CSE 444: Database Internals

Lecture 12

Query Optimization (part 3)

Announcements

Selinger Optimizer History

- 1960's: first database systems
  - Use tree and graph data models
- 1970: Ted Codd proposes relational model
  - E.F. Codd. A relational model of data for large shared data banks. Communications of the ACM, 1970
- 1974: System R from IBM Research
  - One of first systems to implement relational model
- 1979: Seminal query optimizer paper by P. Selinger et. al.
  - Invented cost-based query optimization
  - Dynamic programming algorithm for join order computation

References

- P. Selinger, M. Astrahan, D. Chamberlin, R. Lorie, and T. Price. Access Path Selection in a Relational Database Management System. Proceedings of ACM SIGMOD, **1979**. Pages 22-34.

Selinger Algorithm

Selinger enumeration algorithm considers

- Different logical and physical plans *at the same time*

- Cost of a plan is IO + CPU

- Concept of *interesting order* during plan enumeration
  - A sorted order as that requested by ORDER BY or GROUP GY
  - Or order on attributes that appear in equi-join predicates
    - Because they may enable cheaper sort-merge joins later

More about the Selinger Algorithm

- Step 1: Enumerate all access paths for a single relation
  - File scan or index scan
  - Keep the cheapest for each *interesting order*
- Step 2: Consider all ways to join two relations
  - Use result from step 1 as the outer relation
  - Consider every other possible relation as inner relation
  - Estimate cost when using sort-merge or nested-loop join
  - Keep the cheapest for each *interesting order*
- Steps 3 and later: Repeat for three relations, etc.

Example From Selinger Paper

EMP

NAME	DNO	JOB	SAL
SMITH	50	12	8500
JONES	50	5	15000
DOE	51	5	9500

DEPT

DNO	DNAME	LOC
50	MFG	DENVER
51	BILLING	BOULDER
52	SHIPPING	DENVER

JOB

JOB	TITLE
5	CLERK
6	TYPIST
8	SALES
12	MECHANIC

```
SELECT   NAME,TITLE,SAL,DNAME
FROM     EMP,DEPT,JOB
WHERE    TITLE='CLERK'
AND      LOC='DENVER'
AND      EMP.DNO=DEPT.DNO
AND      EMP.JOB=JOB.JOB
```

“Retrieve the name, salary, job title, and department name of employees who are clerks and work for departments in Denver.”

