

Please refer to the notes of the previous lecture for the new Project3 grammar.

Some updates on the grammar: (LT:less than, GT:greater than, etc.)

Refer also to `/syllabus/projects/lexical_specification.md`:

- New binary operations
- Negation is unary
  - Part of factor
- The `!` character is either
  - Part of the `!=` token
  - The unary negation

```
expression
= expression PLUS expression
| expression MINUS expression
| expression TIMES expression
| expression DIVIDE expression
| expression MOD expression
| expression EQUALS expression
| expression NEQUALS expression
| expression LT expression
| expression GT expression
| expression AND expression
| expression OR expression
| NOT expression
| LPAREN expression RPAREN
| INTEGER
| IDENTIFIER
```

SimpleC now has:

- `==` for equality
- `!=` for inequality
- `<` for less than
- `>` for greater than
- `&&` for conjunction
- `||` for disjunction
- `!` for negation

The `icmp` instruction (integer compare) does integer comparisons. SimpleC variables are always 32-bit signed integers. 1 represents true, zero represents false in LLVM.

Ex:

```
%t2 = icmp slt i32 %t1, 10
```

`%t2` holds a 1-bit integer. `icmp` returns a 1-bit integer. `slt` is less-than for signed integers.

LLVM Instructions for Booleans:

Booleans are 1-bit integers in LLVM.

Ex:

```
%t3 = and i1 %t1, %t2
```

```
%t4 = or i1 %t3, %t1
```

LLVM has no unary negation. Use xor instead.

Ex:

```
%t5 = xor i1 %t4, 1
```

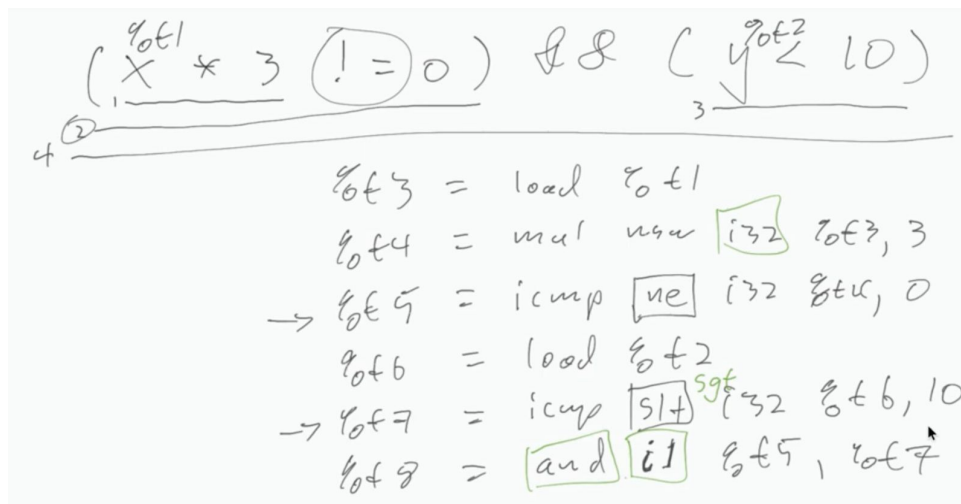
Demo: Boolean Operations and Comparisons

Ex:

```
(x * 3 != 0) && (y < 10)
```

// x \* 3 is executed first, and then first parentheses, next is y < 10 and lastly the and operator.

Let's assume x is allocated to %t1 and y is allocated to %t2.



How do we guarantee booleans are i1 types?

The OR operator takes a signed integer and an integer

We can use a type system to make this formal

Ex: The value zero then it is true. Nonzero then false.

```
#include <stdio.h>

int main(int argc, char **argv) {
    if (argc - 1) {
        printf("hello\n");
    }

    return 0;
}
```

Compile:

```
gcc -o test test.c
```

```
./test // hello
```

To see the llvm:  
gcc -S -O0 test.c  
less test.s

### Operator Types:

- Arithmetic operators: + - \* / %
  - (int, int) -> int
- Comparison operators: == != > <
  - (int, int) -> bool
- Boolean operators: && ||
  - (bool, bool) -> bool
- Reminder:
  - "(bool, bool)" is the list of parameters to the function
  - "-> bool" is the return type

