Question 1 (Relational Algebra) [30 points]

Consider the following relation schema below (primary keys are underlined).

```
project (pno, pname, plocation, dnum) [dnum references dnumber in department]

work_on (essn, pno, hours) [essn references ssn in employee]

department (dnumber, dname, mgrssn, mgrstartdate)

[mgrssn refers to an ssn of a manager]

employee (ssn, fname, lname, bdate, addr, sex, salary, superssn, dno)

[superssn refers to an ssn of a supervisor, dno references dnumber in department]

dependent (essn, dependent_name, sex, bdate, relationship)

[essn references ssn in employee]
```

Write the following queries in Relational Algebra.

1. Find the names of all managers who have at least one dependent. [5 points]

```
\Pi_{\text{fname, lname}} ((\Pi_{\text{mgrssn}} (\text{department}) \cap \Pi_{\text{essn}} (\text{dependent})) \bowtie (\text{employee}))
```

2. Find the last names and ssn of all employees who work on project number 1 and on project number 2. [5 points]

```
\begin{split} &\Pi_{ssn, \ lname} \ (\sigma_{pno=1} \ (works\_on) \bowtie employee) \cap \\ &\Pi_{ssn, \ lname} \ (\sigma_{pno=2} \ (works\_on) \bowtie employee) \end{split}
```

3. Find the names of all employees who earn more than their supervisors. [5 points]

```
Supervisors \leftarrow \Pi superssn (employee)

SupsSals \leftarrow \rho_{ssal=salary} (\Pi superssn, salary (employee \bowtie_{ssn=superssn} Supervisors))

\Pi fname,Iname (\sigma salary > ssal (employee \bowtie_{superssn=superssn} SupsSals))
```

4. List the names of all employees whose department manager earns less than \$25,000. [5 points]

```
DeptsWithPoorMng \leftarrow \Pi_{dnumber} (department \bowtie_{mgrssn=ssn} (\sigma_{salary < 25,000} \text{ (employee)})

\Pi_{fname, lname} (employee \bowtie_{dno=dnumber} DeptsWithPoorMng)
```

5. Find the names of all employees who work on the project that is conducted by department 5. [5 points]

```
\Pi fname, lname ((\Pi pno, essn(works on) \div \Pi pno (\sigma dnum=5(project))) \bowtie employee)
```

Question 2 (SQL) [45 points]

A) Consider the following table schemas:

```
branch (branch_name, branch_city, assets )
customer (customer_name, customer_street, customer_city)
account (account_number, branch_name, balance)
loan (loan_number, branch_name, amount)
depositor (customer_name, account_number)
borrower (customer_name, loan_number)
```

Write SQL statements for the following queries:

1. Find the number of depositors for each branch. [4 points]

```
Select branch_name, count (distinct customer_name) from depositor, account where depositor.account_number = account.account_number group by branch_name
```

2. Find all customers who have both an account and a loan at the Gainesville branch. [5 points]

3. Find the average balance of all customers who live in Gainesville and have more than three accounts. [5 points]

```
Select depositor.customer_name, avg(balance)
From depositor, account, customer
Where depositor.account_number = account.account_number
and depositor.customer_name = customer.customer_name
and customer.customer_city = 'Gainesville'
group by depositor.customer_name
having count(distinct depositor.account_number) >= 3
```

4. Find the name of the branch where the average account balance is greater than the other branches. [6 points]

B) Consider the following table schemas:

```
dept (deptno, dname, loc) emp (empno, ename, job, hiredate, sal, deptno) (:employee's information)
```

Write SQL statements for the following queries:

5. Find the employees (empno, name, job, salary) whose job is the same as the employee with empno = 7777 and who receive less salary than the employee with empno = 8888. [7 points]

```
Select empno, ename, job, sal from emp where job =

(
select job from emp where empno = 7777
)

and sal <
(
select sal from emp where empno = 8888
)
```

