

127 P(4, 4) and the receiver will receive the next probe packet and calculate the burstiness value.
 128 Assume that next time the probe packet is P(4, 7), so the burstiness value also 2, the burstiness time
 129 count will be increased by 1 value to describe that the burstiness value 2 is duplicated. Let the result
 130 of transmission per link at i^{th} sequence if successful is "S" and "F" if failed transmission. With the
 131 above example of node 1 and node 4, we have the result of the transmission after 7 probe packet is
 132 "SFFSFFS". After finish the measure link quality period (MLQ), we have a BDL of a link with 1000
 133 probe packets transmission in Table 1.

134 **Table 1.** Burstiness Distribution List (BDL) of a link after finish MLQ.

Burstiness value	0	1	2	3	4
Burstiness time count	634	129	31	2	1

135
 136 Let make the relationship between the burstiness value and burstiness time count in BDL. The
 137 burstiness value is defined as the number of consecutive loss during probe packets transmission,
 138 and the burstiness time count has defined the time burstiness value happen. For example, in Table 1,
 139 after 1000 probe packets transmission, burstiness value 0 happens 634 times, burstiness value 1
 140 happens 129 times, and so on. It means that in 1000 probe packets transmitted, 634 times happens
 141 the result of transmission is "SS", 129 times happen the result of transmission is "SFS", ...With the
 142 link has good quality, the burstiness value will be small while the burstiness time count will be large
 143 at min burstiness value. On the contrary, the bad link has a higher burstiness value and low
 144 burstiness time count at min burstiness value.

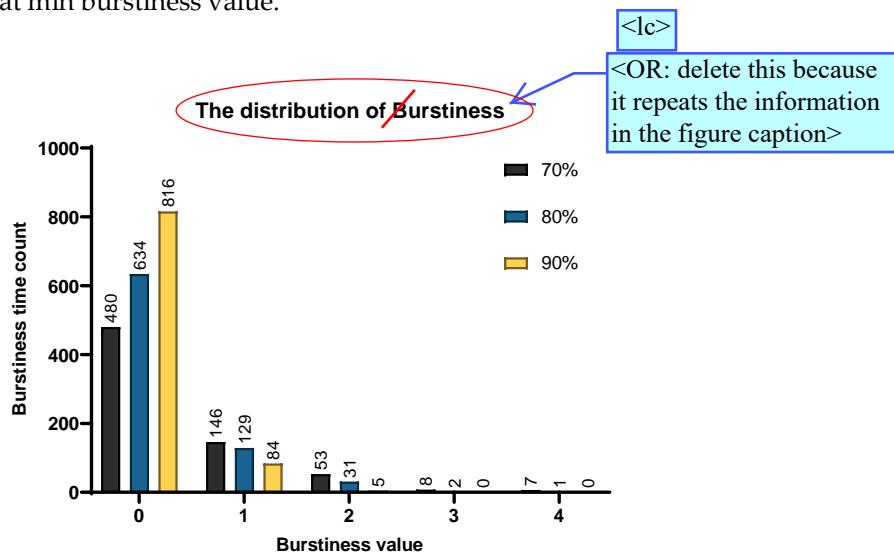


Figure 2. The distribution of Burstiness compares with link PRR.

145 For example, Figure 2 shows the simulation result of one pair of nodes with the different configured
 146 of the PRR value. The distribution of the burstiness value compare with link PRR 70%, 80%, and 90%
 147 with the number of probes is 1000 packets in Figure 2 describe that with the good link (PRR 90%) has
 148 the burstiness time count with min burstiness value is 816 and max burstiness value is 2. On the
 149 other hand, the bad link quality (PRR 70%) has the burstiness time count with a min burstiness value
 150 is 480 and the max burstiness value is 4.

