

COMP 737011 - Memory Safety and Programming Language Design

# Lecture 12: Dynamic Analysis of Rust Code

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# Outline

- 1. Background of Dynamic Analysis
- 2. Blackbox Fuzzing
- 3. Fuzz Target Generation
- 4. Dynamic Symbolic Execution

# 1. Background of Dynamic Analysis

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# Any Bugs in genvec()?

```
fn genvec(s:&mut String)->Vec<u8>{
    let mut v = Vec::new();
    let l = s.len();
    if l%10 == 0 {
        let ptr = s.as_mut_ptr();
        unsafe{
            v = Vec::from_raw_parts(ptr,l,s.len());
        }
    }
    return v;
}

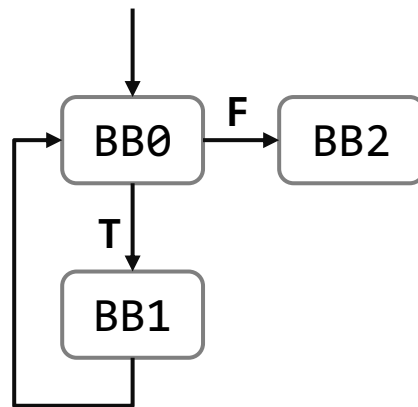
fn main(){
    let args:Vec<String> = env::args().collect();
    let mut s = String::from(&args[1]);
    let v = genvec(&mut s);
    println!("{:?}",v);
}
```

# Challenge of Bug Detection

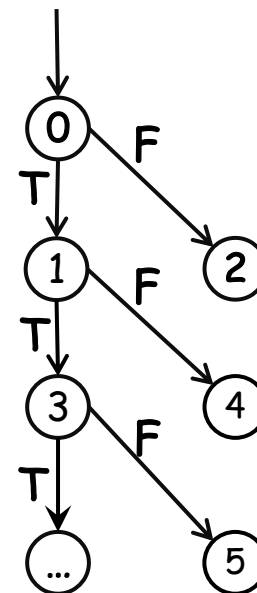
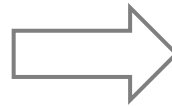
- Testing?
  - Bugs are on uncommon execution paths.
  - How to achieve high coverage (code/path)
    - define test cases manually?
    - automated test generation
  - Executing all paths is almost infeasible
- How to test Library APIs?

# Programs as Computation Trees

- A control-flow graph with loops unrolled.
- Each node is a a conditional statement (state)
- Each edge is a sequence of sequential statements
- Each path represents an equivalence class of inputs.



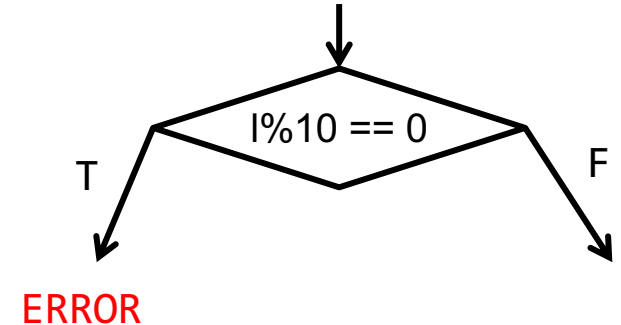
Control-flow graph



Computation tree

# Example of Computation Tree

```
fn genvec(s:&mut String)->Vec<u8>{  
  let mut v = Vec::new();  
  let l = s.len();  
  if l%10 == 0 {  
    let ptr = s.as_mut_ptr();  
    unsafe{  
      v = Vec::from_raw_parts(ptr,l,s.len());  
    }  
  }  
  return v;  
}
```



# How to Traverse the Tree?

- Random testing or fuzzing:
  - generate random inputs to execute the program.
  - probability of reaching error could be very small
- Symbolic execution:
  - Collect the constraints for each path
  - Use constraint solvers to find solutions

```
if x == 94389 {  
    unreachable!();  
}
```

Probability =  $1/2^{32}$



## 2. Blackbox Fuzzing

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# Evolutionary Fuzzing with AFL

- AFL is the most famous fuzzing tool
  - <http://lcamtuf.coredump.cx/afl/>
  - Many followup tools available, e.g., afl++
- Use a genetic approach to learn interesting mutations:
  - trigger code code => keep it
  - cannot trigger new code => discard