Figure 1. JOIN example

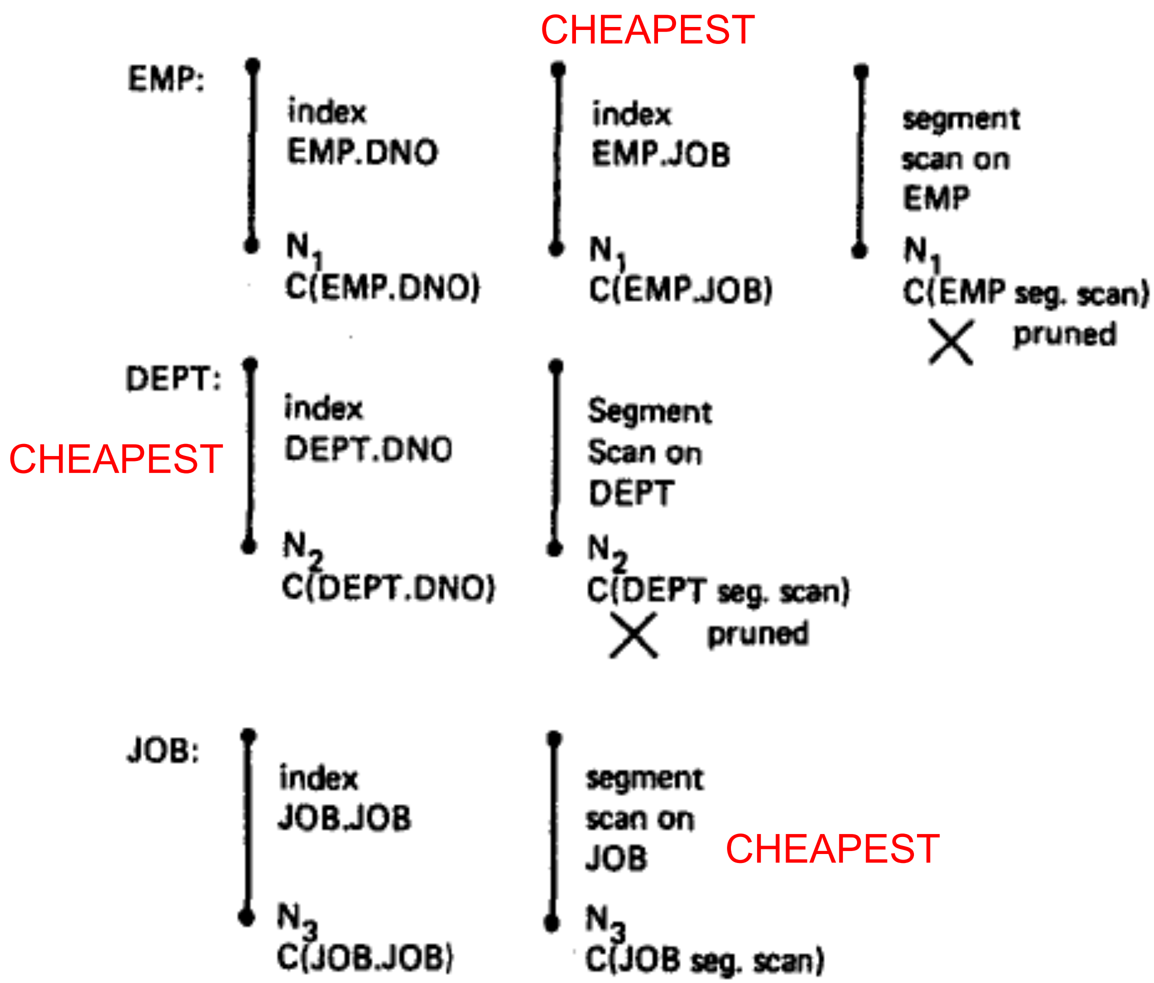
Step1: Access Path Selection for Single Relations

- Eligible Predicates: Local Predicates Only
- “Interesting” Orderings: DNO, JOB

```
SELECT   NAME,TITLE,SAL,DNAME
FROM     EMP,DEPT,JOB
WHERE    TITLE='CLERK'
AND      LOC='DENVER'
AND      EMP.DNO=DEPT.DNO
AND      EMP.JOB=JOB.JOB
```

