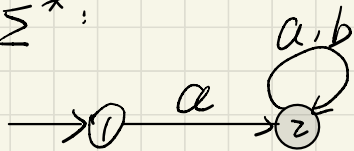
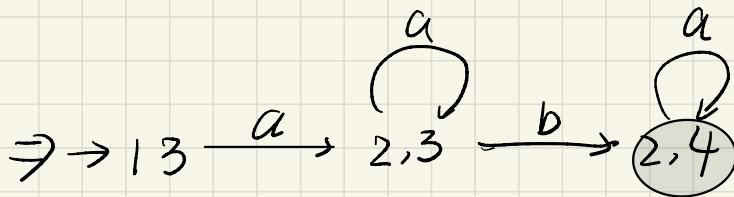
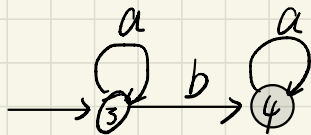


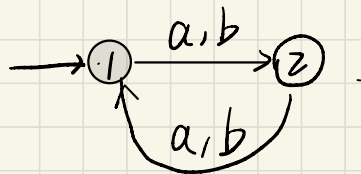
1. (e) ①  $a \Sigma^*$ :



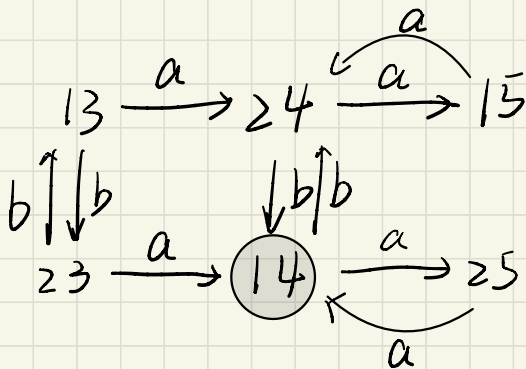
②  $a^* b a^*$



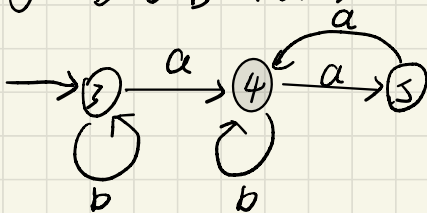
(g) ①  $((a+b)(a+b))^*$



$\Rightarrow$

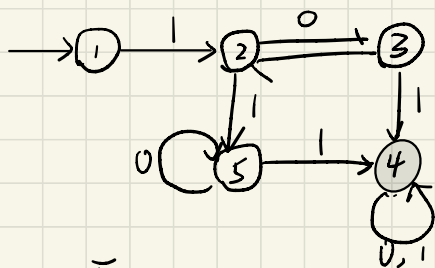


②  $b^* a b^* (a a)^*$



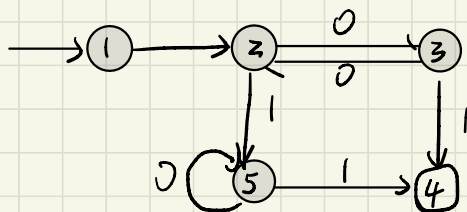
2.

DFA for  $\overline{F}$ :



1 1 1 1

DFA for  $F$ :

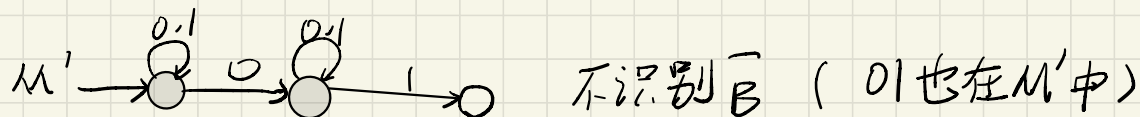
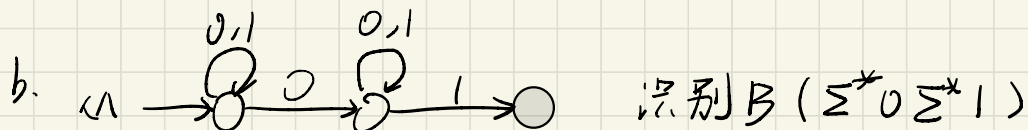


3.

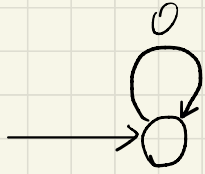
a.  $M = (Q, \Sigma, \delta, q_0, F)$

$M' = (Q, \Sigma, \delta, q_0, Q - F)$

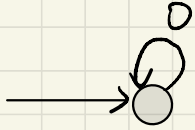
对于  $w = a_1 \dots a_n$  且  $w \in M$ , 则  $w$  一定停在  $q \in F$  上; 对于  $w \notin M$ , 则  $w$  一定停在  $q \in Q - F$  上, 因此  $M'$  识别  $\bar{B} (\Sigma^* - B)$



4. 考虑  $L$  为空语言, 则有  $N$  识别  $L$ ;



有  $N$  为

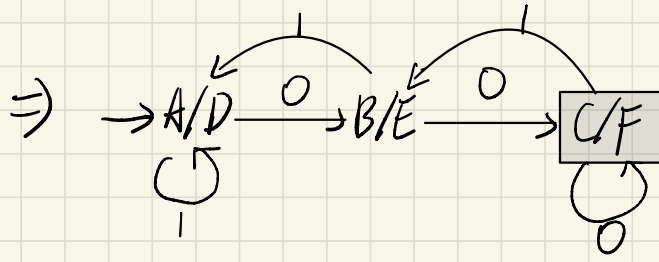


识别  $0^* \neq \emptyset$  得证.

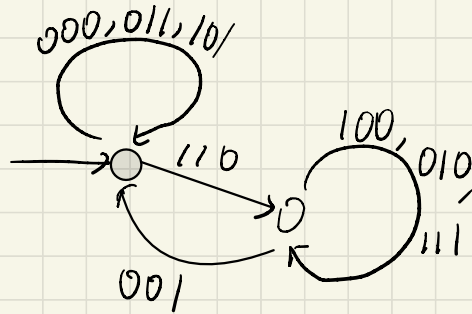
5. initialize:

<del>A</del>	<del>B</del>	0	1	<del>B</del>	<del>C</del>	0	1	CF	0	1	FIF	0	1	<del>D</del>	<del>E</del>	0	1	EIF
A	C	B	I	B	E	C	I	F	D	D								
A	D	B	I	E	A	A												
<del>A</del>	<del>E</del>	B	I	F														

$\therefore A \equiv D, B \equiv E, C \equiv F$



6.  $B$  is regular  $\Leftrightarrow B^R$  is regular. So prove  $B^R$  is regular.  
that is to find a FA  $L(Q, \Sigma, \delta, q_0, F)$  that recognizes  $B$ .



Note that every arc is denoted by transpose of the "letter".

Proven.

7.

$$\begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix}$$

Assume  $D$  to be regular, then there is a pumping length  $n$ .

For any given  $w \in D$  and  $|w| > n$ ,  $w = xyz$  where  $|xy| \leq n$ .

Let  $w$  be  $\begin{bmatrix} 1 \\ 0 \end{bmatrix}^n \begin{bmatrix} 0 \\ 1 \end{bmatrix}^n$ , thus  $y = \begin{bmatrix} 1 \\ 0 \end{bmatrix}^i$  where  $0 < i \leq n$

So  $xy^3z = \begin{bmatrix} 1 \\ 0 \end{bmatrix}^n \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix}^n \notin D$ .

Contradiction.