american fuzzy lop 0.47b (readpng)			
<b>process timing</b>		<b>overall results</b>	
run time : 0 days, 0 hrs, 4 min, 43 sec		cycles done : 0	
last new path : 0 days, 0 hrs, 0 min, 26 sec		total paths : 195	
last uniq crash : none seen yet		uniq crashes : 0	
last uniq hang : 0 days, 0 hrs, 1 min, 51 sec		uniq hangs : 1	
<b>cycle progress</b>		<b>map coverage</b>	
now processing : 38 (19.49%)		map density : 1217 (7.43%)	
paths timed out : 0 (0.00%)		count coverage : 2.55 bits/tuple	
<b>stage progress</b>		<b>findings in depth</b>	
now trying : interest 32/8		favored paths : 128 (65.64%)	
stage execs : 0/9990 (0.00%)		new edges on : 85 (43.59%)	
total execs : 654k		total crashes : 0 (0 unique)	
exec speed : 2306/sec		total hangs : 1 (1 unique)	
<b>fuzzing strategy yields</b>		<b>path geometry</b>	
bit flips : 88/14.4k, 6/14.4k, 6/14.4k		levels : 3	
byte flips : 0/1804, 0/1786, 1/1750		pending : 178	
arithmetics : 31/126k, 3/45.6k, 1/17.8k		pend fav : 114	
known ints : 1/15.8k, 4/65.8k, 6/78.2k		imported : 0	
havoc : 34/254k, 0/0		variable : 0	
trim : 2876 B/931 (61.45% gain)		latent : 0	



# Fuzz Rust with AFL

## 1. Install AFL

<https://rust-fuzz.github.io/book/introduction.html>

## 2. Create a new project

```
#: cargo install afl
```

```
#: cargo new fuzztarget
```

## 3. Add deps in Cargo.toml

```
[dependencies] afl = "*"
```

## 4. fuzz target

```
#[macro_use]
extern crate afl;
fn genvec(s:&mut String)->Vec<u8>{ ...}
fn main(){
    fuzz!(|data: &[u8]| {
        let mut s = String::from(&data);
        let v = genvec(&mut s);
    });
}
```

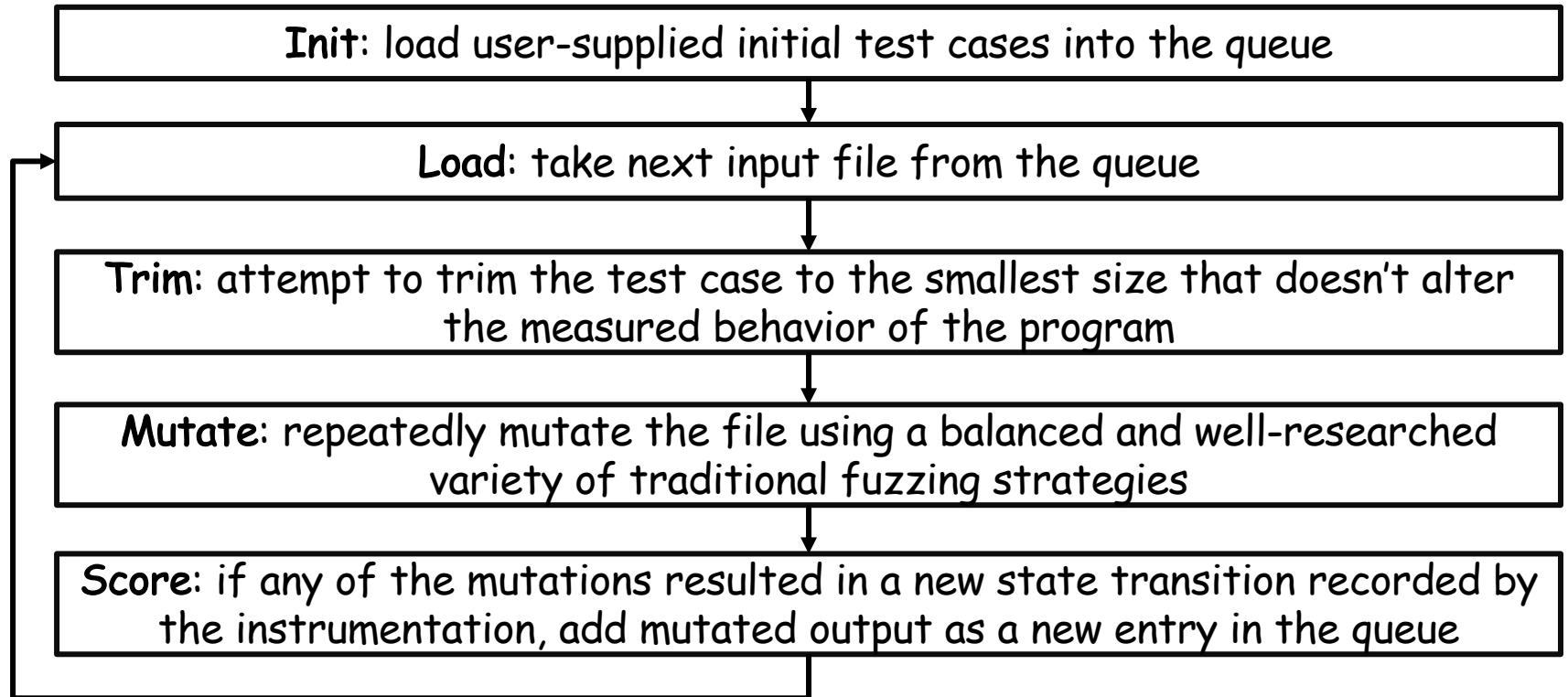
## 5. build the fuzz target

```
#: cargo afl build
```

