NETWORK AUDIT AND SECURITY PROJECT

REALIZATION OF A "BRUTEFORCE" ATTACK FOR ACCESS TO A WPA-PSK PROTECTED WIFI NETWORK

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1 Introduction

Wireless security is a crucial aspect of staying safe online. Connecting to the internet over insecure links or networks is a security risk that could potentially lead to data loss, leaked account credentials, and the installation of malware on your network. Using the proper Wi-Fi security measures is critical – but in doing so, it's important to understand the differences between different wireless encryption standards, including WEP, WPA, WPA2, and WPA3.

Wi-Fi Protected Access (WPA) is a security standard for computing devices with wireless internet connections. It was developed by the Wi-Fi Alliance to provide better data encryption and user authentication than Wired Equivalent Privacy (WEP), which was the original Wi-Fi security standard. Since the late 1990s, Wi-Fi security types have gone through multiple evolutions to improve them.

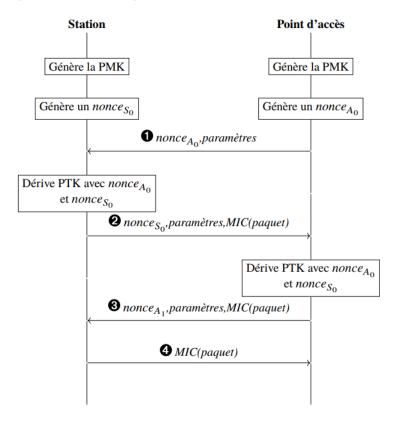


Figure 1: The 4-way handshake WPA/WPA2

In this project, we will learn how to attack WPA-PSK security by "brute force". Through the analysis of the captured packet, we will reconstruct step by step the original WPA-PSK encryption, thereby creating a password dictionary and performing the attack.

2 Analysis and Attack

In this section, we will go into detail about analyzing target packets with two types of encryption WPA (TKIP) and WPA2 (AES) respectively. From there, build an attack method and create a complete tool. And we will use Wireshark as a support tool to analyze and get information from the captured package

2.1 WPA

We'll start with the captured package **capture_wpa.pcap**. This file is taken from the **Audit & Sécurité réseaux** page at http://p-fb.net/.

2.1.1 Detect type of encryption

For the first, we need to detect type of encryption. We will look at the "Key Information" section in the first package to detect the type of encryption in use. And we got $RC4\ Cipher,\ HMAC-MD5\ MIC$

Figure 2: Encryption Type

2.1.2 PMK

The first step of WPA encryption is the calculation of the PMK. So we will build a dictionary for PSK and look up the SSID from the captured package

1. Build PSK dictionary

Based on the suggestion given by the project, we know that all letters are lowercase and start with four 'a' characters, and PSK length is 8. So we will build a dictionary for PSK with any set of 4 lowercase characters. This is Python code:

```
import itertools

letters = 'abcdefghijklmnopqrstuvwxyz'

words = [''.join(x) for x in itertools.product(letters, repeat=4)]

print(len(words))

with open('pwd_list.txt', 'w') as file:
    for word in words:
        file.write(''.join(word) + '\n')
```

Listing 1: PSK generate

2. SSID

After opening the package with Wireshark, we can immediately see the |SSID| = |WPAM1|

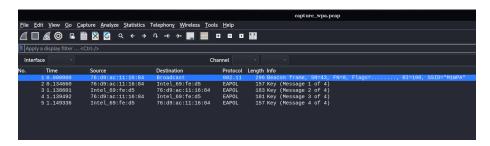


Figure 3: SSID

2.1.3 PTK

After obtaining the PMK, we will proceed to calculate the PTK. To achieve this we need to search for $Nonce_A$, $Nonce_S$, MAC_{Client} , $MAC_{Authenticator}$, Finnaly, use this information to calculate the PTK using the PRF_512 function.

```
PTK = PRF\_512 (PMK, "Pairwise \ key \ expansion", LowerMAC | || HigherMAC || || HigherNonce || || HigherNonce || HigherNonce
```

1. $Nonce_A$, $Nonce_S$, MAC_{Client} , $MAC_{Authenticator}$

Based on analyzing the captured package using Wireshark, we got the following information:

