Database Management Systems (COP 5725)

Spring 2016

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Exam 2 Solutions

	Name:	
	UFID:	
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P	ledge (Must be sig	gned according to UF Honor Code)
	On my honor, I hav ssignment.	re neither given nor received unauthorized aid in doing this
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S	ignature	

	Maximum	Received
Question 1	25	
Question 2	45	
Question 3	10	
Question 4	20	
Total	100	

For scoring use only:

Question 1 (Relational Algebra) [30 points]

Consider the following relation schema below (primary keys are underlined).

```
project (pno, pname, plocation, dnum) [dnum references dnumber in department]

work_on (essn, pno, hours) [essn references ssn in employee]

department (dnumber, dname, mgrssn, mgrstartdate)

[mgrssn refers to an ssn of a manager]

employee (ssn, fname, lname, bdate, addr, sex, salary, superssn, dno)

[superssn refers to an ssn of a supervisor, dno references dnumber in department]

dependent (essn, dependent_name, sex, bdate, relationship)

[essn references ssn in employee]
```

Write the following queries in Relational Algebra.

1. Find the names of all managers who have at least one dependent. [5 points]

```
\Pi_{\text{ fname, lname}}((\Pi_{\text{ mgrssn}}(\text{department}) \cap \Pi_{\text{ essn}}(\text{dependent})) \bowtie (\text{employee}))
```

2. Find the last names and ssn of all employees who work on project number 1 and on project number 2. [5 points]

```
\Pi_{ssn, lname} (\sigma_{pno=1} (works_on) \bowtie employee) \cap \Pi_{ssn, lname} (\sigma_{pno=2} (works_on) \bowtie employee)
```

3. Find the names of all employees who earn more than their supervisors. [5 points]

```
Supervisors \leftarrow \Pi superssn (employee)
SupsSals \leftarrow \rho_{ssal=salary} (\Pi superssn, salary (employee \bowtie_{ssn=superssn} Supervisors))
\Pi fname,lname (\sigma salary > ssal (employee \bowtie_{superssn=superssn} SupsSals))
```

4. List the names of all employees whose department manager earns less than \$25,000. [5 points]

```
DeptsWithPoorMng \leftarrow \Pi_{dnumber} (department \bowtie_{mgrssn=ssn} (\sigma_{salary < 25,000} \text{ (employee)})

\Pi_{fname, lname} (employee \bowtie_{dno=dnumber} DeptsWithPoorMng)
```

5. Find the names of all employees who work on the project that is conducted by department 5. [5 points]

```
\Pi fname, lname ((\Pi pno, essn(works on) \div \Pi pno (\sigma dnum=5(project))) \bowtie employee)
```

Question 2 (SQL) [45 points]

A) Consider the following table schemas:

```
branch (branch_name, branch_city, assets)
customer (customer_name, customer_street, customer_city)
account (account_number, branch_name, balance)
loan (loan_number, branch_name, amount)
depositor (customer_name, account_number)
borrower (customer_name, loan_number)
```

Write SQL statements for the following queries:

1. Find the number of depositors for each branch. [4 points]

```
Select branch_name, count (distinct customer_name) from depositor, account where depositor.account_number = account.account_number group by branch_name
```

2. Find all customers who have both an account and a loan at the Gainesville branch. [5 points]

3. Find the average balance of all customers who live in Gainesville and have more than three accounts. [5 points]

```
Select depositor.customer_name, avg(balance)
From depositor, account, customer
Where depositor.account_number = account.account_number
and depositor.customer_name = customer.customer_name
and customer.customer_city = 'Gainesville'
group by depositor.customer_name
having count(distinct depositor.account_number) >= 3
```

4. Find the name of the branch where the average account balance is greater than the other branches. [6 points]

B) Consider the following table schemas:

```
dept (deptno, dname, loc) emp (empno, ename, job, hiredate, sal, deptno) (:employee's information)
```

Write SQL statements for the following queries:

5. Find the employees (empno, name, job, salary) whose job is the same as the employee with empno = 7777 and who receive less salary than the employee with empno = 8888. [7 points]

```
Select empno, ename, job, sal from emp where job =

(
select job from emp where empno = 7777
)

and sal <
(
select sal from emp where empno = 8888
)
```