## 6. add initial test seeds into a folder and then fuzz

```
#: cargo afl fuzz -i initcases -o output targetbin
```

# Overall Algorithm



The discovered test cases are also periodically culled to eliminate ones that have been obsoleted by newer, higher-coverage finds

# Instrumentation

- Support programs
  - written in C/C++/Objective C
  - variants for Python/Go/Rust/OCaml
- Code instrumented to observe execution paths:
  - if source code is available, using modified compiler
  - if source code is unavailable, running code in an emulator
- Code coverage represented as a 64KB bitmap
  - different executions may collide with a small chance.
  - tradeoff between collision and efficiency

# How to Instrument the Source Code?

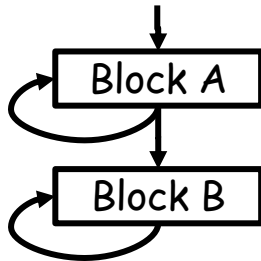
- Objective: Different jumps from src to dest should result in different offsets of the bitmap.
- Where to inject?
  - at every branching point
- What code to inject?
  - use a random number as the id/hash of the block.
  - update the corresponding bitmap cell

## Pseudocode:

```
cur_location = <COMPILE_TIME_RANDOM_FOR_THE_BLOCK>;  
bitmap[cur_location  $\oplus$  (prev_location >> 1)]++;
```

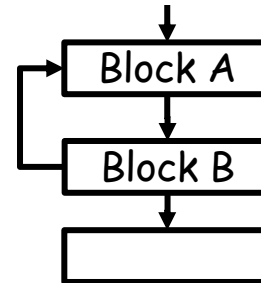
# More about the Design

- Why shift 1 bit? Consider the following cases



How to differentiate

- Block A → Block A
- Block B → Block B



How to differentiate

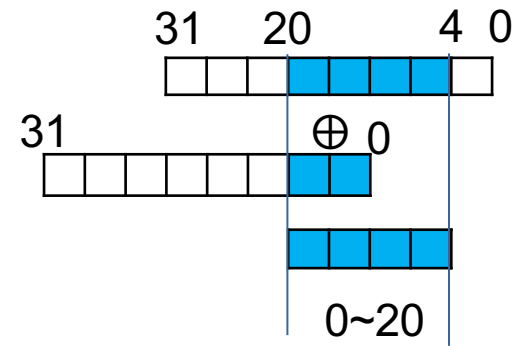
- Block A → Block B
- Block B → Block A

- Edge coverage in nature
- Collision rate?

Branch cnt	Colliding tuples	Example targets
1,000	0.75%	giflib, lzo
2,000	1.5%	zlib, tar, xz
5,000	3.5%	libpng, libwebp
10,000	7%	libxml
20,000	14%	sqlite
50,000	30%	-

# How to Instrument the Binaries

- Use the instruction address?
  - 32/64 bit
  - The bitmap offset is only 16bit;
- Select important bits heuristically.

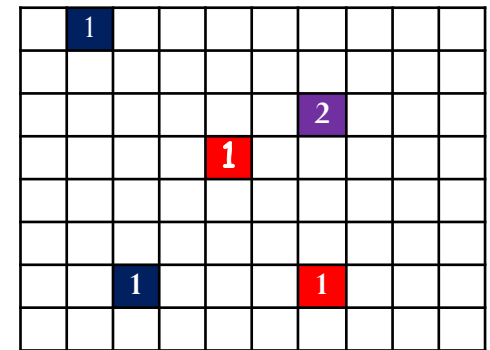
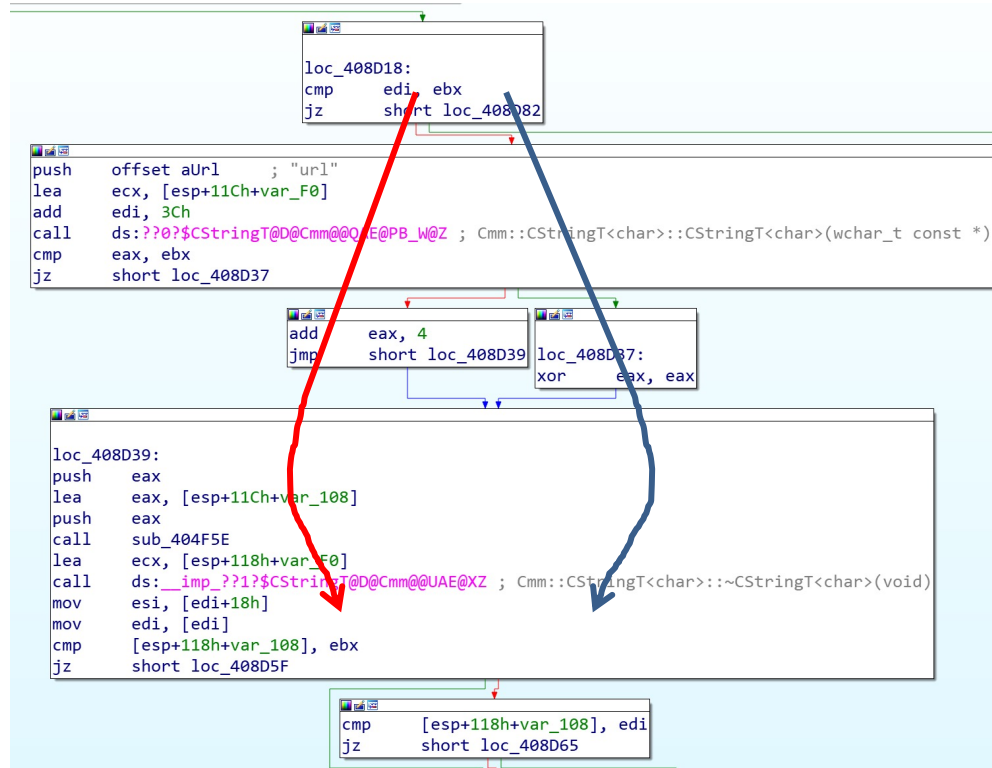


