

1. Let $A_0 = \begin{bmatrix} 10 & 8 & 12 & 15 & 15 \\ 8 & 14 & 8 & 7 & 9 \\ 12 & 8 & 10 & 13 & 9 \\ 15 & 7 & 13 & 4 & 10 \\ 15 & 9 & 9 & 10 & 4 \end{bmatrix}$, $A_1 = \begin{bmatrix} 12 & 11 & 14 & 10 & 3 \\ 11 & 14 & 10 & 14 & 6 \\ 14 & 10 & 16 & 18 & 4 \\ 10 & 14 & 18 & 18 & 8 \\ 3 & 6 & 4 & 8 & 8 \end{bmatrix}$, $A_2 = \begin{bmatrix} 4 & 13 & 12 & 16 & 6 \\ 13 & 4 & 14 & 9 & 15 \\ 12 & 14 & 6 & 5 & 5 \\ 16 & 9 & 5 & 2 & 6 \\ 6 & 15 & 5 & 6 & 8 \end{bmatrix}$.

Suppose $A: \mathbb{R}^2 \rightarrow S^5$ is defined by

$$A(x) = A_0 + x_1 A_1 + x_2 A_2.$$

Let $\lambda_1(x) \geq \lambda_2(x) \geq \lambda_3(x) \geq \lambda_4(x) \geq \lambda_5(x)$ denote the eigenvalues of $A(x)$.

- (a) Formulate the problem of minimizing the spread of the eigenvalues $\lambda_1(x) - \lambda_5(x)$ as an SDP. (10%)
- (b) Solve (a) by using MATLAB with the CVX tool. What are the optimal point and optimal value? (10%)

(a)

$$\begin{aligned} &\text{minimize} \quad \lambda_1(x) - \lambda_5(x) \\ &\text{subject to} \quad A(x) \leq \lambda_1(x) I \\ &\quad \quad \quad A(x) \geq \lambda_5(x) I \end{aligned}$$

(b)

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clear all; close all; clc;
A0 = [10 8 12 15 15; 8 14 8 7 9; 12 8 10 13 9; 15 7 13 4 10; 15 9 9 10 4];
A1 = [12 11 14 10 3; 11 14 10 14 6; 14 10 16 18 4; 10 14 18 18 8; 3 6 4 8 8];
A2 = [4 13 12 16 6; 13 4 14 9 15; 12 14 6 5 5; 16 9 5 2 6; 6 15 5 6 8];

cvx_begin sdp
    variables x1 x2 lamda1 lamda2
    minimize( lamda1 - lamda2 )
    subject to
        A = A0 + x1*A1 + x2*A2;
        A <= lamda1*eye(5);
        A >= lamda2*eye(5);
cvx_end

fprintf('x1 = %f, x2 = %f\n', x1, x2);
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Status: Solved

Optimal value (cvx_optval): +28.1544

x1 = -0.596605, x2 = -0.335843

optimal points

2. Prove that $x^* = (1, 1/2, -1)$ is optimal for the optimization problem

$$\begin{array}{ll} \text{minimize} & f_0(x) = (1/2)x^T P x + q^T x + r \\ \text{subject to} & -1 \leq x_i \leq 1, \quad i = 1, 2, 3 \end{array}$$

where

$$P = \begin{bmatrix} 13 & 12 & -2 \\ 12 & 17 & 6 \\ -2 & 6 & 12 \end{bmatrix}, \quad q = \begin{bmatrix} -22 \\ -14.5 \\ 13 \end{bmatrix}, \quad r = 1.$$