151 3.3. Burstiness Distribution Metric

152 In this section, we present a scheme to calculate the number of retransmissions by Burstiness
 153 Distribution List (BDL) with named Burstiness Distribution Metric (Bdist). Firstly, the neighbor list

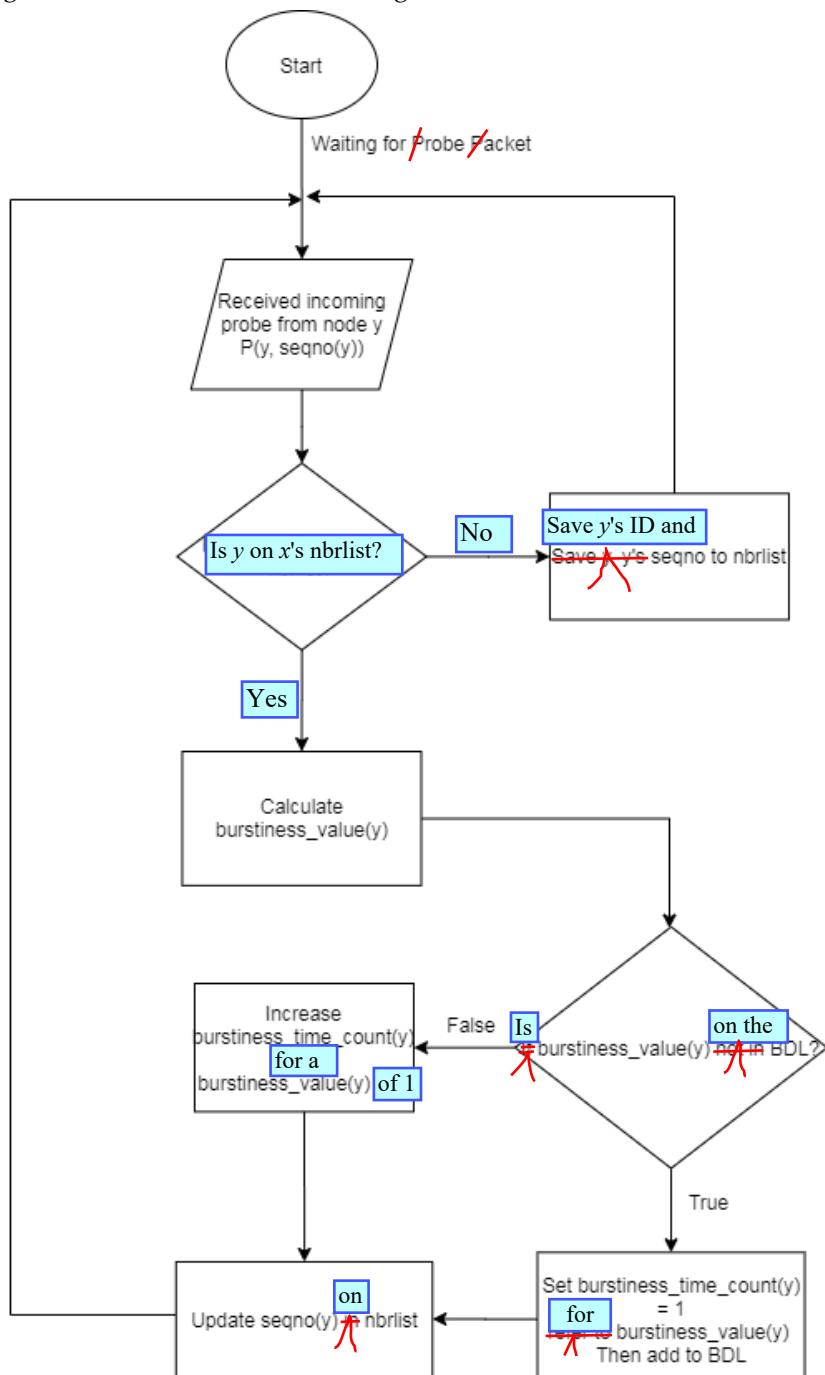
154 (nbrlist) and the BDL are set to empty. When the sensor node, called x , received a probe packet from
 155 node y , it saves node y and node y 's sequence to its nbrlist if node y does not belong to node x 's
 156 nbrlist. If node y has belonged in node x 's nbrlist, node x will calculate the burstiness value refer to
 157 node y as follows:

$$158 \quad \text{burstiness_value}(y) = \text{seqno}(y) - \text{nbrlist}[y].\text{seqno}(y) - 1$$

