Homework 4

4.1

a.

$$b_i = \sum_{k=1}^n p_{ki} \cdot a_k \ \sum_{i=1}^n b_i = \sum_{i=1}^n \sum_{k=1}^n p_{ki} \cdot a_k = \sum_{k=1}^n a_k = 1$$

 \therefore **b** is a probability vector.

Let

$$f(x) = -x \log x$$

$$\therefore f''(x) = -\frac{1}{x} < 0, x \in [0,1]$$

$$\therefore f(x)$$
 is concave on $[0,1]$. (deem $-\frac{1}{0}=-\infty<0$)

Since $\sum_{k=1}^n p_{ki}=1$, by Jensen's inequality:

$$-b_i \log b_i = f(b_i) = f(\sum_{k=1}^n p_{ki} \cdot a_k) \geq \sum_{k=1}^n p_{ki} f(a_k) = -\sum_{k=1}^n p_{ki} a_k \log a_k$$

Then

$$H(b_1,b_2,\ldots,b_n) = -\sum_{i=1}^n b_i \log b_i \geq -\sum_{i=1}^n \sum_{k=1}^n p_{ki} a_k \log a_k = -\sum_{k=1}^n a_k \log a_k = H(a_1,a_2,\ldots,a_n)$$

b.

Since $\mu=(\mu_1,\mu_2,\ldots,\mu_n)$ is a stationary distribution, we have

$$\mu_j = \sum_{k=1}^n \mu_k \cdot P_{kj}$$

This problem asks us to verify that uniform distribution this property, which is true because:

$$rac{1}{n}=\mu_j=rac{1}{n}\sum_{k=1}^n=\sum_{k=1}^n\mu_k\cdot P_{kj}$$

c.

Since $\mu=(\mu_1,\mu_2,\ldots,\mu_n)$ is a stationary distribution, we have

$$\mu_j = \sum_{k=1}^n \mu_k \cdot P_{kj}$$

$$\mu_j = rac{1}{n} = \sum_{k=1}^n \mu_k \cdot P_{kj} = rac{1}{n} \sum_{k=1}^n P_{kj}$$

which means

$$\sum_{k=1}^{n} P_{kj} = 1$$

And since P is Markov transition matrix, we have

$$\sum_{i=1}^{n} P_{kj} = 1$$

 $\therefore P$ is doubly stochastic.

4.3

$$H(TX) = H(TX|T) + I(TX;T) \ \geq H(TX|T)$$

If given T, we can reverse TX since $f(TX) = T^{-1}TX$ is a bijection. Therefore:

$$egin{aligned} H(TX) &= H(TX|T) + I(TX;T) \ &\geq H(TX|T) \ &= H(T^{-1}TX|T) \ &= H(X|T) \ &= H(X) \ , \ ext{given T and X are independent} \end{aligned}$$

4.6

a.

$$egin{aligned} rac{H\left(X_{1}, X_{2}, \ldots, X_{n}
ight)}{n} &= rac{\sum_{i=1}^{n} H\left(X_{i} \mid X^{i-1}
ight)}{n} \ &= rac{H\left(X_{n} \mid X^{n-1}
ight) + \sum_{i=1}^{n-1} H\left(X_{i} \mid X^{i-1}
ight)}{n} \ &= rac{H\left(X_{n} \mid X^{n-1}
ight) + H\left(X_{1}, X_{2}, \ldots, X_{n-1}
ight)}{n} \end{aligned}$$

From the inequalities below

$$H(X_{n+1} | X_n, ..., X_1)$$

 $\leq H(X_{n+1} | X_n, ..., X_2)$
 $= H(X_n | X_{n-1}, ..., X_1)$

We know that $\forall \ 1 \leq i \leq n$, $H(X_i|X^{i-1}) \geq H(X_n|X^{n-1})$, which means

$$H(X_n|X^{n-1}) \leq rac{1}{n-1} \sum_{i=1}^{n-1} H(X_i|X^{i-1}) = rac{1}{n-1} H(X_1,X_2,\dots,X_{n-1})$$

Therefore

$$egin{aligned} rac{H\left(X_{1},X_{2},\ldots,X_{n}
ight)}{n} &= rac{H\left(X_{n}\mid X^{n-1}
ight) + H\left(X_{1},X_{2},\ldots,X_{n-1}
ight)}{n} \ &\leq rac{H\left(X_{1},X_{2},\ldots,X_{n-1}
ight)}{n-1} \end{aligned}$$

b.