(Hint: you may verify the optimality condition.) (15%)

$$\begin{aligned} \nabla f_0(x) &= Px + q \Rightarrow \nabla f_0(x^*) = \begin{bmatrix} 13 & 12 & -2 \\ 12 & 17 & 6 \\ -2 & 6 & 12 \end{bmatrix} \begin{bmatrix} 1 \\ 1/2 \\ -1 \end{bmatrix} + \begin{bmatrix} -22 \\ -14.5 \\ 13 \end{bmatrix} \\ &= \begin{bmatrix} 21 \\ 14.5 \\ -11 \end{bmatrix} + \begin{bmatrix} -22 \\ -14.5 \\ 13 \end{bmatrix} = \begin{bmatrix} -1 \\ 0 \\ 2 \end{bmatrix} \end{aligned}$$

By the optimal condition, we know that if x^* is optimal point

$$\nabla f_0(x^*)^T (y - x^*) \geq 0; \quad \forall y \in \text{feasible set}$$

$$\Rightarrow [-1 \ 0 \ 2] \begin{bmatrix} y_1 - 1 \\ y_2 - 1/2 \\ y_3 + 1 \end{bmatrix} \geq 0$$

$$\Rightarrow -(y_1 - 1) + 2(y_3 + 1) \geq 0$$

$$\Rightarrow -y_1 + 2y_3 + 3 \geq 0$$

for all y satisfying $-1 \leq y_i \leq 1$; $i = 1, 2, 3$ is true $\#$

3. Consider a problem of the form

$$\begin{aligned} & \text{minimize} && f_0(x)/(c^T x + d) \\ & \text{subject to} && f_i(x) \leq 0, \quad i = 1, \dots, m \\ & && Ax = b \end{aligned}$$

where f_0, f_1, \dots, f_m are convex, and the domain of the objective function is defined as $\{x \in \text{dom } f_0 : c^T x + d > 0\}$.

Prove that the above problem is equivalent to the following problem:

$$\begin{aligned} & \text{minimize} && g_0(y, t) \\ & \text{subject to} && g_i(y, t) \leq 0, \quad i = 1, \dots, m \\ & && Ay = bt \\ & && c^T y + dt = 1 \end{aligned}$$

where g_i is the perspective of f_i (see §3.2.6). The variables are $y \in \mathbb{R}^n$ and $t \in \mathbb{R}$.

Further show that this problem is convex. (20%)

① Suppose x is feasible in original problem. Define

$$t = \frac{1}{c^T x + d}, \text{ where } t > 0; \quad y = \frac{x}{c^T x + d}$$

$$\Rightarrow g_0(y, t) = \frac{1}{t} f_0\left(\frac{y}{t}\right) = \frac{1}{t} \cdot f_0(x) = \frac{f_0(x)}{c^T x + d}$$

$$\Rightarrow g_i(y, t) \leq 0 \quad i = 1, \dots, m \quad \because f_i(x) \leq 0 \quad i = 1, \dots, m$$

$$\Rightarrow Ay = bt \Rightarrow \frac{Ax}{c^T x + d} = \frac{b}{c^T x + d} \Rightarrow Ax = b$$

$$\Rightarrow c^T y + dt = 1 \Rightarrow \frac{c^T x}{c^T x + d} + \frac{d}{c^T x + d} = 1$$

Then $t \cdot y$ are feasible in the equivalent problem

② Suppose y, t are feasible for the transformed problem.

We must have $t > 0$. Define $x = \frac{y}{t}$.

We have $x \in \text{dom } f_i$ for $i = 0, \dots, m$

$$\Rightarrow \frac{f_0(x)}{c^T x + d} = \frac{f_0\left(\frac{y}{t}\right)}{c^T \left(\frac{y}{t}\right) + d} = \frac{t f_0\left(\frac{y}{t}\right)}{c^T y + dt} = t f_0\left(\frac{y}{t}\right) = g_0(y, t)$$

$$\Rightarrow f_i(x) = f_i\left(\frac{y}{t}\right) = t g_i(y, t) \leq 0$$

$$\Rightarrow Ax = b \Rightarrow A\left(\frac{x}{t}\right) = b \Rightarrow Ay = bt$$

