Assignment 2 Notes

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1 Read and Writer

- Reader don't modify the data so we can have multiple readers, but only one writer
- Are examples of a common computing problem in concurrency
- Is a part of semaphore problem

2 Product/Consumer

```
producer() {                consumer() {
  while(1) {                  while(1) {

    add_to_buf()                remove_from_buf()

  }                           }
}                           }
```

- Is essentially how `pipes()` are implemented
- Has bounded buffer as a shared variable
  - Bounded buffer is also used when piping the output of one program into another

    Example
    `grep foo file.txt | wc -l`
    * `grep`
      · searches the input files for lines containing a match to a given pattern lis
      · when it finds a match in a line, it copies the line to standard output (by default)
    * `wc -l`
      · stands for word count
      · is used to find the number of lines (in this case)
    * `grep` is the producer
    * `wc` is the consumer
- Single buffer producer/consumer solution

- Is to use two different conditial variables
  - Is nice, trouble free and simple

```
cond_t  empty, fill;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);
        while (count == 1)
            Pthread_cond_wait(&empty, &mutex);
        put(i);
        Pthread_cond_signal(&fill);
        Pthread_mutex_unlock(&mutex);
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);
        while (count == 0)
            Pthread_cond_wait(&fill, &mutex);
        int tmp = get();
        Pthread_cond_signal(&empty);
        Pthread_mutex_unlock(&mutex);
        printf("%d\n", tmp);
    }
}
```

Conditional variable 1

Conditional variable 2

YES

```
int loops;
cond_t  cond;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                    // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                               // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}
```

Same conditional variable

NONO

- The general correct producer/consumer solution

```
int buffer[MAX];
int fill_ptr = 0;
int use_ptr  = 0;
int count    = 0;

void put(int value) {
    buffer[fill_ptr] = value;
    fill_ptr = (fill_ptr + 1) % MAX;
    count++;
}

int get() {
    int tmp = buffer[use_ptr];
    use_ptr = (use_ptr + 1) % MAX;
    count--;
    return tmp;
}
```

Changes made to make solution general

```
cond_t empty, fill;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);             // p1
        while (count == MAX)                    // p2
            Pthread_cond_wait(&empty, &mutex);  // p3
        put(i);                                 // p4
        Pthread_cond_signal(&fill);             // p5
        Pthread_mutex_unlock(&mutex);           // p6
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);             // c1
        while (count == 0)                      // c2
            Pthread_cond_wait(&fill, &mutex);   // c3
        int tmp = get();                        // c4
        Pthread_cond_signal(&empty);            // c5
        Pthread_mutex_unlock(&mutex);           // c6
        printf("%d\n", tmp);
    }
}
```

3 Condtional Variable

```
lock_acquire(lock);
while(condition not true) {
    cv_wait(cond, lock);
}
... // do stuff

cv_signal(cond); //or cv_broadcast(cond)
lock_release(lock);
```

Condition

Conditional variable

- is a queue of waiting threads

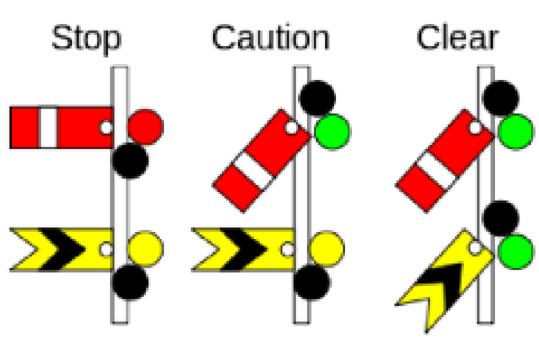
- has two operations associated with it:
  1. `cv_wait(struct cv *cv, struct lock *lock)`
     - Is executed when a thread wishes to put itself to sleep
     - Releases lock, waits, re-acquires lock before return
       * Is to prevent race condtions from occuring when a thread is trying to put itself to sleep
  2. `cv_signal(struct cv *cv, struct lock *lock)`
     - Wakes one enqueued thread
  3. `cv_broadcast(struct cv *cv, struct lock *lock)` [from notes]
     - Wakes all enqueued threads
- If no one is waiting, signal or broadcast has no effect
- has rules
  - always use with `while` loops
    * on waking up, tread checks for condition in while loop
    * if condition is true, then thread goes back to sleep
- is always used together with locks

4 Semaphore

```
Wait(Sem){
  while(Sem <= 0) ;
  Sem--;
}
```

```
Signal(Sem){
   Sem++;
}
```

- Was first invented by Dijkstra
- Is abstract data types that provide synchronization



- Has two <u>atomic</u> operations
  - `wait`
    - Is also called P or decrement
    - waits if value of count is negative
  - `signal`
    - increments the variable, unblock a waiting thread if there are any
    - wakes one thread if there are one or more threads waiting
- Has two types
  1. Binary semaphore (count = 0 or 1)
     - Has single access to a resource
     - Can be used like a lock
       - Provides mutual exclusion to a critical section
     - Needs to have initial value of 1
       - Is for decrementation of `count` in `sem_wait()`

Need to initialize semaphore to 1

```
sem_t m;
sem_init(&m, 0, X); // initialize to X; what should X be?

sem_wait(&m);
// critical section here
sem_post(&m);
```

Decreases the value of **count** by 1

Increases the value of **count** by 1

Value of Semaphore	Thread 0	Thread 1
1		
1	call `sem_wait()`	
0	`sem_wait()` returns	
0	`(crit sect)`	
0	call `sem_post()`	
1	`sem_post()` returns	

  2. Counting semaphore
     - Resource with many units available
     - Is resource that allows certain kinds of unsynchronized concurrent access (e.g. reading)
     - Mutex has count = 1, counting has count = N
  3. Implementation Tips
     - Implement as few constraints as possible (e.g. how many semaphores do you need to solve the problem?)
     -

4.1 Example 1: Producer/Consumer (Bounded Buffer Problem)

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