## Pseudocode:

```
cur_location = (block_addr >> 4)  $\oplus$  (block_addr << 8);  
bitmap[cur_location  $\oplus$  (prev_location >> 1)]++;
```



# Visualize Bitmap



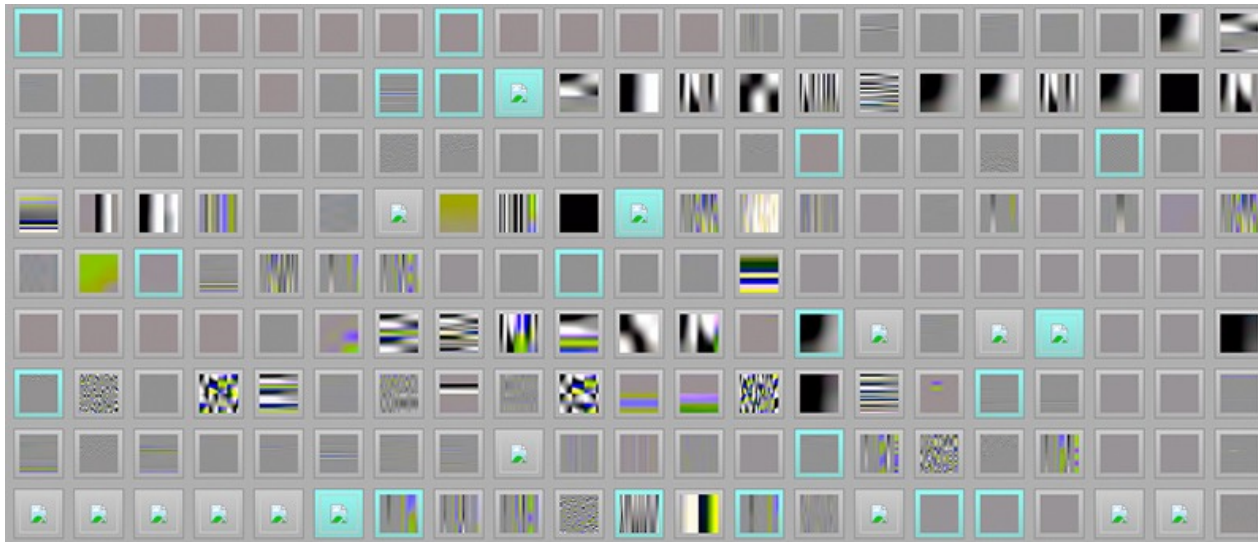
Bitmap (64KB)

# Sample Mutation Strategies

- Bit flip
- Increase/decrease integers
- Use predefined interesting values
  - *eg.*, 0, -1, MAX\_INT for integer
- Delete/combine/zero input block

# Example: Learning the JPG file format

- Task: to fuzz a program that expects a JPG as input
- Experiment: start with 'hello world' as initial test input
- Result:
  - Generate the first image after six hours on an 8-core system
  - Produced many interesting pics by using the image as a seed