159 in which, $\text{seqno}(y)$ is the sequence number get from new incoming probe packet from node y ,
 160 $\text{nbrlist}[y].\text{seqno}(y)$ is the previous sequence number of node y saved in the neighbor list.

161 The burstiness_time_count(y) refer to burstiness_value(y) set to 1 and the pair
 162 (burstiness_value(y), burstiness_time_count(y)) will be saved in node y 's BDL if this value does not
 163 belong to node y 's BDL. If burstiness_value(y) is belong to node y 's BDL, node x will increase the
 164 burstiness_time_count(y) refer to burstiness_value(y) one. Then update the seqno(y) value in the
 165 neighbor list.

166 This process will repeat during the measure link quality period. The detailed algorithm is
 167 described in Algorithm 1 and the flowchart in Figure 3.



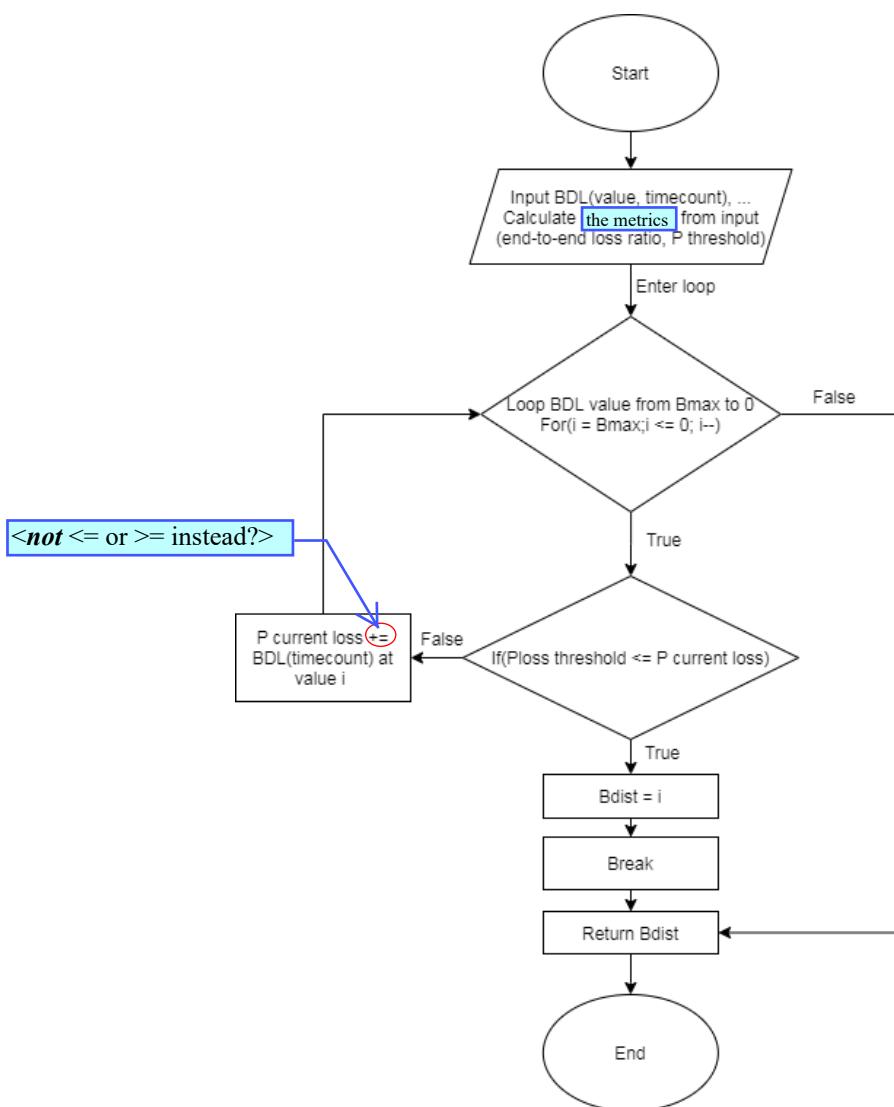


Figure 4. The flowchart of the main procedures of Algorithm 2.

185 4. Evaluation

186 By using Cooja Simulator, the Bdist is compared with some schemes such as ETX, PRR. The key
187 parameter is present in Table 2.

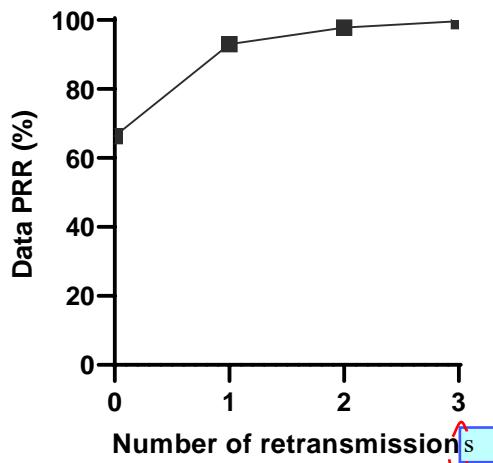
188 **Table 2.** Summary of the simulation set up parameters.

Number of nodes	10
Number of probes	1000
Target PRR	99%
Link quality of the channel (is set)	70% - 90%
Number of the data packets	1000
Network area	100m x 100m
Payload size	26 bytes
Simulation time for each scenario	About 55 minutes

189 In which the link quality of the channel of each link indicates the number percent of the
 190 successful packet at the receiver compared with the number packet of senders. For example, if the
 191 link quality of the channel equals 70 indicate that the receiver will receive 70 packets if the sender
