UNIVERSITÄT
LEIPZIG

# Asymmetric Numeral System

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## Outline

1. Introduction

2. Asymmetric Numeral System

3. rANS

4. Streaming-rANS

5. Summary

6. Further Study

7. Reference

# Introduction

# Outline

1. Introduction
   1.1 Background
   1.2 Information Theory
   1.3 Source Coding Theorem
   1.4 Entropy Coding

- 1948: Claude Shannon introduced the concept of information theory in his paper ”A Mathematical Theory of Communication” [1].
- 1950: David Huffman introduced Huffman coding [3].
- 1977: Arithmetic coding was introduced by Robert M. Fano [4].
- 2014: Asymmetric Numeral System (ANS) was introduced by Jarosław Duda [2].
- 2016: Zstandard (Zstd) was introduced by Yann Collet [5].
- 2018: AV1 Image File Format (AVIF) was introduced by Alliance for Open Media [6].
- 2020: JPEG XL was introduced by JPEG Committee [7].

## Information Theory[1]

- Claude Shannon: ”A Mathematical Theory of Communication”
- Self-information
  a symbol of probability $p$ contains $log_2 \frac{1}{p}$ bits of information.
- Shannon Entropy
  Let $X$ be a random variable, which takes values in the set $\mathcal{X}$ and is distributed according to $p$

$$H(X) = \sum_{x \in \mathcal{X}} p(x) log_2 \frac{1}{p(x)}$$

## Source Coding Theorem (Shannon) [1]

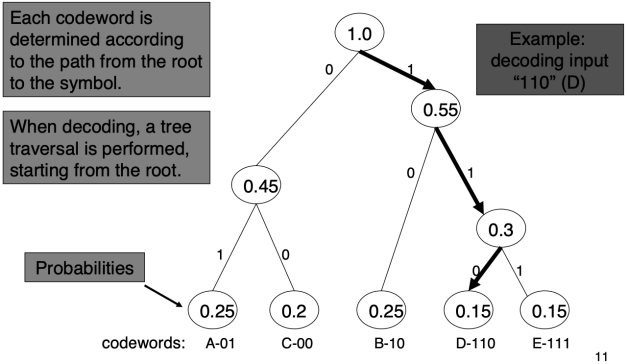
- prerequisite: lossless data compression
- $H(X)$ is compression limit, i.e. lower bound $L_{\text{avg}} \geq H(X)$
  which measures the minimum average number of bits needed to encode a symbol from a data source
- $H(X)$ is achievable, $L_{\text{avg}} \rightarrow H(X)$
  for a sufficiently long sequence of symbols it is possible to achieve an average code length close to $H(X)$
- Shannon's Source Coding Theorem [1]

$$N \cdot H(X)$$

## Entropy Coding

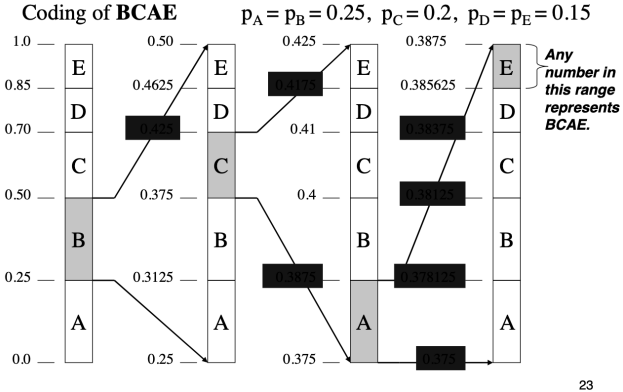
- Implementation of Shannon's Source Coding Theorem
- High-probability symbols get shorter codes
- Low-probability symbols get longer codes.
- A method that attempts to approach the lower bound $H(X)$
- Huffman Coding: ZIP, PNG
- Arithmetic Coding: JPEG, H.264, HEVC.
- Asymmetric Numeral System (ANS): Zstandard, AVIF, JPEG XL

## Huffman Coding




source: http://www.math.tau.ac.il/ dcor/Graphics/adv-slides/entropy.pdf

## Arithmetic Coding




source: http://www.math.tau.ac.il/ dcor/Graphics/adv-slides/entropy.pdf

# Asymmetric Numeral System

# Outline

2. Asymmetric Numeral System
- 2.1 Motivation
- 2.2 Key Concepts
  - 2.2.1 Symbol spread function
  - 2.2.2 Representation of information
  - 2.2.3 Symmetric or Asymmetric
- 2.3 Performance

## Motivation

- Huffman coding is efficient for symbol-based compression but cannot achieve optimal compression in many cases due to its restriction to integer-length codes
- Arithmetic coding achieves near-optimal compression but has a higher computational cost, especially in encoding and decoding ( it uses multiplication and division operations for each encoded symbol)
- ANS combines the compression ratio of arithmetic coding (integer arithmetic, modular arithmetic, and table lookups), with a processing cost similar to that of Huffman coding [4]