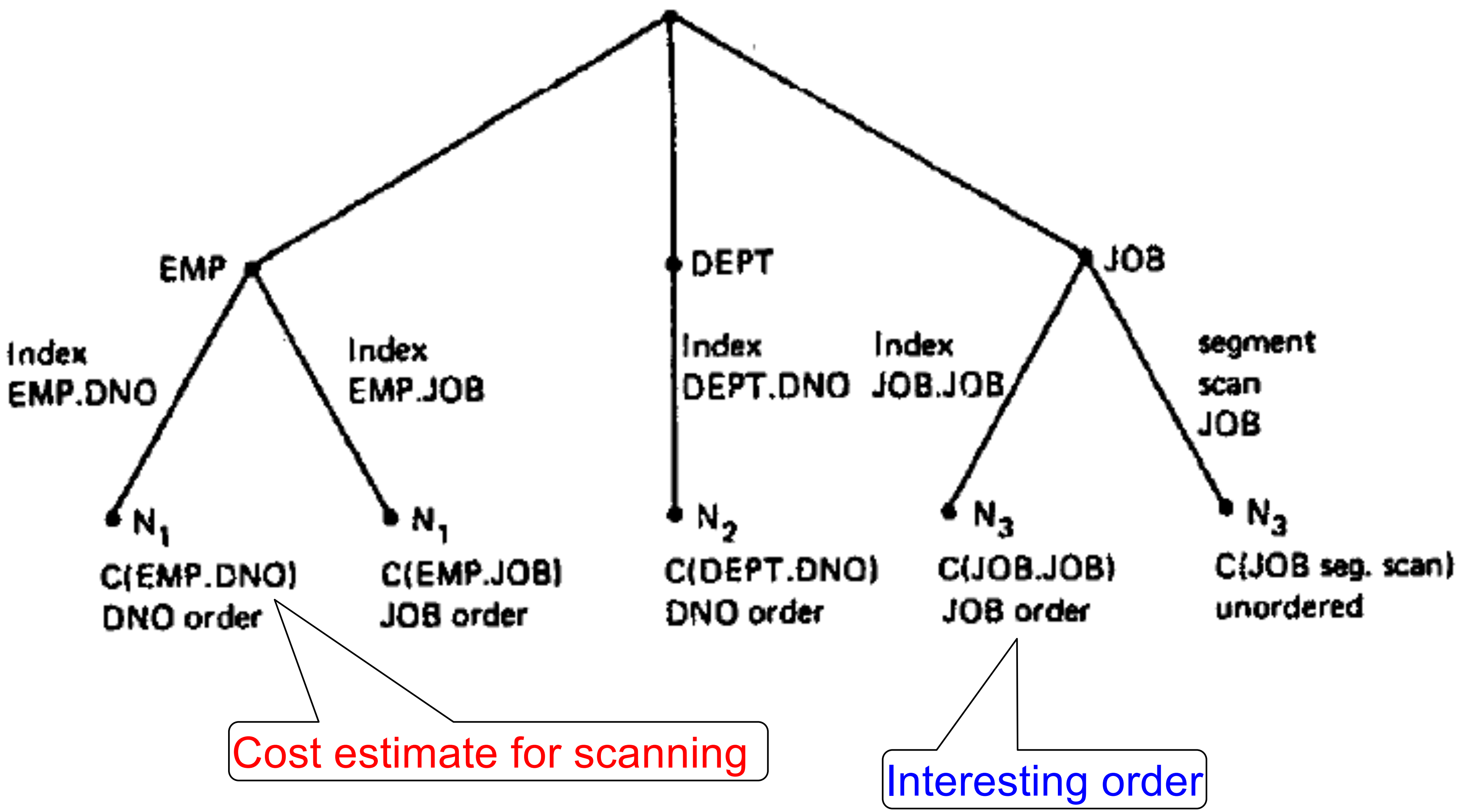
“Retrieve the name, salary, job title, and department name of employees who are clerks and work for departments in Denver.”

