

# Assignment 2 Notes

Hyungmo Gu

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## 1 Read and Writer

- Reader don't modify the data so we can have multiple readers, but only one writer
- Are examples of a common computing problem in concurrency
- Is a part of semaphore problem

## 2 Product/Consumer

```
producer() {                consumer() {
    while(1) {                while(1) {

        add_to_buf()          remove_from_buf()

    }                          }
}
```

- Is essentially how `pipes()` are implemented
- Has bounded buffer as a shared variable
  - Bounded buffer is also used when piping the output of one program into another

### Example

```
grep foo file.txt | wc -l
```

\* `grep`

- searches the input files for lines containing a match to a given pattern
- when it finds a match in a line, it copies the line to standard output (by default)

\* `wc -l`

- stands for word count
- is used to find the number of lines (in this case)

\* `grep` is the producer

\* `wc` is the consumer

- Single buffer producer/consumer solution

- Is to use two different conditinal variables
- \* Is nice, trouble free and simple

```

1  cond_t  empty, fill;
2  mutex_t mutex;
3
4  void *producer(void *arg) {
5      int i;
6      for (i = 0; i < loops; i++) {
7          Pthread_mutex_lock(&mutex);
8          while (count == 1)
9              Pthread_cond_wait(&empty, &mutex);
10         put(i);
11         Pthread_cond_signal(&fill);
12         Pthread_mutex_unlock(&mutex);
13     }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);
20         while (count == 0)
21             Pthread_cond_wait(&fill, &mutex);
22         int tmp = get();
23         Pthread_cond_signal(&empty);
24         Pthread_mutex_unlock(&mutex);
25         printf("%d\n", tmp);
26     }
27 }

```

Conditional variable 1

Conditional variable 2

YES

```

1  int loops;
2  cond_t  cond;
3  mutex_t mutex;
4
5  void *producer(void *arg) {
6      int i;
7      for (i = 0; i < loops; i++) {
8          Pthread_mutex_lock(&mutex);           // p1
9          while (count == 1)                   // p2
10             Pthread_cond_wait(&cond, &mutex); // p3
11         put(i);                               // p4
12         Pthread_cond_signal(&cond);           // p5
13         Pthread_mutex_unlock(&mutex);         // p6
14     }
15 }
16
17 void *consumer(void *arg) {
18     int i;
19     for (i = 0; i < loops; i++) {
20         Pthread_mutex_lock(&mutex);           // c1
21         while (count == 0)                   // c2
22             Pthread_cond_wait(&cond, &mutex); // c3
23         int tmp = get();                     // c4
24         Pthread_cond_signal(&cond);           // c5
25         Pthread_mutex_unlock(&mutex);         // c6
26         printf("%d\n", tmp);
27     }
28 }

```

Same conditional variable

NONO

- The general correct producer/consumer solution

```

1  int buffer[MAX];
2  int fill_ptr = 0;
3  int use_ptr = 0;
4  int count = 0;
5
6  void put(int value) {
7      buffer[fill_ptr] = value;
8      fill_ptr = (fill_ptr + 1) % MAX;
9      count++;
10 }
11
12 int get() {
13     int tmp = buffer[use_ptr];
14     use_ptr = (use_ptr + 1) % MAX;
15     count--;
16     return tmp;
17 }
18
19 ~
20 cond_t empty, fill;
21 mutex_t mutex;
22
23 void *producer(void *arg) {
24     int i;
25     for (i = 0; i < loops; i++) {
26         Pthread_mutex_lock(&mutex);           // p1
27         while (count == MAX)                  // p2
28             Pthread_cond_wait(&empty, &mutex); // p3
29         put(i);                               // p4
30         Pthread_cond_signal(&fill);           // p5
31         Pthread_mutex_unlock(&mutex);         // p6
32     }
33 }
34
35 void *consumer(void *arg) {
36     int i;
37     for (i = 0; i < loops; i++) {
38         Pthread_mutex_lock(&mutex);           // c1
39         while (count == 0)                    // c2
40             Pthread_cond_wait(&fill, &mutex); // c3
41         int tmp = get();                      // c4
42         Pthread_cond_signal(&empty);          // c5
43         Pthread_mutex_unlock(&mutex);         // c6
44         printf("%d\n", tmp);
45     }
46 }

```

Changes made to make solution general

### 3 Conditional Variable

```

lock_acquire(lock);
while(condition not true) {
    cv_wait(cond, lock);
}
... // do stuff
cv_signal(cond); //or cv_broadcast(cond)
lock_release(lock);

```

Condition

Conditional variable

- is a queue of waiting threads

- has two operations associated with it:
  1. `cv_wait(struct cv *cv, struct lock *lock)`
    - Is executed when a thread wishes to put itself to sleep
    - Releases lock, waits, re-acquires lock before return
      - \* Is to prevent race conditions from occurring when a thread is trying to put itself to sleep
  2. `cv_signal(struct cv *cv, struct lock *lock)`
    - Wakes one enqueued thread
  3. `cv_broadcast(struct cv *cv, struct lock *lock)` [from notes]
    - Wakes all enqueued threads
- If no one is waiting, signal or broadcast has no effect
- has rules
  - always use with while loops
    - \* on waking up, thread checks for condition in while loop
    - \* if condition is true, then thread goes back to sleep
- is always used together with locks

## 4 Semaphore

```
Wait(Sem){
    while(Sem <= 0) ;
    Sem--;
}
```

```
Signal(Sem){
    Sem++;
}
```

- Was first invented by Dijkstra
- Is abstract data types that provide synchronization
- Has two atomic operations
  - wait
    - \* Is also called P or decrement
  - signal
    - \* increments the variable, unblock a waiting thread if there are any
- Has two types

1. Binary semaphore (count = 0 or 1)
  - Has single access to a resource
  - Mutual exclusion to a critical section
2. Counting semaphore
  - Resource with many units available
  - Is resource that allows certain kinds of unsynchronized concurrent access (e.g. reading)
  - Mutex has count = 1, counting has count = N