## Key Concepts

- **Symbol spread function** → ”Numeral System”, e.g. standard binary system

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- **Symbol spread function** → ”Numeral System”, e.g. standard binary system
- **Representation of information**, e.g. standard decimal system represents an entire message as a single number
- **Symmetric or ”Asymmetric”**

## Symbol spread function

- the function that determines the split of the $\mathbb{N}$ subsets corresponding to different symbols [3]

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- $\bar{s} : \mathbb{N} \rightarrow \mathcal{A}, \bar{s}(x) = \text{mod}(x, b) = s$

## Symbol spread function

- the function that determines the split of the $\mathbb{N}$ subsets corresponding to different symbols [3]
- $\bar{s} : \mathbb{N} \rightarrow \mathcal{A}, \bar{s}(x) = \text{mod}(x, b) = s$
- In the standard binary system example: $\bar{s}(x) = \text{mod}(x, 2)$
  i.e. splits natural numbers into subsets of even/odd numbers [3]

## Symbol spread function (Exa. 1)

$S = \{0, 1\}$

### Exa. 1: Standard Binary System

| $\mathbb{N}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $s = 0$ | | | | | | | | | | | | | | | | | | | |
| $s = 1$ | | | | | | | | | | | | | | | | | | | |

**Fig.1**

## Symbol spread function (Exa. 2)

$S = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$x' = 10x + s$

| $\mathbb{N}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | | 202 | | 2025 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $s = 0$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 1$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 2$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 3$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 4$ | | | | | | | | | | | | | | | | | | | | | | ... | | ... | |
| $s = 5$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 6$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 7$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 8$ | | | | | | | | | | | | | | | | | | | | | | | | | |
| $s = 9$ | | | | | | | | | | | | | | | | | | | | | | | | | |

**Fig.2**

## Representation of information (Exa. 2)

- a finite sequence of symbols/digits:

$$\mathcal{A} = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$$

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- the sequence of symbols/digits to be encoded:

$$S_{in} = (2, 0, 2, 5)$$

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- a finite sequence of symbols/digits:

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- the sequence of symbols/digits to be encoded:

$$S_{in} = (2, 0, 2, 5)$$

- How can we represent this sequence as a single number?
  The easiest way is to concatenate the symbols/digits:

$$x_{out} = 2025$$

## Representation of Information (Exa. 2)

- encoding the sequence of digits (2,0,2,5,1,8)

$$
\begin{aligned}
x_{\text{init}} &= 0 \\
x_1 &= x_{\text{init}} \times 10 + 2 = 2 \\
x_2 &= x_1 \times 10 + 0 = 20 \\
x_3 &= x_2 \times 10 + 2 = 202 \\
x_{\text{out}} &= x_3 \times 10 + 5 = 2025
\end{aligned}
$$

## Representation of Information (Exa. 2)

– What is the encoding function?

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- Encoding function

$$x_{i+1} = x_i \times 10 + s_i$$

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- What is the encoding function?
- Encoding function

$$x_{i+1} = x_i \times 10 + s_i$$

- $x_{i+1}$ is the $x_i$ -th appearance of symbol $s$ [3]

$$x_{i+1} = C(s, x_i) = x_i \cdot 10 + s$$

## Representation of Information (Exa. 2)

$x' = 10x + s$

$s = 0$ $s = 2$ $s = 2$ $s = 5$

| $\mathbb{N}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | | 202 | | 2025 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $s = 0$ | 0 | | | | | | | | | | 1 | | | | | | | | | | 2 | | | | |
| $s = 1$ | | 0 | | | | | | | | | | 1 | | | | | | | | | | | | | |
| $s = 2$ | | | 0 | | | | | | | | | | 1 | | | | | | | | | | | | |
| $s = 3$ | | | | 0 | | | | | | | | | | 1 | | | | | | | | | | | |
| $s = 4$ | | | | | 0 | | | | | | | | | | 1 | | | | | | | ... | | ... | |
| $s = 5$ | | | | | | 0 | | | | | | | | | | 1 | | | | | | | | | |
| $s = 6$ | | | | | | | 0 | | | | | | | | | | 1 | | | | | | | | |
| $s = 7$ | | | | | | | | 0 | | | | | | | | | | 1 | | | | | | | |
| $s = 8$ | | | | | | | | | 0 | | | | | | | | | | 1 | | | | | | |
| $s = 9$ | | | | | | | | | | 0 | | | | | | | | | | 1 | | | | | |