Figure 1. JOIN example



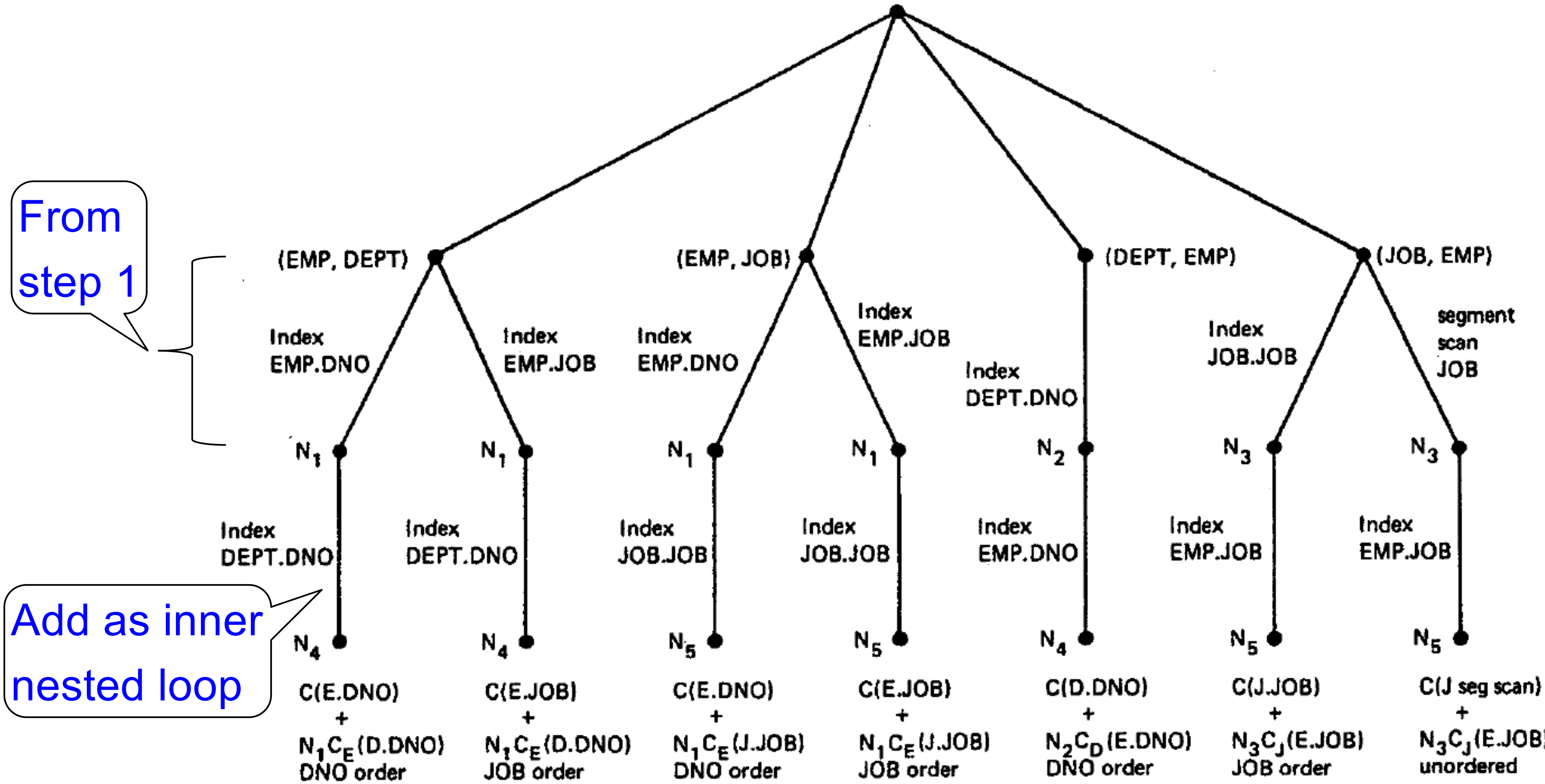
SELECT NAME, TITLE, SAL, DNAME
FROM EMP, DEPT, JOB
WHERE TITLE=‘CLERK’ AND LOC=‘DENVER’ AND EMP.DNO=DEPT.DNO AND EMP.JOB=JOB.JOB

Step1: Access Path Selection for Single Relations
Resulting Plan Search Tree for Single Relations



