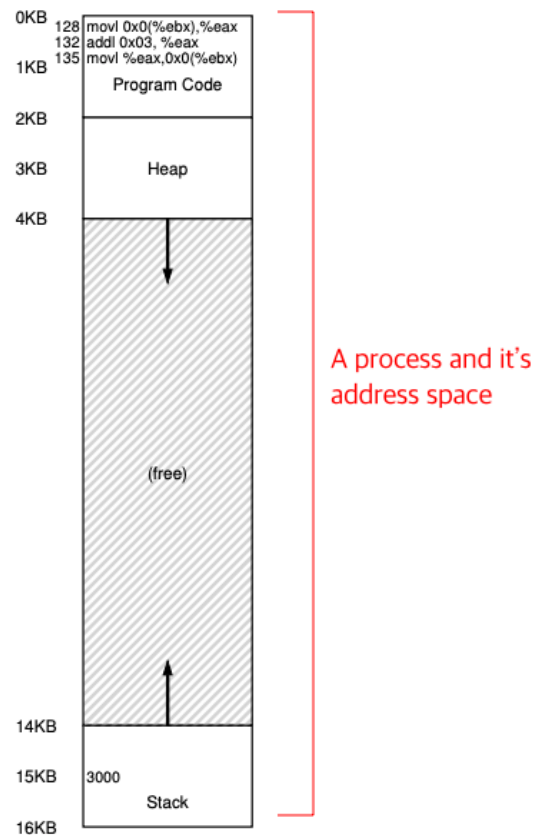


1 Concurrency

Vocabularies

1. Address Space

- Is a range of discrete addresses where each corresponds to a memory cell

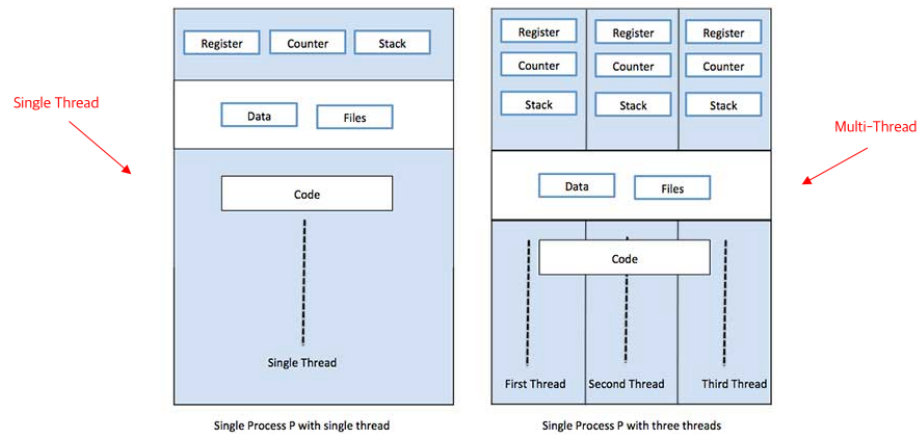


2. Thread

- is the unit of execution within a process.
 - A process can have anywhere from one thread to many threads

3. Multi-thread

- Is a process with more than one thread



4. Context Switch

- is the process of storing the state of a process or thread, so that it can be restored and resume execution at a later point

5. Process Control Block

- Is also called **process descriptor**
- Is a data structure used by computer operating systems to store all the information about a process

```
// the registers xv6 will save and restore
// to stop and subsequently restart a process
struct context {
    int eip;
    int esp;
    int ebx;
    int ecx;
    int edx;
    int esi;
    int edi;
    int ebp;
};

// the different states a process can be in
enum proc_state { UNUSED, EMBRYO, SLEEPING,
    RUNNABLE, RUNNING, ZOMBIE };

// the information xv6 tracks about each process
// including its register context and state
struct proc {
    char *mem;           // Start of process memory
    uint sz;             // Size of process memory
    char *kstack;        // Bottom of kernel stack
                        // for this process
    enum proc_state state; // Process state
    int pid;             // Process ID
    struct proc *parent; // Parent process
    void *chan;          // If !zero, sleeping on chan
    int killed;          // If !zero, has been killed
    struct file *ofile[NOFILE]; // Open files
    struct inode *cwd;    // Current directory
    struct context ctx;   // Switch here to run process
    struct trapframe *tf; // Trap frame for the
                        // current interrupt
};
```

register context data structure

Process Control Block

Where register context is

6. Thread Control Block

- Is also called **process descriptor**
- Is a data structure used by computer operating systems to store all the information about a thread

7. Parallelism

- Is a condition that arises when at least two threads are executing simultaneously.