**Fig.3**

# Representation of Information (Exa. 1)

**Exa. 1: Standard Binary System**

$$x_{i+1} = 2x_i + s$$

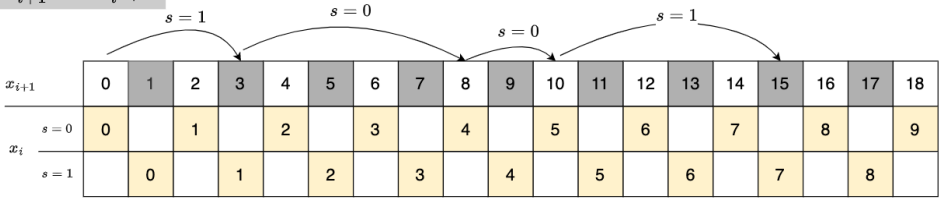


**Fig.4**

$x_{i+1} = 2x_i + s$

$s = 1$ $s = 0$ $s = 0$ $s = 1$

| $x_{i+1}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $x_i$ $s = 0$ | 0 | | 1 | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 | | 9 |
| $x_i$ $s = 1$ | | 0 | | 1 | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 | |

$x_{i+1} = 10x_i + s$

$s = 2$ $s = 0$ $s = 2$ $s = 5$

| $\mathbb{N}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | | 202 | | 2025 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $s = 0$ | 0 | | | | | | | | | | 1 | | | | | | | | | | 2 | | | | |
| $s = 1$ | | 0 | | | | | | | | | | 1 | | | | | | | | | | | | | |
| $s = 2$ | | | 0 | | | | | | | | | | 1 | | | | | | | | | | | | |
| $s = 3$ | | | | 0 | | | | | | | | | | 1 | | | | | | | | | | | |
| $s = 4$ | | | | | 0 | | | | | | | | | | 1 | | | | | | | ... | | ... | |
| $s = 5$ | | | | | | 0 | | | | | | | | | | 1 | | | | | | | | | |
| $s = 6$ | | | | | | | 0 | | | | | | | | | | 1 | | | | | | | | |
| $s = 7$ | | | | | | | | 0 | | | | | | | | | | 1 | | | | | | | |
| $s = 8$ | | | | | | | | | 0 | | | | | | | | | | 1 | | | | | | |
| $s = 9$ | | | | | | | | | | 0 | | | | | | | | | | 1 | | | | | |

**Fig.5**

# What is asymmetric?

$x_{i+1} = 2x_i + s$

$s = 1$ $s = 0$ $s = 0$ $s = 1$

| $x_{i+1}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $x_i$, $s = 0$ | 0 | | 1 | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 | | 9 |
| $x_i$, $s = 1$ | | 0 | | 1 | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 | |

$x_{i+1} = 10x_i + s$

$s = 2$ $s = 0$ $s = 2$ $s = 5$

| $\mathbb{N}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 | | 202 | | 2025 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $s = 0$ | 0 | | | | | | | | | | 1 | | | | | | | | | | 2 | | | | |
| $s = 1$ | | 0 | | | | | | | | | | 1 | | | | | | | | | | | | | |
| $s = 2$ | | | 0 | | | | | | | | | | 1 | | | | | | | | | | | | |
| $s = 3$ | | | | 0 | | | | | | | | | | 1 | | | | | | | | | | | |
| $s = 4$ | | | | | 0 | | | | | | | | | | 1 | | | | | | | ... | | ... | |
| $s = 5$ | | | | | | 0 | | | | | | | | | | 1 | | | | | | | | | |
| $s = 6$ | | | | | | | 0 | | | | | | | | | | 1 | | | | | | | | |
| $s = 7$ | | | | | | | | 0 | | | | | | | | | | 1 | | | | | | | |
| $s = 8$ | | | | | | | | | 0 | | | | | | | | | | 1 | | | | | | |
| $s = 9$ | | | | | | | | | | 0 | | | | | | | | | | 1 | | | | | |

**Fig.6**

## Exa. 3: Asymmetric Binary System [3]

- Let $\mathcal{A} = \{0, 1\}$
- $Pr(0) = 1/4, Pr(1) = 3/4$

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- Let $\mathcal{A} = \{0, 1\}$
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- a cycle of length 4, which contains 1 zero and 3 ones

## Exa. 3: Asymmetric Binary System [3]

- Let $\mathcal{A} = \{0, 1\}$
- $Pr(0) = 1/4, Pr(1) = 3/4$
- a cycle of length 4, which contains 1 zero and 3 ones
- Assumption:

$x_{i+1} = C(s, x_i) \approx x_i / p_s$, where $p_s$ is the probability of symbol $s$