# 3. Fuzz Target Generation

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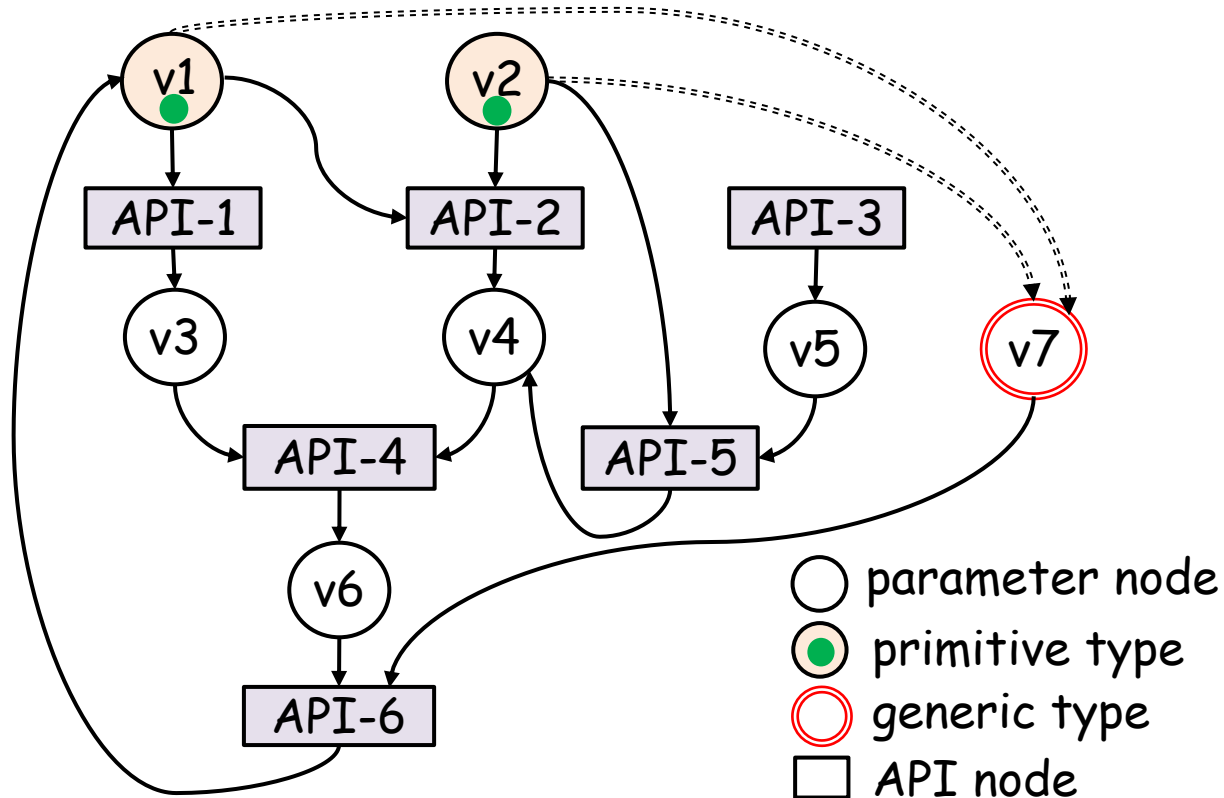
# Problem

- How to fuzz Rust libraries automatically?
- We need fuzz targets for library APIs
- How to generate fuzz target automatically?

```
#[macro_use]
extern crate afl;
fn genvec(s:&mut String)->Vec<u8>{ ...}
fn main(){
    fuzz!(|data: &[u8]| {
        let mut s = String::from(&data);
        let v = genvec(&mut s);
    });
}
```