Then x is feasible in the original problem, with the objective value $f_0(y, t)$

③ So the transformed problem is equivalent to the original problem. #

And this problem is convex #

4. Formulate the following problem as a linear program.

$$\text{minimize} \quad \|Ax - b\|_\infty$$

where $x \in \mathbb{R}^n$, $A \in \mathbb{R}^{m \times n}$, and $b \in \mathbb{R}^m$. Show that your linear program is equivalent to the original problem. (15%)

Equivalent to the linear program

$$\begin{array}{ll} \text{minimize} & t \\ \text{subject to} & Ax - b \leq t \mathbf{1} \\ & Ax - b \geq -t \mathbf{1} \end{array}$$

in the variable x, t . To see the equivalence, assume x is fixed in this problem, and we optimize only over t .

The constraints say that

$$-t \leq a_k^T x - b_k \leq t$$

for each k , i.e., $t \geq |a_k^T x - b_k|$, i.e.,

$$t \geq \max_k |a_k^T x - b_k| = \|Ax - b\|_\infty$$

Clearly, if x is fixed, the optimal value of the LP is $p^*(x) = \|Ax - b\|_\infty$. Therefore optimizing over t and x simultaneously is equivalent to the original problem.

5. (Dual of general LP). Find the dual function of the LP

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Gx \leq h \\ & Ax = b.\end{array}$$

Give the dual problem, and make the implicit equality constraints explicit. (15%)

The Lagrangian is

$$\begin{aligned} \mathcal{L}(x, \lambda, v) &= c^T x + \lambda^T (Gx - h) + v^T (Ax - b) \\ &= (c^T + \lambda^T G + v^T A)x - \lambda^T h - v^T b \end{aligned}$$

which is an affine function of x .

① The dual function is given by

$$g(\lambda, v) = \inf_x \{ \mathcal{L}(x, \lambda, v) \} = \begin{cases} -\lambda^T h - v^T b & ; \text{ if } c^T + \lambda^T G + v^T A = 0 \\ -\infty & ; \text{ otherwise} \end{cases}$$

② The dual problem is

$$\begin{array}{ll}\text{maximize} & g(\lambda, v) \\ \text{subject to} & \lambda \geq 0\end{array}$$

③ After making the implicit constraint explicit, we have

$$\begin{array}{ll}\text{maximize} & -h^T \lambda - b^T v \\ \text{subject to} & c + G^T \lambda + A^T v = 0 \\ & \lambda \geq 0\end{array}$$

6. Derive a dual problem for

$$\text{minimize} \quad - \sum_{i=1}^m \log(b_i - a_i^T x)$$

with domain $\{x : a_i^T x < b_i, i = 1, \dots, m\}$. First introduce new variables y_i and equality constraints $y_i = b_i - a_i^T x$. (15%)

We derive the dual of the problem

$$\text{minimize} \quad - \sum_{i=1}^m \log y_i$$

$$\text{subject to} \quad y = b - Ax$$

where $A \in \mathbb{R}^{m \times n}$ has a_i^T as its i th row. The Lagrangian is

$$\mathcal{L}(x, y, v) = - \sum_{i=1}^m \log y_i + v^T (y - b + Ax)$$

and the dual function is

$$g(v) = \inf_{x, y} \{ \mathcal{L}(x, y, v) \} = \inf_{x, y} \left\{ - \sum_{i=1}^m \log y_i + v^T (y - b + Ax) \right\}$$

The term $v^T Ax$ is unbounded unless $A^T v = 0$ and the term in y are unbounded unless $v > 0$. And achieve their minimum for $y_i = 1/v_i$. The dual function

$$\textcircled{1} \quad g(v) = \begin{cases} \sum_{i=1}^m \log v_i + m - b^T v & ; A^T v = 0, v > 0 \\ -\infty & ; \text{otherwise} \end{cases}$$

And the dual problem

$$\textcircled{2} \quad \begin{aligned} &\text{maximize} \quad \sum_{i=1}^m \log v_i - b^T v + m \\ &\text{subject to} \quad A^T v = 0 \end{aligned}$$