# Exa. 3: Asymmetric Binary System [3]

$x_{i+1} \approx x_i / Pr(s)$ $\quad Pr(0) = 1/4, Pr(1) = 3/4$

| $x_{i+1}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | ... | 52 |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $x_i$ $s=0$ | | | | | | | | | | | | | | | | | | | | ... | |
| $x_i$ $s=1$ | | | | | | | | | | | | | | | | | | | | | |

**Fig.7_0**

# Exa. 3: Asymmetric Binary System [3]

$x_{i+1} \approx x_i / Pr(s)$ $\quad$ $Pr(0) = 1/4, Pr(1) = 3/4$

| | | | | | | | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $x_{i+1}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | ... | 52 |
| $x_i$ $s=0$ | | | | | | | | | | | | | | | | | | | | ... | |
| $x_i$ $s=1$ | | | | | | | | | | | | | | | | | | | | | |

**Fig.7_1**

# Exa. 3: Asymmetric Binary System [3]

$x_{i+1} \approx x_i / Pr(s)$ $\quad Pr(0) = 1/4, Pr(1) = 3/4$

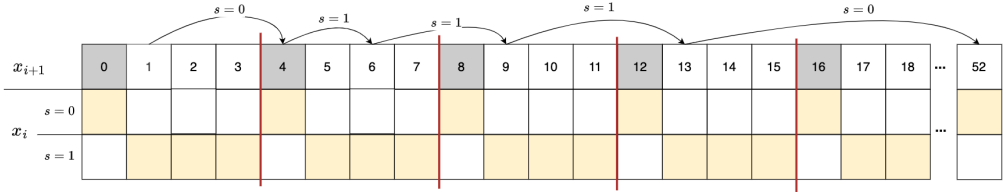



**Fig.7_2**

## Exa. 3: Asymmetric Binary System [3]

$x_{i+1} \approx x_i / Pr(s)$ $\quad Pr(0) = 1/4, Pr(1) = 3/4$

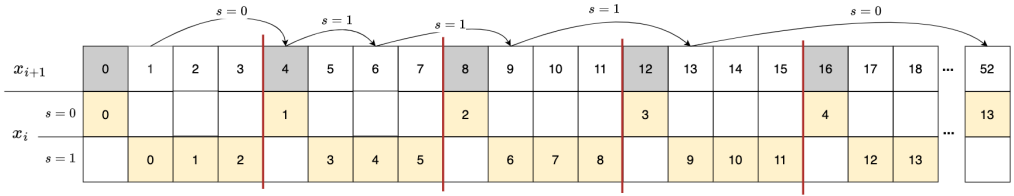



**Fig.7**

We refine even/odd number by changing their densities

## Symmetric or Asymmetric

- Symmetric: uniform distribution
- Asymmetric: non-uniform distribution

## Symmetric or Asymmetric

Assumption: $x_{i+1} = C(s, x_i) \approx x_i / p_s$, where $p_s$ is the probability of symbol $s$

- To change the symmetric behavior so that the information added to $x_i$ by coding a new symbol depends on the probability distribution of the symbols

## Symmetric or Asymmetric

Assumption: $x_{i+1} = C(s, x_i) \approx x_i / p_s$, where $p_s$ is the probability of symbol $s$

- To change the symmetric behavior so that the information added to $x_i$ by coding a new symbol depends on the probability distribution of the symbols
- $log_2(x_i)$ bits

## Symmetric or Asymmetric

Assumption: $x_{i+1} = C(s, x_i) \approx x_i / p_s$, where $p_s$ is the probability of symbol $s$

- To change the symmetric behavior so that the information added to $x_i$ by coding a new symbol depends on the probability distribution of the symbols
- $log_2(x_i)$ bits
- A symbol of probability pcarries $log_2(1/p)$ bits of information (Shannon)

## Symmetric or Asymmetric

Assumption: $x_{i+1} = C(s, x_i) \approx x_i/p_s$, where $p_s$ is the probability of symbol $s$

- To change the symmetric behavior so that the information added to $x_i$ by coding a new symbol depends on the probability distribution of the symbols
- $log_2(x_i)$ bits
- A symbol of probability pcarries $log_2(1/p)$ bits of information (Shannon)
- How many bits of new information? $log_2 x_i + log_2(1/p) \rightarrow log_2 x_{i+1}$

## Symmetric or Asymmetric

Assumption: $x_{i+1} = C(s, x_i) \approx x_i/p_s$, where $p_s$ is the probability of symbol $s$

- To change the symmetric behavior so that the information added to $x_i$ by coding a new symbol depends on the probability distribution of the symbols
- $log_2(x_i)$ bits
- A symbol of probability pcarries $log_2(1/p)$ bits of information (Shannon)
- How many bits of new information? $log_2 x_i + log_2(1/p) \rightarrow log_2 x_{i+1}$
- $log_2 x_{i+1} = log_2 x_i + log_2(1/p) = log_2(x_i/p) \Rightarrow x_{i+1} \approx x_i/p$

## Performance

- f

# rANS

# Outline

3. rANS
   3.1 Key Concepts
   3.2 Example: 101110
   3.3 Notations
   3.4 Encoding Function
   3.5 Implementation of Encoder
   3.6 Decoding Function
   3.7 Implemenation of Decoder

## rANS

- ”r” = ”range”: similar to range coding (RC)
- symbol spread function
- representation of information
- asymmetric

## Exa.1: Standard Binary System

– How to encode 101110? What is the encoding function $C(s, x_i)$?