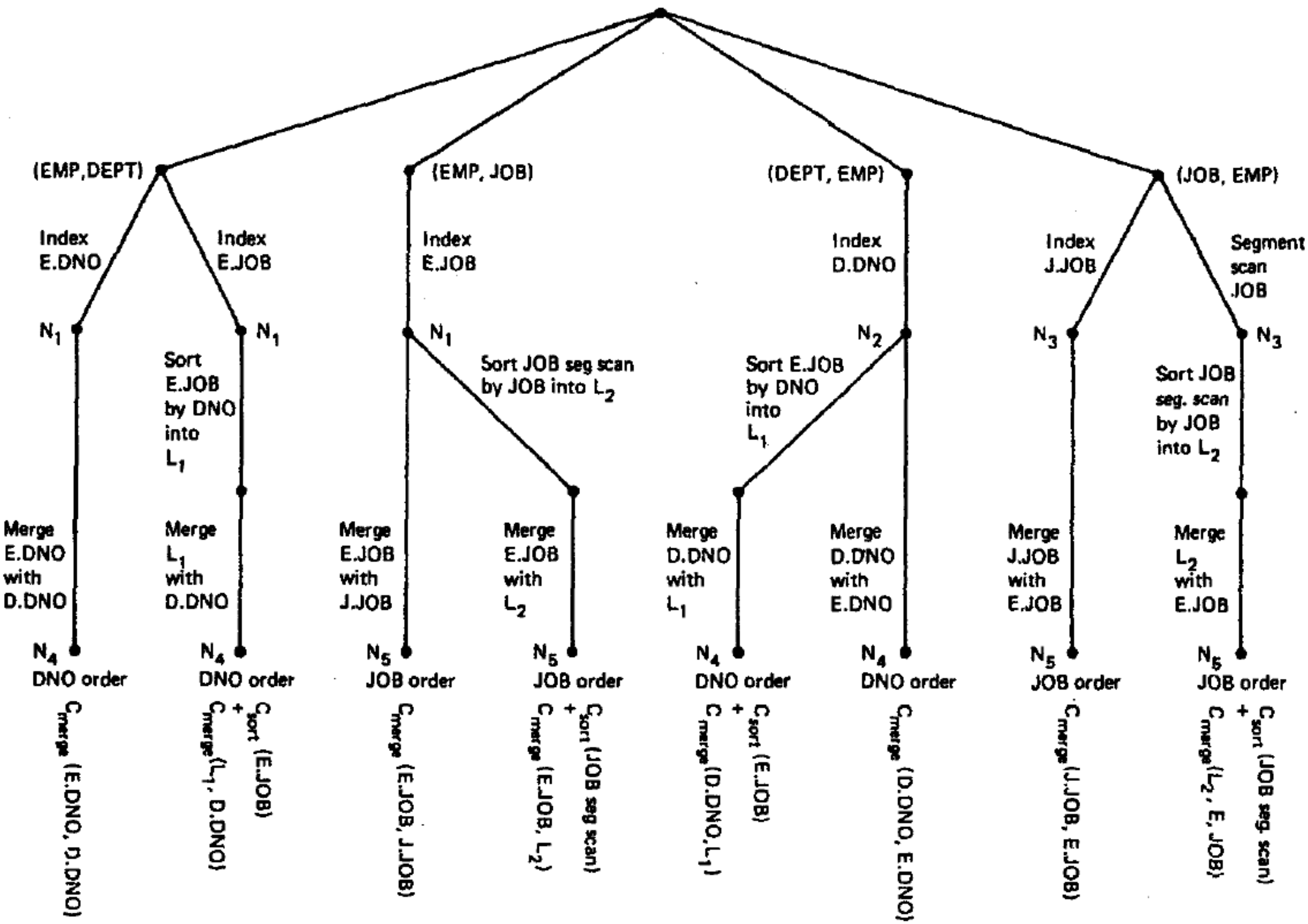
SELECT NAME, TITLE, SAL, DNAME
FROM EMP, DEPT, JOB
WHERE TITLE='CLERK' AND LOC='DENVER' AND EMP.DNO=DEPT.DNO AND EMP.JOB=JOB.JOB

Step2: Pairs of Relations (nested loop joins)