6. Find the employees (that is, output deptno, dname, ename, job, hiredate) who work in the sales department and were hired in 2015. (in ascending order by hiredate) (4 points)

```
select e.deptno, d.dname, e.ename, e.job, e.hiredate from emp e, dept d where e.deptno = d.deptno and e.hiredate >= to_date('2015-01-01') and e.hiredate <= to_date('2015-12-31') and d.dname = 'sales' order by hiredate asc
```

7. Find the departments (that is, output deptno, number of employees) for which more employees work than for other departments. [7 points]

8. Find the employees (that is, output empno, name, deptno, hiredate) who have worked longer than others in each department. [7 points]

```
select empno, ename, deptno, hiredate
from emp
where (deptno, hiredate) in (
select deptno, min(hiredate)
from emp
group by deptno)
```

Question 3 (Relational Algebra and SQL) [10 points]

The division (or quotient) operator \div of the Relational Algebra does not have a direct equivalent in SQL. It identifies attribute values from a relation that are paired with all of the values from another relation. Without loss of generality, let $\mathbf{R} = \{A_1 : C_1, ..., A_n : C_n, B_1 : D_1, ..., B_m : D_m\}$ and $\mathbf{S} = \{B_1 : D_1, ..., B_m : D_m\}$ be two relation schemas. Let R be a relation with respect to \mathbf{R} and S be a relation with respect to \mathbf{S} .

1. Provide the formal definition of the division operator by means of a Relational Algebra expression that only makes use of the basic Relational Algebra operators. [4 points]

```
R \div S = \pi_{R-S}(R) - \pi_{R-S}((\pi_{R-S}(R) \times S) - R)
```

2. A first solution to mapping the definition in 1. into SQL is somewhat difficult to understand. It makes use of a doubly nested and negated SQL statement and is based on the *not exists* predicate. We are here not interested in this solution. Instead, we follow another idea that maps the different parts of the definition in 1. *one by one* into SQL and does not use the *not exists* predicate. Write down the corresponding SQL query. [6 points]

```
SELECT DISTINCT A1, ..., An
FROM R
MINUS
SELECT A1, ..., An
FROM (SELECT * FROM (SELECT A1, ..., An FROM R), S)
          MINUS
          R
        );
```

It is obvious that the "..." have to be replaced by the respective attributes $A_2, ..., A_{n-1}$.

Question 4 (QBE) [20 points]

Consider the following database schema:

Drivers (did, dname, gender, age)
Reserve (did, cid, day, cost)
Cars (cid, cname, model, color, rid)
RentalCompany (rid, rname, revenue, rating)
IsMember(did, rid, join_time, member_type)

Rid2

Primary key attributes are underlined.

Answer the following questions using QBE. Draw tables in your answer.

1. Find the oldest driver who is a member of the company 'Budget' and the company 'Avis'. [7 points]

Drivers	Did	dname	gender	age
P.	_Id			_A
¬ (negation)	_Id2			>_A
IsMember	Did	rid	join_time	member_type
	_Id	_Rid1		
	_Id	_Rid2		
	_Id2	_Rid1		
	_Id2	_Rid2		
	1			
RentalCompany	Rid	rname	revenue	rating
	_Rid1	'Avis'		

2. Find the name of the customer who has reserved a car named 'A6' on '01/03/2016' from the company 'Budget'. [6 points]

'Budget'

Drivers	Did	dname	gender	age
---------	-----	-------	--------	-----

	_Id	PS					
Reserve	Did	cid		day		C	ost
	_Id	_Cid	_Cid		'01/03/2016'		
	·					•	
Cars	cid	cname	model		color		rid
	_Cid	'A6'					_Rid
				1			
RentalCompany	Rid	rname	rname		revenue		ating
	_Rid	'Budget	,				

3. Update the member type to 'VIP' for those drivers who were members of company 'Avis' and have spent more than 1000 in renting (reserving) cars from Avis. [7 points]

Drivers	Did	dname	gender	age
	_Did			
IsMember	Did	rid	join_time	member_type
U.	_Did	_Rid		'VIP'
RentalCompany	Rid	rname	revenue	rating
	_Rid	'AVIS'		
Reserve	Did	cid	day	cost
	GDid			SUM.ALLX
	ı		ı	

Condition (Reserve)	
SUM.ALLX > '1000'	