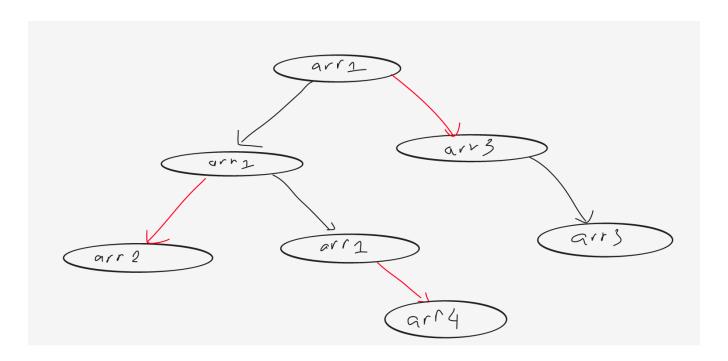


Abstract

You have to maintain a list of arrays. Initially, you have a single array $a_1, a_2, ..., a_n$ (array 1). You have to process three types of queries:

- 1. Change the value at index p of array k to x.
- 2. Return the sum in range [a, b] of array k.
- 3. Create a copy of array k and append as the last array.

Tutorial



Lets check out the image above which is a illustrated tree of *applying queries* in this problem. Where

- Each node illustrate a version of a particular array.
- Each direct edge denotes as a query (red for query 3 and black for the others).

Therefore, for each queries we have to create new version of array. So data structure persitent is what we need to solve this problem.

For more meanings, persistent is a data structure that always preverse version itself when it modified (in this problem, it's applied query).

So how should we apply persistent in this problem?

Let's define a array has a number of version. The array 1 will have version 1. So when applying a query, a new version is created. And we will update the version corresponding to the array. More details,

- If the query is 1 or 2, the last version is the version of the current array k.
- Else, the last version is the version of new array.

Conclusion, we have to store the latest version of each array and for a query we need to create new version from the version of corresponding considered array.

The pseudo code as follow:

```
cur_ver = 2 # version 1 corresponds to array 1
ver = []
latest_ver = []
for q in Q:
    array_k, update_info = q
    update(q, latest_ver[array_k], cur_ver)
    latest_ver[array_k] = cur_ver
    cur_ver += 1
```

The complexity of above alogrithm is O(np) which p the complexity of the updating. We can use a popular data structure Segment Tree to implement the update function in O(logn). But in this problem, it's a bit difficult that we have to combine 2 data structure above together.

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