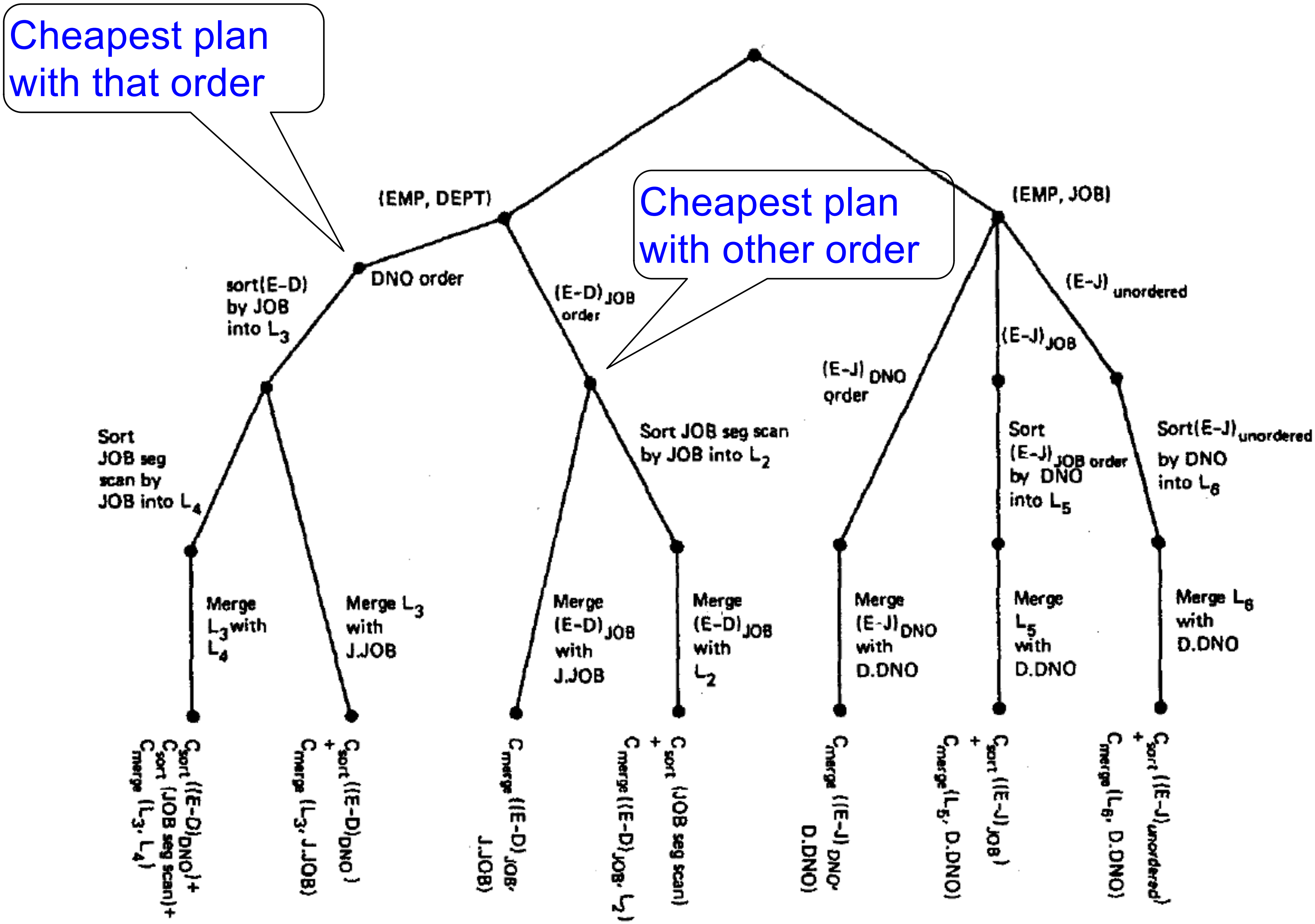


SELECT NAME, TITLE, SAL, DNAME
FROM EMP, DEPT, JOB
WHERE TITLE='CLERK' AND LOC='DENVER' AND EMP.DNO=DEPT.DNO AND EMP.JOB=JOB.JOB

Step2: Pairs of Relations (sort-merge joins)



Step3:Add Third Relation (sort-merge join)



Next Example Acks

Implement variant of Selinger optimizer in SimpleDB

Designed to help you with Lab 5

Many following slides from Sam Madden at MIT

Dynamic Programming

OrderJoins(…):

SimpleDB Lab5:
you implement orderJoins

R = set of relations to join

For d = 1 to N: /* where N = |R| */

For S in {all size-d subsets of R}:

Use: enumerateSubsets

optjoin(S) = (S – a) join a,

where a is the single relation that minimizes:

cost(**optjoin**(S – a)) +

Use:
computeCostAndCardOfSubplan

min.cost to join (S – a) with a +

min.access cost for a

Note: **optjoin**(S-a) is cached from previous iterations

Example

Subplan S	optJoin(S)	Cost(OptJoin(S))
A		

- **orderJoins(A, B, C, D)**
- Assume all joins are Nested Loop

Example

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
C	Seq scan	120
D	B+tree scan	400

- **orderJoins(A, B, C, D)**
- Assume all joins are NL
- d = 1
  - A = best way to access A
    (sequential scan, predicate-pushdown on index, etc)
  - B = best way to access B
  - C = best way to access C
  - D = best way to access D
- Total number of steps: choose(N, 1)