 192 sends 100 packets. The number of probes to measure link quality and the number of data packets to
 193 evaluate the measure link quality metric are 1000 packets.

194 *4.1. Relationship of the number of retransmissions and network performance*

195 To evaluate the effect of the number of retransmissions on the performance of the network, a
 196 simple network with one sink (gateway) and nine sensor nodes are considered. The link quality of
 197 the channel will be fixed while the number of retransmissions will be varying from 0 to 3 times and
 198 we will examine the packet reception rate at the sink with each value of the number of
 199 retransmissions.



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Figure 5. Relationship of the number of retransmissions to packet reception rate.

200 The result in Figure 5 shows that since the sensor nodes are not allowed retransmissions, the
 201 PRR at the sink can achieve a low PRR, 66.3%. when the sensor nodes are allowed retransmission,
 202 the PRR will be improved and increased since the number of retransmissions increased. It is shown
 203 that since the link quality is not very good, the number of retransmissions will decide the network
 204 performance. Therefore, estimate the number of retransmissions is very important.

205 *4.2. Effect of the hop count on network performance*

206 We evaluate the effect of hop count on PRR. A simple linear network with the hop count
 207 varying from 1 to 4 is considered. In which we compare our proposed with some schemes such as
 208 ETX, PRR. Each scheme estimates the number of retransmissions by itself. We will examine the
 209 packet reception rate at the sink with each value of the hop count.

