# 1 File System Operations

## 1.1 Directory Entries

- Has user-readable name and i-number
  - Map files to inodes

**Example**

```
("file-name", i-number)
```

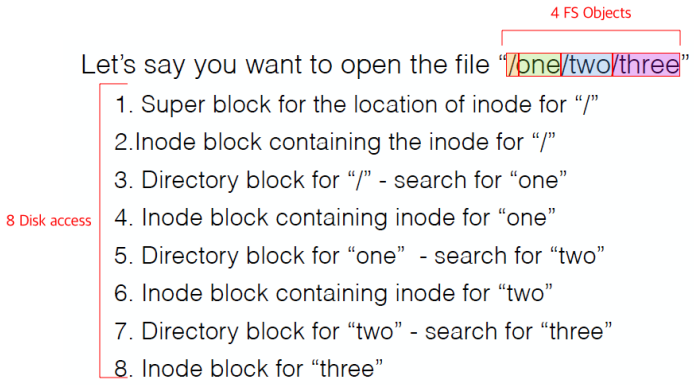
- Linux's ext2 has following attributes in its inode
  - Inode number
  - Directory entry length
  - Name length
  - File type
  - Name

## 1.2 How do we follow a path?

- Example
  - Read the first block of "/somefile"
    - The inode for the root directory "/" is known
    - Read the inode for "/" to file the data block(s) that hold the directory entires for the root directory
    - Read the first data block for the directory
    - Search the entries for "`somefile`"
    - The entry for "somefile" identifies the inode
    - Read the inode for "`somefile`"
    - The inode tells us the location of the first data block for the file
    - Read that block into memory access the data at the beginning of file

## 1.3 How many disks blocks are accessed?

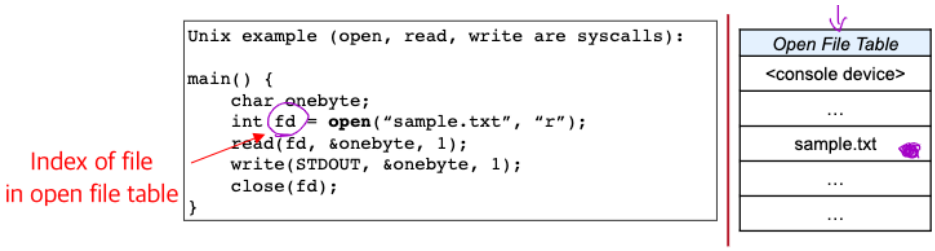
- Example




  - There are 8 disk accesses in total
  - There are 4 file system objects

## 1.4 Handling File Operations

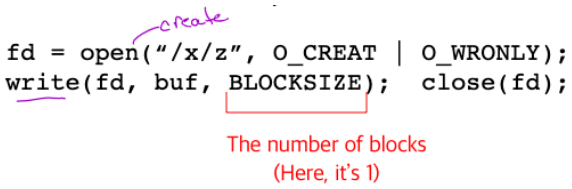
- Open file table
  - On open, file name, and it's inode number is stored system-wide into array
  - Index to entry in open-file table is used in subsequent use, so no searching is required
  - This is what `fd` is for




## 1.5 Possible Operations

- `mkdir("/x")`
  - Create a directory
- `creat("/x/y")`
  - Create an empty file
- `unlink("/x/y")`
  - Remove file or directory

- ```
  fd = open("/x/z", O_CREAT | O_WRONLY);
  write (fd, buf, BLOCKSIZE);
  close(fd);
  ```
  - Open a file for writing, and write one block to it
  - If it does not exist, create it


  ```