Example

- **orderJoins(A, B, C, D)**

- d = 2
  - {A,B} = AB or BA
    use previously computed
    best way to access A and B

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
…		

Example

- **orderJoins(A, B, C, D)**

- d = 2
  - {A,B} = AB or BA
    use previously computed
    best way to access A and B

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
…		
{A, B}	BA	156

Example

- **orderJoins(A, B, C, D)**

- d = 2
  - {A,B} = AB or BA
    use previously computed
    best way to access A and B
  - {B,C} = BC or CB

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
…		
{A, B}	BA	156
{B, C}	BC	98

Example

- **orderJoins(A, B, C, D)**

- d = 2
  - {A,B} = AB or BA
    use previously computed
    best way to access A and B
  - {B,C} = BC or CB

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
…		
{A, B}	BA	156
{B, C}	BC	98

Example

- **orderJoins(A, B, C, D)**

- d = 2
  - {A,B} = AB or BA
    use previously computed
    best way to access A and B
  - $\boxed{\text{\{B,C\}}}$ = BC or CB
  - {C,D} = CD or DC
  - {A,C} = AC or CA
  - {B,D} = BD or DB
  - {A,D} = AD or DA

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
…		
{A, B}	BA	156
$\boxed{\text{\{B, C\}}}$	BC	98
……..		

Example

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
…		
{A, B}	BA	156
{B, C}	BC	98
……..		

- **orderJoins(A, B, C, D)**
- d = 2
  - {A,B} = AB or BA
    use previously computed
    best way to access A and B
  - {B,C} = BC or CB
  - {C,D} = CD or DC
  - {A,C} = AC or CA
  - {B,D} = BD or DB
  - {A,D} = AD or DA
- Total number of steps: choose(N, 2) × 2

Example

- **orderJoins(A, B, C, D)**

- d = 3

  – {A,B,C} =
  Remove A: compare A({B,C}) to ({B,C})A

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
….		
{A, B}	BA	156
{B, C}	BC	98
….		

Example

- **orderJoins(A, B, C, D)**

- d = 3

  – {A,B,C} =
  Remove A: compare A($\boxed{\{B,C\}}$) to ({B,C})A

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
….		
{A, B}	BA	156
{B, C}	BC	98
….		

optJoin(B,C)
and its cost are
already cached
in table

Example

- **orderJoins(A, B, C, D)**
- d = 3
  - {A,B,C} =
    Remove A: compare A({B,C}) to ({B,C})A
    Remove B: compare B({A,C}) to ({A,C})B
    Remove C: compare C({A,B}) to ({A,B})C

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
….		
{A, B}	BA	156
{B, C}	BC	98
….		
{A, B, C}	BAC	500
……..		

optJoin(B,C) and its cost are already cached in table

Example

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
….		
{A, B}	BA	156
{B, C}	BC	98
…		
{A, B, C}	BAC	500
……..		

- **orderJoins(A, B, C, D)**
- d = 3
  - {A,B,C} =
    Remove A: compare A({B,C}) to ({B,C})A
    Remove B: compare B({A,C}) to ({A,C})B
    Remove C: compare C({A,B}) to ({A,B})C

optJoin(B,C) and its cost are already cached in table

Example

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
….		
{A, B}	BA	156
{B, C}	BC	98
…		
{A, B, C}	BAC	500
……..		

- **orderJoins(A, B, C, D)**
- d = 3
  - {A,B,C} =
    Remove A: compare A({B,C}) to ({B,C})A
    Remove B: compare B({A,C}) to ({A,C})B
    Remove C: compare C({A,B}) to ({A,B})C
  - {A,B,D} =
    Remove A: compare A({B,D}) to ({B,D})A
    …
  - {A,C,D} =…
  - {B,C,D} =…
- Total number of steps: choose(N, 3) × 3 × 2

optJoin(B,C) and its cost are already cached in table

Example

- **orderJoins(A, B, C, D)**

Subplan S	optJoin(S)	Cost(OptJoin(S))
A	Index scan	100
B	Seq. scan	50
{A, B}	BA	156
{B, C}	BC	98
{A, B, C}	BAC	500
{B, C, D}	DBC	150
……..		