$x_{i+1} = 2x_i + s$

$s = 0$ $s = 1$ $s = 1$

| | | | | | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|---|
| $x_{i+1}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 |
| $x_i$ $s=0$ | 0 | | 1 | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 | | 9 | |
| $x_i$ $s=1$ | | 0 | | 1 | | 2 | | 3 | | 4 | | 5 | | 6 | | 7 | | 8 | | 9 |
| $x_{i+1}$ | 20 | 21 | 22 | 23 | 24 | 25 | 26 | 27 | 28 | 29 | 30 | 31 | 32 | 33 | 34 | 35 | 36 | 37 | 38 | 39 |
| $x_i$ $s=0$ | 10 | | 11 | | 12 | | 13 | | 14 | | 15 | | 16 | | 17 | | 18 | | 19 | |
| $x_i$ $s=1$ | | 10 | | 11 | | 12 | | 13 | | 14 | | 15 | | 16 | | 17 | | 18 | | 19 |
| $x_{i+1}$ | 40 | 41 | 42 | 43 | 44 | 45 | 46 | 47 | 48 | 49 | 50 | 51 | 52 | 53 | 54 | 55 | 56 | 57 | 58 | 59 |
| $x_i$ $s=0$ | 20 | | 21 | | 22 | | 23 | | 24 | | 25 | | 26 | | 27 | | 28 | | 29 | |
| $x_i$ $s=1$ | | 20 | | 21 | | 22 | | 23 | | 24 | | 25 | | 26 | | 27 | | 28 | | 29 |

$s = 1$

$s = 0$

**Fig.4: 101110**

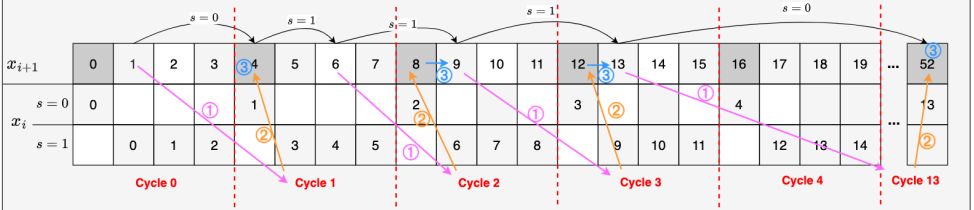
**Exa.3:** $Pr(0) = 1/4, Pr(1) = 3/4$

- How to encode 101110? What is the encoding function $C(s, x_i)$?
- How can we find the $x$ -th appearance of symbol $s$?
- $x_{i+1} = C(s, x_i) \approx x_i / p_s$

## Some Notations

- $\mathcal{A} = \{s_1, s_2, \ldots, s_k\}$ is the alphabet set
- $s$ is the symbol, if it is used in subscript: $s$ is the current symbol, $s-1$ is the previous symbol, $s+1$ is the next symbol
- $F_s$ is the absolute frequency of symbol $s$
- $m = \sum_{s \in \mathcal{A}} F_s$ is the total frequency
- $c_s = \sum_{s \in \mathcal{A}} F_{s-1}$ is the cumulative frequency ($c_{s+1}$ is the cumulative frequency of the next symbol)
- $p_s = F_s / m$ is the probability of symbol $s$
- $x_i$ is the current state (next state: $x_{i+1}$, previous state: $x_{i-1}$)

$x_{i+1} \approx x_i / Pr(s)$ | $Pr(0) = 1/4, Pr(1) = 3/4$




Given current state $x_i$ and next symbol $s$. How to find next state $x_{i+1}$, which is the $x_i$-th appearance of symbol $s$ ?