As mentioned in a.,

$$H(X_n|X^{n-1}) \leq rac{1}{n} \sum_{i=1}^n H(X_i|X^{i-1}) = rac{1}{n} H(X_1,X_2,\ldots,X_n)$$

4.7

a.

To calculate the stationary distribution, let

$$\mu \cdot P = \mu$$

We have

$$egin{align} \mu_1 &= rac{p_{10}}{p_{10} + p_{01}} \ \mu_2 &= rac{p_{01}}{p_{10} + p_{01}} \ \end{align*}$$

The entropy rate:

$$H(\mathcal{X}) = H(X_2|X_1) = \mu_1 H(p_{01}) + \mu_2 H(p_{10}) = rac{p_{10} H(p_{01}) + p_{01} H(p_{10})}{p_{10} + p_{01}}$$

b.

Since H(p) is concave,

$$rac{p_{10}H(p_{01})+p_{01}H(p_{10})}{p_{10}+p_{01}} \leq H(rac{2p_{10}p_{01}}{p_{01}+p_{10}}) \leq H(rac{1}{2})$$

According to the equality-hold condition of Jensen's inequality, the equality is reached if and only if $p_{10}=p_{01}=rac{1}{2}$.

c.

Let $p_{01} = p, p_{10} = 1$

$$H(\mathcal{X}) = H(X_2|X_1) = rac{p_{10}H(p_{01}) + p_{01}H(p_{10})}{p_{10} + p_{01}} = rac{H(p)}{1+p}$$

d.

Let $f(p) = \frac{H(p)}{1+p}$

$$f'(p) = \frac{2\log(1-p) - \log p}{(1+p)^2}$$

Make $f'(p) \geq 0$, we have $p \geq \frac{\sqrt{5}+3}{2}$ (excluded) or $p \leq \frac{3-\sqrt{5}}{2}$

The maximum entropy rate is $H(rac{\sqrt{5}+3}{2})pprox 0.694 bits$

We might as well assume $\mu = (\mu_1, \mu_2)$ respectively denotes two states: 0 and 1.

Then from the transition matrix, we notice that if the current state is 1, then the next state must be 0 while there is no constraint when the current state is 0.

From this observation, we could divide the sequence into two cases:

- 1. Started with 1, then the next state is 0, and the following t-2 states are not fixed.
- 2. Started with 0, then the next t-1 states are not fixed.

Then we have

$$N(t) = N(t-1) + N(t-2), t \ge 3$$

For base cases, we have N(1) = 2, N(2) = 3

This form implies that $\{N(t)\}$ is Fibonacci sequence, which have the expression :

$$N(t) = rac{1}{\sqrt{5}} \left[\left(rac{1+\sqrt{5}}{2}
ight)^{t+1} - \left(rac{1-\sqrt{5}}{2}
ight)^{t+1}
ight]$$

Since $\left|\frac{1-\sqrt{5}}{2}\right| < 1$, we have

$$\lim_{t o\infty}rac{1}{t}{
m log}\,N(t)=\lim_{t o\infty}[rac{t+1}{t}{
m log}(rac{1+\sqrt{5}}{2})+rac{1}{t}{
m log}\,rac{1}{\sqrt{5}}]pprox {
m log}\,rac{1}{\sqrt{5}}pprox 0.694bits$$

Since there are N(t) cases of sequence $\{X_i\}_{i=1}^t$, the upper bound of $H(X_1, X_2, \ldots, X_t)$ is $\log N(t)$. And the entropy rate of Markov chain is $\lim_{t\to\infty} \frac{1}{t} H(X_1, X_2, \ldots, X_t)$, thus $H_0 = \lim_{t\to\infty} \frac{1}{t} \log N(t)$ is an upper bound on the entropy rate of Markov chain.

Actually, we see the upper bound is reached in part (d).

4.10

a.