- d = 4
  - {A,B,C,D} =

Remove A: compare A({B,C,D}) to ({B,C,D})A
Remove B: compare B({A,C,D}) to ({A,C,D})B
Remove C: compare C({A,B,D}) to ({A,B,D})C
Remove D: compare D({A,B,C}) to ({A,B,C})D

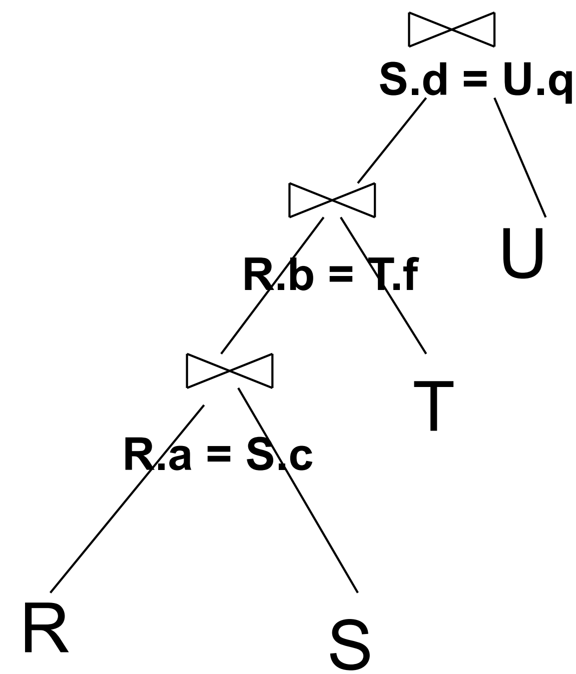
optJoin(B, C, D) and its cost are already cached in table

- Total number of steps: choose(N, 4) × 4 × 2

Implementation in SimpleDB (lab5)

1. JoinOptimizer.java (and the classes used there)

2. Returns vector of "LogicalJoinNode"
   Two base tables, two join attributes, predicate
   e.g. R(a, b), S(c, d), T(a, f), U(p, q)
   (R, S, R.a, S.c, =)
   Recall that SimpleDB keeps all attributes of
   R, S after their join R.a, R.b, S.c, S.d



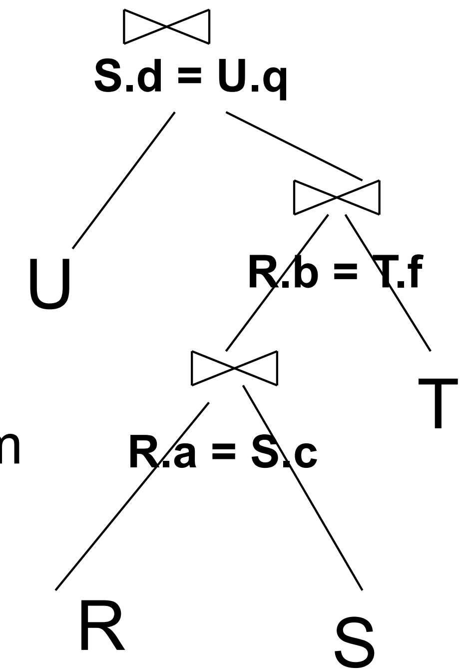
3. Output vector looks like:
   <(R, S, R.a, S.c), (R, T, R.b, T.f), (S, U, S.d, U.q)>

Implementation in SimpleDB (lab5)

Any advantage of returning pairs?

- Flexibility to consider all linear plans

  <(R, S, R.a,S.c), (R, T, R.b, T.f), (U, S, U.q, S.d)>



More Details:

1. You mainly need to implement "orderJoins(..)"
2. "CostCard" data structure stores a plan, its cost and cardinality: you would need to estimate them
3. "PlanCache" stores the table in dyn. Prog:

   Maps a <u>set</u> of LJN to

   a <u>vector</u> of LJN (best plan for the vector),

   its cost, and its cardinality

   LJN = LogicalJoinNode