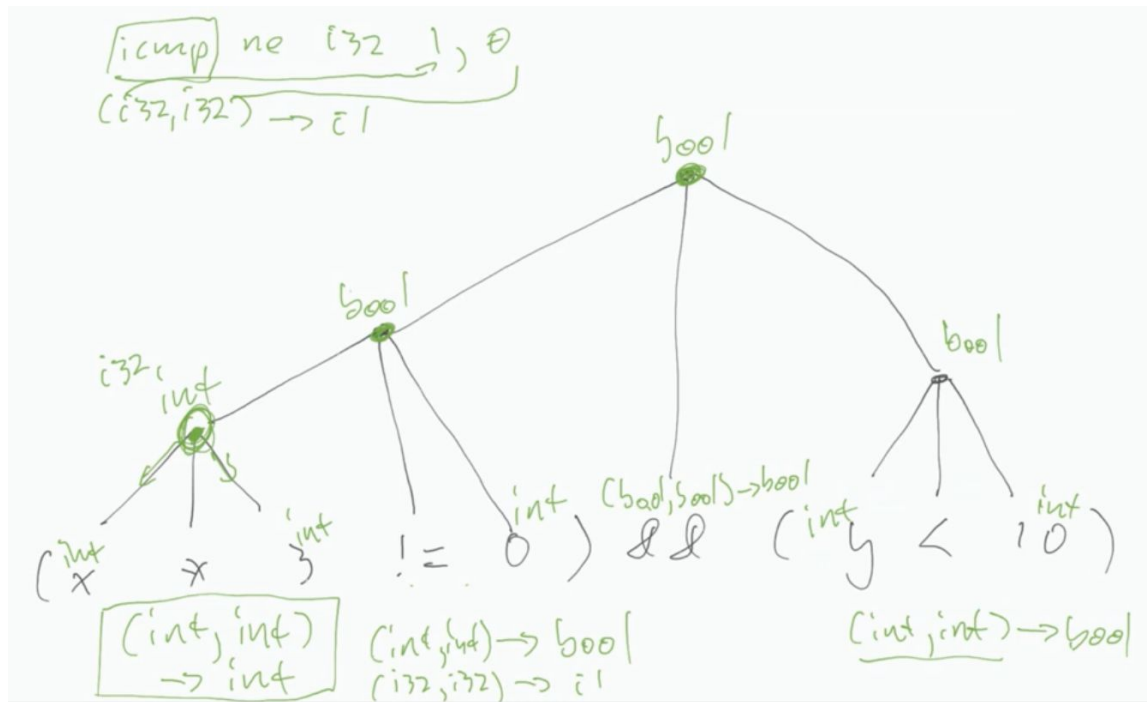
Type checking guarantees correct types, i1 is used only in booleans.  
Type checker can reject the program, require a rewrite

### Demo: Type-checking expressions

Type of x is stored in symbol table as int.

3 is just a number as int;

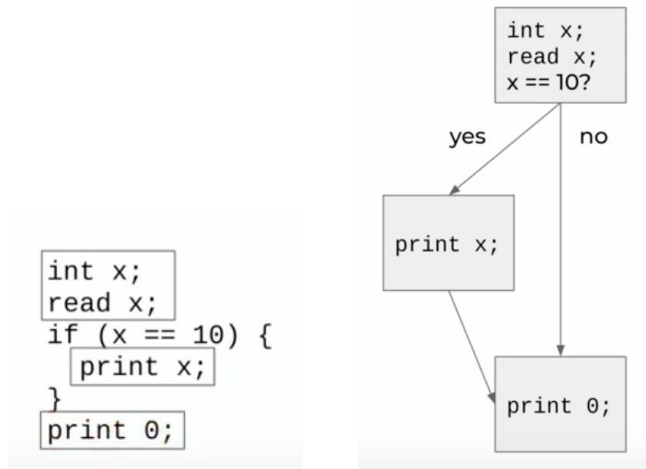
x\*3 (int, int) -> int. Result of operation satisfies the boolean (!=) and zero is number (int)



Implementing if-statements:

Each branch means a mutually exclusive choice

Only one branch to be executed



If x is 11 it will print 0.

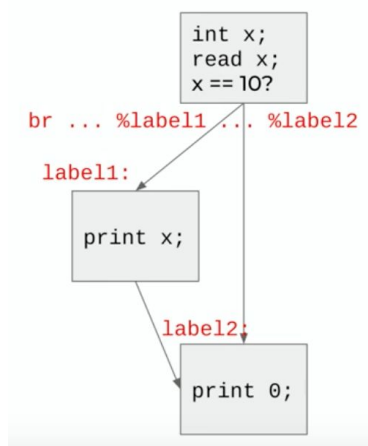
If it is 10, it will print 10 and 0.

Multiple out edges can be a branch

In-edges are targets of the branch

Branching in LLVM:

```
br i1 %cond, label %label1, label %label2  
label1:  
    ; next instruction is %cond is true  
  
label2:  
    ; next instruction is %cond is false
```



The branch instruction goes right after comparison and labels get inserted right before “print x” and “print 0”

```
; "int x;" allocate space for x
%t1 = alloca i32 ; allocate space for x
; "read x;" read x from input
%t2 = call i32 @read_integer() ; read an integer from stdin
store i32 %t2, i32* %t1 ; store the result of read_integer
; compute x == 10
%t3 = load i32, i32* %t1 ; get value of x
%cond = icmp eq i32 %t3, 10 ; do comparison
; if (x == 10)
br i1 %cond, label %label1, label %label2
label1: ; body of the if statement
%t4 = load i32, i32* %t1 ; get value of x
call void @print_integer(i32 %t4) ; print the value of x
br label %label2
label2: ; after the if statement
call void @print_integer(i32 0) ; print 0
ret i32 0
```