- 1. If $i, j \in \{1, 2, ..., n-1\}$, it has been pointed out that X_i, X_j are i.i.d. random variables, thus independent.
- 2. If $i \in \{1, 2, \dots, n-1\}, j = n$.

We want to prove $p(\sum_{i=1}^n X_i \text{ is odd}) = p(\sum_{i=1}^n X_i \text{ is even}) = \frac{1}{2}$ by induction.

Assume that it's true for k-1, then

$$P(\sum_{i=1}^k X_i ext{ is odd}) = p(\sum_{i=1}^{k-1} X_i ext{ is odd}) p(X_k = 0) + p(\sum_{i=1}^{k-1} X_i ext{ is even}) p(X_k = 1) = rac{1}{2}$$

Therefore,

$$p(X_n=1)=p(X_n=0)=\frac{1}{2}$$

WOLG, we just focus on the case below:

$$p(X_i = 1, X_n = 1) = p(X_i = 1, \sum_{k=1, k
eq i}^n X_k ext{ is even}) = rac{1}{2} \cdot rac{1}{2} = p(X_i = 1) p(X_n = 1)$$

Other cases are completely similar, which proves that X_n, X_i are independent.

b.

Since X_i, X_j are independent, we have $H(X_i, X_j) = H(X_i) + H(X_j) = 2H((X_i)) = 2bits$

C.

Since X_n is the function of $X_1, X_2, \ldots, X_{n-1}$, we have:

$$egin{align} H(X_1,X_2,\ldots,X_n) &= \sum_{i=1}^n H(X_i|X_1,X_2,\ldots,X_{i-1}) \ &= H(X_n|X_1,X_2,\ldots,X_{n-1}) + \sum_{i=1}^{n-1} H(X_i) \ &= (n-1)H(X_1)
eq nH(X_1) \end{aligned}$$

4.12

a.

Since this is a Markov chain, we have

$$egin{align} H(X_1,X_2,\dots,X_n) &= H(X_1) + \sum_{i=2}^n H(X_i|X_1,X_2,\dots,X_{i-1}) \ &= \sum_{i=1}^n H(X_i|X_{i-1}) \ &= H(X_1) + (n-1)H(p) \ \end{gathered}$$

Since the first step is equally like to be positive or negative, we have $H(X_1)=1bits$

Therefore,

$$H(X_1,X_2,\ldots,X_n) = H(X_1) + (n-1)H(p) pprox 1 + 0.469(n-1) \quad (bits)$$

b.

Entropy rate is given by

$$\lim_{n o\infty}rac{1}{n}H(X_1,X_2,\ldots,X_n)=H(p)pprox 0.469$$

C.

Let q = 1 - p = 0.9

$$egin{align} E(S) &= p \sum_{k=0}^{\infty} (k+2) q^{k+2} \ &= p \lim_{k o \infty} [rac{q(1-q^k)}{(1-q)^2} + rac{2}{1-q} - rac{(k+2) q^{k+1}}{1-k}] \ &= 11 \end{aligned}$$

$$egin{aligned} &\lim_{n o \infty} rac{1}{n} H(X_n, \dots, X_1 | X_0, X_{-1}, \dots, X_{-k}) \ &= \lim_{n o \infty} rac{1}{n} \sum_{i=1}^n H(X_i | X_{i-1}, X_{i-2}, \dots, X_{-k}) \ &= \lim_{n o \infty} H(X_n | X_{n-1}, X_{n-2}, \dots, X_{-k}) \ &= H(\mathcal{X}) \end{aligned}$$

4.20



守卫

能不能4.20就做一个King啊, 不然rook,bishop(黑) bishop(白)和queen还要把类 似的事情重复5遍偷



警长

确实, 五遍很离谱...