6. Find the employees (that is, output deptno, dname, ename, job, hiredate) who work in the sales department and were hired in 2015. (in ascending order by hiredate) (4 points)

```
select e.deptno, d.dname, e.ename, e.job, e.hiredate from emp e, dept d where e.deptno = d.deptno and e.hiredate >= to_date('2015-01-01') and e.hiredate <= to_date('2015-12-31') and d.dname = 'sales' order by hiredate asc
```

7. Find the departments (that is, output deptno, number of employees) for which more employees work than for other departments. [7 points]

8. Find the employees (that is, output empno, name, deptno, hiredate) who have worked longer than others in each department. [7 points]

```
select empno, ename, deptno, hiredate
from emp
where (deptno, hiredate) in (
select deptno, min(hiredate)
from emp
group by deptno)
```

Question 3 (Relational Algebra and SQL) [10 points]

The division (or quotient) operator \div of the Relational Algebra does not have a direct equivalent in SQL. It identifies attribute values from a relation that are paired with all of the values from another relation. Without loss of generality, let $\mathbf{R} = \{A_1 : C_1, ..., A_n : C_n, B_1 : D_1, ..., B_m : D_m\}$ and $\mathbf{S} = \{B_1 : D_1, ..., B_m : D_m\}$ be two relation schemas. Let R be a relation with respect to \mathbf{R} and S be a relation with respect to \mathbf{S} .

1. Provide the formal definition of the division operator by means of a Relational Algebra expression that only makes use of the basic Relational Algebra operators. [4 points]

```
R \div S = \pi_{R-S}(R) - \pi_{R-S}((\pi_{R-S}(R) \times S) - R)
```

2. A first solution to mapping the definition in 1. into SQL is somewhat difficult to understand. It makes use of a doubly nested and negated SQL statement and is based on the *not exists* predicate. We are here not interested in this solution. Instead, we follow another idea that maps the different parts of the definition in 1. *one by one* into SQL and does not use the *not exists* predicate. Write down the corresponding SQL query. [6 points]

```
SELECT DISTINCT A1, ..., An
FROM R
MINUS
SELECT A1, ..., An
FROM (SELECT * FROM (SELECT A1, ..., An FROM R), S)
MINUS
R
);
```

It is obvious that the "..." have to be replaced by the respective attributes $A_2, ..., A_{n-1}$.

Question 4 (QBE) [20 points]

Consider the following database schema:

Drivers (did, dname, gender, age)
Reserve (did, cid, day, cost)
Cars (cid, cname, model, color, rid)
RentalCompany (rid, rname, revenue, rating)
IsMember(did, rid, join time, member type)

Primary key attributes are underlined.

Answer the following questions using QBE. Draw tables in your answer.

1. Find the oldest driver who is a member of the company 'Budget' and the company 'Avis'. [7 points]

Drivers	Did	dname	gender	age
P.	_Id			_A
っ (negation)	_Id2			>_A
IsMember	Did	rid	join_time	member_type
	_Id	_Rid1		
	_Id	_Rid2		
	_Id2	_Rid1		
	_Id2	_Rid2		
RentalCompany	Rid	rname	revenue	rating
	_Rid1	'Avis'		
	_Rid2	'Budget'		

2. Find the name of the customer who has reserved a car named 'A6' on '01/03/2016' from the company 'Budget'. [6 points]

Drivers D	Did	dname	gender	age
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	_Id	PS					
Reserve	Did	cid		day		co	ost
	_Id	_Cid		' 01/03	3/2016'		
						•	
Cars	cid	cname	model		color		rid
	_Cid	'A6'					_Rid
	1	<u>'</u>			·		
RentalCompany	Rid	rname	rname		revenue		ting
	_Rid	'Budget'					

3. Update the member type to 'VIP' for those drivers who were members of company 'Avis' and have spent more than 1000 in renting (reserving) cars from Avis. [7 points]

Drivers	Did	dname	gender	age
	_Did			
IsMember	Did	rid	join_time	member_type
U.	_Did	_Rid		'VIP'
RentalCompany	Rid	rname	revenue	rating
	_Rid	'AVIS'		
Reserve	Did	cid	day	cost
	GDid			SUM.ALLX
	<u> </u>	<u> </u>	<u> </u>	<u> </u>

Condition (Reserve)	
SUM.ALL. X > '1000'	