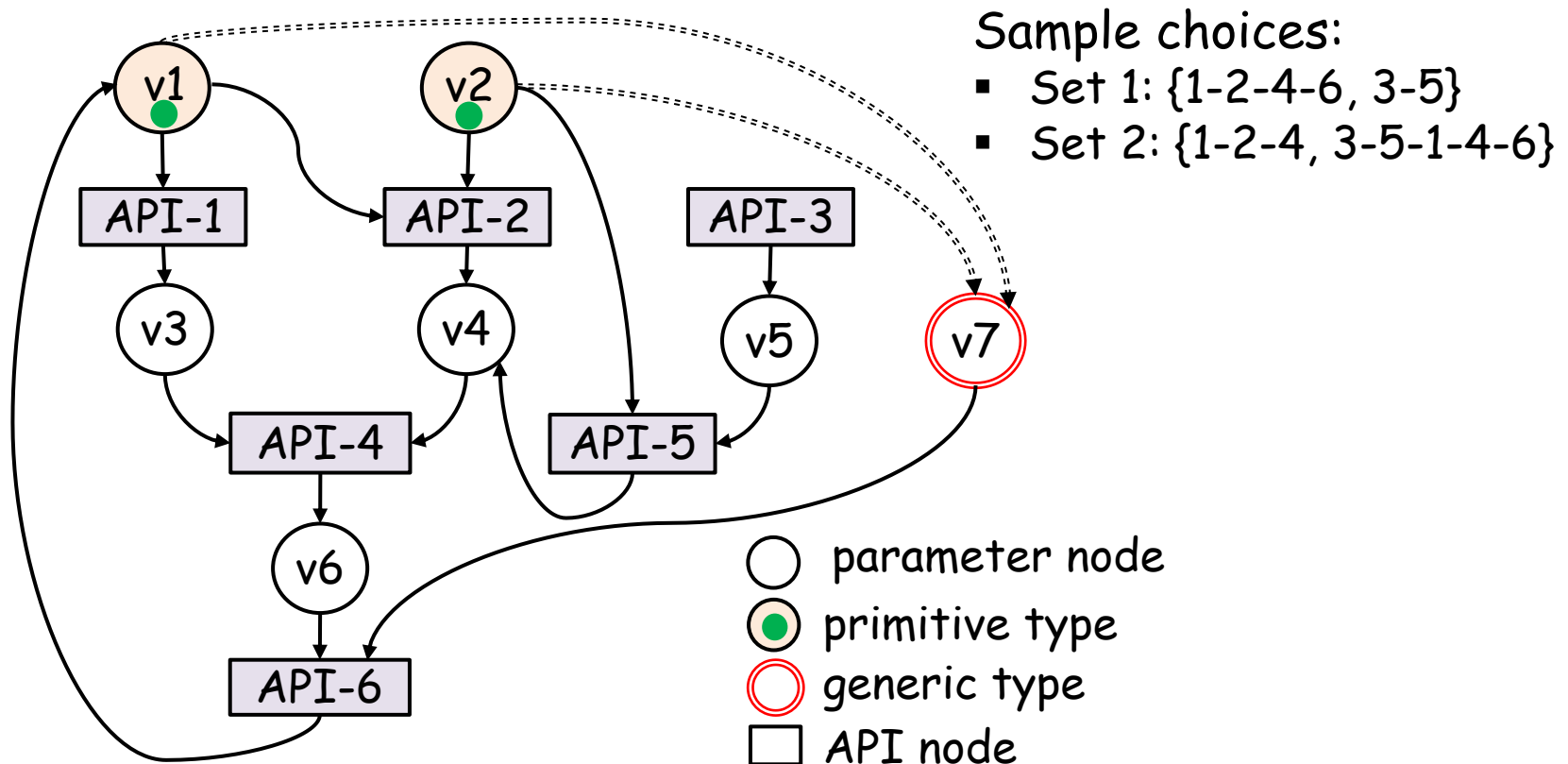
# Modeling the Dependencies of APIs

- Use PetriNet



# How to Traverse the Graph?

- Objective: high API coverage + efficiency
- Topological sort: generate short sequences
- Backward search: generate long sequences