上午7:03



🗤 19 群主 程帆 计算机导论老师 🧳

那就挑两个做吧

下午3.18

King

The stationary distribution μ_i is E_i/E , where E_i is the number of adjacent points that can be reached by King, E is the sum up of E_i . There are several types of E_i :

- $E_i = 3, i \in \{1, 3, 7, 9\}$
- $ullet E_i = 5, i \in \{2,4,6,8\}$
- $E_i = 8, i = 5$

Then E=40. Since it's of same probabilty for the king to choose a possible direction, we have

- $ullet \ H(X_2|X_1=x) = \log 3 \ ext{bits}, x \in \{1,3,7,9\}$
- ullet $H(X_2|X_1=x)=\log 5$ bits, $x\in\{2,4,6,8\}$
- ullet $H(X_2|X_1=x)=\log 8=3$ bits, x=5

$$H(\mathcal{X}) = H(X_2|X_1) = \sum_{} p(x)H(X_2|X_1 = x)$$

$$= \sum_{} \mu H(X_2|X_1 = x)$$

$$= (4 \cdot \frac{3}{40}\log 3 + 4 \cdot \frac{5}{40}\log 5 + \frac{8}{40} \cdot 3) \text{ bits}$$

$$\approx 2.24 \text{ bits}$$

Queen

- $E_i = 6, i \in \{1, 3, 7, 9, 2, 4, 6, 8\}$
- $E_i = 8, i = 5$

Then E=56. Since it's of same probabilty for the queen to choose a possible direction, we have

- $\bullet \ \ H(X_2|X_1=x)=\log 6 \ {
 m bits}, x\in\{1,3,7,9,2,4,6,8\}$
- $H(X_2|X_1=x) = \log 8 = 3$ bits, x=5

$$H(\mathcal{X}) = H(X_2|X_1) = \sum_{} p(x)H(X_2|X_1 = x)$$

$$= \sum_{} \mu H(X_2|X_1 = x)$$

$$= (8 \cdot \frac{6}{56}\log 6 + \frac{8}{56} \cdot 3) \text{ bits}$$

$$\approx 2.644 \text{ bits}$$

4.33

$$\begin{split} &I(X_1;X_4) + I(X_2;X_3) - I(X_1;X_3) - I(X_2;X_4) \\ &= H(X_1) - H(X_1|X_4) + H(X_2) - H(X_2|X_3) - H(X_1) + H(X_1|X_3) - H(X_2) + H(X_2|X_4) \\ &= -H(X_1|X_4) - H(X_2|X_3) + H(X_1|X_3) + H(X_2|X_4) \\ &= H(X_1,X_2|X_3) - H(X_2|X_1,X_3) - (H(X_1,X_2|X_3) - H(X_1|X_2,X_3)) \\ &+ H(X_1,X_2|X_4) - H(X_1|X_2,X_4) - (H(X_1,X_2|X_4) - H(X_2|X_1,X_4)) \\ &= H(X_1|X_2,X_3) - H(X_2|X_1,X_3) + H(X_2|X_1,X_4) - H(X_1|X_2,X_4) \end{split}$$

Since this is a markov chain, we have $H(X_1|X_2,X_3)=H(X_1|X_2,X_4)$, as the information of X_1 only relates to X_2 .

Therefore:

$$egin{aligned} I(X_1;X_4) + I(X_2;X_3) - I(X_1;X_3) - I(X_2;X_4) \ &= H(X_1|X_2,X_3) - H(X_2|X_1,X_3) + H(X_2|X_1,X_4) - H(X_1|X_2,X_4) \ &= -H(X_2|X_1,X_3) + H(X_2|X_1,X_4) \ &= -H(X_2|X_1,X_3,X_4) + H(X_2|X_1,X_4) \ &= I(X_2;X_3|X_1,X_4) > 0 \end{aligned}$$

4.34

From data process inequality, we know that $I(X;Y) \geq I(X;Z,W)$, therefore

$$\begin{split} I(X;Y) + I(Z;W) - I(X;Z) - I(X;W) \\ &\geq I(X;Z,W) + I(Z;W) - I(X;Z) - I(X;W) \\ &= H(Z,W) - H(Z,W|X) + H(Z) - H(Z|W) - H(X) + H(X|Z) - H(X) + H(X|W) \\ &= H(Z,W) - H(Z,W,X) + H(X) + H(Z) + H(W) - H(Z,W) - H(X) - H(Z) \\ &+ H(X,Z) - H(X) - H(W) + H(X,W) \\ &= -H(Z,W,X) + H(X,Z) - H(X) + H(X,W) \\ &= H(W|X) - H(W|X,Z) \\ &= I(W;Z|X) \geq 0 \end{split}$$