2. **PRF_512**

Based on the tutorial we have built the PRF_512 calculation function as follows:

```
def PRFn(self, K, A, B, n = 512):
    i = 0
    R = b''

while len(R)*8 < n:
    data = A + b'\x00' + B + bytes([i])
    r = hmac.new(K, data, sha1).digest()
    R += r
    i += 1
    return R[:n//8]</pre>
```

Listing 2: PRF_512

With:

A ="Pairwise key expansion",

B = LowerMAC |||HigerMAC|||LowerNonce|||HigherNonce||

2.1.4 MIC generation

The next step, once we have the necessary information, we will proceed to calculate the MIC from the corresponding EAPOL. Continuing to use Wireshark and based on the original data exchange diagram, we look at the packages in turn and get the following target MIC values and EAPOL (include MIC):

First target MIC and First EAPOL

From second frame, we have:

Figure 4: First target MIC

```
00 00 aa aa 03 00 00 00
                                 88 8e
         00 20 00 00 00
                         00
                                 00 00 02 ee a4 12 4e 3f
0040
                             00
0050
      ac f8 e0 27 0d b5 87 fc
      6b 93 1f c2 6d 35 b4 c7
0060
                                 db bb ae 00 00 00 00 00
0070
         00 00 00 00 00
                                              00 00
0080
         00 00 00 00 00 00 00
      00
                                    00 00
                                                     ec e5
      24 d3 99 17 9c bc 03 9e
                                    39 e4 00 1a dd 18 00
0090
                                                              P · · · · · P · · · · · P ·
00a0
      50 f2 01 01 00 00 50 f2
                                 02 01 00 00 50 f2 02 01
      00 00 50 f2 02 2a 00
00b0
```

Figure 5: First EAPOL

|MIC1| = |082793ece524d399179cbc039e0239e4|

Second target MIC and Second EAPOL

From third frame, we have:

Figure 6: Second target MIC

```
00 00 18 00 2e 48 00 00
                                00 6c 6c 09 c0 00 be 01
                                                               . Н
                                                                          5i
      00 00 00 00 00 00 00
                               88 02 24 00 00 0e 35 69
                                                                      $
      fe d5 76 d9 ac 11 16 84
                                76 d9 ac 11 16 84 20 00
                                                                    V.
                                   8e 01 03 00 77 fe 01
00 03 7c 67 f2 24 a6
0030
      00 00 aa aa 03 00 00 00
                                88 8e 01
0040
         00
               00
                  00
                     00 00 00
                                00
0050
      e0 81 93 23 0f ee b0 ef
                                f9 a0 7e c6 cb f0 16 3f
      96 2b a3 4d 31 db db 2b
0060
                                c6 9d 8d 00 00 00 00 00
      00 00 00 00 00 00 00
                                00 00 00 00 00 00 00
0080
      00 00 00
               00 00 00
                        00 00
                                   00 00
                                                           89 a8 1d 42 23 82 81 9e
0090
      50 f2 01 01 00 00 50 f2
                                02 01 00 00 50 f2 02 01
00a0
00b0
      00 00 50 f2 02
                                                            · P · ·
```

Figure 7: Second EAPOL

MIC2 = e2180d61d789a81d422382819e3efe4e

Third target MIC and Second EAPOL

From fourth frame, we have:

Figure 8: Third target MIC

-

```
00 00 18 00 2e 48 00 00
                                00 02 6c 09 a0 00 d2 01
      00 00 00 00 00 00 00 00
                                88 01 3a 01 76 d9 ac 11
      16 84 00 0e 35 69 fe d5
                                76 d9 ac 11 16 84 20 00
                                                               ·5i·
                                                                     ٧
0030
      00 00 aa aa 03 00 00 00
                                88 8e <mark>01 03 00</mark>
                                                5f
                                00 00 03 00 00 00 00 00
0040
         00
            20 00 00 00 00 00
0050
      00 00 00 00 00 00 00 00
                                00 00 00 00 00 00 00 00
                                00 00 00 00 00 00 00 00
0060
      00 00 00 00 00 00 00 00
0070
      00 00 00 00 00 00 00
                                00 00 00 00 00 00 00
0080
      00 00 00 00 00 00 00 00
                                                25 cc f2
                        f2 b2
                                      d2 00 00
0090
```