| STEPS / EXAMPLE | $Pr(0) = 1/4, Pr(1) = 3/4$ | | General |
|---|---|---|---|
| | $s = 0$ | $s = 1$ | |
| Step ①: Find cycle number $k$ | $k = \lfloor x_i/1 \rfloor$ | $k = \lfloor x_i/3 \rfloor$ | $k = \lfloor x_i/F_s \rfloor$ |
| Step ②: Navigate to the beginning (first number) of that cycle | $4 \cdot k$ | $4 \cdot k$ | $m \cdot k$ |
| Step ③: Find the position of the symbol $s$ in that cycle | part 1: first of one<br>+ 0 | part 1:<br>+ mod$(x_i, 3)$ | part 1:<br>+ mod$(x_i, F_s)$ |
| - part 1: position of the symbol $s$ in the sequence of the same symbol<br>- part 2: how many other symbols before the symbol $s$ | part 2: no other symbols<br>+ 0 | part 2: only one other symbol (0) and with one of that symbol in one cycle<br>+ 1 | part 2: only one other symbol (0) and with one of that symbol in one cycle<br>$+c_s$ |
| | $x_{i+1} = m \cdot k$ | $x_{i+1} = 4 \cdot \lfloor x_i/3 \rfloor + \text{mod}(x_i, 3) + 1$ | $x_{i+1} = m \cdot \lfloor x/F_s \rfloor + \text{mod}(x, F_s) + c_s$ |

$$x_{i+1} = C(s, x_i) = m \cdot \lfloor x_i/F_s \rfloor + \text{mod}(x_i, F_s) + c_s$$

**Fig.8**

## Encoding Function [3] [5]:

$$x_{i+1} = m \cdot \lfloor x_i / F_s \rfloor + \mathrm{mod}(x_i, F_s) + c_s$$

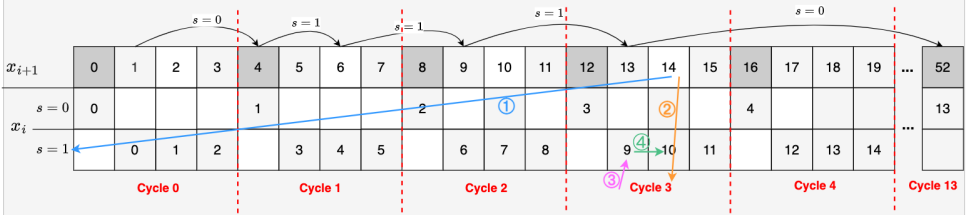
## implementation with python (dash app)

```python
def rANS_encoder(symbol_counts, s_input):
    cumul_freq = [] # c_s
    sum_freq = 0 # m
    x = 0 # initial state
    output = [] # list to store encoded (symbols, state)

    for count in symbol_counts:
        cumul_freq.append(sum_freq)
        sum_freq += count
    for s in s_input:
        F_s = symbol_counts[s]
        c_s = cumul_freq[s]
        x = sum_freq * (x // F_s) + (x % F_s) + c_s
        output.append((s, x))

    L_avg = math.ceil(math.log(x) / math.log(2.0)) / len(s_input)
    return output, x, L_avg, entropy(symbol_counts)
```

$x_{i+1} \approx x_i / Pr(s)$ $\quad Pr(0) = 1/4, Pr(1) = 3/4$




Given current state $x_i$. How to find previous state $x_{i-1}$, where $x_i$ is the $x_{i-1}$-th appearance of symbol $s$ ?

| STEPS / EXAMPLE | $Pr(0) = 1/4, Pr(1) = 3/4$ | | General |
|---|---|---|---|
| | $s = 0$ | $s = 1$ | |
| Step ① : Find the symbol $s$ for the state $x_i$ | if $x_i \bmod 4 = 0$ | otherwise or if $x_i \bmod 4 = 1, 2, 3$ | $s$ such that $c_s \leq \mathrm{mod}(x_i, m) < c_{s+1}$ |
| Step ②: Find the cycle $k$ of the current state $x_i$ | $k = x_i/4$ | $k = \lfloor x_I/4 \rfloor$ | $k = \lfloor x_i/m \rfloor$ |
| Step ③: Find the first number of that symbol $s$ in that cycle $k$ | $x_i/4 \cdot 1$ | $\lfloor x_i/4 \rfloor \cdot 3$ | $\lfloor x_i/m \rfloor \cdot F_s$ |
| Step ④: Allocate to the exact position | $x_{i-1} = x_i/4 \cdot 1 + 0 - 0$ | $x_{i-1} = \lfloor x_i/4 \rfloor \cdot 3 + 2 - 1$ | $x_{i+1} = \lfloor x_i/m \rfloor \cdot F_s + \mathrm{mod}(x_i, F_s) - c_s$ |

1. $s$ such that $c_s \leq \mathrm{mod}(x_1, m) < c_{s+1}$
2. $x_{i-1} = D(x_i) = \lfloor x_i/m \rfloor \cdot F_s + \mathrm{mod}(x_i, F_s) - c_s$

**Fig.9**

## Decoding Function [3] [5]:

- $s$

$$s \text{ such that } c_s \leq \text{mod}(x_i, m) < c_{s+1}$$

- $x_{i-1}$

$$x_{i-1} = D(x_i) = \lfloor x_i / m \rfloor \cdot F_s + \text{mod}(x_i, F_s) - c_s$$