8. Parallelization

- Is the task of transforming a standard **single-threaded** program into a program that does this (sort of) work on multiple CPUs

9. Multiprogramming

- Is a computer running more than one program at a time (like running Excel and Firefox simultaneously).

10. Critical Section

- Is a piece of code that accesses a shared resource, usually a variable or data structure

11. Race Condition (Data race)

- Is a condition that arises if multiple threads of execution enter the critical section at roughly the same time; both attempt to update the shared data structure, leading to a surprising outcome

12. Mutual Exclusion

- Is a concurrency control property which is introduced to prevent race conditions, resulting in deterministic program output.

13. Deterministic Execution

- Means path of execution is fully determined by the specification of computation
- Is guaranteed to produce the same outcome, given the same input

14. Indeterminate

- Means path of execution isn't fully determined by the specification of computation
- Same input can produce different outcomes

15. Synchronization Primitives

- Are simple software mechanisms provided by a platform (e.g. operating system) to its users for the purposes of supporting thread or process synchronization.

Example

Mutex, event, conditional variables and semaphores

16. Atomic

- Means “All or nothing”; no in-between state

17. Transaction

- is the grouping of many actions into a single atomic action

18. Conditional Variable

- Is a synchronization primitives that enable threads to wait until a particular condition occurs

1.1 Why Use Threads?

• Parallelism

- Potential of speeding up a process by using multiple processors
 - * by making 1 CPU → perform a portion of work
- **Parallelization** converts **single-threaded** programs to (sort of) **multi-threaded** program

• Avoid blocking program due to slow I/O

- Switch to another thread while a thread waits (in blocked state) for I/O
- **Thread** is a natural way of getting unstuck
 - * Enables overlap of I/O within a single program

Example

- * Typing on Messenger (thread A) while (thread B) is sending 'sent' message to user
- * Typing on Visual Studio Code (thread A), while (thread B) is saving data to hard drive

1.2 Thread Creation

- Created using function `pthread_create`
- Created thread is in one of two states

1. Ready (wait until something happens)
2. Running (run immediately)

1.3 Why It Gets Worse: Shared Data

- Thread doesn't produce **deterministic** result

Thread 1: Run loop 10,000,000 times
 Thread 2: Run loop 10,000,000 times
 Goal: Return total count shared by both Thread 1 and Thread 2

```

prompt> gcc -o main main.c -Wall -pthread; ./main
main: begin (counter = 0)
A: begin
B: begin
A: done
B: done
main: done with both (counter = 20000000)
  
```

Expected Result: 20,000,000

```

prompt> ./main
main: begin (counter = 0)
A: begin
B: begin
A: done
B: done
main: done with both (counter = 19345221)
  
```

Test 1
 Result: 19,345,221 (오잉?)

```

prompt> ./main
main: begin (counter = 0)
A: begin
B: begin
A: done
B: done
main: done with both (counter = 19221041)
  
```

Test 2
 Result: 19,221,041
 (Result has different value than test 1)

1.4 The Heart of The Problem: Uncontrolled Scheduling

- Race Condition (Data race)
 - Two threads accessed the shared variable at the same time
 - Should increase by 2, but only 1 occurred
- Shared variable is a **critical section**
- **Mutual Exclusion** is needed

1.5 The Wish For Atomicity

- **Mutual Exclusion** guarantees **atomicity** in **critical section**
- Is done using **synchronization primitives**
 - Allows multi-threaded code to access **critical section** in a synchronized and controlled manner and produce correct result

Examples

Mutex, event, conditional variables and semaphores

1.6 One More Problem: Waiting For Another

- There are two common problems in concurrency
 1. Accessing a shared variable
 2. One thread waiting for another to complete some action before it can continue
 - Is solved by **conditional variable**

Example

Put a process to sleep when performing disk I/O; when I/O is complete, wake up from sleep and continue