

This problem has 2 conditions:

- 1) vowels == consonants
- 2) (vowels * consonants) % k == 0

To check the first condition for any sub-array $i-j$, we can compute it by $O(1)$ with prefix sum.

let $\text{presum}[i] = \# \text{ of vowels} - \# \text{ of consonants in } S[1 \rightarrow i]$
(S: 1-indexed style)

so: sub-array $[i, j]$ has (1) $\Leftrightarrow \text{presum}[j] == \text{presum}[i-1]$

\Rightarrow let create $\text{group}[k]$ contain all i index that have $\text{presum}[i] = k$.

\Rightarrow any pairs in each group always satisfy (1) ..

Now, we want to find how many pairs in each group satisfy (2)

Because vowels == consonants in valid sub-array.

$$\Rightarrow (2) \Leftrightarrow \text{vowel}^2 \% k = 0$$

let consider sub array $[i+1, j]$ satisfy both conditions.

$$\text{we kn: } \left(\frac{j-i}{2} \right)^2 \% k = 0 \quad \left(\begin{array}{l} \text{vowel} = \text{consonant} \\ \text{vowel} + \text{consonant} = j - i \end{array} \right)$$

$$\Leftrightarrow \frac{j^2 - 2ij + i^2}{2} \% k = 0$$

The part $i * j$ make the computation $O(n^2)$ to check all i, j .

But we need some other direction to reduce the complexity

If the condition $\left(\frac{j-i}{2}\right)^2 g_k$ be equalled by $\frac{j-i}{2} g_{k_{\text{new}}}$ it will reduce the complexity.
To be simply, let make this condition just related to i, j .

$$\begin{aligned} \left(\frac{j-i}{2}\right)^2 g_k &= 0 \\ \Leftrightarrow \frac{(j-i)^2}{4} g_k &= 0 \\ \Leftrightarrow (j-i)^2 g_{4k} &= 0. \end{aligned}$$

let $k \leftarrow 4k$

so now, make condition (2) equal new condition related to $j-i$ only.

let : $j-i = n$. (Find and for n to $n^2 g_k = 0$)

we need convert $n^2 g_k = 0$ to $n g_{k_{\text{new}}} = 0$

let: $k = \prod_{i=1}^m p_i^{q_i}$ with p_i is prime number.

so $n^2 : k \Leftrightarrow n^2 = \prod_{i=1}^m p_i^{t_i} \times \dots / t_i \geq q_i$

$$n = \prod_{i=1}^m p_i^{\frac{t_i}{2}} \times \dots \quad / \quad \frac{t_i}{2} \geq \frac{q_i}{2}$$

so $\forall n \quad / \quad \frac{t_i}{2} \geq \frac{q_i}{2} \Rightarrow n^2 : k$

$$\Leftrightarrow \forall n : \boxed{k_{\text{new}} = \prod_{i=1}^m p_i^{\frac{q_i}{2}}} \quad (*)$$

To compute k_{new} we have 2 way.
Factor k to find q_i , $i = \overline{1, m}$.

or brute for k :

$$\text{From (1)} \begin{cases} k_{\text{new}} \geq \sqrt{k} \\ k_{\text{new}}^2 \leq k \end{cases}$$

\Rightarrow Check all $k_{\text{new}} = [\sqrt{k}, k]$, find $\min k_{\text{new}}^2 \leq k$.

$$\text{So } (j-i)^2 \leq k \Leftrightarrow j-i \leq k_{\text{new}}$$

$$\Leftrightarrow j \% k_{\text{new}} = i \% k_{\text{new}}$$

So, For each group:

$$\text{any pair } i, j \mid j \% k_{\text{new}} == i \% k_{\text{new}}$$

will create beautiful string.