m \geq 0\}$



Q4: Find a dfa that accepts the language defined by the nfa in Figure 2.

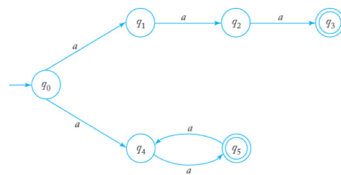
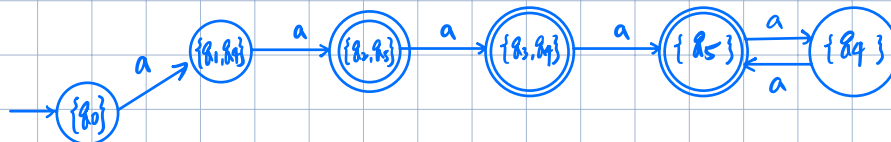


Figure 2: NFA diagram.

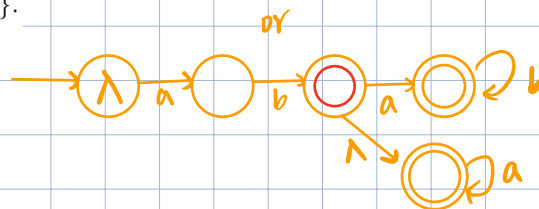
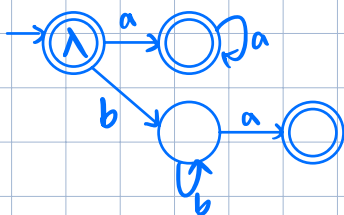
$L = a^+ \cup a^{2n}$
 $\hookrightarrow a^2, a^3, a^4, a^6, \dots$



Q5: Design an nfa with **no more than five states** for the set $\{abab^n : n \geq 0\} \cup \{aba^n : n \geq 0\}$



Q6: Find an nfa with **four states** for $L = \{a^n : n \geq 0\} \cup \{b^n a : n \geq 1\}$.



Q7: Convert the nfa defined by

$$\delta(q_0, a) = \{q_0, q_1\}$$

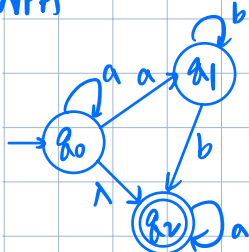
$$\delta(q_1, b) = \{q_1, q_2\}$$

$$\delta(q_2, a) = \{q_2\}$$

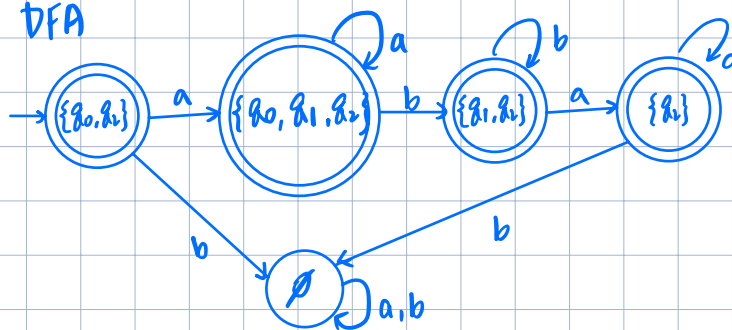
$$\delta(q_0, \lambda) = \{q_2\}$$

with initial state q_0 and final state q_2 into an equivalent dfa.

NFA

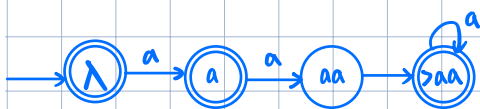


DFA



Q8: Find a minimal dfa for the following language, and prove that the result is minimal.

$$L = \{a^n : n \geq 0, n \neq 2\}$$



$$\{q_0, q_1, q_3\} \in F, \{q_2\} \notin F$$

$$\text{找 } (q_0 \text{ vs } q_2), (q_0 \text{ vs } q_1), (q_1 \text{ vs } q_2)$$

<pf>

$$q_0 \notin F, q_2 \in F \Rightarrow q_0, q_2 \text{ 為 distinguishable}$$

$$\textcircled{2} \delta(q_1, a) = q_2 \notin F, \delta(q_2, a) = q_3 \in F \Rightarrow q_1, q_2 \text{ 為 distinguishable}$$

$$\textcircled{3} \delta(q_0, aa) = q_2 \notin F, \delta(q_1, aa) = q_3 \in F \Rightarrow q_0, q_1 \text{ 為 distinguishable}$$

$$\textcircled{4} \delta(q_0, aa) = q_2 \notin F, \delta(q_2, aa) = q_3 \in F \Rightarrow q_0, q_2 \text{ 為 distinguishable}$$

由以上可知, 4 狀態互相為 distinguishable

\therefore DFA 為最小