Figure 9: Third EAPOL

MIC3 = adda25ccf2fcaecfd18b37f2b2ffafd2

MIC Calculation

We will calculate the MIC from each PSK in the dictionary and EAPOL without MIC, then compare it with the target MIC, thereby finding the correct password.

```
MIC = HMAC_{MD5}(KCK, EAPOL_{noMIC})
```

With KCK:

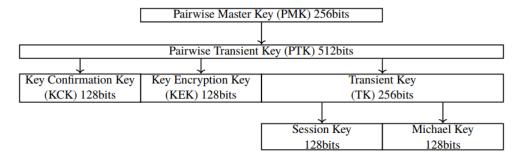


Figure 10: Key in WPA

2.1.5 Bruteforce Attacking

This is last step in our project. After getting all of information, we proceed to program script for doing brute force PSK. We obtained following result:

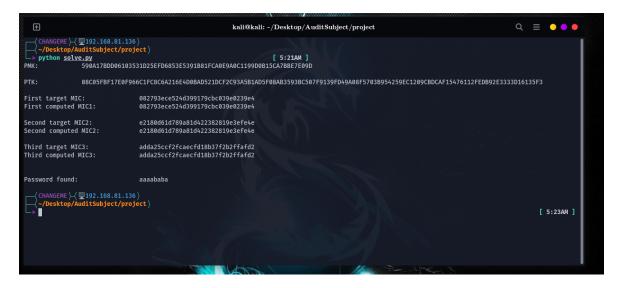


Figure 11: Final Result

As we can see, all of MICs were matched. Then we got |PSK| = |aaaababa|

2.2 WPA2

In this section, we'll start with the captured package **wpa-Induction.pcap**. This file is taken from https://wiki.wireshark.org/SampleCaptures.

Similar to the WPA part, we also follow the same steps

2.2.1 Detect type of encryption

Look at the Key Information section, we got AESCipher, HMAC - SHA1MIC

Figure 12: Encryption Type

2.2.2 PMK

1. Build PSK dictionary

Assume that we knew first 5 characters are "Induc". We will create PSK dictionary with the same above way

2. SSID

```
|SSID| = |Coherer|
```

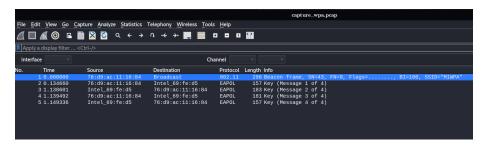


Figure 13: SSID

2.2.3 PTK

 $Nonce_A, Nonce_S, MAC_{Client}, MAC_{Authenticator}$

Based on analyzing the captured package using Wireshark, we got the following information:

2.2.4 MIC generation

Similar to the previous section, we capture using Wireshark:

```
MIC1 = a462a7029ad5ba30b6af0df391988e45

MIC2 = 7d0af6df51e99cde7a187453f0f93537

MIC3 = 10bba3bdfbcfde2bc537509d71f2ecd1
```

```
00 8e 58 00 00
                                 10 6c 6c 09 c0 00 64 00
      00 38
            00 00 8a 0b 2e f7
                                08 01 2c 00 00 0c 41 82
      b2 55 00 0d 93 82 36 3a
                                00 0c 41 82 b2 55 90 01
                                                              U
      aa aa 03 00 00 00 88 8e
9030
                                                    0a
0050
         3d
            4e ac
                      14 b1 d3
                                86 00 00 00 00 00
                                                    00 00
0060
                                                              9N · R ·
0070
        00 00 00 00 00 00 00
                                00 00 00 00 00
                                                00 00 00
0080
                                                    d5 ba
                                    a4
0090
                         98 8e
                                       16 30
                                                    00 00
            02 01 00 00
                                       00 00
                                             0f
                                                    02 00
90a0
00b0
      00 8a 0b 2e f7
```

