### **Referential Integrity (V)**

- □ Standard behavior of the DBMS when violating a referential integrity constraint: rejection of the action that caused the violation, i.e., the transaction performing the update action is rolled back
- Attempt to update or delete a tuple in the referenced relation *R* that has a matching tuple in the referencing relation *S* depends on the referential action specified using the **on delete** or **on update** sub-clauses of the **foreign key** clause
- ☐ Possible referential actions for the on delete sub-clause
  - on delete cascade: The tuple (including its primary key) of the referenced relation *R* as well as all the matching tuples (including their foreign keys) in a referencing relation *S* are deleted.
  - ❖ on delete set null: The tuple of the referenced relation R is deleted, and the foreign key values in all matching tuples of a referencing relation S are set to *null*. This option is only valid if the foreign key attributes in S do not have the *not null* constraint specified.

### **Referential Integrity (VI)**

- ☐ Possible referential actions for the on delete sub-clause (continued)
  - ❖ on delete set default: The tuple of the referenced relation R is deleted, and the foreign key values in all matching tuples of a referencing relation S are set to their specified default values. This option is only valid if all foreign key attributes in S have a default value specified.
  - ❖ on delete no action: If a tuple of a relation S exists that references a tuple of a relation R, the deletion of the tuple in R is rejected. This is the default setting if the on delete rule is omitted.
- ☐ Possible referential actions for the on update sub-clause
  - on update cascade: A change of the value of the primary key attributes in a tuple of the referenced relation R is propagated to all tuples in a referencing relation S with the same value in their foreign key attributes.
  - ❖ on update set null: The value of the primary key attributes in the tuple of the referenced relation *R* is changed. The foreign key values in all matching tuples of a referencing relation *S* are set to *null*. This option is only valid if the foreign key attributes in *S* do not have the *not null* constraint specified.

### **Referential Integrity (VII)**

- ☐ Possible referential actions for the on update sub-clause (continued)
  - ❖ on update set default: The value of the primary key attributes in the tuple of the referenced relation *R* is changed. The foreign key values in all matching tuples of a referencing relation *S* are set to their specified default values. This option is only valid if all foreign key attributes in *S* have a default value specified.
  - ❖ on update no action: If a tuple of a relation S exists that references a tuple of a relation R, the update of the primary key attribute values of the tuple in R is rejected. This is the default setting if the on update rule is omitted.
- Examples

```
create table lectures
(...,
foreign key(held_by) references professors(pers-id)
on delete cascade
on update cascade);
```

# **Referential Integrity (VIII)**

- Examples (continued)
  - The deletion of a tuple with the key *pers-id* in *professors* is permitted and has the effect on *lectures* that there all referencing tuples are deleted. This enables the realization of a dependent relationship.
  - A change of a value of the attribute pers-id in professors is propagated to all matching values of the foreign key attribute held\_by in lectures
  - create table lectures
    (...,
    foreign key(held\_by) references professors(pers-id)
    on delete set null
    on update set null);
    - The value of the attribute held\_by in lectures is set to null, if the referenced tuple with the key pers-id in professors is deleted
    - The value of the attribute held\_by of all referencing tuples in lectures is set to null, if the value of the primary key attribute pers-id of the referenced tuple in professors is changed

# **General Integrity**

An assertion is a predicate that expresses a condition which is to be always satisfied by a database system
Domain constraints and referential ICs are special kinds of assertions
There are also conditions which cannot be expressed with these two kinds like conditions with respect to several relations
Syntax: create assertion <assertion name=""> check <condition></condition></assertion>
Example: There must be at least four professors in order to maintain teaching create assertion AlwaysFourProfessors check (4 <= (select count(*) from professors))
The DBMS tests an assertion for validity first, and if the assertion is valid, future modifications of the database are only allowed if the assertion is not violated
Complex assertions can lead to an overhead; they should be used with care
Most DBS have not implemented assertions

### Triggers (I)

#### □ A trigger

- is a user-defined procedure that the DBMS executes automatically if a certain condition is fulfilled or as a side effect of a modification of the database
- is a general and powerful mechanism for maintaining data(base) consistency, monitoring database updates, and updating database statistics
- cannot only deploy check functions but also computation functions
- ☐ A trigger is based on the ECA (Event-Condition-Action) model and has three components:
  - The event(s) to which it reacts
    - These are usually database manipulation operations (insert, delete, update) that are explicitly applied to the database
    - Person who writes the trigger must ensure that all possible events are accounted for
    - Trigger must specify if it should be executed before or after an event

### **Triggers (II)**

- The condition that determines whether the action should be executed
  - If no condition is specified, the action will be executed once the event occurs
  - If a condition is specified in the when clause of the trigger, it is first evaluated, and only if it evaluates to true, the action will be executed
- The action(s) to be taken
  - The action is usually a sequence of SQL statements
  - It can also be a database transaction or an external program that will be automatically executed
- ☐ Part of the SQL syntax for creating a trigger in Oracle

```
create trigger <trigger name>
{before | after}
{insert | delete | update [of <column list>]} on 
[for each row]
[when <condition>]
  <trigger actions>;
```

#### Triggers (III)

- ☐ In Oracle the trigger actions are implemented in PL/SQL, which is Oracle's database application programming language
- ☐ Example: Trigger preventing that professors can be demoted by a rank

```
create trigger noDemotion
before update on professors
for each row
```

when (:old.rank is not null)

#### begin

- if :old.rank = "C3" and :new.rank = "C2" then :new.rank = "C3" end if;
- if :old.rank = "C4" then :new.rank = "C4" end if;
- if :new.rank is null then :new.rank = :old.rank end if;

#### end

- ☐ The clause **create trigger** provides a trigger specification
- The time of releasing the trigger body can be before or after the data manipulation operation that has released the trigger

#### Triggers (IV)

- □ Trigger events
  - ❖ update [of <column1, column2, ...>] on <relation name>
  - insert on <relation name>
  - delete on <relation name>
- A trigger can be defined for one or several events. In case of several events, a case distinction can be expressed in the body through the clauses
  - ❖ if updating [<column1, column2, ...>] then ...
  - if inserting then ...
  - if deleting then ...
- □ Trigger type
  - ❖ A statement-level trigger (default) is released exactly once either before or after the respective event.
  - A row-level trigger is indicated by the **for each row** clause and signifies that the trigger is executed *for each changed tuple* in the table

# **Triggers (V)**

- ☐ Distinguishing tuples before and after a change in a row-oriented trigger
  - Keyword :old refers to a deleted tuple or to a tuple before it was updated
  - Keyword :new refers to a newly inserted or newly updated tuple
  - This is the only possibility to address tuples before or after a change
  - Another access to the relation is not possible any more, even if the relation name would be used in the respective block
- ☐ Trigger restriction
  - The when predicate specifies a condition that must be satisfied to release the execution of the trigger body
  - ❖ If a row-oriented trigger is used, the new resp. old tuple of the relation can be addressed by the keywords :new resp. :old
- ☐ Trigger body / trigger actions
  - ❖ SQL queries
  - Oracle-specific PL/SQL commands with the aforementioned extensions

#### **Triggers (VI)**

■ Example: Protocol of the changes of the attribute *salary* of a relation *Persons* **create trigger** StoreSalary before update on Persons for each row **when** (:old.salary > 1500) begin insert into diff values (:old.salary, :new.salary, sysdate) end; ☐ Example: Check at insertion time that a salary increase is inapplicable to persons with a monthly salary greater than \$10,000 create trigger CheckSalary before update on Persons for each row when (:new.salary > 10000) begin :new.salary := :old.salary; // assignment only possible for before update end;