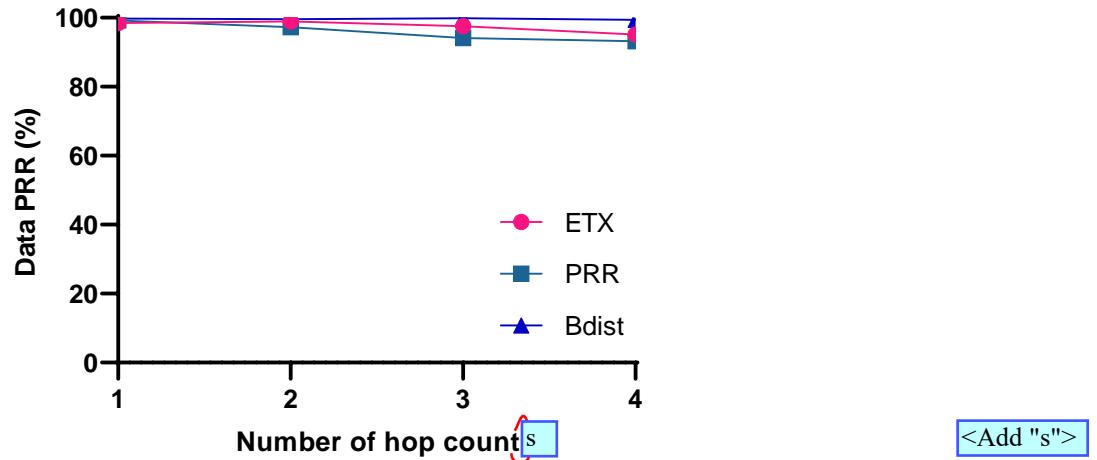


Figure 6. Effect of the hop count on network performance.

210 The result is present in Figure 6. From Figure 6, we can see that our proposed archives very
 211 high PRR even at the highest hop. It is because our proposed have a good estimation of the number
 212 of retransmissions by using the target PRR, estimate link quality method. While other schemes the
 213 PRR will be decreased since the hop count increased. This is because the scheme to estimate the
 214 number of retransmissions of two other schemes is not very good.

215 *4.3. Evaluate the network with some estimation scheme*

216 In this examination, nine sensor nodes are deployed randomly and the link quality of each
 217 between two nodes is set randomly from 70% to 90%. Each sensor node will generate 1000 packets
 218 periodically. The TSCH MAC protocol is used with the PCLLF [24] scheduling algorithm to
 219 guarantee that no collision in the network. We collect the packet reception rate of each sensor node at
 220 the sink.

221 The result of the data PRR of each sensor node at the sink is present in Figure 7. From Figure 7,
 222 the PRR of the proposed scheme almost archives targets PRR with the estimated times of
 223 retransmission for each sensor node. It shows that the algorithm of calculating the number of
 224 retransmission work well. Furthermore, the other archives much lower PRR since the estimation of
 225 the number of retransmissions of ETX and PRR is not very good. This is because ETX and PRR did
 226 not consider the burstiness happen in the link.

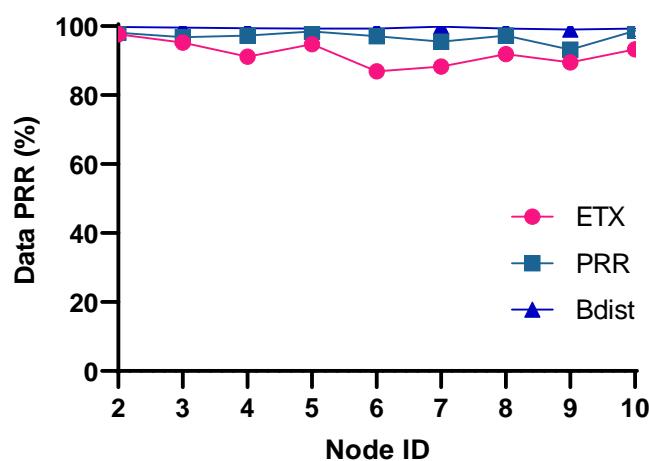


Figure 7. The data PRR of each sensor node at the sink of some estimation scheme.

227 *4.4. Evaluate the network with some type of network*