Figure 14: First target MIC

0000	00	00	18	00	8e	58	00	00	10	6c	6c	09	с0	00	64	00	X	·ll···d·
0010	00	28	00	00	6c	39	91	0c	08	02	2c	00	00	0d	93	82	·(··l9··	,
0020	36	За	00	0с	41	82	b2	55	00	9с	41	82	b2	55	с0	fc	6: A U	· · À · · U · ·
0030	aa	aa	03	00	00	00	88	8e	02	03	00	af	02	13	ca	00		
0040	10	00	00	00	00	00	00	00	01	Зе	8e	96	7d	ac	d9	60		·>··}··`
0050	32	4c	ac	5b	6a	a7	21	23	5b	f5	7b	94	97	71	с8	67	2L · [j · !#	[· { · · q · g
0060	98	9f	49	d0	4e	d4	7c	69	33	f5	7b	94	97	71	с8	67	··I·N· i	3 · { · · q · g
0070	98	9f	49	d0	4e	d4	7c	69	34	cf	02	00	00	00	00	00	··I·N· i	4 · · · · · ·
0080	00	00	00	00	00	00	00	00	00	7d	0a	f6	df	51	е9	9c		· } · · · · Q · ·
0090	de	7a	18	74	53	f⊙	f9	35	37	00	50	cf	a7	2c	de	35	·z·tS··5	7·P··,·5
00a0	b2	c1	e2	31	92	55	80	6a	b3	64	17	9f	d9	67	30	41	···1·U·j	·d···g0A
00b0	b9	a5	93	9f	a1	a2	01	0d	2a	с7	94	e2	51	68	05	5f		* · · · Qh · _
00c0	79	4d	dc	1f	df	ae	35	21	f4	44	6b	fd	11	da	98	34	yM · · · · 5!	$\cdot Dk \cdot \cdots \cdot 4$
00d0	5f	54	3d	f6	се	19	9d	f8	fe	48	f8	cd	d1	7a	dc	a8	_T=	$\cdot H \cdot \cdot \cdot z \cdot \cdot$
00e0	7b	f4	57	11	18	3с	49	6d	41	aa	0c	6c	39	91	9c		$\{\cdot W \cdot \cdot < Im\}$	A··l9··

Figure 15: Second EAPOL

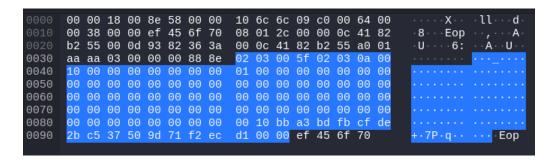


Figure 16: Third EAPOL

MIC Calculation

We will calculate the MIC from each PSK in the dictionary and EAPOL without MIC, then compare it with the target MIC, thereby finding the correct password.

 $MIC = HMAC_{SHA1}(KCK, EAPOL_{noMIC})$

2.2.5 Bruteforce Attacking

We got result:

```
python <u>test_solve.py</u>
                                                                                      [11:01AM ]
                    A288FCF0CAAACDA9A9F58633FF35E8992A01D9C10BA5E02EFDF8CB5D730CE7BC
pmk:
ptk:
                    B1CD792716762903F723424CD7D1651182A644133BFA4E0B75D96D230835843315798D
511BEAE0028313C8AB32F12C7ECB71C893482669DAAF0E9223FE1C0AED
desired mic:
actual mic:
MATCH
                    A462A7029AD5BA30B6AF0DF391988E45
                    A462A7029AD5BA30B6AF0DF391988E45
desired mic:
actual mic:
MATCH
                    7D0AF6DF51E99CDE7A187453F0F93537
                    7D0AF6DF51E99CDE7A187453F0F93537
                    10BBA3BDFBCFDE2BC537509D71F2ECD1
desired mic:
actual mic:
                    10BBA3BDFBCFDE2BC537509D71F2ECD1
MATCH
Password found:
                              Induction
    \langle CHANGEME 
angle\langle f \extstyle 2.168.81.136 <math>
angle\langle \sim /Desktop/AuditSubject/project 
angle
```

Figure 17: Final Result

As we can see, all of MICs were matched. Then we got PSK = Induction