## implementation with python (dash app)

```python
def rANS_decoder(symbol_counts, x):
    cumul_freq = []  # c_s
    sum_freq = 0 # m
    output = []  # List to store decoded symbols and states
    # compute m and c_s
    for count in symbol_counts:
        cumul_freq.append(sum_freq)
        sum_freq += count
    while x > 0:
        slot = x % sum_freq  # Find the slot in the cumulative frequency range
        s = bisect_right(cumul_freq, slot) - 1 # Binary search to find the symbol s (slot).
        F_s = symbol_counts[s]          # Frequency of the symbol
        c_s = cumul_freq[s]             # Cumulative count of the symbol
        output.append((s,x))            # Append decoded symbol to the output
        x = (x // sum_freq) * F_s + x % sum_freq - c_s  # Update the state
    return output, state
```

# Streaming-rANS

# Outline

4. Streaming-rANS
- 4.1 Problem
- 4.2 Idea
- 4.3 Solution
  - 4.3.1 Renormalization
- 4.4 Implementation
  - 4.4.1 Encoding
  - 4.4.2 Dncoding

## Streaming-rANS

To avoid using large number arithmetic, in practice stream variants are used: which enforce by renormalization: sending the least significant bits of x to or from the bitstream (usually L and b are powers of 2). [4]

Example of the problem //TODO: add example

- ANS (Asymmetric Numeral Systems) uses a single integer x, known as the "state," to represent the encoded data.
- Over time, as more symbols are encoded, the state x can grow arbitrarily large. This can lead to:
  - Overflow: If the value exceeds the precision limits of your system (e.g., a 32-bit or 64-bit register).
  - Inefficient Encoding: A larger x means you need more memory or bits to represent it.
- Using the C function from the previous section to encode succeeding symbols, we can encode a sequence into a very large natural number $x$, which can be stored as a length $\lfloor lg(x) \rfloor + 1$ bit sequence. [2]

## Renormalization in Range Coding

- Range Coding
  - 1
  - 2
- we use the similar idea for rANS
  - range $[L, H]$
  - Encoding: shrink the current state $x$ until it is in the range $[L, H]$
  - Decoding: expand the current state $x$ until it is in the range $[L, H]$
- In AC, renormalization is used to allow finite precision, an analogous approach should be used for ANS. Specifically, we will enforce xto remain in some fixed range I by transferring the least significant bits to the stream. For generality, let us transfer a base $b$ digit, for example $b = 2$ means transferring single bits, while $b = 2^8$ means transferring complete bytes.

## Methods from Jarek Duda: Method 1 - Decoding [2]

**ANS decoding step from state *x***

$(s, x) = D(x)$ *# decode the current state x*

useSymbol($s$)

**while** $x < L$ **do** *# if x is too small that not in the range [L,H]*

*# expand the x by reading the outstreamed digits back from the stream (an array)*

expand$(x, b) = b \cdot x +$ readDigit$() \rightarrow x$

**end while**

*# return the new state*

## What is $L, H$ ?

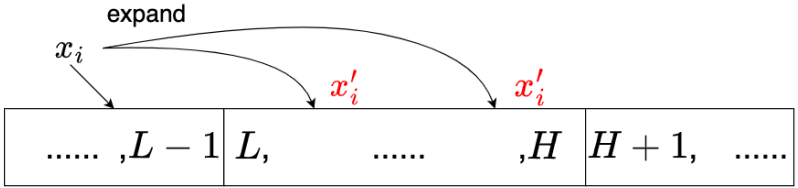



**Fig 11**

## What is $L, H$ ?

- $x_i = b \cdot x_{i-1} +$ bit
- Either $x_i$ in $[L, H]$ or $x_{i-1}$ in $[L, H]$. Assume $x_{i-1}$ in $[L, H]$.
  - $x_{i-1} >= L$
  - $x_i >= b \cdot L + bit >= b \cdot L$
  - $x_i$ can't be in $[L, H]$
  - derive: $H <= b \cdot L - 1$
- $[L, H]$, where $H <= b \cdot L - 1$ (Constraint 1)
- $I = \{L, ..., b \cdot L - 1\}$, which is defined by Jarek Duda $b$-**unique interval** and $L \in \mathbb{N}$
- This will ensure that the loop in Method 1 will always return to the original value

## Methods from Jarek Duda: Method 2 - Encoding [2]

shrink the current state $x$ by stream out the least significant bit ($x$ will be stored in binary and we take out the bit from right side)

