# **Algorithms Exercise #1**

## a. Environment

I. OS: Windows 11

II. Compiler: GNU GCC Compiler

III. IDE: Code::Blocks 20.03

## **b.** Results

- I. Methods:
  - i. Main function:
  - (1) Using *ifstream* and *ofstream* to input and output txt.file.

```
//read in txt
 94
95
              ifstream in;
              ofstream out;
in.open("input.txt");
out.open("output.txt");
if(out.fail()){
 97
 98
                    cout<<"Error opening output\n";
exit(1);</pre>
 99
100
101
102
              if(in.fail()){
   cout<<"Error opening input\n";</pre>
103
104
                    exit(1);
105
106
              int yt num;
107
              in>>yt_num;
```

(2) Using while loop to handle each tableau, and perform certain operation based on 'choice'. Using *getline* and *istringstream* to read in a whole line of input. The inserted numbers are stored in 'insert key' vector.

```
while(vt num!=0){
111
                in>>choice;
113
114
115
                if(choice==1){
                     vector<vector<int>> vt;
116
117
                     vector(int) insert_key;
                     int insert_num;
118
119
                     string 1;
in.get();
                    getline(in,1);
istringstream i(1);
120
                    while(i>>insert_num){
   insert_key.push_back(insert_num);
122
123
124
125
126
                     int insert_size=insert_key.size();
                     out<<"Insert
                     for(int i=0;i<insert_size;i++){
127
128
                         out<<insert_key[i]<<
```

(3) Store the tableau index in 'vector<vector<int>> yt', if the read in index is 'x', we store it as INF, which is define as INT MAX.

## ii. <u>Insert:</u>

(1) First, set a swap function for exchanging two indices. We will use it later.

(2) Then we insert the inserted number aka. 'key' in the last position of tableau (yt [m-1][n-1]), and perform the INSERT function, which reconstruct the tableau that make each index be in the correct position.

(3) In INSERT function, we use recursive concept. We move the bottom-right corner's index upwards and leftwards until it is in the correct position. We handle the first row and column separately, and the base case is when we reach the top-left corner.

```
void INSERT(vector<vector<int>>&yt,int i,int j){
              if(i==0 && j==0){
18
                    return;
19
              if(i==0){
20
                    if(yt[i][j]<yt[i][j-1]){
    swap(yt[i][j],yt[i][j-1]);
    INSERT(yt,i,j-1);</pre>
21
22
23
25
                     return;
26
27
              if(j==0){
                                       //first column
                     if(yt[i][j]<yt[i-1][j]){
    swap(yt[i][j],yt[i-1][j]);
    INSERT(yt,i-1,j);</pre>
28
29
31
32
                     return:
33
              if(yt[i][j]<yt[i-1][j]){
    swap(yt[i][j],yt[i-1][j]);
    INSERT(yt,i-1,j);</pre>
34
35
36
              if(yt[i][j]<yt[i][j-1]){
    swap(yt[i][j],yt[i][j-1]);
    INSERT(yt,i,j-1);</pre>
38
                                                               //left
39
41
42 }
```

## iii. Extract-min:

(1) We know that the min in the tableau is the top-left corner index, we assigned it to 'min' and set its position value to INF. Then use fix function to reconstruct the tableau.

```
76  void EXTRACT_MIN(vector<vector<int>>&yt,ofstream& out){
77  int min=yt[0][0];
78  yt[0][0]=INF;
79  fix(yt,0,0);
80  out<<"Extract-min "<<min<'\n";
81  }</pre>
```

(2) We consider (i, j) as the root node, and find the right and bottom nodes of it. Then we compare its right and bottom node, swap the root node with the smaller node. The function starts from (0, 0) node which value is INF, and implement the fix function recursively until its right and bottom nodes are both INF.

```
void fix(vector<vector<int>>&vt.int i.int i){
         int m=yt.size();
         int n=yt[0].size();
57
                             and right num of cur
         int bottom.right;
58
59
         if(i+1<m){bottom=yt[i+1][j];}</pre>
60
         else{bottom=INF;}
61
         if(j+1<n){right=yt[i][j+1];}</pre>
62
         else{right=INF;}
63
         if (bottom==INF && right==INF) {
64
65
              return;
66
         if(bottom<right){ //down</pre>
              swap(yt[i][j],yt[i+1][j]);
68
69
70
              fix(yt,i+1,j);
71
              swap(yt[i][j],yt[i][j+1]);
72
              fix(yt,i,j+1);
73
75
```

iv. Output: Using output function to output the result in the txt.file.

#### II. Running time analysis:

• m is the number of rows, n is the number of columns.

#### i. Insert:

- (1) 'INSERT' function: Time Complexity is O(m+n), which is the worst case that it needs to go upwards m times and leftwards n times to its correct position.
- (2) 'initial' function: Assume there are k inserted number (key), then it called k times INSERT function. Time Complexity is O (k\* (m+n)).

## ii. Extract-min:

- (1) 'EXTRACT-MIN' function called fix function. We divided the m\*n problem into (m-1)\*n or m\*(n-1) subproblem. The worst case is the index at the top-left corner moves downwards m times and rightwards n times, so the time complexity is O(m+n).
- iii. Output: Time Complexity is O(m\*n).