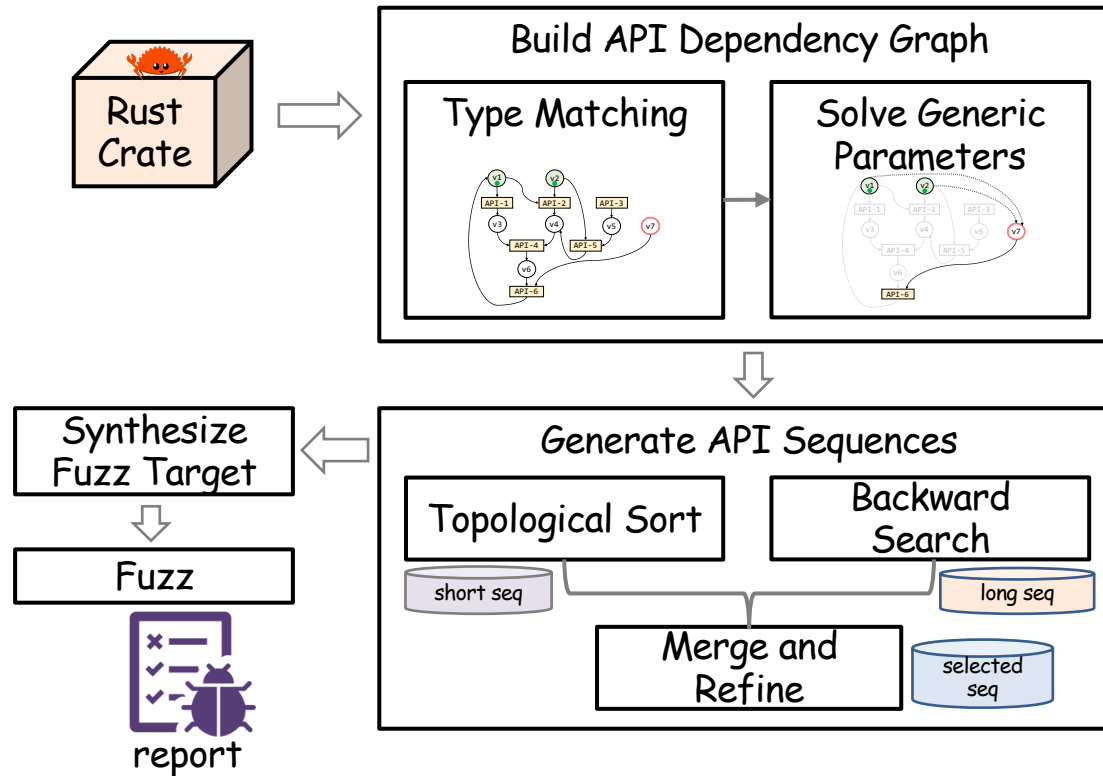


# Sequence Selection

- Selecting the sequence that contains the most uncovered API nodes
- Selecting the sequence that can activate the most uncovered parameter nodes
- Selecting the shortest sequence for those contribute equally using the above two rules
- Randomly picking one sequence if there are multiple candidates.



# Framework Overview



# 4. Dynamic Symbolic Execution

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Concolic (Concrete + Symbolic) Execution

# Dynamic Symbolic Execution

- Stores program state concretely and symbolically
- Step:
  - 1) Start with random input values
  - 2) Keep track of both concrete values and symbolic constraints
  - 3) Use concrete values to simplify symbolic constraints
  - 4) Solve the constraints

# An Illustrative Example

```
fn foo(x:usize)->usize{
    2*x
}

fn test(x:usize, y:usize){
    let z = foo(y);
    if (z == x){
        if (x > y+10){
            unreachable!();
        }
    }
}
```

Concrete  
Execution

Symbolic  
Execution

concrete  
state

symbolic  
state

path  
condition

$x = 22$

$x = x_0$

$2*y_0 \neq x_0$

$y = 7$

$y = y_0$

$z = 14$

$z = 2*y_0$

Solve:  $2*y_0 == x_0$

Solution:  $x_0 = 2, y_0 = 1$

# An Illustrative Example

```
fn foo(x:usize)->usize{
    2*x
}

fn test(x:usize, y:usize){
    let z = foo(y);
    if (z == x){
        if (x > y+10){
            unreachable!();
        }
    }
}
```

Concrete  
Execution

Symbolic  
Execution

concrete  
state

symbolic  
state

path  
condition

$x = 2$

$x = x_0$

$2 * y_0 == x_0$

$y = 1$

$y = y_0$

$x_0 \leq y_0 + 10$

$z = 2$

$z = 2 * y_0$

Solve:  $(2 * y_0 == x_0)$  and  $(x_0 > y_0 + 10)$   
Solution:  $x_0 = 30, y_0 = 15$

# An Illustrative Example

```
fn foo(x:usize)->usize{
    2*x
}

fn test(x:usize, y:usize){
    let z = foo(y);
    if (z == x){
        if (x > y+10){
            unreachable!();
        }
    }
}
```