**ANS encoding step of symbol sfrom state $x$**

*# once the current state x is greater than max value that we define later*

**while** $x > maxX[s]$ **do**

*# write the least significant digit to the stream (stream could be an array in implementaion )*

writeDigit(mod($x$, $b$))

shrink$(x, b) = \lfloor x/b \rfloor \rightarrow x$ *# shrink the x*

**end while**

$x = C(s, x)$ *# encode with the shrunk x*

## What is $maxX[s]$?

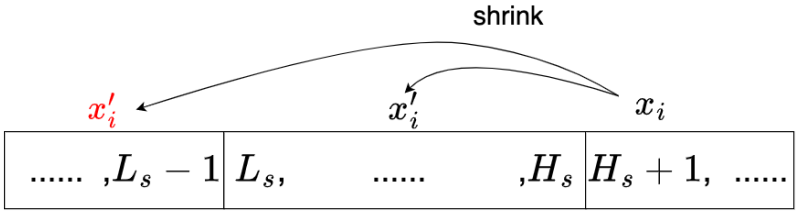



**Fig 10**

## What is $maxX[s]$?

- to make sure not to jump over the range
- Assume that $x_i = H_s + 1$ and it will be shrinked to $x'_i = \lfloor x_i/b \rfloor$
- $x'_i = \lfloor x_i/b \rfloor = \lfloor (H_s + 1)/b \rfloor >= L_s$
- $L_s <= (H_s + 1)/b <= \lfloor (H_s + 1)/b \rfloor < (H_s + 1)/b + 1$
- $H_s >= b \cdot L_s - 1, \forall s \in \mathcal{A}$ (Constraint 2)

## Final Form of $[L, H]$ in practice

- Contraints

$$[L, H]: \text{ where } H <= b \cdot L - 1$$
$$[L_s, H_s] : H_s >= b \cdot L_s - 1, \forall s \in \mathcal{A}$$

- In practice ($b = 2^1$): we choose $L = k \cdot m$, $L_s = k \cdot F_s$ ($k \in \mathbb{N}$)

$$\mathcal{I} = [km, 2 \cdot km - 1]$$
$$\mathcal{I}_s = [kF_s,\ 2 \cdot kF_s - 1] \text{ for } \forall s \in \mathcal{A}$$

## implementation with python (dash app)

```python
def rANS_stream_encoder(symbol_counts, s_input):
    # Init: c_s, m, x, output, bit_stream (omit the code here)
    for s in s_input:
        F_s = symbol_counts[s]
        c_s = cumul_freq[s]
        L_s = F_s # k = 1
        maxX_s = 2 * L_s - 1
        out_bits = ""

        while x > maxX_s:
            out_bits += str(x % 2)
            x = x // 2
        x = (x // F_s) * sum_freq + (x % F_s) + c_s
        bit_stream += out_bits
        output.append((s, x, out_bits, bit_stream))

    return output, x, bit_stream, L_avg, entropy(symbol_counts)
```

## implementation with python (dash app)

```python
def rANS_stream_decoder(symbol_counts, x, bit_stream, num_symbols):
    # Init: c_s, m, x, output, bit_stream (omit the code here)
    bit_stream = list(map(int, bit_stream.split(',')))  # Convert bit_stream to a list of integers
    for _ in range(num_symbols):
        slot = x % sum_freq
        s = bisect_right(cumul_freq, slot) - 1
        F_s = symbol_counts[s]
        c_s = cumul_freq[s]
        output.append((s, x))
        x = (x // sum_freq) * F_s + slot - c_s

        while x < sum_freq and len(bit_stream) > 0:
            x = x * 2 + bit_stream.pop()

    return output, x
```

# Summary

## Summary

- To summarize, this may help you to make a Leipzig university beamer presentation.

# Further Study

# Further Study | |

1. Huffman Codes: An Information Theory Perspective - Reducible
2. What is Asymmetric Numeral Systems? - Kedar Tatwawadi
3. Asymmetric Numeral Systems - EE274: Data Compression, course notes
4. Stanford Compression Library
5. Asymmetric Numeral Systems Toolkit - Jarek Duda
6. Presentation and Implementation - Huo Jiang

# Reference

# Reference


1. C. Shannon, “A mathematical theory of communication,” The Bell Sys- tem Technical Journal, vol. 27, pp. 379–423, 623–656, July, October
2. J. Duda, K. Tahboub, N. J. Gadgil and E. J. Delp, ”The use of asymmetric numeral systems as an accurate replacement for Huffman coding,” 2015 Picture Coding Symposium (PCS), Cairns, QLD, Australia, 2015, pp. 65-69, doi: 10.1109/PCS.2015.7170048.
3. J. Duda, “Asymmetric numeral systems: entropy coding combining speed of huffman coding with compression rate of arithmetic coding,” arXiv:1311.2540.
4. Asymmetric numeral systems - Wikipedia
5. What is Asymmetric Numeral Systems? - Kedar Tatwawadi 1948.

# Thank you!