Concrete  
Execution

Symbolic  
Execution

concrete  
state

symbolic  
state

path  
condition

x = 30

$x = x_0$

$2 * y_0 == x_0$

y = 15

$y = y_0$

$x_0 > y_0 + 10$

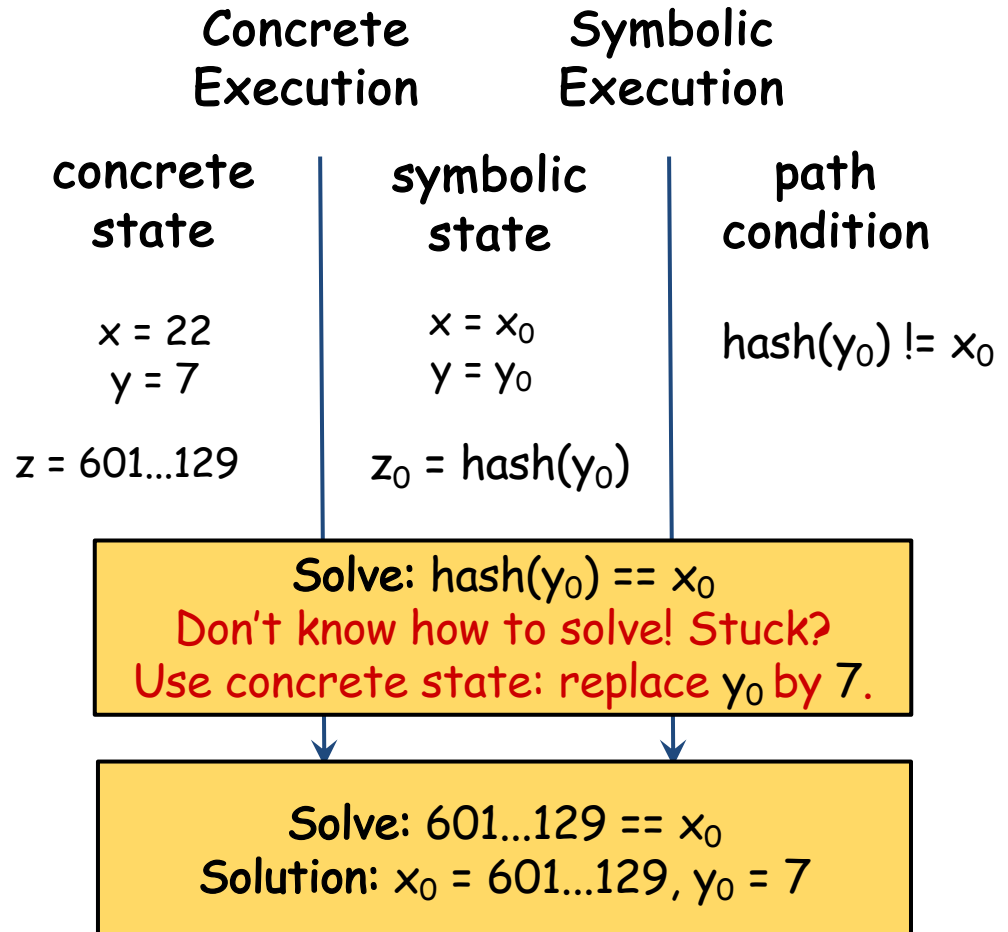
z = 30

$z = 2 * y_0$

# A More Complex Example

```
fn foo(x:usize)->usize{
    hash(x)
}

fn test(x:usize, y:usize){
    → let z = foo(y);
    → if (z == x){
        if (x > y+10){
            unreachable!();
        }
    }
}
```



# A Third Example

```
fn foo(x:usize)->usize{
    hash(x)
}

fn test(x:usize, y:usize){
  → let z = foo(y);
  → if (z == foo(x)){
      if (x > y+10){
          unreachable!();
      }
  }
}
```

Concrete Execution	Symbolic Execution	
concrete state	symbolic state	path condition
$x = 22$ $y = 7$	$x = x_0$ $y = y_0$	$x_0 \neq y_0$ $\text{hash}(x_0) \neq \text{hash}(y_0)$
Solve: $x_0 \neq y_0$ and $\text{hash}(x_0) == \text{hash}(y_0)$ Use concrete state: replace $y_0$ by 7.		
Solve: $x_0 \neq 7$ and $\text{hash}(x_0) == 601\dots129$ Use concrete state: replace $x_0$ by 22.		
Solve: $22 \neq 7$ and $438\dots861 == 601\dots129$ Unsatisfiable!		



# SE Tools Recommended

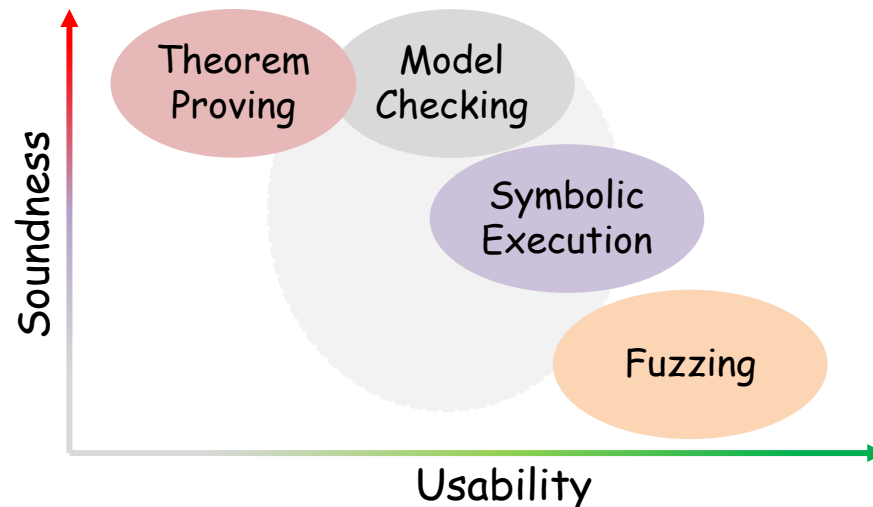
- angr: based on binary
  - <https://angr.io>
- KLEE: require source-code
  - <https://klee.github.io>

# Limitation of Symbolic Execution

- Symbolic reasoning (constraint modeling/solving)
  - array/dynamic memory
  - floating-point arithmetic
  - runtime environment
  - ...
- Path-explosion (scalability): large computation tree
  - loops
  - external functions
  - ...

# Comparison

- Formal methods should be sound
  - Theorem proving: prove some properties based on axioms
  - Model checking: check properties by enumerating each possible execution
    - e.g., via symbolic execution
- Fuzzing is the most easy-to-use approach



# More Reference

- AFL: <https://lcamtuf.coredump.cx/afl>
- angr: <https://angr.io>
- KLEE: <https://klee.github.io>
- "RULF: Rust library fuzzing via API dependency graph traversal", ASE, 2021.
- "Benchmarking the capability of symbolic execution tools with logic bombs", TDSC, 2020.